The Racket Reference

Version 6.1.1

Matthew Flatt
and PLT

November 4, 2014

This manual defines the core Racket language and describes its most prominent libraries. The companion manual ***The Racket Guide*** provides a friendlier (though less precise and less complete) overview of the language.

```
#lang racket/base
```
package: `base`

```
#lang racket
```

Unless otherwise noted, the bindings defined in this manual are exported by the `racket/base` and `racket` languages.

The `racket/base` library is much smaller than the `racket` library and will typically load faster.
The `racket` library combines `racket/base`, `racket/bool`, `racket/bytes`, `racket/class`, `racket/cmdline`, `racket/contract`, `racket/dict`, `racket/file`, `racket/format`, `racket/function`, `racket/future`, `racket/include`, `racket/list`, `racket/local`, `racket/match`, `racket/math`, `racket/path`, `racket/place`, `racket/port`, `racket/pretty`, `racket/promise`, `racket/sequence`, `racket/set`, `racket/shared`, `racket/stream`, `racket/string`, `racket/system`, `racket/tcp`, `racket/udp`, `racket/unit`, and `racket/vector`.

Contents

Index **1082**

1 Language Model

1.1 Evaluation Model

Racket evaluation can be viewed as the simplification of expressions to obtain values. For example, just as an elementary-school student simplifies

```
1 + 1 = 2
```

Racket evaluation simplifies

```
(+ 1 1) → 2
```

The arrow → above replaces the more traditional = to emphasize that evaluation proceeds in a particular direction towards simpler expressions. In particular, a *value* is an expression that evaluation simplifies no further, such as the number 2.

1.1.1 Sub-expression Evaluation and Continuations

Some simplifications require more than one step. For example:

```
(- 4 (+ 1 1)) → (- 4 2) → 2
```

An expression that is not a value can always be partitioned into two parts: a *redex*, which is the part that changed in a single-step simplification (highlighted), and the *continuation*, which is the evaluation context surrounding an expression. In `(- 4 (+ 1 1))`, the redex is `(+ 1 1)`, and the continuation is `(- 4 [])`, where [] takes the place of the redex. That is, the continuation says how to "continue" after the redex is reduced to a value.

Before some things can be evaluated, some sub-expressions must be evaluated; for example, in the application `(- 4 (+ 1 1))`, the application of `-` cannot be reduced until the sub-expression `(+ 1 1)` is reduced.

Thus, the specification of each syntactic form specifies how (some of) its sub-expressions are evaluated, and then how the results are combined to reduce the form away.

The *dynamic extent* of an expression is the sequence of evaluation steps during which the expression contains the redex.

1.1.2 Tail Position

An expression *expr1* is in *tail position* with respect to an enclosing expression *expr2* if, whenever *expr1* becomes a redex, its continuation is the same as was the enclosing *expr2*'s

continuation.

For example, the `(+ 1 1)` expression is *not* in tail position with respect to `(- 4 (+ 1 1))`. To illustrate, we use the notation *C*[*expr*] to mean the expression that is produced by substituting *expr* in place of [] in the continuation *C*:

C[`(- 4 (+ 1 1))`] → *C*[`(- 4 2)`]

In this case, the continuation for reducing `(+ 1 1)` is *C*[`(- 4 [])`], not just *C*.

In contrast, `(+ 1 1)` is in tail position with respect to `(if (zero? 0) (+ 1 1) 3)`, because, for any continuation *C*,

C[`(if (zero? 0) (+ 1 1) 3)`] → *C*[`(if #t (+ 1 1) 3)`] → *C*[`(+ 1 1)`]

The steps in this reduction sequence are driven by the definition of `if`, and they do not depend on the continuation *C*. The "then" branch of an `if` form is always in tail position with respect to the `if` form. Due to a similar reduction rule for `if` and `#f`, the "else" branch of an `if` form is also in tail position.

Tail-position specifications provide a guarantee about the asymptotic space consumption of a computation. In general, the specification of tail positions goes with each syntactic form, like `if`.

1.1.3 Multiple Return Values

A Racket expression can evaluate to *multiple values*, in the same way that a procedure can accept multiple arguments.

Most continuations expect a particular number of result values. Indeed, most continuations, such as `(+ [] 1)` expect a single value. The continuation `(let-values ([(x y) []]) expr)` expects two result values; the first result replaces `x` in the body *expr*, and the second replaces `y` in *expr*. The continuation `(begin [] (+ 1 2))` accepts any number of result values, because it ignores the result(s).

In general, the specification of a syntactic form indicates the number of values that it produces and the number that it expects from each of its sub-expression. In addition, some procedures (notably `values`) produce multiple values, and some procedures (notably `call-with-values`) create continuations internally that accept a certain number of values.

1.1.4 Top-Level Variables

Given

```
x = 10
```

then an algebra student simplifies `x + 1` as follows:

```
x + 1 = 10 + 1 = 11
```

Racket works much the same way, in that a set of top-level variables are available for substitutions on demand during evaluation. For example, given

```
(define x 10)
```

then

```
(+ x 1) → (+ 10 1) → 11
```

In Racket, the way definitions appear is just as important as the way that they are used. Racket evaluation thus keeps track of both definitions and the current expression, and it extends the set of definitions in response to evaluating forms such as `define`.

Each evaluation step, then, takes the current set of definitions and program to a new set of definitions and program. Before a `define` can be moved into the set of definitions, its right-hand expression must be reduced to a value.

```
   defined:
   evaluate: (begin (define x (+ 9 1)) (+ x 1))
→ defined:
   evaluate: (begin (define x 10) (+ x 1))
→ defined: (define x 10)
   evaluate: (begin (void) (+ x 1))
→ defined: (define x 10)
   evaluate: (+ x 1)
→ defined: (define x 10)
   evaluate: (+ 10 1)
→ defined: (define x 10)
   evaluate: 11
```

Using `set!`, a program can change the value associated with an existing top-level variable:

```
   defined: (define x 10)
   evaluate: (begin (set! x 8) x)
→ defined: (define x 8)
   evaluate: (begin (void) x)
→ defined: (define x 8)
   evaluate: x
→ defined: (define x 8)
   evaluate: 8
```

1.1.5 Objects and Imperative Update

In addition to `set!` for imperative update of top-level variables, various procedures enable the modification of elements within a compound data structure. For example, `vector-set!` modifies the content of a vector.

To allow such modifications to data, we must distinguish between values, which are the results of expressions, and *objects*, which hold the data referenced by a value.

A few kinds of objects can serve directly as values, including booleans, `(void)`, and small exact integers. More generally, however, a value is a reference to an object. For example, a value can be a reference to a particular vector that currently holds the value `10` in its first slot. If an object is modified, then the modification is visible through all copies of the value that reference the same object.

In the evaluation model, a set of objects must be carried along with each step in evaluation, just like the definition set. Operations that create objects, such as `vector`, add to the set of objects:

```
   objects:
   defined:
   evaluate: (begin (define x (vector 10 20))
                    (define y x)
                    (vector-set! x 0 11)
                    (vector-ref y 0))
→  objects:  (define <o1> (vector 10 20))
   defined:
   evaluate: (begin (define x <o1>)
                    (define y x)
                    (vector-set! x 0 11)
                    (vector-ref y 0))
→  objects:  (define <o1> (vector 10 20))
   defined:  (define x <o1>)
   evaluate: (begin (void)
                    (define y x)
                    (vector-set! x 0 11)
                    (vector-ref y 0))
→  objects:  (define <o1> (vector 10 20))
   defined:  (define x <o1>)
   evaluate: (begin (define y x)
                    (vector-set! x 0 11)
                    (vector-ref y 0))
→  objects:  (define <o1> (vector 10 20))
   defined:  (define x <o1>)
```

```
  evaluate: (begin (define y <o1>)
                   (vector-set! x 0 11)
                   (vector-ref y 0))
→ objects: (define <o1> (vector 10 20))
  defined: (define x <o1>)
           (define y <o1>)
  evaluate: (begin (void)
                   (vector-set! x 0 11)
                   (vector-ref y 0))
→ objects: (define <o1> (vector 10 20))
  defined: (define x <o1>)
           (define y <o1>)
  evaluate: (begin (vector-set! x 0 11)
                   (vector-ref y 0))
→ objects: (define <o1> (vector 10 20))
  defined: (define x <o1>)
           (define y <o1>)
  evaluate: (begin (vector-set! <o1> 0 11)
                   (vector-ref y 0))
→ objects: (define <o1> (vector 11 20))
  defined: (define x <o1>)
           (define y <o1>)
  evaluate: (begin (void)
                   (vector-ref y 0))
→ objects: (define <o1> (vector 11 20))
  defined: (define x <o1>)
           (define y <o1>)
  evaluate: (vector-ref y 0)
→ objects: (define <o1> (vector 11 20))
  defined: (define x <o1>)
           (define y <o1>)
  evaluate: (vector-ref <o1> 0)
→ objects: (define <o1> (vector 11 20))
  defined: (define x <o1>)
           (define y <o1>)
  evaluate: 11
```

The distinction between a top-level variable and an object reference is crucial. A top-level variable is not a value; each time a variable expression is evaluated, the value is extracted from the current set of definitions. An object reference, in contrast is a value, and therefore needs no further evaluation. The model evaluation steps above use angle-bracketed `<o1>` for an object reference to distinguish it from a variable name.

A direct object reference can never appear in a text-based source program. A program rep-

resentation created with datum->syntax, however, can embed direct references to existing objects.

1.1.6 Object Identity and Comparisons

The eq? operator compares two values, returning #t when the values refer to the same object. This form of equality is suitable for comparing objects that support imperative update (e.g., to determine that the effect of modifying an object through one reference is visible through another reference). Also, an eq? test evaluates quickly, and eq?-based hashing is more lightweight than equal?-based hashing in hash tables.

In some cases, however, eq? is unsuitable as a comparison operator, because the generation of objects is not clearly defined. In particular, two applications of + to the same two exact integers may or may not produce results that are eq?, although the results are always equal?. Similarly, evaluation of a lambda form typically generates a new procedure object, but it may re-use a procedure object previously generated by the same source lambda form.

The behavior of a datatype with respect to eq? is generally specified with the datatype and its associated procedures.

1.1.7 Garbage Collection

See §16 “Memory Management” for functions related to garbage collection.

In the program state

```
objects: (define <o1> (vector 10 20))
         (define <o2> (vector 0))
defined: (define x <o1>)
evaluate: (+ 1 x)
```

evaluation cannot depend on <o2>, because it is not part of the program to evaluate, and it is not referenced by any definition that is accessible in the program. The object <o2> may therefore be removed from the evaluation by *garbage collection*.

A few special compound datatypes hold *weak references* to objects. Such weak references are treated specially by the garbage collector in determining which objects are reachable for the remainder of the computation. If an object is reachable only via a weak reference, then the object can be reclaimed, and the weak reference is replaced by a different value (typically #f).

As a special case, a fixnum is always considered reachable by the garbage collector. Many other values are always reachable due to the way they are implemented and used: A character in the Latin-1 range is always reachable, because equal? Latin-1 characters are always eq?, and all of the Latin-1 characters are referenced by an internal module. Similarly, null, #t, #f, eof, and #<void> and are always reachable. Values produced by quote remain

reachable when the `quote` expression itself is reachable.

1.1.8 Procedure Applications and Local Variables

Given

```
f(x) = x + 10
```

then an algebra student simplifies `f(7)` as follows:

```
f(7) = 7 + 10 = 17
```

The key step in this simplification is take the body of the defined function `f`, and then replace each `x` with the actual value `7`.

Racket procedure application works much the same way. A procedure is an object, so evaluating `(f 7)` starts with a variable lookup:

```
   objects: (define <p1> (lambda (x) (+ x 10)))
   defined: (define f <p1>)
   evaluate: (f 7)
→ objects: (define <p1> (lambda (x) (+ x 10)))
   defined: (define f <p1>)
   evaluate: (<p1> 7)
```

Unlike in algebra, however, the value associated with an argument can be changed in the body of a procedure by using `set!`, as in the example `(lambda (x) (begin (set! x 3) x))`. Since the value associated with `x` can be changed, an actual value cannot be substituted for `x` when the procedure is applied.

Instead, a new *location* is created for each variable on each application. The argument value is placed in the location, and each instance of the variable in the procedure body is replaced with the new location:

```
   objects: (define <p1> (lambda (x) (+ x 10)))
   defined: (define f <p1>)
   evaluate: (<p1> 7)
→ objects: (define <p1> (lambda (x) (+ x 10)))
   defined: (define f <p1>)
            (define xloc 7)
   evaluate: (+ xloc 10)
→ objects: (define <p1> (lambda (x) (+ x 10)))
   defined: (define f <p1>)
            (define xloc 7)
   evaluate: (+ 7 10)
```

```
→ objects:  (define <p1> (lambda (x) (+ x 10)))
  defined:  (define f <p1>)
            (define xloc 7)
  evaluate: 17
```

A location is the same as a top-level variable, but when a location is generated, it (conceptually) uses a name that has not been used before and that cannot be generated again or accessed directly.

Generating a location in this way means that `set!` evaluates for local variables in the same way as for top-level variables, because the local variable is always replaced with a location by the time the `set!` form is evaluated:

```
  objects:  (define <p1> (lambda (x) (begin (set! x 3) x)))
  defined:  (define f <p1>)
  evaluate: (f 7)
→ objects:  (define <p1> (lambda (x) (begin (set! x 3) x)))
  defined:  (define f <p1>)
  evaluate: (<p1> 7)
→ objects:  (define <p1> (lambda (x) (begin (set! x 3) x)))
  defined:  (define f <p1>)
            (define xloc 7)
  evaluate: (begin (set! xloc 3) xloc)
→ objects:  (define <p1> (lambda (x) (begin (set! x 3) x)))
  defined:  (define f <p1>)
            (define xloc 3)
  evaluate: (begin (void) xloc)
→ objects:  (define <p1> (lambda (x) (begin (set! x 3) x)))
  defined:  (define f <p1>)
            (define xloc 3)
  evaluate: xloc
→ objects:  (define <p1> (lambda (x) (begin (set! x 3) x)))
  defined:  (define f <p1>)
            (define xloc 3)
  evaluate: 3
```

The substitution and location-generation step of procedure application requires that the argument is a value. Therefore, in `((lambda (x) (+ x 10)) (+ 1 2))`, the `(+ 1 2)` subexpression must be simplified to the value `3`, and then `3` can be placed into a location for `x`. In other words, Racket is a *call-by-value* language.

Evaluation of a local-variable form, such as `(let ([x (+ 1 2)]) expr)`, is the same as for a procedure call. After `(+ 1 2)` produces a value, it is stored in a fresh location that replaces every instance of `x` in *expr*.

1.1.9 Variables and Locations

A *variable* is a placeholder for a value, and expressions in an initial program refer to variables. A *top-level variable* is both a variable and a location. Any other variable is always replaced by a location at run-time, so that evaluation of expressions involves only locations. A single *local variable* (i.e., a non-top-level, non-module-level variable), such as a procedure argument, can correspond to different locations through different instantiations.

For example, in the program

```
(define y (+ (let ([x 5]) x) 6))
```

both `y` and `x` are variables. The `y` variable is a top-level variable, and the `x` is a local variable. When this code is evaluated, a location is created for `x` to hold the value `5`, and a location is also created for `y` to hold the value `11`.

The replacement of a variable with a location during evaluation implements Racket's *lexical scoping*. For example, when a procedure-argument variable `x` is replaced by the location `xloc`, then it is replaced throughout the body of the procedure, including any nested `lambda` forms. As a result, future references of the variable always access the same location.

1.1.10 Modules and Module-Level Variables

See §3.1 "Modules: `module`, `module*`, ..." for the syntax of modules.

Most definitions in Racket are in *modules*. In terms of evaluation, a module is essentially a prefix on a defined name, so that different modules can define the name. That is, a *module-level variable* is like a top-level variable from the perspective of evaluation.

One difference between a module and a top-level definition is that a module can be declared without instantiating its module-level definitions. Evaluation of a `require` *instantiates* (i.e., triggers the *instantiation* of) a declared module, which creates variables that correspond to its module-level definitions.

For example, given the module declaration

```
(module m racket
  (define x 10))
```

the evaluation of `(require 'm)` creates the variable `x` and installs `10` as its value. This `x` is unrelated to any top-level definition of `x`.

Phases

A module can be instantiated in multiple *phases*. A phase is an integer that, again, is effectively a prefix on the names of module-level definitions. A top-level `require` instantiates a

module at phase 0, if the module is not already instantiated at phase 0. A top-level `(require (for-syntax ....))` instantiates a module at phase 1 (if it is not already instantiated at that level); `for-syntax` also has a different binding effect on further program parsing, as described in §1.2.3.4 “Introducing Bindings”.

Within a module, some definitions are shifted by a phase already; the `begin-for-syntax` form is similar to `begin`, but it shifts expressions and definitions by a relative phase 1. Thus, if the module is instantiated at phase 1, the variables defined with `begin-for-syntax` are created at phase 2, and so on. Moreover, this relative phase acts as another layer of prefixing, so that a `define` of `x` and a `begin-for-syntax`-wrapped `define` of `x` can co-exist in a module without colliding. A `begin-for-syntax` form can be nested within a `begin-for-syntax` form, in which case definitions and expressions are in relative phase 2, and so on. Higher phases are mainly related to program parsing, instead of normal evaluation.

If a module instantiated at phase *n* `requires` another module, then the `required` module is first instantiated at phase *n*, and so on transitively. (Module `requires` cannot form cycles.) If a module instantiated at phase *n* `requires` `for-syntax` another module, the other module becomes *available* at phase *n*+1, and it may later be instantiated at phase *n*+1. If a module that is available at phase *n* for *n*>0 `requires` `for-template` another module, the other module becomes available at phase *n*-1, and so on. Instantiations of available modules above phase 0 are triggered on demand as described in §1.2.3.8 “Module Phases and Visits”.

A final distinction among module instantiations is that multiple instantiations may exist at phase 1 and higher. These instantiations are created by the parsing of module forms (see §1.2.3.8 “Module Phases and Visits”), and are, again, conceptually distinguished by prefixes.

Top-level variables can exist in multiple phases in the same way as within modules. For example, `define` within `begin-for-syntax` creates a phase 1 variable. Furthermore, reflective operations like `make-base-namespace` and `eval` provide access to top-level variables in higher phases, while module instantiations (triggered by `require`) relative to such top-levels are in corresponding higher phases.

The Separate Compilation Guarantee

When a module is compiled, its phase 1 is instantiated. This can, in turn, trigger the transitive instantiation of many other modules at other phases, including phase 1. Racket provides a very strong guarantee about this instantiation called "The Separate Compilation Guarantee":

"Any effects of the instantiation of the module’s phase 1 due to compilation on the Racket runtime system are discarded."

The guarantee concerns *effects*. There are two different kinds of effects: internal and external.

Internal effects are exemplified by mutation. Mutation is the action of a function such as `set-box!`, which changes the value contained in the box. The modified box is not observable outside of Racket, so the effect is said to be "internal". By definition, internal effects is

not detectable outside of the Racket program.

External effects are exemplified by input/output (or I/O). I/O is the action of a function such as `tcp-connect`, which communicates with the operating system to send network packets outside of the machine running Racket. The transmission of these packets is observable outside of Racket, in particular by the receiver computer or any routers in between. External effects exist to be detectable outside of the Racket program and are often detectable using physical processes.

An effect is *discarded* when it is no longer detectable. For instance, a mutation of a box from `3` to `4` would be discarded if it ceases to be detectable that it was ever changed, and thus would still contain `3`. Because external effects are intrinsically observable outside of Racket, they are irreversible and cannot be discarded.

Thus, The Separate Compilation Guarantee only concerns effects like mutation, because they are exclusively effects "on the Racket runtime system" and not "on the physical universe".

There are many things a Racket program can do that appear to be internal effects, but are actually external effects. For instance, `bytes-set!` is typically an internal effect, except when the bytes were created by `make-shared-bytes` which is allocated in space observable by other processes. Thus, effects which modify them are not discardable, so `bytes-set!`, in this case, is an external effect.

The opposite is also true: some things which appear to be external are actually internal. For instance, if a Racket program starts multiple threads and uses mutation to communicate between them, that mutation is purely internal, because Racket's threads are defined entirely internally.

Furthermore, whenever a Racket program calls an unsafe function, the Racket runtime system makes no promises about its effects. For instance, all foreign calls use `ffi/unsafe`, so all foreign calls are unsafe and their effects cannot be discarded by Racket.

Finally, The Separate Compilation Guarantee only concerns instantiations at phase 1 during compilation and not all phase 1 instantiations generally, such as when its phase 1 is required and used for effects via reflective mechanisms.

The practical consequence of this guarantee is that because effects are never visible, no module can detect whether a module it `requires` is already compiled. Thus, it can never change the compilation of one module to have already compiled a different module. In particular, if module A is shared by the phase 1 portion of modules X and Y, then any internal effects while X is compiled are not visible during the compilation of Y, regardless of whether X and Y are compiled during the same execution of Racket's runtime system.

The following set of modules demonstrate this guarantee. First, we define a module with the ability to observe effects via a `box`:

```
(module box racket/base
  (provide (all-defined-out))
```

```
(define b (box 0)))
```

Next, we define two syntax transformers that use and mutate this box:

```
(module transformers racket/base
  (provide (all-defined-out))
  (require (for-syntax racket/base
                       'box))
  (define-syntax (sett stx)
    (set-box! b 2)
    #'(void))
  (define-syntax (gett stx)
    #`#,(unbox b)))
```

Next, we define a module that uses these transformers:

```
(module user racket/base
  (provide (all-defined-out))
  (require 'transformers)
  (sett)
  (define gott (gett)))
```

Finally, we define a second module that uses these transformers:

```
(module test racket/base
  (require 'box 'transformers 'user)
  (displayln gott)
  (displayln (gett))

  (sett)
  (displayln (gett))

  (displayln (unbox b)))
```

This module displays:

- `2`, because the module `'user` expanded to `2`.
- `0`, because the effects of compiling `'user` were discarded.
- `2`, because the effect of `(sett)` inside `'test` is not discarded.
- `0`, because the effects at phase 1 are irrelevant to the phase 0 use of `b`.

Furthermore, this display will never change, regardless of which order these modules are compiled in or whether they are compiled at the same time or separately.

In contrast, if these modules were changed to store the value of `b` in a file on the filesystem, then the program would only display `2`.

The Separate Compilation Guarantee is described in more detail in "Composable and Compilable Macros" [Flatt02], including informative examples. The paper "Advanced Macrology and the implementation of Typed Scheme" [Culpepper07] also contains an extended example of why it is important and how to design effectful syntactic extensions in its presence.

Cross-Phase Persistent Modules

Module declarations that fit a highly constrained form—including a `(#%declare #:cross-phase-persistent)` form in the module body—create *cross-phase persistent* modules. A cross-phase persistent module's instantiations across all phases and module registries share the variables produced by the first instantiation of the module.

The intent of a cross-phase persistent module is to support values that are recognizable after phase crossings. For example, when a macro transformer running in phase 1 raises a syntax error as represented by a `exn:fail:syntax` instance, the instance is recognizable by a phase-0 exception handler wrapping a call to `eval` or `expand` that triggered the syntax error, because the `exn:fail:syntax` structure type is defined by a cross-phase persistent module.

A cross-phase persistent module imports only other cross-phase persistent modules, and it contains only definitions that bind variables to functions, structure types and related functions, or structure-type properties and related functions. A cross-phase persistent module never includes syntax literals (via `quote-syntax`) or variable references (via `#%variable-reference`). See §1.2.7 "Cross-Phase Persistent Module Declarations" for the syntactic specification of a cross-phase persistent module declaration.

A documented module should be assumed non-cross-phase persistent unless it is specified as cross-phase persistent (such as `racket/kernel`).

Module Redeclarations

When a module is declared using a name for which a module is already declared, the new declaration's definitions replace and extend the old declarations. If a variable in the old declaration has no counterpart in the new declaration, the old variable continues to exist, but its binding is not included in the lexical information for the module body. If a new variable definition has a counterpart in the old declaration, it effectively assigns to the old variable.

If a module is instantiated in any phases before it is redeclared, each redeclaration of the module is immediately instantiated in the same phases.

If the current inspector does not manage a module's declaration inspector (see §14.10 "Code Inspectors"), then the module cannot be redeclared. Similarly, a cross-phase persistent mod-

ule cannot be redeclared. Even if redeclrection succeeds, instantiation of a module that is previously instantiated may fail if instantiation for the redeclaration attempts to modify variables that are constant (see `compile-enforce-module-constants`).

Submodules

A `module` or `module*` form within a top-level `module` form declares a *submodule*. A submodule is accessed relative to its enclosing module, usually with a `submod` path. Submodules can be nested to any depth.

Although a submodule is lexically nested within a module, it cannot necessarily access the bindings of its enclosing module directly. More specifically, a submodule declared with `module` cannot `require` from its enclosing module, but the enclosing module can `require` the submodule. In contrast, a submodule declared with `module*` conceptually follows its enclosing module, so can `require` from its enclosing module, but the enclosing module cannot `require` the submodule. Unless a submodule imports from its enclosing module or vice-versa, then visits or instantiations of the two modules are independent, and thier implementations may even be loaded from bytecode at different times.

A submodule declared with `module` can import any preceding submodule declared with `module`. A submodule declared with `module*` can import any preceding module declared with `module*` and any submodule declared with `module`.

When a submodule declaration has the form `(module* name #f ....)`, then all of the bindings of the enclosing module's bodies are visible in the submodule's body, and the submodule implicitly imports the enclosing module. The submodule can `provide` any bindings that it inherits from its enclosing module.

1.1.11 Continuation Frames and Marks

See §10.5 "Continuation Marks" for continuation-mark forms and functions.

Every continuation C can be partitioned into *continuation frames* C_1, C_2, ..., C_n such that $C = C_1[C_2[...[C_n]]]$, and no frame C_i can be itself partitioned into smaller continuations. Evaluation steps add and remove frames to the current continuation, typically one at a time.

Each frame is conceptually annotated with a set of *continuation marks*. A mark consists of a key and its value; the key is an arbitrary value, and each frame includes at most one mark for any key. Various operations set and extract marks from continuations, so that marks can be used to attach information to a dynamic extent. For example, marks can be used to record information for a "stack trace" to be used when an exception is raised, or to implement dynamic scope.

1.1.12 Prompts, Delimited Continuations, and Barriers

See §10.4 "Continuations" for continuation and prompt functions.

A *prompt* is a special kind of continuation frame that is annotated with a specific *prompt tag* (essentially a continuation mark). Various operations allow the capture of frames in the continuation from the redex position out to the nearest enclosing prompt with a particular prompt tag; such a continuation is sometimes called a *delimited continuation*. Other operations allow the current continuation to be extended with a captured continuation (specifically, a *composable continuation*). Yet other operations abort the computation to the nearest enclosing prompt with a particular tag, or replace the continuation to the nearest enclosing prompt with another one. When a delimited continuation is captured, the marks associated with the relevant frames are also captured.

A *continuation barrier* is another kind of continuation frame that prohibits certain replacements of the current continuation with another. Specifically, a continuation can be replaced by another only when the replacement does not introduce any continuation barriers. It may remove continuation barriers only through jumps to continuations that are a tail of the current continuation. A continuation barrier thus prevents "downward jumps" into a continuation that is protected by a barrier. Certain operations install barriers automatically; in particular, when an exception handler is called, a continuation barrier prohibits the continuation of the handler from capturing the continuation past the exception point.

A *escape continuation* is essentially a derived concept. It combines a prompt for escape purposes with a continuation for mark-gathering purposes. As the name implies, escape continuations are used only to abort to the point of capture.

1.1.13 Threads

See §11 "Concurrency and Parallelism" for thread and synchronization functions.

Racket supports multiple *threads* of evaluation. Threads run concurrently, in the sense that one thread can preempt another without its cooperation, but threads currently all run on the same processor (i.e., the same underlying OS process and thread). See also §11.4 "Futures".

Threads are created explicitly by functions such as `thread`. In terms of the evaluation model, each step in evaluation actually consists of multiple concurrent expressions, up to one per thread, rather than a single expression. The expressions all share the same objects and top-level variables, so that they can communicate through shared state. Most evaluation steps involve a single step in a single expression, but certain synchronization primitives require multiple threads to progress together in one step.

In addition to the state that is shared among all threads, each thread has its own private state that is accessed through *thread cells*. A thread cell is similar to a normal mutable object, but a change to the value inside a thread cell is seen only when extracting a value from the cell from the same thread. A thread cell can be *preserved*; when a new thread is created, the creating thread's value for a preserved thread cell serves as the initial value for the cell in the created thread. For a non-preserved thread cell, a new thread sees the same initial value

(specified when the thread cell is created) as all other threads.

1.1.14 Parameters

See §11.3.2 "Parameters" for parameter forms and functions.

Parameters are essentially a derived concept in Racket; they are defined in terms of continuation marks and thread cells. However, parameters are also built in, in the sense that some primitive procedures consult parameter values. For example, the default output stream for primitive output operations is determined by a parameter.

A parameter is a setting that is both thread-specific and continuation-specific. In the empty continuation, each parameter corresponds to a preserved thread cell; a corresponding *parameter procedure* accesses and sets the thread cell's value for the current thread.

In a non-empty continuation, a parameter's value is determined through a *parameterization* that is associated with the nearest enclosing continuation frame though a continuation mark (whose key is not directly accessible). A parameterization maps each parameter to a preserved thread cell, and the combination of thread cell and current thread yields the parameter's value. A parameter procedure sets or accesses the relevant thread cell for its parameter.

Various operations, such as `parameterize` or `call-with-parameterization`, install a parameterization into the current continuation's frame.

1.1.15 Exceptions

See §10.2 "Exceptions" for exception forms, functions, and types.

Exceptions are essentially a derived concept in Racket; they are defined in terms of continuations, prompts, and continuation marks. However, exceptions are also built in, in the sense that primitive forms and procedures may raise exceptions.

An *exception handler* to catch exceptions can be associated with a continuation frame though a continuation mark (whose key is not directly accessible). When an exception is raised, the current continuation's marks determine a chain of exception handler procedures that are consulted to handle the exception. A handler for uncaught exceptions is designated through a built-in parameter.

One potential action of an exception handler is to abort the current continuation up to an enclosing prompt with a particular prompt tag. The default handler for uncaught exceptions, in particular, aborts to a particular tag for which a prompt is always present, because the prompt is installed in the outermost frame of the continuation for any new thread.

1.1.16 Custodians

See §14.7 "Custodians" for custodian functions.

A *custodian* manages a collection of threads, file-stream ports, TCP ports, TCP listeners, UDP sockets, and byte converters. Whenever a thread, etc., is created, it is placed under the management of the *current custodian* as determined by the `current-custodian` parameter.

Custodians also manage eventspaces from `racket/gui/base`.

Except for the root custodian, every custodian itself is managed by a custodian, so that custodians form a hierarchy. Every object managed by a subordinate custodian is also managed by the custodian's owner.

When a custodian is shut down via `custodian-shutdown-all`, it forcibly and immediately closes the ports, TCP connections, etc., that it manages, as well as terminating (or suspending) its threads. A custodian that has been shut down cannot manage new objects. After the current custodian is shut down, if a procedure is called that attempts to create a managed resource (e.g., `open-input-file`, `thread`), then the `exn:fail:contract` exception is raised.

A thread can have multiple managing custodians, and a suspended thread created with `thread/suspend-to-kill` can have zero custodians. Extra custodians become associated with a thread through `thread-resume` (see §11.1.2 "Suspending, Resuming, and Killing Threads"). When a thread has multiple custodians, it is not necessarily killed by a `custodian-shutdown-all`, but shut-down custodians are removed from the thread's managing set, and the thread is killed when its managing set becomes empty.

The values managed by a custodian are only weakly held by the custodian. As a result, a will can be executed for a value that is managed by a custodian. In addition, a custodian only weakly references its subordinate custodians; if a subordinate custodian is unreferenced but has its own subordinates, then the custodian may be collected, at which point its subordinates become immediately subordinate to the collected custodian's superordinate custodian.

In addition to the other entities managed by a custodian, a *custodian box* created with `make-custodian-box` strongly holds onto a value placed in the box until the box's custodian is shut down. The custodian only weakly retains the box itself, however (so the box and its content can be collected if there are no other references to them).

When Racket is compiled with support for per-custodian memory accounting (see `custodian-memory-accounting-available?`), the `current-memory-use` procedure can report a custodian-specific result. This result determines how much memory is occupied by objects that are reachable from the custodian's managed values, especially its threads, and including its sub-custodians' managed values. If an object is reachable from two custodians where neither is an ancestor of the other, an object is arbitrarily charged to one or the other, and the choice can change after each collection; objects reachable from both a custodian and its descendant, however, are reliably charged to the custodian and not to the descendants, unless the custodian can reach the objects only through a descendant custodian or a descendant's thread. Reachability for per-custodian accounting does not include weak references, references to threads managed by other custodians, references to other custodians,

or references to custodian boxes for other custodians.

1.2 Syntax Model

The syntax of a Racket program is defined by

- a *read* pass that processes a character stream into a syntax object; and
- an *expand* pass that processes a syntax object to produce one that is fully parsed.

For details on the read pass, see §1.3 “The Reader”. Source code is normally read in `read-syntax` mode, which produces a syntax object.

The expand pass recursively processes a syntax object to produce a complete parse of the program. Binding information in a syntax object drives the expansion process, and when the expansion process encounters a binding form, it extends syntax objects for sub-expression with new binding information.

1.2.1 Identifiers and Binding

§4.2 “Identifiers and Binding” in ***The Racket Guide*** introduces binding.

An *identifier* is a source-program entity. Parsing (i.e., expanding) a Racket program reveals that some identifiers correspond to variables, some refer to syntactic forms, and some are quoted to produce a symbol or a syntax object.

An identifier *binds* another (i.e., it is a *binding*) when the former is parsed as a variable and the latter is parsed as a reference to the former; the latter is *bound*. The *scope* of a binding is the set of source forms to which it applies. The *environment* of a form is the set of bindings whose scope includes the form. A binding for a sub-expression *shadows* any bindings (i.e., it is *shadowing*) in its environment, so that uses of an identifier refer to the shadowing binding.

For example, as a bit of source, the text

```
(let ([x 5]) x)
```

includes two identifiers: `let` and `x` (which appears twice). When this source is parsed in a typical environment, `x` turns out to represent a variable (unlike `let`). In particular, the first `x` binds the second `x`.

A *top-level binding* is a binding from a definition at the top-level; a *module binding* is a binding from a definition in a module; all other bindings are *local bindings*. There is no difference between an *unbound* identifier and one with a top-level binding; within a module, references to top-level bindings are disallowed, and so such identifiers are called unbound in a module context.

Throughout the documentation, identifiers are typeset to suggest the way that they are parsed. A black, boldface identifier like `lambda` indicates a reference to a syntactic form. A plain blue identifier like `x` is a variable or a reference to an unspecified top-level variable. A hyperlinked identifier `cons` is a reference to a specific top-level variable.

Every binding has a *phase level* in which it can be referenced, where a phase level normally corresponds to an integer (but the special label phase level does not correspond to an integer). Phase level 0 corresponds to the run time of the enclosing module (or the run time of top-level expressions). Bindings in phase level 0 constitute the *base environment*. Phase level 1 corresponds to the time during which the enclosing module (or top-level expression) is expanded; bindings in phase level 1 constitute the *transformer environment*. Phase level -1 corresponds to the run time of a different module for which the enclosing module is imported for use at phase level 1 (relative to the importing module); bindings in phase level -1 constitute the *template environment*. The *label phase level* does not correspond to any execution time; it is used to track bindings (e.g., to identifiers within documentation) without implying an execution dependency.

If an identifier has a local binding, then it is the same for all phase levels, though the reference is allowed only at a particular phase level. Attempting to reference a local binding in a different phase level from the binding's context produces a syntax error. If an identifier has a top-level binding or module binding, then it can have different such bindings in different phase levels.

1.2.2 Syntax Objects

A *syntax object* combines a simpler Racket value, such as a symbol or pair, with *lexical information* about bindings, source-location information, syntax properties, and tamper status. In particular, an identifier is represented as a symbol object that combines a symbol with lexical and other information.

For example, a `car` identifier might have lexical information that designates it as the `car` from the `racket/base` language (i.e., the built-in `car`). Similarly, a `lambda` identifier's lexical information may indicate that it represents a procedure form. Some other identifier's lexical information may indicate that it references a top-level variable.

When a syntax object represents a more complex expression than an identifier or simple constant, its internal components can be extracted. Even for extracted identifiers, detailed information about binding is available mostly indirectly; two identifiers can be compared to determine whether they refer to the same binding (i.e., `free-identifier=?`), or whether each identifier would bind the other if one were in a binding position and the other in an expression position (i.e., `bound-identifier=?`).

For example, when the program written as

```
(let ([x 5]) (+ x 6))
```

is represented as a syntax object, then two syntax objects can be extracted for the two `x`s. Both the `free-identifier=?` and `bound-identifier=?` predicates will indicate that the `x`s are the same. In contrast, the `let` identifier is not `free-identifier=?` or `bound-identifier=?` to either `x`.

The lexical information in a syntax object is independent of the other half, and it can be copied to a new syntax object in combination with an arbitrary other Racket value. Thus, identifier-binding information in a syntax object is predicated on the symbolic name of the identifier as well as the identifier's lexical information; the same question with the same lexical information but different base value can produce a different answer.

For example, combining the lexical information from `let` in the program above to `'x` would not produce an identifier that is `free-identifier=?` to either `x`, since it does not appear in the scope of the `x` binding. Combining the lexical context of the `6` with `'x`, in contrast, would produce an identifier that is `bound-identifier=?` to both `x`s.

The `quote-syntax` form bridges the evaluation of a program and the representation of a program. Specifically, (`quote-syntax` *`datum`*) produces a syntax object that preserves all of the lexical information that *`datum`* had when it was parsed as part of the `quote-syntax` form.

1.2.3 Expansion (Parsing)

Expansion recursively processes a syntax object in a particular phase level, starting with phase level 0. Bindings from the syntax object's lexical information drive the expansion process, and cause new bindings to be introduced for the lexical information of sub-expressions. In some cases, a sub-expression is expanded in a deeper phase than the enclosing expression.

Fully Expanded Programs

A complete expansion produces a syntax object matching the following grammar:

Beware that the symbolic names of identifiers in a fully expanded program may not match the symbolic names in the grammar. Only the binding (according to `free-identifier=?`) matters.

```
top-level-form = general-top-level-form
               | (#%expression expr)
               | (module id module-path
                   (#%plain-module-begin
                    module-level-form ...))
               | (begin top-level-form ...)
               | (begin-for-syntax top-level-form ...)

module-level-form = general-top-level-form
               | (#%provide raw-provide-spec ...)
               | (begin-for-syntax module-level-form ...)
               | submodule-form
               | (#%declare declaration-keyword ...)
```

```
submodule-form = (module id module-path
                   (#%plain-module-begin
                    module-level-form ...))
               | (module* id module-path
                   (#%plain-module-begin
                    module-level-form ...))
               | (module* id #f
                   (#%plain-module-begin
                    module-level-form ...))

general-top-level-form = expr
                       | (define-values (id ...) expr)
                       | (define-syntaxes (id ...) expr)
                       | (#%require raw-require-spec ...)

          expr = id
               | (#%plain-lambda formals expr ...+)
               | (case-lambda (formals expr ...+) ...)
               | (if expr expr expr)
               | (begin expr ...+)
               | (begin0 expr expr ...)
               | (let-values ([(id ...) expr] ...)
                   expr ...+)
               | (letrec-values ([(id ...) expr] ...)
                   expr ...+)
               | (set! id expr)
               | (quote datum)
               | (quote-syntax datum)
               | (with-continuation-mark expr expr expr)
               | (#%plain-app expr ...+)
               | (#%top . id)
               | (#%variable-reference id)
               | (#%variable-reference (#%top . id))
               | (#%variable-reference)

       formals = (id ...)
               | (id ...+ . id)
               | id
```

A *fully-expanded* syntax object corresponds to a *parse* of a program (i.e., a *parsed* program), and lexical information on its identifiers indicates the parse.

More specifically, the typesetting of identifiers in the above grammar is significant. For example, the second case for *expr* is a syntax-object list whose first element is an identifier, where the identifier's lexical information specifies a binding to the `#%plain-lambda` of the

racket/base language (i.e., the identifier is free-identifier=? to one whose binding is #%plain-lambda). In all cases, identifiers above typeset as syntactic-form names refer to the bindings defined in §3 "Syntactic Forms".

In a fully expanded program for a namespace whose base phase is 0, the relevant phase level for a binding in the program is *N* if the bindings has *N* surrounding begin-for-syntax and define-syntaxes forms—not counting any begin-for-syntax forms that wrap a module form for the body of the module. The *datum* in a quote-syntax form, however, always preserves its information for all phase levels.

In addition to the grammar above, letrec-syntaxes+values can appear in a fully local-expanded expression, as can #%expression in any expression position. For example, letrec-syntaxes+values and #%expression can appear in the result from local-expand when the stop list is empty.

Expansion Steps

In a recursive expansion, each single step in expanding a syntax object at a particular phase level depends on the immediate shape of the syntax object being expanded:

- If it is an identifier (i.e., a syntax-object symbol), then a binding is determined by the identifier's lexical information. If the identifier has a binding other than as a top-level variable, that binding is used to continue. If the identifier has no binding, a new syntax-object symbol '#%top is created using the lexical information of the identifier; if this #%top identifier has no binding (other than as a top-level variable), then parsing fails with an exn:fail:syntax exception. Otherwise, the new identifier is combined with the original identifier in a new syntax-object pair (also using the same lexical information as the original identifier), and the #%top binding is used to continue.

- If it is a syntax-object pair whose first element is an identifier, and if the identifier has a binding other than as a top-level variable, then the identifier's binding is used to continue.

- If it is a syntax-object pair of any other form, then a new syntax-object symbol '#%app is created using the lexical information of the pair. If the resulting #%app identifier has no binding, parsing fails with an exn:fail:syntax exception. Otherwise, the new identifier is combined with the original pair to form a new syntax-object pair (also using the same lexical information as the original pair), and the #%app binding is used to continue.

- If it is any other syntax object, then a new syntax-object symbol '#%datum is created using the lexical information of the original syntax object. If the resulting #%datum identifier has no binding, parsing fails with an exn:fail:syntax exception. Otherwise, the new identifier is combined with the original syntax object in a new syntax-object pair (using the same lexical information as the original pair), and the #%datum binding is used to continue.

Thus, the possibilities that do not fail lead to an identifier with a particular binding. This binding refers to one of three things:

- A transformer binding, such as introduced by `define-syntax` or `let-syntax`. If the associated value is a procedure of one argument, the procedure is called as a syntax transformer (described below), and parsing starts again with the syntax-object result. If the transformer binding is to any other kind of value, parsing fails with an `exn:fail:syntax` exception. The call to the syntax transformer is `parameterized` to set `current-namespace` to a namespace that shares bindings and variables with the namespace being used to expand, except that its base phase is one greater.
- A variable binding, such as introduced by a module-level `define` or by `let`. In this case, if the form being parsed is just an identifier, then it is parsed as a reference to the corresponding variable. If the form being parsed is a syntax-object pair, then an `#%app` is added to the front of the syntax-object pair in the same way as when the first item in the syntax-object pair is not an identifier (third case in the previous enumeration), and parsing continues.
- A core *syntactic form*, which is parsed as described for each form in §3 "Syntactic Forms". Parsing a core syntactic form typically involves recursive parsing of sub-forms, and may introduce bindings that determine the parsing of sub-forms.

Expansion Context

Each expansion step occurs in a particular *context*, and transformers and core syntactic forms may expand differently for different contexts. For example, a `module` form is allowed only in a top-level context, and it fails in other contexts. The possible contexts are as follows:

- *top-level context* : outside of any module, definition, or expression, except that sub-expressions of a top-level `begin` form are also expanded as top-level forms.
- *module-begin context* : inside the body of a module, as the only form within the module.
- *module context* : in the body of a module (inside the module-begin layer).
- *internal-definition context* : in a nested context that allows both definitions and expressions.
- *expression context* : in a context where only expressions are allowed.

Different core syntactic forms parse sub-forms using different contexts. For example, a `let` form always parses the right-hand expressions of a binding in an expression context, but it starts parsing the body in an internal-definition context.

Introducing Bindings

Bindings are introduced during expansion when certain core syntactic forms are encountered:

- When a `require` form is encountered at the top level or module level, all lexical information derived from the top level or the specific module's level is extended with bindings from the specified modules. If not otherwise indicated in the `require` form, bindings are introduced at the phase levels specified by the exporting modules: phase level 0 for each normal `provide`, phase level 1 for each `for-syntax provide`, and so on. The `for-meta provide` form allows exports at an arbitrary phase level (as long as a binding exists within the module at the phase level).

  A `for-syntax` sub-form within `require` imports similarly, but the resulting bindings have a phase level that is one more than the exported phase levels, when exports for the label phase level are still imported at the label phase level. More generally, a `for-meta` sub-form within `require` imports with the specified phase level shift; if the specified shift is `#f`, or if `for-label` is used to import, then all bindings are imported into the label phase level.

- When a `define`, `define-values`, `define-syntax`, or `define-syntaxes` form is encountered at the top level or module level, all lexical information derived from the top level or the specific module's level is extended with bindings for the specified identifiers at phase level 0 (i.e., the base environment is extended).

- When a `begin-for-syntax` form is encountered at the top level or module level, bindings are introduced as for `define-values` and `define-syntaxes`, but at phase level 1 (i.e., the transformer environment is extended). More generally, `begin-for-syntax` forms can be nested, an each `begin-for-syntax` shifts its body definition by one phase level.

- When a `let-values` form is encountered, the body of the `let-values` form is extended (by creating new syntax objects) with bindings for the specified identifiers. The same bindings are added to the identifiers themselves, so that the identifiers in binding position are `bound-identifier=?` to uses in the fully expanded form, and so they are not `bound-identifier=?` to other identifiers. The bindings are available for use at the phase level at which the `let-values` form is expanded.

- When a `letrec-values` or `letrec-syntaxes+values` form is encountered, bindings are added as for `let-values`, except that the right-hand-side expressions are also extended with the bindings.

- Definitions in internal-definition contexts introduce bindings as described in §1.2.3.7 "Internal Definitions".

A new binding in lexical information maps to a new variable. The identifiers mapped to this variable are those that currently have the same binding (i.e., that are currently `bound-identifier=?`) to the identifier associated with the binding.

For example, in

```
(let-values ([(x) 10]) (+ x y))
```

the binding introduced for `x` applies to the `x` in the body, but not the `y` in the body, because (at the point in expansion where the `let-values` form is encountered) the binding `x` and the body `y` are not `bound-identifier=?`.

Transformer Bindings

In a top-level context or module context, when the expander encounters a `define-syntaxes` form, the binding that it introduces for the defined identifiers is a *transformer binding*. The value of the binding exists at expansion time, rather than run time (though the two times can overlap), though the binding itself is introduced with phase level 0 (i.e., in the base environment).

The value for the binding is obtained by evaluating the expression in the `define-syntaxes` form. This expression must be expanded (i.e., parsed) before it can be evaluated, and it is expanded at phase level 1 (i.e., in the transformer environment) instead of phase level 0.

If the resulting `value` is a procedure of one argument or the result of `make-set!-transformer` on a procedure, then it is used as a *syntax transformer* (a.k.a. *macro*). The procedure is expected to accept a syntax object and return a syntax object. A use of the binding (at phase level 0) triggers a call of the syntax transformer by the expander; see §1.2.3.2 “Expansion Steps”.

Before the expander passes a syntax object to a transformer, the syntax object is extended with a *syntax mark* (that applies to all sub-syntax objects). The result of the transformer is similarly extended with the same syntax mark. When a syntax object’s lexical information includes the same mark twice in a row, the marks effectively cancel. Otherwise, two identifiers are `bound-identifier=?` (that is, one can bind the other) only if they have the same binding and if they have the same marks—counting only marks that were added after the binding.

This marking process helps keep binding in an expanded program consistent with the lexical structure of the source program. For example, the expanded form of the program

```
(define x 12)
(define-syntax m
  (syntax-rules ()
    [(_ id) (let ([x 10]) id)]))
(m x)
```

is

```
(define x 12)
(define-syntax m
  (syntax-rules ()
```

```
    [(_ id) (let ([x 10]) id)]))
(let-values ([(x) 10]) x)
```

However, the result of the last expression is `12`, not `10`. The reason is that the transformer bound to `m` introduces the binding `x`, but the referencing `x` is present in the argument to the transformer. The introduced `x` is the one left with a mark, and the reference `x` has no mark, so the binding `x` is not `bound-identifier=?` to the body `x`.

The `set!` form works with the `make-set!-transformer` and `prop:set!-transformer` property to support *assignment transformers* that transform `set!` expressions. An assignment transformer contains a procedure that is applied by `set!` in the same way as a normal transformer by the expander.

The `make-rename-transformer` procedure or `prop:rename-transformer` property creates a value that is also handled specially by the expander and by `set!` as a transformer binding's value. When *id* is bound to a *rename transformer* produced by `make-rename-transformer`, it is replaced with the target identifier passed to `make-rename-transformer`. In addition, as long as the target identifier does not have a true value for the `'not-free-identifier=?` syntax property, the lexical information that contains the binding of *id* is also enriched so that *id* is `free-identifier=?` to the target identifier, `identifier-binding` returns the same results for both identifiers, and `provide` exports *id* as the target identifier. Finally, the binding is treated specially by `syntax-local-value`, and `syntax-local-make-delta-introducer` as used by syntax transformers.

In addition to using marks to track introduced identifiers, the expander tracks the expansion history of a form through syntax properties such as `'origin`. See §12.7 "Syntax Object Properties" for more information.

Finally, the expander uses a tamper status to control the way that unexported and protected module bindings are used. See §12.8 "Syntax Taints" for more information on a tamper status.

The expander's handling of `letrec-syntaxes+values` is similar to its handling of `define-syntaxes`. A `letrec-syntaxes+values` can be expanded in an arbitrary phase level *n* (not just 0), in which case the expression for the transformer binding is expanded at phase level *n*+1.

The expressions in a `begin-for-syntax` form are expanded and evaluated in the same way as for `define-syntaxes`. However, any introduced bindings from definition within `begin-for-syntax` are at phase level 1 (not a transformer binding at phase level 0).

Partial Expansion

In certain contexts, such as an internal-definition context or module context, *partial expansion* is used to determine whether forms represent definitions, expressions, or other declaration forms. Partial expansion works by cutting off the normal recursion expansion when the relevant binding is for a primitive syntactic form.

As a special case, when expansion would otherwise add an `#%app`, `#%datum`, or `#%top` identifier to an expression, and when the binding turns out to be the primitive `#%app`, `#%datum`, or `#%top` form, then expansion stops without adding the identifier.

Internal Definitions

An internal-definition context supports local definitions mixed with expressions. Forms that allow internal definitions document such positions using the *body* meta-variable. Definitions in an internal-definition context are equivalent to local binding via `letrec-syntaxes+values`; macro expansion converts internal definitions to a `letrec-syntaxes+values` form.

Expansion of an internal-definition context relies on partial expansion of each *body* in an internal-definition sequence. Partial expansion of each *body* produces a form matching one of the following cases:

- A `define-values` form: The lexical context of all syntax objects for the body sequence is immediately enriched with bindings for the `define-values` form. Further expansion of the definition is deferred, and partial expansion continues with the rest of the body.
- A `define-syntaxes` form: The right-hand side is expanded and evaluated (as for a `letrec-syntaxes+values` form), and a transformer binding is installed for the body sequence before partial expansion continues with the est of the body.
- A primitive expression form other than `begin`: Further expansion of the expression is deferred, and partial expansion continues with the rest of the body.
- A `begin` form: The sub-forms of the `begin` are spliced into the internal-definition sequence, and partial expansion continues with the first of the newly-spliced forms (or the next form, if the `begin` had no sub-forms).

After all body forms are partially expanded, if no definitions were encountered, then the expressions are collected into a `begin` form as the internal-definition context's expansion. Otherwise, at least one expression must appear after the last definition, and any *expr* that appears between definitions is converted to `(define-values () (begin` *expr* `(values)))`; the definitions are then converted to bindings in a `letrec-syntaxes+values` form, and all expressions after the last definition become the body of the `letrec-syntaxes+values` form.

Module Phases and Visits

A `require` form not only introduces bindings at expansion time, but also *visits* the referenced module when it is encountered by the expander. That is, the expander instantiates any variables defined in the module within `begin-for-syntax`, and it also evaluates all expressions for `define-syntaxes` transformer bindings.

Module visits propagate through `requires` in the same way as module instantiation. Moreover, when a module is visited at phase 0, any module that it `requires` `for-syntax` is instantiated at phase 1, while further `requires` `for-template` leading back to phase 0 causes the required module to be visited at phase 0 (i.e., not instantiated).

During compilation, the top-level of module context is itself implicitly visited. Thus, when the expander encounters `(require (for-syntax ....))`, it immediately instantiates the required module at phase 1, in addition to adding bindings at phase level 1 (i.e., the transformer environment). Similarly, the expander immediately evaluates any form that it encounters within `begin-for-syntax`.

Phases beyond 0 are visited on demand. For example, when the right-hand side of a phase-0 `let-syntax` is to be expanded, then modules that are available at phase 1 are visited. More generally, initiating expansion at phase n visits modules at phase n, which in turn instantiates modules at phase $n+1$. These visits and instantiations apply to available modules in the enclosing namespace's module registry; a per-registry lock prevents multiple threads from concurrently instantiating and visiting available modules.

When the expander encounters `require` and `(require (for-syntax ....))` within a module context, the resulting visits and instantiations are specific to the expansion of the enclosing module, and are kept separate from visits and instantiations triggered from a top-level context or from the expansion of a different module. Along the same lines, when a module is attached to a namespace through `namespace-attach-module`, modules that it `requires` are transitively attached, but instances are attached only at phases at or below the namespace's base phase.

Macro-Introduced Bindings

When a top-level definition binds an identifier that originates from a macro expansion, the definition captures only uses of the identifier that are generated by the same expansion. This behavior is consistent with expansion in internal-definition contexts, where the defined identifier turns into a fresh lexical binding.

Examples:

```
> (define-syntax def-and-use-of-x
    (syntax-rules ()
      [(def-and-use-of-x val)
       ; x below originates from this macro:
       (begin (define x val) x)]))

> (define x 1)

> x
1
> (def-and-use-of-x 2)
2
```

```
> x
1
> (define-syntax def-and-use
    (syntax-rules ()
      [(def-and-use x val)
       ; "x" below was provided by the macro use:
       (begin (define x val) x)]))

> (def-and-use x 3)
3
> x
3
```

For a top-level definition (outside of a module), the order of evaluation affects the binding of a generated definition for a generated identifier use. If the use precedes the definition, then the use refers to a non-generated binding, just as if the generated definition were not present. (No such dependency on order occurs within a module, since a module binding covers the entire module body.) To support the declaration of an identifier before its use, the `define-syntaxes` form avoids binding an identifier if the body of the `define-syntaxes` declaration produces zero results.

Examples:

```
> (define bucket-1 0)

> (define bucket-2 0)

> (define-syntax def-and-set!-use-of-x
    (syntax-rules ()
      [(def-and-set!-use-of-x val)
       (begin (set! bucket-1 x) (define x val) (set! bucket-
2 x))]))

> (define x 1)

> (def-and-set!-use-of-x 2)

> x
1
> bucket-1
1
> bucket-2
2
> (define-syntax defs-and-uses/fail
    (syntax-rules ()
      [(def-and-use)
       (begin
```

```
          ; Initial reference to even precedes definition:
          (define (odd x) (if (zero? x) #f (even (sub1 x))))
          (define (even x) (if (zero? x) #t (odd (sub1 x))))
          (odd 17))]))

> (defs-and-uses/fail)
even: undefined;
 cannot reference undefined identifier
> (define-syntax defs-and-uses
    (syntax-rules ()
      [(def-and-use)
       (begin
         ; Declare before definition via no-values define-
syntaxes:
          (define-syntaxes (odd even) (values))
          (define (odd x) (if (zero? x) #f (even (sub1 x))))
          (define (even x) (if (zero? x) #t (odd (sub1 x))))
          (odd 17))]))

> (defs-and-uses)
#t
```

Macro-generated `"require"` and `"provide"` clauses also introduce and reference generation-specific bindings:

- In `require`, for a *`require-spec`* of the form `(rename-in [`*`orig-id bind-id`*`])` or `(only-in .... [`*`orig-id bind-id`*`])`, the *`bind-id`* is bound only for uses of the identifier generated by the same macro expansion as *`bind-id`*. In `require` for other *`require-spec`*s, the generator of the *`require-spec`* determines the scope of the bindings.

- In `provide`, for a *`provide-spec`* of the form *`id`*, the exported identifier is the one that binds *`id`* within the module in a generator-specific way, but the external name is the plain *`id`*. The exceptions for `all-except-out` are similarly determined in a generator-specific way, as is the *`orig-id`* binding of a `rename-out` form, but plain identifiers are used for the external names. For `all-defined-out`, only identifiers with definitions having the same generator as the `(all-defined-out)` form are exported; the external name is the plain identifier from the definition.

1.2.4 Compilation

Before expanded code is evaluated, it is first *compiled*. A compiled form has essentially the same information as the corresponding expanded form, though the internal representation naturally dispenses with identifiers for syntactic forms and local bindings. One significant

difference is that a compiled form is almost entirely opaque, so the information that it contains cannot be accessed directly (which is why some identifiers can be dropped). At the same time, a compiled form can be marshaled to and from a byte string, so it is suitable for saving and re-loading code.

Although individual read, expand, compile, and evaluate operations are available, the operations are often combined automatically. For example, the `eval` procedure takes a syntax object and expands it, compiles it, and evaluates it.

1.2.5 Namespaces

See §14.1 "Namespaces" for functions that manipulate namespaces.

A *namespace* is a top-level mapping from symbols to binding information. It is the starting point for expanding an expression; a syntax object produced by `read-syntax` has no initial lexical context; the syntax object can be expanded after initializing it with the mappings of a particular namespace. A namespace is also the starting point evaluating expanded code, where the first step in evaluation is linking the code to specific module instances and top-level variables.

For expansion purposes, a namespace maps each symbol in each phase level to one of three possible bindings:

- a particular module binding from a particular module
- a top-level transformer binding named by the symbol
- a top-level variable named by the symbol

An "empty" namespace maps all symbols to top-level variables. Certain evaluations extend a namespace for future expansions; importing a module into the top-level adjusts the namespace bindings for all of the imported names, and evaluating a top-level `define` form updates the namespace's mapping to refer to a variable (in addition to installing a value into the variable).

A namespace also has a *module registry* that maps module names to module declarations (see §1.1.10 "Modules and Module-Level Variables"). This registry is shared by all phase levels.

For evaluation, each namespace encapsulates a distinct set of top-level variables at various phases, as well as a potentially distinct set of module instances in each phase. That is, even though module declarations are shared for all phase levels, module instances are distinct for each phase. Each namespace has a *base phase*, which corresponds to the phase used by reflective operations such as `eval` and `dynamic-require`. In particular, using `eval` on a `require` form instantiates a module in the namespace's base phase.

After a namespace is created, module instances from existing namespaces can be attached

to the new namespace. In terms of the evaluation model, top-level variables from different namespaces essentially correspond to definitions with different prefixes, but attaching a module uses the same prefix for the module's definitions in namespaces where it is attached. The first step in evaluating any compiled expression is to link its top-level variable and module-level variable references to specific variables in the namespace.

At all times during evaluation, some namespace is designated as the *current namespace*. The current namespace has no particular relationship, however, with the namespace that was used to expand the code that is executing, or with the namespace that was used to link the compiled form of the currently evaluating code. In particular, changing the current namespace during evaluation does not change the variables to which executing expressions refer. The current namespace only determines the behavior of reflective operations to expand code and to start evaluating expanded/compiled code.

Examples:

```
> (define x 'orig) ; define in the original namespace

; The following let expression is compiled in the original
; namespace, so direct references to x see 'orig.
> (let ([n (make-base-namespace)]) ; make new namespace
    (parameterize ([current-namespace n])
      (eval '(define x 'new)) ; evals in the new namespace
      (display x) ; displays 'orig
      (display (eval 'x)))) ; displays 'new
origneworig
```

A namespace is purely a top-level entity, not to be confused with an environment. In particular, a namespace does not encapsulate the full environment of an expression inside local-binding forms.

If an identifier is bound to syntax or to an import, then defining the identifier as a variable shadows the syntax or import in future uses of the environment. Similarly, if an identifier is bound to a top-level variable, then binding the identifier to syntax or an import shadows the variable; the variable's value remains unchanged, however, and may be accessible through previously evaluated expressions.

Examples:

```
> (define x 5)

> (define (f) x)

> x
5
> (f)
```

```
5
> (define-syntax x (syntax-id-rules () [_ 10]))

> x
10
> (f)
5
> (define x 7)

> x
7
> (f)
7
> (module m racket (define x 8) (provide x))

> (require 'm)

> x
8
> (f)
7
```

1.2.6 Inferred Value Names

To improve error reporting, names are inferred at compile-time for certain kinds of values, such as procedures. For example, evaluating the following expression:

```
(let ([f (lambda () 0)]) (f 1 2 3))
```

produces an error message because too many arguments are provided to the procedure. The error message is able to report `f` as the name of the procedure. In this case, Racket decides, at compile-time, to name as `'f` all procedures created by the `let`-bound `lambda`.

See `procedure-rename` to override a procedure's inferred name at runtime.

Names are inferred whenever possible for procedures. Names closer to an expression take precedence. For example, in

```
(define my-f
  (let ([f (lambda () 0)]) f))
```

the procedure bound to `my-f` will have the inferred name `'f`.

When an `'inferred-name` property is attached to a syntax object for an expression (see §12.7 "Syntax Object Properties"), the property value is used for naming the expression, and it overrides any name that was inferred from the expression's context. Normally, the

property value should be a symbol. A `'inferred-name` property value of `#<void>` hides a name that would otherwise be inferred from context (perhaps because a binding identifier's was automatically generated and should not be exposed).

When an inferred name is not available, but a source location is available, a name is constructed using the source location information. Inferred and property-assigned names are also available to syntax transformers, via `syntax-local-name`.

1.2.7 Cross-Phase Persistent Module Declarations

A module is cross-phase persistent only if it fits the following grammar, which uses non-terminals from §1.2.3.1 "Fully Expanded Programs", only if it includes `(#%declare #:cross-phase-persistent)`, only it includes no uses of `quote-syntax` or `#%variable-reference`, and only if no module-level binding is `set!`ed.

```
cross-module = (module id module-path
                 (#%plain-module-begin
                   cross-form ...))

 cross-form = (#%declare #:cross-phase-persistent)
            | (begin cross-form ...)
            | (#%provide raw-provide-spec ...)
            | submodule-form
            | (define-values (id ...) cross-expr)
            | (#%require raw-require-spec ...)

 cross-expr = id
            | (quote cross-datum)
            | (#%plain-lambda formals expr ...+)
            | (case-lambda (formals expr ...+) ...)
            | (#%plain-app cons cross-expr ...+)
            | (#%plain-app list cross-expr ...+)
            | (#%plain-app make-struct-type cross-expr ...+)
            | (#%plain-app make-struct-type-property
                           cross-expr ...+)
            | (#%plain-app gensym)
            | (#%plain-app gensym string)
            | (#%plain-app string->uninterned-symbol string)

cross-datum = number
            | boolean
            | identifier
            | string
            | bytes
```

This grammar applies after expansion, but because a cross-phase persistent module imports only from other cross-phase persistent modules, the only relevant expansion steps are the implicit introduction of `#%plain-module-begin`, implicit introduction of `#%plain-app`, and implicit introduction and/or expansion of `#%datum`.

1.3 The Reader

Racket's reader is a recursive-descent parser that can be configured through a readtable and various other parameters. This section describes the reader's parsing when using the default readtable.

Reading from a stream produces one *datum*. If the result datum is a compound value, then reading the datum typically requires the reader to call itself recursively to read the component data.

The reader can be invoked in either of two modes: `read` mode, or `read-syntax` mode. In `read-syntax` mode, the result is always a syntax object that includes source-location and (initially empty) lexical information wrapped around the sort of datum that `read` mode would produce. In the case of pairs, vectors, and boxes, the content is also wrapped recursively as a syntax object. Unless specified otherwise, this section describes the reader's behavior in `read` mode, and `read-syntax` mode does the same modulo wrapping of the final result.

Reading is defined in terms of Unicode characters; see §13.1 "Ports" for information on how a byte stream is converted to a character stream.

Symbols, keywords, strings, byte strings, regexps, characters, and numbers produced by the reader in `read-syntax` mode are *interned*, which means that such values in the result of `read-syntax` are always `eq?` when they are `equal?` (whether from the same call or different calls to `read-syntax`). Symbols and keywords are interned in both `read` and `read-syntax` mode. Sending an interned value across a place channel does not necessarily produce an interned value at the receiving place. See also `datum-intern-literal` and `datum->syntax`.

1.3.1 Delimiters and Dispatch

Along with whitespace, the following characters are *delimiters*:

```
( ) [ ] { } " , ' ` ;
```

A delimited sequence that starts with any other character is typically parsed as either a symbol, number, or extflonum, but a few non-delimiter characters play special roles:

- `#` has a special meaning as an initial character in a delimited sequence; its meaning depends on the characters that follow; see below.
- `|` starts a subsequence of characters to be included verbatim in the delimited sequence (i.e., they are never treated as delimiters, and they are not case-folded when case-insensitivity is enabled); the subsequence is terminated by another `|`, and neither the initial nor terminating `|` is part of the subsequence.
- `\` outside of a `|` pair causes the following character to be included verbatim in a delimited sequence.

More precisely, after skipping whitespace, the reader dispatches based on the next character or characters in the input stream as follows:

`(`	starts a pair or list; see §1.3.6 “Reading Pairs and Lists”
`[`	starts a pair or list; see §1.3.6 “Reading Pairs and Lists”
`{`	starts a pair or list; see §1.3.6 “Reading Pairs and Lists”
`)`	matches `(` or raises `exn:fail:read`
`]`	matches `[` or raises `exn:fail:read`
`}`	matches `{` or raises `exn:fail:read`
`"`	starts a string; see §1.3.7 “Reading Strings”
`'`	starts a quote; see §1.3.8 “Reading Quotes”
`` ` ``	starts a quasiquote; see §1.3.8 “Reading Quotes”
`,`	starts a [splicing] unquote; see §1.3.8 “Reading Quotes”
`;`	starts a line comment; see §1.3.9 “Reading Comments”
`#t` or `#T`	true; see §1.3.5 “Reading Booleans”
`#f` or `#F`	false; see §1.3.5 “Reading Booleans”
`#(`	starts a vector; see §1.3.10 “Reading Vectors”
`#[`	starts a vector; see §1.3.10 “Reading Vectors”
`#{`	starts a vector; see §1.3.10 “Reading Vectors”
`#fl(`	starts a flvector; see §1.3.10 “Reading Vectors”
`#fl[`	starts a flvector; see §1.3.10 “Reading Vectors”
`#fl{`	starts a flvector; see §1.3.10 “Reading Vectors”
`#fx(`	starts a fxvector; see §1.3.10 “Reading Vectors”
`#fx[`	starts a fxvector; see §1.3.10 “Reading Vectors”
`#fx{`	starts a fxvector; see §1.3.10 “Reading Vectors”
`#s(`	starts a structure literal; see §1.3.11 “Reading Structures”
`#s[`	starts a structure literal; see §1.3.11 “Reading Structures”
`#s{`	starts a structure literal; see §1.3.11 “Reading Structures”
`#\`	starts a character; see §1.3.14 “Reading Characters”
`#"`	starts a byte string; see §1.3.7 “Reading Strings”
`#%`	starts a symbol; see §1.3.2 “Reading Symbols”
`#:`	starts a keyword; see §1.3.15 “Reading Keywords”
`#&`	starts a box; see §1.3.13 “Reading Boxes”
`#|`	starts a block comment; see §1.3.9 “Reading Comments”
`#;`	starts an S-expression comment; see §1.3.9 “Reading Comments”

`#'`	starts a syntax quote; see §1.3.8 “Reading Quotes”
`#!`	starts a line comment; see §1.3.9 “Reading Comments”
`#!/`	starts a line comment; see §1.3.9 “Reading Comments”
`#!`	may start a reader extension; see §1.3.18 “Reading via an Extension”
`#``	starts a syntax quasiquote; see §1.3.8 “Reading Quotes”
`#,`	starts a syntax [splicing] unquote; see §1.3.8 “Reading Quotes”
`#~`	starts compiled code; see §1.4.16 “Printing Compiled Code”
`#i` or `#I`	starts a number; see §1.3.3 “Reading Numbers”
`#e` or `#E`	starts a number; see §1.3.3 “Reading Numbers”
`#x` or `#X`	starts a number or extflonum; see §1.3.3 “Reading Numbers”
`#o` or `#O`	starts a number or extflonum; see §1.3.3 “Reading Numbers”
`#d` or `#D`	starts a number or extflonum; see §1.3.3 “Reading Numbers”
`#b` or `#B`	starts a number or extflonum; see §1.3.3 “Reading Numbers”
`#<<`	starts a string; see §1.3.7 “Reading Strings”
`#rx`	starts a regular expression; see §1.3.16 “Reading Regular Expressions”
`#px`	starts a regular expression; see §1.3.16 “Reading Regular Expressions”
`#ci`, `#cI`, `#Ci`, or `#CI`	switches case sensitivity; see §1.3.2 “Reading Symbols”
`#cs`, `#cS`, `#Cs`, or `#CS`	switches case sensitivity; see §1.3.2 “Reading Symbols”
`#hash`	starts a hash table; see §1.3.12 “Reading Hash Tables”
`#reader`	starts a reader extension use; see §1.3.18 “Reading via an Extension”
`#lang`	starts a reader extension use; see §1.3.18 “Reading via an Extension”
`#`⟨*digit*$_{10}$⟩$^+$`(`	starts a vector; see §1.3.10 “Reading Vectors”
`#`⟨*digit*$_{10}$⟩$^+$`[`	starts a vector; see §1.3.10 “Reading Vectors”
`#`⟨*digit*$_{10}$⟩$^+$`{`	starts a vector; see §1.3.10 “Reading Vectors”
`#fl`⟨*digit*$_{10}$⟩$^+$`(`	starts a flvector; see §1.3.10 “Reading Vectors”
`#fl`⟨*digit*$_{10}$⟩$^+$`[`	starts a flvector; see §1.3.10 “Reading Vectors”
`#fl`⟨*digit*$_{10}$⟩$^+$`{`	starts a flvector; see §1.3.10 “Reading Vectors”
`#fx`⟨*digit*$_{10}$⟩$^+$`(`	starts a fxvector; see §1.3.10 “Reading Vectors”
`#fx`⟨*digit*$_{10}$⟩$^+$`[`	starts a fxvector; see §1.3.10 “Reading Vectors”
`#fx`⟨*digit*$_{10}$⟩$^+$`{`	starts a fxvector; see §1.3.10 “Reading Vectors”
`#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`=`	binds a graph tag; see §1.3.17 “Reading Graph Structure”
`#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`#`	uses a graph tag; see §1.3.17 “Reading Graph Structure”
otherwise	starts a symbol; see §1.3.2 “Reading Symbols”

1.3.2 Reading Symbols

§3.6 “Symbols” in ***The Racket Guide*** introduces the syntax of symbols.

A sequence that does not start with a delimiter or `#` is parsed as either a symbol, a number (see §1.3.3 “Reading Numbers”), or a extflonum (see §1.3.4 “Reading Extflonums”), except that `.` by itself is never parsed as a symbol or number (unless the `read-accept-dot` parameter is set to `#f`). A `#%` also starts a symbol. The resulting symbol is interned. A successful number or extflonum parse takes precedence over a symbol parse.

When the `read-case-sensitive` parameter is set to `#f`, characters in the sequence that are not quoted by `|` or `\` are first case-normalized. If the reader encounters `#ci`, `#CI`, `#Ci`,

or `#cI`, then it recursively reads the following datum in case-insensitive mode. If the reader encounters `#cs`, `#CS`, `#Cs`, or `#cS`, then it recursively reads the following datum in case-sensitive mode.

Examples:

`Apple`	reads equal to	`(string->symbol "Apple")`
`Ap#ple`	reads equal to	`(string->symbol "Ap#ple")`
`Ap ple`	reads equal to	`(string->symbol "Ap")`
`Ap| |ple`	reads equal to	`(string->symbol "Ap ple")`
`Ap\ ple`	reads equal to	`(string->symbol "Ap ple")`
`#ci Apple`	reads equal to	`(string->symbol "apple")`
`#ci |A|pple`	reads equal to	`(string->symbol "Apple")`
`#ci \Apple`	reads equal to	`(string->symbol "Apple")`
`#ci#cs Apple`	reads equal to	`(string->symbol "Apple")`
`#%Apple`	reads equal to	`(string->symbol "#%Apple")`

1.3.3 Reading Numbers

§3.2 “Numbers” in ***The Racket Guide*** introduces the syntax of numbers.

A sequence that does not start with a delimiter is parsed as a number when it matches the following grammar case-insensitively for ⟨*number*$_{10}$⟩ (decimal), where *n* is a meta-meta-variable in the grammar. The resulting number is interned in `read-syntax` mode.

A number is optionally prefixed by an exactness specifier, `#e` (exact) or `#i` (inexact), which specifies its parsing as an exact or inexact number; see §4.2 “Numbers” for information on number exactness. As the non-terminal names suggest, a number that has no exactness specifier and matches only ⟨*inexact-number*$_n$⟩ is normally parsed as an inexact number, otherwise it is parsed as an exact number. If the `read-decimal-as-inexact` parameter is set to `#f`, then all numbers without an exactness specifier are instead parsed as exact.

If the reader encounters `#b` (binary), `#o` (octal), `#d` (decimal), or `#x` (hexadecimal), it must be followed by a sequence that is terminated by a delimiter or end-of-file, and that is either an extflonum (see §1.3.4 “Reading Extflonums”) or matches the ⟨*general-number*$_2$⟩, ⟨*general-number*$_8$⟩, ⟨*general-number*$_{10}$⟩, or ⟨*general-number*$_{16}$⟩ grammar, respectively.

A `#e` or `#i` followed immediately by `#b`, `#o`, `#d`, or `#x` is treated the same as the reverse order: `#b`, `#o`, `#d`, or `#x` followed by `#e` or `#i`.

An ⟨*exponent-mark*$_n$⟩ in an inexact number serves both to specify an exponent and to specify a numerical precision. If single-precision IEEE floating point is supported (see §4.2 “Numbers”), the marks `f` and `s` specify single-precision. Otherwise, or with any other mark, double-precision IEEE floating point is used. In addition, single- and double-precision specials are distinct; specials with the `.0` suffix, like `+nan.0` are double-precision, whereas specials with the `.f` suffix are single-precision.

A `#` in an ⟨*inexact*$_n$⟩ number is the same as `0`, but `#` can be used to suggest that the digit’s

actual value is unknown.

⟨*number*$_n$⟩ ::= ⟨*exact*$_n$⟩ | ⟨*inexact*$_n$⟩
⟨*exact*$_n$⟩ ::= ⟨*exact-rational*$_n$⟩ | ⟨*exact-complex*$_n$⟩
⟨*exact-rational*$_n$⟩ ::= [⟨*sign*⟩] ⟨*unsigned-rational*$_n$⟩
⟨*unsigned-rational*$_n$⟩ ::= ⟨*unsigned-integer*$_n$⟩
| ⟨*unsigned-integer*$_n$⟩ `/` ⟨*unsigned-integer*$_n$⟩
⟨*exact-integer*$_n$⟩ ::= [⟨*sign*⟩] ⟨*unsigned-integer*$_n$⟩
⟨*unsigned-integer*$_n$⟩ ::= ⟨*digit*$_n$⟩$^+$
⟨*exact-complex*$_n$⟩ ::= ⟨*exact-rational*$_n$⟩ ⟨*sign*⟩ ⟨*unsigned-rational*$_n$⟩ `i`
⟨*inexact*$_n$⟩ ::= ⟨*inexact-real*$_n$⟩ | ⟨*inexact-complex*$_n$⟩
⟨*inexact-real*$_n$⟩ ::= [⟨*sign*⟩] ⟨*inexact-normal*$_n$⟩
| ⟨*sign*⟩ ⟨*inexact-special*$_n$⟩
⟨*inexact-unsigned*$_n$⟩ ::= ⟨*inexact-normal*$_n$⟩ | ⟨*inexact-special*$_n$⟩
⟨*inexact-normal*$_n$⟩ ::= ⟨*inexact-simple*$_n$⟩ [⟨*exp-mark*$_n$⟩ ⟨*exact-integer*$_n$⟩]
⟨*inexact-simple*$_n$⟩ ::= ⟨*digits#*$_n$⟩ [`.`] `#`*
| [⟨*unsigned-integer*$_n$⟩] `.` ⟨*digits#*$_n$⟩
| ⟨*digits#*$_n$⟩ `/` ⟨*digits#*$_n$⟩
⟨*inexact-special*$_n$⟩ ::= `inf.0` | `nan.0` | `inf.f` | `nan.f`
⟨*digits#*$_n$⟩ ::= ⟨*digit*$_n$⟩$^+$ `#`*
⟨*inexact-complex*$_n$⟩ ::= [⟨*inexact-real*$_n$⟩] ⟨*sign*⟩ ⟨*inexact-unsigned*$_n$⟩ `i`
| ⟨*inexact-real*$_n$⟩ `@` ⟨*inexact-real*$_n$⟩
⟨*sign*⟩ ::= `+` | `-`
⟨*digit*$_{16}$⟩ ::= ⟨*digit*$_{10}$⟩ | `a` | `b` | `c` | `d` | `e` | `f`
⟨*digit*$_{10}$⟩ ::= ⟨*digit*$_8$⟩ | `8` | `9`
⟨*digit*$_8$⟩ ::= ⟨*digit*$_2$⟩ | `2` | `3` | `4` | `5` | `6` | `7`
⟨*digit*$_2$⟩ ::= `0` | `1`
⟨*exp-mark*$_{16}$⟩ ::= `s` | `l`
⟨*exp-mark*$_{10}$⟩ ::= ⟨*exp-mark*$_{16}$⟩ | `d` | `e` | `f`
⟨*exp-mark*$_8$⟩ ::= ⟨*exp-mark*$_{10}$⟩
⟨*exp-mark*$_2$⟩ ::= ⟨*exp-mark*$_{10}$⟩
⟨*general-number*$_n$⟩ ::= [⟨*exactness*⟩] ⟨*number*$_n$⟩
⟨*exactness*⟩ ::= `#e` | `#i`

Examples:

`-1`	reads equal to `-1`
`1/2`	reads equal to `(/ 1 2)`
`1.0`	reads equal to `(exact->inexact 1)`
`1+2i`	reads equal to `(make-complex 1 2)`
`1/2+3/4i`	reads equal to `(make-complex (/ 1 2) (/ 3 4))`
`1.0+3.0e7i`	reads equal to `(exact->inexact (make-complex 1 30000000))`
`2e5`	reads equal to `(exact->inexact 200000)`
`#i5`	reads equal to `(exact->inexact 5)`
`#e2e5`	reads equal to `200000`
`#x2e5`	reads equal to `741`
`#b101`	reads equal to `5`

1.3.4 Reading Extflonums

An extflonum has the same syntax as an ⟨*inexact-real*$_n$⟩ that includes an ⟨*exp-mark*$_n$⟩, but with `t` or `T` in place of the ⟨*exp-mark*$_n$⟩. In addition, `+inf.t`, `-inf.t`, `+nan.t`, `-nan.t` are extflonums. A `#b` (binary), `#o` (octal), `#d` (decimal), or `#x` (hexadecimal) radix specification can prefix an extflonum, but `#i` or `#e` cannot, and a extflonum cannot be used to form a complex number. The `read-decimal-as-inexact` parameter has no effect on extflonum reading.

1.3.5 Reading Booleans

A `#true`, `#t`, `#T` followed by a delimiter is the input syntax for the boolean constant "true," and `#false`, `#f`, or `#F` followed by a delimiter is the complete input syntax for the boolean constant "false."

1.3.6 Reading Pairs and Lists

When the reader encounters a `(`, `[`, or `{`, it starts parsing a pair or list; see §4.9 "Pairs and Lists" for information on pairs and lists.

To parse the pair or list, the reader recursively reads data until a matching `)`, `]`, or `}` (respectively) is found, and it specially handles a delimited `.`. Pairs `()`, `[]`, and `{}` are treated the same way, so the remainder of this section simply uses "parentheses" to mean any of these pair.

If the reader finds no delimited `.` among the elements between parentheses, then it produces a list containing the results of the recursive reads.

If the reader finds two data between the matching parentheses that are separated by a delimited `.`, then it creates a pair. More generally, if it finds two or more data where the last datum is preceded by a delimited `.`, then it constructs nested pairs: the next-to-last element is paired with the last, then the third-to-last datum is paired with that pair, and so on.

If the reader finds three or more data between the matching parentheses, and if a pair of delimited `.`s surrounds any other than the first and last elements, the result is a list containing the element surrounded by `.`s as the first element, followed by the others in the read order. This convention supports a kind of infix notation at the reader level.

In `read-syntax` mode, the recursive reads for the pair/list elements are themselves in `read-syntax` mode, so that the result is a list or pair of syntax objects that is itself wrapped as a syntax object. If the reader constructs nested pairs because the input included a single delimited `.`, then only the innermost pair and outermost pair are wrapped as syntax objects. Whether wrapping a pair or list, if the pair or list was formed with `[` and `]`, then a `'paren-`

shape property is attached to the result with the value #\[; if the list or pair was formed with { and }, then a 'paren-shape property is attached to the result with the value #\{.

If a delimited . appears in any other configuration, then the exn:fail:read exception is raised. Similarly, if the reader encounters a),], or } that does not end a list being parsed, then the exn:fail:read exception is raised.

Examples:

`()`	reads equal to `(list)`
`(1 2 3)`	reads equal to `(list 1 2 3)`
`{1 2 3}`	reads equal to `(list 1 2 3)`
`[1 2 3]`	reads equal to `(list 1 2 3)`
`(1 (2) 3)`	reads equal to `(list 1 (list 2) 3)`
`(1 . 3)`	reads equal to `(cons 1 3)`
`(1 . (3))`	reads equal to `(list 1 3)`
`(1 . 2 . 3)`	reads equal to `(list 2 1 3)`

If the read-square-bracket-as-paren parameter is set to #f, then when the reader encounters [and], the "exn:fail:read" exception is raised. Similarly, if the read-curly-brace-as-paren parameter is set to #f, then when the reader encounters { and }, the "exn:fail:read" exception is raised.

If the read-accept-dot parameter is set to #f, then a delimited . triggers an exn:fail:read exception. If the read-accept-infix-dot parameter is set to #f, then multiple delimited .s trigger an exn:fail:read exception, instead of the infix conversion.

1.3.7 Reading Strings

§3.4 “Strings (Unicode)” in ***The Racket Guide*** introduces the syntax of strings.

When the reader encounters ", it begins parsing characters to form a string. The string continues until it is terminated by another " (that is not escaped by \). The resulting string is interned in read-syntax mode.

Within a string sequence, the following escape sequences are recognized:

- \a: alarm (ASCII 7)
- \b: backspace (ASCII 8)
- \t: tab (ASCII 9)
- \n: linefeed (ASCII 10)
- \v: vertical tab (ASCII 11)
- \f: formfeed (ASCII 12)
- \r: return (ASCII 13)

- `\e`: escape (ASCII 27)
- `\"`: double-quotes (without terminating the string)
- `\'`: quote (i.e., the backslash has no effect)
- `\\`: backslash (i.e., the second is not an escaping backslash)
- `\`⟨*digit*$_8$⟩$^{\{1,3\}}$: Unicode for the octal number specified by digit$_8^{\{1,3\}}$ (i.e., 1 to 3 ⟨*digit*$_8$⟩s), where each ⟨*digit*$_8$⟩ is `0`, `1`, `2`, `3`, `4`, `5`, `6`, or `7`. A longer form takes precedence over a shorter form, and the resulting octal number must be between 0 and 255 decimal, otherwise the `exn:fail:read` exception is raised.
- `\x`⟨*digit*$_{16}$⟩$^{\{1,2\}}$: Unicode for the hexadecimal number specified by ⟨*digit*$_{16}$⟩$^{\{1,2\}}$, where each ⟨*digit*$_{16}$⟩ is `0`, `1`, `2`, `3`, `4`, `5`, `6`, `7`, `8`, `9`, `a`, `b`, `c`, `d`, `e`, or `f` (case-insensitive). The longer form takes precedence over the shorter form.
- `\u`⟨*digit*$_{16}$⟩$^{\{1,4\}}$: like `\x`, but with up to four hexadecimal digits (longer sequences take precedence). The resulting hexadecimal number must be a valid argument to `integer->char`, otherwise the `exn:fail:read` exception is raised—unless the encoding continues with another `\u` to form a surrogate-style encoding.
- `\u`⟨*digit*$_{16}$⟩$^{\{4,4\}}$`\u`⟨*digit*$_{16}$⟩$^{\{4,4\}}$: like `\u`, but for two hexadecimal numbers, where the first is in the range `#xD800` to `#xDBFF` and the second is in the range `#xDC00` to `#xDFFF`; the resulting character is the one represented by the numbers as a UTF-16 surrogate pair.
- `\U`⟨*digit*$_{16}$⟩$^{\{1,8\}}$: like `\x`, but with up to eight hexadecimal digits (longer sequences take precedence). The resulting hexadecimal number must be a valid argument to `integer->char`, otherwise the `exn:fail:read` exception is raised.
- `\`⟨*newline*⟩: elided, where ⟨*newline*⟩ is either a linefeed, carriage return, or carriage return–linefeed combination. This convention allows single-line strings to span multiple lines in the source.

If the reader encounters any other use of a backslash in a string constant, the `exn:fail:read` exception is raised.

A string constant preceded by `#` is parsed as a byte string. (That is, `#"` starts a byte-string literal.) See §4.4 “Byte Strings” for information on byte strings. The resulting byte string is interned in `read-syntax` mode. Byte-string constants support the same escape sequences as character strings, except `\u` and `\U`.

§3.5 “Bytes and Byte Strings” in ***The Racket Guide*** introduces the syntax of byte strings.

When the reader encounters `#<<`, it starts parsing a *here string*. The characters following `#<<` until a newline character define a terminator for the string. The content of the string includes all characters between the `#<<` line and a line whose only content is the specified terminator. More precisely, the content of the string starts after a newline following `#<<`, and it ends before a newline that is followed by the terminator, where the terminator is itself followed by either a newline or end-of-file. No escape sequences are recognized between

the starting and terminating lines; all characters are included in the string (and terminator) literally. A return character is not treated as a line separator in this context. If no characters appear between `#<<` and a newline or end-of-file, or if an end-of-file is encountered before a terminating line, the `exn:fail:read` exception is raised.

Examples:

`"Apple"`	reads equal to	`"Apple"`
`"\x41pple"`	reads equal to	`"Apple"`
`"\"Apple\""`	reads equal to	`"\x22Apple\x22"`
`"\\"`	reads equal to	`"\x5C"`
`#"Apple"`	reads equal to	`(bytes 65 112 112 108 101)`

1.3.8 Reading Quotes

When the reader encounters `'`, it recursively reads one datum and forms a new list containing the symbol `'quote` and the following datum. This convention is mainly useful for reading Racket code, where `'s` can be used as a shorthand for `(quote s)`.

Several other sequences are recognized and transformed in a similar way. Longer prefixes take precedence over short ones:

`'`	adds	`quote`
`` ` ``	adds	`quasiquote`
`,`	adds	`unquote`
`,@`	adds	`unquote-splicing`
`#'`	adds	`syntax`
`` #` ``	adds	`quasisyntax`
`#,`	adds	`unsyntax`
`#,@`	adds	`unsyntax-splicing`

Examples:

`'apple`	reads equal to	`(list 'quote 'apple)`
`` `(1 ,2) ``	reads equal to	`(list 'quasiquote (list 1 (list 'unquote 2)))`

The `` ` ``, `,`, and `,@` forms are disabled when the `read-accept-quasiquote` parameter is set to `#f`, in which case the `exn:fail:read` exception is raised instead.

1.3.9 Reading Comments

A `;` starts a line comment. When the reader encounters `;`, it skips past all characters until the next linefeed (ASCII 10), carriage return (ASCII 13), next-line (Unicode 133), line-separator (Unicode 8232), or paragraph-separator (Unicode 8233) character.

A `#|` starts a nestable block comment. When the reader encounters `#|`, it skips past all

characters until a closing `|#`. Pairs of matching `#|` and `|#` can be nested.

A `#;` starts an S-expression comment. When the reader encounters `#;`, it recursively reads one datum, and then discards it (continuing on to the next datum for the read result).

A `#!` (which is `#!` followed by a space) or `#!/` starts a line comment that can be continued to the next line by ending a line with `\`. This form of comment normally appears at the beginning of a Unix script file.

Examples:

`; comment`	reads equal to	*nothing*
`#| a |# 1`	reads equal to	`1`
`#| #| a |# 1 |# 2`	reads equal to	`2`
`#;1 2`	reads equal to	`2`
`#!/bin/sh`	reads equal to	*nothing*
`#! /bin/sh`	reads equal to	*nothing*

1.3.10 Reading Vectors

When the reader encounters a `#(`, `#[`, or `#{`, it starts parsing a vector; see §4.11 "Vectors" for information on vectors. A `#fl` in place of `#` starts an flvector, but is not allowed in `read-syntax` mode; see §4.2.3.2 "Flonum Vectors" for information on flvectors. A `#fx` in place of `#` starts an fxvector, but is not allowed in `read-syntax` mode; see §4.2.4.2 "Fixnum Vectors" for information on fxvectors. The `#[`, `#{`, `#fl[`, `#fl{`, `#fx[`, and `#fx{` forms can be disabled through the `read-square-bracket-as-paren` and `read-curly-brace-as-paren` parameters.

The elements of the vector are recursively read until a matching `)`, `]`, or `}` is found, just as for lists (see §1.3.6 "Reading Pairs and Lists"). A delimited `.` is not allowed among the vector elements. In the case of flvectors, the recursive read for element is implicitly prefixed with `#i` and must produce a flonum. In the case of flvectors, the recursive read for element is implicitly prefixed with `#e` and must produce a fixnum.

An optional vector length can be specified between `#`, `#fl`, `#fx` and `(`, `[`, or `{`. The size is specified using a sequence of decimal digits, and the number of elements provided for the vector must be no more than the specified size. If fewer elements are provided, the last provided element is used for the remaining vector slots; if no elements are provided, then `0` is used for all slots.

In `read-syntax` mode, each recursive read for vector elements is also in `read-syntax` mode, so that the wrapped vector's elements are also wrapped as syntax objects, and the vector is immutable.

Examples:

`#(1 apple 3)`	reads equal to	`(vector 1 'apple 3)`
`#3("apple" "banana")`	reads equal to	`(vector "apple" "banana" "banana")`

`#3()` reads equal to `(vector 0 0 0)`

1.3.11 Reading Structures

When the reader encounters a `#s(`, `#s[`, or `#s{`, it starts parsing an instance of a prefab structure type; see §5 “Structures” for information on structure types. The `#s[` and `#s{` forms can be disabled through the `read-square-bracket-as-paren` and `read-curly-brace-as-paren` parameters.

The elements of the structure are recursively read until a matching `)`, `]`, or `}` is found, just as for lists (see §1.3.6 “Reading Pairs and Lists”). A single delimited `.` is not allowed among the elements, but two `.`s can be used as in a list for an infix conversion.

The first element is used as the structure descriptor, and it must have the form (when quoted) of a possible argument to `make-prefab-struct`; in the simplest case, it can be a symbol. The remaining elements correspond to field values within the structure.

In `read-syntax` mode, the structure type must not have any mutable fields. The structure’s elements are read in `read-syntax` mode, so that the wrapped structure’s elements are also wrapped as syntax objects.

If the first structure element is not a valid prefab structure type key, or if the number of provided fields is inconsistent with the indicated prefab structure type, the `exn:fail:read` exception is raised.

1.3.12 Reading Hash Tables

A `#hash` starts an immutable hash-table constant with key matching based on `equal?`. The characters after `hash` must parse as a list of pairs (see §1.3.6 “Reading Pairs and Lists”) with a specific use of delimited `.`: it must appear between the elements of each pair in the list and nowhere in the sequence of list elements. The first element of each pair is used as the key for a table entry, and the second element of each pair is the associated value.

A `#hasheq` starts a hash table like `#hash`, except that it constructs a hash table based on `eq?` instead of `equal?`.

A `#hasheqv` starts a hash table like `#hash`, except that it constructs a hash table based on `eqv?` instead of `equal?`.

In all cases, the table is constructed by adding each mapping to the hash table from left to right, so later mappings can hide earlier mappings if the keys are equivalent.

Examples, where `make-...` stands for `make-immutable-hash`:
`#hash()` reads equal to `(make-... '())`

`#hasheq()`	reads equal to	`(make-...eq '())`
`#hash(("a" . 5))`	reads equal to	`(make-... '(("a" . 5)))`
`#hasheq((a . 5) (b . 7))`	reads equal to	`(make-...eq '((a . 5) (b . 7)))`
`#hasheq((a . 5) (a . 7))`	reads equal to	`(make-...eq '((a . 7)))`

1.3.13 Reading Boxes

When the reader encounters a `#&`, it starts parsing a box; see §4.12 “Boxes” for information on boxes. The content of the box is determined by recursively reading the next datum.

In `read-syntax` mode, the recursive read for the box content is also in `read-syntax` mode, so that the wrapped box’s content is also wrapped as a syntax object, and the box is immutable.

Examples:
`#&17` reads equal to `(box 17)`

1.3.14 Reading Characters

§3.3 “Characters” in ***The Racket Guide*** introduces the syntax of characters.

A `#\` starts a character constant, which has one of the following forms:

- `#\nul` or `#\null`: NUL (ASCII 0); the next character must not be alphabetic.
- `#\backspace`: backspace (ASCII 8); the next character must not be alphabetic.
- `#\tab`: tab (ASCII 9); the next character must not be alphabetic.
- `#\newline` or `#\linefeed`: linefeed (ASCII 10); the next character must not be alphabetic.
- `#\vtab`: vertical tab (ASCII 11); the next character must not be alphabetic.
- `#\page`: page break (ASCII 12); the next character must not be alphabetic.
- `#\return`: carriage return (ASCII 13); the next character must not be alphabetic.
- `#\space`: space (ASCII 32); the next character must not be alphabetic.
- `#\rubout`: delete (ASCII 127); the next character must not be alphabetic.
- `#\`$\langle digit_8 \rangle^{\{1,3\}}$: Unicode for the octal number specified by $\langle digit_8 \rangle^{\{1,3\}}$, as in string escapes (see §1.3.7 “Reading Strings”).
- `#\u`$\langle digit_{16} \rangle^{\{1,4\}}$: like `#\x`, but with up to four hexadecimal digits.
- `#\U`$\langle digit_{16} \rangle^{\{1,6\}}$: like `#\x`, but with up to six hexadecimal digits.

- `#\`⟨*c*⟩: the character ⟨*c*⟩, as long as `#\`⟨*c*⟩ and the characters following it do not match any of the previous cases, and as long as the character after ⟨*c*⟩ is not alphabetic.

Examples:
```
#\newline reads equal to (integer->char 10)
#\n       reads equal to (integer->char 110)
#\u3BB    reads equal to (integer->char 955)
#\λ       reads equal to (integer->char 955)
```

1.3.15 Reading Keywords

A `#:` starts a keyword. The parsing of a keyword after the `#:` is the same as for a symbol, including case-folding in case-insensitive mode, except that the part after `#:` is never parsed as a number. The resulting keyword is interned.

Examples:
```
#:Apple reads equal to (string->keyword "Apple")
#:1     reads equal to (string->keyword "1")
```

1.3.16 Reading Regular Expressions

A `#rx` or `#px` starts a regular expression. The characters immediately after `#rx` or `#px` must parse as a string or byte string (see §1.3.7 “Reading Strings”). A `#rx` prefix starts a regular expression as would be constructed by `regexp`, `#px` as constructed by `pregexp`, `#rx#` as constructed by `byte-regexp`, and `#px#` as constructed by `byte-pregexp`. The resulting regular expression is interned in `read-syntax` mode.

Examples:
```
#rx".*"       reads equal to (regexp ".*")
#px"[\\s]*"   reads equal to (pregexp "[\\s]*")
#rx#".*"      reads equal to (byte-regexp #".*")
#px#"[\\s]*"  reads equal to (byte-pregexp #"[\\s]*")
```

1.3.17 Reading Graph Structure

A `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`=` tags the following datum for reference via `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`#`, which allows the reader to produce a datum that has graph structure.

For a specific ⟨*digit*$_{10}$⟩$^{\{1,8\}}$ in a single read result, each `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`#` reference is replaced by the datum read for the corresponding `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`=`; the definition `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`=` also produces just the datum after it. A `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`=` definition can appear at most once, and a `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`=` definition must appear before a `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`#` reference appears,

otherwise the `exn:fail:read` exception is raised. If the `read-accept-graph` parameter is set to `#f`, then `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`=` or `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`#` triggers a `exn:fail:read` exception.

Although a comment parsed via `#;` discards the datum afterward, `#`⟨*digit*$_{10}$⟩$^{\{1,8\}}$`=` definitions in the discarded datum still can be referenced by other parts of the reader input, as long as both the comment and the reference are grouped together by some other form (i.e., some recursive read); a top-level `#;` comment neither defines nor uses graph tags for other top-level forms.

Examples:

```
(#1=100 #1# #1#)  reads equal to (list 100 100 100)
#0=(1 . #0#)      reads equal to (let* ([ph (make-placeholder #f)]
                                        [v (cons 1 ph)])
                                   (placeholder-set! ph v)
                                   (make-reader-graph v))
```

1.3.18 Reading via an Extension

§17.2 “Reader Extensions” in ***The Racket Guide*** introduces reader extension.

When the reader encounters `#reader`, it loads an external reader procedure and applies it to the current input stream.

The reader recursively reads the next datum after `#reader`, and passes it to the procedure that is the value of the `current-reader-guard` parameter; the result is used as a module path. The module path is passed to `dynamic-require` with either `'read` or `'read-syntax` (depending on whether the reader is in `read` or `read-syntax` mode).

The arity of the resulting procedure determines whether it accepts extra source-location information: a `read` procedure accepts either one argument (an input port) or five, and a `read-syntax` procedure accepts either two arguments (a name value and an input port) or six. In either case, the four optional arguments are the reader's module path (as a syntax object in `read-syntax` mode) followed by the line (positive exact integer or `#f`), column (non-negative exact integer or `#f`), and position (positive exact integer or `#f`) of the start of the `#reader` form. The input port is the one whose stream contained `#reader`, where the stream position is immediately after the recursively read module path.

The procedure should produce a datum result. If the result is a syntax object in `read` mode, then it is converted to a datum using `syntax->datum`; if the result is not a syntax object in `read-syntax` mode, then it is converted to one using `datum->syntax`. See also §13.7.2 “Reader-Extension Procedures” for information on the procedure's results.

If the `read-accept-reader` parameter is set to `#f`, then if the reader encounters `#reader`, the `exn:fail:read` exception is raised.

§6.2.2 “The `#lang` Shorthand” in ***The Racket Guide*** introduces `#lang`.

The `#lang` reader form is similar to `#reader`, but more constrained: the `#lang` must be followed by a single space (ASCII 32), and then a non-empty sequence of alphanumeric

ASCII, `+`, `-`, `_`, and/or `/` characters terminated by whitespace or an end-of-file. The sequence must not start or end with `/`. A sequence `#lang` ⟨*name*⟩ is equivalent to either `#reader (submod` ⟨*name*⟩ `reader)` or `#reader` ⟨*name*⟩`/lang/reader`, where the former is tried first guarded by a `module-declared?` check (but after filtering by `current-reader-guard`, so both are passed to the value of `current-reader-guard` if the latter is used). Note that the terminating whitespace (if any) is not consumed before the external reading procedure is called.

§17.3 "Defining new `#lang` Languages" in ***The Racket Guide*** introduces the creation languages for `#lang`.

Finally, `#!` is an alias for `#lang` followed by a space when `#!` is followed by alphanumeric ASCII, `+`, `-`, or `_`. Use of this alias is discouraged except as needed to construct programs that conform to certain grammars, such as that of R^{6}RS [Sperber07].

The `syntax/module-reader` library provides a domain-specific language for writing language readers.

By convention, `#lang` normally appears at the beginning of a file, possibly after comment forms, to specify the syntax of a module.

If the `read-accept-reader` or `read-accept-lang` parameter is set to `#f`, then if the reader encounters `#lang` or equivalent `#!`, the `exn:fail:read` exception is raised.

S-Expression Reader Language

`#lang s-exp` package: `base`

§17.1.2 "Using `#lang s-exp`" in ***The Racket Guide*** introduces the `s-exp` meta-language.

The `s-exp` "language" is a kind of meta-language. It `read`s the S-expression that follows `#lang s-exp` and uses it as the language of a `module` form. It also reads all remaining S-expressions until an end-of-file, using them for the body of the generated `module`.

That is,

```
#lang s-exp module-path
form ...
```

is equivalent to

```
(module name-id module-path
  form ...)
```

where *name-id* is derived from the source input port's name: if the port name is a filename path, the filename without its directory path and extension is used for *name-id*, otherwise *name-id* is `anonymous-module`.

Chaining Reader Language

`#lang reader` package: `base`

§17.3.2 "Using `#lang reader`" in ***The Racket Guide*** introduces the `reader` meta-language.

The `reader` "language" is a kind of meta-language. It `read`s the S-expression that follows `#lang reader` and uses it as a module path (relative to the module being read) that effectively takes the place of `reader`. In other words, the `reader` meta-language generalizes the

syntax of the module specified after `#lang` to be a module path, and without the implicit addition of `/lang/reader` to the path.

1.4 The Printer

The Racket printer supports three modes:

- `write` mode prints core datatypes in such a way that using `read` on the output produces a value that is `equal?` to the printed value;
- `display` mode prints core datatypes in a more "end-user" style rather than "programmer" style; for example, a string `display`s as its content characters without surrounding `"`s or escapes;
- `print` mode by default—when `print-as-expression` is `#t`—prints most datatypes in such a way that evaluating the output as an expression produces a value that is `equal?` to the printed value; when `print-as-expression` is set to `#f`, then `print` mode is like `write` mode.

In `print` mode when `print-as-expression` is `#t` (as is the default), a value prints at a *quoting depth* of either `0` (unquoted) or `1` (quoted). The initial quoting depth is accepted as an optional argument by `print`, and printing of some compound datatypes adjusts the print depth for component values. For example, when a list is printed at quoting depth `0` and all of its elements are *quotable*, the list is printed with a `'` prefix, and the list's elements are printed at quoting depth `1`.

When the `print-graph` parameter is set to `#t`, then the printer first scans an object to detect cycles. The scan traverses the components of pairs, mutable pairs, vectors, boxes (when `print-box` is `#t`), hash tables (when `print-hash-table` is `#t`), fields of structures exposed by `struct->vector` (when `print-struct` is `#t`), and fields of structures exposed by printing when the structure's type has the `prop:custom-write` property. If `print-graph` is `#t`, then this information is used to print sharing through graph definitions and references (see §1.3.17 "Reading Graph Structure"). If a cycle is detected in the initial scan, then `print-graph` is effectively set to `#t` automatically.

With the exception of displaying byte strings, printing is defined in terms of Unicode characters; see §13.1 "Ports" for information on how a character stream is written to a port's underlying byte stream.

1.4.1 Printing Symbols

Symbols containing spaces or special characters `write` using escaping `\` and quoting `|`s. When the `read-case-sensitive` parameter is set to `#f`, then symbols containing uppercase characters also use escaping `\` and quoting `|`s. In addition, symbols are quoted with

|s or leading \ when they would otherwise print the same as a numerical constant or as a delimited . (when read-accept-dot is #t).

When read-accept-bar-quote is #t, |s are used in printing when one | at the beginning and one | at the end suffice to correctly print the symbol. Otherwise, \s are always used to escape special characters, instead of quoting them with |s.

When read-accept-bar-quote is #f, then | is not treated as a special character. The following are always special characters:

```
( ) [ ] } } " , ' ` ; \
```

In addition, # is a special character when it appears at the beginning of the symbol, and when it is not followed by %.

Symbols display without escaping or quoting special characters. That is, the display form of a symbol is the same as the display form of symbol->string applied to the symbol.

Symbols print the same as they write, unless print-as-expression is set to #t (as is the default) and the current quoting depth is 0. In that case, the symbol's printed form is prefixed with '. For the purposes of printing enclosing datatypes, a symbol is quotable.

1.4.2 Printing Numbers

A number prints the same way in write, display, and print modes. For the purposes of printing enclosing datatypes, a number is quotable.

A complex number that is not a real number always prints as ⟨*m*⟩+⟨*n*⟩i, where ⟨*m*⟩ and ⟨*n*⟩ are the printed forms of its real and imaginary parts, respectively.

An inexact real number prints with either a . decimal point, an e exponent marker, or both. The form is selected so that the output is as short as possible, with the constraint that reading the printed form back in produces an equal? number.

An exact 0 prints as 0.

A positive, exact integer prints as a sequence of decimal digits that does not start with 0.

A positive, exact, real, non-integer number prints as ⟨*m*⟩/⟨*n*⟩, where ⟨*m*⟩ and ⟨*n*⟩ are the printed forms of the number's numerator and denominator (as determined by numerator and denominator).

A negative exact number prints with a - prefix on the printed form of the number's exact negation.

1.4.3 Printing Extflonums

An extflonum prints the same way in `write`, `display`, and `print` modes. For the purposes of printing enclosing datatypes, an extflonum is quotable.

An extflonum prints in the same way an inexact number, but always with a `t` or `T` exponent marker. When extflonum operations are supported, printing always uses `t`; when extflonum operations are not supported, an extflonum prints the same as its reader (see §1.3 "The Reader") source, since reading is the only way to produce an extflonum.

1.4.4 Printing Booleans

The boolean constant `#t` prints as `#true` or `#t` in all modes (`display`, `write`, and `print`), depending on the value of `print-boolean-long-form`, and the constant `#f` prints as `#false` or `#f`. For the purposes of printing enclosing datatypes, a symbol is quotable.

1.4.5 Printing Pairs and Lists

In `write` and `display` modes, an empty list prints as `()`. A pair normally prints starting with `(` followed by the printed form of its `car`. The rest of the printed form depends on the `cdr`:

- If the `cdr` is a pair or the empty list, then the printed form of the pair completes with the printed form of the `cdr`, except that the leading `(` in the `cdr`'s printed form is omitted.
- Otherwise, the printed for of the pair continues with a space, `.`, another space, the printed form of the `cdr`, and a `)`.

If `print-reader-abbreviations` is set to `#t`, then pair printing in `write` mode is adjusted in the case of a pair that starts a two-element list whose first element is `'quote`, `'quasiquote`, `'unquote`, `'unquote-splicing`, `'syntax`, `'quasisyntax`, `'unsyntax`, or `'unsyntax-splicing`. In that case, the pair is printed with the corresponding reader syntax: `'`, `` ` ``, `,`, `,@`, `#'`, `#` `, `#,`, or `#,@`, respectively. After the reader syntax, the second element of the list is printed. When the list is a tail of an enclosing list, the tail is printed after a `.` in the enclosing list (after which the reader abbreviations work), instead of including the tail as two elements of the enclosing list.

The printed form of a pair is the same in both `write` and `display` modes, except as the printed form of the pair's `car` and `cdr` vary with the mode. The `print` form is also the same if `print-as-expression` is `#f` or the quoting depth is `1`.

For print mode when print-as-expression is #t and the quoting depth is 0, then the empty list prints as '(). For a pair whose car and cdr are quotable, the pair prints in write mode but with a ' prefix; the pair's content is printed with quoting depth 1. Otherwise, when the car or cdr is not quotable, then pair prints with either cons (when the cdr is not a pair), list (when the pair is a list), or list* (otherwise) after the opening (, any . that would otherwise be printed is suppressed, and the pair content is printed at quoting depth 0. In all cases, when print-as-expression is #t for print mode, then the value of print-reader-abbreviations is ignored and reader abbreviations are always used for lists printed at quoting depth 1.

By default, mutable pairs (as created with mcons) print the same as pairs for write and display, except that { and } are used instead of (and). Note that the reader treats {...} and (...) equivalently on input, creating immutable pairs in both cases. Mutable pairs in print mode with print-as-expression as #f or a quoting depth of 1 also use { and }. In print mode with print-as-expression as #t and a quoting depth of 0, a mutable pair prints as (mcons , the mcar and mcdr printed at quoting depth 0 and separated by a space, and a closing).

If the print-pair-curly-braces parameter is set to #t, then pairs print using { and } when not using print mode with print-as-expression as #t and a quoting depth of 0. If the print-mpair-curly-braces parameter is set to #f, then mutable pairs print using (and) in that mode.

For the purposes of printing enclosing datatypes, an empty list is always quotable, a pair is quotable when its car and cdr are quotable, and a mutable list is never quotable.

1.4.6 Printing Strings

All strings display as their literal character sequences.

The write or print form of a string starts with " and ends with another ". Between the "s, each character is represented. Each graphic or blank character is represented as itself, with two exceptions: " is printed as \", and \ is printed as \\. Each non-graphic, non-blank character (according to char-graphic? and char-blank?) is printed using the escape sequences described in §1.3.7 "Reading Strings", using \a, \b, \t, \n, \v, \f, \r, or \e if possible, otherwise using \u with four hexadecimal digits or \U with eight hexadecimal digits (using the latter only if the character value does not fit into four digits).

All byte strings display as their literal byte sequence; this byte sequence may not be a valid UTF-8 encoding, so it may not correspond to a sequence of characters.

The write or print form of a byte string starts with #" and ends with a ". Between the "s, each byte is written using the corresponding ASCII decoding if the byte is between 0 and 127 and the character is graphic or blank (according to char-graphic? and char-blank?). Otherwise, the byte is written using \a, \b, \t, \n, \v, \f, \r, or \e if possible, otherwise

using \ followed by one to three octal digits (only as many as necessary).

For the purposes of printing enclosing datatypes, a string or a byte string is quotable.

1.4.7 Printing Vectors

In `display` mode, the printed form of a vector is `#` followed by the printed form of `vector->list` applied to the vector. In `write` mode, the printed form is the same, except that when the `print-vector-length` parameter is `#t`, a decimal integer is printed after the `#`, and a repeated last element is printed only once.

Vectors `print` the same as they `write`, unless `print-as-expression` is set to `#t` and the current quoting depth is `0`. In that case, if all of the vector's elements are quotable, then the vector's `print`ed form is prefixed with `'` and its elements printed with quoting depth `1`. If its elements are not all quotable, then the vector `print`s as `(vector` , the elements at quoting depth `0`, and a closing `)`. A vector is quotable when all of its elements are quotable.

In `write` or `display` mode, an flvector prints like a vector, but with a `#fl` prefix instead of `#`. A fxvector similarly prints with a `#fx` prefix instead of `#`. The `print-vector-length` parameter affects flvector and fxvector printing the same as vector printing. In `print` mode, flvectors and fxvectors are not quotable, and they print like a vector at quoting depth 0 using a `(flvector` or `(fxvector` prefix, respectively.

1.4.8 Printing Structures

When the `print-struct` parameter is set to `#t`, then the way that structures print depends on details of the structure type for which the structure is an instance:

- If the structure type is a prefab structure type, then it prints in `write` or `display` mode using `#s(` followed by the prefab structure type key, then the printed form of each field in the structure, and then `)`.

  In `print` mode when `print-as-expression` is set to `#t` and the current quoting depth is `0`, if the structure's content is all quotable, then the structure's `print`ed form is prefixed with `'` and its content is printed with quoting depth `1`. If any of its content is not quotable, then the structure type prints the same as a non-prefab structure type.

  An instance of a prefab structure type is quotable when all of its content is quotable.

- If the structure has a `prop:custom-write` property value, then the associated procedure is used to print the structure, unless the `print-unreadable` parameter is set to `#f`.

  For `print` mode, an instance of a structure type with a `prop:custom-write` property is treated as quotable if it has the `prop:custom-print-quotable` property with a

value of 'always. If it has 'maybe as the property value, then the structure is treated as quotable if its content is quotable, where the content is determined by the values recursively printed by the structure's prop:custom-write procedure. Finally, if the structure has 'self as the property value, then it is treated as quotable.

In print mode when print-as-expression is #t, the structure's prop:custom-write procedure is called with either 0 or 1 as the quoting depth, normally depending on the structure's prop:custom-print-quotable property value. If the property value is 'always, the quoting depth is normally 1. If the property value is 'maybe, then the quoting depth is 1 if the structure is quotable, or normally 0 otherwise. If the property value is 'self, then the quoting depth may be 0 or 1; it is normally 0 if the structure is not printed as a part of an enclosing quotable value, even though the structure is treated as quotable. Finally, if the property value is 'never, then the quoting depth is normally 0. The quoting depth can vary from its normal value if the structure is printed with an explicit quoting depth of 1.

- If the structure's type is transparent or if any ancestor is transparent (i.e., struct? on the instance produces #t), then the structure prints as the vector produced by struct->vector in display mode, in write mode, or in print mode when print-as-expression is set to #f or when the quoting depth is 0.

  In print mode with print-as-expression as #t and a quoting depth of 0, the structure content is printed with a (followed by the structure's type name (as determined by object-name) in write mode; the remaining elements are printed at quoting depth 0 and separated by a space, and finally a closing).

  A transparent structure type that is not a prefab structure type is never quotable.

- For any other structure type, the structure prints as an unreadable value; see §1.4.15 "Printing Unreadable Values" for more information.

If the print-struct parameter is set to #f, then all structures without a prop:custom-write property print as unreadable values (see §1.4.15 "Printing Unreadable Values") and count as quotable.

1.4.9 Printing Hash Tables

When the print-hash-table parameter is set to #t, in write and display modes, a hash table prints starting with #hash(, #hasheqv(, or #hasheq(for a table using equal?, eqv?, or eq? key comparisons, respectively. After the prefix, each key–value mapping is shown as (, the printed form of a key, a space, ., a space, the printed form the corresponding value, and), with an additional space if the key–value pair is not the last to be printed. After all key–value pairs, the printed form completes with).

In print mode when print-as-expression is #f or the quoting depth is 1, the printed form is the same as for write. Otherwise, if the hash table's keys and values are all quotable, the table prints with a ' prefix, and the table's key and values are printed at quoting depth

1. If some key or value is not quotable, the hash table prints as `(hash` , `(hasheqv` , or `(hasheq` followed by alternating keys and values printed at quoting depth 1 and separated by spaces, and finally a closing `)`. A hash table is quotable when all of its keys and values are quotable.

When the `print-hash-table` parameter is set to `#f`, a hash table prints as `#<hash>` and counts as quotable.

1.4.10 Printing Boxes

When the `print-box` parameter is set to `#t`, a box prints as `#&` followed by the printed form of its content in `write`, `display`, or `print` mode when `print-as-expression` is `#f` or the quoting depth is 1.

In `print` mode when `print-as-expression` is `#t` and the quoting depth is 0, a box prints with a `'` prefix and its value is printed at quoting depth 1 when its content is quotable, otherwise the box prints a `(box` followed by the content at quoting depth 0 and a closing `)`. A box is quotable when its content is quotable.

When the `print-box` parameter is set to `#f`, a box prints as `#<box>` and counts as quotable.

1.4.11 Printing Characters

Characters with the special names described in §1.3.14 “Reading Characters” `write` and `print` using the same name. (Some characters have multiple names; the `#\newline` and `#\nul` names are used instead of `#\linefeed` and `#\null`.) Other graphic characters (according to `char-graphic?`) `write` as `#\` followed by the single character, and all others characters are written in `#\u` notation with four digits or `#\U` notation with eight digits (using the latter only if the character value does not fit in four digits).

All characters `display` directly as themselves (i.e., a single character).

For the purposes of printing enclosing datatypes, a character is quotable.

1.4.12 Printing Keywords

Keywords `write`, `print`, and `display` the same as symbols (see §1.4.1 “Printing Symbols”) except with a leading `#:` (after any `'` prefix added in `print` mode), and without special handling for an initial `#` or when the printed form would match a number or a delimited `.` (since `#:` distinguishes the keyword).

For the purposes of printing enclosing datatypes, a keyword is quotable.

1.4.13 Printing Regular Expressions

Regexp values `write`, `display`, and `print` starting with `#px` (for `pregexp`-based regexps) or `#rx` (for `regexp`-based regexps) followed by the `write` form of the regexp's source string or byte string.

For the purposes of printing enclosing datatypes, a regexp value is quotable.

1.4.14 Printing Paths

Paths `write` and `print` as `#<path:....>`. A path `display`s the same as the string produced by `path->string`. For the purposes of printing enclosing datatypes, a path counts as quotable.

Although a path can be converted to a string with `path->string` or to a byte string with `path->bytes`, neither is clearly the right choice for printing a path and reading it back. If the path value is meant to be moved among platforms, then a string is probably the right choice, despite the potential for losing information when converting a path to a string. For a path that is intended to be re-read on the same platform, a byte string is probably the right choice, since it preserves information in an unportable way. Paths do not print in a readable way so that programmers are not misled into thinking that either choice is always appropriate.

1.4.15 Printing Unreadable Values

For any value with no other printing specification, assuming that the `print-unreadable` parameter is set to `#t`, the output form is `#<`⟨*something*⟩`>`, where ⟨*something*⟩ is specific to the type of the value and sometimes to the value itself. If `print-unreadable` is set to `#f`, then attempting to print an unreadable value raises `exn:fail`.

For the purposes of printing enclosing datatypes, a value that prints unreadably nevertheless counts as quotable.

1.4.16 Printing Compiled Code

Compiled code as produced by `compile` prints using `#~`. Compiled code printed with `#~` is essentially assembly code for Racket, and reading such a form produces a compiled form when the `read-accept-compiled` parameter is set to `#t`.

When a compiled form contains syntax object constants, they must not be tainted or armed; the `#~`-marshaled form drops source-location information and properties (see §12.7 "Syntax Object Properties") for the syntax objects.

Compiled code parsed from `#~` may contain references to unexported or protected bindings from a module. At read time, such references are associated with the current code inspector (see `current-code-inspector`), and the code will only execute if that inspector controls the relevant module invocation (see §14.10 “Code Inspectors”).

A compiled-form object may contain uninterned symbols (see §4.6 “Symbols”) that were created by `gensym` or `string->uninterned-symbol`. When the compiled object is read via `#~`, each uninterned symbol in the original form is mapped to a new uninterned symbol, where multiple instances of a single symbol are consistently mapped to the same new symbol. The original and new symbols have the same printed representation. Unreadable symbols, which are typically generated indirectly during expansion and compilation, are saved and restored consistently through `#~`.

The dynamic nature of uninterned symbols and their localization within `#~` can cause problems when `gensym` or `string->uninterned-symbol` is used to construct an identifier for a top-level or module binding (depending on how the identifier and its references are compiled). To avoid problems, generate distinct identifiers either with `generate-temporaries` or by applying the result of `make-syntax-introducer` to an existing identifier; those functions lead to top-level and module variables with unreadable symbolic names, and the names are deterministic as long as expansion is otherwise deterministic.

Despite the problems inherent with uninterned symbols as variable names, they are partially supported even across multiple `#~`s: When compiled code contains a reference to a module-defined variable whose name is an uninterned symbol, the relative position of the variable among the module’s definitions is recorded, and the reference can be linked back to the definition based on its position and the characters in its name. This accommodation works only for variable references in compiled code; it does not work for `syntax`-quoted identifiers, for example.

Finally, a compiled form may contain path literals. Although paths are not normally printed in a way that can be read back in, path literals can be written and read as part of compiled code. The `current-write-relative-directory` parameter is used to convert the path to a relative path as is it written, and then `current-load-relative-directory` parameter is used to convert any relative path back as it is read. The relative-path conversion applies on reading whether the path was originally relative or not.

2 Notation for Documentation

This chapter introduces essential terminology and notation that is used throughout Racket documentation.

2.1 Notation for Module Documentation

Since Racket programs are organized into modules, documentation reflects that organization with an annotation at the beginning of a section or subsection that describes the bindings that a particular module provides.

For example, the section that describes the functionality provided by `racket/list` starts

```
(require racket/list)
```
package: `base`

Instead of `require`, some modules are introduced with `#lang`:

```
#lang racket/base
```
package: `base`

Using `#lang` means that the module is normally used as the language of a whole module—that is, by a module that starts `#lang` followed by the language—instead of imported with `require`. Unless otherwise specified, however, a module name documented with `#lang` can also be used with `require` to obtain the language's bindings.

The module annotation also shows the package that the module belongs to on the right-hand side. For more details about packages, see *Package Management in Racket*.

Sometimes, a module specification appears at the beginning of a document or at the start of a section that contains many subsections. The document's section or section's subsections are meant to "inherit" the module declaration of the enclosing document or section. Thus, bindings documented in ***The Racket Reference*** are available from `racket` and `racket/base` unless otherwise specified in a section or subsection.

2.2 Notation for Syntactic Form Documentation

§4.1 "Notation" in ***The Racket Guide*** introduces this notation for syntactic forms.

Syntactic forms are specified with a grammar. Typically, the grammar starts with an open parenthesis followed by the syntactic form's name, as in the grammar for `if`:

```
(if test-expr then-expr else-expr)
```

Since every *form* is expressed in terms of syntax objects, parentheses in a grammar specification indicate a syntax object wrapping a list, and the leading `if` is an identifier that starts the list whose binding is the `if` binding of the module being documented—in this case, `racket/base`. Square brackets in the grammar indicate a syntax-object list in the same way as parentheses, but in places square brackets are normally used by convention in a program's source.

Italic identifiers in the grammar are *metavariables* that correspond to other grammar productions. Certain metavariable names have implicit grammar productions:

- A metavariable that ends in *`id`* stands for an identifier.
- A metavariable that ends in *`keyword`* stands for a syntax-object keyword.
- A metavariable that ends with *`expr`* stands for any form, and the form will be parsed as an expression.
- A metavariable that ends with *`body`* stands for any form; the form will be parsed as either a local definition or an expression. A *`body`* can parse as a definition only if it is not preceded by any expression, and the last *`body`* must be an expression; see also §1.2.3.7 "Internal Definitions".
- A metavariable that ends with *`datum`* stands for any form, and the form is normally uninterpreted (e.g., `quoted`).
- A metavariable that ends with *`number`* or *`boolean`* stands for any syntax-object (i.e., literal) number or boolean, respectively.

In a grammar, *`form`* `...` stands for any number of forms (possibly zero) matching *`form`*, while *`form`* `...+` stands for one or more forms matching *`form`*.

Metavariables without an implicit grammar are defined by productions alongside the syntactic form's overall grammar. For example, in

```
(lambda formals body ...+)

formals = id
        | (id ...)
        | (id ...+ . rest-id)
```

the *`formals`* metavariable stands for either an identifier, zero or more identifiers in a syntax-object list, or a syntax object corresponding to a chain of one or more pairs where the chain ends in an identifier instead of an empty list.

Some syntactic forms have multiple top-level grammars, in which case the documentation of the syntactic forms shows multiple grammars. For example,

```
(init-rest id)
(init-rest)
```

indicates that `init-rest` can either be alone in its syntax-object list or followed by a single identifier.

Finally, a grammar specification that includes *expr* metavariables may be augmented with run-time contracts on some of the metavariables, which indicate a predicate that the result of the expression must satisfy at run time. For example,

```
(parameterize ([parameter-expr value-expr] ...)
  body ...+)

  parameter-expr : parameter?
```

indicates that the result of each *parameter-expr* must be a value *v* for which `(parameter? v)` returns true.

2.3 Notation for Function Documentation

Procedures and other values are described using a notation based on contracts. In essence, these contracts describe the interfaces of the documented library using Racket predicates and expressions.

For example, the following is the header of the definition of a typical procedure:

```
(char->integer char) → exact-integer?
  char : char?
```

The function being defined, `char->integer`, is typeset as if it were being applied. The metavariables that come after the function name stand in for arguments. The white text in the corner identifies the kind of value that is being documented.

Each metavariable is described with a contract. In the preceding example, the metavariable *char* has the contract `char?`. This contract specifies that any argument *char* that answers true to the `char?` predicate is valid. The documented function may or may not actually check this property, but the contract signals the intent of the implementer.

The contract on the right of the arrow, `exact-integer?` in this case, specifies the expected result that is produced by the function.

Contract specifications can be more expressive than just names of predicates. Consider the following header for `argmax`:

```
(argmax proc lst) → any
  proc : (-> any/c real?)
  lst : (and/c pair? list?)
```

The contract `(-> any/c real?)` denotes a function contract specifying that `proc`'s argument can be any single value and the result should be a real number. The contract `(and/c pair? list?)` for *lst* specifies that *lst* should pass both `pair?` and `list?` (i.e., that it is a non-empty list).

Both `->` and `and/c` are examples of contract combinators. Contract combinators such as `or/c`, `cons/c`, `listof`, and others are used throughout the documentation. Clicking on the hyperlinked combinator name will provide more information on its meaning.

A Racket function may be documented as having one or more optional arguments. The `read` function is an example of such a function:

```
(read [in]) → any
  in : input-port? = (current-input-port)
```

The brackets surrounding the *in* argument in the application syntax indicates that it is an optional argument.

The header for `read` specifies a contract for the parameter *in* as usual. To the right of the contract, it also specifies a default value `(current-input-port)` that is used if `read` is called with no arguments.

Functions may also be documented as accepting mandatory or optional keyword-based arguments. For example, the `sort` function has two optional, keyword-based arguments:

```
(sort  lst
       less-than?
      [#:key extract-key
       #:cache-keys? cache-keys?]) → list?
  lst : list?
  less-than? : (any/c any/c . -> . any/c)
  extract-key : (any/c . -> . any/c) = (lambda (x) x)
  cache-keys? : boolean? = #f
```

The brackets around the *extract-key* and *cache-keys?* arguments indicate that they are optional as before. The contract section of the header shows the default values that are provided for these keyword arguments.

2.4 Notation for Structure Type Documentation

A structure type is also documented using contract notation:

```
(struct color (red green blue alpha))
  red : (and/c natural-number/c (<=/c 255))
  green : (and/c natural-number/c (<=/c 255))
  blue : (and/c natural-number/c (<=/c 255))
  alpha : (and/c natural-number/c (<=/c 255))
```

The structure type is typeset as it were declared in the source code of a program using the `struct` form. Each field of the structure is documented with a corresponding contract that specifies the values that are accepted for that field.

In the example above, the structure type *color* has four fields: *red*, *green*, *blue*, and *alpha*. The constructor for the structure type accepts field values that satisfy `(and/c natural-number/c (<=/c 255))`, i.e., non-negative exact integers up to 255.

Additional keywords may appear after the field names in the documentation for a structure type:

```
(struct data-source (connector args extensions)
    #:mutable)
  connector : (or/c 'postgresql 'mysql 'sqlite3 'odbc)
  args : list?
  extensions : (listof (list/c symbol? any/c))
```

Here, the `#:mutable` keyword indicates that the fields of instances of the *data-source* structure type can be mutated with their respective setter functions.

2.5 Notation for Parameter Documentation

A parameter is documented the same way as a function:

```
(current-command-line-arguments) → (vectorof string?)
(current-command-line-arguments argv) → void?
  argv : (vectorof (and/c string? immutable?))
```

Since parameters can be referenced or set, there are two entries in the header above. Calling `current-command-line-arguments` with no arguments accesses the parameter's value,

which must be a vector whose elements pass both `string?` and `immutable?`. Calling `current-command-line-arguments` with a single argument sets the parameter's value, where the value must be a vector whose elements pass `string?` (and a guard on the parameter coerces the strings to immutable form, if necessary).

2.6 Notation for Other Documentation

Some libraries provide bindings to constant values. These values are documented with a separate header:

```
object% : class?
```

The `racket/class` library provides the `object%` value, which is the root of the class hierarchy in Racket. Its documentation header just indicates that it is a value that satisfies the predicate `class?`.

3 Syntactic Forms

This section describes the core syntax forms that appear in a fully expanded expression, plus many closely related non-core forms. See §1.2.3.1 “Fully Expanded Programs” for the core grammar.

Notation

Each syntactic form is described by a BNF-like notation that describes a combination of (syntax-wrapped) pairs, symbols, and other data (not a sequence of characters). These grammatical specifications are shown as in the following specification of a `something` form:

```
(something id thing-expr ...)

  thing-expr : number?
```

Within such specifications,

- `...` indicates zero or more repetitions of the preceding datum; more generally, *N* consecutive `...`s a row indicate a consecutive repetition of the preceding *N* datums.
- `...+` indicates one or more repetitions of the preceding datum.
- Italic meta-identifiers play the role of non-terminals. Some meta-identifier names imply syntactic constraints:
  - A meta-identifier that ends in *id* stands for an identifier.
  - A meta-identifier that ends in *keyword* stands for a keyword.
  - A meta-identifier that ends with *expr* (such as *thing-expr*) stands for a sub-form that is expanded as an expression.
  - A meta-identifier that ends with *body* stands for a sub-form that is expanded in an internal-definition context (see §1.2.3.7 “Internal Definitions”).
- Contracts indicate constraints on sub-expression results. For example, *thing-expr* : `number?` indicates that the expression *thing-expr* must produce a number.

3.1 Modules: `module`, `module*`, ...

```
(module id module-path form ...)
```

§6.2.1 “The `module` Form” in ***The Racket Guide*** introduces `module`.

Declares a top-level module or a submodule. For a top-level module, if the `current-module-declare-name` parameter is set, the parameter value is used for the module name

and *id* is ignored, otherwise (quote *id*) is the name of the declared module. For a submodule, *id* is the name of the submodule to be used as an element within a submod module path.

For a module-like form that works in definitions context other than the top level or a module body, see define-package.

The *module-path* form must be as for require, and it supplies the initial bindings for the body *form*s. That is, it is treated like a (require *module-path*) prefix before the *form*s, except that the bindings introduced by *module-path* can be shadowed by definitions and requires in the module body *form*s.

If a single *form* is provided, then it is partially expanded in a module-begin context. If the expansion leads to #%plain-module-begin, then the body of the #%plain-module-begin is the body of the module. If partial expansion leads to any other primitive form, then the form is wrapped with #%module-begin using the lexical context of the module body; this identifier must be bound by the initial *module-path* import, and its expansion must produce a #%plain-module-begin to supply the module body. Finally, if multiple *form*s are provided, they are wrapped with #%module-begin, as in the case where a single *form* does not expand to #%plain-module-begin.

After such wrapping, if any, and before any expansion, an 'enclosing-module-name property is attached to the #%module-begin syntax object (see §12.7 “Syntax Object Properties”); the property's value is a symbol corresponding to *id*.

Each *form* is partially expanded (see §1.2.3.6 “Partial Expansion”) in a module context. Further action depends on the shape of the form:

- If it is a begin form, the sub-forms are flattened out into the module's body and immediately processed in place of the begin.
- If it is a define-syntaxes form, then the right-hand side is evaluated (in phase 1), and the binding is immediately installed for further partial expansion within the module. Evaluation of the right-hand side is parameterized to set current-namespace as in let-syntax.
- If it is a begin-for-syntax form, then the body is expanded (in phase 1) and evaluated. Expansion within a begin-for-syntax form proceeds with the same partial-expansion process as for a module body, but in a higher phase, and saving all #%provide forms for all phases until the end of the module's expansion. Evaluation of the body is parameterized to set current-namespace as in let-syntax.
- If the form is a #%require form, bindings are introduced immediately, and the imported modules are instantiated or visited as appropriate.
- If the form is a #%provide form, then it is recorded for processing after the rest of the body.
- If the form is a define-values form, then the binding is installed immediately, but the right-hand expression is not expanded further.

- If the form is a `module` form, then it is immediately expanded and declared for the extent of the current top-level enclosing module's expansion.
- If the form is a `module*` form, then it is not expanded further.
- Similarly, if the form is an expression, it is not expanded further.

After all *form*s have been partially expanded this way, then the remaining expression forms (including those on the right-hand side of a definition) are expanded in an expression context. After all expression forms, `#%provide` forms are processed in the order in which they appear (independent of phase) in the expanded module. Finally, all `module*` forms are expanded in order, so that each becomes available for use by subsequent `module*` forms; the enclosing module itself is also available for use by `module*` submodules.

The scope of all imported identifiers covers the entire module body, except for nested `module` and `module*` forms (assuming a non-`#f` *module-path* in the latter case). The scope of any identifier defined within the module body similarly covers the entire module body except for such nested `module` and `module*` forms. The ordering of syntax definitions does not affect the scope of the syntax names; a transformer for `A` can produce expressions containing `B`, while the transformer for `B` produces expressions containing `A`, regardless of the order of declarations for `A` and `B`. However, a syntactic form that produces syntax definitions must be defined before it is used.

No identifier can be imported or defined more than once at any phase level within a single module. Every exported identifier must be imported or defined. No expression can refer to a top-level variable. A `module*` form in which the enclosing module's bindings are visible (i.e., a nested `module*` with `#f` instead of a *module-path*) can define or import bindings that shadow the enclosing module's bindings.

The evaluation of a `module` form does not evaluate the expressions in the body of the module. Evaluation merely declares a module, whose full name depends both on *id* or `(current-module-declare-name)`.

A module body is executed only when the module is explicitly instantiated via `require` or `dynamic-require`. On invocation, imported modules are instantiated in the order in which they are `required` into the module (although earlier instantiations or transitive `requires` can trigger the instantiation of a module before its order within a given module). Then, expressions and definitions are evaluated in order as they appear within the module. Each evaluation of an expression or definition is wrapped with a continuation prompt (see `call-with-continuation-prompt`) for the default prompt tag and using a prompt handler that re-aborts and propagates its argument to the next enclosing prompt. Each evaluation of a definition is followed, outside of the prompt, by a check that each of the definition's variables has a value; if the portion of the prompt-delimited continuation that installs values is skipped, then the `exn:fail:contract:variable?` exception is raised.

Accessing a module-level variable before it is defined signals a run-time error, just like accessing an undefined global variable. If a module (in its fully expanded form) does not con-

tain a `set!` for an identifier that defined within the module, then the identifier is a *constant* after it is defined; its value cannot be changed afterward, not even through reflective mechanisms. The `compile-enforce-module-constants` parameter, however, can be used to disable enforcement of constants.

When a syntax object representing a `module` form has a `'module-language` syntax property attached, and when the property value is a vector of three elements where the first is a module path (in the sense of `module-path?`) and the second is a symbol, then the property value is preserved in the corresponding compiled and/or declared module. The third component of the vector should be printable and `read`able, so that it can be preserved in marshaled bytecode. The `racket/base` and `racket` languages attach `'#(racket/language-info get-info #f)` to a `module` form. See also `module-compiled-language-info`, `module->language-info`, and `racket/language-info`.

If a `module` form has a single body *`form`* and if the form is a `#%plain-module-begin` form, then the body *`form`* is traversed to find `module` and `module*` forms that are either immediate, under `begin`, or under `begin-for-syntax`. (That is, the body is searched before adding any lexical context due to the module's initial *`module-path`* import.) Each such module form is given a `'submodule` syntax property that whose value is the initial module form. Then, when `module` or `module*` is expanded in a submodule position, if the form has a `'submodule` syntax property, the property value is used as the form to expand. This protocol avoids the contamination of submodule lexical scope when re-expanding `module` forms that contain submodules.

See also §1.1.10 "Modules and Module-Level Variables" and §1.2.3.8 "Module Phases and Visits".

Example:

```
> (module duck racket/base
    (provide num-eggs quack)
    (define num-eggs 2)
    (define (quack n)
      (unless (zero? n)
        (printf "quack\n")
        (quack (sub1 n)))))
```

```
(module* id module-path form ...)
(module* id #f form ...)
```

§6.2.3 "Submodules" in ***The Racket Guide*** introduces `module*`.

Like `module`, but only for declaring a submodule within a module, and for submodules that may `require` the enclosing module.

Instead of a *`module-path`* after *`id`*, `#f` indicates that all bindings from the enclosing module are visible in the submodule; `begin-for-syntax` forms that wrap the `module*` form shift the phase level of the enclosing module's bindings relative to the submodule. When a

`module*` form has a *module-path*, the submodule starts with an empty lexical context in the same way as a top-level `module` form, and enclosing `begin-for-syntax` forms have no effect on the submodule.

```
(module+ id form ...)
```

§6.2.4 "Main and Test Submodules" in ***The Racket Guide*** introduces `module+`.

Declares and/or adds to a submodule named *id*.

Each addition for *id* is combined in order to form the entire submodule using `(module* id #f ....)` at the end of the enclosing module. If there is only one `module+` for a given *id*, then `(module+ id form ...)` is equivalent to `(module* id #f form ...)`, but still moved to the end of the enclosing module.

When a module contains multiple submodules declared with `module+`, then the relative order of the initial `module+` declarations for each submodule determines the relative order of the `module*` declarations at the end of the enclosing module.

A submodule must not be defined using `module+` *and* `module` or `module*`. That is, if a submodule is made of `module+` pieces, then it must be made *only* of `module+` pieces.

```
(#%module-begin form ...)
```

Legal only in a module begin context, and handled by the `module` and `module*` forms.

The `#%module-begin` form of `racket/base` wraps every top-level expression to print non-`#<void>` results using `current-print`.

The `#%module-begin` form of `racket/base` also declares a `configure-runtime` submodule (before any other *form*), unless some *form* is either an immediate `module` or `module*` form with the name `configure-runtime`. If a `configure-runtime` submodule is added, the submodule calls the `configure` function of `racket/runtime-config`.

```
(#%printing-module-begin form ...)
```

Legal only in a module begin context.

Like `#%module-begin`, but without adding a `configure-runtime` submodule.

```
(#%plain-module-begin form ...)
```

Legal only in a module begin context, and handled by the `module` and `module*` forms.

```
(#%declare declaration-keyword ...)

declaration-keyword = #:cross-phase-persistent
```

Declarations that affect run-time or reflective properties of the module:

- `#:cross-phase-persistent` — declares the module as cross-phase persistent, and reports a syntax error if the module does not meet the import or syntactic constraints of a cross-phase persistent module.

A `#%declare` form must appear in a module context or a module-begin context. Each *declaration-keyword* can be declared at most once within a `module` body.

3.2 Importing and Exporting: `require` and `provide`

```
(require require-spec ...)
```

§6.4 "Imports: `require`" in ***The Racket Guide*** introduces `require`.

```
      require-spec = module-path
                   | (only-in require-spec id-maybe-renamed ...)
                   | (except-in require-spec id ...)
                   | (prefix-in prefix-id require-spec)
                   | (rename-in require-spec [orig-id bind-id] ...)
                   | (combine-in require-spec ...)
                   | (relative-in module-path require-spec ...)
                   | (only-meta-in phase-level require-spec ...)
                   | (for-syntax require-spec ...)
                   | (for-template require-spec ...)
                   | (for-label require-spec ...)
                   | (for-meta phase-level require-spec ...)
                   | derived-require-spec

       module-path = root-module-path
                   | (submod root-module-path submod-path-element ...)
                   | (submod "." submod-path-element ...)
                   | (submod ".." submod-path-element ...)

  root-module-path = (quote id)
                   | rel-string
                   | (lib rel-string ...+)
                   | id
                   | (file string)
                   | (planet id)
                   | (planet string)
                   | (planet rel-string
                             (user-string pkg-string vers)
                             rel-string ...)

submod-path-element = id
                    | ".."

  id-maybe-renamed = id
                   | [orig-id bind-id]

       phase-level = exact-integer
                   | #f

              vers =
                   | nat
                   | nat minor-vers

        minor-vers = nat
                   | (nat nat)
                   | (= nat)
                   | (+ nat)
                   | (- nat)
```

In a top-level context, `require` instantiates modules (see §1.1.10 “Modules and Module-Level Variables”). In a top-level context or module context, expansion of `require` visits modules (see §1.2.3.8 “Module Phases and Visits”). In both contexts and both evaluation and expansion, `require` introduces bindings into a namespace or a module (see §1.2.3.4 “Introducing Bindings”). A `require` form in a expression context or internal-definition context is a syntax error.

A *require-spec* designates a particular set of identifiers to be bound in the importing context. Each identifier is mapped to a particular export of a particular module; the identifier to bind may be different from the symbolic name of the originally exported identifier. Each identifier also binds at a particular phase level.

No identifier can be bound multiple times in a given phase level by an import, unless all of the bindings refer to the same original definition in the same module. In a module context, an identifier can be either imported or defined for a given phase level, but not both.

The syntax of *require-spec* can be extended via `define-require-syntax`, and when multiple *require-spec*s are specified in a `require`, the bindings of each *require-spec* are visible for expanding later *require-spec*s. The pre-defined forms (as exported by `racket/base`) are as follows:

module-path

Imports all exported bindings from the named module, using the export identifiers as the local identifiers. (See below for information on *module-path*.) The lexical context of the *module-path* form determines the context of the introduced identifiers.

(only-in *require-spec id-maybe-renamed* ...)

Like *require-spec*, but constrained to those exports for which the identifiers to bind match *id-maybe-renamed*: as *id* or as *orig-id* in [*orig-id bind-id*]. If the *id* or *orig-id* of any *id-maybe-renamed* is not in the set that *require-spec* describes, a syntax error is reported.

Examples:

```
> (require (only-in racket/tcp
                    tcp-listen
                    [tcp-accept my-accept]))

> tcp-listen
#<procedure:tcp-listen>
> my-accept
#<procedure:tcp-accept>
```

```
> tcp-accept
tcp-accept: undefined;
 cannot reference undefined identifier
```

(except-in *require-spec id ...*)

Like *require-spec*, but omitting those imports for which *id*s are the identifiers to bind; if any *id* is not in the set that *require-spec* describes, a syntax error is reported.

Examples:

```
> (require (except-in racket/tcp
                      tcp-listen))

> tcp-accept
#<procedure:tcp-accept>
> tcp-listen
tcp-listen: undefined;
 cannot reference undefined identifier
```

(prefix-in *prefix-id require-spec*)

Like *require-spec*, but adjusting each identifier to be bound by prefixing it with *prefix-id*. The lexical context of the *prefix-id* is ignored, and instead preserved from the identifiers before prefixing.

Examples:

```
> (require (prefix-in tcp: racket/tcp))

> tcp:tcp-accept
#<procedure:tcp-accept>
> tcp:tcp-listen
#<procedure:tcp-listen>
```

(rename-in *require-spec* [*orig-id bind-id*] ...)

Like *require-spec*, but replacing the identifier to bind *orig-id* with *bind-id*; if any *orig-id* is not in the set that *require-spec* describes, a syntax error is reported.

Examples:

```
> (require (rename-in racket/tcp
                      (tcp-accept accept)
                      (tcp-listen listen)))

> accept
#<procedure:tcp-accept>
> listen
#<procedure:tcp-listen>
```

(combine-in *require-spec* ...)

The union of the *require-spec*s. If two or more imports from the *require-spec*s have the same identifier name but they do not refer to the same original binding, a syntax error is reported.

Examples:

```
> (require (combine-in (only-in racket/tcp tcp-accept)
                       (only-in racket/tcp tcp-listen)))

> tcp-accept
#<procedure:tcp-accept>
> tcp-listen
#<procedure:tcp-listen>
```

(relative-in *module-path require-spec* ...)

Like the union of the *require-spec*s, but each relative module path in a *require-spec* is treated as relative to *module-path* instead of the enclosing context.

The require transformer that implements relative-in sets current-require-module-path to adjust module paths in the *require-spec*s.

(only-meta-in *phase-level require-spec* ...)

Like the combination of *require-spec*s, but removing any binding that is not for *phase-level*, where #f for *phase-level* corresponds to the label phase level.

The following example imports bindings only at phase level 1, the transform phase:

```
> (module nest racket
    (provide (for-syntax meta-eggs)
             (for-meta 1 meta-chicks)
             num-eggs)
    (define-for-syntax meta-eggs 2)
    (define-for-syntax meta-chicks 3)
    (define num-eggs 2))

> (require (only-meta-in 1 'nest))

> (define-syntax (desc stx)
    (printf "~s ~s\n" meta-eggs meta-chicks)
    #'(void))

> (desc)
2 3

> num-eggs
num-eggs: undefined;
 cannot reference undefined identifier
```

The following example imports only bindings at phase level 0, the normal phase.

```
> (require (only-meta-in 0 'nest))

> num-eggs
2
```

(for-meta *phase-level require-spec* ...)

Like the combination of *require-spec*s, but the binding specified by each *require-spec* is shifted by *phase-level*. The label phase level corresponds to `#f`, and a shifting combination that involves `#f` produces `#f`.

Examples:

```
> (module nest racket
    (provide num-eggs)
    (define num-eggs 2))

> (require (for-meta 0 'nest))

> num-eggs
2
> (require (for-meta 1 'nest))
```

```
> (define-syntax (roost stx)
    (datum->syntax stx num-eggs))

> (roost)
2
```

(for-syntax *require-spec* ...)

Same as (for-meta 1 *require-spec* ...).

(for-template *require-spec* ...)

Same as (for-meta -1 *require-spec* ...).

(for-label *require-spec* ...)

Same as (for-meta #f *require-spec* ...). If an identifier in any of the *require-spec*s is bound at more than one phase level, a syntax error is reported.

derived-require-spec

See define-require-syntax for information on expanding the set of *require-spec* forms.

§6.3 "Module Paths" in ***The Racket Guide*** introduces module paths.

A *module-path* identifies a module, either a root module or a submodule that is declared lexically within another module. A root module is identified either through a concrete name in the form of an identifier, or through an indirect name that can trigger automatic loading of the module declaration. Except for the (quote *id*) case below, the actual resolution of a root module path is up to the current module name resolver (see current-module-name-resolver), and the description below corresponds to the default module name resolver.

(quote *id*)

Refers to a submodule previously declared with the name *id* or a module previously declared interactively with the name *id*. When *id* refers to a submodule, (quote *id*) is equivalent to (submod "." *id*).

Examples:

```
; a module declared interactively as test:
> (require 'test)
```

rel-string

A path relative to the containing source (as determined by `current-load-relative-directory` or `current-directory`). Regardless of the current platform, *rel-string* is always parsed as a Unix-format relative path: `/` is the path delimiter (multiple adjacent `/`s are treated as a single delimiter), `..` accesses the parent directory, and `.` accesses the current directory. The path cannot be empty or contain a leading or trailing slash, path elements before than the last one cannot include a file suffix (i.e., a `.` in an element other than `.` or `..`), and the only allowed characters are ASCII letters, ASCII digits, `-`, `+`, `_`, `.`, `/`, and `%`. Furthermore, a `%` is allowed only when followed by two lowercase hexadecimal digits, and the digits must form a number that is not the ASCII value of a letter, digit, `-`, `+`, or `_`.

The `%` provision is intended to support a one-to-one encoding of arbitrary strings as path elements (after UTF-8 encoding). Such encodings are not decoded to arrive at a filename, but instead preserved in the file access.

If *rel-string* ends with a `".ss"` suffix, it is converted to a `".rkt"` suffix. The compiled-load handler may reverse that conversion if a `".rkt"` file does not exist and a `".ss"` exists.

Examples:

```
; a module named "x.rkt" in the same
; directory as the enclosing module's file:
> (require "x.rkt")

; a module named "x.rkt" in the parent directory
; of the enclosing module file's directory:
> (require "../x.rkt")
```

(lib *rel-string* ...+)

A path to a module installed into a collection (see §18.2 “Libraries and Collections”). The *rel-string*s in `lib` are constrained similar to the plain *rel-string* case, with the additional constraint that a *rel-string* cannot contain `.` or `..` directory indicators.

The specific interpretation of the path depends on the number and shape of the *rel-string*s:

- If a single *rel-string* is provided, and if it consists of a single element (i.e., no `/`) with no file suffix (i.e., no `.`), then *rel-string* names a collection, and `"main.rkt"` is the library file name.

Examples:

```
; the main swindle library:
> (require (lib "swindle"))

; the same:
> (require (lib "swindle/main.rkt"))
```

- If a single *rel-string* is provided, and if it consists of multiple /-separated elements, then each element up to the last names a collection, subcollection, etc., and the last element names a file. If the last element has no file suffix, ".rkt" is added, while a ".ss" suffix is converted to ".rkt".

  Examples:

  ```
  ; "turbo.rkt" from the "swindle" collection:
  > (require (lib "swindle/turbo"))

  ; the same:
  > (require (lib "swindle/turbo.rkt"))

  ; the same:
  > (require (lib "swindle/turbo.ss"))
  ```

- If a single *rel-string* is provided, and if it consists of a single element *with* a file suffix (i.e, with a .), then *rel-string* names a file within the "mzlib" collection. A ".ss" suffix is converted to ".rkt". (This convention is for compatibility with older version of Racket.)

  Examples:

  ```
  ; "tar.rkt" module from the "mzlib" collection:
  > (require (lib "tar.ss"))
  ```

- Otherwise, when multiple *rel-string*s are provided, the first *rel-string* is effectively moved after the others, and all *rel-string*s are appended with / separators. The resulting path names a collection, then subcollection, etc., ending with a file name. No suffix is added automatically, but a ".ss" suffix is converted to ".rkt". (This convention is for compatibility with older version of Racket.)

  Examples:

  ```
  ; "tar.rkt" module from the "mzlib" collection:
  > (require (lib "tar.ss" "mzlib"))
  ```

id

A shorthand for a `lib` form with a single *rel-string* whose characters are the same as in the symbolic form of *id*. In addition to the constraints of a `lib` *rel-string*, *id* must not contain `..`.

Example:

```
> (require racket/tcp)
```

(file *string*)

Similar to the plain *rel-string* case, but *string* is a path—possibly absolute—using the current platform's path conventions and `expand-user-path`. A `".ss"` suffix is converted to `".rkt"`.

Example:

```
> (require (file "~/tmp/x.rkt"))
```

```
(planet id)
(planet string)
(planet rel-string (user-string pkg-string vers)
        rel-string ...)
```

Specifies a library available via the PLaneT server.

The first form is a shorthand for the last one, where the *id*'s character sequence must match the following ⟨*spec*⟩ grammar:

⟨*spec*⟩ ::= ⟨*owner*⟩ `/` ⟨*pkg*⟩ ⟨*lib*⟩
⟨*owner*⟩ ::= ⟨*elem*⟩
⟨*pkg*⟩ ::= ⟨*elem*⟩ | ⟨*elem*⟩ `:` ⟨*version*⟩
⟨*version*⟩ ::= ⟨*int*⟩ | ⟨*int*⟩ `:` ⟨*minor*⟩
⟨*minor*⟩ ::= ⟨*int*⟩ | `<=` ⟨*int*⟩ | `>=` ⟨*int*⟩ | `=` ⟨*int*⟩
| ⟨*int*⟩ `-` ⟨*int*⟩
⟨*lib*⟩ ::= ⟨*empty*⟩ | `/` ⟨*path*⟩
⟨*path*⟩ ::= ⟨*elem*⟩ | ⟨*elem*⟩ `/` ⟨*path*⟩

and where an ⟨*elem*⟩ is a non-empty sequence of characters that are ASCII letters, ASCII digits, `-`, `+`, `_`, or `%` followed by lowercase hexadecimal digits (that do not encode one of the other allowed characters), and an ⟨*int*⟩ is a non-empty sequence of ASCII digits. As this shorthand is expended, a `".plt"` extension is added to ⟨*pkg*⟩, and a `".rkt"` extension is added to ⟨*path*⟩; if no ⟨*path*⟩ is included, `"main.rkt"` is used in the expansion.

A (planet string) form is like a (planet id) form with the identifier converted to a string, except that the string can optionally end with a file extension (i.e., a .) for a ⟨path⟩. A ".ss" file extension is converted to ".rkt".

In the more general last form of a planet module path, the rel-strings are similar to the lib form, except that the (user-string pkg-string vers) names a PLaneT-based package instead of a collection. A version specification can include an optional major and minor version, where the minor version can be a specific number or a constraint: (nat nat) specifies an inclusive range, (= nat) specifies an exact match, (+ nat) specifies a minimum version and is equivalent to just nat, and (- nat) specifies a maximum version. The =, +, and - identifiers in a minor-version constraint are recognized symbolically.

Examples:

```
; "main.rkt" in package "farm" by "mcdonald":
> (require (planet mcdonald/farm))

; "main.rkt" in version >= 2.0 of "farm" by "mcdonald":
> (require (planet mcdonald/farm:2))

; "main.rkt" in version >= 2.5 of "farm" by "mcdonald":
> (require (planet mcdonald/farm:2:5))

; "duck.rkt" in version >= 2.5 of "farm" by "mcdonald":
> (require (planet mcdonald/farm:2:5/duck))
```

```
(submod root-module submod-path-element ...)
(submod "." submod-path-element ...)
(submod ".." submod-path-element ...)
```

Identifies a submodule within the module specified by root-module or relative to the current module in the case of (submod "." ....), where (submod ".." submod-path-element ...) is equivalent to (submod "." ".." submod-path-element ...). Submodules have symbolic names, and a sequence of identifiers as submod-path-elements determine a path of successively nested submodules with the given names. A ".." as a submod-path-element names the enclosing module of a submodule, and it's intended for use in (submod "." ....) and (submod ".." ....) forms.

As require prepares to handle a sequence of require-specs, it logs a "prefetch" message to the current logger at the 'info level, using the name 'module-prefetch, and including message data that is a list of two elements: a list of module paths that appear to be imported, and a directory path to use for relative module paths. The logged list of module paths may be incomplete, but a compilation manager can use approximate prefetch information to start on compilations in parallel.

Changed in version 6.0.1.10 of package `base`: Added prefetch logging.

```
(local-require require-spec ...)
```

Like `require`, but for use in a internal-definition context to import just into the local context. Only bindings from phase level 0 are imported.

Examples:

```
> (let ()
    (local-require racket/control)
    fcontrol)
#<procedure:fcontrol>
> fcontrol
fcontrol: undefined;
 cannot reference undefined identifier
```

§6.5 “Exports: `provide`” in ***The Racket Guide*** introduces `provide`.

```
(provide provide-spec ...)

provide-spec = id
             | (all-defined-out)
             | (all-from-out module-path ...)
             | (rename-out [orig-id export-id] ...)
             | (except-out provide-spec provide-spec ...)
             | (prefix-out prefix-id provide-spec)
             | (struct-out id)
             | (combine-out provide-spec ...)
             | (protect-out provide-spec ...)
             | (for-meta phase-level provide-spec ...)
             | (for-syntax provide-spec ...)
             | (for-template provide-spec ...)
             | (for-label provide-spec ...)
             | derived-provide-spec

 phase-level = exact-integer
             | #f
```

Declares exports from a module. A `provide` form must appear in a module context or a module-begin context.

A *provide-spec* indicates one or more bindings to provide. For each exported binding, the external name is a symbol that can be different from the symbolic form of the identifier that is bound within the module. Also, each export is drawn from a particular phase level and exported at the same phase level; by default, the relevant phase level is the number of `begin-for-syntax` forms that enclose the `provide` form.

The syntax of *provide-spec* can be extended by bindings to provide transformers or provide pre-transformers, such as via `define-provide-syntax`, but the pre-defined forms are as follows.

id

Exports *id*, which must be bound within the module (i.e., either defined or imported) at the relevant phase level. The symbolic form of *id* is used as the external name, and the symbolic form of the defined or imported identifier must match (otherwise, the external name could be ambiguous).

Examples:

```
> (module nest racket
    (provide num-eggs)
    (define num-eggs 2))

> (require 'nest)

> num-eggs
2
```

If *id* has a transformer binding to a rename transformer, then the transformer affects the exported binding. See `make-rename-transformer` for more information.

(all-defined-out)

Exports all identifiers that are defined at the relevant phase level within the exporting module, and that have the same lexical context as the `(all-defined-out)` form, excluding bindings to rename transformers where the target identifier has the `'not-provide-all-defined` syntax property. The external name for each identifier is the symbolic form of the identifier. Only identifiers accessible from the lexical context of the `(all-defined-out)` form are included; that is, macro-introduced imports are not re-exported, unless the `(all-defined-out)` form was introduced at the same time.

Examples:

```
> (module nest racket
    (provide (all-defined-out))
    (define num-eggs 2))

> (require 'nest)
```

```
> num-eggs
2
```

(all-from-out *module-path* ...)

Exports all identifiers that are imported into the exporting module using a require-spec built on each *module-path* (see §3.2 “Importing and Exporting: require and provide”) with no phase-level shift. The symbolic name for export is derived from the name that is bound within the module, as opposed to the symbolic name of the export from each *module-path*. Only identifiers accessible from the lexical context of the *module-path* are included; that is, macro-introduced imports are not re-exported, unless the *module-path* was introduced at the same time.

Examples:

```
> (module nest racket
    (provide num-eggs)
    (define num-eggs 2))

> (module hen-house racket
    (require 'nest)
    (provide (all-from-out 'nest)))

> (require 'hen-house)

> num-eggs
2
```

(rename-out [*orig-id export-id*] ...)

Exports each *orig-id*, which must be bound within the module at the relevant phase level. The symbolic name for each export is *export-id* instead orig-d.

Examples:

```
> (module nest racket
    (provide (rename-out [count num-eggs]))
    (define count 2))

> (require 'nest)

> num-eggs
2
> count
count: undefined;
 cannot reference undefined identifier
```

(except-out *provide-spec provide-spec ...*)

Like the first *provide-spec*, but omitting the bindings listed in each subsequent *provide-spec*. If one of the latter bindings is not included in the initial *provide-spec*, a syntax error is reported. The symbolic export name information in the latter *provide-spec*s is ignored; only the bindings are used.

Examples:

```
> (module nest racket
    (provide (except-out (all-defined-out)
                         num-chicks))
    (define num-eggs 2)
    (define num-chicks 3))

> (require 'nest)

> num-eggs
2
> num-chicks
num-chicks: undefined;
 cannot reference undefined identifier
```

(prefix-out *prefix-id provide-spec*)

Like *provide-spec*, but with each symbolic export name from *provide-spec* prefixed with *prefix-id*.

Examples:

```
> (module nest racket
    (provide (prefix-out chicken: num-eggs))
    (define num-eggs 2))

> (require 'nest)

> chicken:num-eggs
2
```

(struct-out *id*)

Exports the bindings associated with a structure type *id*. Typically, *id* is bound with (struct *id* ....); more generally, *id* must have a transformer binding of structure-type information at the relevant phase level; see §5.7 “Structure Type Transformer Binding”. Furthermore, for each identifier mentioned

in the structure-type information, the enclosing module must define or import one identifier that is free-identifier=?. If the structure-type information includes a super-type identifier, and if the identifier has a transformer binding of structure-type information, the accessor and mutator bindings of the super-type are *not* included by `struct-out` for export.

Examples:

```
> (module nest racket
    (provide (struct-out egg))
    (struct egg (color wt)))

> (require 'nest)

> (egg-color (egg 'blue 10))
'blue
```

(combine-out *provide-spec* ...)

The union of the *provide-spec*s.

Examples:

```
> (module nest racket
    (provide (combine-out num-eggs num-chicks))
    (define num-eggs 2)
    (define num-chicks 1))

> (require 'nest)

> num-eggs
2
> num-chicks
1
```

(protect-out *provide-spec* ...)

Like the union of the *provide-spec*s, except that the exports are protected; requiring modules may refer to these bindings, but may not extract these bindings from macro expansions or access them via eval without access privileges. For more details, see §14.10 “Code Inspectors”. The *provide-spec* must specify only bindings that are defined within the exporting module.

Examples:

```
> (module nest racket
    (provide num-eggs (protect-out num-chicks))
    (define num-eggs 2)
    (define num-chicks 3))
```

```
> (define weak-inspector (make-inspector (current-code-
inspector)))

> (define (weak-eval x)
    (parameterize ([current-code-inspector weak-
inspector])
      (define weak-ns (make-base-namespace))
      (namespace-attach-module (current-namespace)
                               ''nest
                               weak-ns)
      (parameterize ([current-namespace weak-ns])
        (namespace-require ''nest)
        (eval x))))

> (require 'nest)

> (list num-eggs num-chicks)
'(2 3)
> (weak-eval 'num-eggs)
2
> (weak-eval 'num-chicks)
num-chicks: access disallowed by code inspector to protected
variable from module: 'nest
  in: num-chicks
```

(for-meta *phase-level provide-spec* ...)

Like the union of the *provide-spec*s, but adjusted to apply to the phase level specified by *phase-level* relative to the current phase level (where `#f` corresponds to the label phase level). In particular, an *id* or `rename-out` form as a *provide-spec* refers to a binding at *phase-level* relative to the current level, an `all-defined-out` exports only definitions at *phase-level* relative to the current phase level, and an `all-from-out` exports bindings imported with a shift by *phase-level*.

Examples:

```
> (module nest racket
    (begin-for-syntax
     (define eggs 2))
    (define chickens 3)
    (provide (for-syntax eggs)
             chickens))

> (require 'nest)
```

```
> (define-syntax (test-eggs stx)
    (printf "Eggs are ~a\n" eggs)
    #'0)

> (test-eggs)
Eggs are 2
0
> chickens
3
> (module broken-nest racket
    (define eggs 2)
    (define chickens 3)
    (provide (for-syntax eggs)
             chickens))
eval:7:0: module: provided identifier not defined or
imported for phase 1
  at: eggs
  in: (#%module-begin (do-wrapping-module-begin print-result
(module configure-runtime (quote #%kernel) (#%require
racket/runtime-config) (configure #f)))
(do-wrapping-module-begin print-result (define eggs 2))
(do-wrapping-module-begin print-result (define c...
> (module nest2 racket
    (begin-for-syntax
     (define eggs 2))
    (provide (for-syntax eggs)))

> (require (for-meta 2 racket/base)
           (for-syntax 'nest2))

> (define-syntax (test stx)
    (define-syntax (show-eggs stx)
      (printf "Eggs are ~a\n" eggs)
      #'0)
    (begin
      (show-eggs)
      #'0))
Eggs are 2

> (test)
0
```

(for-syntax *provide-spec* ...)

Same as (for-meta 1 *provide-spec* ...).

(for-template *provide-spec* ...)

Same as (for-meta -1 *provide-spec* ...).

(for-label *provide-spec* ...)

Same as (for-meta #f *provide-spec* ...).

derived-provide-spec

See define-provide-syntax for information on expanding the set of *provide-spec* forms.

Each export specified within a module must have a distinct symbolic export name, though the same binding can be specified with the multiple symbolic names.

(for-meta *phase-level require-spec* ...)

See require and provide.

(for-syntax *require-spec* ...)

See require and provide.

(for-template *require-spec* ...)

See require and provide.

(for-label *require-spec* ...)

See require and provide.

(#%require *raw-require-spec* ...)

```
    raw-require-spec = phaseless-spec
                     | (for-meta phase-level phaseless-spec ...)
                     | (for-syntax phaseless-spec ...)
                     | (for-template phaseless-spec ...)
                     | (for-label phaseless-spec ...)
                     | (just-meta phase-level raw-require-spec ...)

        phase-level = exact-integer
                     | #f

      phaseless-spec = raw-module-path
                     | (only raw-module-path id ...)
                     | (prefix prefix-id raw-module-path)
                     | (all-except raw-module-path id ...)
                     | (prefix-all-except prefix-id
                                          raw-module-path id ...)
                     | (rename raw-module-path local-id exported-id)

     raw-module-path = raw-root-module-path
                     | (submod raw-root-module-path id ...+)
                     | (submod "." id ...+)

raw-root-module-path = (quote id)
                     | rel-string
                     | (lib rel-string ...)
                     | id
                     | (file string)
                     | (planet rel-string
                               (user-string pkg-string vers ...))
                     | literal-path
```

The primitive import form, to which `require` expands. A *raw-require-spec* is similar to a *require-spec* in a `require` form, except that the syntax is more constrained, not composable, and not extensible. Also, sub-form names like `for-syntax` and `lib` are recognized symbolically, instead of via bindings. Although not formalized in the grammar above, a `just-meta` form cannot appear within a `just-meta` form.

Each *raw-require-spec* corresponds to the obvious *require-spec*, but the `rename` sub-form has the identifiers in reverse order compared to `rename-in`.

For most *raw-require-spec*s, the lexical context of the *raw-require-spec* determines the context of introduced identifiers. The exception is the `rename` sub-form, where the lexical context of the *local-id* is preserved.

A *literal-path* as a *raw-root-module-path* corresponds to a path in the sense of `path?`. Since path values are never produced by `read-syntax`, they appear only in pro-

grammatically constructed expressions. They also appear naturally as arguments to functions such as `namespace-require`, with otherwise take a quoted `raw-module-spec`.

```
(#%provide raw-provide-spec ...)

raw-provide-spec = phaseless-spec
                 | (for-meta phase-level phaseless-spec)
                 | (for-syntax phaseless-spec)
                 | (for-label phaseless-spec)
                 | (protect raw-provide-spec)

    phase-level = exact-integer
                 | #f

  phaseless-spec = id
                 | (rename local-id export-id)
                 | (struct struct-id (field-id ...))
                 | (all-from raw-module-path)
                 | (all-from-except raw-module-path id ...)
                 | (all-defined)
                 | (all-defined-except id ...)
                 | (prefix-all-defined prefix-id)
                 | (prefix-all-defined-except prefix-id id ...)
                 | (protect phaseless-spec ...)
                 | (expand (id . datum))
```

The primitive export form, to which `provide` expands. A *raw-module-path* is as for `#%require`. A `protect` sub-form cannot appear within a `protect` sub-form.

Like `#%require`, the sub-form keywords for `#%provide` are recognized symbolically, and nearly every *raw-provide-spec* has an obvious equivalent *provide-spec* via `provide`, with the exception of the `struct` and `expand` sub-forms.

A `(struct` *struct-id* `(`*field-id* `...))` sub-form expands to *struct-id*, `make-`*struct-id*, `struct:`*struct-id*, *struct-id*`?`, *struct-id*`-`*field-id* for each *field-id*, and `set-`*struct-id*`-`*field-id*`!` for each *field-id*. The lexical context of the *struct-id* is used for all generated identifiers.

Unlike `#%require`, the `#%provide` form is macro-extensible via an explicit `expand` sub-form; the `(`*id* `.` *datum*`)` part is locally expanded as an expression (even though it is not actually an expression), stopping when a `begin` form is produced; if the expansion result is `(begin` *raw-provide-spec* `...)`, it is spliced in place of the `expand` form, otherwise a syntax error is reported. The `expand` sub-form is not normally used directly; it provides a hook for implementing `provide` and provide transformers.

The `all-from` and `all-from-except` forms re-export only identifiers that are accessible in lexical context of the `all-from` or `all-from-except` form itself. That is, macro-

introduced imports are not re-exported, unless the `all-from` or `all-from-except` form was introduced at the same time. Similarly, `all-defined` and its variants export only definitions accessible from the lexical context of the *phaseless-spec* form.

3.2.1 Additional `require` Forms

`(require racket/require)` package: `base`

The bindings documented in this section are provided by the `racket/require` library, not `racket/base` or `racket`.

The following forms support more complex selection and manipulation of sets of imported identifiers.

`(matching-identifiers-in` *regexp require-spec*`)`

Like *require-spec*, but including only imports whose names match *regexp*. The *regexp* must be a literal regular expression (see §4.7 “Regular Expressions”).

Examples:

```
> (module zoo racket/base
    (provide tunafish swordfish blowfish
             monkey lizard ant)
    (define tunafish 1)
    (define swordfish 2)
    (define blowfish 3)
    (define monkey 4)
    (define lizard 5)
    (define ant 6))

> (require racket/require)

> (require (matching-identifiers-in #rx"\\w*fish" 'zoo))

> tunafish
1
> swordfish
2
> blowfish
3
> monkey
monkey: undefined;
 cannot reference undefined identifier
```

`(subtract-in` *require-spec subtracted-spec ...*`)`

Like *require-spec*, but omitting those imports that would be imported by one of the *subtracted-spec*s.

Examples:

```
> (module earth racket
    (provide land sea air)
    (define land 1)
    (define sea 2)
    (define air 3))

> (module mars racket
    (provide aliens)
    (define aliens 4))

> (module solar-system racket
    (require 'earth 'mars)
    (provide (all-from-out 'earth)
             (all-from-out 'mars)))

> (require racket/require)

> (require (subtract-in 'solar-system 'earth))

> land
land: undefined;
 cannot reference undefined identifier
> aliens
4
```

(filtered-in *proc-expr require-spec*)

Applies an arbitrary transformation on the import names (as strings) of *require-spec*. The *proc-expr* must evaluate at expansion time to a single-argument procedure, which is applied on each of the names from *require-spec*. For each name, the procedure must return either a string for the import's new name or `#f` to exclude the import.

For example,

```
(require (filtered-in
          (lambda (name)
            (and (regexp-match? #rx"^[a-z-]+$" name)
                 (regexp-replace #rx"-" (string-
titlecase name) "")))
          racket/base))
```

imports only bindings from racket/base that match the pattern #rx"^[a-z-]+$", and it converts the names to “camel case.”

(path-up rel-string ...)

Specifies paths to modules named by the rel-strings similar to using the rel-strings directly, except that if a required module file is not found relative to the enclosing source, it is searched for in the parent directory, and then in the grand-parent directory, etc., all the way to the root directory. The discovered path relative to the enclosing source becomes part of the expanded form.

This form is useful in setting up a “project environment.” For example, using the following "config.rkt" file in the root directory of your project:

```
#lang racket/base
(require racket/require-syntax
         (for-syntax "utils/in-here.rkt"))

(provide utils-in)
(define-require-syntax utils-in in-here-transformer)
```

and using "utils/in-here.rkt" under the same root directory:

```
#lang racket/base
(require racket/runtime-path)
(provide in-here-transformer)
(define-runtime-path here ".")
(define (in-here-transformer stx)
  (syntax-case stx ()
    [(_ sym)
     (identifier? #'sym)
     (let ([path (build-path here (format "~a.rkt" (syntax-
e #'sym)))])
       (datum->syntax stx `(file ,(path->string path)) stx))]))
```

then path-up works for any other module under the project directory to find "config.rkt":

```
(require racket/require
         (path-up "config.rkt")
         (utils-in foo))
```

Note that the order of requires in the example is important, as each of the first two bind the identifier used in the following.

An alternative in this scenario is to use `path-up` directly to find the utility module:

```
(require racket/require
         (path-up "utils/foo.rkt"))
```

but then sub-directories that are called `"utils"` override the one in the project's root. In other words, the previous method requires only a single unique name.

```
(multi-in subs ...+)

    subs = sub-path
         | (sub-path ...)

sub-path = rel-string
         | id
```

Specifies multiple files to be required from a hierarchy of directories or collections. The set of required module paths is computed as the Cartesian product of the *subs* groups, where each *sub-path* is combined with other *sub-path*s in order using a `/` separator. A *sub-path* as a *subs* is equivalent to `(`*sub-path*`)`. All *sub-path*s in a given `multi-in` form must be either strings or identifiers.

Examples:

```
(require (multi-in racket (dict list)))
```

is equivalent to `(require racket/dict racket/list)`

```
(require (multi-in "math" "matrix" "utils.rkt"))
```

is equivalent to `(require "math/matrix/utils.rkt")`

```
(require (multi-in "utils" ("math.rkt" "matrix.rkt")))
```

is equivalent to `(require "utils/math.rkt" "utils/matrix.rkt")`

```
(require (multi-in ("math" "matrix") "utils.rkt"))
```

is equivalent to `(require "math/utils.rkt" "matrix/utils.rkt")`

```
(require (multi-in ("math" "matrix") ("utils.rkt" "helpers.rkt")))
```

is equivalent to `(require "math/utils.rkt" "math/helpers.rkt"`
`"matrix/utils.rkt" "matrix/helpers.rkt")`

3.2.2 Additional `provide` Forms

```
(require racket/provide)
```
package: base

The bindings documented in this section are provided by the `racket/provide` library, not `racket/base` or `racket`.

```
(matching-identifiers-out regexp provide-spec)
```

Like *provide-spec*, but including only exports of bindings with an external name that matches *regexp*. The *regexp* must be a literal regular expression (see §4.7 "Regular Expressions").

```
(filtered-out proc-expr provide-spec)
```

Analogous to `filtered-in`, but for filtering and renaming exports.

For example,

```
(provide (filtered-out
          (lambda (name)
            (and (regexp-match? #rx"^[a-z-]+$" name)
                 (regexp-replace
                  #rx"-" (string-titlecase name) "")))
          (all-defined-out)))
```

exports only bindings that match the pattern `#rx"^[a-z-]+$"`, and it converts the names to "camel case."

3.3 Literals: `quote` and `#%datum`

Many forms are implicitly quoted (via `#%datum`) as literals. See §1.2.3.2 "Expansion Steps" for more information.

§4.10 "Quoting: `quote` and `'`" in ***The Racket Guide*** introduces `quote`.

```
(quote datum)
```

Produces a constant value corresponding to *datum* (i.e., the representation of the program fragment) without its lexical information, source location, etc. Quoted pairs, vectors, and boxes are immutable.

Examples:

```
> (quote x)
'x
> (quote (+ 1 2))
'(+ 1 2)
> (+ 1 2)
3
```

```
(#%datum . datum)
```

Expands to `(quote datum)`, as long as *datum* is not a keyword. If *datum* is a keyword, a syntax error is reported.

See also §1.2.3.2 “Expansion Steps” for information on how the expander introduces `#%datum` identifiers.

Examples:

```
> (#%datum . 10)
10
> (#%datum . x)
'x
> (#%datum . #:x)
eval:6:0: #%datum: keyword used as an expression
  in: #:x
```

3.4 Expression Wrapper: `#%expression`

```
(#%expression expr)
```

Produces the same result as *expr*. Using `#%expression` forces the parsing of a form as an expression.

Examples:

```
> (#%expression (+ 1 2))
3
> (#%expression (define x 10))
eval:8:0: define: not allowed in an expression context
  in: (define x 10)
```

The `#%expression` form is helpful in recursive definition contexts where expanding a subsequent definition can provide compile-time information for the current expression. For example, consider a `define-sym-case` macro that simply records some symbols at compile-time in a given identifier.

```
(define-syntax (define-sym-case stx)
  (syntax-case stx ()
    [(_ id sym ...)
     (andmap identifier? (syntax->list #'(sym ...)))
     #'(define-syntax id
         '(sym ...))]))
```

and then a variant of `case` that checks to make sure the symbols used in the expression match those given in the earlier definition:

```
(define-syntax (sym-case stx)
  (syntax-case stx ()
    [(_ id val-expr [(sym) expr] ...)
     (let ()
       (define expected-ids
         (syntax-local-value
          #'id
          (λ ()
            (raise-syntax-error
             'sym-case
             "expected an identifier bound via def-sym-case"
             stx
             #'id))))
       (define actual-ids (syntax->datum #'(sym ...)))
       (unless (equal? expected-ids actual-ids)
         (raise-syntax-error
          'sym-case
          (format "expected the symbols ~s"
                  expected-ids)
          stx))
       #'(case val-expr [(sym) expr] ...))]))
```

If the definition follows the use like this, then the `define-sym-case` macro does not have a chance to bind `id` and the `sym-case` macro signals an error:

```
> (let ()
    (sym-case land-creatures 'bear
              [(bear) 1]
              [(fox) 2])
    (define-sym-case land-creatures bear fox))
eval:11:0: sym-case: expected an identifier bound via
def-sym-case
```

```
  at: land-creatures
  in: (sym-case land-creatures (quote bear) ((bear) 1)
((fox) 2))
```

But if the `sym-case` is wrapped in an `#%expression`, then the expander does not need to expand it to know it is an expression and it moves on to the `define-sym-case` expression.

```
> (let ()
    (#%expression (sym-case sea-creatures 'whale
                            [(whale) 1]
                            [(squid) 2]))
    (define-sym-case sea-creatures whale squid)
    'more...)
'more...
```

Of course, a macro like `sym-case` should not require its clients to add `#%expression`; instead it should check the basic shape of its arguments and then expand to `#%expression` wrapped around a helper macro that calls `syntax-local-value` and finishes the expansion.

3.5 Variable References and `#%top`

id

Refers to a module-level or local binding, when *id* is not bound as a transformer (see §1.2.3 “Expansion”). At run-time, the reference evaluates to the value in the location associated with the binding.

When the expander encounters an *id* that is not bound by a module-level or local binding, it converts the expression to `(#%top . id)` giving `#%top` the lexical context of the *id*; typically, that context refers to `#%top`. See also §1.2.3.2 “Expansion Steps”.

Examples:

```
> (define x 10)

> x
10
> (let ([x 5]) x)
5
> ((lambda (x) x) 2)
2
```

(#%top . *id*)

Refers to a module-level or top-level definition. If *id* has a local binding in its context, then (#%top . *id*) refers to a top-level definition, but a reference to a top-level definition is disallowed within a module.

Within a module form, (#%top . *id*) expands to just *id*—with the obligation that *id* is defined within the module and has no local binding in its context. At phase level 0, (#%top . *id*) is an immediate syntax error if *id* is not bound. At phase level 1 and higher, a syntax error is reported if *id* is not defined at the corresponding phase by the end of module-body partial expansion.

See also §1.2.3.2 "Expansion Steps" for information on how the expander introduces #%top identifiers.

Examples:

```
> (define x 12)

> (let ([x 5]) (#%top . x))
12
```

3.6 Locations: #%variable-reference

```
(#%variable-reference id)
(#%variable-reference (#%top . id))
(#%variable-reference)
```

Produces an opaque *variable reference* value representing the location of *id*, which must be bound as a variable. If no *id* is supplied, the resulting value refers to an "anonymous" variable defined within the enclosing context (i.e., within the enclosing module, or at the top level if the form is not inside a module).

A variable reference can be used with variable-reference->empty-namespace, variable-reference->resolved-module-path, and variable-reference->namespace, but facilities like define-namespace-anchor and namespace-anchor->namespace wrap those to provide a clearer interface. A variable reference is also useful to low-level extensions; see *Inside: Racket C API*.

3.7 Procedure Applications and #%app

```
(proc-expr arg ...)
```

§4.3 "Function Calls" in ***The Racket Guide*** introduces procedure applications.

Applies a procedure, when *proc-expr* is not an identifier that has a transformer binding (see §1.2.3 "Expansion").

More precisely, the expander converts this form to (#%app *proc-expr arg* ...), giving #%app the lexical context that is associated with the original form (i.e., the pair that combines *proc-expr* and its arguments). Typically, the lexical context of the pair indicates the procedure-application #%app that is described next. See also §1.2.3.2 "Expansion Steps".

Examples:

```
> (+ 1 2)
3
> ((lambda (x #:arg y) (list y x)) #:arg 2 1)
'(2 1)
```

(#%app *proc-expr arg* ...)

Applies a procedure. Each *arg* is one of the following:

arg-expr

The resulting value is a non-keyword argument.

keyword arg-expr

The resulting value is a keyword argument using *keyword*. Each *keyword* in the application must be distinct.

The *proc-expr* and *arg-expr*s are evaluated in order, left to right. If the result of *proc-expr* is a procedure that accepts as many arguments as non-*keyword arg-expr*s, if it accepts arguments for all of the *keyword*s in the application, and if all required keyword-based arguments are represented among the *keyword*s in the application, then the procedure is called with the values of the arg-exprs. Otherwise, the exn:fail:contract exception is raised.

The continuation of the procedure call is the same as the continuation of the application expression, so the results of the procedure are the results of the application expression.

The relative order of *keyword*-based arguments matters only for the order of *arg-expr* evaluations; the arguments are associated with argument variables in the applied procedure based on the *keyword*s, and not their positions. The other *arg-expr* values, in contrast, are associated with variables according to their order in the application form.

See also §1.2.3.2 "Expansion Steps" for information on how the expander introduces #%app identifiers.

Examples:

```
> (#%app + 1 2)
3
> (#%app (lambda (x #:arg y) (list y x)) #:arg 2 1)
'(2 1)
> (#%app cons)
cons: arity mismatch;
 the expected number of arguments does not match the given
number
  expected: 2
  given: 0
```

```
(#%plain-app proc-expr arg-expr ...)
(#%plain-app)
```

Like `#%app`, but without support for keyword arguments. As a special case, `(#%plain-app)` produces `'()`.

3.8 Procedure Expressions: `lambda` and `case-lambda`

§4.4 "Functions: `lambda`" in ***The Racket Guide*** introduces procedure expressions.

```
(lambda kw-formals body ...+)
(λ kw-formals body ...+)

kw-formals = (arg ...)
           | (arg ...+ . rest-id)
           | rest-id

       arg = id
           | [id default-expr]
           | keyword id
           | keyword [id default-expr]
```

Produces a procedure. The *kw-formals* determines the number of arguments and which keyword arguments that the procedure accepts.

Considering only the first *arg* case, a simple *kw-formals* has one of the following three forms:

```
(id ...)
```

> The procedure accepts as many non-keyword argument values as the number of *id*s. Each *id* is associated with an argument value by position.

`(id ...+ . rest-id)`

The procedure accepts any number of non-keyword arguments greater or equal to the number of *id*s. When the procedure is applied, the *id*s are associated with argument values by position, and all leftover arguments are placed into a list that is associated to *rest-id*.

rest-id

The procedure accepts any number of non-keyword arguments. All arguments are placed into a list that is associated with *rest-id*.

More generally, an *arg* can include a keyword and/or default value. Thus, the first two cases above are more completely specified as follows:

`(arg ...)`

Each *arg* has the following four forms:

id

Adds one to both the minimum and maximum number of non-keyword arguments accepted by the procedure. The *id* is associated with an actual argument by position.

`[id default-expr]`

Adds one to the maximum number of non-keyword arguments accepted by the procedure. The *id* is associated with an actual argument by position, and if no such argument is provided, the *default-expr* is evaluated to produce a value associated with *id*. No *arg* with a *default-expr* can appear before an *id* without a *default-expr* and without a *keyword*.

keyword id

The procedure requires a keyword-based argument using *keyword*. The *id* is associated with a keyword-based actual argument using *keyword*.

keyword [*id default-expr*]

> The procedure accepts a keyword-based argument using *keyword*. The *id* is associated with a keyword-based actual argument using *keyword*, if supplied in an application; otherwise, the *default-expr* is evaluated to obtain a value to associate with *id*.

The position of a *keyword arg* in *kw-formals* does not matter, but each specified *keyword* must be distinct.

(*arg* ...+ . *rest-id*)

> Like the previous case, but the procedure accepts any number of non-keyword arguments beyond its minimum number of arguments. When more arguments are provided than non-*keyword* arguments among the *arg*s, the extra arguments are placed into a list that is associated to *rest-id*.

The *kw-formals* identifiers are bound in the *body*s. When the procedure is applied, a new location is created for each identifier, and the location is filled with the associated argument value. The locations are created and filled in order, with *default-expr*s evaluated as needed to fill locations.

In other words, argument bindings with default-value expressions are evaluated analogous to `let*`.

If any identifier appears in the *body*s that is not one of the identifiers in *kw-formals*, then it refers to the same location that it would if it appeared in place of the `lambda` expression. (In other words, variable reference is lexically scoped.)

When multiple identifiers appear in a *kw-formals*, they must be distinct according to `bound-identifier=?`.

If the procedure produced by `lambda` is applied to fewer or more by-position or by-keyword arguments than it accepts, to by-keyword arguments that it does not accept, or without required by-keyword arguments, then the `exn:fail:contract` exception is raised.

The last *body* expression is in tail position with respect to the procedure body.

Examples:

```
> ((lambda (x) x) 10)
10
> ((lambda (x y) (list y x)) 1 2)
'(2 1)
> ((lambda (x [y 5]) (list y x)) 1 2)
'(2 1)
> (let ([f (lambda (x #:arg y) (list y x))])
   (list (f 1 #:arg 2)
         (f #:arg 2 1)))
```

```
'((2 1) (2 1))
```

When compiling a `lambda` or `case-lambda` expression, Racket looks for a `'method-arity-error` property attached to the expression (see §12.7 “Syntax Object Properties”). If it is present with a true value, and if no case of the procedure accepts zero arguments, then the procedure is marked so that an `exn:fail:contract:arity` exception involving the procedure will hide the first argument, if one was provided. (Hiding the first argument is useful when the procedure implements a method, where the first argument is implicit in the original source). The property affects only the format of `exn:fail:contract:arity` exceptions, not the result of `procedure-arity`.

When a keyword-accepting procedure is bound to an identifier in certain ways, and when the identifier is used in the function position of an application form, then the application form may be expanded in such a way that the original binding is obscured as the target of the application. To help expose the connection between the function application and function declaration, an identifier in the expansion of the function application is tagged with a syntax property accessible via `syntax-procedure-alias-property` if it is effectively an alias for the original identifier. An identifier in the expansion is tagged with a syntax property accessible via `syntax-procedure-converted-arguments-property` if it is like the original identifier except that the arguments are converted to a flattened form: keyword arguments, required by-position arguments, by-position optional arguments, and rest arguments—all as required, by-position arguments; the keyword arguments are sorted by keyword name, each optional keyword argument is followed by a boolean to indicate whether a value is provided, and `#f` is used for an optional keyword argument whose value is not provided; optional by-position arguments include `#f` for each non-provided argument, and then the sequence of optional-argument values is followed by a parallel sequence of booleans to indicate whether each optional-argument value was provided.

```
(case-lambda [formals body ...+] ...)

formals = (id ...)
        | (id ...+ . rest-id)
        | rest-id
```

Produces a procedure. Each `[formals body ...+]` clause is analogous to a single `lambda` procedure; applying the `case-lambda`-generated procedure is the same as applying a procedure that corresponds to one of the clauses—the first procedure that accepts the given number of arguments. If no corresponding procedure accepts the given number of arguments, the `exn:fail:contract` exception is raised.

Note that a `case-lambda` clause supports only *`formals`*, not the more general *`kw-formals`* of `lambda`. That is, `case-lambda` does not directly support keyword and optional arguments.

Example:

```
> (let ([f (case-lambda
```

```
             [() 10]
             [(x) x]
             [(x y) (list y x)]
             [r r])])
    (list (f)
          (f 1)
          (f 1 2)
          (f 1 2 3)))
'(10 1 (2 1) (1 2 3))
```

```
(#%plain-lambda formals body ...+)
```

Like `lambda`, but without support for keyword or optional arguments.

3.9 Local Binding: `let`, `let*`, `letrec`, ...

§4.6 "Local Binding" in ***The Racket Guide*** introduces local binding.

```
(let ([id val-expr] ...) body ...+)
(let proc-id ([id init-expr] ...) body ...+)
```

The first form evaluates the *val-expr*s left-to-right, creates a new location for each *id*, and places the values into the locations. It then evaluates the *body*s, in which the *id*s are bound. The last *body* expression is in tail position with respect to the `let` form. The *id*s must be distinct according to `bound-identifier=?`.

Examples:

```
> (let ([x 5]) x)
5
> (let ([x 5])
    (let ([x 2]
          [y x])
      (list y x)))
'(5 2)
```

The second form evaluates the *init-expr*s; the resulting values become arguments in an application of a procedure `(lambda (id ...) body ...+)`, where *proc-id* is bound within the *body*s to the procedure itself.

Example:

```
> (let fac ([n 10])
    (if (zero? n)
        1
        (* n (fac (sub1 n)))))
3628800
```

```
(let* ([id val-expr] ...) body ...+)
```

Like `let`, but evaluates the *val-expr*s one by one, creating a location for each *id* as soon as the value is available. The *id*s are bound in the remaining *val-expr*s as well as the *body*s, and the *id*s need not be distinct; later bindings shadow earlier bindings.

Example:

```
> (let* ([x 1]
         [y (+ x 1)])
    (list y x))
'(2 1)
```

```
(letrec ([id val-expr] ...) body ...+)
```

Like `let`, including left-to-right evaluation of the *val-expr*s, but the locations for all *id*s are created first, all *id*s are bound in all *val-expr*s as well as the *body*s, and each *id* is initialized immediately after the corresponding *val-expr* is evaluated. The *id*s must be distinct according to `bound-identifier=?`.

Referencing or assigning to an *id* before its initialization raises `exn:fail:contract:variable`. If an *id* (i.e., the binding instance or *id*) has an `'undefined-error-name` syntax property whose value is a symbol, the symbol is used as the name of the variable for error reporting, instead of the symbolic form of *id*.

Example:

```
> (letrec ([is-even? (lambda (n)
                       (or (zero? n)
                           (is-odd? (sub1 n))))]
           [is-odd? (lambda (n)
                      (and (not (zero? n))
                           (is-even? (sub1 n))))])
    (is-odd? 11))
#t
```

Changed in version 6.0.1.2 of package `base`: Changed reference or assignment of an uninitialized *id* to an error.

```
(let-values ([(id ...) val-expr] ...) body ...+)
```

Like `let`, except that each *val-expr* must produce as many values as corresponding *id*s, otherwise the `exn:fail:contract` exception is raised. A separate location is created for each *id*, all of which are bound in the *body*s.

Example:

```
> (let-values ([(x y) (quotient/remainder 10 3)])
    (list y x))
'(1 3)
```

```
(let*-values ([(id ...) val-expr] ...) body ...+)
```

Like `let*`, except that each *val-expr* must produce as many values as corresponding *id*s. A separate location is created for each *id*, all of which are bound in the later *val-expr*s and in the *body*s.

Example:

```
> (let*-values ([(x y) (quotient/remainder 10 3)]
                [(z) (list y x)])
    z)
'(1 3)
```

```
(letrec-values ([(id ...) val-expr] ...) body ...+)
```

Like `letrec`, except that each *val-expr* must produce as many values as corresponding *id*s. A separate location is created for each *id*, all of which are bound in all *val-expr*s and in the *body*s.

Example:

```
> (letrec-values ([(is-even? is-odd?)
                   (values
                     (lambda (n)
                       (or (zero? n)
                           (is-odd? (sub1 n))))
                     (lambda (n)
                       (or (= n 1)
                           (is-even? (sub1 n))))))])
    (is-odd? 11))
#t
```

```
(let-syntax ([id trans-expr] ...) body ...+)
```

See also `splicing-let-syntax`.

Creates a transformer binding (see §1.2.3.5 “Transformer Bindings”) of each *id* with the value of *trans-expr*, which is an expression at phase level 1 relative to the surrounding context. (See §1.2.1 “Identifiers and Binding” for information on phase levels.)

The evaluation of each *trans-expr* is `parameterized` to set `current-namespace` to a namespace that shares bindings and variables with the namespace being used to expand the `let-syntax` form, except that its base phase is one greater.

Each *id* is bound in the *body*s, and not in other *trans-expr*s.

```
(letrec-syntax ([id trans-expr] ...) body ...+)
```

See also `splicing-letrec-syntax`.

Like `let-syntax`, except that each *id* is also bound within all *trans-expr*s.

```
(let-syntaxes ([(id ...) trans-expr] ...) body ...+)
```

See also `splicing-let-syntaxes`.

Like `let-syntax`, but each *trans-expr* must produce as many values as corresponding *id*s, each of which is bound to the corresponding value.

```
(letrec-syntaxes ([(id ...) trans-expr] ...) body ...+)
```

See also `splicing-letrec-syntaxes`.

Like `let-syntax`, except that each *id* is also bound within all *trans-expr*s.

```
(letrec-syntaxes+values ([(trans-id ...) trans-expr] ...)
                        ([(val-id ...) val-expr] ...)
  body ...+)
```

Combines `letrec-syntaxes` with a variant of `letrec-values`: each *trans-id* and *val-id* is bound in all *trans-expr*s and *val-expr*s.

The `letrec-syntaxes+values` form is the core form for local compile-time bindings, since forms like `letrec-syntax` and internal-definition contexts expand to it. In a fully expanded expression (see §1.2.3.1 “Fully Expanded Programs”), the *trans-id* bindings are discarded and the form reduces to a combination of `letrec-values` or `let-values`, but `letrec-syntaxes+values` can appear in the result of `local-expand` with an empty stop list.

For variables bound by `letrec-syntaxes+values`, the location-creation rules differ slightly from `letrec-values`. The `[(`*val-id* `...) `*val-expr*`]` binding clauses are partitioned into minimal sets of clauses that satisfy the following rule: if a clause has a *val-id* binding that is referenced (in a full expansion) by the *val-expr* of an earlier clause, the two clauses and all in between are in the same set. If a set consists of a single clause whose *val-expr* does not refer to any of the clause’s *val-id*s, then locations for the *val-id*s are created *after* the *val-expr* is evaluated. Otherwise, locations for all *val-id*s in a set are created just before the first *val-expr* in the set is evaluated.

The end result of the location-creation rules is that scoping and evaluation order are the same as for `letrec-values`, but the compiler has more freedom to optimize away location creation. The rules also correspond to a nesting of `let-values` and `letrec-values`, which is how `letrec-syntaxes+values` for a fully-expanded expression.

See also `local`, which supports local bindings with `define`, `define-syntax`, and more.

3.10 Local Definitions: `local`

```
(require racket/local)
```
package: base

The bindings documented in this section are provided by the `racket/local` and `racket` libraries, but not `racket/base`.

```
(local [definition ...] body ...+)
```

Like `letrec-syntaxes+values`, except that the bindings are expressed in the same way as in the top-level or in a module body: using `define`, `define-values`, `define-syntax`, `struct`, etc. Definitions are distinguished from non-definitions by partially expanding *definition* forms (see §1.2.3.6 “Partial Expansion”). As in the top-level or in a module body, a `begin`-wrapped sequence is spliced into the sequence of *definition*s.

3.11 Constructing Graphs: `shared`

```
(require racket/shared)
```
package: base

The bindings documented in this section are provided by the `racket/shared` and `racket` libraries, but not `racket/base`.

```
(shared ([id expr] ...) body ...+)
```

Binds *id*s with shared structure according to `exprs` and then evaluates the `body-exprs`, returning the result of the last expression.

The `shared` form is similar to `letrec`, except that special forms of *expr* are recognized (after partial macro expansion) to construct graph-structured data, where the corresponding `letrec` would instead produce a use-before-initialization error.

Each *expr* (after partial expansion) is matched against the following *shared-expr* grammar, where earlier variants in a production take precedence over later variants:

```
shared-expr = shell-expr
            | plain-expr

 shell-expr = (cons in-immutable-expr in-immutable-expr)
            | (list in-immutable-expr ...)
            | (list* in-immutable-expr ...)
            | (append early-expr ... in-immutable-expr)
            | (vector-immutable in-immutable-expr ...)
            | (box-immutable in-immutable-expr)
            | (mcons patchable-expr patchable-expr)
            | (vector patchable-expr ...)
```

```
                    | (box patchable-expr)
                    | (prefix:make-id patchable-expr ...)

in-immutable-expr = shell-id
                  | shell-expr
                  | early-expr

         shell-id = id

   patchable-expr = expr

       early-expr = expr

       plain-expr = expr
```

The *prefix*`:make-`*id* identifier above matches three kinds of references. The first kind is any binding whose name has `make-` in the middle, and where *prefix*`:`*id* has a transformer binding to structure information with a full set of mutator bindings; see §5.7 “Structure Type Transformer Binding”. The second kind is an identifier that itself has a transformer binding to structure information. The third kind is an identifier that has a `'constructor-for` syntax property whose value is an identifier with a transformer binding to structure information. A *shell-id*, meanwhile, must be one of the *id*s bound by the `shared` form to a *shell-expr*.

When the *expr*s of the `shared` form are parsed as *shared-expr* (taking into account the order of the variants for parsing precedence), the sub-expressions that were parsed via *early-expr* will be evaluated first when the `shared` form is evaluated. Among such expressions, they are evaluated in the order as they appear within the `shared` form. However, any reference to an *id* bound by `shared` produces a use-before-initialization errror, even if the binding for the *id* appears before the corresponding *early-expr* within the `shared` form.

The *shell-id*s and *shell-expr*s (not counting *patchable-expr* and *early-expr* sub-expressions) are effectively evaluated next:

- A *shell-id* reference produces the same value as the corresponding *id* will produce within the *body*s, assuming that *id* is never mutated with `set!`. This special handling of a *shell-id* reference is one way in which `shared` supports the creation of cyclic data, including immutable cyclic data.

- A *shell-expr* of the form `(mcons` *patchable-expr patchable-expr*`)`, `(vector` *patchable-expr* `...)`, `(box` *patchable-expr*`)`, or `(`*prefix*`:make-`*id patchable-expr* `...)` produces a mutable value whose content positions are initialized to `undefined`. Each content position is *patched* (i.e., updated) after the corresponding *patchable-expr* expression is later evaluated.

Next, the *plain-expr*s are evaluated as for `letrec`, where a reference to an *id* raises

`exn:fail:contract:variable` if it is evaluated before the right-hand side of the *id* binding.

Finally, the *patchable-expr*s are evaluated and their values replace `undefined`s in the results of *shell-expr*s. At this point, all *id*s are bound, so *patchable-expr*s can create data cycles (but only with cycles that can be created via mutation).

Examples:

```
> (shared ([a (cons 1 a)])
    a)
#0='(1 . #0#)
> (shared ([a (cons 1 b)]
           [b (cons 2 a)])
    a)
#0='(1 2 . #0#)
> (shared ([a (cons 1 b)]
           [b 7])
    a)
'(1 . 7)
> (shared ([a a]) ; no indirection...
    a)
a: undefined;
 cannot use before initialization
> (shared ([a (cons 1 b)] ; b is early...
           [b a])
    a)
a: undefined;
 cannot use before initialization
> (shared ([a (mcons 1 b)] ; b is patchable...
           [b a])
    a)
#0=(mcons 1 #0#)
> (shared ([a (vector b b b)]
           [b (box 1)])
    (set-box! b 5)
    a)
'#(#&5 #&5 #&5)
> (shared ([a (box b)]
           [b (vector (unbox a)   ; unbox after a is patched
                      (unbox c))] ; unbox before c is patched
           [c (box b)])
    b)
#0='#(#0# #<undefined>)
```

3.12 Conditionals: `if`, `cond`, `and`, and `or`

§4.7 "Conditionals" in ***The Racket Guide*** introduces conditionals.

```
(if test-expr then-expr else-expr)
```

Evaluates *test-expr*. If it produces any value other than `#f`, then *then-expr* is evaluated, and its results are the result for the `if` form. Otherwise, *else-expr* is evaluated, and its results are the result for the `if` form. The *then-expr* and *else-expr* are in tail position with respect to the `if` form.

Examples:

```
> (if (positive? -5) (error "doesn't get here") 2)
2
> (if (positive? 5) 1 (error "doesn't get here"))
1
> (if 'we-have-no-bananas "yes" "no")
"yes"
```

```
(cond cond-clause ...)

cond-clause = [test-expr then-body ...+]
            | [else then-body ...+]
            | [test-expr => proc-expr]
            | [test-expr]
```

§4.7.3 "Chaining Tests: `cond`" in ***The Racket Guide*** introduces `cond`.

A *cond-clause* that starts with `else` must be the last *cond-clause*.

If no *cond-clause*s are present, the result is `#<void>`.

If only a `[else then-body ...+]` is present, then the *then-body*s are evaluated. The results from all but the last *then-body* are ignored. The results of the last *then-body*, which is in tail position with respect to the `cond` form, are the results for the whole `cond` form.

Otherwise, the first *test-expr* is evaluated. If it produces `#f`, then the result is the same as a `cond` form with the remaining *cond-clause*s, in tail position with respect to the original `cond` form. Otherwise, evaluation depends on the form of the *cond-clause*:

```
[test-expr then-body ...+]
```

> The *then-body*s are evaluated in order, and the results from all but the last *then-body* are ignored. The results of the last *then-body*, which is in tail position with respect to the `cond` form, provides the result for the whole `cond` form.

[*test-expr* => *proc-expr*]

The *proc-expr* is evaluated, and it must produce a procedure that accepts one argument, otherwise the exn:fail:contract exception is raised. The procedure is applied to the result of *test-expr* in tail position with respect to the cond expression.

[*test-expr*]

The result of the *test-expr* is returned as the result of the cond form. The *test-expr* is not in tail position.

Examples:

```
> (cond)

> (cond
    [else 5])
5
> (cond
   [(positive? -5) (error "doesn't get here")]
   [(zero? -5) (error "doesn't get here, either")]
   [(positive? 5) 'here])
'here
> (cond
   [(member 2 '(1 2 3)) => (lambda (l) (map - l))])
'(-2 -3)
> (cond
   [(member 2 '(1 2 3))])
'(2 3)
```

else

Recognized specially within forms like cond. An else form as an expression is a syntax error.

=>

Recognized specially within forms like cond. A => form as an expression is a syntax error.

(and *expr* ...)

If no *expr*s are provided, then result is #t.

§4.7.2 "Combining Tests: and and or" in ***The Racket Guide*** introduces and.

If a single *expr* is provided, then it is in tail position, so the results of the `and` expression are the results of the *expr*.

Otherwise, the first *expr* is evaluated. If it produces `#f`, the result of the `and` expression is `#f`. Otherwise, the result is the same as an `and` expression with the remaining *expr*s in tail position with respect to the original `and` form.

Examples:

```
> (and)
#t
> (and 1)
1
> (and (values 1 2))
1
2
> (and #f (error "doesn't get here"))
#f
> (and #t 5)
5
```

`(or expr ...)`

§4.7.2 "Combining Tests: `and` and `or`" in ***The Racket Guide*** introduces `or`.

If no *expr*s are provided, then result is `#f`.

If a single *expr* is provided, then it is in tail position, so the results of the `or` expression are the results of the *expr*.

Otherwise, the first *expr* is evaluated. If it produces a value other than `#f`, that result is the result of the `or` expression. Otherwise, the result is the same as an `or` expression with the remaining *expr*s in tail position with respect to the original `or` form.

Examples:

```
> (or)
#f
> (or 1)
1
> (or (values 1 2))
1
2
> (or 5 (error "doesn't get here"))
5
> (or #f 5)
5
```

3.13 Dispatch: `case`

```
(case val-expr case-clause ...)

case-clause = [(datum ...) then-body ...+]
            | [else then-body ...+]
```

Evaluates *val-expr* and uses the result to select a *case-clause*. The selected clause is the first one with a *datum* whose `quoted` form is `equal?` to the result of *val-expr*. If no such *datum* is present, the `else` *case-clause* is selected; if no `else` *case-clause* is present, either, then the result of the `case` form is `#<void>`.

The `case` form of `racket` differs from that of *R6RS: Scheme* or *R5RS: Legacy Scheme* by being based `equal?` instead of `eqv?` (in addition to allowing internal definitions).

For the selected *case-clause*, the results of the last *then-body*, which is in tail position with respect to the `case` form, are the results for the whole `case` form.

A *case-clause* that starts with `else` must be the last *case-clause*.

The `case` form can dispatch to a matching *case-clause* in $O(\log N)$ time for N *datum*s.

Examples:

```
> (case (+ 7 5)
    [(1 2 3) 'small]
    [(10 11 12) 'big])
'big
> (case (- 7 5)
    [(1 2 3) 'small]
    [(10 11 12) 'big])
'small
> (case (string-append "do" "g")
    [("cat" "dog" "mouse") "animal"]
    [else "mineral or vegetable"])
"animal"
> (case (list 'y 'x)
    [((a b) (x y)) 'forwards]
    [((b a) (y x)) 'backwards])
'backwards
> (case 'x
    [(x) "ex"]
    [('x) "quoted ex"])
"ex"
> (case (list 'quote 'x)
    [(x) "ex"]
    [('x) "quoted ex"])
"quoted ex"


(define (classify c)
  (case (char-general-category c)
```

```
    [(ll lu lt ln lo) "letter"]
    [(nd nl no) "number"]
    [else "other"]))

> (classify #\A)
"letter"
> (classify #\1)
"number"
> (classify #\!)
"other"
```

3.14 Definitions: `define`, `define-syntax`, ...

§4.5 "Definitions: `define`" in ***The Racket Guide*** introduces definitions.

```
(define id expr)
(define (head args) body ...+)

head = id
     | (head args)

args = arg ...
     | arg ... . rest-id

 arg = arg-id
     | [arg-id default-expr]
     | keyword arg-id
     | keyword [arg-id default-expr]
```

The first form binds *id* to the result of *expr*, and the second form binds *id* to a procedure. In the second case, the generated procedure is `(CVT (head args) body ...+)`, using the `CVT` meta-function defined as follows:

```
(CVT (id . kw-formals) . datum)    = (lambda kw-formals . datum)
(CVT (head . kw-formals) . datum) = (lambda kw-formals expr)
                                      if (CVT head . datum) = expr
```

In an internal-definition context, a `define` form introduces a local binding; see §1.2.3.7 "Internal Definitions". At the top level, the top-level binding for *id* is created after evaluating *expr*, if it does not exist already, and the top-level mapping of *id* (in the namespace linked with the compiled definition) is set to the binding at the same time.

In a context that allows liberal expansion of `define`, *id* is bound as syntax if *expr* is an immediate `lambda` form with keyword arguments or *args* include keyword arguments.

Examples:

```
(define x 10)

> x
10


(define (f x)
  (+ x 1))

> (f 10)
11


(define ((f x) [y 20])
  (+ x y))

> ((f 10) 30)
40
> ((f 10))
30
```

```
(define-values (id ...) expr)
```

Evaluates the *expr*, and binds the results to the *id*s, in order, if the number of results matches the number of *id*s; if *expr* produces a different number of results, the `exn:fail:contract` exception is raised.

In an internal-definition context (see §1.2.3.7 “Internal Definitions”), a `define-values` form introduces local bindings. At the top level, the top-level binding for each *id* is created after evaluating *expr*, if it does not exist already, and the top-level mapping of each *id* (in the namespace linked with the compiled definition) is set to the binding at the same time.

Examples:

```
(define-values () (values))

(define-values (x y z) (values 1 2 3))

> z
3
```

If a `define-values` form for a function definition in a module body has a `'compiler-hint:cross-module-inline` syntax property with a true value, then the Racket treats the property as a performance hint. See §19.4 “Function-Call Optimizations” in ***The Racket Guide*** for more information, and see also `begin-encourage-inline`.

```
(define-syntax id expr)
(define-syntax (head args) body ...+)
```

The first form creates a transformer binding (see §1.2.3.5 “Transformer Bindings”) of *id* with the value of *expr*, which is an expression at phase level 1 relative to the surrounding context. (See §1.2.1 “Identifiers and Binding” for information on phase levels.) Evaluation of *expr* side is `parameterized` to set `current-namespace` as in `let-syntax`.

The second form is a shorthand the same as for `define`; it expands to a definition of the first form where the *expr* is a `lambda` form.

In an internal-definition context (see §1.2.3.7 “Internal Definitions”), a `define-syntax` form introduces a local binding.

Examples:

```
> (define-syntax foo
    (syntax-rules ()
      ((_ a ...)
       (printf "~a\n" (list a ...)))))

> (foo 1 2 3 4)
(1 2 3 4)

> (define-syntax (bar syntax-object)
    (syntax-case syntax-object ()
      ((_ a ...)
       #'(printf "~a\n" (list a ...)))))

> (bar 1 2 3 4)
(1 2 3 4)
```

```
(define-syntaxes (id ...) expr)
```

Like `define-syntax`, but creates a transformer binding for each *id*. The *expr* should produce as many values as *id*s, and each value is bound to the corresponding *id*.

When *expr* produces zero values for a top-level `define-syntaxes` (i.e., not in a module or internal-definition position), then the *id*s are effectively declared without binding; see §1.2.3.9 “Macro-Introduced Bindings”.

In an internal-definition context (see §1.2.3.7 “Internal Definitions”), a `define-syntaxes` form introduces local bindings.

Examples:

```
> (define-syntaxes (foo1 foo2 foo3)
    (let ([transformer1 (lambda (syntax-object)
                          (syntax-case syntax-object ()
```

```
                          [(_) #'1]))]
          [transformer2 (lambda (syntax-object)
                          (syntax-case syntax-object ()
                            [(_) #'2]))]
          [transformer3 (lambda (syntax-object)
                          (syntax-case syntax-object ()
                            [(_) #'3]))])
      (values transformer1
              transformer2
              transformer3)))

> (foo1)
1
> (foo2)
2
> (foo3)
3
```

```
(define-for-syntax id expr)
(define-for-syntax (head args) body ...+)
```

Like `define`, except that the binding is at phase level 1 instead of phase level 0 relative to its context. The expression for the binding is also at phase level 1. (See §1.2.1 “Identifiers and Binding” for information on phase levels.) The form is a shorthand for `(begin-for-syntax (define id expr))` or `(begin-for-syntax (define (head args) body ...+))`.

Within a module, bindings introduced by `define-for-syntax` must appear before their uses or in the same `define-for-syntax` form (i.e., the `define-for-syntax` form must be expanded before the use is expanded). In particular, mutually recursive functions bound by `define-for-syntax` must be defined by the same `define-for-syntax` form.

Examples:

```
> (define-for-syntax helper 2)

> (define-syntax (make-two syntax-object)
    (printf "helper is ~a\n" helper)
    #'2)

> (make-two)
helper is 2
2
; ‘helper' is not bound in the runtime phase
> helper
helper: undefined;
 cannot reference undefined identifier
```

```
> (define-for-syntax (filter-ids ids)
    (filter identifier? ids))

> (define-syntax (show-variables syntax-object)
    (syntax-case syntax-object ()
      [(_ expr ...)
       (with-syntax ([(only-ids ...)
                      (filter-ids (syntax->list #'(expr ...)))])
         #'(list only-ids ...))]))

> (let ([a 1] [b 2] [c 3])
    (show-variables a 5 2 b c))
'(1 2 3)
```

```
(define-values-for-syntax (id ...) expr)
```

Like `define-for-syntax`, but *expr* must produce as many values as supplied *id*s, and all of the *id*s are bound (at phase level 1).

Examples:

```
> (define-values-for-syntax (foo1 foo2) (values 1 2))

> (define-syntax (bar syntax-object)
    (printf "foo1 is ~a foo2 is ~a\n" foo1 foo2)
    #'2)

> (bar)
foo1 is 1 foo2 is 2
2
```

3.14.1 `require` Macros

```
(require racket/require-syntax)
```
package: `base`

The bindings documented in this section are provided by the `racket/require-syntax` library, not `racket/base` or `racket`.

```
(define-require-syntax id proc-expr)
(define-require-syntax (id args ...) body ...+)
```

The first form is like `define-syntax`, but for a `require` sub-form. The *proc-expr* must produce a procedure that accepts and returns a syntax object representing a `require` sub-form.

This form expands to `define-syntax` with a use of `make-require-transformer` (see §12.4.1 "`require` Transformers" for more information), and the syntax object passed to and from the macro transformer is marked via `syntax-local-require-introduce`.

The second form is a shorthand the same as for `define-syntax`; it expands to a definition of the first form where the *`proc-expr`* is a `lambda` form.

```
(syntax-local-require-introduce stx) → syntax?
  stx : syntax?
```

Provided `for-syntax` for use only during the application of a `require` sub-form macro transformer: like `syntax-local-introduce`, but for `require` sub-form expansion.

3.14.2 `provide` Macros

```
(require racket/provide-syntax)
```
package: `base`

The bindings documented in this section are provided by the `racket/provide-syntax` library, not `racket/base` or `racket`.

```
(define-provide-syntax id proc-expr)
(define-provide-syntax (id args ...) body ...+)
```

The first form is like `define-syntax`, but for a `provide` sub-form. The *`proc-expr`* must produce a procedure that accepts and returns a syntax object representing a `provide` sub-form.

This form expands to `define-syntax` with a use of `make-provide-transformer` (see §12.4.2 "`provide` Transformers" for more information), and the syntax object passed to and from the macro transformer is marked via `syntax-local-provide-introduce`.

The second form is a shorthand the same as for `define-syntax`; it expands to a definition of the first form where the `expr` is a `lambda` form.

```
(syntax-local-provide-introduce stx) → syntax?
  stx : syntax?
```

Provided `for-syntax` for use only during the application of a `provide` sub-form macro transformer: like `syntax-local-introduce`, but for `provide` sub-form expansion.

3.15 Sequencing: `begin`, `begin0`, and `begin-for-syntax`

§4.8 "Sequencing" in ***The Racket Guide*** introduces `begin` and `begin0`.

```
(begin form ...)
(begin expr ...+)
```

The first form applies when `begin` appears at the top level, at module level, or in an internal-definition position (before any expression in the internal-definition sequence). In that case, the `begin` form is equivalent to splicing the *`form`*s into the enclosing context.

The second form applies for `begin` in an expression position. In that case, the *`expr`*s are evaluated in order, and the results are ignored for all but the last *`expr`*. The last *`expr`* is in tail position with respect to the `begin` form.

Examples:

```
> (begin
    (define x 10)
    x)
10
> (+ 1 (begin
         (printf "hi\n")
         2))
hi
3
> (let-values ([(x y) (begin
                        (values 1 2 3)
                        (values 1 2))])
   (list x y))
'(1 2)
```

```
(begin0 expr ...+)
```

Evaluates the first *`expr`*, then evaluates the other *`expr`*s in order, ignoring their results. The results of the first *`expr`* are the results of the `begin0` form; the first *`expr`* is in tail position only if no other *`expr`*s are present.

Example:

```
> (begin0
    (values 1 2)
    (printf "hi\n"))
hi
1
2
```

```
(begin-for-syntax form ...)
```

Allowed only in a top-level context or module context, shifts the phase level of each *`form`* by one:

- expressions reference bindings at a phase level one greater than in the context of the `begin-for-syntax` form;

- `define`, `define-values`, `define-syntax`, and `define-syntaxes` forms bind at a phase level one greater than in the context of the `begin-for-syntax` form;
- in `require` and `provide` forms, the default phase level is greater, which is roughly like wrapping the content of the `require` form with `for-syntax`;
- expression form *expr*: converted to `(define-values-for-syntax () (begin` *expr* `(values)))`, which effectively evaluates the expression at expansion time and, in the case of a module context, preserves the expression for future visits of the module.

See also `module` for information about expansion order and partial expansion for `begin-for-syntax` within a module context. Evaluation of an `expr` within `begin-for-syntax` is `parameterized` to set `current-namespace` as in `let-syntax`.

3.16 Guarded Evaluation: `when` and `unless`

§4.8.3 "Effects If...: `when` and `unless`" in ***The Racket Guide*** introduces `when` and `unless`.

`(when` *test-expr body* `...+)`

Evaluates *test-expr*. If the result is `#f`, then the result of the `when` expression is `#<void>`. Otherwise, the *body*s are evaluated, and the last *body* is in tail position with respect to the `when` form.

Examples:

```
> (when (positive? -5)
    (display "hi"))

> (when (positive? 5)
    (display "hi")
    (display " there"))
hi there
```

`(unless` *test-expr body* `...+)`

Equivalent to `(when (not` *test-expr*`)` *body* `...+)`.

Examples:

```
> (unless (positive? 5)
    (display "hi"))

> (unless (positive? -5)
    (display "hi")
    (display " there"))
hi there
```

3.17 Assignment: `set!` and `set!-values`

§4.9 “Assignment: `set!`” in ***The Racket Guide*** introduces `set!`.

```
(set! id expr)
```

If *id* has a transformer binding to an assignment transformer, as produced by `make-set!-transformer` or as an instance of a structure type with the `prop:set!-transformer` property, then this form is expanded by calling the assignment transformer with the full expressions. If *id* has a transformer binding to a rename transformer as produced by `make-rename-transformer` or as an instance of a structure type with the `prop:rename-transformer` property, then this form is expanded by replacing *id* with the target identifier (e.g., the one provided to `make-rename-transformer`). If a transformer binding has both `prop:set!-transformer` and `prop:rename-transformer` properties, the latter takes precedence.

Otherwise, evaluates *expr* and installs the result into the location for *id*, which must be bound as a local variable or defined as a top-level variable or module-level variable. If *id* refers to an imported binding, a syntax error is reported. If *id* refers to a top-level variable that has not been defined, the `exn:fail:contract` exception is raised.

See also `compile-allow-set!-undefined`.

Examples:

```
(define x 12)

> (set! x (add1 x))

> x
13
> (let ([x 5])
    (set! x (add1 x))
    x)
6
> (set! i-am-not-defined 10)
set!: assignment disallowed;
 cannot set undefined
  variable: i-am-not-defined
```

```
(set!-values (id ...) expr)
```

Assuming that all *id*s refer to variables, this form evaluates *expr*, which must produce as many values as supplied *id*s. The location of each *id* is filled with the corresponding value from *expr* in the same way as for `set!`.

Example:

```
> (let ([a 1]
        [b 2])
    (set!-values (a b) (values b a))
    (list a b))
'(2 1)
```

More generally, the `set!-values` form is expanded to

```
(let-values ([(tmp-id ...) expr])
  (set! id tmp-id) ...)
```

which triggers further expansion if any *id* has a transformer binding to an assignment transformer.

3.18 Iterations and Comprehensions: `for`, `for/list`, ...

§11 “Iterations and Comprehensions” in ***The Racket Guide*** introduces iterations and comprehensions.

The `for` iteration forms are based on SRFI-42 [SRFI-42].

3.18.1 Iteration and Comprehension Forms

```
(for (for-clause ...) body-or-break ... body)

   for-clause = [id seq-expr]
              | [(id ...) seq-expr]
              | #:when guard-expr
              | #:unless guard-expr
              | break-clause

 break-clause = #:break guard-expr
              | #:final guard-expr

body-or-break = body
              | break-clause

  seq-expr : sequence?
```

Iteratively evaluates *body*s. The *for-clause*s introduce bindings whose scope includes *body* and that determine the number of times that *body* is evaluated. A *break-clause* either among the *for-clause*s of *body*s stops further iteration.

In the simple case, each *for-clause* has one of its first two forms, where `[`*id seq-expr*`]` is a shorthand for `[(`*id*`)` *seq-expr*`]`. In this simple case, the *seq-expr*s are evaluated left-to-right, and each must produce a sequence value (see §4.14.1 “Sequences”).

The `for` form iterates by drawing an element from each sequence; if any sequence is empty, then the iteration stops, and `#<void>` is the result of the `for` expression. Otherwise a location is created for each *`id`* to hold the values of each element; the sequence produced by a *`seq-expr`* must return as many values for each iteration as corresponding *`id`*s.

The *`id`*s are then bound in the *`body`*, which is evaluated, and whose results are ignored. Iteration continues with the next element in each sequence and with fresh locations for each *`id`*.

A `for` form with zero *`for-clause`*s is equivalent to a single *`for-clause`* that binds an unreferenced *`id`* to a sequence containing a single element. All of the *`id`*s must be distinct according to `bound-identifier=?`.

If any *`for-clause`* has the form `#:when` *`guard-expr`*, then only the preceding clauses (containing no `#:when` or `#:unless`) determine iteration as above, and the *`body`* is effectively wrapped as

```
(when guard-expr
  (for (for-clause ...) body ...+))
```

using the remaining `for-clauses`. A *`for-clause`* of the form `#:unless` *`guard-expr`* corresponds to the same transformation with `unless` in place of `when`.

A `#:break` *`guard-expr`* clause is similar to a `#:unless` *`guard-expr`* clause, but when `#:break` avoids evaluation of the *`body`*s, it also effectively ends all sequences within the `for` form. A `#:final` *`guard-expr`* clause is similar to `#:break` *`guard-expr`*, but instead of immediately ending sequences and skipping the *`body`*s, it allows at most one more element from each later sequence and at most one more evaluation of the following *`body`*s. Among the *`body`*s, besides stopping the iteration and preventing later *`body`* evaluations, a `#:break` *`guard-expr`* or `#:final` *`guard-expr`* clause starts a new internal-definition context.

Examples:

```
> (for ([i '(1 2 3)]
        [j "abc"]
        #:when (odd? i)
        [k #(#t #f)])
    (display (list i j k)))
(1 a #t)(1 a #f)(3 c #t)(3 c #f)

> (for ([(i j) #hash(("a" . 1) ("b" . 20))])
    (display (list i j)))
(a 1)(b 20)

> (for ([i '(1 2 3)]
        [j "abc"]
        #:break (not (odd? i))
```

```
        [k #(#t #f)])
    (display (list i j k)))
(1 a #t)(1 a #f)

> (for ([i '(1 2 3)]
        [j "abc"]
        #:final (not (odd? i))
        [k #(#t #f)])
    (display (list i j k)))
(1 a #t)(1 a #f)(2 b #t)

> (for ([i '(1 2 3)]
        [j "abc"]
        [k #(#t #f)])
    #:break (not (or (odd? i) k))
    (display (list i j k)))
(1 a #t)

> (for ()
    (display "here"))
here

> (for ([i '()])
    (error "doesn't get here"))
```

(for/list (*for-clause* ...) *body-or-break* ... *body*)

Iterates like `for`, but that the last expression in the *body*s must produce a single value, and the result of the `for/list` expression is a list of the results in order. When evaluation of a *body* is skipped due to a `#:when` or `#:unless` clause, the result list includes no corresponding element.

Examples:

```
> (for/list ([i '(1 2 3)]
             [j "abc"]
             #:when (odd? i)
             [k #(#t #f)])
    (list i j k))
'((1 #\a #t) (1 #\a #f) (3 #\c #t) (3 #\c #f))
> (for/list ([i '(1 2 3)]
             [j "abc"]
             #:break (not (odd? i))
             [k #(#t #f)])
    (list i j k))
```

```
'((1 #\a #t) (1 #\a #f))
> (for/list () 'any)
'(any)
> (for/list ([i '()])
    (error "doesn't get here"))
'()
```

```
(for/vector maybe-length (for-clause ...) body-or-break ... body)

maybe-length =
             | #:length length-expr
             | #:length length-expr #:fill fill-expr

  length-expr : exact-nonnegative-integer?
```

Iterates like `for/list`, but results are accumulated into a vector instead of a list.

If the optional `#:length` clause is specified, the result of *length-expr* determines the length of the result vector. In that case, the iteration can be performed more efficiently, and it terminates when the vector is full or the requested number of iterations have been performed, whichever comes first. If *length-expr* specifies a length longer than the number of iterations, then the remaining slots of the vector are initialized to the value of *fill-expr*, which defaults to `0` (i.e., the default argument of `make-vector`).

Examples:

```
> (for/vector ([i '(1 2 3)]) (number->string i))
'#("1" "2" "3")
> (for/vector #:length 2 ([i '(1 2 3)]) (number->string i))
'#("1" "2")
> (for/vector #:length 4 ([i '(1 2 3)]) (number->string i))
'#("1" "2" "3" 0)
> (for/vector #:length 4 #:fill "?" ([i '(1 2 3)]) (number-
>string i))
'#("1" "2" "3" "?")
```

The `for/vector` form may allocate a vector and mutate it after each iteration of *body*, which means that capturing a continuation during *body* and applying it multiple times may mutate a shared vector.

```
(for/hash (for-clause ...) body-or-break ... body)
(for/hasheq (for-clause ...) body-or-break ... body)
(for/hasheqv (for-clause ...) body-or-break ... body)
```

Like `for/list`, but the result is an immutable hash table; `for/hash` creates a table using `equal?` to distinguish keys, `for/hasheq` produces a table using `eq?`, and `for/hasheqv` produces a table using `eqv?`. The last expression in the *body*s must return two values: a key and a value to extend the hash table accumulated by the iteration.

Example:

```
> (for/hash ([i '(1 2 3)])
    (values i (number->string i)))
'#hash((1 . "1") (2 . "2") (3 . "3"))
```

```
(for/and (for-clause ...) body-or-break ... body)
```

Iterates like `for`, but when last expression of *body* produces `#f`, then iteration terminates, and the result of the `for/and` expression is `#f`. If the *body* is never evaluated, then the result of the `for/and` expression is `#t`. Otherwise, the result is the (single) result from the last evaluation of *body*.

Examples:

```
> (for/and ([i '(1 2 3 "x")])
    (i . < . 3))
#f
> (for/and ([i '(1 2 3 4)])
    i)
4
> (for/and ([i '(1 2 3 4)])
    #:break (= i 3)
    i)
2
> (for/and ([i '()])
    (error "doesn't get here"))
#t
```

```
(for/or (for-clause ...) body-or-break ... body)
```

Iterates like `for`, but when last expression of *body* produces a value other than `#f`, then iteration terminates, and the result of the `for/or` expression is the same (single) value. If the *body* is never evaluated, then the result of the `for/or` expression is `#f`. Otherwise, the result is `#f`.

Examples:

```
> (for/or ([i '(1 2 3 "x")])
    (i . < . 3))
#t
> (for/or ([i '(1 2 3 4)])
    i)
1
> (for/or ([i '()])
    (error "doesn't get here"))
#f
```

```
(for/sum (for-clause ...) body-or-break ... body)
```

Iterates like `for`, but each result of the last *body* is accumulated into a result with `+`.

Example:

```
> (for/sum ([i '(1 2 3 4)]) i)
10
```

```
(for/product (for-clause ...) body-or-break ... body)
```

Iterates like `for`, but each result of the last *body* is accumulated into a result with `*`.

Example:

```
> (for/product ([i '(1 2 3 4)]) i)
24
```

```
(for/lists (id ...) (for-clause ...) body-or-break ... body)
```

Similar to `for/list`, but the last *body* expression should produce as many values as given *id*s, and the result is as many lists as supplied *id*s. The *id*s are bound to the lists accumulated so far in the *for-clause*s and *body*s.

Examples:

```
> (for/lists (l1 l2 l3)
             ([i '(1 2 3)]
              [j "abc"]
              #:when (odd? i)
              [k #(#t #f)])
    (values i j k))
'(1 1 3 3)
'(#\a #\a #\c #\c)
'(#t #f #t #f)
> (for/lists (acc)
             ([x '(tvp tofu seitan tvp tofu)]
              #:unless (member x acc))
    x)
'(tvp tofu seitan)
```

```
(for/first (for-clause ...) body-or-break ... body)
```

Iterates like `for`, but after *body* is evaluated the first time, then the iteration terminates, and the `for/first` result is the (single) result of *body*. If the *body* is never evaluated, then the result of the `for/first` expression is `#f`.

Examples:

```
> (for/first ([i '(1 2 3 "x")]
              #:when (even? i))
    (number->string i))
"2"
> (for/first ([i '()])
    (error "doesn't get here"))
#f
```

```
(for/last (for-clause ...) body-or-break ... body)
```

Iterates like `for`, but the `for/last` result is the (single) result of the last evaluation of *body*. If the *body* is never evaluated, then the result of the `for/last` expression is `#f`.

Examples:

```
> (for/last ([i '(1 2 3 4 5)]
             #:when (even? i))
    (number->string i))
"4"
> (for/last ([i '()])
    (error "doesn't get here"))
#f
```

```
(for/fold ([accum-id init-expr] ...) (for-clause ...)
  body-or-break ... body)
```

Iterates like `for`. Before iteration starts, the *init-expr*s are evaluated to produce initial accumulator values. At the start of each iteration, a location is generated for each *accum-id*, and the corresponding current accumulator value is placed into the location. The last expression in *body* must produce as many values as *accum-id*s, and those values become the current accumulator values. When iteration terminates, the results of the `for/fold` expression are the accumulator values.

Example:

```
> (for/fold ([sum 0]
             [rev-roots null])
            ([i '(1 2 3 4)])
    (values (+ sum i) (cons (sqrt i) rev-roots)))
10
'(2 1.7320508075688772 1.4142135623730951 1)
```

```
(for* (for-clause ...) body-or-break ... body)
```

Like `for`, but with an implicit `#:when #t` between each pair of `for-clauses`, so that all sequence iterations are nested.

Example:

```
> (for* ([i '(1 2)]
         [j "ab"])
    (display (list i j)))
(1 a)(1 b)(2 a)(2 b)
```

```
(for*/list (for-clause ...) body-or-break ... body)
(for*/lists (id ...) (for-clause ...) body-or-break ... body)
(for*/vector maybe-length (for-clause ...) body-or-break ... body)
(for*/hash (for-clause ...) body-or-break ... body)
(for*/hasheq (for-clause ...) body-or-break ... body)
(for*/hasheqv (for-clause ...) body-or-break ... body)
(for*/and (for-clause ...) body-or-break ... body)
(for*/or (for-clause ...) body-or-break ... body)
(for*/sum (for-clause ...) body-or-break ... body)
(for*/product (for-clause ...) body-or-break ... body)
(for*/first (for-clause ...) body-or-break ... body)
(for*/last (for-clause ...) body-or-break ... body)
(for*/fold ([accum-id init-expr] ...) (for-clause ...)
  body-or-break ... body)
```

Like `for/list`, etc., but with the implicit nesting of `for*`.

Example:

```
> (for*/list ([i '(1 2)]
              [j "ab"])
    (list i j))
'((1 #\a) (1 #\b) (2 #\a) (2 #\b))
```

3.18.2 Deriving New Iteration Forms

```
(for/fold/derived orig-datum
  ([accum-id init-expr] ...) (for-clause ...)
  body-or-break ... body)
```

Like `for/fold`, but the extra *orig-datum* is used as the source for all syntax errors.

Examples:

```
> (define-syntax (for/digits stx)
    (syntax-case stx ()
      [(_ clauses . defs+exprs)
```

```
       (with-syntax ([original stx])
         #'(let-values
             ([(n k)
               (for/fold/derived original ([n 0] [k 1]) clauses
                 (define d (let () . defs+exprs))
                 (values (+ n (* d k)) (* k 10)))])
             n))]))

; If we misuse for/digits, we can get good error reporting
; because the use of orig-datum allows for source correlation:
> (for/digits
      [a (in-list '(1 2 3))]
      [b (in-list '(4 5 6))]
    (+ a b))
eval:3:0: for/digits: bad sequence binding clause
  at: a
  in: (for/digits (a (in-list (quote (1 2 3)))) (b (in-list
(quote (4 5 6)))) (+ a b))
> (for/digits
      ([a (in-list '(1 2 3))]
       [b (in-list '(2 4 6))])
    (+ a b))
963
; Another example: compute the max during iteration:
> (define-syntax (for/max stx)
    (syntax-case stx ()
      [(_ clauses . defs+exprs)
       (with-syntax ([original stx])
         #'(for/fold/derived original
                             ([current-max -inf.0])
                             clauses
             (define maybe-new-max
               (let () . defs+exprs))
             (if (> maybe-new-max current-max)
                 maybe-new-max
                 current-max)))]))

> (for/max ([n '(3.14159 2.71828 1.61803)]
            [s '(-1      1       1)])
    (* n s))
2.71828
(for*/fold/derived orig-datum
  ([accum-id init-expr] ...) (for-clause ...)
  body-or-break ... body)
```

Like `for*/fold`, but the extra *`orig-datum`* is used as the source for all syntax errors.

Examples:

```
> (define-syntax (for*/digits stx)
    (syntax-case stx ()
      [(_ clauses . defs+exprs)
       (with-syntax ([original stx])
         #'(let-values
             ([(n k)
               (for*/fold/derived original ([n 0] [k 1]) clauses
                 (define d (let () . defs+exprs))
                 (values (+ n (* d k)) (* k 10)))])
             n))]))

> (for*/digits
      [ds (in-list '((8 3) (1 1)))]
      [d (in-list ds)]
    d)
eval:8:0: for*/digits: bad sequence binding clause
  at: ds
  in: (for*/digits (ds (in-list (quote ((8 3) (1 1))))) (d
(in-list ds)) d)
> (for*/digits
      ([ds (in-list '((8 3) (1 1)))]
       [d (in-list ds)])
    d)
1138
```

```
(define-sequence-syntax id
  expr-transform-expr
  clause-transform-expr)

                          (or/c (-> identifier?)
  expr-transform-expr :         (syntax? . -> . syntax?))

  clause-transform-expr : (syntax? . -> . syntax?)
```

Defines *id* as syntax. An (*id* . *rest*) form is treated specially when used to generate a sequence in a *clause* of `for` (or one of its variants). In that case, the procedure result of *clause-transform-expr* is called to transform the clause.

When *id* is used in any other expression position, the result of *expr-transform-expr* is used. If it is a procedure of zero arguments, then the result must be an identifier *other-id*, and any use of *id* is converted to a use of *other-id*. Otherwise,*expr-transform-expr* must produce a procedure (of one argument) that is used as a macro transformer.

When the *clause-transform-expr* transformer is used, it is given a *clause* as an argument, where the clause's form is normalized so that the left-hand side is a parenthesized

sequence of identifiers. The right-hand side is of the form `(id . rest)`. The result can be either `#f`, to indicate that the forms should not be treated specially (perhaps because the number of bound identifiers is inconsistent with the `(id . rest)` form), or a new *clause* to replace the given one. The new clause might use `:do-in`. To protect identifiers in the result of *clause-transform-expr*, use `for-clause-syntax-protect` instead of `syntax-protect`.

Examples:

```
> (define-sequence-syntax in-digits
    (lambda () #'in-digits/proc)
    (lambda (stx)
      (syntax-case stx ()
        [[(d) (_ nat)]
         #'[(d)
            (:do-in
              ([(n) nat])
              (unless (exact-nonnegative-integer? n)
                (raise-type-error 'in-digits "exact non-negative
integer" n))
              ([i n])
              (not (zero? i))
              ([(j d) (quotient/remainder i 10)])
              #true
              #true
              [j])]])))

> (define (in-digits/proc n [b 10])
    (for/list ([d (in-digits n)]) d))

> (for/list ([d (in-digits 1138)]) d)
'(8 3 1 1)
> (map in-digits (list 137 216))
'((7 3 1) (6 1 2))
```

```
(:do-in ([(outer-id ...) outer-expr] ...)
        outer-check
        ([loop-id loop-expr] ...)
        pos-guard
        ([(inner-id ...) inner-expr] ...)
        pre-guard
        post-guard
        (loop-arg ...))
```

A form that can only be used as a *seq-expr* in a *clause* of `for` (or one of its variants).

Within a `for`, the pieces of the `:do-in` form are spliced into the iteration essentially as follows:

```
(let-values ([(outer-id ...) outer-expr] ...)
  outer-check
  (let loop ([loop-id loop-expr] ...)
    (if pos-guard
        (let-values ([(inner-id ...) inner-expr] ...)
          (if pre-guard
              (let body-bindings
                   (if post-guard
                       (loop loop-arg ...)
                       done-expr))
              done-expr))
         done-expr)))
```

where *body-bindings* and *done-expr* are from the context of the `:do-in` use. The identifiers bound by the `for` clause are typically part of the `([(inner-id ...) inner-expr] ...)` section.

The actual `loop` binding and call has additional loop arguments to support iterations in parallel with the `:do-in` form, and the other pieces are similarly accompanied by pieces from parallel iterations.

For an example of `:do-in`, see `define-sequence-syntax`.

```
(for-clause-syntax-protect stx) → syntax?
  stx : syntax?
```

Provided `for-syntax`: Like `syntax-protect`, but allows the `for` expander to disarm the result syntax object, and arms the pieces of a clause instead of the entire syntax object.

Use this function to protect the result of a *clause-transform-expr* that is bound by `define-sequence-syntax`.

3.18.3 Do Loops

```
(do ([id init-expr step-expr-maybe] ...)
    (stop?-expr finish-expr ...)
  expr ...)

step-expr-maybe =
                | step-expr
```

Iteratively evaluates the *expr*s for as long as *stop?-expr* returns `#f`.

To initialize the loop, the *init-expr*s are evaluated in order and bound to the corresponding *id*s. The *id*s are bound in all expressions within the form other than the *init-expr*s.

After the *id*s have been bound, the *stop?-expr* is evaluated. If it produces `#f`, each *expr* is evaluated for its side-effect. The *id*s are then effectively updated with the values of the *step-expr*s, where the default *step-expr* for *id* is just *id*; more precisely, iteration continues with fresh locations for the *id*s that are initialized with the values of the corresponding *step-expr*s.

When *stop?-expr* produces a true value, then the *finish-expr*s are evaluated in order, and the last one is evaluated in tail position to produce the overall value for the `do` form. If no *finish-expr* is provided, the value of the `do` form is `#<void>`.

3.19 Continuation Marks: `with-continuation-mark`

```
(with-continuation-mark key-expr val-expr result-expr)
```

The *key-expr*, *val-expr*, and *result-expr* expressions are evaluated in order. After *key-expr* is evaluated to obtain a key and *val-expr* is evaluated to obtain a value, the key is mapped to the value as a continuation mark in the current continuation's initial continuation frame. If the frame already has a mark for the key, the mark is replaced. Finally, the *result-expr* is evaluated; the continuation for evaluating *result-expr* is the continuation of the `with-continuation-mark` expression (so the result of the *result-expr* is the result of the `with-continuation-mark` expression, and *result-expr* is in tail position for the `with-continuation-mark` expression).

§10.5 "Continuation Marks" provides more information on continuation marks.

3.20 Quasiquoting: `quasiquote`, `unquote`, and `unquote-splicing`

```
(quasiquote datum)
```

The same as `'datum` if *datum* does not include `(unquote expr)` or `(unquote-splicing expr)`. An `(unquote expr)` form escapes from the quote, however, and the result of the *expr* takes the place of the `(unquote expr)` form in the `quasiquote` result. An `(unquote-splicing expr)` similarly escapes, but the *expr* must produce a list, and its elements are spliced as multiple values place of the `(unquote-splicing expr)`, which must appear as the `car` or a quoted pair, as an element of a quoted vector, or as an element of a quoted prefab structure; in the case of a pair, if the `cdr` of the relevant quoted pair is empty, then *expr* need not produce a list, and its result is used directly in place of the quoted pair (in the same way that `append` accepts a non-list final argument). In a quoted hash table, an `(unquote expr)` or `(unquote-splicing expr)` expression escapes only in the second element of an entry pair (i.e., the value), while entry keys are always implicitly quoted. If `unquote` or `unquote-splicing` appears within `quasiquote` in any other way than as `(unquote expr)` or `(unquote-splicing expr)`, a syntax error is reported.

§4.11 "Quasiquoting: `quasiquote` and `‘`" in ***The Racket Guide*** introduces `quasiquote`.

Examples:

```
> (quasiquote (0 1 2))
'(0 1 2)
> (quasiquote (0 (unquote (+ 1 2)) 4))
'(0 3 4)
> (quasiquote (0 (unquote-splicing (list 1 2)) 4))
'(0 1 2 4)
> (quasiquote (0 (unquote-splicing 1) 4))
unquote-splicing: contract violation
  expected: list?
  given: 1
> (quasiquote (0 (unquote-splicing 1)))
'(0 . 1)
```

A `quasiquote`, `unquote`, or `unquote-splicing` form is typically abbreviated with `` ` ``, `,`, or `,@`, respectively. See also §1.3.8 “Reading Quotes”.

Examples:

```
> `(0 1 2)
'(0 1 2)
> `(1 ,(+ 1 2) 4)
'(1 3 4)
> `#s(stuff 1 ,(+ 1 2) 4)
'#s(stuff 1 3 4)
> `#hash(("a" . ,(+ 1 2)))
'#hash(("a" . 3))
> `#hash((,(+ 1 2) . "a"))
'#hash((,(+ 1 2) . "a"))
> `(1 ,@(list 1 2) 4)
'(1 1 2 4)
> `#(1 ,@(list 1 2) 4)
'#(1 1 2 4)
```

A `quasiquote` form within the original *datum* increments the level of quasiquotation: within the `quasiquote` form, each `unquote` or `unquote-splicing` is preserved, but a further nested `unquote` or `unquote-splicing` escapes. Multiple nestings of `quasiquote` require multiple nestings of `unquote` or `unquote-splicing` to escape.

Examples:

```
> `(1 `,(+ 1 ,(+ 2 3)) 4)
'(1 `,(+ 1 5) 4)
> `(1 ```,,@,,@(list (+ 1 2)) 4)
'(1 ```,,@,3 4)
```

The `quasiquote` form allocates only as many fresh cons cells, vectors, and boxes as are needed without analyzing `unquote` and `unquote-splicing` expressions. For example, in

```
`(,1 2 3)
```

a single tail `'(2 3)` is used for every evaluation of the `quasiquote` expression.

`unquote`

See `quasiquote`, where `unquote` is recognized as an escape. An `unquote` form as an expression is a syntax error.

`unquote-splicing`

See `quasiquote`, where `unquote-splicing` is recognized as an escape. An `unquote-splicing` form as an expression is a syntax error.

3.21 Syntax Quoting: `quote-syntax`

`(quote-syntax datum)`

Similar to `quote`, but produces a syntax object that preserves the lexical information and source-location information attached to *datum* at expansion time.

Unlike `syntax` (`#'`), `quote-syntax` does not substitute pattern variables bound by `with-syntax`, `syntax-parse`, or `syntax-case`.

Examples:

```
> (syntax? (quote-syntax x))
#t
> (quote-syntax (1 2 3))
#<syntax:72:0 (1 2 3)>
> (with-syntax ([a #'5])
    (quote-syntax (a b c)))
#<syntax:73:0 (a b c)>
```

3.22 Interaction Wrapper: `#%top-interaction`

`(#%top-interaction . form)`

Expands to simply *form*. The `#%top-interaction` form is similar to `#%app` and `#%module-begin`, in that it provides a hook to control interactive evaluation through `load` (more precisely, the default load handler) or `read-eval-print-loop`.

3.23 Blocks: `block`

```
(require racket/block)
```
package: base

The bindings documented in this section are provided by the `racket/block` library, not `racket/base` or `racket`.

```
(block defn-or-expr ...)
```

Supports a mixture of expressions and mutually recursive definitions, as in a `module` body. Unlike an internal-definition context, the last *defn-or-expr* need not be an expression.

The result of the `block` form is the result of the last *defn-or-expr* if it is an expression, `#<void>` otherwise. If no *defn-or-expr* is provided (after flattening `begin` forms), the result is `#<void>`.

The final *defn-or-expr* is executed in tail position, if it is an expression.

Examples:

```
> (define (f x)
    (block
      (define y (add1 x))
      (displayln y)
      (define z (* 2 y))
      (+ 3 z)))

> (f 12)
13
29
```

3.24 Internal-Definition Limiting: `#%stratified-body`

```
(#%stratified-body defn-or-expr ...)
```

Like `(let () defn-or-expr ...)` for an internal-definition context sequence, except that an expression is not allowed to precede a definition, and all definitions are treated as referring to all other definitions (i.e., locations for variables are all allocated first, like `letrec` and unlike `letrec-syntaxes+values`).

The `#%stratified-body` form is useful for implementing syntactic forms or languages that supply a more limited kind of internal-definition context.

3.25 Performance Hints: `begin-encourage-inline`

```
(require racket/performance-hint)
```
package: base

The bindings documented in this section are provided by the `racket/performance-hint` library, not `racket/base` or `racket`.

```
(begin-encourage-inline form ...)
```

Attaches a `'compiler-hint:cross-module-inline` syntax property to each *form*, which is useful when a *form* is a function definition. See `define-values`.

```
(define-inline id expr)
(define-inline (head args) body ...+)

head = id
     | (head args)

args = arg ...
     | arg ... . rest-id

 arg = arg-id
     | [arg-id default-expr]
     | keyword arg-id
     | keyword [arg-id default-expr]
```

Like `define`, but ensures that the definition will be inlined at its call sites. Recursive calls are not inlined, to avoid infinite inlining. Higher-order uses are supported, but also not inlined.

`define-inline` may interfere with the Racket compiler's own inlining heuristics, and should only be used when other inlining attempts (such as `begin-encourage-inline`) fail.

3.26 Importing Modules Lazily: `lazy-require`

```
(require racket/lazy-require)
```
package: base

The bindings documented in this section are provided by the `racket/lazy-require` library, not `racket/base` or `racket`.

```
(lazy-require [module-path (fun-import ...)] ...)

fun-import = fun-id
           | (orig-fun-id fun-id)
```

Defines each *fun-id* as a function that, when called, dynamically requires the export named *orig-fun-id* from the module specified by *module-path* and calls it with the same arguments. If *orig-fun-id* is not given, it defaults to *fun-id*.

If the enclosing relative phase level is not 0, then *module-path* is also placed in a submodule (with a use of `define-runtime-module-path-index` at phase level 0 within the submodule). Introduced submodules have the names `lazy-require-`*n*`-`*m*, where *n* is a phase-level number and *m* is a number.

When the use of a lazily-required function triggers module loading, `register-external-module` declares a potential compilation dependency (in case the function is used in the process of compiling a module).

Examples:

```
> (lazy-require
    [racket/list (partition)])

> (partition even? '(1 2 3 4 5))
'(2 4)
'(1 3 5)
> (module hello racket/base
    (provide hello)
    (printf "starting hello server\n")
    (define (hello) (printf "hello!\n")))

> (lazy-require
    ['hello ([hello greet])])

> (greet)
starting hello server
hello!
```

4 Datatypes

§3 “Built-In Datatypes” in ***The Racket Guide*** introduces Datatypes.

Each pre-defined datatype comes with a set of procedures for manipulating instances of the datatype.

4.1 Booleans and Equality

True and false *booleans* are represented by the values `#t` and `#f`, respectively, though operations that depend on a boolean value typically treat anything other than `#f` as true. The `#t` value is always `eq?` to itself, and `#f` is always `eq?` to itself.

See §1.3.5 “Reading Booleans” for information on `read`ing booleans and §1.4.4 “Printing Booleans” for information on `print`ing booleans.

See also `and`, `or`, `andmap`, and `ormap`.

```
(boolean? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is `#t` or `#f`, `#f` otherwise.

Examples:

```
> (boolean? #f)
#t
> (boolean? #t)
#t
> (boolean? 'true)
#f
```

```
(not v) → boolean?
  v : any/c
```

Returns `#t` if `v` is `#f`, `#f` otherwise.

Examples:

```
> (not #f)
#t
> (not #t)
#f
> (not 'we-have-no-bananas)
#f
```

```
(equal? v1 v2) → boolean?
  v1 : any/c
  v2 : any/c
```

Two values are `equal?` if and only if they are `eqv?`, unless otherwise specified for a particular datatype.

Datatypes with further specification of `equal?` include strings, byte strings, pairs, mutable pairs, vectors, boxes, hash tables, and inspectable structures. In the last six cases, equality is recursively defined; if both *v1* and *v2* contain reference cycles, they are equal when the infinite unfoldings of the values would be equal. See also `gen:equal+hash` and `prop:impersonator-of`.

Examples:

```
> (equal? 'yes 'yes)
#t
> (equal? 'yes 'no)
#f
> (equal? (expt 2 100) (expt 2 100))
#t
> (equal? 2 2.0)
#f
> (equal? (make-string 3 #\z) (make-string 3 #\z))
#t
```

```
(eqv? v1 v2) → boolean?
  v1 : any/c
  v2 : any/c
```

Two values are `eqv?` if and only if they are `eq?`, unless otherwise specified for a particular datatype.

The number and character datatypes are the only ones for which `eqv?` differs from `eq?`.

Examples:

```
> (eqv? 'yes 'yes)
#t
> (eqv? 'yes 'no)
#f
> (eqv? (expt 2 100) (expt 2 100))
#t
> (eqv? 2 2.0)
#f
> (eqv? (integer->char 955) (integer->char 955))
```

```
#t
> (eqv? (make-string 3 #\z) (make-string 3 #\z))
#f
```

```
(eq? v1 v2) → boolean?
  v1 : any/c
  v2 : any/c
```

Return `#t` if *v1* and *v2* refer to the same object, `#f` otherwise. See also §1.1.6 “Object Identity and Comparisons”.

Examples:

```
> (eq? 'yes 'yes)
#t
> (eq? 'yes 'no)
#f
> (let ([v (mcons 1 2)]) (eq? v v))
#t
> (eq? (mcons 1 2) (mcons 1 2))
#f
> (eq? (make-string 3 #\z) (make-string 3 #\z))
#f
```

```
(equal?/recur v1 v2 recur-proc) → boolean?
  v1 : any/c
  v2 : any/c
  recur-proc : (any/c any/c -> any/c)
```

Like `equal?`, but using *recur-proc* for recursive comparisons (which means that reference cycles are not handled automatically). Non-`#f` results from *recur-proc* are converted to `#t` before being returned by `equal?/recur`.

Examples:

```
> (equal?/recur 1 1 (lambda (a b) #f))
#t
> (equal?/recur '(1) '(1) (lambda (a b) #f))
#f
> (equal?/recur '#(1 1 1) '#(1 1.2 3/4)
                (lambda (a b) (<= (abs (- a b)) 0.25)))
#t
```

```
(immutable? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is an immutable string, byte string, vector, hash table, or box, `#f` otherwise.

Note that immutable? is not a general predicate for immutability (despite its name). It works only for a handful of datatypes for which a single predicate—string?, vector?, etc.—recognizes both mutable and immutable variants of the datatype. In particular, immutable? produces #f for a pair, even though pairs are immutable, since pair? implies immutability.

Examples:

```
> (immutable? 'hello)
#f
> (immutable? "a string")
#t
> (immutable? (box 5))
#f
> (immutable? #(0 1 2 3))
#t
> (immutable? (make-hash))
#f
> (immutable? (make-immutable-hash '([a b])))
#t
```

gen:equal+hash : any/c

A generic interface (see §5.4 “Generic Interfaces”) that supplies an equality predicate and hashing functions for a structure type. The following methods must be implemented:

- *equal-proc* : (-> any/c any/c (-> any/c any/c boolean?) any/c) — tests whether the first two arguments are equal, where both values are instances of the structure type to which the generic interface is associated (or a subtype of the structure type).

  The third argument is an equal? predicate to use for recursive equality checks; use the given predicate instead of equal? to ensure that data cycles are handled properly and to work with equal?/recur (but beware that an arbitrary function can be provided to equal?/recur for recursive checks, which means that arguments provided to the predicate might be exposed to arbitrary code).

  The *equal-proc* is called for a pair of structures only when they are not eq?, and only when they both have a gen:equal+hash value inherited from the same structure type. With this strategy, the order in which equal? receives two structures does not matter. It also means that, by default, a structure sub-type inherits the equality predicate of its parent, if any.

- *hash-proc* : (-> any/c (-> any/c exact-integer?) exact-integer?) — computes a hash code for the given structure, like equal-hash-code. The first argument is an instance of the structure type (or one of its subtypes) to which the generic interface is associated.

The second argument is an equal-hash-code-like procedure to use for recursive hash-code computation; use the given procedure instead of equal-hash-code to ensure that data cycles are handled properly.

- *hash2-proc* : (-> any/c (-> any/c exact-integer?) exact-integer?) — computes a secondary hash code for the given structure. This procedure is like *hash-proc*, but analogous to equal-secondary-hash-code.

Take care to ensure that *hash-proc* and *hash2-proc* are consistent with *equal-proc*. Specifically, *hash-proc* and *hash2-proc* should produce the same value for any two structures for which *equal-proc* produces a true value.

When a structure type has no gen:equal+hash implementation, then transparent structures (i.e., structures with an inspector that is controlled by the current inspector) are equal? when they are instances of the same structure type (not counting sub-types), and when they have equal? field values. For transparent structures, equal-hash-code and equal-secondary-hash-code derive hash code using the field values. For opaque structure types, equal? is the same as eq?, and equal-hash-code and equal-secondary-hash-code results are based only on eq-hash-code. If a structure has a prop:impersonator-of property, then the prop:impersonator-of property takes precedence over gen:equal+hash if the property value's procedure returns a non-#f value when applied to the structure.

Examples:

```
> (define (farm=? farm1 farm2 recursive-equal?)
    (and (= (farm-apples farm1)
            (farm-apples farm2))
         (= (farm-oranges farm1)
            (farm-oranges farm2))
         (= (farm-sheep farm1)
            (farm-sheep farm2))))

> (define (farm-hash-1 farm recursive-equal-hash)
    (+ (* 10000 (farm-apples farm))
       (* 100 (farm-oranges farm))
       (* 1 (farm-sheep farm))))

> (define (farm-hash-2 farm recursive-equal-hash)
    (+ (* 10000 (farm-sheep farm))
       (* 100 (farm-apples farm))
       (* 1 (farm-oranges farm))))

> (define-struct farm (apples oranges sheep)
                 #:methods gen:equal+hash
                 [(define equal-proc farm=?)
                  (define hash-proc  farm-hash-1)
                  (define hash2-proc farm-hash-2)])
```

```
> (define east (make-farm 5 2 20))

> (define west (make-farm 18 6 14))

> (define north (make-farm 5 20 20))

> (define south (make-farm 18 6 14))

> (equal? east west)
#f
> (equal? east north)
#f
> (equal? west south)
#t
```

prop:equal+hash : struct-type-property?

A deprecated structure type property (see §5.3 “Structure Type Properties”) that supplies an equality predicate and hashing functions for a structure type. The gen:equal+hash generic interface should be used, instead. A prop:equal+hash property value is a list of three procedures that correspond to the methods of gen:equal+hash.

4.1.1 Boolean Aliases

(require racket/bool) package: base

The bindings documented in this section are provided by the racket/bool and racket libraries, but not racket/base.

true : boolean?

An alias for #t.

false : boolean?

An alias for #f.

(symbol=? a b) → boolean?
  a : symbol?
  b : symbol?

Returns (equal? a b) (if a and b are symbols).

```
(boolean=? a b) → boolean?
  a : boolean?
  b : boolean?
```

Returns `(equal? a b)` (if `a` and `b` are booleans).

```
(false? v) → boolean?
  v : any/c
```

Returns `(not v)`.

```
(nand expr ...)
```

Same as `(not (and expr ...))`.

Examples:

```
> (nand #f #t)
#t
> (nand #f (error 'ack "we don't get here"))
#t
```

```
(nor expr ...)
```

Same as `(not (or expr ...))`.

In the two argument case, returns `#t` if neither of the arguments is a true value.

Examples:

```
> (nor #f #t)
#f
> (nor #t (error 'ack "we don't get here"))
#f
```

```
(implies expr1 expr2)
```

Checks to be sure that the first expression implies the second.

Same as `(if expr1 expr2 #t)`.

Examples:

```
> (implies #f #t)
#t
```

```
> (implies #f #f)
#t
> (implies #t #f)
#f
> (implies #f (error 'ack "we don't get here"))
#t
```

```
(xor b1 b2) → any
  b1 : any/c
  b2 : any/c
```

Returns the exclusive or of *b1* and *b2*.

If exactly one of *b1* and *b2* is not `#f`, then return it. Otherwise, returns `#f`.

Examples:

```
> (xor 11 #f)
11
> (xor #f 22)
22
> (xor 11 22)
#f
> (xor #f #f)
#f
```

4.2 Numbers

§3.2 “Numbers” in ***The Racket Guide*** introduces numbers.

All *numbers* are *complex numbers*. Some of them are *real numbers*, and all of the real numbers that can be represented are also *rational numbers*, except for `+inf.0` (positive infinity), `+inf.f` (single-precision variant), `-inf.0` (negative infinity), `-inf.f` (single-precision variant), `+nan.0` (not-a-number), and `+nan.f` (single-precision variant). Among the rational numbers, some are *integers*, because `round` applied to the number produces the same number.

See §1.3.3 “Reading Numbers” for information on the syntax of number literals.

Orthogonal to those categories, each number is also either an *exact number* or an *inexact number*. Unless otherwise specified, computations that involve an inexact number produce inexact results. Certain operations on inexact numbers, however, produce an exact number, such as multiplying an inexact number with an exact `0`. Some operations, which can produce an irrational number for rational arguments (e.g., `sqrt`), may produce inexact results even for exact arguments.

In the case of complex numbers, either the real and imaginary parts are both exact or inexact, or the number has an exact zero real part and an inexact imaginary part; a complex number with an exact zero imaginary part is a real number.

Inexact real numbers are implemented as either single- or double-precision IEEE floating-point numbers—the latter by default, and the former only when a computation starts with numerical constants specified as single-precision numbers. Inexact real numbers that are represented as double-precision floating-point numbers are *flonums*.

The precision and size of exact numbers is limited only by available memory (and the precision of operations that can produce irrational numbers). In particular, adding, multiplying, subtracting, and dividing exact numbers always produces an exact result.

Inexact numbers can be coerced to exact form, except for the inexact numbers `+inf.0`, `+inf.f`, `-inf.0`, `-inf.f`, `+nan.0`, and `+nan.f`, which have no exact form. Dividing a number by exact zero raises an exception; dividing a non-zero number other than `+nan.0` or `+nan.f` by an inexact zero returns `+inf.0`, `+inf.f`, `-inf.0` or `-inf.f`, depending on the sign and precision of the dividend. The `+nan.0` value is not `=` to itself, but `+nan.0` is `eqv?` to itself, and `+nan.f` is similarly `eqv?` but not `=` to itself. Conversely, `(= 0.0 -0.0)` is `#t`, but `(eqv? 0.0 -0.0)` is `#f`, and the same for `0.0f0` and `-0.0f0` (which are single-precision variants). The datum `-nan.0` refers to the same constant as `+nan.0`, and `-nan.f` is the same as `+nan.f`.

Calculations with infinites produce results consistent with IEEE double- or single-precision floating point where IEEE specifies the result; in cases where IEEE provides no specification, the result corresponds to the limit approaching infinity, or `+nan.0` or `+nan.f` if no such limit exists.

A *fixnum* is an exact integer whose two's complement representation fit into 31 bits on a 32-bit platform or 63 bits on a 64-bit platform; furthermore, no allocation is required when computing with fixnums. See also the `racket/fixnum` module, below.

Two fixnums that are `=` are also the same according to `eq?`. Otherwise, the result of `eq?` applied to two numbers is undefined, except that numbers produced by the default reader in `read-syntax` mode are interned and therefore `eq?` when they are `eqv?`.

Two numbers are `eqv?` when they are both inexact with the same precision or both exact, and when they are `=` (except for `+nan.0`, `+nan.f`, `0.0`, `0.0f0`, `-0.0`, and `-0.0f0`, as noted above). Two numbers are `equal?` when they are `eqv?`.

See §1.3.3 "Reading Numbers" for information on `read`ing numbers and §1.4.2 "Printing Numbers" for information on `print`ing numbers.

4.2.1 Number Types

```
(number? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a number, `#f` otherwise.

Examples:

```
> (number? 1)
#t
> (number? 2+3i)
#t
> (number? "hello")
#f
```

```
(complex? v) → boolean?
  v : any/c
```

Returns `(number? v)`, because all numbers are complex numbers.

```
(real? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a real number, `#f` otherwise.

Examples:

```
> (real? 1)
#t
> (real? +inf.0)
#t
> (real? 2+3i)
#f
> (real? 2.0+0.0i)
#f
> (real? "hello")
#f
```

```
(rational? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a rational number, `#f` otherwise.

Examples:

```
> (rational? 1)
#t
> (rational? +inf.0)
#f
> (rational? "hello")
#f
```

```
(integer? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a number that is an integer, `#f` otherwise.

Examples:

```
> (integer? 1)
#t
> (integer? 2.3)
#f
> (integer? 4.0)
#t
> (integer? +inf.0)
#f
> (integer? 2+3i)
#f
> (integer? "hello")
#f
```

```
(exact-integer? v) → boolean?
  v : any/c
```

Returns `(and (integer? v) (exact? v))`.

Examples:

```
> (exact-integer? 1)
#t
> (exact-integer? 4.0)
#f
```

```
(exact-nonnegative-integer? v) → boolean?
  v : any/c
```

Returns `(and (exact-integer? v) (not (negative? v)))`.

Examples:

```
> (exact-nonnegative-integer? 0)
#t
> (exact-nonnegative-integer? -1)
#f
```

```
(exact-positive-integer? v) → boolean?
  v : any/c
```

Returns `(and (exact-integer? v) (positive? v))`.

Examples:

```
> (exact-positive-integer? 1)
#t
> (exact-positive-integer? 0)
#f
```

```
(inexact-real? v) → boolean?
  v : any/c
```

Returns `(and (real? v) (inexact? v))`.

```
(fixnum? v) → boolean?
  v : any/c
```

Return `#t` if `v` is a fixnum, `#f` otherwise.

Note: the result of this function is platform-dependent, so using it in syntax transformers can lead to platform-dependent bytecode files.

```
(flonum? v) → boolean?
  v : any/c
```

Return `#t` if `v` is a flonum, `#f` otherwise.

```
(double-flonum? v) → boolean?
  v : any/c
```

Identical to `flonum?`.

```
(single-flonum? v) → boolean?
  v : any/c
```

Return `#t` if `v` is a single-precision floating-point number, `#f` otherwise.

```
(zero? z) → boolean?
  z : number?
```

Returns `(= 0 z)`.

Examples:

```
> (zero? 0)
#t
> (zero? -0.0)
#t
```

```
(positive? x) → boolean?
  x : real?
```

Returns `(> x 0)`.

Examples:

```
> (positive? 10)
#t
> (positive? -10)
#f
> (positive? 0.0)
#f
```

```
(negative? x) → boolean?
  x : real?
```

Returns `(< x 0)`.

Examples:

```
> (negative? 10)
#f
> (negative? -10)
#t
> (negative? -0.0)
#f
```

```
(even? n) → boolean?
  n : integer?
```

Returns `(zero? (modulo n 2))`.

Examples:

```
> (even? 10.0)
#t
> (even? 11)
#f
> (even? +inf.0)
even?: contract violation
  expected: integer
  given: +inf.0
```

```
(odd? n) → boolean?
  n : integer?
```

Returns `(not (even? n))`.

Examples:

```
> (odd? 10.0)
#f
> (odd? 11)
#t
> (odd? +inf.0)
odd?: contract violation
  expected: integer
  given: +inf.0
```

```
(exact? z) → boolean?
  z : number?
```

Returns `#t` if `z` is an exact number, `#f` otherwise.

Examples:

```
> (exact? 1)
#t
> (exact? 1.0)
#f
```

```
(inexact? z) → boolean?
  z : number?
```

Returns `#t` if `z` is an inexact number, `#f` otherwise.

Examples:

```
> (inexact? 1)
#f
> (inexact? 1.0)
#t
```

```
(inexact->exact z) → exact?
  z : number?
```

Coerces `z` to an exact number. If `z` is already exact, it is returned. If `z` is `+inf.0`, `-inf.0`, or `+nan.0`, then the `exn:fail:contract` exception is raised.

Examples:

```
> (inexact->exact 1)
1
> (inexact->exact 1.0)
1
```

```
(exact->inexact z) → inexact?
  z : number?
```

Coerces `z` to an inexact number. If `z` is already inexact, it is returned.

Examples:

```
> (exact->inexact 1)
1.0
> (exact->inexact 1.0)
1.0
```

```
(real->single-flonum x) → single-flonum?
  x : real?
```

Coerces `x` to a single-precision floating-point number. If `x` is already a single-precision floating-point number, it is returned.

```
(real->double-flonum x) → flonum?
  x : real?
```

Coerces `x` to a double-precision floating-point number. If `x` is already a double-precision floating-point number, it is returned.

4.2.2 Generic Numerics

Most Racket numeric operations work on any kind of number.

Arithmetic

```
(+ z ...) → number?
  z : number?
```

Returns the sum of the `z`s, adding pairwise from left to right. If no arguments are provided, the result is `0`.

Examples:

```
> (+ 1 2)
3
> (+ 1.0 2+3i 5)
8.0+3.0i
> (+)
0
```

```
(- z) → number?
  z : number?
(- z w ...+) → number?
  z : number?
  w : number?
```

When no `w`s are supplied, returns `(- 0 z)`. Otherwise, returns the subtraction of the `w`s from `z` working pairwise from left to right.

Examples:

```
> (- 5 3.0)
2.0
> (- 1)
-1
> (- 2+7i 1 3)
-2+7i
```

```
(* z ...) → number?
  z : number?
```

Returns the product of the `z`s, multiplying pairwise from left to right. If no arguments are provided, the result is `1`. Multiplying any number by exact `0` produces exact `0`.

Examples:

```
> (* 2 3)
6
> (* 8.0 9)
72.0
> (* 1+2i 3+4i)
-5+10i
```

```
(/ z) → number?
  z : number?
(/ z w ...+) → number?
  z : number?
  w : number?
```

When no `w`s are supplied, returns `(/ 1 z)`. Otherwise, returns the division of `z` by the `w`s working pairwise from left to right.

If `z` is exact `0` and no `w` is exact `0`, then the result is exact `0`. If any `w` is exact `0`, the `exn:fail:contract:divide-by-zero` exception is raised.

Examples:

```
> (/ 3 4)
3/4
> (/ 81 3 3)
9
> (/ 10.0)
0.1
> (/ 1+2i 3+4i)
11/25+2/25i
```

```
(quotient n m) → integer?
  n : integer?
  m : integer?
```

Returns `(truncate (/ n m))`.

Examples:

```
> (quotient 10 3)
3
> (quotient -10.0 3)
-3.0
> (quotient +inf.0 3)
quotient: contract violation
  expected: integer?
  given: +inf.0
  argument position: 1st
  other arguments...:
   3
```

```
(remainder n m) → integer?
  n : integer?
  m : integer?
```

Returns q with the same sign as n such that

- `(abs q)` is between `0` (inclusive) and `(abs m)` (exclusive), and
- `(+ q (* m (quotient n m)))` equals n.

If m is exact `0`, the `exn:fail:contract:divide-by-zero` exception is raised.

Examples:

```
> (remainder 10 3)
1
> (remainder -10.0 3)
-1.0
> (remainder 10.0 -3)
1.0
> (remainder -10 -3)
-1
> (remainder +inf.0 3)
remainder: contract violation
  expected: integer?
```

given: +inf.0
argument position: 1st
other arguments...:
3

```
(quotient/remainder n m) → integer? integer?
  n : integer?
  m : integer?
```

Returns `(values (quotient n m) (remainder n m))`, but the combination may be computed more efficiently than separate calls to `quotient` and `remainder`.

Example:

```
> (quotient/remainder 10 3)
3
1
```

```
(modulo n m) → integer?
  n : integer?
  m : integer?
```

Returns `q` with the same sign as `m` where

- `(abs q)` is between `0` (inclusive) and `(abs m)` (exclusive), and
- the difference between `q` and `(- n (* m (quotient n m)))` is a multiple of `m`.

If `m` is exact `0`, the `exn:fail:contract:divide-by-zero` exception is raised.

Examples:

```
> (modulo 10 3)
1
> (modulo -10.0 3)
2.0
> (modulo 10.0 -3)
-2.0
> (modulo -10 -3)
-1
> (modulo +inf.0 3)
modulo: contract violation
  expected: integer?
  given: +inf.0
  argument position: 1st
  other arguments...:
   3
```

```
(add1 z) → number?
  z : number?
```

Returns `(+ z 1)`.

```
(sub1 z) → number?
  z : number?
```

Returns `(- z 1)`.

```
(abs x) → number?
  x : real?
```

Returns the absolute value of `x`.

Examples:

```
> (abs 1.0)
1.0
> (abs -1)
1
```

```
(max x ...+) → real?
  x : real?
```

Returns the largest of the `x`s, or `+nan.0` if any `x` is `+nan.0`. If any `x` is inexact, the result is coerced to inexact.

Examples:

```
> (max 1 3 2)
3
> (max 1 3 2.0)
3.0
```

```
(min x ...+) → real?
  x : real?
```

Returns the smallest of the `x`s, or `+nan.0` if any `x` is `+nan.0`. If any `x` is inexact, the result is coerced to inexact.

Examples:

```
> (min 1 3 2)
1
> (min 1 3 2.0)
1.0
```

```
(gcd n ...) → rational?
  n : rational?
```

Returns the greatest common divisor (a non-negative number) of the *n*s; for non-integer *n*s, the result is the `gcd` of the numerators divided by the `lcm` of the denominators. If no arguments are provided, the result is `0`. If all arguments are zero, the result is zero.

Examples:

```
> (gcd 10)
10
> (gcd 12 81.0)
3.0
> (gcd 1/2 1/3)
1/6
```

```
(lcm n ...) → rational?
  n : rational?
```

Returns the least common multiple (a non-negative number) of the *n*s; non-integer *n*s, the result is the absolute value of the product divided by the `gcd`. If no arguments are provided, the result is `1`. If any argument is zero, the result is zero; furthermore, if any argument is exact `0`, the result is exact `0`.

Examples:

```
> (lcm 10)
10
> (lcm 3 4.0)
12.0
> (lcm 1/2 2/3)
2
```

```
(round x) → (or/c integer? +inf.0 -inf.0 +nan.0)
  x : real?
```

Returns the integer closest to `x`, resolving ties in favor of an even number, but `+inf.0`, `-inf.0`, and `+nan.0` round to themselves.

Examples:

```
> (round 17/4)
4
> (round -17/4)
-4
```

```
> (round 2.5)
2.0
> (round -2.5)
-2.0
> (round +inf.0)
+inf.0
```

```
(floor x) → (or/c integer? +inf.0 -inf.0 +nan.0)
  x : real?
```

Returns the largest integer that is no more than `x`, but `+inf.0`, `-inf.0`, and `+nan.0` floor to themselves.

Examples:

```
> (floor 17/4)
4
> (floor -17/4)
-5
> (floor 2.5)
2.0
> (floor -2.5)
-3.0
> (floor +inf.0)
+inf.0
```

```
(ceiling x) → (or/c integer? +inf.0 -inf.0 +nan.0)
  x : real?
```

Returns the smallest integer that is at least as large as `x`, but `+inf.0`, `-inf.0`, and `+nan.0` ceiling to themselves.

Examples:

```
> (ceiling 17/4)
5
> (ceiling -17/4)
-4
> (ceiling 2.5)
3.0
> (ceiling -2.5)
-2.0
> (ceiling +inf.0)
+inf.0
```

```
(truncate x) → (or/c integer? +inf.0 -inf.0 +nan.0)
  x : real?
```

Returns the integer farthest from 0 that is not farther from 0 than *x*, but +inf.0, -inf.0, and +nan.0 truncate to themselves.

Examples:

```
> (truncate 17/4)
4
> (truncate -17/4)
-4
> (truncate 2.5)
2.0
> (truncate -2.5)
-2.0
> (truncate +inf.0)
+inf.0
```

```
(numerator q) → integer?
  q : rational?
```

Coerces *q* to an exact number, finds the numerator of the number expressed in its simplest fractional form, and returns this number coerced to the exactness of *q*.

Examples:

```
> (numerator 5)
5
> (numerator 17/4)
17
> (numerator 2.3)
2589569785738035.0
```

```
(denominator q) → integer?
  q : rational?
```

Coerces *q* to an exact number, finds the numerator of the number expressed in its simplest fractional form, and returns this number coerced to the exactness of *q*.

Examples:

```
> (denominator 5)
1
> (denominator 17/4)
4
> (denominator 2.3)
1125899906842624.0
```

```
(rationalize x tolerance) → real?
  x : real?
  tolerance : real?
```

Among the real numbers within (abs *tolerance*) of *x*, returns the one corresponding to an exact number whose denominator is the smallest. If multiple integers are within *tolerance* of *x*, the one closest to 0 is used.

Examples:

```
> (rationalize 1/4 1/10)
1/3
> (rationalize -1/4 1/10)
-1/3
> (rationalize 1/4 1/4)
0
> (rationalize 11/40 1/4)
1/2
```

Number Comparison

```
(= z w ...+) → boolean?
  z : number?
  w : number?
```

Returns #t if all of the arguments are numerically equal, #f otherwise. An inexact number is numerically equal to an exact number when the exact coercion of the inexact number is the exact number. Also, 0.0 and -0.0 are numerically equal, but +nan.0 is not numerically equal to itself.

Examples:

```
> (= 1 1.0)
#t
> (= 1 2)
#f
> (= 2+3i 2+3i 2+3i)
#t
```

```
(< x y ...+) → boolean?
  x : real?
  y : real?
```

Returns #t if the arguments in the given order are strictly increasing, #f otherwise.

Examples:

```
> (< 1 1)
#f
> (< 1 2 3)
#t
```

```
> (< 1 +inf.0)
#t
> (< 1 +nan.0)
#f
```

```
(<= x y ...+) → boolean?
  x : real?
  y : real?
```

Returns `#t` if the arguments in the given order are non-decreasing, `#f` otherwise.

Examples:

```
> (<= 1 1)
#t
> (<= 1 2 1)
#f
```

```
(> x y ...+) → boolean?
  x : real?
  y : real?
```

Returns `#t` if the arguments in the given order are strictly decreasing, `#f` otherwise.

Examples:

```
> (> 1 1)
#f
> (> 3 2 1)
#t
> (> +inf.0 1)
#t
> (> +nan.0 1)
#f
```

```
(>= x y ...+) → boolean?
  x : real?
  y : real?
```

Returns `#t` if the arguments in the given order are non-increasing, `#f` otherwise.

Examples:

```
> (>= 1 1)
#t
> (>= 1 2 1)
#f
```

Powers and Roots

```
(sqrt z) → number?
  z : number?
```

Returns the principal square root of `z`. The result is exact if `z` is exact and `z`'s square root is rational. See also `integer-sqrt`.

Examples:

```
> (sqrt 4/9)
2/3
> (sqrt 2)
1.4142135623730951
> (sqrt -1)
0+1i
```

```
(integer-sqrt n) → complex?
  n : integer?
```

Returns `(floor (sqrt n))` for positive `n`. The result is exact if `n` is exact. For negative `n`, the result is `(* (integer-sqrt (- n)) 0+1i)`.

Examples:

```
> (integer-sqrt 4.0)
2.0
> (integer-sqrt 5)
2
> (integer-sqrt -4.0)
0+2.0i
> (integer-sqrt -4)
0+2i
```

```
(integer-sqrt/remainder n) → complex? integer?
  n : integer?
```

Returns `(integer-sqrt n)` and `(- n (expt (integer-sqrt n) 2))`.

Examples:

```
> (integer-sqrt/remainder 4.0)
2.0
0.0
> (integer-sqrt/remainder 5)
2
1
```

```
(expt z w) → number?
  z : number?
  w : number?
```

Returns `z` raised to the power of `w`.

If `w` is exact `0`, the result is exact `1`. If `w` is `0.0` or `-0.0` and `z` is a real number, the result is `1.0` (even if `z` is `+nan.0`).

If `z` is exact `1`, the result is exact `1`. If `z` is `1.0` and `w` is a real number, the result is `1.0` (even if `w` is `+nan.0`).

If `z` is exact `0` and `w` is negative, the `exn:fail:contract:divide-by-zero` exception is raised.

Further special cases when `w` is a real number:

These special cases correspond to `pow` in C99 [C99], except when `z` is negative and `w` is a not an integer.

- `(expt 0.0 w)`:
  - `w` is negative — `+inf.0`
  - `w` is positive — `0.0`
- `(expt -0.0 w)`:
  - `w` is negative:
    - `w` is an odd integer — `-inf.0`
    - `w` otherwise rational — `+inf.0`
  - `w` is positive:
    - `w` is an odd integer — `-0.0`
    - `w` otherwise rational — `0.0`
- `(expt z -inf.0)` for positive `z`:
  - `z` is less than `1.0` — `+inf.0`
  - `z` is greater than `1.0` — `0.0`
- `(expt z +inf.0)` for positive `z`:
  - `z` is less than `1.0` — `0.0`
  - `z` is greater than `1.0` — `+inf.0`
- `(expt -inf.0 w)` for integer `w`:
  - `w` is negative:
    - `w` is odd — `-0.0`

        * `w` is even — `0.0`
    - `w` is positive:
        * `w` is odd — `-inf.0`
        * `w` is even — `+inf.0`

- `(expt +inf.0 w)`:
    - `w` is negative — `0.0`
    - `w` is positive — `+inf.0`

Examples:

```
> (expt 2 3)
8
> (expt 4 0.5)
2.0
> (expt +inf.0 0)
1
```

`(exp z) → number?`
  `z` : `number?`

Returns Euler's number raised to the power of `z`. The result is normally inexact, but it is exact `1` when `z` is an exact `0`.

Examples:

```
> (exp 1)
2.718281828459045
> (exp 2+3i)
-7.315110094901103+1.0427436562359045i
> (exp 0)
1
```

`(log z) → number?`
  `z` : `number?`

Returns the natural logarithm of `z`. The result is normally inexact, but it is exact `0` when `z` is an exact `1`. When `z` is exact `0`, `exn:fail:contract:divide-by-zero` exception is raised.

Examples:

```
> (log (exp 1))
1.0
> (log 2+3i)
1.2824746787307684+0.9827937232473291i
> (log 1)
0
```

Trigonometric Functions

```
(sin z) → number?
  z : number?
```

Returns the sine of `z`, where `z` is in radians. The result is normally inexact, but it is exact `0` if `z` is exact `0`.

Examples:

```
> (sin 3.14159)
2.65358979335273e-06
> (sin 1.0+5.0i)
62.44551846769653+40.0921657779984i
```

```
(cos z) → number?
  z : number?
```

Returns the cosine of `z`, where `z` is in radians.

Examples:

```
> (cos 3.14159)
-0.9999999999964793
> (cos 1.0+5.0i)
40.095806306298826-62.439848680799631
```

```
(tan z) → number?
  z : number?
```

Returns the tangent of `z`, where `z` is in radians. The result is normally inexact, but it is exact `0` if `z` is exact `0`.

Examples:

```
> (tan 0.7854)
1.0000036732118496
> (tan 1.0+5.0i)
8.256719834227411e-05+1.0000377833796008i
```

```
(asin z) → number?
  z : number?
```

Returns the arcsine in radians of `z`. The result is normally inexact, but it is exact `0` if `z` is exact `0`.

Examples:

```
> (asin 0.25)
0.25268025514207865
> (asin 1.0+5.0i)
0.1937931365549321+2.3309746530493123i
```

```
(acos z) → number?
  z : number?
```

Returns the arccosine in radians of `z`.

Examples:

```
> (acos 0.25)
1.318116071652818
> (acos 1.0+5.0i)
1.3770031902399644-2.3309746530493123i
```

```
(atan z) → number?
  z : number?
(atan y x) → number?
  y : real?
  x : real?
```

In the one-argument case, returns the arctangent of the inexact approximation of `z`, except that the result is an exact `0` for an exact `0` argument.

In the two-argument case, the result is roughly the same as `(atan (/ (exact->inexact y)) (exact->inexact x))`, but the signs of `y` and `x` determine the quadrant of the result. Moreover, a suitable angle is returned when `y` divided by `x` produces `+nan.0` in the case that neither `y` nor `x` is `+nan.0`. Finally, if `y` is exact `0` and `x` is an exact positive number, the result is exact `0`. If both `x` and `y` are exact `0`, the `exn:fail:contract:divide-by-zero` exception is raised.

Examples:

```
> (atan 0.5)
0.4636476090008061
> (atan 2 1)
1.1071487177940904
> (atan -2 -1)
-2.0344439357957027
> (atan 1.0+5.0i)
1.530881333938778+0.19442614214700213i
> (atan +inf.0 -inf.0)
2.356194490192345
```

Complex Numbers

```
(make-rectangular x y) → number?
  x : real?
  y : real?
```

Returns `(+ x (* y 0+1i))`.

Example:

```
> (make-rectangular 3 4.0)
3.0+4.0i
```

```
(make-polar magnitude angle) → number?
  magnitude : real?
  angle : real?
```

Returns `(+ (* magnitude (cos angle)) (* magnitude (sin angle) 0+1i))`.

Examples:

```
> (make-polar 10 (* pi 1/2))
6.123233995736766e-16+10.0i
> (make-polar 10 (* pi 1/4))
7.0710678118654755+7.071067811865475i
```

```
(real-part z) → real?
  z : number?
```

Returns the real part of the complex number `z` in rectangle coordinates.

Examples:

```
> (real-part 3+4i)
3
> (real-part 5.0)
5.0
```

```
(imag-part z) → real?
  z : number?
```

Returns the imaginary part of the complex number `z` in rectangle coordinates.

Examples:

```
> (imag-part 3+4i)
4
> (imag-part 5.0)
0
> (imag-part 5.0+0.0i)
0.0
```

```
(magnitude z) → (and/c real? (not/c negative?))
  z : number?
```

Returns the magnitude of the complex number `z` in polar coordinates.

Examples:

```
> (magnitude -3)
3
> (magnitude 3.0)
3.0
> (magnitude 3+4i)
5
```

```
(angle z) → real?
  z : number?
```

Returns the angle of the complex number `z` in polar coordinates.

The result is guaranteed to be between `(- pi)` and `pi`, possibly equal to `pi` (but never equal to `(- pi)`).

Examples:

```
> (angle -3)
3.141592653589793
> (angle 3.0)
0
> (angle 3+4i)
0.9272952180016122
> (angle +inf.0+inf.0i)
0.7853981633974483
> (angle -1)
3.141592653589793
```

Bitwise Operations

```
(bitwise-ior n ...) → exact-integer?
  n : exact-integer?
```

Returns the bitwise "inclusive or" of the `n`s in their (semi-infinite) two's complement representation. If no arguments are provided, the result is `0`.

Examples:

```
> (bitwise-ior 1 2)
3
> (bitwise-ior -32 1)
-31
```

```
(bitwise-and n ...) → exact-integer?
  n : exact-integer?
```

Returns the bitwise “and” of the *n*s in their (semi-infinite) two’s complement representation. If no arguments are provided, the result is `-1`.

Examples:

```
> (bitwise-and 1 2)
0
> (bitwise-and -32 -1)
-32
```

```
(bitwise-xor n ...) → exact-integer?
  n : exact-integer?
```

Returns the bitwise “exclusive or” of the *n*s in their (semi-infinite) two’s complement representation. If no arguments are provided, the result is `0`.

Examples:

```
> (bitwise-xor 1 5)
4
> (bitwise-xor -32 -1)
31
```

```
(bitwise-not n) → exact-integer?
  n : exact-integer?
```

Returns the bitwise “not” of *n* in its (semi-infinite) two’s complement representation.

Examples:

```
> (bitwise-not 5)
-6
> (bitwise-not -1)
0
```

```
(bitwise-bit-set? n m) → boolean?
  n : exact-integer?
  m : exact-nonnegative-integer?
```

Returns #t when the *m*th bit of *n* is set in *n*'s (semi-infinite) two's complement representation.

This operation is equivalent to (not (zero? (bitwise-and *n* (arithmetic-shift 1 *m*)))), but it is faster and runs in constant time when *n* is positive.

Examples:

```
> (bitwise-bit-set? 5 0)
#t
> (bitwise-bit-set? 5 2)
#t
> (bitwise-bit-set? -5 (expt 2 700))
#t
```

```
(bitwise-bit-field n start end) → exact-integer?
  n : exact-integer?
  start : exact-nonnegative-integer?
  end : (and/c exact-nonnegative-integer?
               (start . <= . end))
```

Extracts the bits between position *start* and (- *end* 1) (inclusive) from *n* and shifts them down to the least significant portion of the number.

This operation is equivalent to the computation

```
(bitwise-and (sub1 (arithmetic-shift 1 (- end start)))
             (arithmetic-shift n (- start)))
```

but it runs in constant time when *n* is positive, *start* and *end* are fixnums, and (- *end* *start*) is no more than the maximum width of a fixnum.

Each pair of examples below uses the same numbers, showing the result both in binary and as integers.

Examples:

```
> (format "~b" (bitwise-bit-field (string->number "1101" 2) 1 1))
"0"
> (bitwise-bit-field 13 1 1)
0
> (format "~b" (bitwise-bit-field (string->number "1101" 2) 1 3))
"10"
> (bitwise-bit-field 13 1 3)
2
> (format "~b" (bitwise-bit-field (string->number "1101" 2) 1 4))
"110"
> (bitwise-bit-field 13 1 4)
6
```

```
(arithmetic-shift n m) → exact-integer?
  n : exact-integer?
  m : exact-integer?
```

Returns the bitwise "shift" of *n* in its (semi-infinite) two's complement representation. If *m* is non-negative, the integer *n* is shifted left by *m* bits; i.e., *m* new zeros are introduced as rightmost digits. If *m* is negative, *n* is shifted right by `(- m)` bits; i.e., the rightmost *m* digits are dropped.

Examples:

```
> (arithmetic-shift 1 10)
1024
> (arithmetic-shift 255 -3)
31
```

```
(integer-length n) → exact-integer?
  n : exact-integer?
```

Returns the number of bits in the (semi-infinite) two's complement representation of *n* after removing all leading zeros (for non-negative *n*) or ones (for negative *n*).

Examples:

```
> (integer-length 8)
4
> (integer-length -8)
3
```

Random Numbers

```
(random k [rand-gen]) → exact-nonnegative-integer?
  k : (integer-in 1 4294967087)
  rand-gen : pseudo-random-generator?
           = (current-pseudo-random-generator)
(random [rand-gen]) → (and/c real? inexact? (>/c 0) (</c 1))
  rand-gen : pseudo-random-generator?
           = (current-pseudo-random-generator)
```

When called with an integer argument *k*, returns a random exact integer in the range `0` to *k*-1. When called with zero arguments, returns a random inexact number between `0` and `1`, exclusive.

In each case, the number is provided by the given pseudo-random number generator (which defaults to the current one, as produced by `current-pseudo-random-generator`). The generator maintains an internal state for generating numbers. The random number generator uses a 54-bit version of L'Ecuyer's MRG32k3a algorithm [L'Ecuyer02].

```
(random-seed k) → void?
  k : (integer-in 1 (sub1 (expt 2 31)))
```

Seeds the current pseudo-random number generator with *k*. Seeding a generator sets its internal state deterministically; that is, seeding a generator with a particular number forces it to produce a sequence of pseudo-random numbers that is the same across runs and across platforms.

The random-seed function is convenient for some purposes, but note that the space of states for a pseudo-random number generator is much larger that the space of allowed values for *k*. Use vector->pseudo-random-generator! to set a pseudo-random number generator to any of its possible states.

```
(make-pseudo-random-generator) → pseudo-random-generator?
```

Returns a new pseudo-random number generator. The new generator is seeded with a number derived from (current-milliseconds).

```
(pseudo-random-generator? v) → boolean?
  v : any/c
```

Returns #t if *v* is a pseudo-random number generator, #f otherwise.

```
(current-pseudo-random-generator) → pseudo-random-generator?
(current-pseudo-random-generator rand-gen) → void?
  rand-gen : pseudo-random-generator?
```

A parameter that determines the pseudo-random number generator used by random.

```
(pseudo-random-generator->vector rand-gen)
 → pseudo-random-generator-vector?
  rand-gen : pseudo-random-generator?
```

Produces a vector that represents the complete internal state of *rand-gen*. The vector is suitable as an argument to vector->pseudo-random-generator to recreate the generator in its current state (across runs and across platforms).

```
(vector->pseudo-random-generator vec)
 → pseudo-random-generator?
  vec : pseudo-random-generator-vector?
```

Produces a pseudo-random number generator whose internal state corresponds to *vec*.

```
(vector->pseudo-random-generator! rand-gen
                                  vec)      → void?
  rand-gen : pseudo-random-generator?
  vec : pseudo-random-generator-vector?
```

Like vector->pseudo-random-generator, but changes *rand-gen* to the given state, instead of creating a new generator.

```
(pseudo-random-generator-vector? v) → boolean?
  v : any/c
```

Returns #t if *v* is a vector of six exact integers, where the first three integers are in the range 0 to 4294967086, inclusive; the last three integers are in the range 0 to 4294944442, inclusive; at least one of the first three integers is non-zero; and at least one of the last three integers is non-zero. Otherwise, the result is #f.

Number–String Conversions

```
(number->string z [radix]) → string?
  z : number?
  radix : (or/c 2 8 10 16) = 10
```

Returns a string that is the printed form of *z* in the base specified by *radix*. If *z* is inexact, *radix* must be 10, otherwise the exn:fail:contract exception is raised.

Examples:

```
> (number->string 3.0)
"3.0"
> (number->string 255 8)
"377"
```

```
(string->number s [radix]) → (or/c number? #f)
  s : string?
  radix : (integer-in 2 16) = 10
```

Reads and returns a number datum from *s* (see §1.3.3 “Reading Numbers”), returning #f if *s* does not parse exactly as a number datum (with no whitespace). The optional *radix* argument specifies the default base for the number, which can be overridden by #b, #o, #d, or #x in the string. The read-decimal-as-inexact parameter affects string->number in the same as way as read.

Examples:

```
> (string->number "3.0+2.5i")
3.0+2.5i
> (string->number "hello")
#f
> (string->number "111" 7)
57
> (string->number "#b111" 7)
7
```

```
(real->decimal-string n [decimal-digits]) → string?
  n : real?
  decimal-digits : exact-nonnegative-integer? = 2
```

Prints *n* into a string and returns the string. The printed form of *n* shows exactly *decimal-digits* digits after the decimal point. The printed form uses a minus sign if *n* is negative, and it does not use a plus sign if *n* is positive.

Before printing, *n* is converted to an exact number, multiplied by (expt 10 *decimal-digits*), rounded, and then divided again by (expt 10 *decimal-digits*). The result of this process is an exact number whose decimal representation has no more than *decimal-digits* digits after the decimal (and it is padded with trailing zeros if necessary).

Examples:

```
> (real->decimal-string pi)
"3.14"
> (real->decimal-string pi 5)
"3.14159"
```

```
(integer-bytes->integer  bstr
                         signed?
                        [big-endian?
                         start
                         end])        → exact-integer?
  bstr : bytes?
  signed? : any/c
  big-endian? : any/c = (system-big-endian?)
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (bytes-length bstr)
```

Converts the machine-format number encoded in *bstr* to an exact integer. The *start* and *end* arguments specify the substring to decode, where (- *end* *start*) must be 2, 4, or 8. If *signed?* is true, then the bytes are decoded as a two's-complement number, otherwise it is decoded as an unsigned integer. If *big-endian?* is true, then the first character's ASCII value provides the most significant eight bits of the number, otherwise the first character provides the least-significant eight bits, and so on.

```
(integer->integer-bytes  n
                         size-n
                         signed?
                        [big-endian?
                         dest-bstr
                         start])      → bytes?
  n : exact-integer?
  size-n : (or/c 2 4 8)
  signed? : any/c
  big-endian? : any/c = (system-big-endian?)
  dest-bstr :  (and/c bytes? (not/c immutable?))
            = (make-bytes size-n)
  start : exact-nonnegative-integer? = 0
```

Converts the exact integer *n* to a machine-format number encoded in a byte string of length *size-n*, which must be 2, 4, or 8. If *signed?* is true, then the number is encoded as two's complement, otherwise it is encoded as an unsigned bit stream. If *big-endian?* is true, then the most significant eight bits of the number are encoded in the first character of the resulting byte string, otherwise the least-significant bits are encoded in the first byte, and so on.

The *dest-bstr* argument must be a mutable byte string of length *size-n*. The encoding of *n* is written into *dest-bstr* starting at offset *start*, and *dest-bstr* is returned as the result.

If *n* cannot be encoded in a string of the requested size and format, the `exn:fail:contract` exception is raised. If *dest-bstr* is not of length *size-n*, the `exn:fail:contract` exception is raised.

```
(floating-point-bytes->real  bstr
                            [big-endian?
                             start
                             end])       → flonum?
  bstr : bytes?
  big-endian? : any/c = (system-big-endian?)
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (bytes-length bstr)
```

Converts the IEEE floating-point number encoded in *bstr* from position *start* (inclusive) to *end* (exclusive) to an inexact real number. The difference between *start* an *end* must be either 4 or 8 bytes. If *big-endian?* is true, then the first byte's ASCII value provides the most significant eight bits of the IEEE representation, otherwise the first byte provides the least-significant eight bits, and so on.

```
(real->floating-point-bytes  x
                             size-n
                            [big-endian?
                             dest-bstr
                             start])      → bytes?
  x : real?
  size-n : (or/c 4 8)
  big-endian? : any/c = (system-big-endian?)
  dest-bstr : (and/c bytes? (not/c immutable?))
            = (make-bytes size-n)
  start : exact-nonnegative-integer? = 0
```

Converts the real number *x* to its IEEE representation in a byte string of length *size-n*, which must be 4 or 8. If *big-endian?* is true, then the most significant eight bits of the number are encoded in the first byte of the resulting byte string, otherwise the least-significant bits are encoded in the first character, and so on.

The *dest-bstr* argument must be a mutable byte string of length *size-n*. The encoding of n is written into *dest-bstr* starting with byte *start*, and *dest-bstr* is returned as the result.

If *dest-bstr* is provided and it has less than *start* plus *size-n* bytes, the exn:fail:contract exception is raised.

```
(system-big-endian?) → boolean?
```

Returns #t if the native encoding of numbers is big-endian for the machine running Racket, #f if the native encoding is little-endian.

Extra Constants and Functions

(require racket/math) package: base

The bindings documented in this section are provided by the racket/math and racket libraries, but not racket/base.

```
pi : flonum?
```

An approximation of π, the ratio of a circle's circumference to its diameter.

Examples:

```
> pi
3.141592653589793
> (cos pi)
-1.0
```

```
pi.f : single-flonum?
```

Like pi, but in single precision.

Examples:

```
> pi.f
3.1415927f0
> (* 2.0f0 pi)
6.283185307179586
> (* 2.0f0 pi.f)
6.2831855f0
```

```
(degrees->radians x) → real?
  x : real?
```

Converts an *x*-degree angle to radians.

Examples:

```
> (degrees->radians 180)
3.141592653589793
> (sin (degrees->radians 45))
0.7071067811865475
```

```
(radians->degrees x) → real?
  x : real?
```

Converts *x* radians to degrees.

Examples:

```
> (radians->degrees pi)
180.0
> (radians->degrees (* 1/4 pi))
45.0
```

```
(sqr z) → number?
  z : number?
```

Returns (* *z* *z*).

```
(sgn x) → (or/c 1 0 -1 1.0 0.0 -1.0)
  x : real?
```

Returns the sign of *x* as either -1, 0, or 1.

Examples:

```
> (sgn 10)
1
> (sgn -10.0)
-1.0
> (sgn 0)
0
```

(conjugate *z*) → number?
  z : number?

Returns the complex conjugate of *z*.

Examples:

```
> (conjugate 1)
1
> (conjugate 3+4i)
3-4i
```

(sinh *z*) → number?
  z : number?

Returns the hyperbolic sine of *z*.

(cosh *z*) → number?
  z : number?

Returns the hyperbolic cosine of *z*.

(tanh *z*) → number?
  z : number?

Returns the hyperbolic tangent of *z*.

(exact-round *x*) → exact-integer?
  x : rational?

Equivalent to (inexact->exact (round *x*)).

(exact-floor *x*) → exact-integer?
  x : rational?

Equivalent to (inexact->exact (floor *x*)).

```
(exact-ceiling x) → exact-integer?
  x : rational?
```

Equivalent to `(inexact->exact (ceiling x))`.

```
(exact-truncate x) → exact-integer?
  x : rational?
```

Equivalent to `(inexact->exact (truncate x))`.

```
(order-of-magnitude r) → (and/c exact? integer?)
  r : (and/c real? positive?)
```

Computes the greatest exact integer `m` such that:

```
(<= (expt 10 m)
    (inexact->exact r))
```

Hence also:

```
(< (inexact->exact r)
   (expt 10 (add1 m)))
```

Examples:

```
> (order-of-magnitude 999)
2
> (order-of-magnitude 1000)
3
> (order-of-magnitude 1/100)
-2
> (order-of-magnitude 1/101)
-3
```

```
(nan? x) → boolean?
  x : real?
```

Returns `#t` if `x` is `eqv?` to `+nan.0` or `+nan.f`; otherwise `#f`.

```
(infinite? x) → boolean?
  x : real?
```

Returns `#t` if `z` is `+inf.0`, `-inf.0`, `+inf.f`, `-inf.f`; otherwise `#f`.

4.2.3 Flonums

```
(require racket/flonum)
```
package: base

The racket/flonum library provides operations like fl+ that consume and produce only flonums. Flonum-specific operations can provide better performance when used consistently, and they are as safe as generic operations like +.

See also §19.7 "Fixnum and Flonum Optimizations" in ***The Racket Guide***.

Flonum Arithmetic

```
(fl+ a b) → flonum?
  a : flonum?
  b : flonum?
(fl- a b) → flonum?
  a : flonum?
  b : flonum?
(fl* a b) → flonum?
  a : flonum?
  b : flonum?
(fl/ a b) → flonum?
  a : flonum?
  b : flonum?
(flabs a) → flonum?
  a : flonum?
```

Like +, -, *, /, and abs, but constrained to consume flonums. The result is always a flonum.

```
(fl= a b) → boolean?
  a : flonum?
  b : flonum?
(fl< a b) → boolean?
  a : flonum?
  b : flonum?
(fl> a b) → boolean?
  a : flonum?
  b : flonum?
(fl<= a b) → boolean?
  a : flonum?
  b : flonum?
(fl>= a b) → boolean?
  a : flonum?
  b : flonum?
(flmin a b) → flonum?
  a : flonum?
  b : flonum?
```

```
(flmax a b) → flonum?
  a : flonum?
  b : flonum?
```

Like `=`, `<`, `>`, `<=`, `>=`, `min`, and `max`, but constrained to consume flonums.

```
(flround a) → flonum?
  a : flonum?
(flfloor a) → flonum?
  a : flonum?
(flceiling a) → flonum?
  a : flonum?
(fltruncate a) → flonum?
  a : flonum?
```

Like `round`, `floor`, `ceiling`, and `truncate`, but constrained to consume flonums.

```
(flsin a) → flonum?
  a : flonum?
(flcos a) → flonum?
  a : flonum?
(fltan a) → flonum?
  a : flonum?
(flasin a) → flonum?
  a : flonum?
(flacos a) → flonum?
  a : flonum?
(flatan a) → flonum?
  a : flonum?
(fllog a) → flonum?
  a : flonum?
(flexp a) → flonum?
  a : flonum?
(flsqrt a) → flonum?
  a : flonum?
```

Like `sin`, `cos`, `tan`, `asin`, `acos`, `atan`, `log`, `exp`, and `sqrt`, but constrained to consume and produce flonums. The result is `+nan.0` when a number outside the range `-1.0` to `1.0` is given to `flasin` or `flacos`, or when a negative number is given to `fllog` or `flsqrt`.

```
(flexpt a b) → flonum?
  a : flonum?
  b : flonum?
```

Like `expt`, but constrained to consume and produce flonums.

Due to the result constraint, the results compared to `expt` differ in the following cases:

These special cases correspond to `pow` in C99 [C99].

- `(flexpt -1.0 +inf.0)` — `1.0`
- `(flexpt a +inf.0)` where `a` is negative — `(expt (abs a) +inf.0)`
- `(flexpt a -inf.0)` where `a` is negative — `(expt (abs a) -inf.0)`
- `(expt -inf.0 b)` where *b* is a non-integer:
  - *b* is negative — `0.0`
  - *b* is positive — `+inf.0`
- `(flexpt a b)` where `a` is negative and *b* is not an integer — `+nan.0`

```
(->fl a) → flonum?
  a : exact-integer?
```

Like `exact->inexact`, but constrained to consume exact integers, so the result is always a flonum.

```
(fl->exact-integer a) → exact-integer?
  a : flonum?
```

Like `inexact->exact`, but constrained to consume an integer flonum, so the result is always an exact integer.

```
(make-flrectangular a b)
    (and/c complex?
 →         (lambda (c) (flonum? (real-part c)))
           (lambda (c) (flonum? (imag-part c))))
  a : flonum?
  b : flonum?
(flreal-part a) → flonum?
      (and/c complex?
  a :         (lambda (c) (flonum? (real-part c)))
              (lambda (c) (flonum? (imag-part c))))
(flimag-part a) → flonum?
      (and/c complex?
  a :         (lambda (c) (flonum? (real-part c)))
              (lambda (c) (flonum? (imag-part c))))
```

Like `make-rectangular`, `real-part`, and `imag-part`, but both parts of the complex number must be inexact.

```
(flrandom rand-gen) → (and flonum? (>/c 0) (</c 1))
  rand-gen : pseudo-random-generator?
```

Equivalent to (random rand-gen).

Flonum Vectors

A *flvector* is like a vector, but it holds only inexact real numbers. This representation can be more compact, and unsafe operations on flvectors (see racket/unsafe/ops) can execute more efficiently than unsafe operations on vectors of inexact reals.

An f64vector as provided by ffi/vector stores the same kinds of values as a flvector, but with extra indirections that make f64vectors more convenient for working with foreign libraries. The lack of indirections makes unsafe flvector access more efficient.

Two flvectors are equal? if they have the same length, and if the values in corresponding slots of the flvectors are equal?.

```
(flvector? v) → boolean?
  v : any/c
```

Returns #t if v is a flvector, #f otherwise.

```
(flvector x ...) → flvector?
  x : flonum?
```

Creates a flvector containing the given inexact real numbers.

Example:

```
> (flvector 2.0 3.0 4.0 5.0)
(flvector 2.0 3.0 4.0 5.0)
```

```
(make-flvector size [x]) → flvector?
  size : exact-nonnegative-integer?
  x : flonum? = 0.0
```

Creates a flvector with *size* elements, where every slot in the flvector is filled with x.

Example:

```
> (make-flvector 4 3.0)
(flvector 3.0 3.0 3.0 3.0)
```

```
(flvector-length vec) → exact-nonnegative-integer?
  vec : flvector?
```

Returns the length of *vec* (i.e., the number of slots in the flvector).

```
(flvector-ref vec pos) → flonum?
  vec : flvector?
  pos : exact-nonnegative-integer?
```

Returns the inexact real number in slot *pos* of *vec*. The first slot is position 0, and the last slot is one less than (flvector-length *vec*).

```
(flvector-set! vec pos x) → flonum?
  vec : flvector?
  pos : exact-nonnegative-integer?
  x : flonum?
```

Sets the inexact real number in slot *pos* of *vec*. The first slot is position 0, and the last slot is one less than (flvector-length *vec*).

```
(flvector-copy vec [start end]) → flvector?
  vec : flvector?
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (vector-length v)
```

Creates a fresh flvector of size (- *end start*), with all of the elements of *vec* from *start* (inclusive) to *end* (exclusive).

```
(in-flvector vec [start stop step]) → sequence?
  vec : flvector?
  start : exact-nonnegative-integer? = 0
  stop : (or/c exact-integer? #f) = #f
  step : (and/c exact-integer? (not/c zero?)) = 1
```

Returns a sequence equivalent to *vec* when no optional arguments are supplied.

The optional arguments *start*, *stop*, and *step* are as in in-vector.

A in-flvector application can provide better performance for flvector iteration when it appears directly in a for clause.

```
(for/flvector maybe-length (for-clause ...) body ...)
(for*/flvector maybe-length (for-clause ...) body ...)

maybe-length =
             | #:length length-expr
             | #:length length-expr #:fill fill-expr

  length-expr : exact-nonnegative-integer?
  fill-expr : flonum?
```

Like `for/vector` or `for*/vector`, but for flvectors. The default *fill-expr* produces `0.0`.

```
(shared-flvector x ...) → flvector?
  x : flonum?
```

Creates a flvector containing the given inexact real numbers. For communication among places, the new flvector is allocated in the shared memory space.

Example:

```
> (shared-flvector 2.0 3.0 4.0 5.0)
(flvector 2.0 3.0 4.0 5.0)
```

```
(make-shared-flvector size [x]) → flvector?
  size : exact-nonnegative-integer?
  x : flonum? = 0.0
```

Creates a flvector with *size* elements, where every slot in the flvector is filled with *x*. For communication among places, the new flvector is allocated in the shared memory space.

Example:

```
> (make-shared-flvector 4 3.0)
(flvector 3.0 3.0 3.0 3.0)
```

4.2.4 Fixnums

`(require racket/fixnum)` package: `base`

The `racket/fixnum` library provides operations like `fx+` that consume and produce only fixnums. The operations in this library are meant to be safe versions of unsafe operations like `unsafe-fx+`. These safe operations are generally no faster than using generic primitives like `+`.

The expected use of the `racket/fixnum` library is for code where the `require` of `racket/fixnum` is replaced with

```
(require (filtered-in
          (λ (name) (regexp-replace #rx"unsafe-" name ""))
          racket/unsafe/ops))
```

to drop in unsafe versions of the library. Alternately, when encountering crashes with code that uses unsafe fixnum operations, use the `racket/fixnum` library to help debug the problems.

Fixnum Arithmetic

```
(fx+ a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fx- a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fx* a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fxquotient a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fxremainder a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fxmodulo a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fxabs a) → fixnum?
  a : fixnum?
```

Safe versions of unsafe-fx+, unsafe-fx-, unsafe-fx*, unsafe-fxquotient, unsafe-fxremainder, unsafe-fxmodulo, and unsafe-fxabs. The exn:fail:contract:non-fixnum-result exception is raised if the arithmetic result would not be a fixnum.

```
(fxand a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fxior a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fxxor a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fxnot a) → fixnum?
  a : fixnum?
(fxlshift a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fxrshift a b) → fixnum?
  a : fixnum?
  b : fixnum?
```

Safe versions of unsafe-fxand, unsafe-fxior, unsafe-fxxor, unsafe-fxnot,

unsafe-fxlshift, and unsafe-fxrshift. The exn:fail:contract:non-fixnum-result exception is raised if the arithmetic result would not be a fixnum.

```
(fx= a b) → boolean?
  a : fixnum?
  b : fixnum?
(fx< a b) → boolean?
  a : fixnum?
  b : fixnum?
(fx> a b) → boolean?
  a : fixnum?
  b : fixnum?
(fx<= a b) → boolean?
  a : fixnum?
  b : fixnum?
(fx>= a b) → boolean?
  a : fixnum?
  b : fixnum?
(fxmin a b) → fixnum?
  a : fixnum?
  b : fixnum?
(fxmax a b) → fixnum?
  a : fixnum?
  b : fixnum?
```

Safe versions of unsafe-fx=, unsafe-fx<, unsafe-fx>, unsafe-fx<=, unsafe-fx>=, unsafe-fxmin, and unsafe-fxmax.

```
(fx->fl a) → flonum?
  a : fixnum?
(fl->fx a) → fixnum?
  a : flonum?
```

Safe versions of unsafe-fx->fl and unsafe-fl->fx.

Fixnum Vectors

A *fxvector* is like a vector, but it holds only fixnums. The only advantage of a fxvector over a vector is that a shared version can be created with functions like shared-fxvector.

Two fxvectors are equal? if they have the same length, and if the values in corresponding slots of the fxvectors are equal?.

```
(fxvector? v) → boolean?
  v : any/c
```

Returns #t if v is a fxvector, #f otherwise.

```
(fxvector x ...) → fxvector?
  x : fixnum?
```

Creates a fxvector containing the given fixnums.

Example:

```
> (fxvector 2 3 4 5)
(fxvector 2 3 4 5)
```

```
(make-fxvector size [x]) → fxvector?
  size : exact-nonnegative-integer?
  x : fixnum? = 0
```

Creates a fxvector with *size* elements, where every slot in the fxvector is filled with *x*.

Example:

```
> (make-fxvector 4 3)
(fxvector 3 3 3 3)
```

```
(fxvector-length vec) → exact-nonnegative-integer?
  vec : fxvector?
```

Returns the length of *vec* (i.e., the number of slots in the fxvector).

```
(fxvector-ref vec pos) → fixnum?
  vec : fxvector?
  pos : exact-nonnegative-integer?
```

Returns the fixnum in slot *pos* of *vec*. The first slot is position 0, and the last slot is one less than (fxvector-length *vec*).

```
(fxvector-set! vec pos x) → fixnum?
  vec : fxvector?
  pos : exact-nonnegative-integer?
  x : fixnum?
```

Sets the fixnum in slot *pos* of *vec*. The first slot is position 0, and the last slot is one less than (fxvector-length *vec*).

```
(fxvector-copy vec [start end]) → fxvector?
  vec : fxvector?
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (vector-length v)
```

Creates a fresh fxvector of size `(- end start)`, with all of the elements of *vec* from *start* (inclusive) to *end* (exclusive).

```
(in-fxvector vec [start stop step]) → sequence?
  vec : fxvector?
  start : exact-nonnegative-integer? = 0
  stop : (or/c exact-integer? #f) = #f
  step : (and/c exact-integer? (not/c zero?)) = 1
```

Returns a sequence equivalent to *vec* when no optional arguments are supplied.

The optional arguments *start*, *stop*, and *step* are as in `in-vector`.

An `in-fxvector` application can provide better performance for fxvector iteration when it appears directly in a `for` clause.

```
(for/fxvector maybe-length (for-clause ...) body ...)
(for*/fxvector maybe-length (for-clause ...) body ...)

maybe-length =
             | #:length length-expr
             | #:length length-expr #:fill fill-expr

  length-expr : exact-nonnegative-integer?
  fill-expr : fixnum?
```

Like `for/vector` or `for*/vector`, but for fxvectors. The default *fill-expr* produces `0`.

```
(shared-fxvector x ...) → fxvector?
  x : fixnum?
```

Creates a fxvector containing the given fixnums. For communication among places, the new fxvector is allocated in the shared memory space.

Example:

```
> (shared-fxvector 2 3 4 5)
(fxvector 2 3 4 5)
```

```
(make-shared-fxvector size [x]) → fxvector?
  size : exact-nonnegative-integer?
  x : fixnum? = 0
```

Creates a fxvector with *size* elements, where every slot in the fxvector is filled with *x*. For communication among places, the new fxvector is allocated in the shared memory space.

Example:

```
> (make-shared-fxvector 4 3)
(fxvector 3 3 3 3)
```

4.2.5 Extflonums

```
(require racket/extflonum)
```
package: base

An *extflonum* is an extended-precision (80-bit) floating-point number. extflonum arithmetic is supported on platforms with extended-precision hardware and where the extflonum implementation does not conflict with normal double-precision arithmetic (i.e., on x86 and x86_64 platforms when Racket is compiled to use SSE instructions for floating-point operations, and on Windows when `"longdouble.dll"` is available).

A extflonum is **not** a number in the sense of `number?`. Only extflonum-specific operations such as `extfl+` perform extflonum arithmetic.

A literal extflonum is written like an inexact number, but using an explicit `t` or `T` exponent marker (see §1.3.4 "Reading Extflonums"). For example, `3.5t0` is an extflonum. The extflonum infinities and non-a-number values are `+inf.t`, `-inf.t`, and `+nan.t`.

If `(extflonum-available?)` produces `#f`, then all operations exported by `racket/extflonum` raise `exn:fail:unsupported`, except for `extflonum?`, `extflonum-available?`, and `extflvector?` (which always work). The reader (see §1.3 "The Reader") always accepts extflonum input; when extflonum operations are not supported, printing an extflonum from the reader uses its source notation (as opposed to normalizing the format).

Two extflonums are `equal?` if `extfl=` produces `#t` for the extflonums. If extflonums are not supported in a platform, extflonums are `equal?` only if they are `eq?`.

```
(extflonum? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is an extflonum, `#f` otherwise.

```
(extflonum-available?) → boolean?
```

Returns `#t` if extflonum operations are supported on the current platform, `#f` otherwise.

Extflonum Arithmetic

```
(extfl+ a b) → extflonum?
  a : extflonum?
  b : extflonum?
```

```
(extfl- a b) → extflonum?
  a : extflonum?
  b : extflonum?
(extfl* a b) → extflonum?
  a : extflonum?
  b : extflonum?
(extfl/ a b) → extflonum?
  a : extflonum?
  b : extflonum?
(extflabs a) → extflonum?
  a : extflonum?
```

Like `fl+`, `fl-`, `fl*`, `fl/`, and `flabs`, but for extflonums.

```
(extfl= a b) → boolean?
  a : extflonum?
  b : extflonum?
(extfl< a b) → boolean?
  a : extflonum?
  b : extflonum?
(extfl> a b) → boolean?
  a : extflonum?
  b : extflonum?
(extfl<= a b) → boolean?
  a : extflonum?
  b : extflonum?
(extfl>= a b) → boolean?
  a : extflonum?
  b : extflonum?
(extflmin a b) → extflonum?
  a : extflonum?
  b : extflonum?
(extflmax a b) → extflonum?
  a : extflonum?
  b : extflonum?
```

Like `fl=`, `fl<`, `fl>`, `fl<=`, `fl>=`, `flmin`, and `flmax`, but for extflonums.

```
(extflround a) → extflonum?
  a : extflonum?
(extflfloor a) → extflonum?
  a : extflonum?
(extflceiling a) → extflonum?
  a : extflonum?
(extfltruncate a) → extflonum?
  a : extflonum?
```

Like `flround`, `flfloor`, `flceiling`, and `fltruncate`, but for extflonums.

```
(extflsin a) → extflonum?
  a : extflonum?
(extflcos a) → extflonum?
  a : extflonum?
(extfltan a) → extflonum?
  a : extflonum?
(extflasin a) → extflonum?
  a : extflonum?
(extflacos a) → extflonum?
  a : extflonum?
(extflatan a) → extflonum?
  a : extflonum?
(extfllog a) → extflonum?
  a : extflonum?
(extflexp a) → extflonum?
  a : extflonum?
(extflsqrt a) → extflonum?
  a : extflonum?
(extflexpt a b) → extflonum?
  a : extflonum?
  b : extflonum?
```

Like `flsin`, `flcos`, `fltan`, `flasin`, `flacos`, `flatan`, `fllog`, `flexp`, and `flsqrt`, and `flexpt`, but for extflonums.

```
(->extfl a) → extflonum?
  a : exact-integer?
(extfl->exact-integer a) → exact-integer?
  a : extflonum?
(real->extfl a) → extflonum?
  a : real?
(extfl->exact a) → (and/c real? exact?)
  a : real?
(extfl->inexact a) → flonum?
  a : real?
```

The first four are like `->fl`, `fl->exact`, `fl->real`, `inexact->exact`, but for extflonums. The `extfl->inexact` function converts a extflonum to its closest `"flonum"` approximation.

Extflonum Constants

```
pi.t : extflonum?
```

Like `pi`, but with 80 bits precision.

Extflonum Vectors

An *extflvector* is like an flvector, but it holds only extflonums. See also §17.3 “Unsafe Extflonum Operations”.

Two extflvectors are `equal?` if they have the same length, and if the values in corresponding slots of the extflvectors are `equal?`.

```
(extflvector? v) → boolean?
  v : any/c
(extflvector x ...) → extflvector?
  x : extflonum?
(make-extflvector size [x]) → extflvector?
  size : exact-nonnegative-integer?
  x : extflonum? = 0.0
(extflvector-length vec) → exact-nonnegative-integer?
  vec : extflvector?
(extflvector-ref vec pos) → extflonum?
  vec : extflvector?
  pos : exact-nonnegative-integer?
(extflvector-set! vec pos x) → extflonum?
  vec : extflvector?
  pos : exact-nonnegative-integer?
  x : extflonum?
(extflvector-copy vec [start end]) → extflvector?
  vec : extflvector?
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (vector-length v)
```

Like `flvector?`, `flvector`, `make-flvector`, `flvector-length`, `flvector-ref`, `flvector-set`, and `flvector-copy`, but for extflvectors.

```
(in-extflvector vec [start stop step]) → sequence?
  vec : extflvector?
  start : exact-nonnegative-integer? = 0
  stop : (or/c exact-integer? #f) = #f
  step : (and/c exact-integer? (not/c zero?)) = 1
(for/extflvector maybe-length (for-clause ...) body ...)
(for*/extflvector maybe-length (for-clause ...) body ...)

maybe-length =
             | #:length length-expr
             | #:length length-expr #:fill fill-expr

  length-expr : exact-nonnegative-integer?
  fill-expr : extflonum?
```

Like in-flvector, for/flvector, and for*/flvector, but for extflvectors.

```
(make-shared-extflvector size [x]) → extflvector?
  size : exact-nonnegative-integer?
  x : extflonum? = 0.0
```

Like make-shared-flvector, but for extflvectors.

Extflonum Byte Strings

```
(floating-point-bytes->extfl  bstr
                              [big-endian?
                               start
                               end])        → extflonum?
  bstr : bytes?
  big-endian? : any/c = (system-big-endian?)
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (bytes-length bstr)
```

Like floating-point-bytes->real, but for extflonums: Converts the extended-precision floating-point number encoded in *bstr* from position *start* (inclusive) to *end* (exclusive) to an extflonum. The difference between *start* an *end* must be 10 bytes.

```
(extfl->floating-point-bytes  x
                              [big-endian?
                               dest-bstr
                               start])      → bytes?
  x : extflonum?
  big-endian? : any/c = (system-big-endian?)
  dest-bstr :  (and/c bytes? (not/c immutable?))
            = (make-bytes 10)
  start : exact-nonnegative-integer? = 0
```

Like real->floating-point-bytes, but for extflonums: Converts *x* to its representation in a byte string of length 10.

4.3 Strings

§3.4 “Strings (Unicode)” in ***The Racket Guide*** introduces strings.

A *string* is a fixed-length array of characters.

A string can be *mutable* or *immutable*. When an immutable string is provided to a procedure like string-set!, the exn:fail:contract exception is raised. String constants generated by the default reader (see §1.3.7 “Reading Strings”) are immutable, and they are interned in read-syntax mode.

Two strings are equal? when they have the same length and contain the same sequence of characters.

A string can be used as a single-valued sequence (see §4.14.1 “Sequences”). The characters of the string serve as elements of the sequence. See also in-string.

See §1.3.7 “Reading Strings” for information on reading strings and §1.4.6 “Printing Strings” for information on printing strings.

See also: immutable?, symbol->string, bytes->string/utf-8.

4.3.1 String Constructors, Selectors, and Mutators

```
(string? v) → boolean?
  v : any/c
```

Returns #t if *v* is a string, #f otherwise.

Examples:

```
> (string? "Apple")
#t
> (string? 'apple)
#f
```

```
(make-string k [char]) → string?
  k : exact-nonnegative-integer?
  char : char? = #\nul
```

Returns a new mutable string of length *k* where each position in the string is initialized with the character *char*.

Example:

```
> (make-string 5 #\z)
"zzzzz"
```

```
(string char ...) → string?
  char : char?
```

Returns a new mutable string whose length is the number of provided *char*s, and whose positions are initialized with the given *char*s.

Example:

```
> (string #\A #\p #\p #\l #\e)
"Apple"
```

```
(string->immutable-string str) → (and/c string? immutable?)
  str : string?
```

Returns an immutable string with the same content as *str*, returning *str* itself if *str* is immutable.

```
(string-length str) → exact-nonnegative-integer?
  str : string?
```

Returns the length of *str*.

Example:

```
> (string-length "Apple")
5
```

```
(string-ref str k) → char?
  str : string?
  k : exact-nonnegative-integer?
```

Returns the character at position *k* in *str*. The first position in the string corresponds to 0, so the position *k* must be less than the length of the string, otherwise the `exn:fail:contract` exception is raised.

Example:

```
> (string-ref "Apple" 0)
#\A
```

```
(string-set! str k char) → void?
  str : (and/c string? (not/c immutable?))
  k : exact-nonnegative-integer?
  char : char?
```

Changes the character position *k* in *str* to *char*. The first position in the string corresponds to 0, so the position *k* must be less than the length of the string, otherwise the `exn:fail:contract` exception is raised.

Examples:

```
> (define s (string #\A #\p #\p #\l #\e))

> (string-set! s 4 #\y)
```

```
> s
"Apply"
```

```
(substring str start [end]) → string?
  str : string?
  start : exact-nonnegative-integer?
  end : exact-nonnegative-integer? = (string-length str)
```

Returns a new mutable string that is `(- end start)` characters long, and that contains the same characters as *str* from *start* inclusive to *end* exclusive. The first position in a string corresponds to `0`, so the *start* and *end* arguments so they must be less than or equal to the length of *str*, and *end* must be greater than or equal to *start*, otherwise the `exn:fail:contract` exception is raised.

Examples:

```
> (substring "Apple" 1 3)
"pp"
> (substring "Apple" 1)
"pple"
```

```
(string-copy str) → string?
  str : string?
```

Returns `(substring str 0)`.

Examples:

```
> (define s1 "Yui")

> (define pilot (string-copy s1))

> (list s1 pilot)
'("Yui" "Yui")
> (for ([i (in-naturals)] [ch '(#\R #\e #\i)])
    (string-set! pilot i ch))

> (list s1 pilot)
'("Yui" "Rei")
```

```
(string-copy! dest
              dest-start
              src
             [src-start
              src-end])   → void?
  dest : (and/c string? (not/c immutable?))
```

```
  dest-start : exact-nonnegative-integer?
  src : string?
  src-start : exact-nonnegative-integer? = 0
  src-end : exact-nonnegative-integer? = (string-length src)
```

Changes the characters of *dest* starting at position *dest-start* to match the characters in *src* from *src-start* (inclusive) to *src-end* (exclusive), where the first position in a string corresponds to 0. The strings *dest* and *src* can be the same string, and in that case the destination region can overlap with the source region; the destination characters after the copy match the source characters from before the copy. If any of *dest-start*, *src-start*, or *src-end* are out of range (taking into account the sizes of the strings and the source and destination regions), the exn:fail:contract exception is raised.

Examples:

```
> (define s (string #\A #\p #\p #\l #\e))

> (string-copy! s 4 "y")

> (string-copy! s 0 s 3 4)

> s
"lpply"
```

```
(string-fill! dest char) → void?
  dest : (and/c string? (not/c immutable?))
  char : char?
```

Changes *dest* so that every position in the string is filled with *char*.

Examples:

```
> (define s (string #\A #\p #\p #\l #\e))

> (string-fill! s #\q)

> s
"qqqqq"
```

```
(string-append str ...) → string?
  str : string?
```

Returns a new mutable string that is as long as the sum of the given *str*s' lengths, and that contains the concatenated characters of the given *str*s. If no *str*s are provided, the result is a zero-length string.

Example:

```
> (string-append "Apple" "Banana")
"AppleBanana"
```

```
(string->list str) → (listof char?)
  str : string?
```

Returns a new list of characters corresponding to the content of *str*. That is, the length of the list is (string-length *str*), and the sequence of characters in *str* is the same sequence in the result list.

Example:

```
> (string->list "Apple")
'(#\A #\p #\p #\l #\e)
```

```
(list->string lst) → string?
  lst : (listof char?)
```

Returns a new mutable string whose content is the list of characters in *lst*. That is, the length of the string is (length *lst*), and the sequence of characters in *lst* is the same sequence in the result string.

Example:

```
> (list->string (list #\A #\p #\p #\l #\e))
"Apple"
```

```
(build-string n proc) → string?
  n : exact-nonnegative-integer?
  proc : (exact-nonnegative-integer? . -> . char?)
```

Creates a string of *n* characters by applying *proc* to the integers from 0 to (sub1 *n*) in order. If *str* is the resulting string, then (string-ref *str i*) is the character produced by (*proc i*).

Example:

```
> (build-string 5 (lambda (i) (integer->char (+ i 97))))
"abcde"
```

4.3.2 String Comparisons

```
(string=? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Returns `#t` if all of the arguments are `equal?`.

Examples:

```
> (string=? "Apple" "apple")
#f
> (string=? "a" "as" "a")
#f
```

```
(string<? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Returns `#t` if the arguments are lexicographically sorted increasing, where individual characters are ordered by `char<?`, `#f` otherwise.

Examples:

```
> (string<? "Apple" "apple")
#t
> (string<? "apple" "Apple")
#f
> (string<? "a" "b" "c")
#t
```

```
(string<=? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string<?`, but checks whether the arguments are nondecreasing.

Examples:

```
> (string<=? "Apple" "apple")
#t
> (string<=? "apple" "Apple")
#f
> (string<=? "a" "b" "b")
#t
```

```
(string>? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string<?`, but checks whether the arguments are decreasing.

Examples:

```
> (string>? "Apple" "apple")
#f
> (string>? "apple" "Apple")
#t
> (string>? "c" "b" "a")
#t
```

```
(string>=? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string<?`, but checks whether the arguments are nonincreasing.

Examples:

```
> (string>=? "Apple" "apple")
#f
> (string>=? "apple" "Apple")
#t
> (string>=? "c" "b" "b")
#t
```

```
(string-ci=? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Returns `#t` if all of the arguments are `eqv?` after locale-insensitive case-folding via `string-foldcase`.

Examples:

```
> (string-ci=? "Apple" "apple")
#t
> (string-ci=? "a" "a" "a")
#t
```

```
(string-ci<? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string<?`, but checks whether the arguments would be in increasing order if each was first case-folded using `string-foldcase` (which is locale-insensitive).

Examples:

```
> (string-ci<? "Apple" "apple")
#f
```

```
> (string-ci<? "apple" "banana")
#t
> (string-ci<? "a" "b" "c")
#t
```

```
(string-ci<=? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like string-ci<?, but checks whether the arguments would be nondecreasing after case-folding.

Examples:

```
> (string-ci<=? "Apple" "apple")
#t
> (string-ci<=? "apple" "Apple")
#t
> (string-ci<=? "a" "b" "b")
#t
```

```
(string-ci>? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like string-ci<?, but checks whether the arguments would be decreasing after case-folding.

Examples:

```
> (string-ci>? "Apple" "apple")
#f
> (string-ci>? "banana" "Apple")
#t
> (string-ci>? "c" "b" "a")
#t
```

```
(string-ci>=? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like string-ci<?, but checks whether the arguments would be nonincreasing after case-folding.

Examples:

```
> (string-ci>=? "Apple" "apple")
#t
> (string-ci>=? "apple" "Apple")
#t
> (string-ci>=? "c" "b" "b")
#t
```

4.3.3 String Conversions

(string-upcase *str*) → string?
  str : string?

Returns a string whose characters are the upcase conversion of the characters in *str*. The conversion uses Unicode's locale-independent conversion rules that map code-point sequences to code-point sequences (instead of simply mapping a 1-to-1 function on code points over the string), so the string produced by the conversion can be longer than the input string.

Examples:

```
> (string-upcase "abc!")
"ABC!"
> (string-upcase "Straße")
"STRASSE"
```

(string-downcase *string*) → string?
  string : string?

Like string-upcase, but the downcase conversion.

Examples:

```
> (string-downcase "aBC!")
"abc!"
> (string-downcase "Straße")
"straße"
> (string-downcase "KAOΣ")
"καος"
> (string-downcase "Σ")
"σ"
```

(string-titlecase *string*) → string?
  string : string?

Like string-upcase, but the titlecase conversion only for the first character in each sequence of cased characters in str (ignoring case-ignorable characters).

Examples:

```
> (string-titlecase "aBC  twO")
"Abc  Two"
> (string-titlecase "y2k")
"Y2K"
> (string-titlecase "main straße")
"Main Straße"
> (string-titlecase "stra ße")
"Stra Sse"
```

(string-foldcase *string*) → string?
  string : string?

Like string-upcase, but the case-folding conversion.

Examples:

```
> (string-foldcase "aBC!")
"abc!"
> (string-foldcase "Straße")
"strasse"
> (string-foldcase "KAOΣ")
"καοσ"
```

(string-normalize-nfd *string*) → string?
  string : string?

Returns a string that is the Unicode normalized form D of *string*. If the given string is already in the corresponding Unicode normal form, the string may be returned directly as the result (instead of a newly allocated string).

(string-normalize-nfkd *string*) → string?
  string : string?

Like string-normalize-nfd, but for normalized form KD.

(string-normalize-nfc *string*) → string?
  string : string?

Like string-normalize-nfd, but for normalized form C.

(string-normalize-nfkc *string*) → string?
  string : string?

Like string-normalize-nfd, but for normalized form KC.

4.3.4 Locale-Specific String Operations

```
(string-locale=? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string=?`, but the strings are compared in a locale-specific way, based on the value of `current-locale`. See §13.1.1 “Encodings and Locales” for more information on locales.

```
(string-locale<? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string<?`, but the sort order compares strings in a locale-specific way, based on the value of `current-locale`. In particular, the sort order may not be simply a lexicographic extension of character ordering.

```
(string-locale>? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string>?`, but locale-specific like `string-locale<?`.

```
(string-locale-ci=? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string-locale=?`, but strings are compared using rules that are both locale-specific and case-insensitive (depending on what “case-insensitive” means for the current locale).

```
(string-locale-ci<? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string<?`, but both locale-sensitive and case-insensitive like `string-locale-ci=?`.

```
(string-locale-ci>? str1 str2 ...+) → boolean?
  str1 : string?
  str2 : string?
```

Like `string>?`, but both locale-sensitive and case-insensitive like `string-locale-ci=?`.

```
(string-locale-upcase string) → string?
  string : string?
```

Like string-upcase, but using locale-specific case-conversion rules based on the value of current-locale.

```
(string-locale-downcase string) → string?
  string : string?
```

Like string-downcase, but using locale-specific case-conversion rules based on the value of current-locale.

4.3.5 Additional String Functions

```
(require racket/string)
```
package: base

The bindings documented in this section are provided by the racket/string and racket libraries, but not racket/base.

```
(string-append* str ... strs) → string?
  str : string?
  strs : (listof string?)
```

Like string-append, but the last argument is used as a list of arguments for string-append, so (string-append* str ... strs) is the same as (apply string-append str ... strs). In other words, the relationship between string-append and string-append* is similar to the one between list and list*.

Examples:

```
> (string-append* "a" "b" '("c" "d"))
"abcd"
> (string-append* (cdr (append* (map (lambda (x) (list ", " x))
                                     '("Alpha" "Beta" "Gamma")))))
"Alpha, Beta, Gamma"
```

```
(string-join strs
             [sep
              #:before-first before-first
              #:before-last before-last
              #:after-last after-last])   → string?
  strs : (listof string?)
  sep : string? = " "
  before-first : string? = ""
  before-last : string? = sep
  after-last : string? = ""
```

Appends the strings in strs, inserting sep between each pair of strings in strs. before-last, before-first, and after-last are analogous to the inputs of add-between: they

specify an alternate separator between the last two strings, a prefix string, and a suffix string respectively.

Examples:

```
> (string-join '("one" "two" "three" "four"))
"one two three four"
> (string-join '("one" "two" "three" "four") ", ")
"one, two, three, four"
> (string-join '("one" "two" "three" "four") " potato ")
"one potato two potato three potato four"
> (string-join '("x" "y" "z") ", "
               #:before-first "Todo: "
               #:before-last " and "
               #:after-last ".")
"Todo: x, y and z."
```

```
(string-normalize-spaces str
                         [sep
                          space
                          #:trim? trim?
                          #:repeat? repeat?]) → string?
  str : string?
  sep : (or/c string? regexp?) = #px"\\s+"
  space : string? = " "
  trim? : any/c = #t
  repeat? : any/c = #f
```

Normalizes spaces in the input *str* by trimming it (using string-trim and *sep*) and replacing all whitespace sequences in the result with *space*, which defaults to a single space.

Example:

```
> (string-normalize-spaces "  foo bar  baz \r\n\t")
"foo bar baz"
```

The result of (string-normalize-spaces *str sep space*) is the same as (string-join (string-split *str sep* ....) *space*).

```
(string-replace str from to [#:all? all?]) → string?
  str : string?
  from : (or/c string? regexp?)
  to : string?
  all? : any/c = #t
```

Returns *str* with all occurrences of *from* replaced with by *to*. If *from* is a string, it is matched literally (as opposed to being used as a regular expression).

By default, all occurrences are replaced, but only the first match is replaced if *all?* is #f.

Example:

```
> (string-replace "foo bar baz" "bar" "blah")
"foo blah baz"
```

```
(string-split  str
               [sep
                #:trim? trim?
                #:repeat? repeat?]) → (listof string?)
  str : string?
  sep : (or/c string? regexp?) = #px"\\s+"
  trim? : any/c = #t
  repeat? : any/c = #f
```

Splits the input *str* on whitespaces, returning a list of substrings of *str* that are separated by *sep*. The input is first trimmed using *sep* (see string-trim), unless *trim?* is #f. Empty matches are handled in the same way as for regexp-split. As a special case, if *str* is the empty string after trimming, the result is '() instead of '("").

Like string-trim, provide *sep* to use a different separator, and *repeat?* controls matching repeated sequences.

Examples:

```
> (string-split "  foo bar  baz \r\n\t")
'("foo" "bar" "baz")
> (string-split "  ")
'()
> (string-split "  " #:trim? #f)
'("" "")
```

```
(string-trim  str
              [sep
               #:left? left?
               #:right? right?
               #:repeat? repeat?]) → string?
  str : string?
  sep : (or/c string? regexp?) = #px"\\s+"
  left? : any/c = #t
  right? : any/c = #t
  repeat? : any/c = #f
```

Trims the input *str* by removing prefix and suffix *sep*, which defaults to whitespace. A string *sep* is matched literally (as opposed to being used as a regular expression).

Use `#:left? #f` or `#:right? #f` to suppress trimming the corresponding side. When *repeat?* is `#f` (the default), only one match is removed from each side; when *repeat?* is true, all initial or trailing matches are trimmed (which is an alternative to using a regular expression *sep* that contains `+`).

Examples:

```
> (string-trim "  foo bar  baz \r\n\t")
"foo bar  baz"
> (string-trim "  foo bar  baz \r\n\t" " " #:repeat? #t)
"foo bar  baz \r\n\t"
> (string-trim "aaaxaayaa" "aa")
"axaay"
```

4.3.6 Converting Values to Strings

`(require racket/format)` package: `base`

The bindings documented in this section are provided by the `racket/format` and `racket` libraries, but not `racket/base`.

The `racket/format` library provides functions for converting Racket values to strings. In addition to features like padding and numeric formatting, the functions have the virtue of being shorter than `format` (with format string), `number->string`, or `string-append`.

```
(~a  v
     ...
    [#:separator separator
     #:width width
     #:max-width max-width
     #:min-width min-width
     #:limit-marker limit-marker
     #:align align
     #:pad-string pad-string
     #:left-pad-string left-pad-string
     #:right-pad-string right-pad-string]) → string?
  v : any/c
  separator : string? = ""
  width : (or/c exact-nonnegative-integer? #f) = #f
  max-width : (or/c exact-nonnegative-integer? +inf.0)
            = (or width +inf.0)
  min-width : exact-nonnegative-integer? = (or width 0)
  limit-marker : string? = ""
  align : (or/c 'left 'center 'right) = 'left
  pad-string : non-empty-string? = " "
  left-pad-string : non-empty-string? = pad-string
```

right-pad-string : non-empty-string? = *pad-string*

Converts each *v* to a string in `display` mode—that is, like `(format "~a" v)`—then concatentates the results with *separator* between consecutive items, and then pads or truncates the string to be at least *min-width* characters and at most *max-width* characters.

```
> (~a "north")
"north"
> (~a 'south)
"south"
> (~a #"east")
"east"
> (~a #\w "e" 'st)
"west"
> (~a (list "red" 'green #"blue"))
"(red green blue)"
> (~a 17)
"17"
> (~a #e1e20)
"100000000000000000000"
> (~a pi)
"3.141592653589793"
> (~a (expt 6.1 87))
"2.1071509386211452e+68"
```

The `~a` function is primarily useful for strings, numbers, and other atomic data. The `~v` and `~s` functions are better suited to compound data.

Let *s* be the concatenated string forms of the *v*s plus separators. If *s* is longer than *max-width* characters, it is truncated to exactly *max-width* characters. If *s* is shorter than *min-width* characters, it is padded to exactly *min-width* characters. Otherwise *s* is returned unchanged. If *min-width* is greater than *max-width*, an exception is raised.

If *s* is longer than *max-width* characters, it is truncated and the end of the string is replaced with *limit-marker*. If *limit-marker* is longer than *max-width*, an exception is raised.

```
> (~a "abcde" #:max-width 5)
"abcde"
> (~a "abcde" #:max-width 4)
"abcd"
> (~a "abcde" #:max-width 4 #:limit-marker "*")
"abc*"
> (~a "abcde" #:max-width 4 #:limit-marker "...")
"a..."
> (~a "The quick brown fox" #:max-width 15 #:limit-marker "")
"The quick brown"
```

```
> (~a "The quick brown fox" #:max-width 15 #:limit-marker "...")
"The quick br..."
```

If *s* is shorter than *min-width*, it is padded to at least *min-width* characters. If *align* is 'left, then only right padding is added; if *align* is 'right, then only left padding is added; and if *align* is 'center, then roughly equal amounts of left padding and right padding are added.

Padding is specified as a non-empty string. Left padding consists of *left-pad-string* repeated in its entirety as many times as possible followed by a *prefix* of *left-pad-string* to fill the remaining space. In contrast, right padding consists of a *suffix* of *right-pad-string* followed by a number of copies of *right-pad-string* in its entirety. Thus left padding starts with the start of *left-pad-string* and right padding ends with the end of *right-pad-string*.

```
> (~a "apple" #:min-width 20 #:align 'left)
"apple               "
> (~a "pear" #:min-width 20 #:align 'left #:right-pad-string "
.")
"pear . . . . . ."
> (~a "plum" #:min-width 20 #:align 'right #:left-pad-string ".
")
". . . . . plum"
> (~a "orange" #:min-width 20 #:align 'center
               #:left-pad-string "- " #:right-pad-string " -")
"- - - -orange- - - -"
```

Use *width* to set both *max-width* and *min-width* simultaneously, ensuring that the resulting string is exactly *width* characters long:

```
> (~a "terse" #:width 6)
"terse "
> (~a "loquacious" #:width 6)
"loquac"
```

```
(~v  v
     ...
    [#:separator separator
     #:width width
     #:max-width max-width
     #:min-width min-width
     #:limit-marker limit-marker
     #:align align
     #:pad-string pad-string
     #:left-pad-string left-pad-string
     #:right-pad-string right-pad-string]) → string?
```

```
v : any/c
separator : string? = " "
width : (or/c exact-nonnegative-integer? #f) = #f
max-width : (or/c exact-nonnegative-integer? +inf.0)
          = (or width +inf.0)
min-width : exact-nonnegative-integer? = (or width 0)
limit-marker : string? = "..."
align : (or/c 'left 'center 'right) = 'left
pad-string : non-empty-string? = " "
left-pad-string : non-empty-string? = pad-string
right-pad-string : non-empty-string? = pad-string
```

Like ~a, but each value is converted like (format "~v" v), the default separator is " ", and the default limit marker is "...".

```
> (~v "north")
"\"north\""
> (~v 'south)
"'south"
> (~v #"east")
"#\"east\""
> (~v #\w)
"#\\w"
> (~v (list "red" 'green #"blue"))
"'(\"red\" green #\"blue\")"
```

Use ~v to produce text that talks about Racket values.

```
> (let ([nums (for/list ([i 10]) i)])
    (~a "The even numbers in " (~v nums)
        " are " (~v (filter even? nums)) "."))
"The even numbers in '(0 1 2 3 4 5 6 7 8 9) are '(0 2 4 6 8)."
```

```
(~s  v
     ...
    [#:separator separator
     #:width width
     #:max-width max-width
     #:min-width min-width
     #:limit-marker limit-marker
     #:align align
     #:pad-string pad-string
     #:left-pad-string left-pad-string
     #:right-pad-string right-pad-string]) → string?
  v : any/c
```

```
  separator : string? = " "
  width : (or/c exact-nonnegative-integer? #f) = #f
  max-width : (or/c exact-nonnegative-integer? +inf.0)
            = (or width +inf.0)
  min-width : exact-nonnegative-integer? = (or width 0)
  limit-marker : string? = "..."
  align : (or/c 'left 'center 'right) = 'left
  pad-string : non-empty-string? = " "
  left-pad-string : non-empty-string? = pad-string
  right-pad-string : non-empty-string? = pad-string
```

Like ~a, but each value is converted like (format "~s" v), the default separator is " ", and the default limit marker is "...".

```
> (~s "north")
"\"north\""
> (~s 'south)
"south"
> (~s #"east")
"#\"east\""
> (~s #\w)
"#\\w"
> (~s (list "red" 'green #"blue"))
"(\"red\" green #\"blue\")"
```

```
(~e  v
     ...
    [#:separator separator
     #:width width
     #:max-width max-width
     #:min-width min-width
     #:limit-marker limit-marker
     #:align align
     #:pad-string pad-string
     #:left-pad-string left-pad-string
     #:right-pad-string right-pad-string]) → string?
  v : any/c
  separator : string? = " "
  width : (or/c exact-nonnegative-integer? #f) = #f
  max-width : (or/c exact-nonnegative-integer? +inf.0)
            = (or width +inf.0)
  min-width : exact-nonnegative-integer? = (or width 0)
  limit-marker : string? = "..."
  align : (or/c 'left 'center 'right) = 'left
  pad-string : non-empty-string? = " "
```

```
  left-pad-string : non-empty-string? = pad-string
  right-pad-string : non-empty-string? = pad-string
```

Like ~a, but each value is converted like (format "~e" v), the default separator is " ", and the default limit marker is "...".

```
> (~e "north")
"\"north\""
> (~e 'south)
"'south"
> (~e #"east")
"#\"east\""
> (~e #\w)
"#\\w"
> (~e (list "red" 'green #"blue"))
"'(\"red\" green #\"blue\")"
```

```
(~r  x
     [#:sign sign
      #:base base
      #:precision precision
      #:notation notation
      #:format-exponent format-exponent
      #:min-width min-width
      #:pad-string pad-string])          → string?
  x : rational?
          (or/c #f '+ '++ 'parens
  sign :        (let ([ind (or/c string? (list/c string? string?))])
                  (list/c ind ind ind)))
       = #f
  base : (or/c (integer-in 2 36) (list/c 'up (integer-in 2 36)))
       = 10
             (or/c exact-nonnegative-integer?
  precision :       (list/c '= exact-nonnegative-integer?))  = 6
            (or/c 'positional 'exponential
  notation :              (-> rational? (or/c 'positional 'exponential)))
           = 'positional
  format-exponent : (or/c #f string? (-> exact-integer? string?))
                  = #f
  min-width : exact-positive-integer? = 1
  pad-string : non-empty-string? = " "
```

Converts the rational number *x* to a string in either positional or exponential notation, depending on *notation*. The exactness or inexactness of *x* does not affect its formatting.

The optional arguments control number formatting:

- *notation* — determines whether the number is printed in positional or exponential notation. If *notation* is a function, it is applied to *x* to get the notation to be used.

```
> (~r 12345)
"12345"
> (~r 12345 #:notation 'exponential)
"1.2345e+04"
> (let ([pick-notation
         (lambda (x)
           (if (or (< (abs x) 0.001) (> (abs x) 1000))
               'exponential
               'positional))])
    (for/list ([i (in-range 1 5)])
      (~r (expt 17 i) #:notation pick-notation)))
'("17" "289" "4.913e+03" "8.3521e+04")
```

- *precision* — controls the number of digits after the decimal point (or more accurately, the radix point). When *x* is formatted in exponential form, *precision* applies to the significand.

  If *precision* is a natural number, then up to *precision* digits are displayed, but trailing zeroes are dropped, and if all digits after the decimal point are dropped the decimal point is also dropped. If *precision* is `(list '= digits)`, then exactly *digits* digits after the decimal point are used, and the decimal point is never dropped.

```
> (~r pi)
"3.141593"
> (~r pi #:precision 4)
"3.1416"
> (~r pi #:precision 0)
"3"
> (~r 1.5 #:precision 4)
"1.5"
> (~r 1.5 #:precision '(= 4))
"1.5000"
> (~r 50 #:precision 2)
"50"
> (~r 50 #:precision '(= 2))
"50.00"
> (~r 50 #:precision '(= 0))
"50."
```

- *min-width* — if *x* would normally be printed with fewer than *min-width* digits (including the decimal point but not including the sign indicator), the digits are left-padded using *pad-string*.

```
> (~r 17)
```

```
"17"
> (~r 17 #:min-width 4)
"  17"
> (~r -42 #:min-width 4)
"-  42"
> (~r 1.5 #:min-width 4)
" 1.5"
> (~r 1.5 #:precision 4 #:min-width 10)
"       1.5"
> (~r 1.5 #:precision '(= 4) #:min-width 10)
"    1.5000"
> (~r #e1e10 #:min-width 6)
"10000000000"
```

- *pad-string* — specifies the string used to pad the number to at least *min-width* characters (not including the sign indicator). The padding is placed between the sign and the normal digits of *x*.

```
> (~r 17 #:min-width 4 #:pad-string "0")
"0017"
> (~r -42 #:min-width 4 #:pad-string "0")
"-0042"
```

- *sign* — controls how the sign of the number is indicated:
  - If *sign* is `#f` (the default), no sign output is generated if *x* is either positive or zero, and a minus sign is prefixed if *x* is negative.

```
> (for/list ([x '(17 0 -42)]) (~r x))
'("17" "0" "-42")
```

  - If *sign* is `'+`, no sign output is generated if *x* is zero, a plus sign is prefixed if *x* is positive, and a minus sign is prefixed if *x* is negative.

```
> (for/list ([x '(17 0 -42)]) (~r x #:sign '+))
'("+17" "0" "-42")
```

  - If *sign* is `'++`, a plus sign is prefixed if *x* is zero or positive, and a minus sign is prefixed if *x* is negative.

```
> (for/list ([x '(17 0 -42)]) (~r x #:sign '++))
'("+17" "+0" "-42")
```

  - If *sign* is `'parens`, no sign output is generated if *x* is zero or positive, and the number is enclosed in parentheses if *x* is negative.

```
> (for/list ([x '(17 0 -42)]) (~r x #:sign 'parens))
'("17" "0" "(42)")
```

– If *sign* is (list *pos-ind zero-ind neg-ind*), then *pos-ind*, *zero-ind*, and *neg-ind* are used to indicate positive, zero, and negative numbers, respectively. Each indicator is either a string to be used as a prefix or a list containing two strings: a prefix and a suffix.

```
> (let ([sign-table '(("" " up") "an even " ("" "
down"))])
    (for/list ([x '(17 0 -42)]) (~r x #:sign sign-
table)))
'("17 up" "an even 0" "42 down")
```

The default behavior is equivalent to '("" "" "-"); the 'parens mode is equivalent to '("" "" ("(" ")")).

- *base* — controls the base that *x* is formatted in. If *base* is a number greater than 10, then lower-case letters are used. If *base* is (list 'up *base**) and *base** is greater than 10, then upper-case letters are used.

```
> (~r 100 #:base 7)
"202"
> (~r 4.5 #:base 2)
"100.1"
> (~r 3735928559 #:base 16)
"deadbeef"
> (~r 3735928559 #:base '(up 16))
"DEADBEEF"
> (~r 3735928559 #:base '(up 16) #:notation 'exponential)
"D.EADBEF*16^+07"
```

- *format-exponent* — determines how the exponent is displayed.

  If *format-exponent* is a string, the exponent is displayed with an explicit sign (as with a *sign* of '++) and at least two digits, separated from the significand by the "exponent marker" *format-exponent*:

```
> (~r 1234 #:notation 'exponential #:format-exponent "E")
"1.234E+03"
```

  If *format-exponent* is #f, the "exponent marker" is "e" if *base* is 10 and a string involving *base* otherwise:

```
> (~r 1234 #:notation 'exponential)
"1.234e+03"
> (~r 1234 #:notation 'exponential #:base 8)
"2.322*8^+03"
```

  If *format-exponent* is a procedure, it is applied to the exponent and the resulting string is appended to the significand:

```
> (~r 1234 #:notation 'exponential
           #:format-exponent (lambda (e) (format "E~a" e)))
"1.234E3"
```

```
(~.a  v
      ...
      [#:separator separator
       #:width width
       #:max-width max-width
       #:min-width min-width
       #:limit-marker limit-marker
       #:align align
       #:pad-string pad-string
       #:left-pad-string left-pad-string
       #:right-pad-string right-pad-string]) → string?
  v : any/c
  separator : string? = ""
  width : (or/c exact-nonnegative-integer? #f) = #f
  max-width : (or/c exact-nonnegative-integer? +inf.0)
            = (or width +inf.0)
  min-width : exact-nonnegative-integer? = (or width 0)
  limit-marker : string? = ""
  align : (or/c 'left 'center 'right) = 'left
  pad-string : non-empty-string? = " "
  left-pad-string : non-empty-string? = pad-string
  right-pad-string : non-empty-string? = pad-string
```

```
(~.v  v
      ...
     [#:separator separator
      #:width width
      #:max-width max-width
      #:min-width min-width
      #:limit-marker limit-marker
      #:align align
      #:pad-string pad-string
      #:left-pad-string left-pad-string
      #:right-pad-string right-pad-string]) → string?
  v : any/c
  separator : string? = " "
  width : (or/c exact-nonnegative-integer? #f) = #f
  max-width : (or/c exact-nonnegative-integer? +inf.0)
            = (or width +inf.0)
  min-width : exact-nonnegative-integer? = (or width 0)
  limit-marker : string? = "..."
  align : (or/c 'left 'center 'right) = 'left
  pad-string : non-empty-string? = " "
  left-pad-string : non-empty-string? = pad-string
  right-pad-string : non-empty-string? = pad-string
(~.s  v
      ...
     [#:separator separator
      #:width width
      #:max-width max-width
      #:min-width min-width
      #:limit-marker limit-marker
      #:align align
      #:pad-string pad-string
      #:left-pad-string left-pad-string
      #:right-pad-string right-pad-string]) → string?
  v : any/c
  separator : string? = " "
  width : (or/c exact-nonnegative-integer? #f) = #f
  max-width : (or/c exact-nonnegative-integer? +inf.0)
            = (or width +inf.0)
  min-width : exact-nonnegative-integer? = (or width 0)
  limit-marker : string? = "..."
  align : (or/c 'left 'center 'right) = 'left
  pad-string : non-empty-string? = " "
  left-pad-string : non-empty-string? = pad-string
  right-pad-string : non-empty-string? = pad-string
```

Like ~a, ~v, and ~s, but each v is formatted like (format "~.a" v), (format "~.v" v), and (format "~.s" v), respectively.

4.4 Byte Strings

§3.5 “Bytes and Byte Strings” in ***The Racket Guide*** introduces byte strings.

A *byte string* is a fixed-length array of bytes. A *byte* is an exact integer between 0 and 255 inclusive.

A byte string can be *mutable* or *immutable*. When an immutable byte string is provided to a procedure like bytes-set!, the exn:fail:contract exception is raised. Byte-string constants generated by the default reader (see §1.3.7 “Reading Strings”) are immutable, and they are interned in read-syntax mode.

Two byte strings are equal? when they have the same length and contain the same sequence of bytes.

A byte string can be used as a single-valued sequence (see §4.14.1 “Sequences”). The bytes of the string serve as elements of the sequence. See also in-bytes.

See §1.3.7 “Reading Strings” for information on reading byte strings and §1.4.6 “Printing Strings” for information on printing byte strings.

See also: immutable?.

4.4.1 Byte String Constructors, Selectors, and Mutators

```
(bytes? v) → boolean?
  v : any/c
```

Returns #t if v is a byte string, #f otherwise.

Examples:

```
> (bytes? #"Apple")
#t
> (bytes? "Apple")
#f
```

```
(make-bytes k [b]) → bytes?
  k : exact-nonnegative-integer?
  b : byte? = 0
```

Returns a new mutable byte string of length k where each position in the byte string is initialized with the byte b.

Example:

```
> (make-bytes 5 65)
#"AAAAA"
```

```
(bytes b ...) → bytes?
  b : byte?
```

Returns a new mutable byte string whose length is the number of provided `b`s, and whose positions are initialized with the given `b`s.

Example:

```
> (bytes 65 112 112 108 101)
#"Apple"
```

```
(bytes->immutable-bytes bstr) → (and/c bytes? immutable?)
  bstr : bytes?
```

Returns an immutable byte string with the same content as *bstr*, returning *bstr* itself if *bstr* is immutable.

Examples:

```
> (bytes->immutable-bytes (bytes 65 65 65))
#"AAA"
> (define b (bytes->immutable-bytes (make-bytes 5 65)))

> (bytes->immutable-bytes b)
#"AAAAA"
> (eq? (bytes->immutable-bytes b) b)
#t
```

```
(byte? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a byte (i.e., an exact integer between `0` and `255` inclusive), `#f` otherwise.

Examples:

```
> (byte? 65)
#t
> (byte? 0)
#t
> (byte? 256)
#f
> (byte? -1)
#f
```

```
(bytes-length bstr) → exact-nonnegative-integer?
  bstr : bytes?
```

Returns the length of *bstr*.

Example:

```
> (bytes-length #"Apple")
5
```

```
(bytes-ref bstr k) → byte?
  bstr : bytes?
  k : exact-nonnegative-integer?
```

Returns the character at position *k* in *bstr*. The first position in the bytes coore-sponds to 0, so the position *k* must be less than the length of the bytes, otherwise the exn:fail:contract exception is raised.

Example:

```
> (bytes-ref #"Apple" 0)
65
```

```
(bytes-set! bstr k b) → void?
  bstr : (and/c bytes? (not/c immutable?))
  k : exact-nonnegative-integer?
  b : byte?
```

Changes the character position *k* in *bstr* to *b*. The first position in the byte string coore-sponds to 0, so the position *k* must be less than the length of the bytes, otherwise the exn:fail:contract exception is raised.

Examples:

```
> (define s (bytes 65 112 112 108 101))

> (bytes-set! s 4 121)

> s
#"Apply"
```

```
(subbytes bstr start [end]) → bytes?
  bstr : bytes?
  start : exact-nonnegative-integer?
  end : exact-nonnegative-integer? = (bytes-length str)
```

Returns a new mutable byte string that is `(- end start)` bytes long, and that contains the same bytes as *bstr* from *start* inclusive to *end* exclusive. The *start* and *end* arguments must be less than or equal to the length of *bstr*, and *end* must be greater than or equal to *start*, otherwise the `exn:fail:contract` exception is raised.

Examples:

```
> (subbytes #"Apple" 1 3)
#"pp"
> (subbytes #"Apple" 1)
#"pple"
```

```
(bytes-copy bstr) → bytes?
  bstr : bytes?
```

Returns `(subbytes str 0)`.

```
(bytes-copy!  dest
              dest-start
              src
             [src-start
              src-end])   → void?
  dest : (and/c bytes? (not/c immutable?))
  dest-start : exact-nonnegative-integer?
  src : bytes?
  src-start : exact-nonnegative-integer? = 0
  src-end : exact-nonnegative-integer? = (bytes-length src)
```

Changes the bytes of *dest* starting at position *dest-start* to match the bytes in *src* from *src-start* (inclusive) to *src-end* (exclusive). The bytes strings *dest* and *src* can be the same byte string, and in that case the destination region can overlap with the source region; the destination bytes after the copy match the source bytes from before the copy. If any of *dest-start*, *src-start*, or *src-end* are out of range (taking into account the sizes of the bytes strings and the source and destination regions), the `exn:fail:contract` exception is raised.

Examples:

```
> (define s (bytes 65 112 112 108 101))

> (bytes-copy! s 4 #"y")

> (bytes-copy! s 0 s 3 4)

> s
#"lpply"
```

```
(bytes-fill! dest b) → void?
  dest : (and/c bytes? (not/c immutable?))
  b : byte?
```

Changes *dest* so that every position in the bytes is filled with *b*.

Examples:

```
> (define s (bytes 65 112 112 108 101))

> (bytes-fill! s 113)

> s
#"qqqqq"
```

```
(bytes-append bstr ...) → bytes?
  bstr : bytes?
```

Returns a new mutable byte string that is as long as the sum of the given *bstr*s' lengths, and that contains the concatenated bytes of the given *bstr*s. If no *bstr*s are provided, the result is a zero-length byte string.

Example:

```
> (bytes-append #"Apple" #"Banana")
#"AppleBanana"
```

```
(bytes->list bstr) → (listof byte?)
  bstr : bytes?
```

Returns a new list of bytes corresponding to the content of *bstr*. That is, the length of the list is `(bytes-length bstr)`, and the sequence of bytes in *bstr* is the same sequence in the result list.

Example:

```
> (bytes->list #"Apple")
'(65 112 112 108 101)
```

```
(list->bytes lst) → bytes?
  lst : (listof byte?)
```

Returns a new mutable byte string whose content is the list of bytes in *lst*. That is, the length of the byte string is `(length lst)`, and the sequence of bytes in *lst* is the same sequence in the result byte string.

Example:

```
> (list->bytes (list 65 112 112 108 101))
#"Apple"
```

```
(make-shared-bytes k [b]) → bytes?
  k : exact-nonnegative-integer?
  b : byte? = 0
```

Returns a new mutable byte string of length *k* where each position in the byte string is initialized with the byte *b*. For communication among places, the new byte string is allocated in the shared memory space.

Example:

```
> (make-shared-bytes 5 65)
#"AAAAA"
```

```
(shared-bytes b ...) → bytes?
  b : byte?
```

Returns a new mutable byte string whose length is the number of provided *b*s, and whose positions are initialized with the given *b*s. For communication among places, the new byte string is allocated in the shared memory space.

Example:

```
> (shared-bytes 65 112 112 108 101)
#"Apple"
```

4.4.2 Byte String Comparisons

```
(bytes=? bstr1 bstr2 ...+) → boolean?
  bstr1 : bytes?
  bstr2 : bytes?
```

Returns `#t` if all of the arguments are `eqv?`.

Examples:

```
> (bytes=? #"Apple" #"apple")
#f
> (bytes=? #"a" #"as" #"a")
#f
```

```
(bytes<? bstr1 bstr2 ...+) → boolean?
  bstr1 : bytes?
  bstr2 : bytes?
```

Returns `#t` if the arguments are lexicographically sorted increasing, where individual bytes are ordered by `<`, `#f` otherwise.

Examples:

```
> (bytes<? #"Apple" #"apple")
#t
> (bytes<? #"apple" #"Apple")
#f
> (bytes<? #"a" #"b" #"c")
#t
```

```
(bytes>? bstr1 bstr2 ...+) → boolean?
  bstr1 : bytes?
  bstr2 : bytes?
```

Like `bytes<?`, but checks whether the arguments are decreasing.

Examples:

```
> (bytes>? #"Apple" #"apple")
#f
> (bytes>? #"apple" #"Apple")
#t
> (bytes>? #"c" #"b" #"a")
#t
```

4.4.3 Bytes to/from Characters, Decoding and Encoding

```
(bytes->string/utf-8 bstr [err-char start end]) → string?
  bstr : bytes?
  err-char : (or/c #f char?) = #f
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (bytes-length bstr)
```

Produces a string by decoding the *start* to *end* substring of *bstr* as a UTF-8 encoding of Unicode code points. If *err-char* is not `#f`, then it is used for bytes that fall in the range `128` to `255` but are not part of a valid encoding sequence. (This rule is consistent with reading characters from a port; see §13.1.1 “Encodings and Locales” for more details.) If *err-char* is `#f`, and if the *start* to *end* substring of *bstr* is not a valid UTF-8 encoding overall, then the `exn:fail:contract` exception is raised.

Example:

```
> (bytes->string/utf-8 (bytes 195 167 195 176 195 182 194 163))
"çðö£"
```

```
(bytes->string/locale bstr
                      [err-char
                       start
                       end])     → string?
  bstr : bytes?
  err-char : (or/c #f char?) = #f
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (bytes-length bstr)
```

Produces a string by decoding the *start* to *end* substring of *bstr* using the current locale's encoding (see also §13.1.1 "Encodings and Locales"). If *err-char* is not `#f`, it is used for each byte in *bstr* that is not part of a valid encoding; if *err-char* is `#f`, and if the *start* to *end* substring of *bstr* is not a valid encoding overall, then the `exn:fail:contract` exception is raised.

```
(bytes->string/latin-1 bstr
                       [err-char
                        start
                        end])     → string?
  bstr : bytes?
  err-char : (or/c #f char?) = #f
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (bytes-length bstr)
```

Produces a string by decoding the *start* to *end* substring of *bstr* as a Latin-1 encoding of Unicode code points; i.e., each byte is translated directly to a character using `integer->char`, so the decoding always succeeds. The *err-char* argument is ignored, but present for consistency with the other operations.

Example:

```
> (bytes->string/latin-1 (bytes 254 211 209 165))
"þÓÑ¥"
```

```
(string->bytes/utf-8 str [err-byte start end]) → bytes?
  str : string?
  err-byte : (or/c #f byte?) = #f
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (string-length str)
```

Produces a byte string by encoding the *start* to *end* substring of *str* via UTF-8 (always succeeding). The *err-byte* argument is ignored, but included for consistency with the other operations.

Examples:

```
> (define b
    (bytes->string/utf-8
     (bytes 195 167 195 176 195 182 194 163)))

> (string->bytes/utf-8 b)
#"\303\247\303\260\303\266\302\243"
> (bytes->string/utf-8 (string->bytes/utf-8 b))
"çðö£"
```

```
(string->bytes/locale str [err-byte start end]) → bytes?
  str : string?
  err-byte : (or/c #f byte?) = #f
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (string-length str)
```

Produces a string by encoding the *start* to *end* substring of *str* using the current locale's encoding (see also §13.1.1 "Encodings and Locales"). If *err-byte* is not `#f`, it is used for each character in *str* that cannot be encoded for the current locale; if *err-byte* is `#f`, and if the *start* to *end* substring of *str* cannot be encoded, then the `exn:fail:contract` exception is raised.

```
(string->bytes/latin-1  str
                        [err-byte
                         start
                         end])     → bytes?
  str : string?
  err-byte : (or/c #f byte?) = #f
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (string-length str)
```

Produces a string by encoding the *start* to *end* substring of *str* using Latin-1; i.e., each character is translated directly to a byte using `char->integer`. If *err-byte* is not `#f`, it is used for each character in *str* whose value is greater than 255. If *err-byte* is `#f`, and if the *start* to *end* substring of *str* has a character with a value greater than 255, then the `exn:fail:contract` exception is raised.

Examples:

```
> (define b
    (bytes->string/latin-1 (bytes 254 211 209 165)))

> (string->bytes/latin-1 b)
#"\376\323\321\245"
> (bytes->string/latin-1 (string->bytes/latin-1 b))
"þÓÑ¥"
```

```
(string-utf-8-length str [start end]) → exact-nonnegative-integer?
  str : string?
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (string-length str)
```

Returns the length in bytes of the UTF-8 encoding of *str*'s substring from *start* to *end*, but without actually generating the encoded bytes.

Examples:

```
> (string-utf-8-length
    (bytes->string/utf-8 (bytes 195 167 195 176 195 182 194 163)))
8
> (string-utf-8-length "hello")
5
```

```
(bytes-utf-8-length bstr [err-char start end])
 → exact-nonnegative-integer?
  bstr : bytes?
  err-char : (or/c #f char?) = #f
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (bytes-length bstr)
```

Returns the length in characters of the UTF-8 decoding of *bstr*'s substring from *start* to *end*, but without actually generating the decoded characters. If *err-char* is `#f` and the substring is not a UTF-8 encoding overall, the result is `#f`. Otherwise, *err-char* is used to resolve decoding errors as in `bytes->string/utf-8`.

Examples:

```
> (bytes-utf-8-length (bytes 195 167 195 176 195 182 194 163))
4
> (bytes-utf-8-length (make-bytes 5 65))
5
```

```
(bytes-utf-8-ref bstr [skip err-char start end]) → char?
  bstr : bytes?
  skip : exact-nonnegative-integer? = 0
  err-char : (or/c #f char?) = #f
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (bytes-length bstr)
```

Returns the *skip*th character in the UTF-8 decoding of *bstr*'s substring from *start* to *end*, but without actually generating the other decoded characters. If the substring is not a UTF-8 encoding up to the *skip*th character (when *err-char* is `#f`), or if the substring decoding produces fewer than *skip* characters, the result is `#f`. If *err-char* is not `#f`, it is used to resolve decoding errors as in `bytes->string/utf-8`.

Examples:

```
> (bytes-utf-8-ref (bytes 195 167 195 176 195 182 194 163) 0)
#\ç
> (bytes-utf-8-ref (bytes 195 167 195 176 195 182 194 163) 1)
#\ð
> (bytes-utf-8-ref (bytes 195 167 195 176 195 182 194 163) 2)
#\ö
> (bytes-utf-8-ref (bytes 65 66 67 68) 0)
#\A
> (bytes-utf-8-ref (bytes 65 66 67 68) 1)
#\B
> (bytes-utf-8-ref (bytes 65 66 67 68) 2)
#\C
```

```
(bytes-utf-8-index  bstr
                   [skip
                    err-char
                    start
                    end])     → exact-nonnegative-integer?
  bstr : bytes?
  skip : exact-nonnegative-integer? = 0
  err-char : (or/c #f char?) = #f
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (bytes-length bstr)
```

Returns the offset in bytes into *bstr* at which the *skip*th character's encoding starts in the UTF-8 decoding of *bstr*'s substring from *start* to *end* (but without actually generating the other decoded characters). The result is relative to the start of *bstr*, not to *start*. If the substring is not a UTF-8 encoding up to the *skip*th character (when *err-char* is `#f`), or if the substring decoding produces fewer than *skip* characters, the result is `#f`. If *err-char* is not `#f`, it is used to resolve decoding errors as in `bytes->string/utf-8`.

Examples:

```
> (bytes-utf-8-index (bytes 195 167 195 176 195 182 194 163) 0)
0
> (bytes-utf-8-index (bytes 195 167 195 176 195 182 194 163) 1)
2
> (bytes-utf-8-index (bytes 195 167 195 176 195 182 194 163) 2)
4
> (bytes-utf-8-index (bytes 65 66 67 68) 0)
0
> (bytes-utf-8-index (bytes 65 66 67 68) 1)
1
> (bytes-utf-8-index (bytes 65 66 67 68) 2)
2
```

4.4.4 Bytes to Bytes Encoding Conversion

```
(bytes-open-converter from-name to-name)
 → (or/c bytes-converter? #f)
  from-name : string?
  to-name : string?
```

Produces a *byte converter* to go from the encoding named by `from-name` to the encoding named by `to-name`. If the requested conversion pair is not available, `#f` is returned instead of a converter.

Certain encoding combinations are always available:

- `(bytes-open-converter "UTF-8" "UTF-8")` — the identity conversion, except that encoding errors in the input lead to a decoding failure.
- `(bytes-open-converter "UTF-8-permissive" "UTF-8")` — the identity conversion, except that any input byte that is not part of a valid encoding sequence is effectively replaced by the UTF-8 encoding sequence for `#\uFFFD`. (This handling of invalid sequences is consistent with the interpretation of port bytes streams into characters; see §13.1 “Ports”.)
- `(bytes-open-converter "" "UTF-8")` — converts from the current locale’s default encoding (see §13.1.1 “Encodings and Locales”) to UTF-8.
- `(bytes-open-converter "UTF-8" "")` — converts from UTF-8 to the current locale’s default encoding (see §13.1.1 “Encodings and Locales”).
- `(bytes-open-converter "platform-UTF-8" "platform-UTF-16")` — converts UTF-8 to UTF-16 on Unix and Mac OS X, where each UTF-16 code unit is a sequence of two bytes ordered by the current platform’s endianness. On Windows, the input can include encodings that are not valid UTF-8, but which naturally extend the UTF-8 encoding to support unpaired surrogate code units, and the output is a sequence of UTF-16 code units (as little-endian byte pairs), potentially including unpaired surrogates.
- `(bytes-open-converter "platform-UTF-8-permissive" "platform-UTF-16")` — like `(bytes-open-converter "platform-UTF-8" "platform-UTF-16")`, but an input byte that is not part of a valid UTF-8 encoding sequence (or valid for the unpaired-surrogate extension on Windows) is effectively replaced with `(char->integer #\?)`.
- `(bytes-open-converter "platform-UTF-16" "platform-UTF-8")` — converts UTF-16 (bytes ordered by the current platform’s endianness) to UTF-8 on Unix and Mac OS X. On Windows, the input can include UTF-16 code units that are unpaired surrogates, and the corresponding output includes an encoding of each surrogate in a natural extension of UTF-8. On Unix and Mac OS X, surrogates are assumed

to be paired: a pair of bytes with the bits 55296 starts a surrogate pair, and the 1023 bits are used from the pair and following pair (independent of the value of the 56320 bits). On all platforms, performance may be poor when decoding from an odd offset within an input byte string.

A newly opened byte converter is registered with the current custodian (see §14.7 “Custodians”), so that the converter is closed when the custodian is shut down. A converter is not registered with a custodian (and does not need to be closed) if it is one of the guaranteed combinations not involving "" on Unix, or if it is any of the guaranteed combinations (including "") on Windows and Mac OS X.

In the Racket software distributions for Windows, a suitable "iconv.dll" is included with "libmzsch*VERS*.dll".

The set of available encodings and combinations varies by platform, depending on the iconv library that is installed; the *from-name* and *to-name* arguments are passed on to iconv_open. On Windows, "iconv.dll" or "libiconv.dll" must be in the same directory as "libmzsch*VERS*.dll" (where *VERS* is a version number), in the user’s path, in the system directory, or in the current executable’s directory at run time, and the DLL must either supply _errno or link to "msvcrt.dll" for _errno; otherwise, only the guaranteed combinations are available.

Use bytes-convert with the result to convert byte strings.

```
(bytes-close-converter converter) → void
  converter : bytes-converter?
```

Closes the given converter, so that it can no longer be used with bytes-convert or bytes-convert-end.

```
(bytes-convert  converter
                src-bstr
               [src-start-pos
                src-end-pos
                dest-bstr
                dest-start-pos
                dest-end-pos])
   (or/c bytes? exact-nonnegative-integer?)
 → exact-nonnegative-integer?
   (or/c 'complete 'continues 'aborts 'error)
  converter : bytes-converter?
  src-bstr : bytes?
  src-start-pos : exact-nonnegative-integer? = 0
  src-end-pos : exact-nonnegative-integer?
              = (bytes-length src-bstr)
  dest-bstr : (or/c bytes? #f) = #f
  dest-start-pos : exact-nonnegative-integer? = 0
  dest-end-pos : (or/c exact-nonnegative-integer? #f)
               = (and dest-bstr
                      (bytes-length dest-bstr))
```

Converts the bytes from *src-start-pos* to *src-end-pos* in *src-bstr*.

If *dest-bstr* is not `#f`, the converted bytes are written into *dest-bstr* from *dest-start-pos* to *dest-end-pos*. If *dest-bstr* is `#f`, then a newly allocated byte string holds the conversion results, and if *dest-end-pos* is not `#f`, the size of the result byte string is no more than `(- dest-end-pos dest-start-pos)`.

The result of `bytes-convert` is three values:

- *result-bstr* or *dest-wrote-amt* — a byte string if *dest-bstr* is `#f` or not provided, or the number of bytes written into *dest-bstr* otherwise.
- *src-read-amt* — the number of bytes successfully converted from *src-bstr*.
- `'complete`, `'continues`, `'aborts`, or `'error` — indicates how conversion terminated:
  - `'complete`: The entire input was processed, and *src-read-amt* will be equal to `(- src-end-pos src-start-pos)`.
  - `'continues`: Conversion stopped due to the limit on the result size or the space in *dest-bstr*; in this case, fewer than `(- dest-end-pos dest-start-pos)` bytes may be returned if more space is needed to process the next complete encoding sequence in *src-bstr*.
  - `'aborts`: The input stopped part-way through an encoding sequence, and more input bytes are necessary to continue. For example, if the last byte of input is `195` for a `"UTF-8-permissive"` decoding, the result is `'aborts`, because another byte is needed to determine how to use the `195` byte.
  - `'error`: The bytes starting at `(+ src-start-pos src-read-amt)` bytes in *src-bstr* do not form a legal encoding sequence. This result is never produced for some encodings, where all byte sequences are valid encodings. For example, since `"UTF-8-permissive"` handles an invalid UTF-8 sequence by dropping characters or generating "?," every byte sequence is effectively valid.

Applying a converter accumulates state in the converter (even when the third result of `bytes-convert` is `'complete`). This state can affect both further processing of input and further generation of output, but only for conversions that involve "shift sequences" to change modes within a stream. To terminate an input sequence and reset the converter, use `bytes-convert-end`.

Examples:

```
> (define convert (bytes-open-converter "UTF-8" "UTF-16"))

> (bytes-convert convert (bytes 65 66 67 68))
#"\377\376A\0B\0C\0D\0"
4
'complete
```

```
> (bytes 195 167 195 176 195 182 194 163)
#"\303\247\303\260\303\266\302\243"
> (bytes-convert convert (bytes 195 167 195 176 195 182 194 163))
#"\347\0\360\0\366\0\243\0"
8
'complete
> (bytes-close-converter convert)
```

```
(bytes-convert-end  converter
                   [dest-bstr
                    dest-start-pos
                    dest-end-pos])
 → (or/c bytes? exact-nonnegative-integer?)
   (or/c 'complete 'continues)
  converter : bytes-converter?
  dest-bstr : (or/c bytes? #f) = #f
  dest-start-pos : exact-nonnegative-integer? = 0
  dest-end-pos : (or/c exact-nonnegative-integer? #f)
               = (and dest-bstr
                      (bytes-length dest-bstr))
```

Like `bytes-convert`, but instead of converting bytes, this procedure generates an ending sequence for the conversion (sometimes called a "shift sequence"), if any. Few encodings use shift sequences, so this function will succeed with no output for most encodings. In any case, successful output of a (possibly empty) shift sequence resets the converter to its initial state.

The result of `bytes-convert-end` is two values:

- *result-bstr* or *dest-wrote-amt* — a byte string if *dest-bstr* is `#f` or not provided, or the number of bytes written into *dest-bstr* otherwise.
- `'complete` or `'continues` — indicates whether conversion completed. If `'complete`, then an entire ending sequence was produced. If `'continues`, then the conversion could not complete due to the limit on the result size or the space in *dest-bstr*, and the first result is either an empty byte string or `0`.

```
(bytes-converter? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a byte converter produced by `bytes-open-converter`, `#f` otherwise.

Examples:

```
> (bytes-converter? (bytes-open-converter "UTF-8" "UTF-16"))
#t
> (bytes-converter? (bytes-open-converter "whacky" "not likely"))
#f
> (define b (bytes-open-converter "UTF-8" "UTF-16"))

> (bytes-close-converter b)

> (bytes-converter? b)
#t
```

(locale-string-encoding) → any

Returns a string for the current locale's encoding (i.e., the encoding normally identified by ""). See also system-language+country.

4.4.5 Additional Byte String Functions

(require racket/bytes) package: base

The bindings documented in this section are provided by the racket/bytes and racket libraries, but not racket/base.

(bytes-append* *str* ... *strs*) → bytes?
  str : bytes?
  strs : (listof bytes?)

Like bytes-append, but the last argument is used as a list of arguments for bytes-append, so (bytes-append* *str* ... *strs*) is the same as (apply bytes-append *str* ... *strs*). In other words, the relationship between bytes-append and bytes-append* is similar to the one between list and list*.

Examples:

```
> (bytes-append* #"a" #"b" '(#"c" #"d"))
#"abcd"
> (bytes-append* (cdr (append* (map (lambda (x) (list #", " x))
                                    '(#"Alpha" #"Beta" #"Gamma")))))
#"Alpha, Beta, Gamma"
```

(bytes-join *strs* *sep*) → bytes?
  strs : (listof bytes?)
  sep : bytes?

Appends the byte strings in *strs*, inserting *sep* between each pair of bytes in *strs*.

Example:

```
> (bytes-join '(#"one" #"two" #"three" #"four") #" potato ")
#"one potato two potato three potato four"
```

4.5 Characters

§3.3 “Characters” in ***The Racket Guide*** introduces characters.

Characters range over Unicode scalar values, which includes characters whose values range from `#x0` to `#x10FFFF`, but not including `#xD800` to `#xDFFF`. The scalar values are a subset of the Unicode code points.

Two characters are `eqv?` if they correspond to the same scalar value. For each scalar value less than 256, character values that are `eqv?` are also `eq?`. Characters produced by the default reader are interned in `read-syntax` mode.

See §1.3.14 “Reading Characters” for information on `read`ing characters and §1.4.11 “Printing Characters” for information on `print`ing characters.

4.5.1 Characters and Scalar Values

```
(char? v) → boolean?
  v : any/c
```

Return `#t` if *v* is a character, `#f` otherwise.

```
(char->integer char) → exact-integer?
  char : char?
```

Returns a character’s code-point number.

Example:

```
> (char->integer #\A)
65
```

```
(integer->char k) → char?
      (and/c exact-integer?
  k :        (or/c (integer-in 0 55295)
                   (integer-in 57344 1114111)))
```

Return the character whose code-point number is *k*. For *k* less than `256`, the result is the same object for the same *k*.

Example:

```
> (integer->char 65)
#\A
```

```
(char-utf-8-length char) → (integer-in 1 6)
  char : char?
```

Produces the same result as `(bytes-length (string->bytes/utf-8 (string char)))`.

4.5.2 Character Comparisons

```
(char=? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Returns `#t` if all of the arguments are `eqv?`.

Examples:

```
> (char=? #\a #\a)
#t
> (char=? #\a #\A #\a)
#f
```

```
(char<? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Returns `#t` if the arguments are sorted increasing, where two characters are ordered by their scalar values, `#f` otherwise.

Examples:

```
> (char<? #\A #\a)
#t
> (char<? #\a #\A)
#f
> (char<? #\a #\b #\c)
#t
```

```
(char<=? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Like `char<?`, but checks whether the arguments are nondecreasing.

Examples:

```
> (char<=? #\A #\a)
#t
> (char<=? #\a #\A)
#f
> (char<=? #\a #\b #\b)
#t
```

```
(char>? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Like `char<?`, but checks whether the arguments are decreasing.

Examples:

```
> (char>? #\A #\a)
#f
> (char>? #\a #\A)
#t
> (char>? #\c #\b #\a)
#t
```

```
(char>=? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Like `char<?`, but checks whether the arguments are nonincreasing.

Examples:

```
> (char>=? #\A #\a)
#f
> (char>=? #\a #\A)
#t
> (char>=? #\c #\b #\b)
#t
```

```
(char-ci=? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Returns `#t` if all of the arguments are `eqv?` after locale-insensitive case-folding via `char-foldcase`.

Examples:

```
> (char-ci=? #\A #\a)
#t
> (char-ci=? #\a #\a #\a)
#t
```

```
(char-ci<? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Like `char<?`, but checks whether the arguments would be in increasing order if each was first case-folded using `char-foldcase` (which is locale-insensitive).

Examples:

```
> (char-ci<? #\A #\a)
#f
> (char-ci<? #\a #\b)
#t
> (char-ci<? #\a #\b #\c)
#t
```

```
(char-ci<=? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Like `char-ci<?`, but checks whether the arguments would be nondecreasing after case-folding.

Examples:

```
> (char-ci<=? #\A #\a)
#t
> (char-ci<=? #\a #\A)
#t
> (char-ci<=? #\a #\b #\b)
#t
```

```
(char-ci>? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Like `char-ci<?`, but checks whether the arguments would be decreasing after case-folding.

Examples:

```
> (char-ci>? #\A #\a)
#f
```

```
> (char-ci>? #\b #\A)
#t
> (char-ci>? #\c #\b #\a)
#t
```

```
(char-ci>=? char1 char2 ...+) → boolean?
  char1 : char?
  char2 : char?
```

Like `char-ci<?`, but checks whether the arguments would be nonincreasing after case-folding.

Examples:

```
> (char-ci>=? #\A #\a)
#t
> (char-ci>=? #\a #\A)
#t
> (char-ci>=? #\c #\b #\b)
#t
```

4.5.3 Classifications

```
(char-alphabetic? char) → boolean?
  char : char?
```

Returns `#t` if *char* has the Unicode "Alphabetic" property.

```
(char-lower-case? char) → boolean?
  char : char?
```

Returns `#t` if *char* has the Unicode "Lowercase" property.

```
(char-upper-case? char) → boolean?
  char : char?
```

Returns `#t` if *char* has the Unicode "Uppercase" property.

```
(char-title-case? char) → boolean?
  char : char?
```

Returns `#t` if *char*'s Unicode general category is Lt, `#f` otherwise.

```
(char-numeric? char) → boolean?
  char : char?
```

Returns `#t` if *char* has the Unicode "Numeric" property.

```
(char-symbolic? char) → boolean?
  char : char?
```

Returns `#t` if *char*'s Unicode general category is Sm, Sc, Sk, or So, `#f` otherwise.

```
(char-punctuation? char) → boolean?
  char : char?
```

Returns `#t` if *char*'s Unicode general category is Pc, Pd, Ps, Pe, Pi, Pf, or Po, `#f` otherwise.

```
(char-graphic? char) → boolean?
  char : char?
```

Returns `#t` if *char*'s Unicode general category is Ll, Lm, Lo, Lt, Lu, Nd, Nl, No, Mn, Mc, or Me, or if one of the following produces `#t` when applied to *char*: `char-alphabetic?`, `char-numeric?`, `char-symbolic?`, or `char-punctuation?`.

```
(char-whitespace? char) → boolean?
  char : char?
```

Returns `#t` if *char* has the Unicode "White_Space" property.

```
(char-blank? char) → boolean?
  char : char?
```

Returns `#t` if *char*'s Unicode general category is Zs or if *char* is `#\tab`. (These correspond to horizontal whitespace.)

```
(char-iso-control? char) → boolean?
  char : char?
```

Return `#t` if *char* is between `#\nul` and `#\u001F` inclusive or `#\rubout` and `#\u009F` inclusive.

```
(char-general-category char) → symbol?
  char : char?
```

Returns a symbol representing the character's Unicode general category, which is `'lu`, `'ll`, `'lt`, `'lm`, `'lo`, `'mn`, `'mc`, `'me`, `'nd`, `'nl`, `'no`, `'ps`, `'pe`, `'pi`, `'pf`, `'pd`, `'pc`, `'po`, `'sc`, `'sm`, `'sk`, `'so`, `'zs`, `'zp`, `'zl`, `'cc`, `'cf`, `'cs`, `'co`, or `'cn`.

```
(make-known-char-range-list)
    (listof (list/c exact-nonnegative-integer?
 →                  exact-nonnegative-integer?
                    boolean?))
```

Produces a list of three-element lists, where each three-element list represents a set of consecutive code points for which the Unicode standard specifies character properties. Each three-element list contains two integers and a boolean; the first integer is a starting code-point value (inclusive), the second integer is an ending code-point value (inclusive), and the boolean is `#t` when all characters in the code-point range have identical results for all of the character predicates above. The three-element lists are ordered in the overall result list such that later lists represent larger code-point values, and all three-element lists are separated from every other by at least one code-point value that is not specified by Unicode.

4.5.4 Character Conversions

```
(char-upcase char) → char?
  char : char?
```

Produces a character consistent with the 1-to-1 code point mapping defined by Unicode. If *char* has no upcase mapping, `char-upcase` produces *char*.

String procedures, such as `string-upcase`, handle the case where Unicode defines a locale-independent mapping from the code point to a code-point sequence (in addition to the 1-1 mapping on scalar values).

Examples:

```
> (char-upcase #\a)
#\A
> (char-upcase #\λ)
#\Λ
> (char-upcase #\space)
#\space
```

```
(char-downcase char) → char?
  char : char?
```

Like `char-upcase`, but for the Unicode downcase mapping.

Examples:

```
> (char-downcase #\A)
#\a
> (char-downcase #\Λ)
#\λ
> (char-downcase #\space)
#\space
```

```
(char-titlecase char) → char?
  char : char?
```

Like `char-upcase`, but for the Unicode titlecase mapping.

Examples:

```
> (char-upcase #\a)
#\A
> (char-upcase #\λ)
#\Λ
> (char-upcase #\space)
#\space
```

(char-foldcase *char*) → char?
  char : char?

Like char-upcase, but for the Unicode case-folding mapping.

Examples:

```
> (char-foldcase #\A)
#\a
> (char-foldcase #\Σ)
#\σ
> (char-foldcase #\ς)
#\σ
> (char-foldcase #\space)
#\space
```

4.6 Symbols

§3.6 “Symbols” in ***The Racket Guide*** introduces symbols.

A *symbol* is like an immutable string, but symbols are normally interned, so that two symbols with the same character content are normally eq?. All symbols produced by the default reader (see §1.3.2 “Reading Symbols”) are interned.

The two procedures string->uninterned-symbol and gensym generate *uninterned* symbols, i.e., symbols that are not eq?, eqv?, or equal? to any other symbol, although they may print the same as other symbols.

The procedure string->unreadable-symbol returns an *unreadable symbol* that is partially interned. The default reader (see §1.3.2 “Reading Symbols”) never produces a unreadable symbol, but two calls to string->unreadable-symbol with equal? strings produce eq? results. An unreadable symbol can print the same as an interned or uninterned symbol. Unreadable symbols are useful in expansion and compilation to avoid collisions with symbols that appear in the source; they are usually not generated directly, but they can appear in the result of functions like identifier-binding.

Interned and unreadable symbols are only weakly held by the internal symbol table. This weakness can never affect the result of an eq?, eqv?, or equal? test, but a symbol may

disappear when placed into a weak box (see §16.1 “Weak Boxes”) used as the key in a weak hash table (see §4.13 “Hash Tables”), or used as an ephemeron key (see §16.2 “Ephemerons”).

See §1.3.2 “Reading Symbols” for information on reading symbols and §1.4.1 “Printing Symbols” for information on printing symbols.

```
(symbol? v) → boolean?
  v : any/c
```

Returns #t if *v* is a symbol, #f otherwise.

Examples:

```
> (symbol? 'Apple)
#t
> (symbol? 10)
#f
```

```
(symbol-interned? sym) → boolean?
  sym : symbol?
```

Returns #t if *sym* is interned, #f otherwise.

Examples:

```
> (symbol-interned? 'Apple)
#t
> (symbol-interned? (gensym))
#f
> (symbol-interned? (string->unreadable-symbol "Apple"))
#f
```

```
(symbol-unreadable? sym) → boolean?
  sym : symbol?
```

Returns #t if *sym* is an unreadable symbol, #f otherwise.

Examples:

```
> (symbol-unreadable? 'Apple)
#f
> (symbol-unreadable? (gensym))
#f
> (symbol-unreadable? (string->unreadable-symbol "Apple"))
#t
```

```
(symbol->string sym) → string?
  sym : symbol?
```

Returns a freshly allocated mutable string whose characters are the same as in *sym*.

Example:

```
> (symbol->string 'Apple)
"Apple"
```

```
(string->symbol str) → symbol?
  str : string?
```

Returns an interned symbol whose characters are the same as in *str*.

Examples:

```
> (string->symbol "Apple")
'Apple
> (string->symbol "1")
'|1|
```

```
(string->uninterned-symbol str) → symbol?
  str : string?
```

Like `(string->symbol str)`, but the resulting symbol is a new uninterned symbol. Calling `string->uninterned-symbol` twice with the same *str* returns two distinct symbols.

Examples:

```
> (string->uninterned-symbol "Apple")
'Apple
> (eq? 'a (string->uninterned-symbol "a"))
#f
> (eq? (string->uninterned-symbol "a")
       (string->uninterned-symbol "a"))
#f
```

```
(string->unreadable-symbol str) → symbol?
  str : string?
```

Like `(string->symbol str)`, but the resulting symbol is a new unreadable symbol. Calling `string->unreadable-symbol` twice with equivalent *str*s returns the same symbol, but `read` never produces the symbol.

Examples:

```
> (string->unreadable-symbol "Apple")
'Apple
> (eq? 'a (string->unreadable-symbol "a"))
#f
> (eq? (string->unreadable-symbol "a")
       (string->unreadable-symbol "a"))
#t
```

```
(gensym [base]) → symbol?
  base : (or/c string? symbol?) = "g"
```

Returns a new uninterned symbol with an automatically-generated name. The optional *base* argument is a prefix symbol or string.

Example:

```
> (gensym "apple")
'apple1065
```

```
(symbol<? a-sym b-sym ...) → boolean?
  a-sym : symbol?
  b-sym : symbol?
```

Returns `#t` if the arguments are sorted, where the comparison for each pair of symbols is the same as using `symbol->string` with `string->bytes/utf-8` and `bytes<?`.

4.7 Regular Expressions

§9 “Regular Expressions” in ***The Racket Guide*** introduces regular expressions.

Regular expressions are specified as strings or byte strings, using the same pattern language as either the Unix utility `egrep` or Perl. A string-specified pattern produces a character regexp matcher, and a byte-string pattern produces a byte regexp matcher. If a character regexp is used with a byte string or input port, it matches UTF-8 encodings (see §13.1.1 “Encodings and Locales”) of matching character streams; if a byte regexp is used with a character string, it matches bytes in the UTF-8 encoding of the string.

A regular expression that is represented as a string or byte string can be compiled to a *regexp value*, which can be used more efficiently by functions such as `regexp-match` compared to the string or byte string form. The `regexp` and `byte-regexp` procedures convert a string or byte string (respectively) into a regexp value using a syntax of regular expressions that is most compatible to `egrep`. The `pregexp` and `byte-pregexp` procedures produce a regexp value using a slightly different syntax of regular expressions that is more compatible with Perl.

Two regexp values are `equal?` if they have the same source, use the same pattern language, and are both character regexps or both byte regexps.

A literal or printed regexp value starts with `#rx` or `#px`. See §1.3.16 “Reading Regular Expressions” for information on `read`ing regular expressions and §1.4.13 “Printing Regular Expressions” for information on `print`ing regular expressions. Regexp values produced by the default reader are interned in `read-syntax` mode.

The internal size of a regexp value is limited to 32 kilobytes; this limit roughly corresponds to a source string with 32,000 literal characters or 5,000 operators.

4.7.1 Regexp Syntax

The following syntax specifications describe the content of a string that represents a regular expression. The syntax of the corresponding string may involve extra escape characters. For example, the regular expression `(.*)\1` can be represented with the string `"(.*)\\1"` or the regexp constant `#rx"(.*)\\1"`; the `\` in the regular expression must be escaped to include it in a string or regexp constant.

The `regexp` and `pregexp` syntaxes share a common core:

⟨*regexp*⟩	::=	⟨*pces*⟩	Match ⟨*pces*⟩	
	|	⟨*regexp*⟩`|`⟨*regexp*⟩	Match either ⟨*regexp*⟩, try left first	ex1
⟨*pces*⟩	::=	⟨*pce*⟩	Match ⟨*pce*⟩	
	|	⟨*pce*⟩⟨*pces*⟩	Match ⟨*pce*⟩ followed by ⟨*pces*⟩	
⟨*pce*⟩	::=	⟨*repeat*⟩	Match ⟨*repeat*⟩, longest possible	ex3
	|	⟨*repeat*⟩`?`	Match ⟨*repeat*⟩, shortest possible	ex6
	|	⟨*atom*⟩	Match ⟨*atom*⟩ exactly once	
⟨*repeat*⟩	::=	⟨*atom*⟩`*`	Match ⟨*atom*⟩ 0 or more times	ex3
	|	⟨*atom*⟩`+`	Match ⟨*atom*⟩ 1 or more times	ex4
	|	⟨*atom*⟩`?`	Match ⟨*atom*⟩ 0 or 1 times	ex5
⟨*atom*⟩	::=	`(`⟨*regexp*⟩`)`	Match sub-expression ⟨*regexp*⟩ and report	ex11
	|	`[`⟨*rng*⟩`]`	Match any character in ⟨*rng*⟩	ex2
	|	`[^`⟨*rng*⟩`]`	Match any character not in ⟨*rng*⟩	ex12
	|	`.`	Match any (except newline in multi mode)	ex13
	|	`^`	Match start (or after newline in multi mode)	ex14
	|	`$`	Match end (or before newline in multi mode)	ex15
	|	⟨*literal*⟩	Match a single literal character	ex1
	|	`(?`⟨*mode*⟩`:`⟨*regexp*⟩`)`	Match ⟨*regexp*⟩ using ⟨*mode*⟩	ex35
	|	`(?>`⟨*regexp*⟩`)`	Match ⟨*regexp*⟩, only first possible	
	|	⟨*look*⟩	Match empty if ⟨*look*⟩ matches	
	|	`(?`⟨*tst*⟩⟨*pces*⟩`|`⟨*pces*⟩`)`	Match 1st ⟨*pces*⟩ if ⟨*tst*⟩, else 2nd ⟨*pces*⟩	ex36
	|	`(?`⟨*tst*⟩⟨*pces*⟩`)`	Match ⟨*pces*⟩ if ⟨*tst*⟩, empty if not ⟨*tst*⟩	
⟨*rng*⟩	::=	`]`	⟨*rng*⟩ contains `]` only	ex27
	|	`-`	⟨*rng*⟩ contains `-` only	ex28
	|	⟨*mrng*⟩	⟨*rng*⟩ contains everything in ⟨*mrng*⟩	
	|	⟨*mrng*⟩`-`	⟨*rng*⟩ contains `-` and everything in ⟨*mrng*⟩	
⟨*mrng*⟩	::=	`]`⟨*lrng*⟩	⟨*mrng*⟩ contains `]` and everything in ⟨*lrng*⟩	ex29

	`-`⟨*lrng*⟩	⟨*mrng*⟩ contains `-` and everything in ⟨*lrng*⟩	ex29
	⟨*lirng*⟩	⟨*mrng*⟩ contains everything in ⟨*lirng*⟩	
⟨*lirng*⟩ ::=	⟨*riliteral*⟩	⟨*lirng*⟩ contains a literal character	
	⟨*riliteral*⟩`-`⟨*rliteral*⟩	⟨*lirng*⟩ contains Unicode range inclusive	ex22
	⟨*lirng*⟩⟨*lrng*⟩	⟨*lirng*⟩ contains everything in both	
⟨*lrng*⟩ ::=	`^`	⟨*lrng*⟩ contains `^`	ex30
	⟨*rliteral*⟩`-`⟨*rliteral*⟩	⟨*lrng*⟩ contains Unicode range inclusive	
	`^`⟨*lrng*⟩	⟨*lrng*⟩ contains `^` and more	
	⟨*lirng*⟩	⟨*lrng*⟩ contains everything in ⟨*lirng*⟩	
⟨*look*⟩ ::=	`(?=`⟨*regexp*⟩`)`	Match if ⟨*regexp*⟩ matches	ex31
	`(?!`⟨*regexp*⟩`)`	Match if ⟨*regexp*⟩ doesn't match	ex32
	`(?<=`⟨*regexp*⟩`)`	Match if ⟨*regexp*⟩ matches preceding	ex33
	`(?<!`⟨*regexp*⟩`)`	Match if ⟨*regexp*⟩ doesn't match preceding	ex34
⟨*tst*⟩ ::=	`(`⟨*n*⟩`)`	True if Nth `(` has a match	
	⟨*look*⟩	True if ⟨*look*⟩ matches	ex36
⟨*mode*⟩ ::=		Like the enclosing mode	
	⟨*mode*⟩`i`	Like ⟨*mode*⟩, but case-insensitive	ex35
	⟨*mode*⟩`-i`	Like ⟨*mode*⟩, but sensitive	
	⟨*mode*⟩`s`	Like ⟨*mode*⟩, but not in multi mode	
	⟨*mode*⟩`-s`	Like ⟨*mode*⟩, but in multi mode	
	⟨*mode*⟩`m`	Like ⟨*mode*⟩, but in multi mode	
	⟨*mode*⟩`-m`	Like ⟨*mode*⟩, but not in multi mode	

The following completes the grammar for `regexp`, which treats `{` and `}` as literals, `\` as a literal within ranges, and `\` as a literal producer outside of ranges.

⟨*literal*⟩ ::=	Any character except `(`, `)`, `*`, `+`, `?`, `[`, `.`, `^`, `\`, or `|`		
	`\`⟨*aliteral*⟩	Match ⟨*aliteral*⟩	ex21
⟨*aliteral*⟩ ::=	Any character		
⟨*riliteral*⟩ ::=	Any character except `]`, `-`, or `^`		
⟨*rliteral*⟩ ::=	Any character except `]` or `-`		

The following completes the grammar for `pregexp`, which uses `{` and `}` bounded repetition and uses `\` for meta-characters both inside and outside of ranges.

⟨*repeat*⟩ ::=	...	...	
	⟨*atom*⟩`{`⟨*n*⟩`}`	Match ⟨*atom*⟩ exactly ⟨*n*⟩ times	ex7
	⟨*atom*⟩`{`⟨*n*⟩`,}`	Match ⟨*atom*⟩ ⟨*n*⟩ or more times	ex8
	⟨*atom*⟩`{,`⟨*m*⟩`}`	Match ⟨*atom*⟩ between 0 and ⟨*m*⟩ times	ex9
	⟨*atom*⟩`{`⟨*n*⟩`,`⟨*m*⟩`}`	Match ⟨*atom*⟩ between ⟨*n*⟩ and ⟨*m*⟩ times	ex10
⟨*atom*⟩ ::=	...	...	
	`\`⟨*n*⟩	Match latest reported match for ⟨*n*⟩th `(`	ex16
	⟨*class*⟩	Match any character in ⟨*class*⟩	
	`\b`	Match `\w*` boundary	ex17
	`\B`	Match where `\b` does not	ex18
	`\p{`⟨*property*⟩`}`	Match (UTF-8 encoded) in ⟨*property*⟩	ex19

	|	`\P{`⟨*property*⟩`}`	Match (UTF-8 encoded) not in ⟨*property*⟩	ex20
⟨*literal*⟩	::=	Any character except `(`, `)`, `*`, `+`, `?`, `[`, `]`, `{`, `}`, `.`, `^`, `\`, or `|`	
	|	`\`⟨*aliteral*⟩	Match ⟨*aliteral*⟩	ex21
⟨*aliteral*⟩	::=	Any character except `a-z`, `A-Z`, `0-9`	
⟨*lirng*⟩	::=	...	...	
	|	⟨*class*⟩	⟨*lirng*⟩ contains all characters in ⟨*class*⟩	
	|	⟨*posix*⟩	⟨*lirng*⟩ contains all characters in ⟨*posix*⟩	ex26
	|	`\`⟨*eliteral*⟩	⟨*lirng*⟩ contains ⟨*eliteral*⟩	
⟨*riliteral*⟩	::=	Any character except `]`, `\`, `-`, or `^`	
⟨*rliteral*⟩	::=	Any character except `]`, `\`, or `-`	
⟨*eliteral*⟩	::=	Any character except `a-z`, `A-Z`	
⟨*class*⟩	::=	`\d`	Contains `0-9`	ex23
	|	`\D`	Contains ASCII other than those in `\d`	
	|	`\w`	Contains `a-z`, `A-Z`, `0-9`, `_`	ex24
	|	`\W`	Contains ASCII other than those in `\w`	
	|	`\s`	Contains space, tab, newline, formfeed, return	ex25
	|	`\S`	Contains ASCII other than those in `\s`	
⟨*posix*⟩	::=	`[:alpha:]`	Contains `a-z`, `A-Z`	
	|	`[:upper:]`	Contains `A-Z`	
	|	`[:lower:]`	Contains `a-z`	ex26
	|	`[:digit:]`	Contains `0-9`	
	|	`[:xdigit:]`	Contains `0-9`, `a-f`, `A-F`	
	|	`[:alnum:]`	Contains `a-z`, `A-Z`, `0-9`	
	|	`[:word:]`	Contains `a-z`, `A-Z`, `0-9`, `_`	
	|	`[:blank:]`	Contains space and tab	
	|	`[:space:]`	Contains space, tab, newline, formfeed, return	
	|	`[:graph:]`	Contains all ASCII characters that use ink	
	|	`[:print:]`	Contains space, tab, and ASCII ink users	
	|	`[:cntrl:]`	Contains all characters with scalar value < 32	
	|	`[:ascii:]`	Contains all ASCII characters	
⟨*property*⟩	::=	⟨*category*⟩	Includes all characters in ⟨*category*⟩	
	|	`^`⟨*category*⟩	Includes all characters not in ⟨*category*⟩	

The Unicode categories follow.

⟨*category*⟩	::=	`Ll`	Letter, lowercase	ex19
	|	`Lu`	Letter, uppercase	
	|	`Lt`	Letter, titlecase	
	|	`Lm`	Letter, modifier	
	|	`L&`	Union of `Ll`, `Lu`, `Lt`, and `Lm`	
	|	`Lo`	Letter, other	
	|	`L`	Union of `L&` and `Lo`	
	|	`Nd`	Number, decimal digit	
	|	`Nl`	Number, letter	
	|	`No`	Number, other	
	|	`N`	Union of `Nd`, `Nl`, and `No`	

`Ps`	Punctuation, open
`Pe`	Punctuation, close
`Pi`	Punctuation, initial quote
`Pf`	Punctuation, final quote
`Pc`	Punctuation, connector
`Pd`	Punctuation, dash
`Po`	Punctuation, other
`P`	Union of `Ps`, `Pe`, `Pi`, `Pf`, `Pc`, `Pd`, and `Po`
`Mn`	Mark, non-spacing
`Mc`	Mark, spacing combining
`Me`	Mark, enclosing
`M`	Union of `Mn`, `Mc`, and `Me`
`Sc`	Symbol, currency
`Sk`	Symbol, modifier
`Sm`	Symbol, math
`So`	Symbol, other
`S`	Union of `Sc`, `Sk`, `Sm`, and `So`
`Zl`	Separator, line
`Zp`	Separator, paragraph
`Zs`	Separator, space
`Z`	Union of `Zl`, `Zp`, and `Zs`
`Cc`	Other, control
`Cf`	Other, format
`Cs`	Other, surrogate
`Cn`	Other, not assigned
`Co`	Other, private use
`C`	Union of `Cc`, `Cf`, `Cs`, `Cn`, and `Co`
`.`	Union of all Unicode categories

Examples:

```
> (regexp-match #rx"a|b" "cat") ; ex1
'("a")
> (regexp-match #rx"[at]" "cat") ; ex2
'("a")
> (regexp-match #rx"ca*[at]" "caaat") ; ex3
'("caaat")
> (regexp-match #rx"ca+[at]" "caaat") ; ex4
'("caaat")
> (regexp-match #rx"ca?t?" "ct") ; ex5
'("ct")
> (regexp-match #rx"ca*?[at]" "caaat") ; ex6
'("ca")
> (regexp-match #px"ca{2}" "caaat") ; ex7, uses #px
'("caa")
> (regexp-match #px"ca{2,}t" "catcaat") ; ex8, uses #px
```

```
'("caat")
> (regexp-match #px"ca{,2}t" "caaatcat") ; ex9, uses #px
'("cat")
> (regexp-match #px"ca{1,2}t" "caaatcat") ; ex10, uses #px
'("cat")
> (regexp-match #rx"(c*)(a*)" "caat") ; ex11
'("caa" "c" "aa")
> (regexp-match #rx"[^ca]" "caat") ; ex12
'("t")
> (regexp-match #rx".(.)." "cat") ; ex13
'("cat" "a")
> (regexp-match #rx"^a|^c" "cat") ; ex14
'("c")
> (regexp-match #rx"a$|t$" "cat") ; ex15
'("t")
> (regexp-match #px"c(.)\\1t" "caat") ; ex16, uses #px
'("caat" "a")
> (regexp-match #px".\\b." "cat in hat") ; ex17, uses #px
'("t ")
> (regexp-match #px".\\B." "cat in hat") ; ex18, uses #px
'("ca")
> (regexp-match #px"\\p{Ll}" "Cat") ; ex19, uses #px
'("a")
> (regexp-match #px"\\P{Ll}" "cat!") ; ex20, uses #px
'("!")
> (regexp-match #rx"\\|" "c|t") ; ex21
'("|")
> (regexp-match #rx"[a-f]*" "cat") ; ex22
'("ca")
> (regexp-match #px"[a-f\\d]*" "1cat") ; ex23, uses #px
'("1ca")
> (regexp-match #px" [\\w]" "cat hat") ; ex24, uses #px
'(" h")
> (regexp-match #px"t[\\s]" "cat\nhat") ; ex25, uses #px
'("t\n")
> (regexp-match #px"[[:lower:]]+" "Cat") ; ex26, uses #px
'("at")
> (regexp-match #rx"[]]" "c]t") ; ex27
'("]")
> (regexp-match #rx"[-]" "c-t") ; ex28
'("-")
> (regexp-match #rx"[]a[]+" "c[a]t") ; ex29
'("[a]")
> (regexp-match #rx"[a^]+" "ca^t") ; ex30
'("a^")
> (regexp-match #rx".a(?=p)" "cat nap") ; ex31
```

```
'("na")
> (regexp-match #rx".a(?!t)" "cat nap") ; ex32
'("na")
> (regexp-match #rx"(?<=n)a." "cat nap") ; ex33
'("ap")
> (regexp-match #rx"(?<!c)a." "cat nap") ; ex34
'("ap")
> (regexp-match #rx"(?i:a)[tp]" "cAT nAp") ; ex35
'("Ap")
> (regexp-match #rx"(?(?<=c)a|b)+" "cabal") ; ex36
'("ab")
```

4.7.2 Additional Syntactic Constraints

In addition to matching a grammar, regular expressions must meet two syntactic restrictions:

- In a ⟨*repeat*⟩ other than ⟨*atom*⟩`?`, the ⟨*atom*⟩ must not match an empty sequence.
- In a `(?<=`⟨*regexp*⟩`)` or `(?<!`⟨*regexp*⟩`)`, the ⟨*regexp*⟩ must match a bounded sequence only.

These constraints are checked syntactically by the following type system. A type $[n, m]$ corresponds to an expression that matches between n and m characters. In the rule for `(`⟨*Regexp*⟩`)`, N means the number such that the opening parenthesis is the Nth opening parenthesis for collecting match reports. Non-emptiness is inferred for a backreference pattern, `\`⟨N⟩, so that a backreference can be used for repetition patterns; in the case of mutual dependencies among backreferences, the inference chooses the fixpoint that maximizes non-emptiness. Finiteness is not inferred for backreferences (i.e., a backreference is assumed to match an arbitrarily large sequence).

$$\frac{\langle regexp\rangle_1 \ : \ [n_1, m_1] \qquad \langle regexp\rangle_2 \ : \ [n_2, m_2]}{\langle regexp\rangle_1\texttt{|}\langle regexp\rangle_2 \ : \ [\min(n_1, n_2), \max(m_1, m_2)]}$$

$$\frac{\langle pce\rangle \ : \ [n_1, m_1] \qquad \langle pces\rangle \ : \ [n_2, m_2]}{\langle pce\rangle\langle pces\rangle \ : \ [n_1{+}n_2, m_1{+}m_2]}$$

$$\frac{\langle repeat\rangle \ : \ [n, m]}{\langle repeat\rangle\texttt{?} \ : \ [0, m]} \qquad \frac{\langle atom\rangle \ : \ [n, m] \qquad n > 0}{\langle atom\rangle\texttt{*} \ : \ [0, \infty]}$$

$$\frac{\langle atom\rangle \ : \ [n, m] \qquad n > 0}{\langle atom\rangle\texttt{+} \ : \ [1, \infty]} \qquad \frac{\langle atom\rangle \ : \ [n, m]}{\langle atom\rangle\texttt{?} \ : \ [0, m]}$$

$$\frac{\langle atom\rangle \ : \ [n, m] \qquad n > 0}{\langle atom\rangle\texttt{\{}\langle n\rangle\texttt{\}} \ : \ [n{*}\langle n\rangle, m{*}\langle n\rangle]}$$

$$\frac{\langle atom\rangle \;:\; [n,m] \qquad n>0}{\langle atom\rangle\texttt{\{}\langle n\rangle\texttt{,\}} \;:\; [n{*}\langle n\rangle,\infty]}$$

$$\frac{\langle atom\rangle \;:\; [n,m] \qquad n>0}{\langle atom\rangle\texttt{\{,}\langle m\rangle\texttt{\}} \;:\; [0,m{*}\langle m\rangle]}$$

$$\frac{\langle atom\rangle \;:\; [n,m] \qquad n>0}{\langle atom\rangle\texttt{\{}\langle n\rangle\texttt{,}\langle m\rangle\texttt{\}} \;:\; [n{*}\langle n\rangle,m{*}\langle m\rangle]}$$

$$\frac{\langle regexp\rangle \;:\; [n,m]}{\texttt{(}\langle regexp\rangle\texttt{)} \;:\; [n,m] \qquad \alpha_{\mathrm{N}}{=}n}$$

$$\frac{\langle regexp\rangle \;:\; [n,m]}{\texttt{(?}\langle mode\rangle\texttt{:}\langle regexp\rangle\texttt{)} \;:\; [n,m]}$$

$$\frac{\langle regexp\rangle \;:\; [n,m]}{\texttt{(?=}\langle regexp\rangle\texttt{)} \;:\; [0,0]} \qquad \frac{\langle regexp\rangle \;:\; [n,m]}{\texttt{(?!}\langle regexp\rangle\texttt{)} \;:\; [0,0]}$$

$$\frac{\langle regexp\rangle \;:\; [n,m] \qquad m<\infty}{\texttt{(?<=}\langle regexp\rangle\texttt{)} \;:\; [0,0]} \qquad \frac{\langle regexp\rangle \;:\; [n,m] \qquad m<\infty}{\texttt{(?<!}\langle regexp\rangle\texttt{)} \;:\; [0,0]}$$

$$\frac{\langle regexp\rangle \;:\; [n,m]}{\texttt{(?>}\langle regexp\rangle\texttt{)} \;:\; [n,m]}$$

$$\frac{\langle tst\rangle \;:\; [n_0,m_0] \qquad \langle pces\rangle_1 \;:\; [n_1,m_1] \qquad \langle pces\rangle_2 \;:\; [n_2,m_2]}{\texttt{(?}\langle tst\rangle\langle pces\rangle_1\texttt{|}\langle pces\rangle_2\texttt{)} \;:\; [\min(n_1,n_2),\max(m_1,m_2)]}$$

$$\frac{\langle tst\rangle \;:\; [n_0,m_0] \qquad \langle pces\rangle \;:\; [n_1,m_1]}{\texttt{(?}\langle tst\rangle\langle pces\rangle\texttt{)} \;:\; [0,m_1]}$$

$$\texttt{(}\langle n\rangle\texttt{)} \;:\; [\alpha_{\mathrm{N}},\infty] \qquad \texttt{[}\langle rng\rangle\texttt{]} \;:\; [1,1] \qquad \texttt{[\^{}}\langle rng\rangle\texttt{]} \;:\; [1,1]$$

$$\texttt{.} \;:\; [1,1] \qquad \texttt{\^{}} \;:\; [0,0] \qquad \texttt{\$} \;:\; [0,0]$$

$$\langle literal\rangle \;:\; [1,1] \qquad \texttt{\textbackslash}\langle n\rangle \;:\; [\alpha_{\mathrm{N}},\infty] \qquad \langle class\rangle \;:\; [1,1]$$

$$\texttt{\textbackslash b} \;:\; [0,0] \qquad \texttt{\textbackslash B} \;:\; [0,0]$$

$$\texttt{\textbackslash p\{}\langle property\rangle\texttt{\}} \;:\; [1,6] \qquad \texttt{\textbackslash P\{}\langle property\rangle\texttt{\}} \;:\; [1,6]$$

4.7.3 Regexp Constructors

```
(regexp? v) → boolean?
```

```
  v : any/c
```

Returns `#t` if *v* is a regexp value created by `regexp` or `pregexp`, `#f` otherwise.

```
(pregexp? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a regexp value created by `pregexp` (not `regexp`), `#f` otherwise.

```
(byte-regexp? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a regexp value created by `byte-regexp` or `byte-pregexp`, `#f` otherwise.

```
(byte-pregexp? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a regexp value created by `byte-pregexp` (not `byte-regexp`), `#f` otherwise.

```
(regexp str) → regexp?
  str : string?
```

Takes a string representation of a regular expression (using the syntax in §4.7.1 "Regexp Syntax") and compiles it into a regexp value. Other regular expression procedures accept either a string or a regexp value as the matching pattern. If a regular expression string is used multiple times, it is faster to compile the string once to a regexp value and use it for repeated matches instead of using the string each time.

The `object-name` procedure returns the source string for a regexp value.

Examples:

```
> (regexp "ap*le")
#rx"ap*le"
> (object-name #rx"ap*le")
"ap*le"
```

```
(pregexp string) → pregexp?
  string : string?
```

Like `regexp`, except that it uses a slightly different syntax (see §4.7.1 "Regexp Syntax"). The result can be used with `regexp-match`, etc., just like the result from `regexp`.

Examples:

```
> (pregexp "ap*le")
#px"ap*le"
> (regexp? #px"ap*le")
#t
```

```
(byte-regexp bstr) → byte-regexp?
  bstr : bytes?
```

Takes a byte-string representation of a regular expression (using the syntax in §4.7.1 “Regexp Syntax”) and compiles it into a byte-regexp value.

The `object-name` procedure returns the source byte string for a regexp value.

Examples:

```
> (byte-regexp #"ap*le")
#rx#"ap*le"
> (object-name #rx#"ap*le")
#"ap*le"
> (byte-regexp "ap*le")
byte-regexp: contract violation
  expected: byte?
  given: "ap*le"
```

```
(byte-pregexp bstr) → byte-pregexp?
  bstr : bytes?
```

Like `byte-regexp`, except that it uses a slightly different syntax (see §4.7.1 “Regexp Syntax”). The result can be used with `regexp-match`, etc., just like the result from `byte-regexp`.

Example:

```
> (byte-pregexp #"ap*le")
#px#"ap*le"
```

```
(regexp-quote str [case-sensitive?]) → string?
  str : string?
  case-sensitive? : any/c = #t
(regexp-quote bstr [case-sensitive?]) → bytes?
  bstr : bytes?
  case-sensitive? : any/c = #t
```

Produces a string or byte string suitable for use with `regexp` to match the literal sequence of characters in *str* or sequence of bytes in *bstr*. If *case-sensitive?* is true (the default), the resulting regexp matches letters in *str* or `bytes` case-sensitively, otherwise it matches case-insensitively.

Examples:

```
> (regexp-match "." "apple.scm")
'("a")
> (regexp-match (regexp-quote ".") "apple.scm")
'(".")
```

```
(regexp-max-lookbehind pattern) → exact-nonnegative-integer?
  pattern : (or/c regexp? byte-regexp?)
```

Returns the maximum number of bytes that *pattern* may consult before the starting position of a match to determine the match. For example, the pattern `(?<=abc)d` consults three bytes preceding a matching `d`, while `e(?<=a..)d` consults two bytes before a matching `ed`. A `^` pattern may consult a preceding byte to determine whether the current position is the start of the input or of a line.

4.7.4 Regexp Matching

```
(regexp-match  pattern
               input
              [start-pos
               end-pos
               output-port
               input-prefix])
   (if (and (or (string? pattern) (regexp? pattern))
            (or (string? input) (path? input)))
→      (or/c #f (cons/c string? (listof (or/c string? #f))))
       (or/c #f (cons/c bytes?  (listof (or/c bytes?  #f)))))
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input : (or/c string? bytes? path? input-port?)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : (or/c exact-nonnegative-integer? #f) = #f
  output-port : (or/c output-port? #f) = #f
  input-prefix : bytes? = #""
```

Attempts to match *pattern* (a string, byte string, regexp value, or byte-regexp value) once to a portion of *input*. The matcher finds a portion of *input* that matches and is closest to the start of the input (after *start-pos*).

If *input* is a path, it is converted to a byte string with `path->bytes` if *pattern* is a byte string or a byte-based regexp. Otherwise, *input* is converted to a string with `path->string`.

The optional *start-pos* and *end-pos* arguments select a portion of *input* for matching; the default is the entire string or the stream up to an end-of-file. When *input* is a string,

start-pos is a character position; when *input* is a byte string, then *start-pos* is a byte position; and when *input* is an input port, *start-pos* is the number of bytes to skip before starting to match. The *end-pos* argument can be `#f`, which corresponds to the end of the string or an end-of-file in the stream; otherwise, it is a character or byte position, like *start-pos*. If *input* is an input port, and if an end-of-file is reached before *start-pos* bytes are skipped, then the match fails.

In *pattern*, a start-of-string `^` refers to the first position of *input* after *start-pos*, assuming that *input-prefix* is `#""`. The end-of-input `$` refers to the *end-pos*th position or (in the case of an input port) an end-of-file, whichever comes first.

The *input-prefix* specifies bytes that effectively precede *input* for the purposes of `^` and other look-behind matching. For example, a `#""` prefix means that `^` matches at the beginning of the stream, while a `#"\n"` *input-prefix* means that a start-of-line `^` can match the beginning of the input, while a start-of-file `^` cannot.

If the match fails, `#f` is returned. If the match succeeds, a list containing strings or byte string, and possibly `#f`, is returned. The list contains strings only if *input* is a string and *pattern* is not a byte regexp. Otherwise, the list contains byte strings (substrings of the UTF-8 encoding of *input*, if *input* is a string).

The first [byte] string in a result list is the portion of *input* that matched *pattern*. If two portions of *input* can match *pattern*, then the match that starts earliest is found.

Additional [byte] strings are returned in the list if *pattern* contains parenthesized sub-expressions (but not when the opening parenthesis is followed by `?`). Matches for the sub-expressions are provided in the order of the opening parentheses in *pattern*. When sub-expressions occur in branches of an `|` "or" pattern, in a `*` "zero or more" pattern, or other places where the overall pattern can succeed without a match for the sub-expression, then a `#f` is returned for the sub-expression if it did not contribute to the final match. When a single sub-expression occurs within a `*` "zero or more" pattern or other multiple-match positions, then the rightmost match associated with the sub-expression is returned in the list.

If the optional *output-port* is provided as an output port, the part of *input* from its beginning (not *start-pos*) that precedes the match is written to the port. All of *input* up to *end-pos* is written to the port if no match is found. This functionality is most useful when *input* is an input port.

When matching an input port, a match failure reads up to *end-pos* bytes (or end-of-file), even if *pattern* begins with a start-of-string `^`; see also `regexp-try-match`. On success, all bytes up to and including the match are eventually read from the port, but matching proceeds by first peeking bytes from the port (using `peek-bytes-avail!`), and then (re-)reading matching bytes to discard them after the match result is determined. Non-matching bytes may be read and discarded before the match is determined. The matcher peeks in blocking mode only as far as necessary to determine a match, but it may peek extra bytes to fill an internal buffer if immediately available (i.e., without blocking). Greedy repeat operators in *pattern*, such as `*` or `+`, tend to force reading the entire content of the port (up

to *end-pos*) to determine a match.

If the input port is read simultaneously by another thread, or if the port is a custom port with inconsistent reading and peeking procedures (see §13.1.9 “Custom Ports”), then the bytes that are peeked and used for matching may be different than the bytes read and discarded after the match completes; the matcher inspects only the peeked bytes. To avoid such interleaving, use `regexp-match-peek` (with a `progress-evt` argument) followed by `port-commit-peeked`.

Examples:

```
> (regexp-match #rx"x." "12x4x6")
'("x4")
> (regexp-match #rx"y." "12x4x6")
#f
> (regexp-match #rx"x." "12x4x6" 3)
'("x6")
> (regexp-match #rx"x." "12x4x6" 3 4)
#f
> (regexp-match #rx#"x." "12x4x6")
'(#"x4")
> (regexp-match #rx"x." "12x4x6" 0 #f (current-output-port))
12
'("x4")
> (regexp-match #rx"(-[0-9]*)+" "a-12--345b")
'("-12--345" "-345")
```

```
(regexp-match* pattern
               input
              [start-pos
               end-pos
               input-prefix
               #:match-select match-select
               #:gap-select? gap-select])
  (if (and (or (string? pattern) (regexp? pattern))
           (or (string? input) (path? input)))
→      (listof (or/c string? (listof (or/c #f string?))))
       (listof (or/c bytes? (listof (or/c #f bytes?)))))
 pattern : (or/c string? bytes? regexp? byte-regexp?)
 input : (or/c string? bytes? path? input-port?)
 start-pos : exact-nonnegative-integer? = 0
 end-pos : (or/c exact-nonnegative-integer? #f) = #f
 input-prefix : bytes? = #""
 match-select : (or/c (list? . -> . (or/c any/c list?))
                      #f)                                  = car
 gap-select : any/c = #f
```

Like regexp-match, but the result is a list of strings or byte strings corresponding to a sequence of matches of *pattern* in *input*.

The *pattern* is used in order to find matches, where each match attempt starts at the end of the last match, and ^ is allowed to match the beginning of the input (if *input-prefix* is #"") only for the first match. Empty matches are handled like other matches, returning a zero-length string or byte sequence (they are more useful in making this a complement of regexp-split), but *pattern* is restricted from matching an empty sequence immediately after an empty match.

If *input* contains no matches (in the range *start-pos* to *end-pos*), null is returned. Otherwise, each item in the resulting list is a distinct substring or byte sequence from *input* that matches *pattern*. The *end-pos* argument can be #f to match to the end of *input* (which corresponds to an end-of-file if *input* is an input port).

Examples:

```
> (regexp-match* #rx"x." "12x4x6")
'("x4" "x6")
> (regexp-match* #rx"x*" "12x4x6")
'("" "" "x" "" "x" "" "")
```

match-select specifies the collected results. The default of car means that the result is the list of matches without returning parenthesized sub-patterns. It can be given as a 'selector' function which chooses an item from a list, or it can choose a list of items. For example, you can use cdr to get a list of lists of parenthesized sub-patterns matches, or values (as an identity function) to get the full matches as well. (Note that the selector must choose an element of its input list or a list of elements, but it must not inspect its input as they can be either a list of strings or a list of position pairs. Furthermore, the selector must be consistent in its choice(s).)

Examples:

```
> (regexp-match* #rx"x(.)" "12x4x6" #:match-select cadr)
'("4" "6")
> (regexp-match* #rx"x(.)" "12x4x6" #:match-select values)
'(("x4" "4") ("x6" "6"))
```

In addition, specifying *gap-select* as a non-#f value will make the result an interleaved list of the matches as well as the separators between them matches, starting and ending with a separator. In this case, *match-select* can be given as #f to return *only* the separators, making such uses equivalent to regexp-split.

Examples:

```
> (regexp-match* #rx"x(.)" "12x4x6" #:match-select cadr #:gap-
select? #t)
'("12" "4" "" "6" "")
```

```
> (regexp-match* #rx"x(.)" "12x4x6" #:match-select #f #:gap-
select? #t)
'("12" "" "")
```

```
(regexp-try-match  pattern
                   input
                  [start-pos
                   end-pos
                   output-port
                   input-prefix])
   (if (and (or (string? pattern) (regexp? pattern))
            (string? input))
→
       (or/c #f (cons/c string? (listof (or/c string? #f))))
       (or/c #f (cons/c bytes?  (listof (or/c bytes?  #f)))))
 pattern : (or/c string? bytes? regexp? byte-regexp?)
 input : input-port?
 start-pos : exact-nonnegative-integer? = 0
 end-pos : (or/c exact-nonnegative-integer? #f) = #f
 output-port : (or/c output-port? #f) = #f
 input-prefix : bytes? = #""
```

Like `regexp-match` on input ports, except that if the match fails, no characters are read and discarded from `in`.

This procedure is especially useful with a *pattern* that begins with a start-of-string `^` or with a non-`#f` *end-pos*, since each limits the amount of peeking into the port. Otherwise, beware that a large portion of the stream may be peeked (and therefore pulled into memory) before the match succeeds or fails.

```
(regexp-match-positions  pattern
                         input
                        [start-pos
                         end-pos
                         output-port
                         input-prefix])
   (or/c (cons/c (cons/c exact-nonnegative-integer?
                         exact-nonnegative-integer?)
                 (listof (or/c (cons/c exact-nonnegative-integer?
→                                       exact-nonnegative-integer?)
                               #f)))
         #f)
 pattern : (or/c string? bytes? regexp? byte-regexp?)
 input : (or/c string? bytes? path? input-port?)
 start-pos : exact-nonnegative-integer? = 0
 end-pos : (or/c exact-nonnegative-integer? #f) = #f
 output-port : (or/c output-port? #f) = #f
 input-prefix : bytes? = #""
```

Like regexp-match, but returns a list of number pairs (and #f) instead of a list of strings. Each pair of numbers refers to a range of characters or bytes in *input*. If the result for the same arguments with regexp-match would be a list of byte strings, the resulting ranges correspond to byte ranges; in that case, if *input* is a character string, the byte ranges correspond to bytes in the UTF-8 encoding of the string.

Range results are returned in a substring- and subbytes-compatible manner, independent of *start-pos*. In the case of an input port, the returned positions indicate the number of bytes that were read, including *start-pos*, before the first matching byte.

Examples:

```
> (regexp-match-positions #rx"x." "12x4x6")
'((2 . 4))
> (regexp-match-positions #rx"x." "12x4x6" 3)
'((4 . 6))
> (regexp-match-positions #rx"(-[0-9]*)+" "a-12--345b")
'((1 . 9) (5 . 9))
```

```
(regexp-match-positions* pattern
                         input
                        [start-pos
                         end-pos
                         input-prefix
                         #:match-select match-select])
   (or/c (listof (cons/c exact-nonnegative-integer?
                         exact-nonnegative-integer?))
→        (listof (listof (or/c #f (cons/c exact-nonnegative-integer?
                                          exact-nonnegative-integer?)))))
 pattern : (or/c string? bytes? regexp? byte-regexp?)
 input : (or/c string? bytes? path? input-port?)
 start-pos : exact-nonnegative-integer? = 0
 end-pos : (or/c exact-nonnegative-integer? #f) = #f
 input-prefix : bytes? = #""
 match-select : (list? . -> . (or/c any/c list?)) = car
```

Like regexp-match-positions, but returns multiple matches like regexp-match*.

Examples:

```
> (regexp-match-positions* #rx"x." "12x4x6")
'((2 . 4) (4 . 6))
> (regexp-match-positions* #rx"x(.)" "12x4x6" #:match-select cadr)
'((3 . 4) (5 . 6))
```

Note that unlike regexp-match*, there is no #:gap-select? input keyword, as this information can be easily inferred from the resulting matches.

```
(regexp-match?  pattern
                input
               [start-pos
                end-pos
                output-port
                input-prefix]) → boolean?
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input : (or/c string? bytes? path? input-port?)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : (or/c exact-nonnegative-integer? #f) = #f
  output-port : (or/c output-port? #f) = #f
  input-prefix : bytes? = #""
```

Like `regexp-match`, but returns merely `#t` when the match succeeds, `#f` otherwise.

Examples:

```
> (regexp-match? #rx"x." "12x4x6")
#t
> (regexp-match? #rx"y." "12x4x6")
#f
```

```
(regexp-match-exact? pattern input) → boolean?
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input : (or/c string? bytes? path?)
```

Like `regexp-match?`, but `#t` is only returned when the entire content of *input* matches *pattern*.

Examples:

```
> (regexp-match-exact? #rx"x." "12x4x6")
#f
> (regexp-match-exact? #rx"1.*x." "12x4x6")
#t
```

```
(regexp-match-peek  pattern
                    input
                   [start-pos
                    end-pos
                    progress
                    input-prefix])
  → (or/c (cons/c bytes? (listof (or/c bytes? #f)))
          #f)
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input : input-port?
```

```
  start-pos : exact-nonnegative-integer? = 0
  end-pos : (or/c exact-nonnegative-integer? #f) = #f
  progress : (or/c evt #f) = #f
  input-prefix : bytes? = #""
```

Like `regexp-match` on input ports, but only peeks bytes from *input* instead of reading them. Furthermore, instead of an output port, the last optional argument is a progress event for *input* (see `port-progress-evt`). If *progress* becomes ready, then the match stops peeking from *input* and returns `#f`. The *progress* argument can be `#f`, in which case the peek may continue with inconsistent information if another process meanwhile reads from *input*.

Examples:

```
> (define p (open-input-string "a abcd"))

> (regexp-match-peek ".*bc" p)
'(#"a abc")
> (regexp-match-peek ".*bc" p 2)
'(#"abc")
> (regexp-match ".*bc" p 2)
'(#"abc")
> (peek-char p)
#\d
> (regexp-match ".*bc" p)
#f
> (peek-char p)
#<eof>
```

```
(regexp-match-peek-positions pattern
                             input
                            [start-pos
                             end-pos
                             progress
                             input-prefix])
   (or/c (cons/c (cons/c exact-nonnegative-integer?
                         exact-nonnegative-integer?)
                 (listof (or/c (cons/c exact-nonnegative-integer?
→                                       exact-nonnegative-integer?)
                               #f)))
         #f)
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input : input-port?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : (or/c exact-nonnegative-integer? #f) = #f
  progress : (or/c evt #f) = #f
  input-prefix : bytes? = #""
```

Like regexp-match-positions on input ports, but only peeks bytes from *input* instead of reading them, and with a *progress* argument like regexp-match-peek.

```
(regexp-match-peek-immediate  pattern
                              input
                             [start-pos
                              end-pos
                              progress
                              input-prefix])
  (or/c (cons/c bytes? (listof (or/c bytes? #f)))
→       #f)
 pattern : (or/c string? bytes? regexp? byte-regexp?)
 input : input-port?
 start-pos : exact-nonnegative-integer? = 0
 end-pos : (or/c exact-nonnegative-integer? #f) = #f
 progress : (or/c evt #f) = #f
 input-prefix : bytes? = #""
```

Like regexp-match-peek, but it attempts to match only bytes that are available from *input* without blocking. The match fails if not-yet-available characters might be used to match *pattern*.

```
(regexp-match-peek-positions-immediate  pattern
                                        input
                                       [start-pos
                                        end-pos
                                        progress
                                        input-prefix])
  (or/c (cons/c (cons/c exact-nonnegative-integer?
                        exact-nonnegative-integer?)
                (listof (or/c (cons/c exact-nonnegative-integer?
→                                     exact-nonnegative-integer?)
                              #f)))
        #f)
 pattern : (or/c string? bytes? regexp? byte-regexp?)
 input : input-port?
 start-pos : exact-nonnegative-integer? = 0
 end-pos : (or/c exact-nonnegative-integer? #f) = #f
 progress : (or/c evt #f) = #f
 input-prefix : bytes? = #""
```

Like regexp-match-peek-positions, but it attempts to match only bytes that are available from *input* without blocking. The match fails if not-yet-available characters might be used to match *pattern*.

```
(regexp-match-peek-positions*  pattern
                               input
                              [start-pos
                               end-pos
                               input-prefix
                               #:match-select match-select])
   (or/c (listof (cons/c exact-nonnegative-integer?
                         exact-nonnegative-integer?))
→        (listof (listof (or/c #f (cons/c exact-nonnegative-integer?
                                          exact-nonnegative-integer?)))))
 pattern : (or/c string? bytes? regexp? byte-regexp?)
 input : input-port?
 start-pos : exact-nonnegative-integer? = 0
 end-pos : (or/c exact-nonnegative-integer? #f) = #f
 input-prefix : bytes? = #""
 match-select : (list? . -> . (or/c any/c list?)) = car
```

Like `regexp-match-peek-positions`, but returns multiple matches like `regexp-match-positions*`.

```
(regexp-match/end  pattern
                   input
                  [start-pos
                   end-pos
                   output-port
                   input-prefix
                   count])
   (if (and (or (string? pattern) (regexp? pattern))
            (or/c (string? input) (path? input)))
→     (or/c #f (cons/c string? (listof (or/c string? #f))))
       (or/c #f (cons/c bytes?  (listof (or/c bytes?  #f)))))
   (or/c #f bytes?)
 pattern : (or/c string? bytes? regexp? byte-regexp?)
 input : (or/c string? bytes? path? input-port?)
 start-pos : exact-nonnegative-integer? = 0
 end-pos : (or/c exact-nonnegative-integer? #f) = #f
 output-port : (or/c output-port? #f) = #f
 input-prefix : bytes? = #""
 count : nonnegative-exact-integer? = 1
```

Like `regexp-match`, but with a second result: a byte string of up to *count* bytes that correspond to the input (possibly including the *input-prefix*) leading to the end of the match; the second result is `#f` if no match is found.

The second result can be useful as an *input-prefix* for attempting a second match on *input* starting from the end of the first match. In that case, use `regexp-max-lookbehind`

to determine an appropriate value for *count*.

```
(regexp-match-positions/end  pattern
                             input
                            [start-pos
                             end-pos
                             input-prefix
                             count])
   (listof (cons/c exact-nonnegative-integer?
 →                 exact-nonnegative-integer?))
   (or/c #f bytes?)
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input : (or/c string? bytes? path? input-port?)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : (or/c exact-nonnegative-integer? #f) = #f
  input-prefix : bytes? = #""
  count : exact-nonnegative-integer? = 1
(regexp-match-peek-positions/end  pattern
                                  input
                                 [start-pos
                                  end-pos
                                  progress
                                  input-prefix
                                  count])
   (or/c (cons/c (cons/c exact-nonnegative-integer?
                         exact-nonnegative-integer?)
                 (listof (or/c (cons/c exact-nonnegative-integer?
 →                                     exact-nonnegative-integer?)
                               #f)))
         #f)
   (or/c #f bytes?)
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input : input-port?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : (or/c exact-nonnegative-integer? #f) = #f
  progress : (or/c evt #f) = #f
  input-prefix : bytes? = #""
  count : exact-nonnegative-integer? = 1
```

```
(regexp-match-peek-positions-immediate/end  pattern
                                            input
                                           [start-pos
                                            end-pos
                                            progress
                                            input-prefix
                                            count])
   (or/c (cons/c (cons/c exact-nonnegative-integer?
                         exact-nonnegative-integer?)
                 (listof (or/c (cons/c exact-nonnegative-integer?
→                                       exact-nonnegative-integer?)
                               #f)))
         #f)
   (or/c #f bytes?)
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input : input-port?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : (or/c exact-nonnegative-integer? #f) = #f
  progress : (or/c evt #f) = #f
  input-prefix : bytes? = #""
  count : exact-nonnegative-integer? = 1
```

Like `regexp-match-positions`, etc., but with a second result like `regexp-match/end`.

4.7.5 Regexp Splitting

```
(regexp-split  pattern
               input
              [start-pos
               end-pos
               input-prefix])
   (if (and (or (string? pattern) (regexp? pattern))
            (string? input))
→       (cons/c string? (listof string?))
        (cons/c bytes? (listof bytes?)))
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input : (or/c string? bytes? input-port?)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : (or/c exact-nonnegative-integer? #f) = #f
  input-prefix : bytes? = #""
```

The complement of `regexp-match*`: the result is a list of strings (if *pattern* is a string or character regexp and *input* is a string) or byte strings (otherwise) from *input* that are

separated by matches to *pattern*. Adjacent matches are separated with "" or #"". Zero-length matches are treated the same as for regexp-match*.

If *input* contains no matches (in the range *start-pos* to *end-pos*), the result is a list containing *input*'s content (from *start-pos* to *end-pos*) as a single element. If a match occurs at the beginning of *input* (at *start-pos*), the resulting list will start with an empty string or byte string, and if a match occurs at the end (at *end-pos*), the list will end with an empty string or byte string. The *end-pos* argument can be #f, in which case splitting goes to the end of *input* (which corresponds to an end-of-file if *input* is an input port).

Examples:

```
> (regexp-split #rx" +" "12  34")
'("12" "34")
> (regexp-split #rx"." "12  34")
'("" "" "" "" "" "" "")
> (regexp-split #rx"" "12  34")
'("" "1" "2" " " " " " " "3" "4" "")
> (regexp-split #rx" *" "12  34")
'("" "1" "2" "" "3" "4" "")
> (regexp-split #px"\\b" "12, 13 and 14.")
'("" "12" ", " "13" " " "and" " " "14" ".")
> (regexp-split #rx" +" "")
'("")
```

4.7.6 Regexp Substitution

```
(regexp-replace  pattern
                 input
                 insert
                [input-prefix])
   (if (and (or (string? pattern) (regexp? pattern))
            (string? input))
→       string?
        bytes?)
 pattern : (or/c string? bytes? regexp? byte-regexp?)
 input : (or/c string? bytes?)
          (or/c string? bytes?
 insert :       ((string?) () #:rest (listof string?) . ->* . string?)
                ((bytes?) () #:rest (listof bytes?) . ->* . bytes?))
 input-prefix : bytes? = #""
```

Performs a match using *pattern* on *input*, and then returns a string or byte string in which the matching portion of *input* is replaced with *insert*. If *pattern* matches no part of *input*, then *input* is returned unmodified.

The *insert* argument can be either a (byte) string, or a function that returns a (byte) string. In the latter case, the function is applied on the list of values that regexp-match would return (i.e., the first argument is the complete match, and then one argument for each parenthesized sub-expression) to obtain a replacement (byte) string.

If *pattern* is a string or character regexp and *input* is a string, then *insert* must be a string or a procedure that accept strings, and the result is a string. If *pattern* is a byte string or byte regexp, or if *input* is a byte string, then *insert* as a string is converted to a byte string, *insert* as a procedure is called with a byte string, and the result is a byte string.

If *insert* contains &, then & is replaced with the matching portion of *input* before it is substituted into the match's place. If *insert* contains \⟨*n*⟩ for some integer ⟨*n*⟩, then it is replaced with the ⟨*n*⟩th matching sub-expression from *input*. A & and \0 are aliases. If the ⟨*n*⟩th sub-expression was not used in the match, or if ⟨*n*⟩ is greater than the number of sub-expressions in *pattern*, then \⟨*n*⟩ is replaced with the empty string.

To substitute a literal & or \, use \& and \\, respectively, in *insert*. A \$ in *insert* is equivalent to an empty sequence; this can be used to terminate a number ⟨*n*⟩ following \. If a \ in *insert* is followed by anything other than a digit, &, \, or $, then the \ by itself is treated as \0.

Note that the \ described in the previous paragraphs is a character or byte of *input*. To write such an *input* as a Racket string literal, an escaping \ is needed before the \. For example, the Racket constant "\\1" is \1.

Examples:

```
> (regexp-replace #rx"mi" "mi casa" "su")
"su casa"
> (regexp-replace #rx"mi" "mi casa" string-upcase)
"MI casa"
> (regexp-replace #rx"([Mm])i ([a-zA-Z]*)" "Mi Casa" "\\1y \\2")
"My Casa"
> (regexp-replace #rx"([Mm])i ([a-zA-Z]*)" "mi cerveza Mi Mi Mi"
                  "\\1y \\2")
"my cerveza Mi Mi Mi"
> (regexp-replace #rx"x" "12x4x6" "\\\\")
"12\\4x6"
> (display (regexp-replace #rx"x" "12x4x6" "\\\\"))
12\4x6
```

```
(regexp-replace* pattern
                 input
                 insert
                 [start-pos
                  end-pos
                  input-prefix]) → (or/c string? bytes?)
```

```
pattern : (or/c string? bytes? regexp? byte-regexp?)
input : (or/c string? bytes?)
         (or/c string? bytes?
insert :       ((string?) () #:rest (listof string?) . ->* . string?)
               ((bytes?) () #:rest (listof bytes?) . ->* . bytes?))
start-pos : exact-nonnegative-integer? = 0
end-pos : (or/c exact-nonnegative-integer? #f) = #f
input-prefix : bytes? = #""
```

Like `regexp-replace`, except that every instance of *pattern* in *input* is replaced with *insert*, instead of just the first match. Only non-overlapping instances of *pattern* in *input* are replaced, so instances of *pattern* within inserted strings are *not* replaced recursively. Zero-length matches are treated the same as in `regexp-match*`.

The optional *start-pos* and *end-pos* arguments select a portion of *input* for matching; the default is the entire string or the stream up to an end-of-file.

Examples:

```
> (regexp-replace* #rx"([Mm])i ([a-zA-Z]*)" "mi cerveza Mi Mi Mi"
                   "\\1y \\2")
"my cerveza My Mi Mi"
> (regexp-replace* #rx"([Mm])i ([a-zA-Z]*)" "mi cerveza Mi Mi Mi"
                   (lambda (all one two)
                     (string-append (string-downcase one) "y"
                                    (string-upcase two))))
"myCERVEZA myMI Mi"
> (regexp-replace* #px"\\w" "hello world" string-upcase 0 5)
"HELLO world"
> (display (regexp-replace* #rx"x" "12x4x6" "\\\\"))
12\4\6
```

```
(regexp-replaces input replacements) → (or/c string? bytes?)
  input : (or/c string? bytes?)
                (listof
                 (list/c (or/c string? bytes? regexp? byte-regexp?)
  replacements :          (or/c string? bytes?
                              ((string?) () #:rest (listof string?) . ->* . string?)
                              ((bytes?) () #:rest (listof bytes?) . ->* . bytes?))))
```

Performs a chain of `regexp-replace*` operations, where each element in *replacements* specifies a replacement as a `(list pattern replacement)`. The replacements are done in order, so later replacements can apply to previous insertions.

Examples:

```
> (regexp-replaces "zero-or-more?"
                   '([#rx"-" "_"] [#rx"(.*)\\?$" "is_\\1"]))
"is_zero_or_more"
> (regexp-replaces "zero-or-more?"
                   '([#rx"e" "o"] [#rx"o" "oo"]))
"zooroo-oor-mooroo?"
```

```
(regexp-replace-quote str) → string?
  str : string?
(regexp-replace-quote bstr) → bytes?
  bstr : bytes?
```

Produces a string suitable for use as the third argument to `regexp-replace` to insert the literal sequence of characters in *str* or bytes in *bstr* as a replacement. Concretely, every `\` and `&` in *str* or *bstr* is protected by a quoting `\`.

Examples:

```
> (regexp-replace #rx"UT" "Go UT!" "A&M")
"Go AUTM!"
> (regexp-replace #rx"UT" "Go UT!" (regexp-replace-quote "A&M"))
"Go A&M!"
```

4.8 Keywords

§3.7 “Keywords” in ***The Racket Guide*** introduces keywords.

A *keyword* is like an interned symbol, but its printed form starts with `#:`, and a keyword cannot be used as an identifier. Furthermore, a keyword by itself is not a valid expression, though a keyword can be `quoted` to form an expression that produces the symbol.

Two keywords are `eq?` if and only if they print the same (i.e., keywords are always interned).

Like symbols, keywords are only weakly held by the internal keyword table; see §4.6 “Symbols” for more information.

See §1.3.15 “Reading Keywords” for information on `read`ing keywords and §1.4.12 “Printing Keywords” for information on `print`ing keywords.

```
(keyword? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a keyword, `#f` otherwise.

```
(keyword->string keyword) → string?
  keyword : keyword?
```

Returns a string for the `display`ed form of *keyword*, not including the leading `#:`.

```
(string->keyword str) → keyword?
  str : string?
```

Returns a keyword whose displayed form is the same as that of *str*, but with a leading `#:`.

```
(keyword<? a-keyword b-keyword ...+) → boolean?
  a-keyword : keyword?
  b-keyword : keyword?
```

Returns `#t` if the arguments are sorted, where the comparison for each pair of keywords is the same as using `keyword->string` with `string->bytes/utf-8` and `bytes<?`.

4.9 Pairs and Lists

§3.8 “Pairs and Lists” in ***The Racket Guide*** introduces pairs and lists.

A *pair* combines exactly two values. The first value is accessed with the `car` procedure, and the second value is accessed with the `cdr` procedure. Pairs are not mutable (but see §4.10 “Mutable Pairs and Lists”).

A *list* is recursively defined: it is either the constant `null`, or it is a pair whose second value is a list.

A list can be used as a single-valued sequence (see §4.14.1 “Sequences”). The elements of the list serve as elements of the sequence. See also `in-list`.

Cyclic data structures can be created using only immutable pairs via `read` or `make-reader-graph`. If starting with a pair and using some number of `cdr`s returns to the starting pair, then the pair is not a list.

See §1.3.6 “Reading Pairs and Lists” for information on reading pairs and lists and §1.4.5 “Printing Pairs and Lists” for information on printing pairs and lists.

4.9.1 Pair Constructors and Selectors

```
(pair? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a pair, `#f` otherwise.

Examples:

```
> (pair? 1)
```

```
#f
> (pair? (cons 1 2))
#t
> (pair? (list 1 2))
#t
> (pair? '(1 2))
#t
> (pair? '())
#f
```

```
(null? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is the empty list, `#f` otherwise.

Examples:

```
> (null? 1)
#f
> (null? '(1 2))
#f
> (null? '())
#t
> (null? (cdr (list 1)))
#t
```

```
(cons a d) → pair?
  a : any/c
  d : any/c
```

Returns a newly allocated pair whose first element is `a` and second element is *d*.

Examples:

```
> (cons 1 2)
'(1 . 2)
> (cons 1 '())
'(1)
```

```
(car p) → any/c
  p : pair?
```

Returns the first element of the pair *p*.

Examples:

```
> (car '(1 2))
1
> (car (cons 2 3))
2
```

```
(cdr p) → any/c
  p : pair?
```

Returns the second element of the pair *p*.

Examples:

```
> (cdr '(1 2))
'(2)
> (cdr '(1))
'()
```

```
null : null?
```

The empty list.

Examples:

```
> null
'()
> '()
'()
> (eq? '() null)
#t
```

```
(list? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a list: either the empty list, or a pair whose second element is a list. This procedure effectively takes constant time due to internal caching (so that any necessary traversals of pairs can in principle count as an extra cost of allocating the pairs).

Examples:

```
> (list? '(1 2))
#t
> (list? (cons 1 (cons 2 '())))
#t
> (list? (cons 1 2))
#f
```

```
(list v ...) → list?
  v : any/c
```

Returns a newly allocated list containing the `v`s as its elements.

Examples:

```
> (list 1 2 3 4)
'(1 2 3 4)
> (list (list 1 2) (list 3 4))
'((1 2) (3 4))
```

```
(list* v ... tail) → any/c
  v : any/c
  tail : any/c
```

Like `list`, but the last argument is used as the tail of the result, instead of the final element. The result is a list only if the last argument is a list.

Examples:

```
> (list* 1 2)
'(1 . 2)
> (list* 1 2 (list 3 4))
'(1 2 3 4)
```

```
(build-list n proc) → list?
  n : exact-nonnegative-integer?
  proc : (exact-nonnegative-integer? . -> . any)
```

Creates a list of `n` elements by applying `proc` to the integers from `0` to `(sub1 n)` in order. If `lst` is the resulting list, then `(list-ref lst i)` is the value produced by `(proc i)`.

Examples:

```
> (build-list 10 values)
'(0 1 2 3 4 5 6 7 8 9)
> (build-list 5 (lambda (x) (* x x)))
'(0 1 4 9 16)
```

4.9.2 List Operations

```
(length lst) → exact-nonnegative-integer?
  lst : list?
```

Returns the number of elements in `lst`.

Examples:

```
> (length (list 1 2 3 4))
4
> (length '())
0
```

```
(list-ref lst pos) → any/c
  lst : pair?
  pos : exact-nonnegative-integer?
```

Returns the element of `lst` at position `pos`, where the list's first element is position `0`. If the list has `pos` or fewer elements, then the `exn:fail:contract` exception is raised.

The `lst` argument need not actually be a list; `lst` must merely start with a chain of at least `(add1 pos)` pairs.

Examples:

```
> (list-ref (list 'a 'b 'c) 0)
'a
> (list-ref (list 'a 'b 'c) 1)
'b
> (list-ref (list 'a 'b 'c) 2)
'c
> (list-ref (cons 1 2) 0)
1
```

```
(list-tail lst pos) → any/c
  lst : any/c
  pos : exact-nonnegative-integer?
```

Returns the list after the first `pos` elements of `lst`. If the list has fewer than `pos` elements, then the `exn:fail:contract` exception is raised.

The `lst` argument need not actually be a list; `lst` must merely start with a chain of at least `pos` pairs.

Examples:

```
> (list-tail (list 1 2 3 4) 2)
'(3 4)
> (list-ref (cons 1 2) 1)
list-ref: index reaches a non-pair
  index: 1
  in: '(1 . 2)
> (list-ref 'not-a-pair 0)
list-ref: contract violation
  expected: pair?
```

given: 'not-a-pair
argument position: 1st
other arguments...:
0

```
(append lst ...) → list?
  lst : list?
(append lst ... v) → any/c
  lst : list?
  v : any/c
```

When given all list arguments, the result is a list that contains all of the elements of the given lists in order. The last argument is used directly in the tail of the result.

The last argument need not be a list, in which case the result is an "improper list."

Examples:

```
> (append (list 1 2) (list 3 4))
'(1 2 3 4)
> (append (list 1 2) (list 3 4) (list 5 6) (list 7 8))
'(1 2 3 4 5 6 7 8)
```

```
(reverse lst) → list?
  lst : list?
```

Returns a list that has the same elements as *lst*, but in reverse order.

Example:

```
> (reverse (list 1 2 3 4))
'(4 3 2 1)
```

4.9.3 List Iteration

```
(map proc lst ...+) → list?
  proc : procedure?
  lst : list?
```

Applies *proc* to the elements of the *lst*s from the first elements to the last. The *proc* argument must accept the same number of arguments as the number of supplied *lst*s, and all *lst*s must have the same number of elements. The result is a list containing each result of *proc* in order.

Examples:

```
> (map (lambda (number)
         (+ 1 number))
       '(1 2 3 4))
'(2 3 4 5)
> (map (lambda (number1 number2)
         (+ number1 number2))
       '(1 2 3 4)
       '(10 100 1000 10000))
'(11 102 1003 10004)
```

(andmap *proc lst ...+*) → any
  proc : procedure?
  lst : list?

Similar to map in the sense that *proc* is applied to each element of *lst*, but

- the result is #f if any application of *proc* produces #f, in which case *proc* is not applied to later elements of the *lst*s; and
- the result is that of *proc* applied to the last elements of the *lst*s; more specifically, the application of *proc* to the last elements in the *lst*s is in tail position with respect to the andmap call.

If the *lst*s are empty, then #t is returned.

The andmap function is actually closer to foldl than map, since andmap doesn't produce a list. Still, (andmap f (list x y z)) is equivalent to (and (f x) (f y) (f z)) in the same way that (map f (list x y z)) is equivalent to (list (f x) (f y) (f z)).

Examples:

```
> (andmap positive? '(1 2 3))
#t
> (andmap positive? '(1 2 a))
positive?: contract violation
  expected: real?
  given: 'a
> (andmap positive? '(1 -2 a))
#f
> (andmap + '(1 2 3) '(4 5 6))
9
```

(ormap *proc lst ...+*) → any
  proc : procedure?
  lst : list?

Similar to map in the sense that *proc* is applied to each element of *lst*, but

- the result is #f if every application of *proc* produces #f; and

To continue the andmap note above, (ormap f (list x y z)) is equivalent to (or (f x) (f y) (f z)).

- the result is that of the first application of *proc* producing a value other than #f, in which case *proc* is not applied to later elements of the *lst*s; the application of *proc* to the last elements of the *lst*s is in tail position with respect to the ormap call.

If the *lst*s are empty, then #f is returned.

Examples:

```
> (ormap eq? '(a b c) '(a b c))
#t
> (ormap positive? '(1 2 a))
#t
> (ormap + '(1 2 3) '(4 5 6))
5
```

```
(for-each proc lst ...+) → void?
  proc : procedure?
  lst : list?
```

Similar to map, but *proc* is called only for its effect, and its result (which can be any number of values) is ignored.

Example:

```
> (for-each (lambda (arg)
              (printf "Got ~a\n" arg)
              23)
            '(1 2 3 4))
Got 1
Got 2
Got 3
Got 4
```

```
(foldl proc init lst ...+) → any/c
  proc : procedure?
  init : any/c
  lst : list?
```

Like map, foldl applies a procedure to the elements of one or more lists. Whereas map combines the return values into a list, foldl combines the return values in an arbitrary way that is determined by *proc*.

If foldl is called with n lists, then *proc* must take $n+1$ arguments. The extra argument is the combined return values so far. The *proc* is initially invoked with the first item of each list, and the final argument is *init*. In subsequent invocations of *proc*, the last argument

is the return value from the previous invocation of *proc*. The input *lst*s are traversed from left to right, and the result of the whole foldl application is the result of the last application of *proc*. If the *lst*s are empty, the result is *init*.

Unlike foldr, foldl processes the *lst*s in constant space (plus the space for each call to *proc*).

Examples:

```
> (foldl cons '() '(1 2 3 4))
'(4 3 2 1)
> (foldl + 0 '(1 2 3 4))
10
> (foldl (lambda (a b result)
           (* result (- a b)))
         1
         '(1 2 3)
         '(4 5 6))
-27
```

```
(foldr proc init lst ...+) → any/c
  proc : procedure?
  init : any/c
  lst : list?
```

Like foldl, but the lists are traversed from right to left. Unlike foldl, foldr processes the *lst*s in space proportional to the length of *lst*s (plus the space for each call to *proc*).

Examples:

```
> (foldr cons '() '(1 2 3 4))
'(1 2 3 4)
> (foldr (lambda (v l) (cons (add1 v) l)) '() '(1 2 3 4))
'(2 3 4 5)
```

4.9.4 List Filtering

```
(filter pred lst) → list?
  pred : procedure?
  lst : list?
```

Returns a list with the elements of *lst* for which *pred* produces a true value. The *pred* procedure is applied to each element from first to last.

Example:

```
> (filter positive? '(1 -2 3 4 -5))
'(1 3 4)
```

```
(remove v lst [proc]) → list?
  v : any/c
  lst : list?
  proc : procedure? = equal?
```

Returns a list that is like *lst*, omitting the first element of *lst* that is equal to *v* using the comparison procedure *proc* (which must accept two arguments).

Examples:

```
> (remove 2 (list 1 2 3 2 4))
'(1 3 2 4)
> (remove 2 (list 1 2 3 2 4) =)
'(1 3 2 4)
> (remove '(2) (list '(1) '(2) '(3)))
'((1) (3))
> (remove "2" (list "1" "2" "3"))
'("1" "3")
> (remove #\c (list #\a #\b #\c))
'(#\a #\b)
```

```
(remq v lst) → list?
  v : any/c
  lst : list?
```

Returns (remove *v* *lst* eq?).

Examples:

```
> (remq 2 (list 1 2 3 4 5))
'(1 3 4 5)
> (remq '(2) (list '(1) '(2) '(3)))
'((1) (2) (3))
> (remq "2" (list "1" "2" "3"))
'("1" "3")
> (remq #\c (list #\a #\b #\c))
'(#\a #\b)
```

```
(remv v lst) → list?
  v : any/c
  lst : list?
```

Returns (remove *v* *lst* eqv?).

Examples:

```
> (remv 2 (list 1 2 3 4 5))
'(1 3 4 5)
> (remv '(2) (list '(1) '(2) '(3)))
'((1) (2) (3))
> (remv "2" (list "1" "2" "3"))
'("1" "3")
> (remv #\c (list #\a #\b #\c))
'(#\a #\b)
```

```
(remove* v-lst lst [proc]) → list?
  v-lst : list?
  lst : list?
  proc : procedure? = equal?
```

Like `remove`, but removes from *lst* every instance of every element of *v-lst*.

Example:

```
> (remove* (list 1 2) (list 1 2 3 2 4 5 2))
'(3 4 5)
```

```
(remq* v-lst lst) → list?
  v-lst : list?
  lst : list?
```

Returns `(remove* v-lst lst eq?)`.

Example:

```
> (remq* (list 1 2) (list 1 2 3 2 4 5 2))
'(3 4 5)
```

```
(remv* v-lst lst) → list?
  v-lst : list?
  lst : list?
```

Returns `(remove* v-lst lst eqv?)`.

Example:

```
> (remv* (list 1 2) (list 1 2 3 2 4 5 2))
'(3 4 5)
```

```
(sort lst
      less-than?
     [#:key extract-key
      #:cache-keys? cache-keys?]) → list?
```

```
lst : list?
less-than? : (any/c any/c . -> . any/c)
extract-key : (any/c . -> . any/c) = (lambda (x) x)
cache-keys? : boolean? = #f
```

Returns a list sorted according to the *less-than?* procedure, which takes two elements of *lst* and returns a true value if the first is less (i.e., should be sorted earlier) than the second.

The sort is stable; if two elements of *lst* are "equal" (i.e., `proc` does not return a true value when given the pair in either order), then the elements preserve their relative order from *lst* in the output list. To preserve this guarantee, use `sort` with a strict comparison functions (e.g., `<` or `string<?`; not `<=` or `string<=?`).

The `#:key` argument *extract-key* is used to extract a key value for comparison from each list element. That is, the full comparison procedure is essentially

```
(lambda (x y)
  (less-than? (extract-key x) (extract-key y)))
```

By default, *extract-key* is applied to two list elements for every comparison, but if *cache-keys?* is true, then the *extract-key* function is used exactly once for each list item. Supply a true value for *cache-keys?* when *extract-key* is an expensive operation; for example, if `file-or-directory-modify-seconds` is used to extract a timestamp for every file in a list, then *cache-keys?* should be `#t` to minimize file-system calls, but if *extract-key* is `car`, then *cache-keys?* should be `#f`. As another example, providing *extract-key* as `(lambda (x) (random))` and `#t` for *cache-keys?* effectively shuffles the list.

Examples:

```
> (sort '(1 3 4 2) <)
'(1 2 3 4)
> (sort '("aardvark" "dingo" "cow" "bear") string<?)
'("aardvark" "bear" "cow" "dingo")
> (sort '(("aardvark") ("dingo") ("cow") ("bear"))
        #:key car string<?)
'(("aardvark") ("bear") ("cow") ("dingo"))
```

4.9.5 List Searching

```
(member v lst [is-equal?]) → (or/c list? #f)
  v : any/c
  lst : list?
  is-equal? : (any/c any/c -> any/c) = equal?
```

Locates the first element of *lst* that is equal? to *v*. If such an element exists, the tail of *lst* starting with that element is returned. Otherwise, the result is #f.

Examples:

```
> (member 2 (list 1 2 3 4))
'(2 3 4)
> (member 9 (list 1 2 3 4))
#f
> (member #'x (list #'x #'y) free-identifier=?)
'(#<syntax:491:0 x> #<syntax:491:0 y>)
> (member #'a (list #'x #'y) free-identifier=?)
#f
```

```
(memv v lst) → (or/c list? #f)
  v : any/c
  lst : list?
```

Like member, but finds an element using eqv?.

Examples:

```
> (memv 2 (list 1 2 3 4))
'(2 3 4)
> (memv 9 (list 1 2 3 4))
#f
```

```
(memq v lst) → (or/c list? #f)
  v : any/c
  lst : list?
```

Like member, but finds an element using eq?.

Examples:

```
> (memq 2 (list 1 2 3 4))
'(2 3 4)
> (memq 9 (list 1 2 3 4))
#f
```

```
(memf proc lst) → (or/c list? #f)
  proc : procedure?
  lst : list?
```

Like member, but finds an element using the predicate *proc*; an element is found when *proc* applied to the element returns a true value.

Example:

```
> (memf (lambda (arg)
          (> arg 9))
        '(7 8 9 10 11))
'(10 11)
```

```
(findf proc lst) → any/c
  proc : procedure?
  lst : list?
```

Like memf, but returns the element or #f instead of a tail of lst or #f.

Example:

```
> (findf (lambda (arg)
           (> arg 9))
         '(7 8 9 10 11))
10
```

```
(assoc v lst [is-equal?]) → (or/c pair? #f)
  v : any/c
  lst : (listof pair?)
  is-equal? : (any/c any/c -> any/c) = equal?
```

Locates the first element of lst whose car is equal to v according to is-equal?. If such an element exists, the pair (i.e., an element of lst) is returned. Otherwise, the result is #f.

Examples:

```
> (assoc 3 (list (list 1 2) (list 3 4) (list 5 6)))
'(3 4)
> (assoc 9 (list (list 1 2) (list 3 4) (list 5 6)))
#f
> (assoc 3.5
         (list (list 1 2) (list 3 4) (list 5 6))
         (lambda (a b) (< (abs (- a b)) 1)))
'(3 4)
```

```
(assv v lst) → (or/c pair? #f)
  v : any/c
  lst : (listof pair?)
```

Like assoc, but finds an element using eqv?.

Example:

```
> (assv 3 (list (list 1 2) (list 3 4) (list 5 6)))
'(3 4)
```

```
(assq v lst) → (or/c pair? #f)
  v : any/c
  lst : (listof pair?)
```

Like assoc, but finds an element using eq?.

Example:

```
> (assq 3 (list (list 1 2) (list 3 4) (list 5 6)))
'(3 4)
```

```
(assf proc lst) → (or/c list? #f)
  proc : procedure?
  lst : list?
```

Like assoc, but finds an element using the predicate *proc*; an element is found when *proc* applied to the car of an *lst* element returns a true value.

Example:

```
> (assf (lambda (arg)
          (> arg 2))
        (list (list 1 2) (list 3 4) (list 5 6)))
'(3 4)
```

4.9.6 Pair Accessor Shorthands

```
(caar v) → any/c
  v : (cons/c pair? any/c)
```

Returns (car (car p))

Example:

```
> (caar '((1 2) 3 4))
1
```

```
(cadr v) → any/c
  v : (cons/c any/c pair?)
```

Returns (car (cdr p))

Example:

```
> (cadr '((1 2) 3 4))
3
```

```
(cdar v) → any/c
  v : (cons/c pair? any/c)
```

Returns `(cdr (car p))`

Example:

```
> (cdar '((7 6 5 4 3 2 1) 8 9))
'(6 5 4 3 2 1)
```

```
(cddr v) → any/c
  v : (cons/c any/c pair?)
```

Returns `(cdr (cdr p))`

Example:

```
> (cddr '(2 1))
'()
```

```
(caaar v) → any/c
  v : (cons/c (cons/c pair? any/c) any/c)
```

Returns `(car (car (car p)))`

Example:

```
> (caaar '(((6 5 4 3 2 1) 7) 8 9))
6
```

```
(caadr v) → any/c
  v : (cons/c any/c (cons/c pair? any/c))
```

Returns `(car (car (cdr p)))`

Example:

```
> (caadr '(9 (7 6 5 4 3 2 1) 8))
7
```

```
(cadar v) → any/c
  v : (cons/c (cons/c any/c pair?) any/c)
```

Returns `(car (cdr (car p)))`

Example:

```
> (cadar '((7 6 5 4 3 2 1) 8 9))
6
```

```
(caddr v) → any/c
  v : (cons/c any/c (cons/c any/c pair?))
```

Returns `(car (cdr (cdr p)))`

Example:

```
> (caddr '(3 2 1))
1
```

```
(cdaar v) → any/c
  v : (cons/c (cons/c pair? any/c) any/c)
```

Returns `(cdr (car (car p)))`

Example:

```
> (cdaar '(((6 5 4 3 2 1) 7) 8 9))
'(5 4 3 2 1)
```

```
(cdadr v) → any/c
  v : (cons/c any/c (cons/c pair? any/c))
```

Returns `(cdr (car (cdr p)))`

Example:

```
> (cdadr '(9 (7 6 5 4 3 2 1) 8))
'(6 5 4 3 2 1)
```

```
(cddar v) → any/c
  v : (cons/c (cons/c any/c pair?) any/c)
```

Returns `(cdr (cdr (car p)))`

Example:

```
> (cddar '((7 6 5 4 3 2 1) 8 9))
'(5 4 3 2 1)
```

```
(cdddr v) → any/c
  v : (cons/c any/c (cons/c any/c pair?))
```

Returns `(cdr (cdr (cdr p)))`

Example:

```
> (cdddr '(3 2 1))
'()
```

```
(caaaar v) → any/c
  v : (cons/c (cons/c (cons/c pair? any/c) any/c) any/c)
```

Returns `(car (car (car (car p))))`

Example:

```
> (caaaar '((((5 4 3 2 1) 6) 7) 8 9))
5
```

```
(caaadr v) → any/c
  v : (cons/c any/c (cons/c (cons/c pair? any/c) any/c))
```

Returns `(car (car (car (cdr p))))`

Example:

```
> (caaadr '(9 ((6 5 4 3 2 1) 7) 8))
6
```

```
(caadar v) → any/c
  v : (cons/c (cons/c any/c (cons/c pair? any/c)) any/c)
```

Returns `(car (car (cdr (car p))))`

Example:

```
> (caadar '((7 (5 4 3 2 1) 6) 8 9))
5
```

```
(caaddr v) → any/c
  v : (cons/c any/c (cons/c any/c (cons/c pair? any/c)))
```

Returns `(car (car (cdr (cdr p))))`

Example:

```
> (caaddr '(9 8 (6 5 4 3 2 1) 7))
6
```

```
(cadaar v) → any/c
  v : (cons/c (cons/c (cons/c any/c pair?) any/c) any/c)
```

Returns `(car (cdr (car (car p))))`

Example:

```
> (cadaar '(((6 5 4 3 2 1) 7) 8 9))
5
```

```
(cadadr v) → any/c
  v : (cons/c any/c (cons/c (cons/c any/c pair?) any/c))
```

Returns `(car (cdr (car (cdr p))))`

Example:

```
> (cadadr '(9 (7 6 5 4 3 2 1) 8))
6
```

```
(caddar v) → any/c
  v : (cons/c (cons/c any/c (cons/c any/c pair?)) any/c)
```

Returns `(car (cdr (cdr (car p))))`

Example:

```
> (caddar '((7 6 5 4 3 2 1) 8 9))
5
```

```
(cadddr v) → any/c
  v : (cons/c any/c (cons/c any/c (cons/c any/c pair?)))
```

Returns `(car (cdr (cdr (cdr p))))`

Example:

```
> (cadddr '(4 3 2 1))
1
```

```
(cdaaar v) → any/c
  v : (cons/c (cons/c (cons/c pair? any/c) any/c) any/c)
```

Returns `(cdr (car (car (car p))))`

Example:

```
> (cdaaar '((((5 4 3 2 1) 6) 7) 8 9))
'(4 3 2 1)
```

```
(cdaadr v) → any/c
  v : (cons/c any/c (cons/c (cons/c pair? any/c) any/c))
```

Returns `(cdr (car (car (cdr p))))`

Example:

```
> (cdaadr '(9 ((6 5 4 3 2 1) 7) 8))
'(5 4 3 2 1)
```

```
(cdadar v) → any/c
  v : (cons/c (cons/c any/c (cons/c pair? any/c)) any/c)
```

Returns `(cdr (car (cdr (car p))))`

Example:

```
> (cdadar '((7 (5 4 3 2 1) 6) 8 9))
'(4 3 2 1)
```

```
(cdaddr v) → any/c
  v : (cons/c any/c (cons/c any/c (cons/c pair? any/c)))
```

Returns `(cdr (car (cdr (cdr p))))`

Example:

```
> (cdaddr '(9 8 (6 5 4 3 2 1) 7))
'(5 4 3 2 1)
```

```
(cddaar v) → any/c
  v : (cons/c (cons/c (cons/c any/c pair?) any/c) any/c)
```

Returns `(cdr (cdr (car (car p))))`

Example:

```
> (cddaar '(((6 5 4 3 2 1) 7) 8 9))
'(4 3 2 1)
```

```
(cddadr v) → any/c
  v : (cons/c any/c (cons/c (cons/c any/c pair?) any/c))
```

Returns `(cdr (cdr (car (cdr p))))`

Example:

```
> (cddadr '(9 (7 6 5 4 3 2 1) 8))
'(5 4 3 2 1)
```

```
(cdddar v) → any/c
  v : (cons/c (cons/c any/c (cons/c any/c pair?)) any/c)
```

Returns `(cdr (cdr (cdr (car p))))`

Example:

```
> (cdddar '((7 6 5 4 3 2 1) 8 9))
'(4 3 2 1)
```

```
(cddddr v) → any/c
  v : (cons/c any/c (cons/c any/c (cons/c any/c pair?)))
```

Returns `(cdr (cdr (cdr (cdr p))))`

Example:

```
> (cddddr '(4 3 2 1))
'()
```

4.9.7 Additional List Functions and Synonyms

```
(require racket/list)
```
package: base

The bindings documented in this section are provided by the `racket/list` and `racket` libraries, but not `racket/base`.

```
empty : null?
```

The empty list.

Examples:

```
> empty
'()
> (eq? empty null)
#t
```

```
(cons? v) → boolean?
  v : any/c
```

The same as `(pair? v)`.

Example:

```
> (cons? '(1 2))
#t
```

```
(empty? v) → boolean?
  v : any/c
```

The same as `(null? v)`.

Examples:

```
> (empty? '(1 2))
#f
> (empty? '())
#t
```

```
(first lst) → any/c
  lst : list?
```

The same as `(car lst)`, but only for lists (that are not empty).

Example:

```
> (first '(1 2 3 4 5 6 7 8 9 10))
1
```

```
(rest lst) → list?
  lst : list?
```

The same as `(cdr lst)`, but only for lists (that are not empty).

Example:

```
> (rest '(1 2 3 4 5 6 7 8 9 10))
'(2 3 4 5 6 7 8 9 10)
```

```
(second lst) → any
  lst : list?
```

Returns the second element of the list.

Example:

```
> (second '(1 2 3 4 5 6 7 8 9 10))
2
```

```
(third lst) → any
  lst : list?
```

Returns the third element of the list.

Example:

```
> (third '(1 2 3 4 5 6 7 8 9 10))
3
```

```
(fourth lst) → any
  lst : list?
```

Returns the fourth element of the list.

Example:

```
> (fourth '(1 2 3 4 5 6 7 8 9 10))
4
```

```
(fifth lst) → any
  lst : list?
```

Returns the fifth element of the list.

Example:

```
> (fifth '(1 2 3 4 5 6 7 8 9 10))
5
```

```
(sixth lst) → any
  lst : list?
```

Returns the sixth element of the list.

Example:

```
> (sixth '(1 2 3 4 5 6 7 8 9 10))
6
```

```
(seventh lst) → any
  lst : list?
```

Returns the seventh element of the list.

Example:

```
> (seventh '(1 2 3 4 5 6 7 8 9 10))
7
```

```
(eighth lst) → any
  lst : list?
```

Returns the eighth element of the list.

Example:

```
> (eighth '(1 2 3 4 5 6 7 8 9 10))
8
```

```
(ninth lst) → any
  lst : list?
```

Returns the ninth element of the list.

Example:

```
> (ninth '(1 2 3 4 5 6 7 8 9 10))
9
```

```
(tenth lst) → any
  lst : list?
```

Returns the tenth element of the list.

Example:

```
> (tenth '(1 2 3 4 5 6 7 8 9 10))
10
```

```
(last lst) → any
  lst : list?
```

Returns the last element of the list.

Example:

```
> (last '(1 2 3 4 5 6 7 8 9 10))
10
```

```
(last-pair p) → pair?
  p : pair?
```

Returns the last pair of a (possibly improper) list.

Example:

```
> (last-pair '(1 2 3 4))
'(4)
```

```
(make-list k v) → list?
  k : exact-nonnegative-integer?
  v : any/c
```

Returns a newly constructed list of length *k*, holding *v* in all positions.

Example:

```
> (make-list 7 'foo)
'(foo foo foo foo foo foo foo)
```

```
(take lst pos) → list?
  lst : any/c
  pos : exact-nonnegative-integer?
```

Returns a fresh list whose elements are the first *pos* elements of *lst*. If *lst* has fewer than *pos* elements, the `exn:fail:contract` exception is raised.

The *lst* argument need not actually be a list; *lst* must merely start with a chain of at least *pos* pairs.

Examples:

```
> (take '(1 2 3 4) 2)
'(1 2)
> (take 'non-list 0)
'()
```

```
(drop lst pos) → any/c
  lst : any/c
  pos : exact-nonnegative-integer?
```

Just like `list-tail`.

```
(split-at lst pos) → list? any/c
  lst : any/c
  pos : exact-nonnegative-integer?
```

Returns the same result as

```
(values (take lst pos) (drop lst pos))
```

except that it can be faster.

```
(takef lst pred) → list?
  lst : any/c
  pred : procedure?
```

Returns a fresh list whose elements are taken successively from *lst* as long as they satisfy *pred*. The returned list includes up to, but not including, the first element in *lst* for which *pred* returns `#f`.

The *lst* argument need not actually be a list; the chain of pairs in *lst* will be traversed until a non-pair is encountered.

Examples:

```
> (takef '(2 4 5 8) even?)
'(2 4)
> (takef '(2 4 6 8) odd?)
'()
> (takef '(2 4 . 6) even?)
'(2 4)
```

```
(dropf lst pred) → any/c
  lst : any/c
  pred : procedure?
```

Drops elements from the front of *lst* as long as they satisfy *pred*.

Examples:

```
> (dropf '(2 4 5 8) even?)
'(5 8)
> (dropf '(2 4 6 8) odd?)
'(2 4 6 8)
```

```
(splitf-at lst pred) → list? any/c
  lst : any/c
  pred : procedure?
```

Returns the same result as

```
(values (takef lst pred) (dropf lst pred))
```

except that it can be faster.

```
(take-right lst pos) → any/c
  lst : any/c
  pos : exact-nonnegative-integer?
```

Returns the `list`'s *pos*-length tail. If *lst* has fewer than *pos* elements, then the `exn:fail:contract` exception is raised.

The *lst* argument need not actually be a list; *lst* must merely end with a chain of at least *pos* pairs.

Examples:

```
> (take-right '(1 2 3 4) 2)
'(3 4)
> (take-right 'non-list 0)
'non-list
```

```
(drop-right lst pos) → list?
  lst : any/c
  pos : exact-nonnegative-integer?
```

Returns a fresh list whose elements are the prefix of *lst*, dropping its *pos*-length tail. If *lst* has fewer than *pos* elements, then the `exn:fail:contract` exception is raised.

The *lst* argument need not actually be a list; *lst* must merely end with a chain of at least *pos* pairs.

Examples:

```
> (drop-right '(1 2 3 4) 2)
'(1 2)
> (drop-right 'non-list 0)
'()
```

```
(split-at-right lst pos) → list? any/c
  lst : any/c
  pos : exact-nonnegative-integer?
```

Returns the same result as

```
(values (drop-right lst pos) (take-right lst pos))
```

except that it can be faster.

Examples:

```
> (split-at-right '(1 2 3 4 5 6) 3)
'(1 2 3)
'(4 5 6)
> (split-at-right '(1 2 3 4 5 6) 4)
'(1 2)
'(3 4 5 6)
```

```
(takef-right lst pred) → list?
  lst : any/c
  pred : procedure?
(dropf-right lst pred) → any/c
  lst : any/c
  pred : procedure?
(splitf-at-right lst pred) → list? any/c
  lst : any/c
  pred : procedure?
```

Like `takef`, `dropf`, and `splitf-at`, but combined with the from-right functionality of `take-right`, `drop-right`, and `split-at-right`.

```
(add-between  lst
              v
             [#:before-first before-first
              #:before-last before-last
              #:after-last after-last
              #:splice? splice?])          → list?
  lst : list?
  v : any/c
  before-first : list? = '()
```

```
  before-last : any/c = v
  after-last : list? = '()
  splice? : any/c = #f
```

Returns a list with the same elements as *lst*, but with *v* between each pair of elements in *lst*; the last pair of elements will have *before-last* between them, instead of *v* (but *before-last* defaults to *v*).

If *splice?* is true, then *v* and *before-last* should be lists, and the list elements are spliced into the result. In addition, when *splice?* is true, *before-first* and *after-last* are inserted before the first element and after the last element respectively.

Examples:

```
> (add-between '(x y z) 'and)
'(x and y and z)
> (add-between '(x) 'and)
'(x)
> (add-between '("a" "b" "c" "d") "," #:before-last "and")
'("a" "," "b" "," "c" "and" "d")
> (add-between '(x y z) '(-) #:before-last '(- -)
               #:before-first '(begin) #:after-last '(end LF)
               #:splice? #t)
'(begin x - y - - z end LF)
```

```
(append* lst ... lsts) → list?
  lst : list?
  lsts : (listof list?)
(append* lst ... lsts) → any/c
  lst : list?
  lsts : list?
```

Like append, but the last argument is used as a list of arguments for append, so (append* *lst* ... *lsts*) is the same as (apply append *lst* ... *lsts*). In other words, the relationship between append and append* is similar to the one between list and list*.

Examples:

```
> (append* '(a) '(b) '((c) (d)))
'(a b c d)
> (cdr (append* (map (lambda (x) (list ", " x))
                     '("Alpha" "Beta" "Gamma"))))
'("Alpha" ", " "Beta" ", " "Gamma")
```

```
(flatten v) → list?
  v : any/c
```

Flattens an arbitrary S-expression structure of pairs into a list. More precisely, *v* is treated as a binary tree where pairs are interior nodes, and the resulting list contains all of the non-`null` leaves of the tree in the same order as an inorder traversal.

Examples:

```
> (flatten '((a) b (c (d) . e) ()))
'(a b c d e)
> (flatten 'a)
'(a)
```

```
(remove-duplicates lst
                   [same?
                    #:key extract-key]) → list?
  lst : list?
  same? : (any/c any/c . -> . any/c) = equal?
  extract-key : (any/c . -> . any/c) = (lambda (x) x)
```

Returns a list that has all items in *lst*, but without duplicate items, where *same?* determines whether two elements of the list are equivalent. The resulting list is in the same order as *lst*, and for any item that occurs multiple times, the first one is kept.

The `#:key` argument *extract-key* is used to extract a key value from each list element, so two items are considered equal if `(same? (extract-key x) (extract-key y))` is true.

Examples:

```
> (remove-duplicates '(a b b a))
'(a b)
> (remove-duplicates '(1 2 1.0 0))
'(1 2 1.0 0)
> (remove-duplicates '(1 2 1.0 0) =)
'(1 2 0)
```

```
(filter-map proc lst ...+) → list?
  proc : procedure?
  lst : list?
```

Returns `(filter (lambda (x) x) (map proc lst ...))`, but without building the intermediate list.

Example:

```
> (filter-map (lambda (x) (and (positive? x) x)) '(1 2 3 -2 8))
'(1 2 3 8)
```

```
(count proc lst ...+) → exact-nonnegative-integer?
  proc : procedure?
  lst : list?
```

Returns `(length (filter proc lst ...))`, but without building the intermediate list.

Example:

```
> (count positive? '(1 -1 2 3 -2 5))
4
```

```
(partition pred lst) → list? list?
  pred : procedure?
  lst : list?
```

Similar to `filter`, except that two values are returned: the items for which *pred* returns a true value, and the items for which *pred* returns `#f`.

The result is the same as

```
(values (filter pred lst) (filter (negate pred) lst))
```

but *pred* is applied to each item in *lst* only once.

Example:

```
> (partition even? '(1 2 3 4 5 6))
'(2 4 6)
'(1 3 5)
```

```
(range end) → list?
  end : real?
(range start end [step]) → list?
  start : real?
  end : real?
  step : real? = 1
```

Similar to `in-range`, but returns lists.

The resulting list holds numbers starting at *start* and whose successive elements are computed by adding *step* to their predecessor until *end* (excluded) is reached. If no starting point is provided, `0` is used. If no *step* argument is provided, `1` is used.

Examples:

```
> (range 10)
'(0 1 2 3 4 5 6 7 8 9)
> (range 10 20)
'(10 11 12 13 14 15 16 17 18 19)
> (range 20 40 2)
'(20 22 24 26 28 30 32 34 36 38)
> (range 20 10 -1)
'(20 19 18 17 16 15 14 13 12 11)
> (range 10 15 1.5)
'(10 11.5 13.0 14.5)
```

```
(append-map proc lst ...+) → list?
  proc : procedure?
  lst : list?
```

Returns `(append* (map proc lst ...))`.

Example:

```
> (append-map vector->list '(#(1) #(2 3) #(4)))
'(1 2 3 4)
```

```
(filter-not pred lst) → list?
  pred : (any/c . -> . any/c)
  lst : list?
```

Like `filter`, but the meaning of the *pred* predicate is reversed: the result is a list of all items for which *pred* returns `#f`.

Example:

```
> (filter-not even? '(1 2 3 4 5 6))
'(1 3 5)
```

```
(shuffle lst) → list?
  lst : list?
```

Returns a list with all elements from *lst*, randomly shuffled.

Example:

```
> (shuffle '(1 2 3 4 5 6))
'(3 5 4 6 2 1)
```

```
(permutations lst) → list?
  lst : list?
```

Returns a list of all permutations of the input list. Note that this function works without inspecting the elements, and therefore it ignores repeated elements (which will result in repeated permutations).

Examples:

```
> (permutations '(1 2 3))
'((1 2 3) (2 1 3) (1 3 2) (3 1 2) (2 3 1) (3 2 1))
> (permutations '(x x))
'((x x) (x x))
```

```
(in-permutations lst) → sequence?
  lst : list?
```

Returns a sequence of all permutations of the input list. It is equivalent to `(in-list (permutations l))` but much faster since it builds the permutations one-by-one on each iteration

```
(argmin proc lst) → any/c
  proc : (-> any/c real?)
  lst : (and/c pair? list?)
```

Returns the first element in the list *lst* that minimizes the result of *proc*. Signals an error on an empty list.

Examples:

```
> (argmin car '((3 pears) (1 banana) (2 apples)))
'(1 banana)
> (argmin car '((1 banana) (1 orange)))
'(1 banana)
```

```
(argmax proc lst) → any/c
  proc : (-> any/c real?)
  lst : (and/c pair? list?)
```

Returns the first element in the list *lst* that maximizes the result of *proc*. Signals an error on an empty list.

Examples:

```
> (argmax car '((3 pears) (1 banana) (2 apples)))
'(3 pears)
> (argmax car '((3 pears) (3 oranges)))
'(3 pears)
```

4.9.8 Immutable Cyclic Data

```
(make-reader-graph v) → any/c
  v : any/c
```

Returns a value like `v`, with placeholders created by `make-placeholder` replaced with the values that they contain, and with placeholders created by `make-hash-placeholder` with an immutable hash table. No part of `v` is mutated; instead, parts of `v` are copied as necessary to construct the resulting graph, where at most one copy is created for any given value.

Since the copied values can be immutable, and since the copy is also immutable, `make-reader-graph` can create cycles involving only immutable pairs, vectors, boxes, and hash tables.

Only the following kinds of values are copied and traversed to detect placeholders:

- pairs
- vectors, both mutable and immutable
- boxes, both mutable and immutable
- hash tables, both mutable and immutable
- instances of a prefab structure type
- placeholders created by `make-placeholder` and `make-hash-placeholder`

Due to these restrictions, `make-reader-graph` creates exactly the same sort of cyclic values as `read`.

Example:

```
> (let* ([ph (make-placeholder #f)]
         [x (cons 1 ph)])
    (placeholder-set! ph x)
    (make-reader-graph x))
#0='(1 . #0#)
```

```
(placeholder? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a placeholder created by `make-placeholder`, `#f` otherwise.

```
(make-placeholder v) → placeholder?
  v : any/c
```

Returns a placeholder for use with `placeholder-set!` and `make-reader-graph`. The *v* argument supplies the initial value for the placeholder.

```
(placeholder-set! ph datum) → void?
  ph : placeholder?
  datum : any/c
```

Changes the value of *ph* to `v`.

```
(placeholder-get ph) → any/c
  ph : placeholder?
```

Returns the value of *ph*.

```
(hash-placeholder? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a placeholder created by `make-hash-placeholder`, `#f` otherwise.

```
(make-hash-placeholder assocs) → hash-placeholder?
  assocs : (listof pair?)
```

Like `make-immutable-hash`, but produces a table placeholder for use with `make-reader-graph`.

```
(make-hasheq-placeholder assocs) → hash-placeholder?
  assocs : (listof pair?)
```

Like `make-immutable-hasheq`, but produces a table placeholder for use with `make-reader-graph`.

```
(make-hasheqv-placeholder assocs) → hash-placeholder?
  assocs : (listof pair?)
```

Like `make-immutable-hasheqv`, but produces a table placeholder for use with `make-reader-graph`.

4.10 Mutable Pairs and Lists

A *mutable pair* is like a pair created by `cons`, but it supports `set-mcar!` and `set-mcdr!` mutation operations to change the parts of the mutable pair (like traditional Lisp and Scheme pairs).

A *mutable list* is analogous to a list created with pairs, but instead created with mutable pairs.

A mutable pair is not a pair; they are completely separate datatypes. Similarly, a mutable list is not a list, except that the empty list is also the empty mutable list. Instead of programming with mutable pairs and mutable lists, data structures such as pairs, lists, and hash tables are practically always better choices.

A mutable list can be used as a single-valued sequence (see §4.14.1 “Sequences”). The elements of the mutable list serve as elements of the sequence. See also `in-mlist`.

4.10.1 Mutable Pair Constructors and Selectors

```
(mpair? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a mutable pair, `#f` otherwise.

```
(mcons a d) → mpair?
  a : any/c
  d : any/c
```

Returns a newly allocated mutable pair whose first element is `a` and second element is `d`.

```
(mcar p) → any/c
  p : mpair?
```

Returns the first element of the mutable pair `p`.

```
(mcdr p) → any/c
  p : mpair?
```

Returns the second element of the mutable pair `p`.

```
(set-mcar! p v) → void?
  p : mpair?
  v : any/v
```

Changes the mutable pair `p` so that its first element is `v`.

```
(set-mcdr! p v) → void?
  p : mpair?
  v : any/v
```

Changes the mutable pair `p` so that its second element is `v`.

4.11 Vectors

§3.9 “Vectors” in ***The Racket Guide*** introduces vectors.

A *vector* is a fixed-length array with constant-time access and update of the vector slots, which are numbered from 0 to one less than the number of slots in the vector.

Two vectors are `equal?` if they have the same length, and if the values in corresponding slots of the vectors are `equal?`.

A vector can be *mutable* or *immutable*. When an immutable vector is provided to a procedure like `vector-set!`, the `exn:fail:contract` exception is raised. Vectors generated by the default reader (see §1.3.7 “Reading Strings”) are immutable.

A vector can be used as a single-valued sequence (see §4.14.1 “Sequences”). The elements of the vector serve as elements of the sequence. See also `in-vector`.

A literal or printed vector starts with `#(`, optionally with a number between the `#` and `(`. See §1.3.10 “Reading Vectors” for information on `read`ing vectors and §1.4.7 “Printing Vectors” for information on `print`ing vectors.

```
(vector? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a vector, `#f` otherwise.

```
(make-vector size [v]) → vector?
  size : exact-nonnegative-integer?
  v : any/c = 0
```

Returns a mutable vector with `size` slots, where all slots are initialized to contain `v`.

```
(vector v ...) → vector?
  v : any/c
```

Returns a newly allocated mutable vector with as many slots as provided `v`s, where the slots are initialized to contain the given `v`s in order.

```
(vector-immutable v ...) → (and/c vector?
                                  immutable?)
  v : any/c
```

Returns a newly allocated immutable vector with as many slots as provided `v`s, where the slots are contain the given `v`s in order.

```
(vector-length vec) → exact-nonnegative-integer?
  vec : vector?
```

Returns the length of *vec* (i.e., the number of slots in the vector).

```
(vector-ref vec pos) → any/c
  vec : vector?
  pos : exact-nonnegative-integer?
```

Returns the element in slot *pos* of *vec*. The first slot is position 0, and the last slot is one less than (vector-length *vec*).

```
(vector-set! vec pos v) → void?
  vec : (and/c vector? (not/c immutable?))
  pos : exact-nonnegative-integer?
  v : any/c
```

Updates the slot *pos* of *vec* to contain *v*.

```
(vector->list vec) → list?
  vec : vector?
```

Returns a list with the same length and elements as *vec*.

```
(list->vector lst) → vector?
  lst : list?
```

Returns a mutable vector with the same length and elements as *lst*.

```
(vector->immutable-vector vec) → (and/c vector? immutable?)
  vec : vector?
```

Returns an immutable vector with the same length and elements as *vec*. If *vec* is itself immutable, then it is returned as the result.

```
(vector-fill! vec v) → void?
  vec : (and/c vector? (not/c immutable?))
  v : any/c
```

Changes all slots of *vec* to contain *v*.

```
(vector-copy!  dest
               dest-start
               src
              [src-start
               src-end])   → void?
  dest : (and/c vector? (not/c immutable?))
  dest-start : exact-nonnegative-integer?
  src : vector?
  src-start : exact-nonnegative-integer? = 0
  src-end : exact-nonnegative-integer? = (vector-length src)
```

Changes the elements of *dest* starting at position *dest-start* to match the elements in *src* from *src-start* (inclusive) to *src-end* (exclusive). The vectors *dest* and *src* can be the same vector, and in that case the destination region can overlap with the source region; the destination elements after the copy match the source elements from before the copy. If any of *dest-start*, *src-start*, or *src-end* are out of range (taking into account the sizes of the vectors and the source and destination regions), the `exn:fail:contract` exception is raised.

Examples:

```
> (define v (vector 'A 'p 'p 'l 'e))

> (vector-copy! v 4 #(y))

> (vector-copy! v 0 v 3 4)

> v
'#(l p p l y)
```

```
(vector->values vec [start-pos end-pos]) → any
  vec : vector?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (vector-length vec)
```

Returns *end-pos* - *start-pos* values, which are the elements of *vec* from *start-pos* (inclusive) to *end-pos* (exclusive). If *start-pos* or *end-pos* are greater than `(vector-length vec)`, or if *end-pos* is less than *start-pos*, the `exn:fail:contract` exception is raised.

```
(build-vector n proc) → vector?
  n : exact-nonnegative-integer?
  proc : (exact-nonnegative-integer? . -> . any/c)
```

Creates a vector of *n* elements by applying *proc* to the integers from `0` to `(sub1 n)` in order. If *vec* is the resulting vector, then `(vector-ref vec i)` is the value produced by `(proc i)`.

Example:

```
> (build-vector 5 add1)
'#(1 2 3 4 5)
```

4.11.1 Additional Vector Functions

`(require racket/vector)` package: `base`

The bindings documented in this section are provided by the racket/vector and racket libraries, but not racket/base.

```
(vector-set*! vec pos v ... ...) → void?
  vec : (and/c vector? (not/c immutable?))
  pos : exact-nonnegative-integer?
  v : any/c
```

Updates each slot *pos* of *vec* to contain each *v*. The update takes place from the left so later updates overwrite earlier updates.

```
(vector-map proc vec ...+) → vector?
  proc : procedure?
  vec : vector?
```

Applies *proc* to the elements of the *vec*s from the first elements to the last. The *proc* argument must accept the same number of arguments as the number of supplied *vec*s, and all *vec*s must have the same number of elements. The result is a fresh vector containing each result of *proc* in order.

Example:

```
> (vector-map + #(1 2) #(3 4))
'#(4 6)
```

```
(vector-map! proc vec ...+) → vector?
  proc : procedure?
  vec : (and/c vector? (not/c immutable?))
```

Like vector-map, but result of *proc* is inserted into the first *vec* at the index that the arguments to *proc* were taken from. The result is the first *vec*.

Examples:

```
> (define v (vector 1 2 3 4))

> (vector-map! add1 v)
'#(2 3 4 5)
> v
'#(2 3 4 5)
```

```
(vector-append vec ...) → vector?
  vec : vector?
```

Creates a fresh vector that contains all of the elements of the given vectors in order.

Example:

```
> (vector-append #(1 2) #(3 4))
'#(1 2 3 4)
```

```
(vector-take vec pos) → vector?
  vec : vector?
  pos : exact-nonnegative-integer?
```

Returns a fresh vector whose elements are the first *pos* elements of *vec*. If *vec* has fewer than *pos* elements, then the `exn:fail:contract` exception is raised.

Example:

```
> (vector-take #(1 2 3 4) 2)
'#(1 2)
```

```
(vector-take-right vec pos) → vector?
  vec : vector?
  pos : exact-nonnegative-integer?
```

Returns a fresh vector whose elements are the last *pos* elements of *vec*. If *vec* has fewer than *pos* elements, then the `exn:fail:contract` exception is raised.

Example:

```
> (vector-take-right #(1 2 3 4) 2)
'#(3 4)
```

```
(vector-drop vec pos) → vector?
  vec : vector?
  pos : exact-nonnegative-integer?
```

Returns a fresh vector whose elements are the elements of *vec* after the first *pos* elements. If *vec* has fewer than *pos* elements, then the `exn:fail:contract` exception is raised.

Example:

```
> (vector-drop #(1 2 3 4) 2)
'#(3 4)
```

```
(vector-drop-right vec pos) → vector?
  vec : vector?
  pos : exact-nonnegative-integer?
```

Returns a fresh vector whose elements are the elements of *vec* before the first *pos* elements. If *vec* has fewer than *pos* elements, then the `exn:fail:contract` exception is raised.

Example:

```
> (vector-drop-right #(1 2 3 4) 2)
'#(1 2)
```

```
(vector-split-at vec pos) → vector? vector?
  vec : vector?
  pos : exact-nonnegative-integer?
```

Returns the same result as

```
(values (vector-take vec pos) (vector-drop vec pos))
```

except that it can be faster.

Example:

```
> (vector-split-at #(1 2 3 4 5) 2)
'#(1 2)
'#(3 4 5)
```

```
(vector-split-at-right vec pos) → vector? vector?
  vec : vector?
  pos : exact-nonnegative-integer?
```

Returns the same result as

```
(values (vector-take-right vec pos) (vector-drop-right vec pos))
```

except that it can be faster.

Example:

```
> (vector-split-at-right #(1 2 3 4 5) 2)
'#(1 2 3)
'#(4 5)
```

```
(vector-copy vec [start end]) → vector?
  vec : vector?
  start : exact-nonnegative-integer? = 0
  end : exact-nonnegative-integer? = (vector-length v)
```

Creates a fresh vector of size `(- end start)`, with all of the elements of *vec* from *start* (inclusive) to *end* (exclusive).

Examples:

```
> (vector-copy #(1 2 3 4))
'#(1 2 3 4)
> (vector-copy #(1 2 3 4) 3)
'#(4)
> (vector-copy #(1 2 3 4) 2 3)
'#(3)
```

```
(vector-filter pred vec) → vector?
  pred : procedure?
  vec : vector?
```

Returns a fresh vector with the elements of *vec* for which *pred* produces a true value. The *pred* procedure is applied to each element from first to last.

Example:

```
> (vector-filter even? #(1 2 3 4 5 6))
'#(2 4 6)
```

```
(vector-filter-not pred vec) → vector?
  pred : procedure?
  vec : vector?
```

Like vector-filter, but the meaning of the *pred* predicate is reversed: the result is a vector of all items for which *pred* returns #f.

Example:

```
> (vector-filter-not even? #(1 2 3 4 5 6))
'#(1 3 5)
```

```
(vector-count proc vec ...+) → exact-nonnegative-integer?
  proc : procedure?
  vec : vector?
```

Returns the number of elements of the *vec* ... (taken in parallel) on which *proc* does not evaluate to #f.

Examples:

```
> (vector-count even? #(1 2 3 4 5))
2
> (vector-count = #(1 2 3 4 5) #(5 4 3 2 1))
1
```

```
(vector-argmin proc vec) → any/c
  proc : (-> any/c real?)
  vec : vector?
```

This returns the first element in the non-empty vector *vec* that minimizes the result of *proc*.

Examples:

```
> (vector-argmin car #((3 pears) (1 banana) (2 apples)))
'(1 banana)
> (vector-argmin car #((1 banana) (1 orange)))
'(1 banana)
```

```
(vector-argmax proc vec) → any/c
  proc : (-> any/c real?)
  vec : vector?
```

This returns the first element in the non-empty vector *vec* that maximizes the result of *proc*.

Examples:

```
> (vector-argmax car #((3 pears) (1 banana) (2 apples)))
'(3 pears)
> (vector-argmax car #((3 pears) (3 oranges)))
'(3 pears)
```

```
(vector-member v lst) → (or/c natural-number/c #f)
  v : any/c
  lst : vector?
```

Locates the first element of vec that is equal? to v. If such an element exists, the index of that element in vec is returned. Otherwise, the result is #f.

Examples:

```
> (vector-member 2 (vector 1 2 3 4))
1
> (vector-member 9 (vector 1 2 3 4))
#f
```

```
(vector-memv v vec) → (or/c natural-number/c #f)
  v : any/c
  vec : vector?
```

Like vector-member, but finds an element using eqv?.

Examples:

```
> (vector-memv 2 (vector 1 2 3 4))
1
> (vector-memv 9 (vector 1 2 3 4))
#f
```

```
(vector-memq v vec) → (or/c natural-number/c #f)
  v : any/c
  vec : vector?
```

Like vector-member, but finds an element using eq?.

Examples:

```
> (vector-memq 2 (vector 1 2 3 4))
1
> (vector-memq 9 (vector 1 2 3 4))
#f
```

4.12 Boxes

§3.11 “Boxes” in ***The Racket Guide*** introduces boxes.

A *box* is like a single-element vector, normally used as minimal mutable storage.

A literal or printed box starts with #&. See §1.3.13 “Reading Boxes” for information on reading boxes and §1.4.10 “Printing Boxes” for information on printing boxes.

```
(box? v) → boolean?
  v : any/c
```

Returns #t if v is a box, #f otherwise.

```
(box v) → box?
  v : any/c
```

Returns a new mutable box that contains v.

```
(box-immutable v) → (and/c box? immutable?)
  v : any/c
```

Returns a new immutable box that contains v.

```
(unbox box) → any/c
  box : box?
```

Returns the content of box.

For any v, (unbox (box v)) returns v.

```
(set-box! box v) → void?
  box : (and/c box? (not/c immutable?))
  v : any/c
```

Sets the content of *box* to *v*.

```
(box-cas! box old new) → boolean?
  box : (and/c box? (not/c immutable?) (not/c impersonator?))
  old : any/c
  new : any/c
```

Atomically updates the contents of *box* to *new*, provided that *box* currently contains a value that is `eq?` to *old*, and returns `#t` in that case. If *box* does not contain *old*, then the result is `#f`.

If no other threads or futures attempt to access *box*, the operation is equivalent to

```
(and (eq? old (unbox loc)) (set-box! loc new) #t)
```

When Racket is compiled with support for futures, `box-cas!` uses a hardware *compare and set* operation. Uses of `box-cas!` be performed safely in a future (i.e., allowing the future thunk to continue in parallel).

4.13 Hash Tables

§3.10 “Hash Tables” in ***The Racket Guide*** introduces hash tables.

A *hash table* (or simply *hash*) maps each of its keys to a single value. For a given hash table, keys are equivalent via `equal?`, `eqv?`, or `eq?`, and keys are retained either strongly or weakly (see §16.1 “Weak Boxes”). A hash table is also either mutable or immutable. Immutable hash tables support effectively constant-time access and update, just like mutable hash tables; the constant on immutable operations is usually larger, but the functional nature of immutable hash tables can pay off in certain algorithms.

Immutable hash tables actually provide $O(\log N)$ access and update. Since N is limited by the address space so that $\log N$ is limited to less than 30 or 62 (depending on the platform), $\log N$ can be treated reasonably as a constant.

For `equal?`-based hashing, the built-in hash functions on strings, pairs, lists, vectors, prefab or transparent structures, etc., take time proportional to the size of the value. The hash code for a compound data structure, such as a list or vector, depends on hashing each item of the container, but the depth of such recursive hashing is limited (to avoid potential problems with cyclic data). For a non-list pair, both `car` and `cdr` hashing is treated as a deeper hash, but the `cdr` of a list is treated as having the same hashing depth as the list.

A hash table can be used as a two-valued sequence (see §4.14.1 “Sequences”). The keys and values of the hash table serve as elements of the sequence (i.e., each element is a key and its associated value). If a mapping is added to or removed from the hash table during iteration, then an iteration step may fail with `exn:fail:contract`, or the iteration may skip or duplicate keys and values. See also `in-hash`, `in-hash-keys`, `in-hash-values`, and `in-hash-pairs`.

Two hash tables cannot be `equal?` unless they use the same key-comparison procedure (`equal?`, `eqv?`, or `eq?`), both hold keys strongly or weakly, and have the same mutability.

Caveats concerning concurrent modification: A mutable hash table can be manipulated with `hash-ref`, `hash-set!`, and `hash-remove!` concurrently by multiple threads, and the operations are protected by a table-specific semaphore as needed. Three caveats apply, however:

- If a thread is terminated while applying `hash-ref`, `hash-set!`, `hash-remove!`, `hash-ref!`, or `hash-update!` to a hash table that uses `equal?` or `eqv?` key comparisons, all current and future operations on the hash table may block indefinitely.
- The `hash-map`, `hash-for-each`, and `hash-clear!` procedures do not use the table's semaphore to guard the traversal as a whole. Changes by one thread to a hash table can affect the keys and values seen by another thread part-way through its traversal of the same hash table.
- The `hash-update!` and `hash-ref!` functions use a table's semaphore independently for the `hash-ref` and `hash-set!` parts of their functionality, which means that the update as a whole is not "atomic."

Caveat concerning mutable keys: If a key in an `equal?`-based hash table is mutated (e.g., a key string is modified with `string-set!`), then the hash table's behavior for insertion and lookup operations becomes unpredictable.

A literal or printed hash table starts with `#hash`, `#hasheqv`, or `#hasheq`. See §1.3.12 "Reading Hash Tables" for information on `read`ing hash tables and §1.4.9 "Printing Hash Tables" for information on `print`ing hash tables.

```
(hash? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a hash table, `#f` otherwise.

```
(hash-equal? hash) → boolean?
  hash : hash?
```

Returns `#t` if *hash* compares keys with `equal?`, `#f` if it compares with `eq?` or `eqv?`.

```
(hash-eqv? hash) → boolean?
  hash : hash?
```

Returns `#t` if *hash* compares keys with `eqv?`, `#f` if it compares with `equal?` or `eq?`.

```
(hash-eq? hash) → boolean?
  hash : hash?
```

Returns `#t` if *hash* compares keys with `eq?`, `#f` if it compares with `equal?` or `eqv?`.

```
(hash-weak? hash) → boolean?
  hash : hash?
```

Returns `#t` if *hash* retains its keys weakly, `#f` if it retains keys strongly.

```
(hash key val ... ...) → (and/c hash? hash-equal? immutable?)
  key : any/c
  val : any/c
(hasheq key val ... ...) → (and/c hash? hash-eq? immutable?)
  key : any/c
  val : any/c
(hasheqv key val ... ...) → (and/c hash? hash-eqv? immutable?)
  key : any/c
  val : any/c
```

Creates an immutable hash table with each given *key* mapped to the following *val*; each *key* must have a *val*, so the total number of arguments to `hash` must be even.

The `hash` procedure creates a table where keys are compared with `equal?`, `hasheq` procedure creates a table where keys are compared with `eq?`, and `hasheqv` procedure creates a table where keys are compared with `eqv?`.

The *key* to *val* mappings are added to the table in the order that they appear in the argument list, so later mappings can hide earlier mappings if the *key*s are equal.

```
(make-hash [assocs]) → (and/c hash? hash-equal?)
  assocs : (listof pair?) = null
(make-hasheqv [assocs]) → (and/c hash? hash-eqv?)
  assocs : (listof pair?) = null
(make-hasheq [assocs]) → (and/c hash? hash-eq?)
  assocs : (listof pair?) = null
```

Creates a mutable hash table that holds keys strongly.

The `make-hash` procedure creates a table where keys are compared with `equal?`, `make-hasheq` procedure creates a table where keys are compared with `eq?`, and `make-hasheqv` procedure creates a table where keys are compared with `eqv?`.

The table is initialized with the content of *assocs*. In each element of *assocs*, the `car` is a key, and the `cdr` is the corresponding value. The mappings are added to the table in the order that they appear in *assocs*, so later mappings can hide earlier mappings.

See also `make-custom-hash`.

```
(make-weak-hash [assocs]) → (and/c hash? hash-equal? hash-weak?)
  assocs : (listof pair?) = null
```

```
(make-weak-hasheqv [assocs]) → (and/c hash? hash-eqv? hash-weak?)
  assocs : (listof pair?) = null
(make-weak-hasheq [assocs]) → (and/c hash? hash-eq? hash-weak?)
  assocs : (listof pair?) = null
```

Like make-hash, make-hasheq, and make-hasheqv, but creates a mutable hash table that holds keys weakly.

```
(make-immutable-hash [assocs])
 → (and/c hash? hash-equal? immutable?)
  assocs : (listof pair?) = null
(make-immutable-hasheqv [assocs])
 → (and/c hash? hash-eqv? immutable?)
  assocs : (listof pair?) = null
(make-immutable-hasheq [assocs])
 → (and/c hash? hash-eq? immutable?)
  assocs : (listof pair?) = null
```

Like hash, hasheq, and hasheqv, but accepts the key–value mapping in association-list form like make-hash, make-hasheq, and make-hasheqv.

```
(hash-set! hash key v) → void?
  hash : (and/c hash? (not/c immutable?))
  key : any/c
  v : any/c
```

Maps *key* to *v* in *hash*, overwriting any existing mapping for *key*.

See also the caveats concerning concurrent modification and the caveat concerning mutable keys above.

```
(hash-set*! hash key v ... ...) → void?
  hash : (and/c hash? (not/c immutable?))
  key : any/c
  v : any/c
```

Maps each *key* to each *v* in *hash*, overwriting any existing mapping for each *key*. Mappings are added from the left, so later mappings overwrite earlier mappings.

See also the caveats concerning concurrent modification and the caveat concerning mutable keys above.

```
(hash-set hash key v) → (and/c hash? immutable?)
  hash : (and/c hash? immutable?)
  key : any/c
  v : any/c
```

Functionally extends *hash* by mapping *key* to *v*, overwriting any existing mapping for *key*, and returning the extended hash table.

See also the caveat concerning mutable keys above.

```
(hash-set* hash key v ... ...) → (and/c hash? immutable?)
  hash : (and/c hash? immutable?)
  key : any/c
  v : any/c
```

Functionally extends *hash* by mapping each *key* to *v*, overwriting any existing mapping for each *key*, and returning the extended hash table. Mappings are added from the left, so later mappings overwrite earlier mappings.

See also the caveat concerning mutable keys above.

```
(hash-ref hash key [failure-result]) → any
  hash : hash?
  key : any/c
  failure-result : any/c
                 = (lambda ()
                     (raise (make-exn:fail:contract ....)))
```

Returns the value for *key* in *hash*. If no value is found for *key*, then *failure-result* determines the result:

- If *failure-result* is a procedure, it is called (through a tail call) with no arguments to produce the result.
- Otherwise, *failure-result* is returned as the result.

See also the caveats concerning concurrent modification and the caveat concerning mutable keys above.

```
(hash-ref! hash key to-set) → any
  hash : hash?
  key : any/c
  to-set : any/c
```

Returns the value for *key* in *hash*. If no value is found for *key*, then *to-set* determines the result as in hash-ref (i.e., it is either a thunk that computes a value or a plain value), and this result is stored in *hash* for the *key*. (Note that if *to-set* is a thunk, it is not invoked in tail position.)

See also the caveats concerning concurrent modification and the caveat concerning mutable keys above.

```
(hash-has-key? hash key) → boolean?
  hash : hash?
  key : any/c
```

Returns `#t` if *hash* contains a value for the given *key*, `#f` otherwise.

```
(hash-update! hash
              key
              updater
             [failure-result]) → void?
  hash : (and/c hash? (not/c immutable?))
  key : any/c
  updater : (any/c . -> . any/c)
  failure-result : any/c
                 = (lambda ()
                     (raise (make-exn:fail:contract ....)))
```

Composes `hash-ref` and `hash-set!` to update an existing mapping in *hash*, where the optional *failure-result* argument is used as in `hash-ref` when no mapping exists for *key* already. See the caveat above about concurrent updates.

See also the caveats concerning concurrent modification and the caveat concerning mutable keys above.

```
(hash-update hash key updater [failure-result])
 → (and/c hash? immutable?)
  hash : (and/c hash? immutable?)
  key : any/c
  updater : (any/c . -> . any/c)
  failure-result : any/c
                 = (lambda ()
                     (raise (make-exn:fail:contract ....)))
```

Composes `hash-ref` and `hash-set` to functionally update an existing mapping in *hash*, where the optional *failure-result* argument is used as in `hash-ref` when no mapping exists for *key* already.

See also the caveat concerning mutable keys above.

```
(hash-remove! hash key) → void?
  hash : (and/c hash? (not/c immutable?))
  key : any/c
```

Removes any existing mapping for *key* in *hash*.

See also the caveats concerning concurrent modification and the caveat concerning mutable keys above.

```
(hash-remove hash key) → (and/c hash? immutable?)
  hash : (and/c hash? immutable?)
  key : any/c
```

Functionally removes any existing mapping for *key* in *hash*, returning the fresh hash table.

See also the caveat concerning mutable keys above.

```
(hash-clear! hash) → void?
  hash : (and/c hash? (not/c immutable?))
```

Removes all mappings from *hash*.

If *hash* is not an impersonator, then all mappings are removed in constant time. If *hash* is an impersonator, then each key is removed one-by-one using `hash-remove!`.

See also the caveats concerning concurrent modification and the caveat concerning mutable keys above.

```
(hash-clear hash) → (and/c hash? immutable?)
  hash : (and/c hash? immutable?)
```

Functionally removes all mappings from *hash*.

If *hash* is not a chaperone, then clearing is equivalent to creating a new hash table, and the operation is performed in constant time. If *hash* is a chaperone, then each key is removed one-by-one using `hash-remove`.

```
(hash-copy-clear hash) → hash?
  hash : hash?
```

Produces an empty hash table with the same key-comparison procedure and mutability of *hash*.

```
(hash-map hash proc) → (listof any/c)
  hash : hash?
  proc : (any/c any/c . -> . any/c)
```

Applies the procedure *proc* to each element in *hash* in an unspecified order, accumulating the results into a list. The procedure *proc* is called each time with a key and its value.

If a hash table is extended with new keys (either through *proc* or by another thread) while a `hash-map` or `hash-for-each` traversal is in process, arbitrary key–value pairs can be

dropped or duplicated in the traversal. Key mappings can be deleted or remapped (by any thread) with no adverse affects; the change does not affect a traversal if the key has been seen already, otherwise the traversal skips a deleted key or uses the remapped key's new value.

See also the caveats concerning concurrent modification above.

```
(hash-keys hash) → (listof any/c)
  hash : hash?
```

Returns a list of the keys of *hash* in an unspecified order.

See hash-map for information about modifying *hash* during hash-keys.

See also the caveats concerning concurrent modification above.

```
(hash-values hash) → (listof any/c)
  hash : hash?
```

Returns a list of the values of *hash* in an unspecified order.

See hash-map for information about modifying *hash* during hash-values.

See also the caveats concerning concurrent modification above.

```
(hash->list hash) → (listof (cons/c any/c any/c))
  hash : hash?
```

Returns a list of the key–value pairs of *hash* in an unspecified order.

See hash-map for information about modifying *hash* during hash->list.

See also the caveats concerning concurrent modification above.

```
(hash-for-each hash proc) → void?
  hash : hash?
  proc : (any/c any/c . -> . any)
```

Applies *proc* to each element in *hash* (for the side-effects of *proc*) in an unspecified order. The procedure *proc* is called each time with a key and its value.

See hash-map for information about modifying *hash* within *proc*.

See also the caveats concerning concurrent modification above.

```
(hash-count hash) → exact-nonnegative-integer?
  hash : hash?
```

Returns the number of keys mapped by *hash*. Unless *hash* retains keys weakly, the result is computed in constant time and atomically. If *hash* retains it keys weakly, a traversal is required to count the keys.

```
(hash-empty? hash) → boolean?
  hash : hash?
```

Equivalent to `(zero? (hash-count hash))`.

```
(hash-iterate-first hash)
 → (or/c #f exact-nonnegative-integer?)
  hash : hash?
```

Returns `#f` if *hash* contains no elements, otherwise it returns an integer that is an index to the first element in the hash table; "first" refers to an unspecified ordering of the table elements, and the index values are not necessarily consecutive integers. For a mutable *hash*, this index is guaranteed to refer to the first item only as long as no items are added to or removed from *hash*.

```
(hash-iterate-next hash pos)
 → (or/c #f exact-nonnegative-integer?)
  hash : hash?
  pos : exact-nonnegative-integer?
```

Returns either an integer that is an index to the element in *hash* after the element indexed by *pos* (which is not necessarily one more than *pos*) or `#f` if *pos* refers to the last element in *hash*. If *pos* is not a valid index, then the `exn:fail:contract` exception is raised. For a mutable *hash*, the result index is guaranteed to refer to its item only as long as no items are added to or removed from *hash*.

```
(hash-iterate-key hash pos) → any
  hash : hash?
  pos : exact-nonnegative-integer?
```

Returns the key for the element in *hash* at index *pos*. If *pos* is not a valid index for *hash*, the `exn:fail:contract` exception is raised.

```
(hash-iterate-value hash pos) → any
  hash : hash?
  pos : exact-nonnegative-integer?
```

Returns the value for the element in *hash* at index *pos*. If *pos* is not a valid index for *hash*, the `exn:fail:contract` exception is raised.

```
(hash-copy hash) → (and/c hash? (not/c immutable?))
  hash : hash?
```

Returns a mutable hash table with the same mappings, same key-comparison mode, and same key-holding strength as *hash*.

```
(eq-hash-code v) → fixnum?
  v : any/c
```

Returns a fixnum; for any two calls with `eq?` values, the returned number is the same.

Equal fixnums are always `eq?`.

```
(eqv-hash-code v) → fixnum?
  v : any/c
```

Returns a fixnum; for any two calls with `eqv?` values, the returned number is the same.

```
(equal-hash-code v) → fixnum?
  v : any/c
```

Returns a fixnum; for any two calls with `equal?` values, the returned number is the same. A hash code is computed even when *v* contains a cycle through pairs, vectors, boxes, and/or inspectable structure fields. See also `gen:equal+hash`.

```
(equal-secondary-hash-code v) → fixnum?
  v : any/c
```

Like `equal-hash-code`, but computes a secondary value suitable for use in double hashing.

4.14 Sequences and Streams

Sequences and streams abstract over iteration of elements in a collection. Sequences allow iteration with `for` macros or with sequence operations such as `sequence-map`. Streams are functional sequences that can be used either in a generic way or a stream-specific way. Generators are closely related stateful objects that can be converted to a sequence and vice-versa.

4.14.1 Sequences

§11.1 “Sequence Constructors” in ***The Racket Guide*** introduces sequences.

A *sequence* encapsulates an ordered collection of values. The elements of a sequence can be extracted with one of the `for` syntactic forms, with the procedures returned by `sequence-generate`, or by converting the sequence into a stream.

The sequence datatype overlaps with many other datatypes. Among built-in datatypes, the sequence datatype includes the following:

- exact nonnegative integers (see below)
- strings (see §4.3 “Strings”)
- byte strings (see §4.4 “Byte Strings”)
- lists (see §4.9 “Pairs and Lists”)
- mutable lists (see §4.10 “Mutable Pairs and Lists”)
- vectors (see §4.11 “Vectors”)
- flvectors (see §4.2.3.2 “Flonum Vectors”)
- fxvectors (see §4.2.4.2 “Fixnum Vectors”)
- hash tables (see §4.13 “Hash Tables”)
- dictionaries (see §4.15 “Dictionaries”)
- sets (see §4.16 “Sets”)
- input ports (see §13.1 “Ports”)
- streams (see §4.14.2 “Streams”)

An exact number *k* that is a non-negative integer acts as a sequence similar to `(in-range k)`, except that *k* by itself is not a stream.

Custom sequences can be defined using structure type properties. The easiest method to define a custom sequence is to use the `gen:stream` generic interface. Streams are a suitable abstraction for data structures that are directly iterable. For example, a list is directly iterable with `first` and `rest`. On the other hand, vectors are not directly iterable: iteration has to go through an index. For data structures that are not directly iterable, the *iterator* for the data structure can be defined to be a stream (e.g., a structure containing the index of a vector).

For example, unrolled linked lists (represented as a list of vectors) themselves do not fit the stream abstraction, but have index-based iterators that can be represented as streams:

Examples:

```
> (struct unrolled-list-iterator (idx lst)
    #:methods gen:stream
    [(define (stream-empty? iter)
       (define lst (unrolled-list-iterator-lst iter))
       (or (null? lst)
           (and (>= (unrolled-list-iterator-idx iter)
                    (vector-length (first lst)))
                (null? (rest lst)))))
     (define (stream-first iter)
```

```
        (vector-ref (first (unrolled-list-iterator-lst iter))
                    (unrolled-list-iterator-idx iter)))
      (define (stream-rest iter)
        (define idx (unrolled-list-iterator-idx iter))
        (define lst (unrolled-list-iterator-lst iter))
        (if (>= idx (sub1 (vector-length (first lst))))
            (unrolled-list-iterator 0 (rest lst))
            (unrolled-list-iterator (add1 idx) lst))]))

> (define (make-unrolled-list-iterator ul)
    (unrolled-list-iterator 0 (unrolled-list-lov ul)))

> (struct unrolled-list (lov)
    #:property prop:sequence
    make-unrolled-list-iterator)

> (define ul1 (unrolled-list '(#(cracker biscuit) #(cookie scone))))

> (for/list ([x ul1]) x)
'(cracker biscuit cookie scone)
```

The `prop:sequence` property provides more flexibility in specifying iteration, such as when a pre-processing step is needed to prepare the data for iteration. The `make-do-sequence` function creates a sequence given a thunk that returns procedures to implement a sequence, and the `prop:sequence` property can be associated with a structure type to implement its implicit conversion to a sequence.

For most sequence types, extracting elements from a sequence has no side-effect on the original sequence value; for example, extracting the sequence of elements from a list does not change the list. For other sequence types, each extraction implies a side effect; for example, extracting the sequence of bytes from a port causes the bytes to be read from the port. A sequence's state may either span all uses of the sequence, as for a port, or it may be confined to each distinct time that a sequence is *initiate*d by a `for` form, `sequence->stream`, `sequence-generate`, or `sequence-generate*`. Concretely, the thunk passed to `make-do-sequence` is called to initiate the sequence each time the sequence is used.

Individual elements of a sequence typically correspond to single values, but an element may also correspond to multiple values. For example, a hash table generates two values—a key and its value—for each element in the sequence.

Sequence Predicate and Constructors

```
(sequence? v) → boolean?
  v : any/c
```

Returns `#t` if `v` can be used as a sequence, `#f` otherwise.

```
> (sequence? 42)
#t
> (sequence? '(a b c))
#t
> (sequence? "word")
#t
> (sequence? #\x)
#f
```

```
(in-range end) → stream?
  end : number?
(in-range start end [step]) → stream?
  start : number?
  end : number?
  step : number? = 1
```

Returns a sequence (that is also a stream) whose elements are numbers. The single-argument case `(in-range end)` is equivalent to `(in-range 0 end 1)`. The first number in the sequence is *start*, and each successive element is generated by adding *step* to the previous element. The sequence stops before an element that would be greater or equal to *end* if *step* is non-negative, or less or equal to *end* if *step* is negative.

An `in-range` application can provide better performance for number iteration when it appears directly in a `for` clause.

Example: gaussian sum

```
> (for/sum ([x (in-range 10)]) x)
45
```

Example: sum of even numbers

```
> (for/sum ([x (in-range 0 100 2)]) x)
2450
```

```
(in-naturals [start]) → stream?
  start : exact-nonnegative-integer? = 0
```

Returns an infinite sequence (that is also a stream) of exact integers starting with *start*, where each element is one more than the preceding element.

An `in-naturals` application can provide better performance for integer iteration when it appears directly in a `for` clause.

```
> (for/list ([k (in-naturals)]
             [x (in-range 10)])
    (list k x))
'((0 0) (1 1) (2 2) (3 3) (4 4) (5 5) (6 6) (7 7) (8 8) (9 9))
```

(in-list lst) → stream?
  lst : list?

Returns a sequence (that is also a stream) that is equivalent to using *lst* directly as a sequence.

An in-list application can provide better performance for list iteration when it appears directly in a for clause.

See §4.9 “Pairs and Lists” for information on using lists as sequences.

```
> (for/list ([x (in-list '(3 1 4))])
    `(,x ,(* x x)))
'((3 9) (1 1) (4 16))
```

(in-mlist mlst) → sequence?
  mlst : mlist?

Returns a sequence equivalent to *mlst*.

An in-mlist application can provide better performance for mutable list iteration when it appears directly in a for clause.

See §4.10 “Mutable Pairs and Lists” for information on using mutable lists as sequences.

```
> (for/list ([x (in-mlist (mcons "RACKET" (mcons "LANG" '())))])
    (string-length x))
'(6 4)
```

(in-vector vec [start stop step]) → sequence?
  vec : vector?
  start : exact-nonnegative-integer? = 0
  stop : (or/c exact-integer? #f) = #f
  step : (and/c exact-integer? (not/c zero?)) = 1

Returns a sequence equivalent to *vec* when no optional arguments are supplied.

The optional arguments *start*, *stop*, and *step* are analogous to in-range, except that a #f value for *stop* is equivalent to (vector-length *vec*). That is, the first element in the sequence is (vector-ref *vec* *start*), and each successive element is generated by adding *step* to index of the previous element. The sequence stops before an index that would be greater or equal to end if *step* is non-negative, or less or equal to end if *step* is negative.

See §4.11 “Vectors” for information on using vectors as sequences.

If *start* is not a valid index, or *stop* is not in [-1, (vector-length *vec*)] then the exn:fail:contract exception is raised. If *start* is less than *stop* and *step* is negative, then the exn:fail:contract:mismatch exception is raised. Similarly, if *start* is more than *stop* and *step* is positive, then the exn:fail:contract:mismatch exception is raised.

An in-vector application can provide better performance for vector iteration when it appears directly in a for clause.

```
> (define (histogram vector-of-words)
    (define a-hash (hash))
    (for ([word (in-vector vector-of-words)])
      (hash-set! a-hash word (add1 (hash-ref a-hash word 0))))
    a-hash)

> (histogram #("hello" "world" "hello" "sunshine"))
hash-set!: contract violation
  expected: (and/c hash? (not/c immutable?))
  given: '#hash()
  argument position: 1st
  other arguments...:
   "hello"
   1
```

```
(in-string str [start stop step]) → sequence?
  str : string?
  start : exact-nonnegative-integer? = 0
  stop : (or/c exact-integer? #f) = #f
  step : (and/c exact-integer? (not/c zero?)) = 1
```

Returns a sequence equivalent to *str* when no optional arguments are supplied.

See §4.3 “Strings” for information on using strings as sequences.

The optional arguments *start*, *stop*, and *step* are as in in-vector.

An in-string application can provide better performance for string iteration when it appears directly in a for clause.

```
> (define (line-count str)
    (for/sum ([ch (in-string str)])
      (if (char=? #\newline ch) 1 0)))

> (line-count "this string\nhas\nthree \nnewlines")
3
```

```
(in-bytes bstr [start stop step]) → sequence?
  bstr : bytes?
```

```
  start : exact-nonnegative-integer? = 0
  stop : (or/c exact-integer? #f) = #f
  step : (and/c exact-integer? (not/c zero?)) = 1
```

Returns a sequence equivalent to *bstr* when no optional arguments are supplied.

See §4.4 “Byte Strings” for information on using byte strings as sequences.

The optional arguments *start*, *stop*, and *step* are as in `in-vector`.

An `in-bytes` application can provide better performance for byte string iteration when it appears directly in a `for` clause.

```
> (define (has-eof? bs)
    (for/or ([ch (in-bytes bs)])
      (= ch 0)))

> (has-eof? #"this byte string has an \0embedded zero byte")
#t
> (has-eof? #"this byte string does not")
#f
```

```
(in-port [r in]) → sequence?
  r : (input-port? . -> . any/c) = read
  in : input-port? = (current-input-port)
```

Returns a sequence whose elements are produced by calling *r* on *in* until it produces `eof`.

```
(in-input-port-bytes in) → sequence?
  in : input-port?
```

Returns a sequence equivalent to `(in-port read-byte in)`.

```
(in-input-port-chars in) → sequence?
  in : input-port?
```

Returns a sequence whose elements are read as characters from *in* (equivalent to `(in-port read-char in)`).

```
(in-lines [in mode]) → sequence?
  in : input-port? = (current-input-port)
  mode : (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
       = 'any
```

Returns a sequence equivalent to `(in-port (lambda (p) (read-line p mode)) in)`. Note that the default mode is `'any`, whereas the default mode of `read-line` is `'linefeed`.

```
(in-bytes-lines [in mode]) → sequence?
  in : input-port? = (current-input-port)
  mode : (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
       = 'any
```

Returns a sequence equivalent to `(in-port (lambda (p) (read-bytes-line p mode)) in)`. Note that the default mode is `'any`, whereas the default mode of `read-bytes-line` is `'linefeed`.

```
(in-hash hash) → sequence?
  hash : hash?
```

Returns a sequence equivalent to *hash*.

Examples:

```
> (define table (hash 'a 1 'b 2))

> (for ([(key value) (in-hash table)])
    (printf "key: ~a value: ~a\n" key value))
key: a value: 1
key: b value: 2
```

```
(in-hash-keys hash) → sequence?
  hash : hash?
```

See §4.13 “Hash Tables” for information on using hash tables as sequences.

Returns a sequence whose elements are the keys of *hash*.

Examples:

```
> (define table (hash 'a 1 'b 2))

> (for ([key (in-hash-keys table)])
    (printf "key: ~a\n" key))
key: a
key: b
```

```
(in-hash-values hash) → sequence?
  hash : hash?
```

Returns a sequence whose elements are the values of *hash*.

Examples:

```
> (define table (hash 'a 1 'b 2))

> (for ([value (in-hash-values table)])
    (printf "value: ~a\n" value))
value: 1
value: 2
```

```
(in-hash-pairs hash) → sequence?
  hash : hash?
```

Returns a sequence whose elements are pairs, each containing a key and its value from *hash* (as opposed to using *hash* directly as a sequence to get the key and value as separate values for each element).

Examples:

```
> (define table (hash 'a 1 'b 2))

> (for ([key+value (in-hash-pairs table)])
    (printf "key and value: ~a\n" key+value))
key and value: (a . 1)
key and value: (b . 2)
```

```
(in-directory [dir use-dir?]) → sequence?
  dir : (or/c #f path-string?) = #f
  use-dir? : ((and/c path? complete-path?) . -> . any/c)
           = (lambda (dir-path) #t)
```

Returns a sequence that produces all of the paths for files, directories, and links within *dir*, except for the contents of any directory for which *use-dir?* returns `#f`. If *dir* is not `#f`, then every produced path starts with *dir* as its prefix. If *dir* is `#f`, then paths in and relative to the current directory are produced.

An `in-directory` sequence traverses nested subdirectories recursively (filtered by *use-dir?*). To generate a sequence that includes only the immediate content of a directory, use the result of `directory-list` as a sequence.

Changed in version 6.0.0.1 of package `base`: Added *use-dir?* argument.

```
(in-producer producer) → sequence?
  producer : procedure?
(in-producer producer stop arg ...) → sequence?
  producer : procedure?
  stop : any/c
  arg : any/c
```

Returns a sequence that contains values from sequential calls to *producer*, which would usually use some state to do its work.

If a *stop* value is not given, the sequence goes on infinitely, and therefore it common to use it with a finite sequence or using #:break etc. If a *stop* value is given, it is used to identify a value that marks the end of the sequence (and the *stop* value is not included in the sequence); *stop* can be a predicate that is applied to the results of *producer*, or it can be a value that is tested against the result of with eq?. (The *stop* argument must be a predicate if the stop value is itself a function or if *producer* returns multiple values.)

If additional *arg*s are specified, they are passed to every call to *producer*.

Examples:

```
> (define (counter)
    (define n 0)
    (lambda ([d 1]) (set! n (+ d n)) n))

> (for/list ([x (in-producer (counter))] [y (in-range 4)]) x)
'(1 2 3 4)
> (for/list ([x (in-producer (counter))] #:break (= x 5)) x)
'(1 2 3 4)
> (for/list ([x (in-producer (counter) 5)]) x)
'(1 2 3 4)
> (for/list ([x (in-producer (counter) 5 1/2)]) x)
'(1/2 1 3/2 2 5/2 3 7/2 4 9/2)
> (for/list ([x (in-producer read eof (open-input-string "1 2
3"))]) x)
'(1 2 3)
```

```
(in-value v) → sequence?
  v : any/c
```

Returns a sequence that produces a single value: *v*. This form is mostly useful for `let`-like bindings in forms such as `for*/list`.

```
(in-indexed seq) → sequence?
  seq : sequence?
```

Returns a sequence where each element has two values: the value produced by *seq*, and a non-negative exact integer starting with 0. The elements of *seq* must be single-valued.

```
(in-sequences seq ...) → sequence?
  seq : sequence?
```

Returns a sequence that is made of all input sequences, one after the other. Each *seq* is initiated only after the preceding *seq* is exhausted. If a single *seq* is provided, then *seq* is returned; otherwise, the elements of each *seq* must all have the same number of values.

```
(in-cycle seq ...) → sequence?
  seq : sequence?
```

Similar to in-sequences, but the sequences are repeated in an infinite cycle, where each *seq* is initiated afresh in each iteration. Beware that if no *seq*s are provided or if all *seq*s become empty, then the sequence produced by in-cycle never returns when an element is demanded—or even when the sequence is initiated, if all *seq*s are initially empty.

```
(in-parallel seq ...) → sequence?
  seq : sequence?
```

Returns a sequence where each element has as many values as the number of supplied *seq*s; the values, in order, are the values of each *seq*. The elements of each *seq* must be single-valued.

```
(in-values-sequence seq) → sequence?
  seq : sequence?
```

Returns a sequence that is like *seq*, but it combines multiple values for each element from *seq* as a list of elements.

```
(in-values*-sequence seq) → sequence?
  seq : sequence?
```

Returns a sequence that is like *seq*, but when an element of *seq* has multiple values or a single list value, then the values are combined in a list. In other words, in-values*-sequence is like in-values-sequence, except that non-list, single-valued elements are not wrapped in a list.

```
(stop-before seq pred) → sequence?
  seq : sequence?
  pred : (any/c . -> . any)
```

Returns a sequence that contains the elements of *seq* (which must be single-valued), but only until the last element for which applying *pred* to the element produces #t, after which the sequence ends.

```
(stop-after seq pred) → sequence?
  seq : sequence?
  pred : (any/c . -> . any)
```

Returns a sequence that contains the elements of *seq* (which must be single-valued), but only until the element (inclusive) for which applying *pred* to the element produces #t, after which the sequence ends.

```
(make-do-sequence thunk) → sequence?
          (-> (values (any/c . -> . any)
                      (any/c . -> . any/c)
                      any/c
  thunk :             (or/c (any/c . -> . any/c) #f)
                      (or/c (() () #:rest list? . ->* . any/c) #f)
                      (or/c ((any/c) () #:rest list? . ->* . any/c) #f)))
```

Returns a sequence whose elements are generated by the procedures and initial value returned by the thunk, which is called to initiate the sequence. The initiated sequence is defined in terms of a *position*, which is initialized to the third result of the thunk, and the *element*, which may consist of multiple values.

The *thunk* results define the generated elements as follows:

- The first result is a *pos->element* procedure that takes the current position and returns the value(s) for the current element.
- The second result is a *next-pos* procedure that takes the current position and returns the next position.
- The third result is the initial position.
- The fourth result is a *continue-with-pos?* function that takes the current position and returns a true result if the sequence includes the value(s) for the current position, and false if the sequence should end instead of including the value(s). Alternatively, the fourth result can be `#f` to indicate that the sequence should always include the current value(s).
- The fifth result is a *continue-with-val?* function that is like the fourth result, but it takes the current element value(s) instead of the current position. Alternatively, the fifth result can be `#f` to indicate that the sequence should always include the value(s) at the current position.
- The sixth result is a *continue-after-pos+val?* procedure that takes both the current position and the current element value(s) and determines whether the sequence ends after the current element is already included in the sequence. Alternatively, the sixth result can be `#f` to indicate that the sequence can always continue after the current value(s).

Each of the procedures listed above is called only once per position. Among the last three procedures, as soon as one of the procedures returns `#f`, the sequence ends, and none are called again. Typically, one of the functions determines the end condition, and `#f` is used in place of the other two functions.

```
prop:sequence : struct-type-property?
```

Associates a procedure to a structure type that takes an instance of the structure and returns a sequence. If `v` is an instance of a structure type with this property, then `(sequence? v)` produces `#t`.

Using a pre-existing sequence:

Examples:

```
> (struct my-set (table)
    #:property prop:sequence
    (lambda (s)
      (in-hash-keys (my-set-table s))))

> (define (make-set . xs)
    (my-set (for/hash ([x (in-list xs)])
              (values x #t))))

> (for/list ([c (make-set 'celeriac 'carrot 'potato)])
    c)
'(celeriac carrot potato)
```

Using `make-do-sequence`:

Examples:

```
> (define-struct train (car next)
    #:property prop:sequence
    (lambda (t)
      (make-do-sequence
       (lambda ()
         (values train-car train-next t
                 (lambda (t) t)
                 (lambda (v) #t)
                 (lambda (t v) #t))))))

> (for/list ([c (make-train 'engine
                            (make-train 'boxcar
                                        (make-train 'caboose
                                                    #f)))])
    c)
'(engine boxcar caboose)
```

Sequence Conversion

```
(sequence->stream seq) → stream?
  seq : sequence?
```

Coverts a sequence to a stream, which supports the `stream-first` and `stream-rest` operations. Creation of the stream eagerly initiates the sequence, but the stream lazily draws

elements from the sequence, caching each element so that `stream-first` produces the same result each time is applied to a stream.

If extracting an element from *seq* involves a side-effect, then the effect is performed each time that either `stream-first` or `stream-rest` is first used to access or skip an element.

```
(sequence-generate seq) → (-> boolean?) (-> any)
  seq : sequence?
```

Initiates a sequence and returns two thunks to extract elements from the sequence. The first returns `#t` if more values are available for the sequence. The second returns the next element (which may be multiple values) from the sequence; if no more elements are available, the `exn:fail:contract` exception is raised.

```
(sequence-generate* seq)
  →  (or/c list? #f)
     (-> (values (or/c list? #f) procedure?))
  seq : sequence?
```

Like `sequence-generate`, but avoids state (aside from any inherent in the sequence) by returning a list of values for the sequence's first element—or `#f` if the sequence is empty—and a thunk to continue with the sequence; the result of the thunk is the same as the result of `sequence-generate*`, but for the second element of the sequence, and so on. If the thunk is called when the element result is `#f` (indicating no further values in the sequence), the `exn:fail:contract` exception is raised.

Sequence Combinations

```
(require racket/sequence)
```
package: base

The bindings documented in this section are provided by the `racket/sequence` and `racket` libraries, but not `racket/base`.

```
empty-sequence : sequence?
```

A sequence with no elements.

```
(sequence->list s) → list?
  s : sequence?
```

Returns a list whose elements are the elements of *s*, each of which must be a single value. If *s* is infinite, this function does not terminate.

```
(sequence-length s) → exact-nonnegative-integer?
  s : sequence?
```

Returns the number of elements of `s` by extracting and discarding all of them. If `s` is infinite, this function does not terminate.

```
(sequence-ref s i) → any
  s : sequence?
  i : exact-nonnegative-integer?
```

Returns the `i` th element of `s` (which may be multiple values).

```
(sequence-tail s i) → sequence?
  s : sequence?
  i : exact-nonnegative-integer?
```

Returns a sequence equivalent to `s`, except that the first `i` elements are omitted.

In case initiating `s` involves a side effect, the sequence `s` is not initiated until the resulting sequence is initiated, at which point the first `i` elements are extracted from the sequence.

```
(sequence-append s ...) → sequence?
  s : sequence?
```

Returns a sequence that contains all elements of each sequence in the order they appear in the original sequences. The new sequence is constructed lazily.

If all given `s`s are streams, the result is also a stream.

```
(sequence-map f s) → sequence?
  f : procedure?
  s : sequence?
```

Returns a sequence that contains `f` applied to each element of `s`. The new sequence is constructed lazily.

If `s` is a stream, then the result is also a stream.

```
(sequence-andmap f s) → boolean?
  f : (-> any/c ... boolean?)
  s : sequence?
```

Returns `#t` if `f` returns a true result on every element of `s`. If `s` is infinite and `f` never returns a false result, this function does not terminate.

```
(sequence-ormap f s) → boolean?
  f : (-> any/c ... boolean?)
  s : sequence?
```

Returns `#t` if `f` returns a true result on some element of `s`. If `s` is infinite and `f` never returns a true result, this function does not terminate.

```
(sequence-for-each f s) → void?
  f : (-> any/c ... any)
  s : sequence?
```

Applies `f` to each element of `s`. If `s` is infinite, this function does not terminate.

```
(sequence-fold f i s) → any/c
  f : (-> any/c any/c ... any/c)
  i : any/c
  s : sequence?
```

Folds `f` over each element of `s` with `i` as the initial accumulator. If `s` is infinite, this function does not terminate. The `f` function takes the accumulator as its first argument and the next sequence element as its second.

```
(sequence-count f s) → exact-nonnegative-integer?
  f : procedure?
  s : sequence?
```

Returns the number of elements in `s` for which `f` returns a true result. If `s` is infinite, this function does not terminate.

```
(sequence-filter f s) → sequence?
  f : (-> any/c ... boolean?)
  s : sequence?
```

Returns a sequence whose elements are the elements of `s` for which `f` returns a true result. Although the new sequence is constructed lazily, if `s` has an infinite number of elements where `f` returns a false result in between two elements where `f` returns a true result, then operations on this sequence will not terminate during the infinite sub-sequence.

If `s` is a stream, then the result is also a stream.

```
(sequence-add-between s e) → sequence?
  s : sequence?
  e : any/c
```

Returns a sequence whose elements are the elements of `s`, but with `e` between each pair of elements in `s`. The new sequence is constructed lazily.

If `s` is a stream, then the result is also a stream.

Examples:

```
> (let* ([all-reds (in-cycle '("red"))]
         [red-and-blues (sequence-add-between all-reds "blue")])
    (for/list ([n (in-range 10)]
               [elt red-and-blues])
      elt))
'("red" "blue" "red" "blue" "red" "blue" "red" "blue" "red"
"blue")
> (for ([text (sequence-add-between '("veni" "vidi" "duci") ",
")])
    (display text))
veni, vidi, duci
```

4.14.2 Streams

A *stream* is a kind of sequence that supports functional iteration via `stream-first` and `stream-rest`. The `stream-cons` form constructs a lazy stream, but plain lists can be used as streams, and functions such as `in-range` and `in-naturals` also create streams.

`(require racket/stream)` package: `base`

The bindings documented in this section are provided by the `racket/stream` and `racket` libraries, but not `racket/base`.

```
(stream? v) → boolean?
  v : any/c
```

Returns `#t` if *v* can be used as a stream, `#f` otherwise.

```
(stream-empty? s) → boolean?
  s : stream?
```

Returns `#f` if *s* has no elements, `#f` otherwise.

```
(stream-first s) → any
  s : (and/c stream? (not/c stream-empty?))
```

Returns the value(s) of the first element in *s*.

```
(stream-rest s) → stream?
  s : (and/c stream? (not/c stream-empty?))
```

Returns a stream that is equivalent to *s* without its first element.

```
(stream-cons first-expr rest-expr)
```

Produces a lazy stream for which `stream-first` forces the evaluation of *first-expr* to produce the first element of the stream, and `stream-rest` forces the evaluation of *rest-expr* to produce a stream for the rest of the returned stream.

The first element of the stream as produced by *first-expr* must be a single value. The *rest-expr* must produce a stream when it is evaluated, otherwise the `exn:fail:contract?` exception is raised.

```
(stream expr ...)
```

A shorthand for nested `stream-cons`es ending with `empty-stream`.

```
(in-stream s) → sequence?
  s : stream?
```

Returns a sequence that is equivalent to *s*.

An `in-stream` application can provide better performance for streams iteration when it appears directly in a `for` clause.

```
empty-stream : stream?
```

A stream with no elements.

```
(stream->list s) → list?
  s : stream?
```

Returns a list whose elements are the elements of *s*, each of which must be a single value. If *s* is infinite, this function does not terminate.

```
(stream-length s) → exact-nonnegative-integer?
  s : stream?
```

Returns the number of elements of *s*. If *s* is infinite, this function does not terminate.

In the case of lazy streams, this function forces evaluation only of the sub-streams, and not the stream's elements.

```
(stream-ref s i) → any
  s : stream?
  i : exact-nonnegative-integer?
```

Returns the *i* th element of *s* (which may be multiple values).

```
(stream-tail s i) → stream?
  s : stream?
  i : exact-nonnegative-integer?
```

Returns a stream equivalent to `s`, except that the first `i` elements are omitted.

In case extracting elements from `s` involves a side effect, they will not be extracted until the first element is extracted from the resulting stream.

```
(stream-append s ...) → stream?
  s : stream?
```

Returns a stream that contains all elements of each stream in the order they appear in the original streams. The new stream is constructed lazily, while the last given stream is used in the tail of the result.

```
(stream-map f s) → stream?
  f : procedure?
  s : stream?
```

Returns a stream that contains `f` applied to each element of `s`. The new stream is constructed lazily.

```
(stream-andmap f s) → boolean?
  f : (-> any/c ... boolean?)
  s : stream?
```

Returns `#t` if `f` returns a true result on every element of `s`. If `s` is infinite and `f` never returns a false result, this function does not terminate.

```
(stream-ormap f s) → boolean?
  f : (-> any/c ... boolean?)
  s : stream?
```

Returns `#t` if `f` returns a true result on some element of `s`. If `s` is infinite and `f` never returns a true result, this function does not terminate.

```
(stream-for-each f s) → void?
  f : (-> any/c ... any)
  s : stream?
```

Applies `f` to each element of `s`. If `s` is infinite, this function does not terminate.

```
(stream-fold f i s) → any/c
  f : (-> any/c any/c ... any/c)
  i : any/c
  s : stream?
```

Folds `f` over each element of `s` with `i` as the initial accumulator. If `s` is infinite, this function does not terminate.

```
(stream-count f s) → exact-nonnegative-integer?
  f : procedure?
  s : stream?
```

Returns the number of elements in `s` for which `f` returns a true result. If `s` is infinite, this function does not terminate.

```
(stream-filter f s) → stream?
  f : (-> any/c ... boolean?)
  s : stream?
```

Returns a stream whose elements are the elements of `s` for which `f` returns a true result. Although the new stream is constructed lazily, if `s` has an infinite number of elements where `f` returns a false result in between two elements where `f` returns a true result, then operations on this stream will not terminate during the infinite sub-stream.

```
(stream-add-between s e) → stream?
  s : stream?
  e : any/c
```

Returns a stream whose elements are the elements of `s`, but with `e` between each pair of elements in `s`. The new stream is constructed lazily.

```
gen:stream : any/c
```

Associates three methods to a structure type to implement the generic interface (see §5.4 “Generic Interfaces”) for streams.

To supply method implementations, the `#:methods` keyword should be used in a structure type definition. The following three methods should be implemented:

- `stream-empty?` : accepts one argument
- `stream-first` : accepts one argument
- `stream-rest` : accepts one argument

Examples:

```
> (define-struct list-stream (v)
    #:methods gen:stream
    [(define (stream-empty? stream)
       (empty? (list-stream-v stream)))
     (define (stream-first stream)
       (first (list-stream-v stream)))
     (define (stream-rest stream)
       (rest (list-stream-v stream)))])
```

```
> (define l1 (list-stream '(1 2)))

> (stream? l1)
#t
> (stream-first l1)
1
```

```
prop:stream : struct-type-property?
```

A deprecated structure type property used to define custom extensions to the stream API. Use `gen:stream` instead. Accepts a vector of three procedures taking the same arguments as the methods in `gen:stream`.

4.14.3 Generators

A *generator* is a procedure that returns a sequence of values, incrementing the sequence each time that the generator is called. In particular, the `generator` form implements a generator by evaluating a body that calls `yield` to return values from the generator.

```
(require racket/generator)
```
package: `base`

```
(generator? v) → boolean?
  v : any/c
```

Return `#t` if `v` is a generator, `#f` otherwise.

```
(generator formals body ...+)
```

Creates a generator, where *`formals`* is like the *`formals`* of `case-lambda` (i.e., the *`kw-formals`* of `lambda` restricted to non-optional and non-keyword arguments).

For the first call to a generator, the arguments are bound to the *`formals`* and evaluation of *`body`* starts. During the dynamic extent of *`body`*, the generator can return immediately using the `yield` function. A second call to the generator resumes at the `yield` call, producing the arguments of the second call as the results of the `yield`, and so on. The eventual results of *`body`* are supplied to an implicit final `yield`; after that final `yield`, calling the generator again returns the same values, but all such calls must provide 0 arguments to the generator.

Examples:

```
> (define g (generator ()
              (let loop ([x '(a b c)])
                (if (null? x)
```

```
                    0
                    (begin
                      (yield (car x))
                      (loop (cdr x)))))))

> (g)
'a
> (g)
'b
> (g)
'c
> (g)
0
> (g)
0
```

(yield v ...) → any
  v : any/c

Returns vs from a generator, saving the point of execution inside a generator (i.e., within the dynamic extent of a generator body) to be resumed by the next call to the generator. The results of yield are the arguments that are provided to the next call of the generator.

When not in the dynamic extent of a generator, infinite-generator, or in-generator body, yield raises exn:fail after evaluating its exprs.

Examples:

```
> (define my-generator (generator () (yield 1) (yield 2 3 4)))

> (my-generator)
1
> (my-generator)
2
3
4
```

Examples:

```
> (define pass-values-generator
    (generator ()
      (let* ([from-user (yield 2)]
             [from-user-again (yield (add1 from-user))])
        (yield from-user-again))))

> (pass-values-generator)
2
```

```
> (pass-values-generator 5)
6
> (pass-values-generator 12)
12
```

```
(infinite-generator body ...+)
```

Like `generator`, but repeats evaluation of the *body*s when the last *body* completes without implicitly `yield`ing.

Examples:

```
> (define welcome
    (infinite-generator
      (yield 'hello)
      (yield 'goodbye)))

> (welcome)
'hello
> (welcome)
'goodbye
> (welcome)
'hello
> (welcome)
'goodbye
```

```
(in-generator maybe-arity body ...+)

maybe-arity =
            | #:arity arity-k
```

Produces a sequence that encapsulates the generator formed by `(generator () body ...+)`. The values produced by the generator form the elements of the sequence, except for the last value produced by the generator (i.e., the values produced by returning).

Example:

```
> (for/list ([i (in-generator
                 (let loop ([x '(a b c)])
                   (when (not (null? x))
                     (yield (car x))
                     (loop (cdr x)))))])
    i)
'(a b c)
```

If `in-generator` is used immediately with a `for` (or `for/list`, etc.) binding's right-hand side, then its result arity (i.e., the number of values in each element of the sequence) can

be inferred. Otherwise, if the generator produces multiple values for each element, its arity should be declared with an `#:arity arity-k` clause; the `arity-k` must be a literal, exact, non-negative integer.

Examples:

```
> (let ([g (in-generator
            (let loop ([n 3])
              (unless (zero? n) (yield n (add1 n)) (loop (sub1 n)))))])
    (let-values ([(not-empty? next) (sequence-generate g)])
      (let loop () (when (not-empty?) (next) (loop))) 'done))
stop?: arity mismatch;
 the expected number of arguments does not match the given
number
  expected: 1
  given: 2
  arguments...:
   3
   4
> (let ([g (in-generator #:arity 2
            (let loop ([n 3])
              (unless (zero? n) (yield n (add1 n)) (loop (sub1 n)))))])
    (let-values ([(not-empty? next) (sequence-generate g)])
      (let loop () (when (not-empty?) (next) (loop))) 'done))
'done
```

To use an existing generator as a sequence, use `in-producer` with a stop-value known for the generator:

```
> (define abc-generator (generator ()
                          (for ([x '(a b c)])
                             (yield x))))

> (for/list ([i (in-producer abc-generator (void))])
    i)
'(a b c)
> (define my-stop-value (gensym))

> (define my-generator (generator ()
                         (let loop ([x (list 'a (void) 'c)])
                           (if (null? x)
                               my-stop-value
                               (begin
                                 (yield (car x))
                                 (loop (cdr x)))))))

> (for/list ([i (in-producer my-generator my-stop-value)])
```

```
    i)
'(a #<void> c)
```

```
(generator-state g) → symbol?
  g : generator?
```

Returns a symbol that describes the state of the generator.

- `'fresh` — The generator has been freshly created and has not been called yet.
- `'suspended` — Control within the generator has been suspended due to a call to `yield`. The generator can be called.
- `'running` — The generator is currently executing.
- `'done` — The generator has executed its entire body and will continue to produce the same result as from the last call.

Examples:

```
> (define my-generator (generator () (yield 1) (yield 2)))

> (generator-state my-generator)
'fresh
> (my-generator)
1
> (generator-state my-generator)
'suspended
> (my-generator)
2
> (generator-state my-generator)
'suspended
> (my-generator)

> (generator-state my-generator)
'done
> (define introspective-generator (generator () ((yield 1))))

> (introspective-generator)
1
> (introspective-generator
   (lambda () (generator-state introspective-generator)))
'running
> (generator-state introspective-generator)
'done
> (introspective-generator)
'running
```

```
(sequence->generator s) → (-> any)
  s : sequence?
```

Converts a sequence to a generator. The generator returns the next element of the sequence each time the generator is invoked, where each element of the sequence must be a single value. When the sequence ends, the generator returns `#<void>` as its final result.

```
(sequence->repeated-generator s) → (-> any)
  s : sequence?
```

Like `sequence->generator`, but when `s` has no further values, the generator starts the sequence again (so that the generator never stops producing values).

4.15 Dictionaries

A *dictionary* is an instance of a datatype that maps keys to values. The following datatypes are all dictionaries:

- hash tables;
- vectors (using only exact integers as keys);
- lists of pairs (an *association list* using `equal?` to compare keys); and
- structures whose types implement the `gen:dict` generic interface.

```
(require racket/dict)
```
package: `base`

The bindings documented in this section are provided by the `racket/dict` and `racket` libraries, but not `racket/base`.

4.15.1 Dictionary Predicates and Contracts

```
(dict? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a dictionary, `#f` otherwise.

Beware that `dict?` is not a constant-time test on pairs, since checking that `v` is an association list may require traversing the list.

Examples:

```
> (dict? #hash((a . "apple")))
#t
> (dict? '#("apple" "banana"))
#t
> (dict? '("apple" "banana"))
#f
> (dict? '((a . "apple") (b . "banana")))
#t
```

```
(dict-implements? d sym ...) → boolean?
  d : dict?
  sym : symbol?
```

Returns `#t` if *d* implements all of the methods from `gen:dict` named by the *sym*s; returns `#f` otherwise. Fallback implementations do not affect the result; *d* may support the given methods via fallback implementations yet produce `#f`.

Examples:

```
> (dict-implements? (hash 'a "apple") 'dict-set!)
#f
> (dict-implements? (make-hash '((a . "apple") (b . "ba-
nana"))) 'dict-set!)
#t
> (dict-implements? (make-hash '((b . "banana") (a . "ap-
ple"))) 'dict-remove!)
#t
> (dict-implements? (vector "apple" "banana") 'dict-set!)
#t
> (dict-implements? (vector 'a 'b) 'dict-remove!)
#f
> (dict-implements? (vector 'a "apple") 'dict-set! 'dict-remove!)
#f
```

```
(dict-implements/c sym ...) → flat-contract?
  sym : symbol?
```

Recognizes dictionaries that support all of the methods from `gen:dict` named by the *sym*s.

```
(dict-mutable? d) → boolean?
  d : dict?
```

Returns `#t` if *d* is mutable via `dict-set!`, `#f` otherwise.

Equivalent to `(dict-implements? d 'dict-set!)`.

Examples:

```
> (dict-mutable? #hash((a . "apple")))
#f
> (dict-mutable? (make-hash))
#t
> (dict-mutable? '#("apple" "banana"))
#f
> (dict-mutable? (vector "apple" "banana"))
#t
> (dict-mutable? '((a . "apple") (b . "banana")))
#f
```

```
(dict-can-remove-keys? d) → boolean?
  d : dict?
```

Returns `#t` if *d* supports removing mappings via `dict-remove!` and/or `dict-remove`, `#f` otherwise.

Equivalent to `(or (dict-implements? d 'dict-remove!) (dict-implements? d 'dict-remove))`.

Examples:

```
> (dict-can-remove-keys? #hash((a . "apple")))
#t
> (dict-can-remove-keys? '#("apple" "banana"))
#f
> (dict-can-remove-keys? '((a . "apple") (b . "banana")))
#t
```

```
(dict-can-functional-set? d) → boolean?
  d : dict?
```

Returns `#t` if *d* supports functional update via `dict-set`, `#f` otherwise.

Equivalent to `(dict-implements? d 'dict-set)`.

Examples:

```
> (dict-can-functional-set? #hash((a . "apple")))
#t
> (dict-can-functional-set? (make-hash))
#f
> (dict-can-functional-set? '#("apple" "banana"))
#f
> (dict-can-functional-set? '((a . "apple") (b . "banana")))
#t
```

4.15.2 Generic Dictionary Interface

`gen:dict`

A generic interface (see §5.4 “Generic Interfaces”) that supplies dictionary method implementations for a structure type via the `#:methods` option of `struct` definitions. This interface can be used to implement any of the methods documented as §4.15.2.1 “Primitive Dictionary Methods” and §4.15.2.2 “Derived Dictionary Methods”.

Examples:

```
> (struct alist (v)
    #:methods gen:dict
    [(define (dict-ref dict key
                       [default (lambda () (error "key not
found" key))])
       (cond [(assoc key (alist-v dict)) => cdr]
             [else (if (procedure? default) (default) default)]))
     (define (dict-set dict key val)
       (alist (cons (cons key val) (alist-v dict))))
     (define (dict-remove dict key)
       (define al (alist-v dict))
       (alist (remove* (filter (λ (p) (equal? (car p) key)) al) al)))
     (define (dict-count dict)
       (length (remove-duplicates (alist-v dict) #:key car)))])

; etc. other methods
> (define d1 (alist '((1 . a) (2 . b))))

> (dict? d1)
#t
> (dict-ref d1 1)
'a
> (dict-remove d1 1)
#<alist>
```

`prop:dict` : `struct-type-property?`

A deprecated structure type property used to define custom extensions to the dictionary API. Use `gen:dict` instead. Accepts a vector of 10 method implementations:

- `dict-ref`
- `dict-set!`, or `#f` if unsupported
- `dict-set`, or `#f` if unsupported

- `dict-remove!`, or `#f` if unsupported
- `dict-remove`, or `#f` if unsupported
- `dict-count`
- `dict-iterate-first`
- `dict-iterate-next`
- `dict-iterate-key`
- `dict-iterate-value`

Primitive Dictionary Methods

These methods of `gen:dict` have no fallback implementations; they are only supported for dictionary types that directly implement them.

```
(dict-ref dict key [failure-result]) → any
  dict : dict?
  key : any/c
  failure-result : any/c
                 = (lambda () (raise (make-exn:fail ....)))
```

Returns the value for *key* in *dict*. If no value is found for *key*, then *failure-result* determines the result:

- If *failure-result* is a procedure, it is called (through a tail call) with no arguments to produce the result.
- Otherwise, *failure-result* is returned as the result.

Examples:

```
> (dict-ref #hash((a . "apple") (b . "beer")) 'a)
"apple"
> (dict-ref #hash((a . "apple") (b . "beer")) 'c)
hash-ref: no value found for key
  key: 'c
> (dict-ref #hash((a . "apple") (b . "beer")) 'c #f)
#f
> (dict-ref '((a . "apple") (b . "banana")) 'b)
"banana"
> (dict-ref #("apple" "banana") 1)
"banana"
> (dict-ref #("apple" "banana") 3 #f)
#f
```

```
> (dict-ref #("apple" "banana") -3 #f)
dict-ref: contract violation
  expected: exact-nonnegative-integer?
  given: -3
  in: the k argument of
      (->i
       ((d dict?) (k (d) (dict-key-contract d)))
       ((default any/c))
       any)
  contract from: <collects>/racket/dict.rkt
  blaming: top-level
   (assuming the contract is correct)
  at: <collects>/racket/dict.rkt:181.2
```

```
(dict-set! dict key v) → void?
  dict : (and/c dict? (not/c immutable?))
  key : any/c
  v : any/c
```

Maps *key* to *v* in *dict*, overwriting any existing mapping for *key*. The update can fail with a exn:fail:contract exception if *dict* is not mutable or if *key* is not an allowed key for the dictionary (e.g., not an exact integer in the appropriate range when *dict* is a vector).

Examples:

```
> (define h (make-hash))

> (dict-set! h 'a "apple")

> h
'#hash((a . "apple"))
> (define v (vector #f #f #f))

> (dict-set! v 0 "apple")

> v
'#("apple" #f #f)
```

```
(dict-set dict key v) → (and/c dict? immutable?)
  dict : (and/c dict? immutable?)
  key : any/c
  v : any/c
```

Functionally extends *dict* by mapping *key* to *v*, overwriting any existing mapping for *key*, and returning an extended dictionary. The update can fail with a exn:fail:contract exception if *dict* does not support functional extension or if *key* is not an allowed key for the dictionary.

Examples:

```
> (dict-set #hash() 'a "apple")
'#hash((a . "apple"))
> (dict-set #hash((a . "apple") (b . "beer")) 'b "banana")
'#hash((a . "apple") (b . "banana"))
> (dict-set '() 'a "apple")
'((a . "apple"))
> (dict-set '((a . "apple") (b . "beer")) 'b "banana")
'((a . "apple") (b . "banana"))
```

```
(dict-remove! dict key) → void?
  dict : (and/c dict? (not/c immutable?))
  key : any/c
```

Removes any existing mapping for *key* in *dict*. The update can fail if *dict* is not mutable or does not support removing keys (as is the case for vectors, for example).

Examples:

```
> (define h (make-hash))

> (dict-set! h 'a "apple")

> h
'#hash((a . "apple"))
> (dict-remove! h 'a)

> h
'#hash()
```

```
(dict-remove dict key) → (and/c dict? immutable?)
  dict : (and/c dict? immutable?)
  key : any/c
```

Functionally removes any existing mapping for *key* in *dict*, returning the fresh dictionary. The update can fail if *dict* does not support functional update or does not support removing keys.

Examples:

```
> (define h #hash())

> (define h (dict-set h 'a "apple"))
```

```
> h
'#hash((a . "apple"))
> (dict-remove h 'a)
'#hash()
> h
'#hash((a . "apple"))
> (dict-remove h 'z)
'#hash((a . "apple"))
> (dict-remove '((a . "apple") (b . "banana")) 'a)
'((b . "banana"))
```

(dict-iterate-first *dict*) → any/c
  dict : dict?

Returns #f if *dict* contains no elements, otherwise it returns a non-#f value that is an index to the first element in the dict table; "first" refers to an unspecified ordering of the dictionary elements. For a mutable *dict*, this index is guaranteed to refer to the first item only as long as no mappings are added to or removed from *dict*.

Examples:

```
> (dict-iterate-first #hash((a . "apple") (b . "banana")))
1
> (dict-iterate-first #hash())
#f
> (dict-iterate-first #("apple" "banana"))
0
> (dict-iterate-first '((a . "apple") (b . "banana")))
#<assoc-iter>
```

(dict-iterate-next *dict pos*) → any/c
  dict : dict?
  pos : any/c

Returns either a non-#f that is an index to the element in *dict* after the element indexed by *pos* or #f if *pos* refers to the last element in *dict*. If *pos* is not a valid index, then the exn:fail:contract exception is raised. For a mutable *dict*, the result index is guaranteed to refer to its item only as long as no items are added to or removed from *dict*. The dict-iterate-next operation should take constant time.

Examples:

```
> (define h #hash((a . "apple") (b . "banana")))

> (define i (dict-iterate-first h))
```

```
> i
1
> (dict-iterate-next h i)
3
> (dict-iterate-next h (dict-iterate-next h i))
#f
```

```
(dict-iterate-key dict pos) → any
  dict : dict?
  pos : any/c
```

Returns the key for the element in *dict* at index *pos*. If *pos* is not a valid index for *dict*, the exn:fail:contract exception is raised. The dict-iterate-key operation should take constant time.

Examples:

```
> (define h '((a . "apple") (b . "banana")))

> (define i (dict-iterate-first h))

> (dict-iterate-key h i)
'a
> (dict-iterate-key h (dict-iterate-next h i))
'b
```

```
(dict-iterate-value dict pos) → any
  dict : dict?
  pos : any/c
```

Returns the value for the element in *dict* at index *pos*. If *pos* is not a valid index for *dict*, the exn:fail:contract exception is raised. The dict-iterate-key operation should take constant time.

Examples:

```
> (define h '((a . "apple") (b . "banana")))

> (define i (dict-iterate-first h))

> (dict-iterate-value h i)
"apple"
> (dict-iterate-value h (dict-iterate-next h i))
"banana"
```

Derived Dictionary Methods

These methods of gen:dict have fallback implementations in terms of the other methods; they may be supported even by dictionary types that do not directly implement them.

```
(dict-has-key? dict key) → boolean?
  dict : dict?
  key : any/c
```

Returns #t if dict contains a value for the given key, #f otherwise.

Supported for any dict that implements dict-ref.

Examples:

```
> (dict-has-key? #hash((a . "apple") (b . "beer")) 'a)
#t
> (dict-has-key? #hash((a . "apple") (b . "beer")) 'c)
#f
> (dict-has-key? '((a . "apple") (b . "banana")) 'b)
#t
> (dict-has-key? #("apple" "banana") 1)
#t
> (dict-has-key? #("apple" "banana") 3)
#f
> (dict-has-key? #("apple" "banana") -3)
#f
```

```
(dict-set*! dict key v ... ...) → void?
  dict : (and/c dict? (not/c immutable?))
  key : any/c
  v : any/c
```

Maps each key to each v in dict, overwriting any existing mapping for each key. The update can fail with a exn:fail:contract exception if dict is not mutable or if any key is not an allowed key for the dictionary (e.g., not an exact integer in the appropriate range when dict is a vector). The update takes place from the left, so later mappings overwrite earlier mappings.

Supported for any dict that implements dict-set!.

Examples:

```
> (define h (make-hash))

> (dict-set*! h 'a "apple" 'b "banana")

> h
'#hash((a . "apple") (b . "banana"))
> (define v1 (vector #f #f #f))
```

```
> (dict-set*! v1 0 "apple" 1 "banana")

> v1
'#("apple" "banana" #f)
> (define v2 (vector #f #f #f))

> (dict-set*! v2 0 "apple" 0 "banana")

> v2
'#("banana" #f #f)
```

```
(dict-set* dict key v ... ...) → (and/c dict? immutable?)
  dict : (and/c dict? immutable?)
  key : any/c
  v : any/c
```

Functionally extends *dict* by mapping each *key* to each *v*, overwriting any existing mapping for each *key*, and returning an extended dictionary. The update can fail with a exn:fail:contract exception if *dict* does not support functional extension or if any *key* is not an allowed key for the dictionary. The update takes place from the left, so later mappings overwrite earlier mappings.

Supported for any *dict* that implements dict-set.

Examples:

```
> (dict-set* #hash() 'a "apple" 'b "beer")
'#hash((a . "apple") (b . "beer"))
> (dict-set* #hash((a . "apple") (b . "beer")) 'b "banana" 'a "anchor")
'#hash((a . "anchor") (b . "banana"))
> (dict-set* '() 'a "apple" 'b "beer")
'((a . "apple") (b . "beer"))
> (dict-set* '((a . "apple") (b . "beer")) 'b "banana" 'a "anchor")
'((a . "anchor") (b . "banana"))
> (dict-set* '((a . "apple") (b . "beer")) 'b "banana" 'b "ballistic")
'((a . "apple") (b . "ballistic"))
```

```
(dict-ref! dict key to-set) → any
  dict : dict?
  key : any/c
  to-set : any/c
```

Returns the value for *key* in *dict*. If no value is found for *key*, then *to-set* determines the result as in dict-ref (i.e., it is either a thunk that computes a value or a plain value), and this result is stored in *dict* for the *key*. (Note that if *to-set* is a thunk, it is not invoked in tail position.)

Supported for any *dict* that implements dict-ref and dict-set!.

Examples:

```
> (dict-ref! (make-hasheq '((a . "apple") (b . "beer"))) 'a)
dict-ref!: arity mismatch;
 the expected number of arguments does not match the given
number
  expected: 3
  given: 2
  arguments...:
   '#hasheq((b . "beer") (a . "apple"))
   'a
> (dict-ref! (make-hasheq '((a . "apple") (b .
"beer"))) 'c 'cabbage)
'cabbage
> (define h (make-hasheq '((a . "apple") (b . "beer"))))

> (dict-ref h 'c)
hash-ref: no value found for key
  key: 'c
> (dict-ref! h 'c (λ () 'cabbage))
'cabbage
> (dict-ref h 'c)
'cabbage
```

```
(dict-update!  dict
               key
               updater
              [failure-result]) → void?
  dict : (and/c dict? (not/c immutable?))
  key : any/c
  updater : (any/c . -> . any/c)
  failure-result :  any/c
                 = (lambda () (raise (make-exn:fail ....)))
```

Composes `dict-ref` and `dict-set!` to update an existing mapping in *`dict`*, where the optional *`failure-result`* argument is used as in `dict-ref` when no mapping exists for *`key`* already.

Supported for any *`dict`* that implements `dict-ref` and `dict-set!`.

Examples:

```
> (define h (make-hash))

> (dict-update! h 'a add1)
hash-update!: no value found for key: 'a
> (dict-update! h 'a add1 0)
```

```
> h
'#hash((a . 1))
> (define v (vector #f #f #f))

> (dict-update! v 0 not)

> v
'#(#t #f #f)
```

```
(dict-update dict key updater [failure-result])
 → (and/c dict? immutable?)
  dict : dict?
  key : any/c
  updater : (any/c . -> . any/c)
  failure-result : any/c
                 = (lambda () (raise (make-exn:fail ....)))
```

Composes `dict-ref` and `dict-set` to functionally update an existing mapping in *dict*, where the optional *failure-result* argument is used as in `dict-ref` when no mapping exists for *key* already.

Supported for any *dict* that implements `dict-ref` and `dict-set`.

Examples:

```
> (dict-update #hash() 'a add1)
hash-update: no value found for key: 'a
> (dict-update #hash() 'a add1 0)
'#hash((a . 1))
> (dict-update #hash((a . "apple") (b . "beer")) 'b string-length)
'#hash((a . "apple") (b . 4))
```

```
(dict-map dict proc) → (listof any/c)
  dict : dict?
  proc : (any/c any/c . -> . any/c)
```

Applies the procedure *proc* to each element in *dict* in an unspecified order, accumulating the results into a list. The procedure *proc* is called each time with a key and its value.

Supported for any *dict* that implements `dict-iterate-first`, `dict-iterate-next`, `dict-iterate-key`, and `dict-iterate-value`.

Example:

```
> (dict-map #hash((a . "apple") (b . "banana")) vector)
'(#(a "apple") #(b "banana"))
```

```
(dict-for-each dict proc) → void?
  dict : dict?
  proc : (any/c any/c . -> . any)
```

Applies *proc* to each element in *dict* (for the side-effects of *proc*) in an unspecified order. The procedure *proc* is called each time with a key and its value.

Supported for any *dict* that implements `dict-iterate-first`, `dict-iterate-next`, `dict-iterate-key`, and `dict-iterate-value`.

Example:

```
> (dict-for-each #hash((a . "apple") (b . "banana"))
                 (lambda (k v)
                   (printf "~a = ~s\n" k v)))
a = "apple"
b = "banana"
```

```
(dict-empty? dict) → boolean?
  dict : dict?
```

Reports whether *dict* is empty.

Supported for any *dict* that implements `dict-iterate-first`.

Examples:

```
> (dict-empty? #hash((a . "apple") (b . "banana")))
#f
> (dict-empty? (vector))
#t
```

```
(dict-count dict) → exact-nonnegative-integer?
  dict : dict?
```

Returns the number of keys mapped by *dict*, usually in constant time.

Supported for any *dict* that implements `dict-iterate-first` and `dict-iterate-next`.

Examples:

```
> (dict-count #hash((a . "apple") (b . "banana")))
2
> (dict-count #("apple" "banana"))
2
```

```
(dict-copy dict) → dict?
  dict : dict?
```

Produces a new, mutable dictionary of the same type as *dict* and with the same key/value associations.

Supported for any *dict* that implements dict-clear, dict-set!, dict-iterate-first, dict-iterate-next, dict-iterate-key, and dict-iterate-value.

Examples:

```
> (define original (vector "apple" "banana"))

> (define copy (dict-copy original))

> original
'#("apple" "banana")
> copy
'#("apple" "banana")
> (dict-set! copy 1 "carrot")

> original
'#("apple" "banana")
> copy
'#("apple" "carrot")
```

```
(dict-clear dict) → dict?
  dict : dict?
```

Produces an empty dictionary of the same type as *dict*. If *dict* is mutable, the result must be a new dictionary.

Supported for any *dict* that supports dict-remove, dict-iterate-first, dict-iterate-next, and dict-iterate-key.

Examples:

```
> (dict-clear #hash((a . "apple") ("banana" . b)))
'#hash()
> (dict-clear '((1 . two) (three . "four")))
'()
```

```
(dict-clear! dict) → dict?
  dict : dict?
```

Removes all of the key/value associations in *dict*.

Supported for any *dict* that supports dict-remove!, dict-iterate-first, and dict-iterate-key.

Examples:

```
> (define table (make-hash))

> (dict-set! table 'a "apple")

> (dict-set! table "banana" 'b)

> table
'#hash(("banana" . b) (a . "apple"))
> (dict-clear! table)

> table
'#hash()
```

(dict-keys *dict*) → list?
  dict : dict?

Returns a list of the keys from *dict* in an unspecified order.

Supported for any *dict* that implements dict-iterate-first, dict-iterate-next, and dict-iterate-key.

Examples:

```
> (define h #hash((a . "apple") (b . "banana")))

> (dict-keys h)
'(a b)
```

(dict-values *dict*) → list?
  dict : dict?

Returns a list of the values from *dict* in an unspecified order.

Supported for any *dict* that implements dict-iterate-first, dict-iterate-next, and dict-iterate-value.

Examples:

```
> (define h #hash((a . "apple") (b . "banana")))

> (dict-values h)
'("apple" "banana")
```

```
(dict->list dict) → list?
  dict : dict?
```

Returns a list of the associations from *dict* in an unspecified order.

Supported for any *dict* that implements dict-iterate-first, dict-iterate-next, dict-iterate-key, and dict-iterate-value.

Examples:

```
> (define h #hash((a . "apple") (b . "banana")))

> (dict->list h)
'((a . "apple") (b . "banana"))
```

4.15.3 Dictionary Sequences

```
(in-dict dict) → sequence?
  dict : dict?
```

Returns a sequence whose each element is two values: a key and corresponding value from *dict*.

Supported for any *dict* that implements dict-iterate-first, dict-iterate-next, dict-iterate-key, and dict-iterate-value.

Examples:

```
> (define h #hash((a . "apple") (b . "banana")))

> (for/list ([(k v) (in-dict h)])
    (format "~a = ~s" k v))
'("a = \"apple\"" "b = \"banana\"")
```

```
(in-dict-keys dict) → sequence?
  dict : dict?
```

Returns a sequence whose elements are the keys of *dict*.

Supported for any *dict* that implements dict-iterate-first, dict-iterate-next, and dict-iterate-key.

Examples:

```
> (define h #hash((a . "apple") (b . "banana")))

> (for/list ([k (in-dict-keys h)])
    k)
'(a b)
```

(in-dict-values *dict*) → sequence?
  dict : dict?

Returns a sequence whose elements are the values of *dict*.

Supported for any *dict* that implements dict-iterate-first, dict-iterate-next, and dict-iterate-value.

Examples:

```
> (define h #hash((a . "apple") (b . "banana")))

> (for/list ([v (in-dict-values h)])
    v)
'("apple" "banana")
```

(in-dict-pairs *dict*) → sequence?
  dict : dict?

Returns a sequence whose elements are pairs, each containing a key and its value from *dict* (as opposed to using in-dict, which gets the key and value as separate values for each element).

Supported for any *dict* that implements dict-iterate-first, dict-iterate-next, dict-iterate-key, and dict-iterate-value.

Examples:

```
> (define h #hash((a . "apple") (b . "banana")))

> (for/list ([p (in-dict-pairs h)])
    p)
'((a . "apple") (b . "banana"))
```

4.15.4 Contracted Dictionaries

prop:dict/contract : struct-type-property?

A structure type property for defining dictionaries with contracts. The value associated with prop:dict/contract must be a list of two immutable vectors:

```
(list dict-vector
      (vector type-key-contract
              type-value-contract
              type-iter-contract
              instance-key-contract
              instance-value-contract
              instance-iter-contract))
```

The first vector must be a vector of 10 procedures which match the `gen:dict` generic interface (in addition, it must be an immutable vector). The second vector must contain six elements; each of the first three is a contract for the dictionary type's keys, values, and positions, respectively. Each of the second three is either `#f` or a procedure used to extract the contract from a dictionary instance.

```
(dict-key-contract d) → contract?
  d : dict?
(dict-value-contract d) → contract?
  d : dict?
(dict-iter-contract d) → contract?
  d : dict?
```

Returns the contract that *d* imposes on its keys, values, or iterators, respectively, if *d* implements the `prop:dict/contract` interface.

4.15.5 Custom Hash Tables

```
(define-custom-hash-types name
                          optional-predicate
                          comparison-expr
                          optional-hash-functions)

     optional-predicate =
                        | #:key? predicate-expr

optional-hash-functions =
                        | hash1-expr
                        | hash1-expr hash2-expr
```

Creates a new dictionary type based on the given comparison *comparison-expr*, hash functions *hash1-expr* and *hash2-expr*, and key predicate *predicate-expr*; the interfaces for these functions are the same as in `make-custom-hash-types`. The new dictionary type has three variants: immutable, mutable with strongly-held keys, and mutable with weakly-held keys.

Defines seven names:

- *name*? recognizes instances of the new type,
- immutable-*name*? recognizes immutable instances of the new type,
- mutable-*name*? recognizes mutable instances of the new type with strongly-held keys,
- weak-*name*? recognizes mutable instances of the new type with weakly-held keys,
- make-immutable-*name* constructs immutable instances of the new type,
- make-mutable-*name* constructs mutable instances of the new type with strongly-held keys, and
- make-weak-*name* constructs mutable instances of the new type with weakly-held keys.

The constructors all accept a dictionary as an optional argument, providing initial key/value pairs.

Examples:

```
> (define-custom-hash-types string-hash
                            #:key? string?
                            string=?
                            string-length)

> (define imm
    (make-immutable-string-hash
     '(("apple" . a) ("banana" . b))))

> (define mut
    (make-mutable-string-hash
     '(("apple" . a) ("banana" . b))))

> (dict? imm)
#t
> (dict? mut)
#t
> (string-hash? imm)
#t
> (string-hash? mut)
#t
> (immutable-string-hash? imm)
#t
> (immutable-string-hash? mut)
#f
> (dict-ref imm "apple")
```

```
'a
> (dict-ref mut "banana")
'b
> (dict-set! mut "banana" 'berry)

> (dict-ref mut "banana")
'berry
> (equal? imm mut)
#f
> (equal? (dict-remove (dict-remove imm "apple") "banana")
          (make-immutable-string-hash))
#t
```

```
(make-custom-hash-types  eql?
                         [hash1
                          hash2
                          #:key? key?
                          #:name name
                                        (any/c . -> . boolean?)
                                        (any/c . -> . boolean?)
                                        (any/c . -> . boolean?)
                          #:for who]) → (any/c . -> . boolean?)
                                        (->* [] [dict?] dict?)
                                        (->* [] [dict?] dict?)
                                        (->* [] [dict?] dict?)
  eql? : (or/c (any/c any/c . -> . any/c)
               (any/c any/c (any/c any/c . -> . any/c) . -> . any/c))
  hash1 : (or/c (any/c . -> . exact-integer?)
                (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
        = (const 1)
  hash2 : (or/c (any/c . -> . exact-integer?)
                (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
        = (const 1)
  key? : (any/c . -> . boolean?) = (const #true)
  name : symbol? = 'custom-hash
  who : symbol? = 'make-custom-hash-types
```

Creates a new dictionary type based on the given comparison function *eql?*, hash functions *hash1* and *hash2*, and predicate *key?*. The new dictionary type has variants that are immutable, mutable with strongly-held keys, and mutable with weakly-held keys. The given *name* is used when printing instances of the new dictionary type, and the symbol *who* is used for reporting errors.

The comparison function *eql?* may accept 2 or 3 arguments. If it accepts 2 arguments, it given two keys to compare them. If it accepts 3 arguments and does not accept 2 arguments, it is also given a recursive comparison function that handles data cycles when comparing sub-parts of the keys.

The hash functions *hash1* and *hash2* may accept 1 or 2 arguments. If either hash function accepts 1 argument, it is applied to a key to compute the corresponding hash value. If either hash function accepts 2 arguments and does not accept 1 argument, it is also given a recursive hash function that handles data cycles when computing hash values of sub-parts of the keys.

The predicate *key?* must accept 1 argument and is used to recognize valid keys for the new dictionary type.

Produces seven values:

- a predicate recognizing all instances of the new dictionary type,
- a predicate recognizing immutable instances,
- a predicate recognizing mutable instances,
- a predicate recognizing weak instances,
- a constructor for immutable instances,
- a constructor for mutable instances, and
- a constructor for weak instances.

See `define-custom-hash-types` for an example.

```
(make-custom-hash  eql?
                   [hash1
                    hash2
                    #:key? key?]) → dict?
  eql? : (or/c (any/c any/c . -> . any/c)
               (any/c any/c (any/c any/c . -> . any/c) . -> . any/c))
  hash1 : (or/c (any/c . -> . exact-integer?)
                (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
        = (const 1)
  hash2 : (or/c (any/c . -> . exact-integer?)
                (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
        = (const 1)
  key? : (any/c . -> . boolean?) = (const #true)
```

```
(make-weak-custom-hash  eql?
                        [hash1
                         hash2
                         #:key? key?]) → dict?
  eql? : (or/c (any/c any/c . -> . any/c)
               (any/c any/c (any/c any/c . -> . any/c) . -> . any/c))
  hash1 : (or/c (any/c . -> . exact-integer?)
                (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
        = (const 1)
  hash2 : (or/c (any/c . -> . exact-integer?)
                (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
        = (const 1)
  key? : (any/c . -> . boolean?) = (const #true)
(make-immutable-custom-hash  eql?
                             [hash1
                              hash2
                              #:key? key?]) → dict?
  eql? : (or/c (any/c any/c . -> . any/c)
               (any/c any/c (any/c any/c . -> . any/c) . -> . any/c))
  hash1 : (or/c (any/c . -> . exact-integer?)
                (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
        = (const 1)
  hash2 : (or/c (any/c . -> . exact-integer?)
                (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
        = (const 1)
  key? : (any/c . -> . boolean?) = (const #true)
```

Constructs an instance of a new dictionary type based on the given comparison function *eql?*, hash functions *hash1* and *hash2*, and key predicate *key?*.

These procedures are deprecated; use `define-custom-hash-types` instead.

4.16 Sets

A *set* represents a collection of distinct elements. The following datatypes are all sets:

- hash sets;
- lists using `equal?` to compare elements; and
- structures whose types implement the `gen:set` generic interface.

`(require racket/set)` package: `base`

The bindings documented in this section are provided by the `racket/set` and `racket` libraries, but not `racket/base`.

4.16.1 Hash Sets

A *hash set* is a set whose elements are compared via `equal?`, `eqv?`, or `eq?` and partitioned via `equal-hash-code`, `eqv-hash-code`, or `eq-hash-code`. A hash set is either immutable or mutable; mutable hash sets retain their elements either strongly or weakly.

Like operations on immutable hash tables, "constant time" hash set operations actually require $O(\log N)$ time for a set of size N.

A hash set can be used as a stream (see §4.14.2 "Streams") and thus as a single-valued sequence (see §4.14.1 "Sequences"). The elements of the set serve as elements of the stream or sequence. If an element is added to or removed from the hash set during iteration, then an iteration step may fail with `exn:fail:contract`, or the iteration may skip or duplicate elements. See also `in-set`.

Two hash sets are `equal?` when they use the same element-comparison procedure (`equal?`, `eqv?`, or `eq?`), both hold elements strongly or weakly, have the same mutability, and have equivalent elements. Immutable hash sets support effectively constant-time access and update, just like mutable hash sets; the constant on immutable operations is usually larger, but the functional nature of immutable hash sets can pay off in certain algorithms.

All hash sets implement `set->stream`, `set-empty?`, `set-member?`, `set-count`, `subset?`, `proper-subset?`, `set-map`, `set-for-each`, `set-copy`, `set-copy-clear`, `set->list`, and `set-first`. Immutable hash sets in addition implement `set-add`, `set-remove`, `set-clear`, `set-union`, `set-intersect`, `set-subtract`, and `set-symmetric-difference`. Mutable hash sets in addition implement `set-add!`, `set-remove!`, `set-clear!`, `set-union!`, `set-intersect!`, `set-subtract!`, and `set-symmetric-difference!`.

Operations on sets that contain elements that are mutated are unpredictable in much the same way that hash table operations are unpredictable when keys are mutated.

```
(set-equal? x) → boolean?
  x : any/c
(set-eqv? x) → boolean?
  x : any/c
(set-eq? x) → boolean?
  x : any/c
```

Returns `#t` if `x` is a hash set that compares elements with `equal?`, `eqv?`, or `eq?`, respectively; returns `#f` otherwise.

```
(set? x) → boolean?
  x : any/c
(set-mutable? x) → boolean?
  x : any/c
```

```
(set-weak? x) → boolean?
  x : any/c
```

Returns `#t` if *x* is a hash set that is respectively immutable, mutable with strongly-held keys, or mutable with weakly-held keys; returns `#f` otherwise.

```
(set v ...) → (and/c generic-set? set-equal? set?)
  v : any/c
(seteqv v ...) → (and/c generic-set? set-eqv? set?)
  v : any/c
(seteq v ...) → (and/c generic-set? set-eq? set?)
  v : any/c
(mutable-set v ...)
 → (and/c generic-set? set-equal? set-mutable?)
  v : any/c
(mutable-seteqv v ...)
 → (and/c generic-set? set-eqv? set-mutable?)
  v : any/c
(mutable-seteq v ...)
 → (and/c generic-set? set-eq? set-mutable?)
  v : any/c
(weak-set v ...) → (and/c generic-set? set-equal? set-weak?)
  v : any/c
(weak-seteqv v ...) → (and/c generic-set? set-eqv? set-weak?)
  v : any/c
(weak-seteq v ...) → (and/c generic-set? set-eq? set-weak?)
  v : any/c
```

Creates a hash set with the given `v`s as elements. The elements are added in the order that they appear as arguments, so in the case of sets that use `equal?` or `eqv?`, an earlier element may be replaced by a later element that is `equal?` or `eqv?` but not `eq?`.

```
(list->set lst) → (and/c generic-set? set-equal? set?)
  lst : list?
(list->seteqv lst) → (and/c generic-set? set-eqv? set?)
  lst : list?
(list->seteq lst) → (and/c generic-set? set-eq? set?)
  lst : list?
(list->mutable-set lst)
 → (and/c generic-set? set-equal? set-mutable?)
  lst : list?
(list->mutable-seteqv lst)
 → (and/c generic-set? set-eqv? set-mutable?)
  lst : list?
(list->mutable-seteq lst)
 → (and/c generic-set? set-eq? set-mutable?)
  lst : list?
```

```
(list->weak-set lst)
 → (and/c generic-set? set-equal? set-weak?)
  lst : list?
(list->weak-seteqv lst)
 → (and/c generic-set? set-eqv? set-weak?)
  lst : list?
(list->weak-seteq lst) → (and/c generic-set? set-eq? set-weak?)
  lst : list?
```

Creates a hash set with the elements of the given `lst` as the elements of the set. Equivalent to `(apply set lst)`, `(apply seteqv lst)`, `(apply seteq lst)`, and so on, respectively.

```
(for/set (for-clause ...) body ...+)
(for/seteq (for-clause ...) body ...+)
(for/seteqv (for-clause ...) body ...+)
(for*/set (for-clause ...) body ...+)
(for*/seteq (for-clause ...) body ...+)
(for*/seteqv (for-clause ...) body ...+)
(for/mutable-set (for-clause ...) body ...+)
(for/mutable-seteq (for-clause ...) body ...+)
(for/mutable-seteqv (for-clause ...) body ...+)
(for*/mutable-set (for-clause ...) body ...+)
(for*/mutable-seteq (for-clause ...) body ...+)
(for*/mutable-seteqv (for-clause ...) body ...+)
(for/weak-set (for-clause ...) body ...+)
(for/weak-seteq (for-clause ...) body ...+)
(for/weak-seteqv (for-clause ...) body ...+)
(for*/weak-set (for-clause ...) body ...+)
(for*/weak-seteq (for-clause ...) body ...+)
(for*/weak-seteqv (for-clause ...) body ...+)
```

Analogous to `for/list` and `for*/list`, but to construct a hash set instead of a list.

4.16.2 Set Predicates and Contracts

```
(generic-set? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a set; returns `#f` otherwise.

Examples:

```
> (generic-set? (list 1 2 3))
```

```
#t
> (generic-set? (set 1 2 3))
#t
> (generic-set? (mutable-seteq 1 2 3))
#t
> (generic-set? (vector 1 2 3))
#f
```

```
(set-implements? st sym ...) → boolean?
  st : generic-set?
  sym : symbol?
```

Returns `#t` if *st* implements all of the methods from `gen:set` named by the *sym*s; returns `#f` otherwise. Fallback implementations do not affect the result; *st* may support the given methods via fallback implementations yet produce `#f`.

Examples:

```
> (set-implements? (list 1 2 3) 'set-add)
#t
> (set-implements? (list 1 2 3) 'set-add!)
#f
> (set-implements? (set 1 2 3) 'set-add)
#t
> (set-implements? (set 1 2 3) 'set-add!)
#f
> (set-implements? (mutable-seteq 1 2 3) 'set-add)
#f
> (set-implements? (mutable-seteq 1 2 3) 'set-add!)
#t
> (set-implements? (weak-seteqv 1 2 3) 'set-remove 'set-remove!)
#f
```

```
(set-implements/c sym ...) → flat-contract?
  sym : symbol?
```

Recognizes sets that support all of the methods from `gen:set` named by the *sym*s.

```
(set/c elem/c [#:cmp cmp #:kind kind]) → contract?
  elem/c : chaperone-contract?
  cmp : (or/c 'dont-care 'equal 'eqv 'eq) = 'dont-care
  kind : (or/c 'dont-care 'immutable 'mutable 'weak 'mutable-or-weak)
       = 'dont-care
```

Constructs a contract that recognizes sets whose elements match `contract`.

If *kind* is 'immutable, 'mutable, or 'weak, the resulting contract accepts only hash sets that are respectively immutable, mutable with strongly-held keys, or mutable with weakly-held keys. If *kind* is 'mutable-or-weak, the resulting contract accepts any mutable "hash sets", regardless of key-holding strength.

If *cmp* is 'equal, 'eqv, or 'eq, the resulting contract accepts only hash sets that compare elements using equal?, eqv?, or eq?, respectively.

If *cmp* is 'eqv or 'eq, then *elem/c* must be a flat contract.

If *cmp* and *kind* are both 'dont-care, then the resulting contract will accept any kind of set, not just hash sets.

4.16.3 Generic Set Interface

gen:set

A generic interface (see §5.4 "Generic Interfaces") that supplies set method implementations for a structure type via the #:methods option of struct definitions. This interface can be used to implement any of the methods documented as §4.16.3.1 "Set Methods".

Examples:

```
> (struct binary-set [integer]
    #:transparent
    #:methods gen:set
    [(define (set-member? st i)
       (bitwise-bit-set? (binary-set-integer st) i))
     (define (set-add st i)
       (binary-set (bitwise-ior (binary-set-integer st)
                                (arithmetic-shift 1 i))))
     (define (set-remove st i)
       (binary-set (bitwise-and (binary-set-integer st)
                                (bitwise-not (arithmetic-
shift 1 i))))))])

> (define bset (binary-set 5))

> bset
(binary-set 5)
> (generic-set? bset)
#t
> (set-member? bset 0)
#t
> (set-member? bset 1)
```

```
#f
> (set-member? bset 2)
#t
> (set-add bset 4)
(binary-set 21)
> (set-remove bset 2)
(binary-set 1)
```

Set Methods

The methods of `gen:set` can be classified into three categories, as determined by their fallback implementations:

1. methods with no fallbacks,
2. methods whose fallbacks depend on other, non-fallback methods,
3. and methods whose fallbacks can depend on either fallback or non-fallback methods.

As an example, implementing the following methods would guarantee that all the methods in `gen:set` would at least have a fallback method:

- `set-member?`
- `set-add`
- `set-add!`
- `set-remove`
- `set-remove!`
- `set-first`
- `set-empty?`

There may be other such subsets of methods that would guarantee at least a fallback for every method.

```
(set-member? st v) → boolean?
  st : generic-set?
  v : any/c
```

Returns `#t` if `v` is in `st`, `#f` otherwise. Has no fallback.

```
(set-add st v) → generic-set?
  st : generic-set?
  v : any/c
```

Produces a set that includes `v` plus all elements of `st`. This operation runs in constant time for hash sets. Has no fallback.

```
(set-add! st v) → void?
  st : generic-set?
  v : any/c
```

Adds the element `v` to `st`. This operation runs in constant time for hash sets. Has no fallback.

```
(set-remove st v) → generic-set?
  st : generic-set?
  v : any/c
```

Produces a set that includes all elements of `st` except `v`. This operation runs in constant time for hash sets. Has no fallback.

```
(set-remove! st v) → void?
  st : generic-set?
  v : any/c
```

Removes the element `v` from `st`. This operation runs in constant time for hash sets. Has no fallback.

```
(set-empty? st) → boolean?
  st : generic-set?
```

Returns `#t` if `st` has no members; returns `#f` otherwise.

Supported for any `st` that implements `set->stream` or `set-count`.

```
(set-count st) → exact-nonnegative-integer?
  st : generic-set?
```

Returns the number of elements in `st`.

Supported for any `st` that supports `set->stream`.

```
(set-first st) → any/c
  st : (and/c generic-set? (not/c set-empty?))
```

Produces an unspecified element of `st`. Multiple uses of `set-first` on `st` produce the same result.

Supported for any `st` that implements `set->stream`.

```
(set-rest st) → generic-set?
  st : (and/c generic-set? (not/c set-empty?))
```

Produces a set that includes all elements of *st* except (set-first *st*).

Supported for any *st* that implements set-remove and either set-first or set->stream.

```
(set->stream st) → stream?
  st : generic-set?
```

Produces a stream containing the elements of *st*.

Supported for any *st* that implements:

- set->list
- in-set
- set-empty?, set-first, set-rest
- set-empty?, set-first, set-remove
- set-count, set-first, set-rest
- set-count, set-first, set-remove

```
(set-copy st) → generic-set?
  st : generic-set?
```

Produces a new, mutable set of the same type and with the same elements as *st*.

Supported for any *st* that supports set->stream and either implements set-copy-clear and set-add!.

```
(set-copy-clear st) → (and/c generic-set? set-empty?)
  st : generic-set?
```

Produces a new, empty set of the same type, mutability, and key strength as *st*.

A difference between set-copy-clear and set-clear is that the latter conceptually iterates set-remove on the given set, and so it preserves any contract on the given set. The set-copy-clear function produces a new set without any contracts.

Supported for any *st* that implements set-remove and supports set->stream.

```
(set-clear st) → (and/c generic-set? set-empty?)
  st : generic-set?
```

Produces set by removing all elements of *st*.

Supported for any *st* that implements `set-remove` and supports `set->stream`.

```
(set-clear! st) → void?
  st : generic-set?
```

Removes all elements from *st*.

Supported for any *st* that implements `set-remove!` and either supports `set->stream` or implements `set-first` and either `set-count` or `set-empty?`.

```
(set-union st0 st ...) → generic-set?
  st0 : generic-set?
  st : generic-set?
```

Produces a set of the same type as *st0* that includes the elements from *st0* and all of the *st*s.

If *st0* is a list, each *st* must also be a list. This operation runs on lists in time proportional to the total size of the *st*s times the size of the result.

If *st0* is a hash set, each *st* must also be a hash set that uses the same comparison function (`equal?`, `eqv?`, or `eq?`). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the total size of all of the sets except the largest immutable set.

At least one set must be provided to `set-union` to determine the type of the resulting set (list, hash set, etc.). If there is a case where `set-union` may be applied to zero arguments, instead pass an empty set of the intended type as the first argument.

Supported for any *st* that implements `set-add` and supports `set->stream`.

Examples:

```
> (set-union (set))
(set)
> (set-union (seteq))
(seteq)
> (set-union (set 1 2) (set 2 3))
(set 1 2 3)
> (set-union (list 1 2) (list 2 3))
```

```
'(3 1 2)
> (set-union (set 1 2) (seteq 2 3))
set-union: set arguments have incompatible equivalence
predicates
  first set: (set 1 2)
  incompatible set: (seteq 2 3)
; Sets of different types cannot be unioned.
```

(set-union! st0 st ...) → generic-set?
  st0 : generic-set?
  st : generic-set?

Adds the elements from all of the sts to st0.

If st0 is a hash set, each st must also be a hash set that uses the same comparison function (equal?, eqv?, or eq?). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the total size of the sts.

Supported for any st that implements set-add! and supports set->stream.

(set-intersect st0 st ...) → generic-set?
  st0 : generic-set?
  st : generic-set?

Produces a set of the same type as st0 that includes the elements from st0 that are also contained by all of the sts.

If st0 is a list, each st must also be a list. This operation runs on lists in time proportional to the total size of the sts times the size of st0.

If st0 is a hash set, each st must also be a hash set that uses the same comparison function (equal?, eqv?, or eq?). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the size of the smallest immutable set.

Supported for any st that implements either set-remove or both set-clear and set-add, and supports set->stream.

(set-intersect! st0 st ...) → generic-set?
  st0 : generic-set?
  st : generic-set?

Removes every element from st0 that is not contained by all of the sts.

If st0 is a hash set, each st must also be a hash set that uses the same comparison function (equal?, eqv?, or eq?). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the size of st0.

Supported for any st that implements set-remove! and supports set->stream.

```
(set-subtract st0 st ...) → generic-set?
  st0 : generic-set?
  st : generic-set?
```

Produces a set of the same type as *st0* that includes the elements from *st0* that not contained by any of the *st*s.

If *st0* is a list, each *st* must also be a list. This operation runs on lists in time proportional to the total size of the *st*s times the size of *st0*.

If *st0* is a hash set, each *st* must also be a hash set that uses the same comparison function (`equal?`, `eqv?`, or `eq?`). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the size of *st0*.

Supported for any *st* that implements either `set-remove` or both `set-clear` and `set-add`, and supports `set->stream`.

```
(set-subtract! st0 st ...) → generic-set?
  st0 : generic-set?
  st : generic-set?
```

Removes every element from *st0* that is contained by any of the *st*s.

If *st0* is a hash set, each *st* must also be a hash set that uses the same comparison function (`equal?`, `eqv?`, or `eq?`). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the size of *st0*.

Supported for any *st* that implements `set-remove!` and supports `set->stream`.

```
(set-symmetric-difference st0 st ...) → generic-set?
  st0 : generic-set?
  st : generic-set?
```

Produces a set of the same type as *st0* that includes all of the elements contained an even number of times in *st0* and the *st*s.

If *st0* is a list, each *st* must also be a list. This operation runs on lists in time proportional to the total size of the *st*s times the size of *st0*.

If *st0* is a hash set, each *st* must also be a hash set that uses the same comparison function (`equal?`, `eqv?`, or `eq?`). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the total size of all of the sets except the largest immutable set.

Supported for any *st* that implements `set-remove` or both `set-clear` and `set-add`, and supports `set->stream`.

Example:

```
> (set-symmetric-difference (set 1) (set 1 2) (set 1 2 3))
(set 1 3)
```

```
(set-symmetric-difference! st0 st ...) → generic-set?
  st0 : generic-set?
  st : generic-set?
```

Adds and removes elements of *st0* so that it includes all of the elements contained an even number of times in the *st*s and the original contents of *st0*.

If *st0* is a hash set, each *st* must also be a hash set that uses the same comparison function (equal?, eqv?, or eq?). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the total size of the *st*s.

Supported for any *st* that implements set-remove! and supports set->stream.

```
(set=? st st2) → boolean?
  st : generic-set?
  st2 : generic-set?
```

Returns #t if *st* and *st2* contain the same members; returns #f otherwise.

If st0 is a list, each *st* must also be a list. This operation runs on lists in time proportional to the size of *st* times the size of *st2*.

If st0 is a hash set, each *st* must also be a hash set that uses the same comparison function (equal?, eqv?, or eq?). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the size of *st* plus the size of *st2*.

Supported for any *st* and *st2* that both support subset?; also supported for any if *st2* that implements set=? regardless of *st*.

Examples:

```
> (set=? (list 1 2) (list 2 1))
#t
> (set=? (set 1) (set 1 2 3))
#f
> (set=? (set 1 2 3) (set 1))
#f
> (set=? (set 1 2 3) (set 1 2 3))
#t
> (set=? (seteq 1 2) (mutable-seteq 2 1))
#t
> (set=? (seteq 1 2) (seteqv 1 2))
```

```
set=?: set arguments have incompatible equivalence
predicates
  first set: (seteq 1 2)
  incompatible set: (seteqv 1 2)
; Sets of different types cannot be compared.
```

```
(subset? st st2) → boolean?
  st : generic-set?
  st2 : generic-set?
```

Returns `#t` if *st2* contains every member of *st*; returns `#f` otherwise.

If `st0` is a list, each *st* must also be a list. This operation runs on lists in time proportional to the size of *st* times the size of *st2*.

If `st0` is a hash set, each *st* must also be a hash set that uses the same comparison function (`equal?`, `eqv?`, or `eq?`). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the size of *st*.

Supported for any *st* that supports `set->stream`.

Examples:

```
> (subset? (set 1) (set 1 2 3))
#t
> (subset? (set 1 2 3) (set 1))
#f
> (subset? (set 1 2 3) (set 1 2 3))
#t
```

```
(proper-subset? st st2) → boolean?
  st : generic-set?
  st2 : generic-set?
```

Returns `#t` if *st2* contains every member of *st* and at least one additional element; returns `#f` otherwise.

If `st0` is a list, each *st* must also be a list. This operation runs on lists in time proportional to the size of *st* times the size of *st2*.

If `st0` is a hash set, each *st* must also be a hash set that uses the same comparison function (`equal?`, `eqv?`, or `eq?`). The mutability and key strength of the hash sets may differ. This operation runs on hash sets in time proportional to the size of *st* plus the size of *st2*.

Supported for any *st* and *st2* that both support `subset?`.

Examples:

```
> (proper-subset? (set 1) (set 1 2 3))
#t
> (proper-subset? (set 1 2 3) (set 1))
#f
> (proper-subset? (set 1 2 3) (set 1 2 3))
#f
```

```
(set->list st) → list?
  st : generic-set?
```

Produces a list containing the elements of *st*.

Supported for any *st* that supports set->stream.

```
(set-map st proc) → (listof any/c)
  st : generic-set?
  proc : (any/c . -> . any/c)
```

Applies the procedure *proc* to each element in *st* in an unspecified order, accumulating the results into a list.

Supported for any *st* that supports set->stream.

```
(set-for-each st proc) → void?
  st : generic-set?
  proc : (any/c . -> . any)
```

Applies *proc* to each element in *st* (for the side-effects of *proc*) in an unspecified order.

Supported for any *st* that supports set->stream.

```
(in-set st) → sequence?
  st : generic-set?
```

Explicitly converts a set to a sequence for use with `for` and other forms.

Supported for any *st* that supports set->stream.

4.16.4 Custom Hash Sets

```
(define-custom-set-types name
                         optional-predicate
                         comparison-expr
                         optional-hash-functions)
```

```
     optional-predicate =
                        | #:elem? predicate-expr

optional-hash-functions =
                        | hash1-expr
                        | hash1-expr hash2-expr
```

Creates a new set type based on the given comparison *comparison-expr*, hash functions *hash1-expr* and *hash2-expr*, and element predicate *predicate-expr*; the interfaces for these functions are the same as in `make-custom-set-types`. The new set type has three variants: immutable, mutable with strongly-held elements, and mutable with weakly-held elements.

Defines seven names:

- *name*`?` recognizes instances of the new type,
- `immutable-`*name*`?` recognizes immutable instances of the new type,
- `mutable-`*name*`?` recognizes mutable instances of the new type with strongly-held elements,
- `weak-`*name*`?` recognizes mutable instances of the new type with weakly-held elements,
- `make-immutable-`*name* constructs immutable instances of the new type,
- `make-mutable-`*name* constructs mutable instances of the new type with strongly-held elements, and
- `make-weak-`*name* constructs mutable instances of the new type with weakly-held elements.

The constructors all accept a stream as an optional argument, providing initial elements.

Examples:

```
> (define-custom-set-types string-set
                           #:elem? string?
                           string=?
                           string-length)

> (define imm
    (make-immutable-string-set '("apple" "banana")))
```

```
> (define mut
    (make-mutable-string-set '("apple" "banana")))

> (generic-set? imm)
#t
> (generic-set? mut)
#t
> (string-set? imm)
#t
> (string-set? mut)
#t
> (immutable-string-set? imm)
#t
> (immutable-string-set? mut)
#f
> (set-member? imm "apple")
#t
> (set-member? mut "banana")
#t
> (equal? imm mut)
#f
> (set=? imm mut)
#t
> (set-remove! mut "banana")

> (set-member? mut "banana")
#f
> (equal? (set-remove (set-remove imm "apple") "banana")
          (make-immutable-string-set))
#t
```

```
(make-custom-set-types  eql?
                       [hash1
                        hash2
                        #:elem? elem?
                        #:name name
                        #:for who])
   (any/c . -> . boolean?)
   (any/c . -> . boolean?)
   (any/c . -> . boolean?)
 → (any/c . -> . boolean?)
   (->* [] [stream?] generic-set?)
   (->* [] [stream?] generic-set?)
   (->* [] [stream?] generic-set?)
  eql? : (or/c (any/c any/c . -> . any/c)
               (any/c any/c (any/c any/c . -> . any/c) . -> . any/c))
```

```
hash1 : (or/c (any/c . -> . exact-integer?)
              (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
      = (const 1)
hash2 : (or/c (any/c . -> . exact-integer?)
              (any/c (any/c . -> . exact-integer?) . -> . exact-integer?))
      = (const 1)
elem? : (any/c . -> . boolean?) = (const #true)
name : symbol? = 'custom-set
who : symbol? = 'make-custom-set-types
```

Creates a new set type based on the given comparison function *eql?*, hash functions *hash1* and *hash2*, and predicate *elem?*. The new set type has variants that are immutable, mutable with strongly-held elements, and mutable with weakly-held elements. The given *name* is used when printing instances of the new set type, and the symbol *who* is used for reporting errors.

The comparison function *eql?* may accept 2 or 3 arguments. If it accepts 2 arguments, it given two elements to compare them. If it accepts 3 arguments and does not accept 2 arguments, it is also given a recursive comparison function that handles data cycles when comparing sub-parts of the elements.

The hash functions *hash1* and *hash2* may accept 1 or 2 arguments. If either hash function accepts 1 argument, it is applied to a element to compute the corresponding hash value. If either hash function accepts 2 arguments and does not accept 1 argument, it is also given a recursive hash function that handles data cycles when computing hash values of sub-parts of the elements.

The predicate *elem?* must accept 1 argument and is used to recognize valid elements for the new set type.

Produces seven values:

- a predicate recognizing all instances of the new set type,
- a predicate recognizing immutable instances,
- a predicate recognizing mutable instances,
- a predicate recognizing weak instances,
- a constructor for immutable instances,
- a constructor for mutable instances, and
- a constructor for weak instances.

See `define-custom-hash-types` for an example.

4.17 Procedures

```
(procedure? v) → boolean?
  v : any/c
```

Returns #t if v is a procedure, #f otherwise.

```
(apply proc v ... lst #:<kw> kw-arg ...) → any
  proc : procedure?
  v : any/c
  lst : list?
  kw-arg : any/c
```

§4.3.3 "The apply Function" in ***The Racket Guide*** introduces apply.

Applies proc using the content of (list* v ... lst) as the (by-position) arguments. The #:<kw> kw-arg sequence is also supplied as keyword arguments to proc, where #:<kw> stands for any keyword.

The given proc must accept as many arguments as the number of vs plus length of lst, it must accept the supplied keyword arguments, and it must not require any other keyword arguments; otherwise, the exn:fail:contract exception is raised. The given proc is called in tail position with respect to the apply call.

Examples:

```
> (apply + '(1 2 3))
6
> (apply + 1 2 '(3))
6
> (apply + '())
0
> (apply sort (list (list '(2) '(1)) <) #:key car)
'((1) (2))
```

```
(compose proc ...) → procedure?
  proc : procedure?
(compose1 proc ...) → procedure?
  proc : procedure?
```

Returns a procedure that composes the given functions, applying the last proc first and the first proc last. The compose function allows the given functions to consume and produce any number of values, as long as each function produces as many values as the preceding function consumes, while compose1 restricts the internal value passing to a single value. In both cases, the input arity of the last function and the output arity of the first are unrestricted, and they become the corresponding arity of the resulting composition (including keyword arguments for the input side).

When no *proc* arguments are given, the result is `values`. When exactly one is given, it is returned.

Examples:

```
> ((compose1 - sqrt) 10)
-3.1622776601683795
> ((compose1 sqrt -) 10)
0+3.1622776601683795i
> ((compose list split-path) (bytes->path #"/a" 'unix))
'(#<path:/> #<path:a> #f)
```

Note that in many cases, `compose1` is preferred. For example, using `compose` with two library functions may lead to problems when one function is extended to return two values, and the preceding one has an optional input with different semantics. In addition, `compose1` may create faster compositions.

```
(procedure-rename proc name) → procedure?
  proc : procedure?
  name : symbol?
```

Returns a procedure that is like *proc*, except that its name as returned by `object-name` (and as printed for debugging) is *name*.

The given *name* is used for printing an error message if the resulting procedure is applied to the wrong number of arguments. In addition, if *proc* is an accessor or mutator produced by `struct`, `make-struct-field-accessor`, or `make-struct-field-mutator`, the resulting procedure also uses *name* when its (first) argument has the wrong type. More typically, however, *name* is not used for reporting errors, since the procedure name is typically hard-wired into an internal check.

```
(procedure->method proc) → procedure?
  proc : procedure?
```

Returns a procedure that is like *proc* except that, when applied to the wrong number of arguments, the resulting error hides the first argument as if the procedure had been compiled with the `'method-arity-error` syntax property.

```
(procedure-closure-contents-eq? proc1
                                proc2) → boolean?
  proc1 : procedure?
  proc2 : procedure?
```

Compares the contents of the closures of *proc1* and *proc2* for equality by comparing closure elements pointwise using `eq?`

4.17.1 Keywords and Arity

```
(keyword-apply proc
               kw-lst
               kw-val-lst
               v ...
               lst
               #:<kw> kw-arg ...) → any
  proc : procedure?
  kw-lst : (listof keyword?)
  kw-val-lst : list?
  v : any/c
  lst : list?
  kw-arg : any/c
```

§4.3.3 "The apply Function" in ***The Racket Guide*** introduces keyword-apply.

Like apply, but *kw-lst* and *kw-val-lst* supply by-keyword arguments in addition to the by-position arguments of the *v*s and *lst*, and in addition to the directly supplied keyword arguments in the #:<kw> *kw-arg* sequence, where #:<kw> stands for any keyword.

The given *kw-lst* must be sorted using keyword<?. No keyword can appear twice in *kw-lst* or in both kw-list and as a #:<kw>, otherwise, the exn:fail:contract exception is raised. The given *kw-val-lst* must have the same length as *kw-lst*, otherwise, the exn:fail:contract exception is raised. The given *proc* must accept all of the keywords in *kw-lst* plus the #:<kw>s, it must not require any other keywords, and it must accept as many by-position arguments as supplied via the *v*s and *lst*; otherwise, the exn:fail:contract exception is raised.

Examples:

```
(define (f x #:y y #:z [z 10])
  (list x y z))

> (keyword-apply f '(#:y) '(2) '(1))
'(1 2 10)
> (keyword-apply f '(#:y #:z) '(2 3) '(1))
'(1 2 3)
> (keyword-apply f #:z 7 '(#:y) '(2) '(1))
'(1 2 7)
```

```
(procedure-arity proc) → normalized-arity?
  proc : procedure?
```

Returns information about the number of by-position arguments accepted by *proc*. See also procedure-arity? and normalized-arity?.

```
(procedure-arity? v) → boolean?
  v : any/c
```

A valid arity `a` is one of the following:

- An exact non-negative integer, which means that the procedure accepts `a` arguments, only.
- A `arity-at-least` instance, which means that the procedure accepts `(arity-at-least-value a)` or more arguments.
- A list containing integers and `arity-at-least` instances, which means that the procedure accepts any number of arguments that can match one of the elements of `a`.

The result of `procedure-arity` is always normalized in the sense of `normalized-arity?`.

Examples:

```
> (procedure-arity cons)
2
> (procedure-arity list)
(arity-at-least 0)
> (arity-at-least? (procedure-arity list))
#t
> (arity-at-least-value (procedure-arity list))
0
> (arity-at-least-value (procedure-arity (lambda (x . y) x)))
1
> (procedure-arity (case-lambda [(x) 0] [(x y) 1]))
'(1 2)
```

```
(procedure-arity-includes? proc k [kws-ok?]) → boolean?
  proc : procedure?
  k : exact-nonnegative-integer?
  kws-ok? : any/c = #f
```

Returns `#t` if the procedure can accept `k` by-position arguments, `#f` otherwise. If `kws-ok?` is `#f`, the result is `#t` only if `proc` has no required keyword arguments.

Examples:

```
> (procedure-arity-includes? cons 2)
#t
> (procedure-arity-includes? display 3)
#f
> (procedure-arity-includes? (lambda (x #:y y) x) 1)
#f
> (procedure-arity-includes? (lambda (x #:y y) x) 1 #t)
#t
```

```
(procedure-reduce-arity proc arity) → procedure?
  proc : procedure?
  arity : procedure-arity?
```

Returns a procedure that is the same as *proc* (including the same name returned by `object-name`), but that accepts only arguments consistent with *arity*. In particular, when `procedure-arity` is applied to the generated procedure, it returns a value that is `equal?` to *arity*.

If the *arity* specification allows arguments that are not in `(procedure-arity proc)`, the `exn:fail:contract` exception is raised. If *proc* accepts keyword argument, either the keyword arguments must be all optional (and they are not accepted in by the arity-reduced procedure) or *arity* must be the empty list (which makes a procedure that cannot be called); otherwise, the `exn:fail:contract` exception is raised.

Examples:

```
> (define my+ (procedure-reduce-arity + 2))

> (my+ 1 2)
3
> (my+ 1 2 3)
+: arity mismatch;
 the expected number of arguments does not match the given
number
  expected: 2
  given: 3
  arguments...:
   1
   2
   3
```

```
(procedure-keywords proc) → (listof keyword?)
                            (or/c (listof keyword?) #f)
  proc : procedure?
```

Returns information about the keyword arguments required and accepted by a procedure. The first result is a list of distinct keywords (sorted by `keyword<?`) that are required when applying *proc*. The second result is a list of distinct accepted keywords (sorted by `keyword<?`), or `#f` to mean that any keyword is accepted. When the second result is a list, every element in the first list is also in the second list.

Examples:

```
> (procedure-keywords +)
```

```
'()
'()
> (procedure-keywords (lambda (#:tag t #:mode m) t))
'(#:mode #:tag)
'(#:mode #:tag)
> (procedure-keywords (lambda (#:tag t #:mode [m #f]) t))
'(#:tag)
'(#:mode #:tag)
```

```
(make-keyword-procedure proc [plain-proc]) → procedure?
  proc : (((listof keyword?) list?) () #:rest list? . ->* . any)
  plain-proc : procedure?
             = (lambda args (apply proc null null args))
```

Returns a procedure that accepts all keyword arguments (without requiring any keyword arguments).

When the procedure returned by `make-keyword-procedure` is called with keyword arguments, then *proc* is called; the first argument is a list of distinct keywords sorted by `keyword<?`, the second argument is a parallel list containing a value for each keyword, and the remaining arguments are the by-position arguments.

When the procedure returned by `make-keyword-procedure` is called without keyword arguments, then *plain-proc* is called—possibly more efficiently than dispatching through *proc*. Normally, *plain-proc* should have the same behavior as calling *proc* with empty lists as the first two arguments, but that correspondence is in no way enforced.

The result of `procedure-arity` and `object-name` on the new procedure is the same as for *plain-proc*. See also `procedure-reduce-keyword-arity` and `procedure-rename`.

Examples:

```
(define show
  (make-keyword-procedure (lambda (kws kw-args . rest)
                            (list kws kw-args rest))))

> (show 1)
'(() () (1))
> (show #:init 0 1 2 3 #:extra 4)
'((#:extra #:init) (4 0) (1 2 3))
(define show2
  (make-keyword-procedure (lambda (kws kw-args . rest)
                            (list kws kw-args rest))
                          (lambda args
                            (list->vector args))))
```

```
> (show2 1)
'#(1)
> (show2 #:init 0 1 2 3 #:extra 4)
'((#:extra #:init) (4 0) (1 2 3))
```

```
(procedure-reduce-keyword-arity proc
                                arity
                                required-kws
                                allowed-kws) → procedure?
  proc : procedure?
  arity : procedure-arity?
  required-kws : (listof keyword?)
  allowed-kws : (or/c (listof keyword?)
                      #f)
```

Like `procedure-reduce-arity`, but constrains the keyword arguments according to *required-kws* and *allowed-kws*, which must be sorted using `keyword<?` and contain no duplicates. If *allowed-kws* is `#f`, then the resulting procedure still accepts any keyword, otherwise the keywords in *required-kws* must be a subset of those in *allowed-kws*. The original *proc* must require no more keywords than the ones listed in *required-kws*, and it must allow at least the keywords in *allowed-kws* (or it must allow all keywords if *allowed-kws* is `#f`).

Examples:

```
(define orig-show
  (make-keyword-procedure (lambda (kws kw-args . rest)
                            (list kws kw-args rest))))

(define show (procedure-reduce-keyword-arity
              orig-show 3 '(#:init) '(#:extra #:init)))

> (show #:init 0 1 2 3 #:extra 4)
'((#:extra #:init) (4 0) (1 2 3))
> (show 1)
...t/private/kw.rkt:211:14: arity mismatch;
 the expected number of arguments does not match the given
number
  expected: 3 plus an argument with keyword #:init plus an
optional argument with keyword #:extra
  given: 1
  arguments...:
   1
> (show #:init 0 1 2 3 #:extra 4 #:more 7)
application: procedure does not expect an argument with
given keyword
```

```
procedure: ...t/private/kw.rkt:211:14
given keyword: #:more
arguments...:
 1
 2
 3
 #:extra 4
 #:init 0
 #:more 7
```

```
(struct arity-at-least (value)
    #:extra-constructor-name make-arity-at-least)
  value : exact-nonnegative-integer?
```

A structure type used for the result of `procedure-arity`. See also `procedure-arity?`.

```
prop:procedure : struct-type-property?
```

A structure type property to identify structure types whose instances can be applied as procedures. In particular, when `procedure?` is applied to the instance, the result will be `#t`, and when an instance is used in the function position of an application expression, a procedure is extracted from the instance and used to complete the procedure call.

If the `prop:procedure` property value is an exact non-negative integer, it designates a field within the structure that should contain a procedure. The integer must be between `0` (inclusive) and the number of non-automatic fields in the structure type (exclusive, not counting supertype fields). The designated field must also be specified as immutable, so that after an instance of the structure is created, its procedure cannot be changed. (Otherwise, the arity and name of the instance could change, and such mutations are generally not allowed for procedures.) When the instance is used as the procedure in an application expression, the value of the designated field in the instance is used to complete the procedure call. (This procedure can be another structure that acts as a procedure; the immutability of procedure fields disallows cycles in the procedure graph, so that the procedure call will eventually continue with a non-structure procedure.) That procedure receives all of the arguments from the application expression. The procedure's name (see `object-name`), arity (see `procedure-arity`), and keyword protocol (see `procedure-keywords`) are also used for the name, arity, and keyword protocol of the structure. If the value in the designated field is not a procedure, then the instance behaves like `(case-lambda)` (i.e., a procedure which does not accept any number of arguments). See also `procedure-extract-target`.

Providing an integer `proc-spec` argument to `make-struct-type` is the same as both supplying the value with the `prop:procedure` property and designating the field as immutable (so that a property binding or immutable designation is redundant and disallowed).

Examples:

```
> (struct annotated-proc (base note)
```

```
    #:property prop:procedure
               (struct-field-index base))

> (define plus1 (annotated-proc
                 (lambda (x) (+ x 1))
                 "adds 1 to its argument"))

> (procedure? plus1)
#t
> (annotated-proc? plus1)
#t
> (plus1 10)
11
> (annotated-proc-note plus1)
"adds 1 to its argument"
```

When the `prop:procedure` value is a procedure, it should accept at least one non-keyword argument. When an instance of the structure is used in an application expression, the property-value procedure is called with the instance as the first argument. The remaining arguments to the property-value procedure are the arguments from the application expression (including keyword arguments). Thus, if the application expression provides five non-keyword arguments, the property-value procedure is called with six non-keyword arguments. The name of the instance (see `object-name`) and its keyword protocol (see `procedure-keywords`) are unaffected by the property-value procedure, but the instance's arity is determined by subtracting one from every possible non-keyword argument count of the property-value procedure. If the property-value procedure cannot accept at least one argument, then the instance behaves like `(case-lambda)`.

Providing a procedure `proc-spec` argument to `make-struct-type` is the same as supplying the value with the `prop:procedure` property (so that a specific property binding is disallowed).

Examples:

```
> (struct fish (weight color)
    #:mutable
    #:property
    prop:procedure
    (lambda (f n)
      (let ([w (fish-weight f)])
        (set-fish-weight! f (+ n w)))))

> (define wanda (fish 12 'red))

> (fish? wanda)
#t
> (procedure? wanda)
```

```
#t
> (fish-weight wanda)
12
> (for-each wanda '(1 2 3))

> (fish-weight wanda)
18
```

If the value supplied for the prop:procedure property is not an exact non-negative integer or a procedure, the exn:fail:contract exception is raised.

(procedure-struct-type? *type*) → boolean?
  type : struct-type?

Returns #t if instances of the structure type represented by *type* are procedures (according to procedure?), #f otherwise.

(procedure-extract-target *proc*) → (or/c #f procedure?)
  proc : procedure?

If *proc* is an instance of a structure type with property prop:procedure, and if the property value indicates a field of the structure, and if the field value is a procedure, then procedure-extract-target returns the field value. Otherwise, the result is #f.

When a prop:procedure property value is a procedure, the procedure is *not* returned by procedure-extract-target. Such a procedure is different from one accessed through a structure field, because it consumes an extra argument, which is always the structure that was applied as a procedure. Keeping the procedure private ensures that is it always called with a suitable first argument.

prop:arity-string : struct-type-property?

A structure type property that is used for reporting arity-mismatch errors when a structure type with the prop:procedure property is applied to the wrong number of arguments. The value of the prop:arity-string property must be a procedure that takes a single argument, which is the misapplied structure, and returns a string. The result string is used after the word "expects," and it is followed in the error message by the number of actual arguments.

Arity-mismatch reporting automatically uses procedure-extract-target when the prop:arity-string property is not associated with a procedure structure type.

Examples:

```
> (struct evens (proc)
    #:property prop:procedure (struct-field-index proc)
```

```
    #:property prop:arity-string
    (lambda (p)
      "an even number of arguments"))

> (define pairs
    (evens
     (case-lambda
      [() null]
      [(a b . more)
       (cons (cons a b)
             (apply pairs more))])))

> (pairs 1 2 3 4)
'((1 . 2) (3 . 4))
> (pairs 5)
#<procedure>: arity mismatch;
 the expected number of arguments does not match the given
number
  expected: an even number of arguments
  given: 1
  arguments...:
   5
```

```
prop:checked-procedure : struct-type-property?
```

A structure type property that is used with `checked-procedure-check-and-extract`, which is a hook to allow the compiler to improve the performance of keyword arguments. The property can only be attached to a structure type without a supertype and with at least two fields.

```
(checked-procedure-check-and-extract type
                                     v
                                     proc
                                     v1
                                     v2)  → any/c
  type : struct-type?
  v : any/c
  proc : (any/c any/c any/c . -> . any/c)
  v1 : any/c
  v2 : any/c
```

Extracts a value from *v* if it is an instance of *type*, which must have the property `prop:checked-procedure`. If *v* is such an instance, then the first field of *v* is extracted and applied to *v1* and *v2*; if the result is a true value, the result is the value of the second field of *v*.

If *v* is not an instance of *type*, or if the first field of *v* applied to *v1* and *v2* produces #f, then *proc* is applied to *v*, *v1*, and *v2*, and its result is returned by checked-procedure-check-and-extract.

4.17.2 Reflecting on Primitives

A *primitive procedure* is a built-in procedure that is implemented in low-level language. Not all procedures of racket/base are primitives, but many are. The distinction is mainly useful to other low-level code.

```
(primitive? v) → boolean?
  v : any/c
```

Returns #t if *v* is a primitive procedure, #f otherwise.

```
(primitive-closure? v) → boolean
  v : any/c
```

Returns #t if *v* is internally implemented as a primitive closure rather than a simple primitive procedure, #f otherwise.

```
(primitive-result-arity prim) → procedure-arity?
  prim : primitive?
```

Returns the arity of the result of the primitive procedure *prim* (as opposed to the procedure's input arity as returned by procedure-arity). For most primitives, this procedure returns 1, since most primitives return a single value when applied.

4.17.3 Additional Higher-Order Functions

```
(require racket/function)
```
package: base

The bindings documented in this section are provided by the racket/function and racket libraries, but not racket/base.

```
(identity v) → any/c
  v : any/c
```

Returns *v*.

```
(const v) → procedure?
  v : any
```

Returns a procedure that accepts any arguments (including keyword arguments) and returns *v*.

Examples:

```
> ((const 'foo) 1 2 3)
'foo
> ((const 'foo))
'foo
```

```
(thunk  body ...+)
(thunk* body ...+)
```

The `thunk` form creates a nullary function that evaluates the given body. The `thunk*` form is similar, except that the resulting function accepts any arguments (including keyword arguments).

Examples:

```
(define th1 (thunk (define x 1) (printf "~a\n" x)))

> (th1)
1

> (th1 'x)
th1: arity mismatch;
 the expected number of arguments does not match the given
number
  expected: 0
  given: 1
  arguments...:
   'x
> (th1 #:y 'z)
application: procedure does not accept keyword arguments
  procedure: th1
  arguments...:
   #:y 'z
(define th2 (thunk* (define x 1) (printf "~a\n" x)))

> (th2)
1

> (th2 'x)
1

> (th2 #:y 'z)
1
```

```
(negate proc) → procedure?
  proc : procedure?
```

Returns a procedure that is just like *proc*, except that it returns the `not` of *proc*'s result.

Examples:

```
> (filter (negate symbol?) '(1 a 2 b 3 c))
'(1 2 3)
> (map (negate =) '(1 2 3) '(1 1 1))
'(#f #t #t)
```

```
(curry proc) → procedure?
  proc : procedure?
(curry proc v ...+) → any/c
  proc : procedure?
  v : any/c
```

Returns a procedure that is a curried version of *proc*. When the resulting procedure is first applied, unless it is given the maximum number of arguments that it can accept, the result is a procedure to accept additional arguments.

Examples:

```
> ((curry list) 1 2)
#<procedure:curried>
> ((curry cons) 1)
#<procedure:curried>
> ((curry cons) 1 2)
'(1 . 2)
```

After the first application of the result of `curry`, each further application accumulates arguments until an acceptable number of arguments have been accumulated, at which point the original *proc* is called.

Examples:

```
> (((curry list) 1 2) 3)
'(1 2 3)
> (((curry list) 1) 3)
'(1 3)
> ((((curry foldl) +) 0) '(1 2 3))
6
```

A function call `(curry proc v ...)` is equivalent to `((curry proc) v ...)`. In other words, `curry` itself is curried.

The curry function provides limited support for keyworded functions: only the curry call itself can receive keyworded arguments to be propagated eventually to *proc*.

Examples:

```
> (map ((curry +) 10) '(1 2 3))
'(11 12 13)
> (map (curry + 10) '(1 2 3))
'(11 12 13)
> (map (compose (curry * 2) (curry + 10)) '(1 2 3))
'(22 24 26)
> (define foo (curry (lambda (x y z) (list x y z))))

> (foo 1 2 3)
'(1 2 3)
> (((((foo) 1) 2)) 3)
'(1 2 3)
```

```
(curryr proc) → procedure?
  proc : procedure?
(curryr proc v ...+) → any/c
  proc : procedure?
  v : any/c
```

Like curry, except that the arguments are collected in the opposite direction: the first step collects the rightmost group of arguments, and following steps add arguments to the left of these.

Example:

```
> (map (curryr list 'foo) '(1 2 3))
'((1 foo) (2 foo) (3 foo))
```

```
(normalized-arity? arity) → boolean?
  arity : any/c
```

A normalized arity has one of the following forms:

- the empty list;
- an exact non-negative integer;
- an arity-at-least instance;
- a list of two or more strictly increasing, exact non-negative integers; or
- a list of one or more strictly increasing, exact non-negative integers followed by a single arity-at-least instance whose value is greater than the preceding integer by at least 2.

Every normalized arity is a valid procedure arity and satisfies `procedure-arity?`. Any two normalized arity values that are `arity=?` must also be `equal?`.

Examples:

```
> (normalized-arity? (arity-at-least 1))
#t
> (normalized-arity? (list (arity-at-least 1)))
#f
> (normalized-arity? (list 0 (arity-at-least 2)))
#t
> (normalized-arity? (list (arity-at-least 2) 0))
#f
> (normalized-arity? (list 0 2 (arity-at-least 3)))
#f
```

```
(normalize-arity arity)
 → (and/c normalized-arity? (lambda (x) (arity=? x arity)))
  arity : procedure-arity?
```

Produces a normalized form of *arity*. See also `normalized-arity?` and `arity=?`.

Examples:

```
> (normalize-arity 1)
1
> (normalize-arity (list 1))
1
> (normalize-arity (arity-at-least 2))
(arity-at-least 2)
> (normalize-arity (list (arity-at-least 2)))
(arity-at-least 2)
> (normalize-arity (list 1 (arity-at-least 2)))
(arity-at-least 1)
> (normalize-arity (list (arity-at-least 2) 1))
(arity-at-least 1)
> (normalize-arity (list (arity-at-least 2) 3))
(arity-at-least 2)
> (normalize-arity (list 3 (arity-at-least 2)))
(arity-at-least 2)
> (normalize-arity (list (arity-at-least 6) 0 2 (arity-at-
least 4)))
(list 0 2 (arity-at-least 4))
```

```
(arity=? a b) → boolean?
  a : procedure-arity?
  b : procedure-arity?
```

Returns `#true` if procedures with arity `a` and `b` accept the same numbers of arguments, and `#false` otherwise. Equivalent to both `(and (arity-includes? a b) (arity-includes? b a))` and `(equal? (normalize-arity a) (normalize-arity b))`.

Examples:

```
> (arity=? 1 1)
#t
> (arity=? (list 1) 1)
#t
> (arity=? 1 (list 1))
#t
> (arity=? 1 (arity-at-least 1))
#f
> (arity=? (arity-at-least 1) 1)
#f
> (arity=? 1 (arity-at-least 1))
#f
> (arity=? (arity-at-least 1) (list 1 (arity-at-least 2)))
#t
> (arity=? (list 1 (arity-at-least 2)) (arity-at-least 1))
#t
> (arity=? (arity-at-least 1) (list 1 (arity-at-least 3)))
#f
> (arity=? (list 1 (arity-at-least 3)) (arity-at-least 1))
#f
> (arity=? (list 0 1 2 (arity-at-least 3)) (list (arity-at-
least 0)))
#t
> (arity=? (list (arity-at-least 0)) (list 0 1 2 (arity-at-
least 3)))
#t
> (arity=? (list 0 2 (arity-at-least 3)) (list (arity-at-
least 0)))
#f
> (arity=? (list (arity-at-least 0)) (list 0 2 (arity-at-
least 3)))
#f
```

```
(arity-includes? a b) → boolean?
  a : procedure-arity?
  b : procedure-arity?
```

Returns `#true` if procedures with arity `a` accept any number of arguments that procedures with arity `b` accept.

Examples:

```
> (arity-includes? 1 1)
#t
> (arity-includes? (list 1) 1)
#t
> (arity-includes? 1 (list 1))
#t
> (arity-includes? 1 (arity-at-least 1))
#f
> (arity-includes? (arity-at-least 1) 1)
#t
> (arity-includes? 1 (arity-at-least 1))
#f
> (arity-includes? (arity-at-least 1) (list 1 (arity-at-least 2)))
#t
> (arity-includes? (list 1 (arity-at-least 2)) (arity-at-least 1))
#t
> (arity-includes? (arity-at-least 1) (list 1 (arity-at-least 3)))
#t
> (arity-includes? (list 1 (arity-at-least 3)) (arity-at-least 1))
#f
> (arity-includes? (list 0 1 2 (arity-at-least 3)) (list (arity-
at-least 0)))
#t
> (arity-includes? (list (arity-at-least 0)) (list 0 1 2 (arity-
at-least 3)))
#t
> (arity-includes? (list 0 2 (arity-at-least 3)) (list (arity-at-
least 0)))
#f
> (arity-includes? (list (arity-at-least 0)) (list 0 2 (arity-at-
least 3)))
#t
```

4.18 Void

The constant `#<void>` is returned by most forms and procedures that have a side-effect and no useful result.

The `#<void>` value is always `eq?` to itself.

```
(void? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is the constant `#<void>`, `#f` otherwise.

```
(void v ...) → void?
  v : any/c
```

Returns the constant `#<void>`. Each `v` argument is ignored.

4.19 Undefined

`(require racket/undefined)` package: `base`

The bindings documented in this section are provided by the `racket/undefined` library, not `racket/base` or `racket`.

The constant `undefined` can be used as a placeholder value for a value to be installed later, especially for cases where premature access of the value is either difficult or impossible to detect or prevent.

The `undefined` value is always `eq?` to itself.

Added in version 6.0.0.6 of package `base`.

```
undefined : any/c
```

The “undefined” constant.

5 Structures

§5 “Programmer-Defined Datatypes” in ***The Racket Guide*** introduces structure types via `struct`.

A *structure type* is a record datatype composing a number of *fields*. A *structure*, an instance of a structure type, is a first-class value that contains a value for each field of the structure type. A structure instance is created with a type-specific constructor procedure, and its field values are accessed and changed with type-specific accessor and mutator procedures. In addition, each structure type has a predicate procedure that answers `#t` for instances of the structure type and `#f` for any other value.

A structure type's fields are essentially unnamed, though names are supported for error-reporting purposes. The constructor procedure takes one value for each field of the structure type, except that some of the fields of a structure type can be *automatic fields*; the automatic fields are initialized to a constant that is associated with the structure type, and the corresponding arguments are omitted from the constructor procedure. All automatic fields in a structure type follow the non-automatic fields.

A structure type can be created as a *structure subtype* of an existing base structure type. An instance of a structure subtype can always be used as an instance of the base structure type, but the subtype gets its own predicate procedure, and it may have its own fields in addition to the fields of the base type.

A structure subtype “inherits” the fields of its base type. If the base type has m fields, and if n fields are specified for the new structure subtype, then the resulting structure type has $m+n$ fields. The value for automatic fields can be different in a subtype than in its base type.

If m' of the original m fields are non-automatic (where $m'<m$), and n' of the new fields are non-automatic (where $n'<n$), then $m'+n'$ field values must be provided to the subtype's constructor procedure. Values for the first m fields of a subtype instance are accessed with selector procedures for the original base type (or its supertypes), and the last n are accessed with subtype-specific selectors. Subtype-specific accessors and mutators for the first m fields do not exist.

The `struct` form and `make-struct-type` procedure typically create a new structure type, but they can also access *prefab* (i.e., previously fabricated) structure types that are globally shared, and whose instances can be parsed and written by the default reader (see §1.3 “The Reader”) and printer (see §1.4 “The Printer”). Prefab structure types can inherit only from other prefab structure types, and they cannot have guards (see §5.2 “Creating Structure Types”) or properties (see §5.3 “Structure Type Properties”). Exactly one prefab structure type exists for each combination of name, supertype, field count, automatic field count, automatic field value (when there is at least one automatic field), and field mutability.

§13.9 “Serialization” also provides information on reading and writing structures.

Two structure values are `eqv?` if and only if they are `eq?`. Two structure values are `equal?` if they are `eq?`. By default, two structure values are also `equal?` if they are instances of the same structure type, no fields are opaque, and the results of applying `struct->vector` to the structs are `equal?`. (Consequently, `equal?` testing for structures may depend on the current inspector.) A structure type can override the default `equal?` definition through the

gen:equal+hash generic interface.

5.1 Defining Structure Types: struct

§5 “Programmer-Defined Datatypes” in ***The Racket Guide*** introduces struct.

```
(struct id maybe-super (field ...)
        struct-option ...)

  maybe-super =
              | super-id

        field = field-id
              | [field-id field-option ...]

struct-option = #:mutable
              | #:super super-expr
              | #:inspector inspector-expr
              | #:auto-value auto-expr
              | #:guard guard-expr
              | #:property prop-expr val-expr
              | #:transparent
              | #:prefab
              | #:constructor-name constructor-id
              | #:extra-constructor-name constructor-id
              | #:reflection-name symbol-expr
              | #:methods gen:name method-defs
              | #:omit-define-syntaxes
              | #:omit-define-values

 field-option = #:mutable
              | #:auto

  method-defs = (definition ...)

  gen:name : identifier?
```

Creates a new structure type (or uses a pre-existing structure type if #:prefab is specified), and binds transformers and variables related to the structure type.

A struct form with *n* field s defines up to 4+2*n* names:

- struct:id, a *structure type descriptor* value that represents the structure type.
- *constructor-id* (which defaults to *id*), a *constructor* procedure that takes *m* arguments and returns a new instance of the structure type, where *m* is the number of field s that do not include an #:auto option.

- *id*, a transformer binding that encapsulates information about the structure type declaration. This binding is used to define subtypes, and it also works with the `shared` and `match` forms. For detailed information about the binding of *id*, see §5.7 “Structure Type Transformer Binding”.

  The *constructor-id* and *id* can be the same, in which case *id* performs both roles. In that case, the expansion of *id* as an expression produces an otherwise inaccessible identifier that is bound to the constructor procedure; the expanded identifier has a `'constructor-for` property whose value is an identifier that is `free-identifier=?` to *id* as well as a syntax property accessible via `syntax-procedure-alias-property` with an identifier that is `free-identifier=?` to *id*.

- *id*`?`, a *predicate* procedure that returns `#t` for instances of the structure type (constructed by *constructor-id* or the constructor for a subtype) and `#f` for any other value.

- *id-field-id*, for each *field*; an *accessor* procedure that takes an instance of the structure type and extracts the value for the corresponding field.

- `set-`*id-field-id*`!`, for each *field* that includes a `#:mutable` option, or when the `#:mutable` option is specified as a *struct-option*; a *mutator* procedure that takes an instance of the structure type and a new field value. The structure is destructively updated with the new value, and `#<void>` is returned.

If *super-id* is provided, it must have a transformer binding of the same sort bound to *id* (see §5.7 “Structure Type Transformer Binding”), and it specifies a supertype for the structure type. Alternately, the `#:super` option can be used to specify an expression that must produce a structure type descriptor. See §5 “Structures” for more information on structure subtypes and supertypes. If both *super-id* and `#:super` are provided, a syntax error is reported.

Examples:

```
> (struct document (author title content))

> (struct book document (publisher))

> (struct paper (journal) #:super struct:document)
```

If the `#:mutable` option is specified for an individual field, then the field can be mutated in instances of the structure type, and a mutator procedure is bound. Supplying `#:mutable` as a *struct-option* is the same as supplying it for all *field*`s`. If `#:mutable` is specified as both a *field-option* and *struct-option*, a syntax error is reported.

Examples:

```
> (struct cell ([content #:mutable]) #:transparent)
```

```
> (define a-cell (cell 0))

> (set-cell-content! a-cell 1)
```

The `#:inspector`, `#:auto-value`, and `#:guard` options specify an inspector, value for automatic fields, and guard procedure, respectively. See `make-struct-type` for more information on these attributes of a structure type. The `#:property` option, which is the only one that can be supplied multiple times, attaches a property value to the structure type; see §5.3 “Structure Type Properties” for more information on properties. The `#:transparent` option is a shorthand for `#:inspector #f`.

Examples:

```
> (struct point (x y) #:inspector #f)

> (point 3 5)
(point 3 5)
> (struct celcius (temp)
    #:guard (λ (temp name)
              (unless (and (real? temp) (>= temp -273.15))
                (error "not a valid temperature"))
              temp))

> (celcius -275)
not a valid temperature
```

The `#:prefab` option obtains a prefab (pre-defined, globally shared) structure type, as opposed to creating a new structure type. Such a structure type is inherently transparent and cannot have a guard or properties, so using `#:prefab` with `#:transparent`, `#:inspector`, `#:guard`, or `#:property` is a syntax error. If a supertype is specified, it must also be a prefab structure type.

Use the `prop:procedure` property to implement an applicable structure, use `prop:evt` to create a structure type whose instances are synchronizable events, and so on. By convention, property names start with `prop:`.

Examples:

```
> (struct prefab-point (x y) #:prefab)

> (prefab-point 1 2)
'#s(prefab-point 1 2)
> (prefab-point? #s(prefab-point 1 2))
#t
```

If *`constructor-id`* is supplied, then the transformer binding of *`id`* records *`constructor-id`* as the constructor binding; as a result, for example, `struct-out` includes *`constructor-id`* as an export. If *`constructor-id`* is supplied via `#:extra-constructor-name` and it is not *`id`*, applying `object-name` on the constructor produces

the symbolic form of *id* rather than *constructor-id*. If *constructor-id* is supplied via #:constructor-name and it is not the same as *id*, then *id* does not serve as a constructor, and object-name on the constructor produces the symbolic form of *constructor-id*. Only one of #:extra-constructor-name and #:constructor-name can be provided within a struct form.

Examples:

```
> (struct color (r g b) #:constructor-name -color)

> (struct rectangle (w h color) #:extra-constructor-name rect)

> (rectangle 13 50 (-color 192 157 235))
#<rectangle>
> (rect 50 37 (-color 35 183 252))
#<rectangle>
```

If #:reflection-name *symbol-expr* is provided, then *symbol-expr* must produce a symbol that is used to identify the structure type in reflective operations such as struct-type-info. It corresponds to the first argument of make-struct-type. Structure printing uses the reflective name, as do the various procedures that are bound by struct.

Examples:

```
> (struct circle (radius) #:reflection-name '<circle>)

> (circle 15)
#<|<circle>|>
> (circle-radius "bad")
<circle>-radius: contract violation
  expected: <circle>?
  given: "bad"
```

If #:methods *gen:name method-defs* is provided, then *gen:name* must be a transformer binding for the static information about a generic interface produced by define-generics. The *method-defs* define the methods of the *gen:name* interface. A define/generic form or auxiliary definitions and expressions may also appear in *method-defs*.

Examples:

```
> (struct constant-stream (val)
    #:methods gen:stream
    [(define (stream-empty? stream) #f)
     (define (stream-first stream)
       (constant-stream-val stream))
     (define (stream-rest stream) stream)])
```

```
> (stream-ref (constant-stream 'forever) 0)
'forever
> (stream-ref (constant-stream 'forever) 50)
'forever
```

If the `#:omit-define-syntaxes` option is supplied, then *`id`* is not bound as a transformer. If the `#:omit-define-values` option is supplied, then none of the usual variables are bound, but *`id`* is bound. If both are supplied, then the `struct` form is equivalent to `(begin)`.

Examples:

```
> (struct square (side) #:omit-define-syntaxes)

> (match (square 5)
    ; fails to match because syntax is omitted
    [(struct square x) x])
eval:26:0: match: square does not refer to a structure
definition
  at: square
  in: (struct square x)
> (struct ellipse (width height) #:omit-define-values)

> ellipse-width
ellipse-width: undefined;
 cannot reference undefined identifier
```

If `#:auto` is supplied as a *`field-option`*, then the constructor procedure for the structure type does not accept an argument corresponding to the field. Instead, the structure type's automatic value is used for the field, as specified by the `#:auto-value` option, or as defaults to `#f` when `#:auto-value` is not supplied. The field is mutable (e.g., through reflective operations), but a mutator procedure is bound only if `#:mutable` is specified.

If a *`field`* includes the `#:auto` option, then all fields after it must also include `#:auto`, otherwise a syntax error is reported. If any *`field-option`* or *`struct-option`* keyword is repeated, other than `#:property`, a syntax error is reported.

Examples:

```
> (struct posn (x y [z #:auto])
    #:auto-value 0
    #:transparent)

> (posn 1 2)
(posn 1 2 0)
> (posn? (posn 1 2))
#t
> (posn-y (posn 1 2))
2
```

```
(struct color-posn posn (hue) #:mutable)
(define cp (color-posn 1 2 "blue"))

> (color-posn-hue cp)
"blue"
> cp
(color-posn 1 2 0 ...)
> (set-posn-z! cp 3)
set-posn-z!: undefined;
 cannot reference undefined identifier
```

For serialization, see `define-serializable-struct`.

```
(struct-field-index field-id)
```

This form can only appear as an expression within a `struct` form; normally, it is used with `#:property`, especially for a property like `prop:procedure`. The result of a `struct-field-index` expression is an exact, non-negative integer that corresponds to the position within the structure declaration of the field named by *`field-id`*.

Examples:

```
> (struct mood-procedure (base rating)
    #:property prop:procedure (struct-field-index base))

(define happy+ (mood-procedure add1 10))

> (happy+ 2)
3
> (mood-procedure-rating happy+)
10
```

```
(define-struct id-maybe-super (field ...)
               struct-option ...)

id-maybe-super = id
               | (id super-id)
```

Like `struct`, except that the syntax for supplying a *`super-id`* is different, and a *`constructor-id`* that has a `make-` prefix on *`id`* is implicitly supplied via `#:extra-constructor-name` if neither `#:extra-constructor-name` nor `#:constructor-name` is provided.

This form is provided for backwards compatibility; `struct` is preferred.

Examples:

```
(define-struct posn (x y [z #:auto])
   #:auto-value 0
   #:transparent)

> (make-posn 1 2)
(posn 1 2 0)
> (posn? (make-posn 1 2))
#t
> (posn-y (make-posn 1 2))
2
```

```
(define-struct/derived (id . rest-form)
  id-maybe-super (field ...) struct-option ...)
```

The same as `define-struct`, but with an extra `(id . rest-form)` sub-form that is treated as the overall form for syntax-error reporting and otherwise ignored. The only constraint on the sub-form for error reporting is that it starts with `id`. The `define-struct/derived` form is intended for use by macros that expand to `define-struct`.

Examples:

```
> (define-syntax (define-xy-struct stx)
    (syntax-case stx ()
     [(ds name . rest)
      (with-syntax ([orig stx])
        #'(define-struct/derived orig name (x y) . rest))]))

> (define-xy-struct posn)

> (posn-x (make-posn 1 2))
1
> (define-xy-struct posn #:mutable)

> (set-posn-x! (make-posn 1 2) 0)

; this next line will cause an error due to a bad keyword
> (define-xy-struct posn #:bad-option)
eval:51:0: define-xy-struct: unrecognized
struct-specification keyword
  at: #:bad-option
  in: (define-xy-struct posn #:bad-option)
```

5.2 Creating Structure Types

```
(make-struct-type  name
                   super-type
                   init-field-cnt
                   auto-field-cnt
                  [auto-v
                   props
                   inspector
                   proc-spec
                   immutables
                   guard
                   constructor-name])
   struct-type?
   struct-constructor-procedure?
 → struct-predicate-procedure?
   struct-accessor-procedure?
   struct-mutator-procedure?
  name : symbol?
  super-type : (or/c struct-type? #f)
  init-field-cnt : exact-nonnegative-integer?
  auto-field-cnt : exact-nonnegative-integer?
  auto-v : any/c = #f
  props : (listof (cons/c struct-type-property?
                          any/c))                  = null
  inspector : (or/c inspector? #f 'prefab) = (current-inspector)
              (or/c procedure?
  proc-spec :       exact-nonnegative-integer? = #f
                    #f)
  immutables : (listof exact-nonnegative-integer?) = null
  guard : (or/c procedure? #f) = #f
  constructor-name : (or/c symbol? #f) = #f
```

Creates a new structure type, unless *inspector* is `'prefab`, in which case `make-struct-type` accesses a prefab structure type. The *name* argument is used as the type name. If *super-type* is not `#f`, the resulting type is a subtype of the corresponding structure type.

The resulting structure type has *init-field-cnt*+*auto-field-cnt* fields (in addition to any fields from *super-type*), but only *init-field-cnt* constructor arguments (in addition to any constructor arguments from *super-type*). The remaining fields are initialized with *auto-v*. The total field count (including *super-type* fields) must be no more than 32768.

The *props* argument is a list of pairs, where the `car` of each pair is a structure type property descriptor, and the `cdr` is an arbitrary value. A property can be specified multiple times in *props* (including properties that are automatically added by properties that are directly included in *props*) only if the associated values are `eq?`, otherwise the `exn:fail:contract` exception is raised. See §5.3 “Structure Type Properties” for more information about prop-

erties. When *inspector* is 'prefab, then *props* must be null.

The *inspector* argument normally controls access to reflective information about the structure type and its instances; see §14.9 “Structure Inspectors” for more information. If *inspector* is 'prefab, then the resulting prefab structure type and its instances are always transparent.

If *proc-spec* is an integer or procedure, instances of the structure type act as procedures. See prop:procedure for further information. Providing a non-#f value for *proc-spec* is the same as pairing the value with prop:procedure at the end of *props*, plus including *proc-spec* in *immutables* when *proc-spec* is an integer.

The *immutables* argument provides a list of field positions. Each element in the list must be unique, otherwise exn:fail:contract exception is raised. Each element must also fall in the range 0 (inclusive) to *init-field-cnt* (exclusive), otherwise exn:fail:contract exception is raised.

The *guard* argument is either a procedure of $n+1$ arguments or #f, where n is the number of arguments for the new structure type's constructor (i.e., *init-field-cnt* plus constructor arguments implied by *super-type*, if any). If *guard* is a procedure, then the procedure is called whenever an instance of the type is constructed, or whenever an instance of a subtype is created. The arguments to *guard* are the values provided for the structure's first n fields, followed by the name of the instantiated structure type (which is *name*, unless a subtype is instantiated). The *guard* result must be n values, which become the actual values for the structure's fields. The *guard* can raise an exception to prevent creation of a structure with the given field values. If a structure subtype has its own guard, the subtype guard is applied first, and the first n values produced by the subtype's guard procedure become the first n arguments to *guard*. When *inspector* is 'prefab, then *guard* must be #f.

If *constructor-name* is not #f, it is used as the name of the generated constructor procedure as returned by object-name or in the printed form of the constructor value.

The result of make-struct-type is five values:

- a structure type descriptor,
- a constructor procedure,
- a predicate procedure,
- an accessor procedure, which consumes a structure and a field index between 0 (inclusive) and *init-field-cnt*+*auto-field-cnt* (exclusive), and
- a mutator procedure, which consumes a structure, a field index, and a field value.

Examples:

```
> (define-values (struct:a make-a a? a-ref a-set!)
    (make-struct-type 'a #f 2 1 'uninitialized))
```

```
> (define an-a (make-a 'x 'y))

> (a-ref an-a 1)
'y
> (a-ref an-a 2)
'uninitialized
> (define a-first (make-struct-field-accessor a-ref 0))

> (a-first an-a)
'x
> (define-values (struct:b make-b b? b-ref b-set!)
    (make-struct-type 'b struct:a 1 2 'b-uninitialized))

> (define a-b (make-b 'x 'y 'z))

> (a-ref a-b 1)
'y
> (a-ref a-b 2)
'uninitialized
> (b-ref a-b 0)
'z
> (b-ref a-b 1)
'b-uninitialized
> (b-ref a-b 2)
'b-uninitialized

> (define-values (struct:c make-c c? c-ref c-set!)
    (make-struct-type
     'c struct:b 0 0 #f null (make-inspector) #f null
     ; guard checks for a number, and makes it inexact
     (lambda (a1 a2 b1 name)
       (unless (number? a2)
         (error (string->symbol (format "make-~a" name))
                "second field must be a number"))
       (values a1 (exact->inexact a2) b1))))

> (make-c 'x 'y 'z)
make-c: second field must be a number
> (define a-c (make-c 'x 2 'z))

> (a-ref a-c 1)
2.0

> (define p1 #s(p a b c))
```

```
> (define-values (struct:p make-p p? p-ref p-set!)
    (make-struct-type 'p #f 3 0 #f null 'prefab #f '(0 1 2)))

> (p? p1)
#t
> (p-ref p1 0)
'a
> (make-p 'x 'y 'z)
'#s(p x y z)
```

```
(make-struct-field-accessor accessor-proc
                            field-pos
                            [field-name]) → procedure?
  accessor-proc : struct-accessor-procedure?
  field-pos : exact-nonnegative-integer?
  field-name : (or/c symbol? #f)
             = (symbol->string (format "field~a" field-pos))
```

Returns a field accessor that is equivalent to `(lambda (s) (accessor-proc s field-pos))`. The *accessor-proc* must be an accessor returned by `make-struct-type`. The name of the resulting procedure for debugging purposes is derived from *field-name* and the name of *accessor-proc*'s structure type if *field-name* is a symbol.

For examples, see `make-struct-type`.

```
(make-struct-field-mutator mutator-proc
                           field-pos
                           [field-name]) → procedure?
  mutator-proc : struct-mutator-procedure?
  field-pos : exact-nonnegative-integer?
  field-name : (or/c symbol? #f)
             = (symbol->string (format "field~a" field-pos))
```

Returns a field mutator that is equivalent to `(lambda (s v) (mutator-proc s field-pos v))`. The *mutator-proc* must be a mutator returned by `make-struct-type`. The name of the resulting procedure for debugging purposes is derived from *field-name* and the name of *mutator-proc*'s structure type if *field-name* is a symbol.

For examples, see `make-struct-type`.

5.3 Structure Type Properties

§5.4 "Generic Interfaces" provide a high-level API on top of structure type properties.

A *structure type property* allows per-type information to be associated with a structure type (as opposed to per-instance information associated with a structure value). A property value

is associated with a structure type through the make-struct-type procedure (see §5.2 “Creating Structure Types”) or through the #:property option of struct. Subtypes inherit the property values of their parent types, and subtypes can override an inherited property value with a new value.

```
(make-struct-type-property  name
                           [guard
                            supers
                            can-impersonate?])
   struct-type-property?
 → procedure?
   procedure?
  name : symbol?
  guard : (or/c procedure? #f 'can-impersonate) = #f
  supers : (listof (cons/c struct-type-property?
                           (any/c . -> . any/c))) = null
  can-impersonate? : any/c = #f
```

Creates a new structure type property and returns three values:

- a *structure type property descriptor*, for use with make-struct-type and struct;
- a *property predicate* procedure, which takes an arbitrary value and returns #t if the value is a descriptor or instance of a structure type that has a value for the property, #f otherwise;
- a *property accessor* procedure, which returns the value associated with the structure type given its descriptor or one of its instances; if the structure type does not have a value for the property, or if any other kind of value is provided, the exn:fail:contract exception is raised unless a second argument, *failure-result*, is supplied to the procedure. In that case, if *failure-result* is a procedure, it is called (through a tail call) with no arguments to produce the result of the property accessor procedure; otherwise, *failure-result* is itself returned as the result.

If the optional *guard* is supplied as a procedure, it is called by make-struct-type before attaching the property to a new structure type. The *guard* must accept two arguments: a value for the property supplied to make-struct-type, and a list containing information about the new structure type. The list contains the values that struct-type-info would return for the new structure type if it skipped the immediate current-inspector control check (but not the check for exposing an ancestor structure type, if any; see §14.9 “Structure Inspectors”).

The result of calling *guard* is associated with the property in the target structure type, instead of the value supplied to make-struct-type. To reject a property association (e.g., because the value supplied to make-struct-type is inappropriate for the property), the *guard*

can raise an exception. Such an exception prevents `make-struct-type` from returning a structure type descriptor.

If *guard* is `'can-impersonate`, then the property's accessor can be redirected through `impersonate-struct`. This option is identical to supplying `#t` as the *can-impersonate?* argument and is provided for backwards compatibility.

The optional *supers* argument is a list of properties that are automatically associated with some structure type when the newly created property is associated to the structure type. Each property in *supers* is paired with a procedure that receives the value supplied for the new property (after it is processed by *guard*) and returns a value for the associated property (which is then sent to that property's guard, of any).

The optional *can-impersonate?* argument determines if the structure type property can be redirected through `impersonate-struct`. If the argument is `#f`, then redirection is not allowed. Otherwise, the property accessor may be redirected by a struct impersonator.

Examples:

```
> (define-values (prop:p p? p-ref) (make-struct-type-property 'p))

> (define-values (struct:a make-a a? a-ref a-set!)
    (make-struct-type 'a #f 2 1 'uninitialized
                      (list (cons prop:p 8))))

> (p? struct:a)
#t
> (p? 13)
#f
> (define an-a (make-a 'x 'y))

> (p? an-a)
#t
> (p-ref an-a)
8
> (define-values (struct:b make-b b? b-ref b-set!)
    (make-struct-type 'b #f 0 0 #f))

> (p? struct:b)
#f
> (define-values (prop:q q? q-ref) (make-struct-type-property
                                    'q (lambda (v si) (add1 v))
                                    (list (cons prop:p sqrt))))

> (define-values (struct:c make-c c? c-ref c-set!)
    (make-struct-type 'c #f 0 0 'uninit
                      (list (cons prop:q 8))))
```

```
> (q-ref struct:c)
9
> (p-ref struct:c)
3
```

```
(struct-type-property? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a structure type property descriptor value, `#f` otherwise.

```
(struct-type-property-accessor-procedure? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is an accessor procedure produced by `make-struct-type-property`, `#f` otherwise.

5.4 Generic Interfaces

```
(require racket/generic)
```
package: `base`

A *generic interface* allows per-type methods to be associated with generic functions. Generic functions are defined using a `define-generics` form. Method implementations for a structure type are defined using the `#:methods` keyword (see §5.1 “Defining Structure Types: `struct`”).

```
(define-generics id
  generics-opt ...
  [method-id . kw-formals*] ...
  generics-opt ...)

generics-opt = #:defaults ([default-pred? default-impl ...] ...)
             | #:fast-defaults ([fast-pred? fast-impl ...] ...)
             | #:fallbacks [fallback-impl ...]
             | #:defined-predicate defined-pred-id
             | #:defined-table defined-table-id
             | #:derive-property prop-expr prop-value-expr

 kw-formals* = (arg* ...)
             | (arg* ...+ . rest-id)
             | rest-id

       arg* = arg-id
            | [arg-id]
            | keyword arg-id
            | keyword [arg-id]
```

Defines the following names, plus any specified by keyword options.

- `gen:id` as a transformer binding for the static information about a new generic interface;
- `id?` as a predicate identifying instances of structure types that implement this generic group; and
- each `method-id` as a *generic method* that calls the corresponding method on values where `id?` is true. Each `method-id`'s `kw-formals*` must include a required by-position argument that is `free-identifier=?` to `id`. That argument is used in the generic definition to locate the specialization.
- `id/c` as a contract combinator that recognizes instances of structure types which implement the `gen:id` generic interface. The combinator takes pairs of `method-id`s and contracts. The contracts will be applied to each of the corresponding method implementations. The `id/c` combinator is intended to be used to contract the range of a constructor procedure for a struct type that implements the generic interface.

The `#:defaults` option may be provided at most once. When it is provided, each generic function uses `default-pred?`s to dispatch to the given default implementations, `default-impl`s, if dispatching to the generic method table fails. The syntax of the `default-impl`s is the same as the methods provided for the `#:methods` keyword for `struct`.

The `#:fast-defaults` option may be provided at most once. It works the same as `#:defaults`, except the `fast-pred?`s are checked before dispatching to the generic method table. This option is intended to provide a fast path for dispatching to built-in datatypes, such as lists and vectors, that do not overlap with structures implementing `gen:id`.

The `#:fallbacks` option may be provided at most once. When it is provided, the `fallback-impl`s define method implementations that are used for any instance of the generic interface that does not supply a specific implementation. The syntax of the `fallback-impl`s is the same as the methods provided for the `#:methods` keyword for `struct`.

The `#:defined-predicate` option may be provided at most once. When it is provided, `defined-pred-id` is defined as a procedure that reports whether a specific instance of the generic interface implements a given set of methods. Specifically, `(defined-pred-id v 'name ...)` produces `#t` if `v` has implementations for each method `name`, not counting `#:fallbacks` implementations, and produces `#f` otherwise. This procedure is intended for use by higher-level APIs to adapt their behavior depending on method availability.

The `#:defined-table` option may be provided at most once. When it is provided, `defined-table-id` is defined as a procedure that takes an instance of the generic interface and returns an immutable hash table that maps symbols corresponding to method names to booleans representing whether or not that method is implemented by the instance. This option is deprecated; use `#:defined-predicate` instead.

The `#:derive-property` option may be provided any number of times. Each time it is provided, it specifies a structure type property via *prop-expr* and a value for the property via *prop-value-expr*. All structures implementing the generic interface via `#:methods` automatically implement this structure type property using the provided values. When *prop-value-expr* is executed, each *method-id* is bound to its specific implementation for the structure type.

If a value `v` satisfies *id*`?`, then `v` is a *generic instance* of `gen:`*id*.

If a generic instance `v` has a corresponding implementation for some *method-id* provided via `#:methods` in `struct` or via `#:defaults` or `#:fast-defaults` in `define-generics`, then *method-id* is an *implemented generic method* of `v`.

If *method-id* is not an implemented generic method of a generic instance `v`, and *method-id* has a fallback implementation that does not raise an `exn:fail:support` exception when given `v`, then *method-id* is a *supported generic method* of `v`.

```
(raise-support-error name v) → none/c
  name : symbol?
  v : any/c
```

Raises an `exn:fail:support` exception for a generic method called *name* that does not support the generic instance `v`.

Example:

```
> (raise-support-error 'some-method-name '("arbitrary" "instance" "value"))
some-method-name: not implemented for '("arbitrary"
"instance" "value")
```

```
(struct exn:fail:support exn:fail ()
    #:transparent)
```

Raised for generic methods that do not support the given generic instance.

```
(define/generic local-id method-id)

  local-id : identifier?
  method-id : identifier?
```

When used inside the method definitions associated with the `#:methods` keyword, binds *local-id* to the generic for *method-id*. This form is useful for method specializations to use generic methods (as opposed to the local specialization) on other values.

Using the `define/generic` form outside a `#:methods` specification in `struct` (or `define-struct`) is an syntax error.

Examples:

```
> (define-generics printable
    (gen-print printable [port])
    (gen-port-print port printable)
    (gen-print* printable [port] #:width width #:height [height])
    #:defaults ([string?
                 (define/generic super-print gen-print)
                 (define (gen-print s [port (current-output-
port)])
                   (fprintf port "String: ~a" s))
                 (define (gen-port-print port s)
                   (super-print s port))
                 (define (gen-print* s [port (current-output-
port)]
                                       #:width w #:height [h 0])
                   (fprintf port "String (~ax~a): ~a" w h s))]))

> (define-struct num (v)
    #:methods gen:printable
    [(define/generic super-print gen-print)
     (define (gen-print n [port (current-output-port)])
       (fprintf port "Num: ~a" (num-v n)))
     (define (gen-port-print port n)
       (super-print n port))
     (define (gen-print* n [port (current-output-port)]
                         #:width w #:height [h 0])
       (fprintf port "Num (~ax~a): ~a" w h (num-v n)))])

> (define-struct bool (v)
    #:methods gen:printable
    [(define/generic super-print gen-print)
     (define (gen-print b [port (current-output-port)])
       (fprintf port "Bool: ~a"
                (if (bool-v b) "Yes" "No")))
     (define (gen-port-print port b)
       (super-print b port))
     (define (gen-print* b [port (current-output-port)]
                         #:width w #:height [h 0])
       (fprintf port "Bool (~ax~a): ~a" w h
                (if (bool-v b) "Yes" "No")))])

> (define x (make-num 10))

> (gen-print x)
Num: 10
```

```
> (gen-port-print (current-output-port) x)
Num: 10

> (gen-print* x #:width 100 #:height 90)
Num (100x90): 10

> (gen-print "Strings are printable too!")
String: Strings are printable too!

> (define y (make-bool #t))

> (gen-print y)
Bool: Yes

> (gen-port-print (current-output-port) y)
Bool: Yes

> (gen-print* y #:width 100 #:height 90)
Bool (100x90): Yes

> (define/contract make-num-contracted
    (-> number?
        (printable/c
          [gen-print (->* (printable?) (output-port?) void?)]
          [gen-port-print (-> output-port? printable? void?)]
          [gen-print* (->* (printable? #:width exact-nonnegative-
integer?)
                           (output-port? #:height exact-
nonnegative-integer?)
                           void?)]))
     make-num)

> (define z (make-num-contracted 10))

> (gen-print* z #:width "not a number" #:height 5)
make-num-contracted: contract violation
  expected: exact-nonnegative-integer?
  given: "not a number"
  in: the #:width argument of
      method gen-print*
      the range of
      (->
       number?
       (printable/c
        (gen-print
```

```
          (->* (printable?) (output-port?) void?))
         (gen-port-print
          (-> output-port? printable? void?))
         (gen-print*
          (->*
           (printable?
            #:width
            exact-nonnegative-integer?)
           (output-port?
            #:height
            exact-nonnegative-integer?)
           void?))))
  contract from:
      (definition make-num-contracted)
  blaming: top-level
   (assuming the contract is correct)
  at: eval:15.0
```

```
(generic-instance/c gen-id [method-id method-ctc] ...)

  method-ctc : contract?
```

Creates a contract that recognizes structures that implement the generic interface *gen-id*, and constrains their implementations of the specified *method-id*s with the corresponding *method-ctc*s.

```
(impersonate-generics gen-id val-expr [method-id method-proc] ...)

  method-proc : (any/c . -> . any/c)
```

Creates an impersonator of *val-expr*, which must be a structure that implements the generic interface *gen-id*. The impersonator applies the specified *method-proc*s to the structure's implementation of the corresponding *method-id*s, and replaces the method implementation with the result.

```
(chaperone-generics gen-id val-expr [method-id method-proc] ...)

  method-proc : (any/c . -> . any/c)
```

Creates a chaperone of *val-expr*, which must be a structure that implements the generic interface *gen-id*. The chaperone applies the specified *method-proc*s to the structure's implementation of the corresponding *method-id*s, and replaces the method implementation with the result, which must be a chaperone of the original.

```
(redirect-generics mode gen-id val-expr [method-id method-proc] ...)

  method-proc : (any/c . -> . any/c)
```

Creates an impersonator of *val-expr*, like `impersonate-generics`, if *mode* evaluates to `#f`. Creates a chaperone of *val-expr*, like `chaperone-generics`, otherwise.

5.5 Copying and Updating Structures

```
(struct-copy id struct-expr fld-id ...)

fld-id = [field-id expr]
       | [field-id #:parent parent-id expr]
```

Creates a new instance of the structure type *id* with the same field values as the structure produced by *struct-expr*, except that the value of each supplied *field-id* is instead determined by the corresponding *expr*. If `#:parent` is specified, the *parent-id* must be bound to a parent structure type of *id*.

The *id* must have a transformer binding that encapsulates information about a structure type (i.e., like the initial identifier bound by `struct`), and the binding must supply a constructor, a predicate, and all field accessors.

Each *field-id* is combined with *id* (or *parent-id*, if present) to form *id-field-id* (using the lexical context of *field-id*), which must be one of the accessor bindings in *id*. The accessor bindings determined by different *field-id*s must be distinct. The order of the *field-id*s need not match the order of the corresponding fields in the structure type.

The *struct-expr* is evaluated first. The result must be an instance of the *id* structure type, otherwise the `exn:fail:contract` exception is raised. Next, the field *expr*s are evaluated in order (even if the fields that correspond to the *field-id*s are in a different order). Finally, the new structure instance is created.

The result of *struct-expr* can be an instance of a sub-type of *id*, but the resulting copy is an immediate instance of *id* (not the sub-type).

Examples:

```
> (struct fish (color weight) #:transparent)

> (define marlin (fish 'orange-and-white 11))

> (define dory (struct-copy fish marlin
                            [color 'blue]))

> dory
(fish 'blue 11)
> (struct shark fish (weeks-since-eating-fish) #:transparent)
```

```
> (define bruce (shark 'grey 110 3))

> (define chum (struct-copy shark bruce
                            [weight #:parent fish 90]
                            [weeks-since-eating-fish 0]))

> chum
(shark 'grey 90 0)
; subtypes can be copied as if they were supertypes,
; but the result is an instance of the supertype
> (define not-really-chum
    (struct-copy fish bruce
                 [weight 90]))

> not-really-chum
(fish 'grey 90)
```

5.6 Structure Utilities

```
(struct->vector v [opaque-v]) → vector?
  v : any/c
  opaque-v : any/c = '...
```

Creates a vector representing `v`. The first slot of the result vector contains a symbol whose printed name has the form `struct:id`. Each remaining slot contains either the value of a field in `v`, if it is accessible via the current inspector, or `opaque-v` for a field that is not accessible. A single `opaque-v` value is used in the vector for contiguous inaccessible fields. (Consequently, the size of the vector does not match the size of the `struct` if more than one field is inaccessible.)

```
(struct? v) → any
  v : any/c
```

Returns `#t` if `struct-info` exposes any structure types of `v` with the current inspector, `#f` otherwise.

Typically, when `(struct? v)` is true, then `(struct->vector v)` exposes at least one field value. It is possible, however, for the only visible types of `v` to contribute zero fields.

```
(struct-type? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a structure type descriptor value, `#f` otherwise.

```
(struct-constructor-procedure? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a constructor procedure generated by `struct` or `make-struct-type`, `#f` otherwise.

```
(struct-predicate-procedure? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a predicate procedure generated by `struct` or `make-struct-type`, `#f` otherwise.

```
(struct-accessor-procedure? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is an accessor procedure generated by `struct`, `make-struct-type`, or `make-struct-field-accessor`, `#f` otherwise.

```
(struct-mutator-procedure? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a mutator procedure generated by `struct`, `make-struct-type`, or `make-struct-field-mutator`, `#f` otherwise.

```
(prefab-struct-key v) → (or/c #f symbol? list?)
  v : any/c
```

Returns `#f` if *v* is not an instance of a prefab structure type. Otherwise, the result is the shorted key that could be used with `make-prefab-struct` to create an instance of the structure type.

Examples:

```
> (prefab-struct-key #s(cat "Garfield"))
'cat
> (struct cat (name) #:prefab)

> (struct cute-cat cat (shipping-dest) #:prefab)

> (cute-cat "Nermel" "Abu Dhabi")
'#s((cute-cat cat 1) "Nermel" "Abu Dhabi")
> (prefab-struct-key (cute-cat "Nermel" "Abu Dhabi"))
'(cute-cat cat 1)
```

```
(make-prefab-struct key v ...) → struct?
  key : prefab-key?
  v : any/c
```

Creates an instance of a prefab structure type, using the *v*s as field values. The *key* and the number of *v*s determine the prefab structure type.

A *key* identifies a structure type based on a list with the following items:

- A symbol for the structure type's name.
- An exact, nonnegative integer representing the number of non-automatic fields in the structure type, not counting fields from the supertype (if any).
- A list of two items, where the first is an exact, nonnegative integer for the number of automatic fields in the structure type that are not from the supertype (if any), and the second element is an arbitrary value that is the value for the automatic fields.
- A vector of exact, nonnegative integers that indicate mutable non-automatic fields in the structure type, counting from `0` and not including fields from the supertype (if any).
- Nothing else, if the structure type has no supertype. Otherwise, the rest of the list is the key for the supertype.

An empty vector and an auto-field list that starts with `0` can be omitted. Furthermore, the first integer (which indicates the number of non-automatic fields) can be omitted, since it can be inferred from the number of supplied *v*s. Finally, a single symbol can be used instead of a list that contains only a symbol (in the case that the structure type has no supertype, no automatic fields, and no mutable fields).

The total field count must be no more than 32768. If the number of fields indicated by *key* is inconsistent with the number of supplied *v*s, the `exn:fail:contract` exception is raised.

Examples:

```
> (make-prefab-struct 'clown "Binky" "pie")
'#s(clown "Binky" "pie")
> (make-prefab-struct '(clown 2) "Binky" "pie")
'#s(clown "Binky" "pie")
> (make-prefab-struct '(clown 2 (0 #f) #()) "Binky" "pie")
'#s(clown "Binky" "pie")
> (make-prefab-struct '(clown 1 (1 #f) #()) "Binky" "pie")
'#s((clown (1 #f)) "Binky" "pie")
> (make-prefab-struct '(clown 1 (1 #f) #(0)) "Binky" "pie")
'#s((clown (1 #f) #(0)) "Binky" "pie")
```

```
(prefab-key->struct-type key field-count) → struct-type?
  key : prefab-key?
  field-count : (integer-in 0 32768)
```

Returns a structure type descriptor for the prefab structure type specified by the combination of *key* and *field-count*.

If the number of fields indicated by *key* is inconsistent with *field-count*, the exn:fail:contract exception is raised.

```
(prefab-key? v) → boolean?
  v : any/c
```

Return #t if *v* can be a prefab structure type key, #f otherwise.

See make-prefab-struct for a description of valid key shapes.

5.7 Structure Type Transformer Binding

The struct form binds the name of a structure type as a transformer binding that records the other identifiers bound to the structure type, the constructor procedure, the predicate procedure, and the field accessor and mutator procedures. This information can be used during the expansion of other expressions via syntax-local-value.

For example, the struct variant for subtypes uses the base type name t to find the variable struct:t containing the base type's descriptor; it also folds the field accessor and mutator information for the base type into the information for the subtype. As another example, the match form uses a type name to find the predicates and field accessors for the structure type. The struct form in an imported signature for unit causes the unit transformer to generate information about imported structure types, so that match and subtyping struct forms work within the unit.

The expansion-time information for a structure type can be represented directly as a list of six elements (of the same sort that the encapsulated procedure must return):

- an identifier that is bound to the structure type's descriptor, or #f if none is known;
- an identifier that is bound to the structure type's constructor, or #f if none is known;
- an identifier that is bound to the structure type's predicate, or #f if none is known;
- a list of identifiers bound to the field accessors of the structure type, optionally with #f as the list's last element. A #f as the last element indicates that the structure type may have additional fields, otherwise the list is a reliable indicator of the number of

fields in the structure type. Furthermore, the accessors are listed in reverse order for the corresponding constructor arguments. (The reverse order enables sharing in the lists for a subtype and its base type.)

- a list of identifiers bound to the field mutators of the structure type, or `#f` for each field that has no known mutator, and optionally with an extra `#f` as the list's last element (if the accessor list has such a `#f`). The list's order and the meaning of a final `#f` are the same as for the accessor identifiers, and the length of the mutator list is the same as the accessor list's length.

- an identifier that determines a super-type for the structure type, `#f` if the super-type (if any) is unknown, or `#t` if there is no super-type. If a super-type is specified, the identifier is also bound to structure-type expansion-time information.

Instead of this direct representation, the representation can be a structure created by `make-struct-info` (or an instance of a subtype of `struct:struct-info`), which encapsulates a procedure that takes no arguments and returns a list of six elements. Alternately, the representation can be a structure whose type has the `prop:struct-info` structure type property. Finally, the representation can be an instance of a structure type derived from `struct:struct-info` or with the `prop:struct-info` property that also implements `prop:procedure`, and where the instance is further is wrapped by `make-set!-transformer`. In addition, the representation may implement the `prop:struct-auto-info` property.

Use `struct-info?` to recognize all allowed forms of the information, and use `extract-struct-info` to obtain a list from any representation.

The implementor of a syntactic form can expect users of the form to know what kind of information is available about a structure type. For example, the `match` implementation works with structure information containing an incomplete set of accessor bindings, because the user is assumed to know what information is available in the context of the `match` expression. In particular, the `match` expression can appear in a `unit` form with an imported structure type, in which case the user is expected to know the set of fields that are listed in the signature for the structure type.

`(require racket/struct-info)` package: `base`

The bindings documented in this section are provided by the `racket/struct-info` library, not `racket/base` or `racket`.

```
(struct-info? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is either a six-element list with the correct shape for representing structure-type information, a procedure encapsulated by `make-struct-info`, a structure with the `prop:struct-info` property, or a structure type derived from `struct:struct-info` or with `prop:struct-info` and wrapped with `make-set!-transformer`.

```
(checked-struct-info? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a procedure encapsulated by `make-struct-info` and produced by `struct`, but only when no parent type is specified or the parent type is also specified through a transformer binding to such a value.

```
(make-struct-info thunk) → struct-info?
  thunk : (-> (and/c struct-info? list?))
```

Encapsulates a thunk that returns structure-type information in list form. Note that accessors are listed in reverse order, as mentioned in §5.7 “Structure Type Transformer Binding”.

Examples:

```
> (define (new-pair? x) (displayln "new pair?") (pair? x))

> (define (new-car x) (displayln "new car") (car x))

> (define (new-cdr x) (displayln "new cdr") (cdr x))

> (define-syntax new-list
    (make-struct-info
     (λ () (list #f
                 #'cons
                 #'new-pair?
                 (list #'new-cdr #'new-car)
                 (list #f #f)
                 #t))))

> (match (list 1 2 3)
    [(new-list hd tl) (append tl (list hd))])
new pair?
new car
new cdr
'(2 3 1)
```

Examples:

```
> (struct A (x y))

> (define (new-A-x a) (displayln "A-x") (A-x a))

> (define (new-A-y a) (displayln "A-y") (A-y a))
```

```
> (define (new-A? a) (displayln "A?") (A? a))

> (define-syntax A-info
    (make-struct-info
     (λ () (list #'A
                 #'A
                 #'new-A?
                 (list #'new-A-y #'new-A-x)
                 (list #f #f)
                 #t))))

> (define-match-expander B
    (syntax-rules () [(_ x ...) (A-info x ...)]))

> (match (A 10 20)
    [(B x y) (list y x)])
A?
A-x
A-y
'(20 10)
```

```
(extract-struct-info v) → (and/c struct-info? list?)
  v : struct-info?
```

Extracts the list form of the structure type information represented by *v*.

```
struct:struct-info : struct-type?
```

The structure type descriptor for the structure type returned by `make-struct-info`. This structure type descriptor is mostly useful for creating structure subtypes. The structure type includes a guard that checks an instance's first field in the same way as `make-struct-info`.

```
prop:struct-info : struct-type-property?
```

The structure type property for creating new structure types like `struct:struct-info`. The property value must be a procedure of one argument that takes an instance structure and returns structure-type information in list form.

```
prop:struct-auto-info : struct-type-property?
(struct-auto-info? v) → boolean?
  v : any/c
(struct-auto-info-lists sai)
 → (list/c (listof identifier?) (listof identifier?))
  sai : struct-auto-info?
```

The `prop:struct-auto-info` property is implemented to provide static information about which of the accessor and mutator identifiers for a structure type correspond to `#:auto` fields

(so that they have no corresponding argument in the constructor). The property value must be a procedure that accepts an instance structure to which the property is given, and the result must be two lists of identifiers suitable as a result from `struct-auto-info-lists`.

The `struct-auto-info?` predicate recognizes values that implement the `prop:struct-auto-info` property.

The `struct-auto-info-lists` function extracts two lists of identifiers from a value that implements the `prop:struct-auto-info` property. The first list should be a subset of the accessor identifiers for the structure type described by *sai*, and the second list should be a subset of the mutator identifiers. The two subsets correspond to `#:auto` fields.

6 Classes and Objects

§13 “Classes and Objects” in ***The Racket Guide*** introduces classes and objects.

```
(require racket/class)
```
package: base

The bindings documented in this section are provided by the `racket/class` and `racket` libraries, but not `racket/base`.

A *class* specifies

- a collection of fields;
- a collection of methods;
- initial value expressions for the fields; and
- initialization variables that are bound to initialization arguments.

In the context of the class system, an *object* is a collection of bindings for fields that are instantiated according to a class description.

The class system allows a program to define a new class (a *derived class*) in terms of an existing class (the *superclass*) using inheritance, overriding, and augmenting:

- *inheritance*: An object of a derived class supports methods and instantiates fields declared by the derived class’s superclass, as well as methods and fields declared in the derived class expression.
- *overriding*: Some methods declared in a superclass can be replaced in the derived class. References to the overridden method in the superclass use the implementation in the derived class.
- *augmenting*: Some methods declared in a superclass can be merely extended in the derived class. The superclass method specifically delegates to the augmenting method in the derived class.

An *interface* is a collection of method names to be implemented by a class, combined with a derivation requirement. A class *implements* an interface when it

- declares (or inherits) a public method for each variable in the interface;
- is derived from the class required by the interface, if any; and
- specifically declares its intention to implement the interface.

A class can implement any number of interfaces. A derived class automatically implements any interface that its superclass implements. Each class also implements an implicitly-defined interface that is associated with the class. The implicitly-defined interface contains all of the class's public method names, and it requires that all other implementations of the interface are derived from the class.

A new interface can *extend* one or more interfaces with additional method names; each class that implements the extended interface also implements the original interfaces. The derivation requirements of the original interface must be consistent, and the extended interface inherits the most specific derivation requirement from the original interfaces.

Classes, objects, and interfaces are all values. However, a class or interface is not an object (i.e., there are no "meta-classes" or "meta-interfaces").

6.1 Creating Interfaces

§13 "Classes and Objects" in ***The Racket Guide*** introduces classes, objects, and interfaces.

```
(interface (super-interface-expr ...) name-clause ...)

name-clause = id
            | (id contract-expr)
```

Produces an interface. The *id*s must be mutually distinct.

Each *super-interface-expr* is evaluated (in order) when the `interface` expression is evaluated. The result of each *super-interface-expr* must be an interface value, otherwise the `exn:fail:object` exception is raised. The interfaces returned by the *super-interface-expr*s are the new interface's superinterfaces, which are all extended by the new interface. Any class that implements the new interface also implements all of the superinterfaces.

The result of an `interface` expression is an interface that includes all of the specified *id*s, plus all identifiers from the superinterfaces. Duplicate identifier names among the superinterfaces are ignored, but if a superinterface contains one of the *id*s in the `interface` expression, the `exn:fail:object` exception is raised. A given *id* may be paired with a corresponding *contract-expr*.

If no *super-interface-expr*s are provided, then the derivation requirement of the resulting interface is trivial: any class that implements the interface must be derived from `object%`. Otherwise, the implementation requirement of the resulting interface is the most specific requirement from its superinterfaces. If the superinterfaces specify inconsistent derivation requirements, the `exn:fail:object` exception is raised.

Examples:

```
(define file-interface<%>
  (interface () open close read-byte write-byte))
```

```
(define directory-interface<%>
  (interface (file-interface<%>)
    [file-list (->m (listof (is-a?/c file-interface<%>)))]
    parent-directory))
```

```
(interface* (super-interface-expr ...)
            ([property-expr val-expr] ...)
  name-clause ...)

name-clause = id
            | (id contract-expr)
```

Like `interface`, but also associates to the interface the structure-type properties produced by the *property-expr*s with the corresponding *val-expr*s.

Whenever the resulting interface (or a sub-interface derived from it) is explicitly implemented by a class through the `class*` form, each property is attached with its value to a structure type that instantiated by instances of the class. Specifically, the property is attached to a structure type with zero immediate fields, which is extended to produce the internal structure type for instances of the class (so that no information about fields is accessible to the structure type property's guard, if any).

Example:

```
(define i<%> (interface* () ([prop:custom-write
                              (lambda (obj port mode) (void))])
               method1 method2 method3))
```

6.2 Creating Classes

§13 "Classes and Objects" in ***The Racket Guide*** introduces classes and objects.

```
object% : class?
```

A built-in class that has no methods fields, implements only its own interface (`class->interface object%`), and is transparent (i.e,. its inspector is `#f`, so all immediate instances are `equal?`). All other classes are derived from `object%`.

```
(class* superclass-expr (interface-expr ...)
  class-clause
  ...)
```

```
     class-clause = (inspect inspector-expr)
                  | (init init-decl ...)
                  | (init-field init-decl ...)
                  | (field field-decl ...)
                  | (inherit-field maybe-renamed ...)
                  | (init-rest id)
                  | (init-rest)
                  | (public maybe-renamed ...)
                  | (pubment maybe-renamed ...)
                  | (public-final maybe-renamed ...)
                  | (override maybe-renamed ...)
                  | (overment maybe-renamed ...)
                  | (override-final maybe-renamed ...)
                  | (augment maybe-renamed ...)
                  | (augride maybe-renamed ...)
                  | (augment-final maybe-renamed ...)
                  | (private id ...)
                  | (abstract id ...)
                  | (inherit maybe-renamed ...)
                  | (inherit/super maybe-renamed ...)
                  | (inherit/inner maybe-renamed ...)
                  | (rename-super renamed ...)
                  | (rename-inner renamed ...)
                  | method-definition
                  | definition
                  | expr
                  | (begin class-clause ...)

        init-decl = id
                  | (renamed)
                  | (maybe-renamed default-value-expr)

       field-decl = (maybe-renamed default-value-expr)

    maybe-renamed = id
                  | renamed

          renamed = (internal-id external-id)

method-definition = (define-values (id) method-procedure)

 method-procedure = (lambda kw-formals expr ...+)
                  | (case-lambda (formals expr ...+) ...)
                  | (#%plain-lambda formals expr ...+)
                  | (let-values ([(id) method-procedure] ...)
                      method-procedure)
                  | (letrec-values ([(id) method-procedure] ...)
                      method-procedure)
                  | (let-values ([(id) method-procedure] ...+)
                      id)
                  | (letrec-values ([(id) method-procedure] ...+)
                      id)
```

Produces a class value.

The *superclass-expr* expression is evaluated when the `class*` expression is evaluated. The result must be a class value (possibly `object%`), otherwise the `exn:fail:object` exception is raised. The result of the *superclass-expr* expression is the new class's superclass.

The *interface-expr* expressions are also evaluated when the `class*` expression is evaluated, after *superclass-expr* is evaluated. The result of each *interface-expr* must be an interface value, otherwise the `exn:fail:object` exception is raised. The interfaces returned by the *interface-expr*s are all implemented by the class. For each identifier in each interface, the class (or one of its ancestors) must declare a public method with the same name, otherwise the `exn:fail:object` exception is raised. The class's superclass must satisfy the implementation requirement of each interface, otherwise the `exn:fail:object` exception is raised.

An `inspect` *class-clause* selects an inspector (see §14.9 "Structure Inspectors") for the class extension. The *inspector-expr* must evaluate to an inspector or `#f` when the `class*` form is evaluated. Just as for structure types, an inspector controls access to the class's fields, including private fields, and also affects comparisons using `equal?`. If no `inspect` clause is provided, access to the class is controlled by the parent of the current inspector (see §14.9 "Structure Inspectors"). A syntax error is reported if more than one `inspect` clause is specified.

The other *class-clause*s define initialization arguments, public and private fields, and public and private methods. For each *id* or *maybe-renamed* in a `public`, `override`, `augment`, `pubment`, `overment`, `augride`, `public-final`, `override-final`, `augment-final`, or `private` clause, there must be one *method-definition*. All other definition *class-clause*s create private fields. All remaining *expr*s are initialization expressions to be evaluated when the class is instantiated (see §6.3 "Creating Objects").

The result of a `class*` expression is a new class, derived from the specified superclass and implementing the specified interfaces. Instances of the class are created with the `instantiate` form or `make-object` procedure, as described in §6.3 "Creating Objects".

Each *class-clause* is (partially) macro-expanded to reveal its shapes. If a *class-clause* is a `begin` expression, its sub-expressions are lifted out of the `begin` and treated as *class-clause*s, in the same way that `begin` is flattened for top-level and embedded definitions.

Within a `class*` form for instances of the new class, `this` is bound to the object itself; `this%` is bound to the class of the object; `super-instantiate`, `super-make-object`, and `super-new` are bound to forms to initialize fields in the superclass (see §6.3 "Creating Objects"); `super` is available for calling superclass methods (see §6.2.3.1 "Method Definitions"); and `inner` is available for calling subclass augmentations of methods (see §6.2.3.1 "Method Definitions").

```
(class superclass-expr class-clause ...)
```

Like `class*`, but omits the *`interface-expr`*s, for the case that none are needed.

Example:

```
(define book-class%
  (class object%
    (field (pages 5))
    (define/public (letters)
      (* pages 500))
    (super-new)))
```

this

Within a `class*` form, `this` refers to the current object (i.e., the object being initialized or whose method was called). Use outside the body of a `class*` form is a syntax error.

Examples:

```
(define (describe obj)
  (printf "Hello ~a\n" obj))

(define table%
  (class object%
    (define/public (describe-self)
      (describe this))
    (super-new)))

> (send (new table%) describe-self)
Hello #(struct:object:table% ...)
```

this%

Within a `class*` form, `this%` refers to the class of the current object (i.e., the object being initialized or whose method was called). Use outside the body of a `class*` form is a syntax error.

Examples:

```
(define account%
  (class object%
    (super-new)
    (init-field balance)
    (define/public (add n)
      (new this% [balance (+ n balance)]))))
```

```
(define savings%
  (class account%
    (super-new)
    (inherit-field balance)
    (define interest 0.04)
    (define/public (add-interest)
      (send this add (* interest balance)))))

> (let* ([acct (new savings% [balance 500])]
         [acct (send acct add 500)]
         [acct (send acct add-interest)])
    (printf "Current balance: ~a\n" (get-field balance acct)))
Current balance: 1040.0
```

(inspect *inspector-expr*)

See `class*`; use outside the body of a `class*` form is a syntax error.

(init *init-decl* ...)

See `class*` and §6.2.1 “Initialization Variables”; use outside the body of a `class*` form is a syntax error.

Example:

```
> (class object%
    (super-new)
    (init turnip
          [(internal-potato potato)]
          [carrot 'good]
          [(internal-rutabaga rutabaga) 'okay]))
#<class:eval:10:0>
```

(init-field *init-decl* ...)

See `class*`, §6.2.1 “Initialization Variables”, and §6.2.2 “Fields”; use outside the body of a `class*` form is a syntax error.

Example:

```
> (class object%
    (super-new)
    (init-field turkey
```

```
                [(internal-ostrich ostrich)]
                [chicken 7]
                [(internal-emu emu) 13]))
#<class:eval:11:0>
```

```
(field field-decl ...)
```

See `class*` and §6.2.2 “Fields”; use outside the body of a `class*` form is a syntax error.

Example:

```
> (class object%
    (super-new)
    (field [minestrone 'ready]
           [(internal-coq-au-vin coq-au-vin) 'stewing]))
#<class:eval:12:0>
```

```
(inherit-field maybe-renamed ...)
```

See `class*` and §6.2.2 “Fields”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define cookbook%
  (class object%
    (super-new)
    (field [recipes '(caldo-verde oyakodon eggs-benedict)]
           [pages 389])))

> (class cookbook%
    (super-new)
    (inherit-field recipes
                   [internal-pages pages]))
#<class:eval:14:0>
```

```
(init-rest id)
(init-rest)
```

See `class*` and §6.2.1 “Initialization Variables”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define fruit-basket%
```

```
  (class object%
    (super-new)
    (init-rest fruits)
    (displayln fruits)))

> (make-object fruit-basket% 'kiwi 'lychee 'melon)
(kiwi lychee melon)
(object:fruit-basket% ...)
```

```
(public maybe-renamed ...)
```

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define jumper%
  (class object%
    (super-new)
    (define (skip) 'skip)
    (define (hop) 'hop)
    (public skip [hop jump])))

> (send (new jumper%) skip)
'skip
> (send (new jumper%) jump)
'hop
```

```
(pubment maybe-renamed ...)
```

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define runner%
  (class object%
    (super-new)
    (define (run) 'run)
    (define (trot) 'trot)
    (pubment run [trot jog])))

> (send (new runner%) run)
'run
> (send (new runner%) jog)
'trot
```

```
(public-final maybe-renamed ...)
```

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define point%
  (class object%
    (super-new)
    (init-field [x 0] [y 0])
    (define (get-x) x)
    (define (do-get-y) y)
    (public-final get-x [do-get-y get-y])))

> (send (new point% [x 1] [y 3]) get-y)
3
> (class point%
    (super-new)
    (define (get-x) 3.14)
    (override get-x))
class*: cannot override or augment final method
  method name: get-x
  class name: eval:25:0
```

```
(override maybe-renamed ...)
```

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define sheep%
  (class object%
    (super-new)
    (define/public (bleat)
      (displayln "baaaaaaaaah"))))

(define confused-sheep%
  (class sheep%
    (super-new)
    (define (bleat)
      (super bleat)
      (displayln "???"))
    (override bleat)))
```

```
> (send (new sheep%) bleat)
baaaaaaaaah

> (send (new confused-sheep%) bleat)
baaaaaaaaah
???
```

```
(overment maybe-renamed ...)
```

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define turkey%
  (class object%
    (super-new)
    (define/public (gobble)
      (displayln "gobble gobble"))))

(define extra-turkey%
  (class turkey%
    (super-new)
    (define (gobble)
      (super gobble)
      (displayln "gobble gobble gobble")
      (inner (void) gobble))
    (overment gobble)))

(define cyborg-turkey%
  (class extra-turkey%
    (super-new)
    (define/augment (gobble)
      (displayln "1100111110111111000101100010110110011001 01"))))

> (send (new extra-turkey%) gobble)
gobble gobble
gobble gobble gobble

> (send (new cyborg-turkey%) gobble)
gobble gobble
gobble gobble gobble
1100111110111111000101100010110110011001 01
```

(override-final *maybe-renamed* ...)

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define meeper%
  (class object%
    (super-new)
    (define/public (meep)
      (displayln "meep"))))

(define final-meeper%
  (class meeper%
    (super-new)
    (define (meep)
      (super meep)
      (displayln "This meeping ends with me"))
    (override-final meep)))

> (send (new meeper%) meep)
meep

> (send (new final-meeper%) meep)
meep
This meeping ends with me
```

(augment *maybe-renamed* ...)

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define buzzer%
  (class object%
    (super-new)
    (define/pubment (buzz)
      (displayln "bzzzt")
      (inner (void) buzz))))

(define loud-buzzer%
```

```
  (class buzzer%
    (super-new)
    (define (buzz)
      (displayln "BZZZZZZZZZT"))
    (augment buzz)))

> (send (new buzzer%) buzz)
bzzzt

> (send (new loud-buzzer%) buzz)
bzzzt
BZZZZZZZZZT
```

(augride *maybe-renamed* ...)

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

(augment-final *maybe-renamed* ...)

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

(private *id* ...)

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define light%
  (class object%
    (super-new)
    (define on? #t)
    (define (toggle) (set! on? (not on?)))
    (private toggle)
    (define (flick) (toggle))
    (public flick)))

> (send (new light%) toggle)
send: no such method
  method name: toggle
  class name: light%
> (send (new light%) flick)
```

```
(abstract id ...)
```

See `class*` and §6.2.3.1 “Method Definitions”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define train%
  (class object%
    (super-new)
    (abstract get-speed)
    (init-field [position 0])
    (define/public (move)
      (new this% [position (+ position (get-speed))]))))

(define acela%
  (class train%
    (super-new)
    (define/override (get-speed) 241)))

(define talgo-350%
  (class train%
    (super-new)
    (define/override (get-speed) 330)))

> (new train%)
instantiate: cannot instantiate class with abstract methods
  class: #<class:train%>
  abstract methods:
   get-speed
> (send (new acela%) move)
(object:acela% ...)
```

```
(inherit maybe-renamed ...)
```

See `class*` and §6.2.3.2 “Inherited and Superclass Methods”; use outside the body of a `class*` form is a syntax error.

Examples:

```
(define alarm%
  (class object%
    (super-new)
    (define/public (alarm)
      (displayln "beeeeeeeep"))))
```

```
(define car-alarm%
  (class alarm%
    (super-new)
    (init-field proximity)
    (inherit alarm)
    (when (< proximity 10)
      (alarm))))

> (new car-alarm% [proximity 5])
beeeeeeeep
(object:car-alarm% ...)
```

(inherit/super *maybe-renamed* ...)

See `class*` and §6.2.3.2 “Inherited and Superclass Methods”; use outside the body of a `class*` form is a syntax error.

(inherit/inner *maybe-renamed* ...)

See `class*` and §6.2.3.2 “Inherited and Superclass Methods”; use outside the body of a `class*` form is a syntax error.

(rename-super *renamed* ...)

See `class*` and §6.2.3.2 “Inherited and Superclass Methods”; use outside the body of a `class*` form is a syntax error.

(rename-inner *renamed* ...)

See `class*` and §6.2.3.2 “Inherited and Superclass Methods”; use outside the body of a `class*` form is a syntax error.

(public* (*id expr*) ...)

Shorthand for `(begin (public id) ... (define id expr) ...)`.

(pubment* (*id expr*) ...)

Shorthand for `(begin (pubment id) ... (define id expr) ...)`.

```
(public-final* (id expr) ...)
```

Shorthand for `(begin (public-final id) ... (define id expr) ...)`.

```
(override* (id expr) ...)
```

Shorthand for `(begin (override id) ... (define id expr) ...)`.

```
(overment* (id expr) ...)
```

Shorthand for `(begin (overment id) ... (define id expr) ...)`.

```
(override-final* (id expr) ...)
```

Shorthand for `(begin (override-final id) ... (define id expr) ...)`.

```
(augment* (id expr) ...)
```

Shorthand for `(begin (augment id) ... (define id expr) ...)`.

```
(augride* (id expr) ...)
```

Shorthand for `(begin (augride id) ... (define id expr) ...)`.

```
(augment-final* (id expr) ...)
```

Shorthand for `(begin (augment-final id) ... (define id expr) ...)`.

```
(private* (id expr) ...)
```

Shorthand for `(begin (private id) ... (define id expr) ...)`.

```
(define/public id expr)
(define/public (id . formals) body ...+)
```

Shorthand for `(begin (public id) (define id expr))` or `(begin (public id) (define (id . formals) body ...+))`

```
(define/pubment id expr)
(define/pubment (id . formals) body ...+)
```

Shorthand for (begin (pubment *id*) (define *id expr*)) or (begin (pubment *id*) (define (*id* . *formals*) *body* ...+))

```
(define/public-final id expr)
(define/public-final (id . formals) body ...+)
```

Shorthand for (begin (public-final *id*) (define *id expr*)) or (begin (public-final *id*) (define (*id* . *formals*) *body* ...+))

```
(define/override id expr)
(define/override (id . formals) body ...+)
```

Shorthand for (begin (override *id*) (define *id expr*)) or (begin (override *id*) (define (*id* . *formals*) *body* ...+))

```
(define/overment id expr)
(define/overment (id . formals) body ...+)
```

Shorthand for (begin (overment *id*) (define *id expr*)) or (begin (overment *id*) (define (*id* . *formals*) *body* ...+))

```
(define/override-final id expr)
(define/override-final (id . formals) body ...+)
```

Shorthand for (begin (override-final *id*) (define *id expr*)) or (begin (override-final *id*) (define (*id* . *formals*) *body* ...+))

```
(define/augment id expr)
(define/augment (id . formals) body ...+)
```

Shorthand for (begin (augment *id*) (define *id expr*)) or (begin (augment *id*) (define (*id* . *formals*) *body* ...+))

```
(define/augride id expr)
(define/augride (id . formals) body ...+)
```

Shorthand for (begin (augride id) (define id expr)) or (begin (augride id) (define (id . formals) body ...+))

```
(define/augment-final id expr)
(define/augment-final (id . formals) body ...+)
```

Shorthand for (begin (augment-final id) (define id expr)) or (begin (augment-final id) (define (id . formals) body ...+))

```
(define/private id expr)
(define/private (id . formals) body ...+)
```

Shorthand for (begin (private id) (define id expr)) or (begin (private id) (define (id . formals) body ...+))

```
(class/derived original-datum
  (name-id super-expr (interface-expr ...) deserialize-id-expr)
  class-clause
  ...)
```

Like class*, but includes a sub-expression to be used as the source for all syntax errors within the class definition. For example, define-serializable-class expands to class/derived so that errors in the body of the class are reported in terms of define-serializable-class instead of class.

The *original-datum* is the original expression to use for reporting errors.

The *name-id* is used to name the resulting class; if it is #f, the class name is inferred.

The *super-expr*, *interface-expr*s, and *class-clause*s are as for class*.

If the *deserialize-id-expr* is not literally #f, then a serializable class is generated, and the result is two values instead of one: the class and a deserialize-info structure produced by make-deserialize-info. The *deserialize-id-expr* should produce a value suitable as the second argument to make-serialize-info, and it should refer to an export whose value is the deserialize-info structure.

Future optional forms may be added to the sequence that currently ends with *deserialize-id-expr*.

6.2.1 Initialization Variables

A class's initialization variables, declared with init, init-field, and init-rest, are instantiated for each object of a class. Initialization variables can be used in the initial

value expressions of fields, default value expressions for initialization arguments, and in initialization expressions. Only initialization variables declared with `init-field` can be accessed from methods; accessing any other initialization variable from a method is a syntax error.

The values bound to initialization variables are

- the arguments provided with `instantiate` or passed to `make-object`, if the object is created as a direct instance of the class; or,
- the arguments passed to the superclass initialization form or procedure, if the object is created as an instance of a derived class.

If an initialization argument is not provided for an initialization variable that has an associated *`default-value-expr`*, then the *`default-value-expr`* expression is evaluated to obtain a value for the variable. A *`default-value-expr`* is only evaluated when an argument is not provided for its variable. The environment of *`default-value-expr`* includes all of the initialization variables, all of the fields, and all of the methods of the class. If multiple *`default-value-expr`*s are evaluated, they are evaluated from left to right. Object creation and field initialization are described in detail in §6.3 "Creating Objects".

If an initialization variable has no *`default-value-expr`*, then the object creation or superclass initialization call must supply an argument for the variable, otherwise the `exn:fail:object` exception is raised.

Initialization arguments can be provided by name or by position. The external name of an initialization variable can be used with `instantiate` or with the superclass initialization form. Those forms also accept by-position arguments. The `make-object` procedure and the superclass initialization procedure accept only by-position arguments.

Arguments provided by position are converted into by-name arguments using the order of `init` and `init-field` clauses and the order of variables within each clause. When an `instantiate` form provides both by-position and by-name arguments, the converted arguments are placed before by-name arguments. (The order can be significant; see also §6.3 "Creating Objects".)

Unless a class contains an `init-rest` clause, when the number of by-position arguments exceeds the number of declared initialization variables, the order of variables in the superclass (and so on, up the superclass chain) determines the by-name conversion.

If a class expression contains an `init-rest` clause, there must be only one, and it must be last. If it declares a variable, then the variable receives extra by-position initialization arguments as a list (similar to a dotted "rest argument" in a procedure). An `init-rest` variable can receive by-position initialization arguments that are left over from a by-name conversion for a derived class. When a derived class's superclass initialization provides even more by-position arguments, they are prefixed onto the by-position arguments accumulated so far.

If too few or too many by-position initialization arguments are provided to an object creation or superclass initialization, then the exn:fail:object exception is raised. Similarly, if extra by-position arguments are provided to a class with an init-rest clause, the exn:fail:object exception is raised.

Unused (by-name) arguments are to be propagated to the superclass, as described in §6.3 "Creating Objects". Multiple initialization arguments can use the same name if the class derivation contains multiple declarations (in different classes) of initialization variables with the name. See §6.3 "Creating Objects" for further details.

See also §6.2.3.3 "Internal and External Names" for information about internal and external names.

6.2.2 Fields

Each field, init-field, and non-method define-values clause in a class declares one or more new fields for the class. Fields declared with field or init-field are public. Public fields can be accessed and mutated by subclasses using inherit-field. Public fields are also accessible outside the class via class-field-accessor and mutable via class-field-mutator (see §6.4 "Field and Method Access"). Fields declared with define-values are accessible only within the class.

A field declared with init-field is both a public field and an initialization variable. See §6.2.1 "Initialization Variables" for information about initialization variables.

An inherit-field declaration makes a public field defined by a superclass directly accessible in the class expression. If the indicated field is not defined in the superclass, the exn:fail:object exception is raised when the class expression is evaluated. Every field in a superclass is present in a derived class, even if it is not declared with inherit-field in the derived class. The inherit-field clause does not control inheritance, but merely controls lexical scope within a class expression.

When an object is first created, all of its fields have the #<undefined> value (see §4.18 "Void"). The fields of a class are initialized at the same time that the class's initialization expressions are evaluated; see §6.3 "Creating Objects" for more information.

See also §6.2.3.3 "Internal and External Names" for information about internal and external names.

6.2.3 Methods

Method Definitions

Each public, override, augment, pubment, overment, augride, public-final,

override-final, augment-final, and private clause in a class declares one or more method names. Each method name must have a corresponding *method-definition*. The order of public, etc., clauses and their corresponding definitions (among themselves, and with respect to other clauses in the class) does not matter.

As shown in the grammar for class*, a method definition is syntactically restricted to certain procedure forms, as defined by the grammar for *method-procedure*; in the last two forms of *method-procedure*, the body id must be one of the ids bound by let-values or letrec-values. A *method-procedure* expression is not evaluated directly. Instead, for each method, a class-specific method procedure is created; it takes an initial object argument, in addition to the arguments the procedure would accept if the *method-procedure* expression were evaluated directly. The body of the procedure is transformed to access methods and fields through the object argument.

A method declared with public, pubment, or public-final introduces a new method into a class. The method must not be present already in the superclass, otherwise the exn:fail:object exception is raised when the class expression is evaluated. A method declared with public can be overridden in a subclass that uses override, overment, or override-final. A method declared with pubment can be augmented in a subclass that uses augment, augride, or augment-final. A method declared with public-final cannot be overridden or augmented in a subclass.

A method declared with override, overment, or override-final overrides a definition already present in the superclass. If the method is not already present, the exn:fail:object exception is raised when the class expression is evaluated. A method declared with override can be overridden again in a subclass that uses override, overment, or override-final. A method declared with overment can be augmented in a subclass that uses augment, augride, or augment-final. A method declared with override-final cannot be overridden further or augmented in a subclass.

A method declared with augment, augride, or augment-final augments a definition already present in the superclass. If the method is not already present, the exn:fail:object exception is raised when the class expression is evaluated. A method declared with augment can be augmented further in a subclass that uses augment, augride, or augment-final. A method declared with augride can be overridden in a subclass that uses override, overment, or override-final. (Such an override merely replaces the augmentation, not the method that is augmented.) A method declared with augment-final cannot be overridden or augmented further in a subclass.

A method declared with private is not accessible outside the class expression, cannot be overridden, and never overrides a method in the superclass.

When a method is declared with override, overment, or override-final, then the superclass implementation of the method can be called using super form.

When a method is declared with pubment, augment, or overment, then a subclass augmenting method can be called using the inner form. The only difference between public-

`final` and `pubment` without a corresponding `inner` is that `public-final` prevents the declaration of augmenting methods that would be ignored.

A method declared with `abstract` must be declared without an implementation. Subclasses may implement abstract methods via the `override`, `overment`, or `override-final` forms. Any class that contains or inherits any abstract methods is considered abstract and cannot be instantiated.

```
(super id arg ...)
(super id arg ... . arg-list-expr)
```

Always accesses the superclass method, independent of whether the method is overridden again in subclasses. Using the `super` form outside of `class*` is a syntax error. Each *arg* is as for `#%app`: either *arg-expr* or *keyword arg-expr*.

The second form is analogous to using `apply` with a procedure; the *arg-list-expr* must not be a parenthesized expression.

```
(inner default-expr id arg ...)
(inner default-expr id arg ... . arg-list-expr)
```

If the object's class does not supply an augmenting method, then *default-expr* is evaluated, and the *arg* expressions are not evaluated. Otherwise, the augmenting method is called with the *arg* results as arguments, and *default-expr* is not evaluated. If no `inner` call is evaluated for a particular method, then augmenting methods supplied by subclasses are never used. Using the `inner` form outside of `class*` is an syntax error.

The second form is analogous to using `apply` with a procedure; the *arg-list-expr* must not be a parenthesized expression.

Inherited and Superclass Methods

Each `inherit`, `inherit/super`, `inherit/inner`, `rename-super`, and `rename-inner` clause declares one or more methods that are defined in the class, but must be present in the superclass. The `rename-super` and `rename-inner` declarations are rarely used, since `inherit/super` and `inherit/inner` provide the same access. Also, superclass and augmenting methods are typically accessed through `super` and `inner` in a class that also declares the methods, instead of through `inherit/super`, `inherit/inner`, `rename-super`, or `rename-inner`.

Method names declared with `inherit`, `inherit/super`, or `inherit/inner` access overriding declarations, if any, at run time. Method names declared with `inherit/super` can also be used with the `super` form to access the superclass implementation, and method names declared with `inherit/inner` can also be used with the `inner` form to access an augmenting method, if any.

Method names declared with `rename-super` always access the superclass's implementation

at run-time. Methods declared with `rename-inner` access a subclass's augmenting method, if any, and must be called with the form

```
(id (lambda () default-expr) arg ...)
```

so that a `default-expr` is available to evaluate when no augmenting method is available. In such a form, `lambda` is a literal identifier to separate the `default-expr` from the `arg`. When an augmenting method is available, it receives the results of the `arg` expressions as arguments.

Methods that are present in the superclass but not declared with `inherit`, `inherit/super`, or `inherit/inner` or `rename-super` are not directly accessible in the class (though they can be called with `send`). Every public method in a superclass is present in a derived class, even if it is not declared with `inherit` in the derived class; the `inherit` clause does not control inheritance, but merely controls lexical scope within a class expression.

If a method declared with `inherit`, `inherit/super`, `inherit/inner`, `rename-super`, or `rename-inner` is not present in the superclass, the `exn:fail:object` exception is raised when the class expression is evaluated.

Internal and External Names

Each method declared with `public`, `override`, `augment`, `pubment`, `overment`, `augride`, `public-final`, `override-final`, `augment-final`, `inherit`, `inherit/super`, `inherit/inner`, `rename-super`, and `rename-inner` can have separate internal and external names when `(internal-id external-id)` is used for declaring the method. The internal name is used to access the method directly within the class expression (including within `super` or `inner` forms), while the external name is used with `send` and `generic` (see §6.4 "Field and Method Access"). If a single `id` is provided for a method declaration, the identifier is used for both the internal and external names.

Method inheritance, overriding, and augmentation are based on external names only. Separate internal and external names are required for `rename-super` and `rename-inner` (for historical reasons, mainly).

Each `init`, `init-field`, `field`, or `inherit-field` variable similarly has an internal and an external name. The internal name is used within the class to access the variable, while the external name is used outside the class when providing initialization arguments (e.g., to `instantiate`), inheriting a field, or accessing a field externally (e.g., with `class-field-accessor`). As for methods, when inheriting a field with `inherit-field`, the external name is matched to an external field name in the superclass, while the internal name is bound in the `class` expression.

A single identifier can be used as an internal identifier and an external identifier, and it is possible to use the same identifier as internal and external identifiers for different bindings. Furthermore, within a single class, a single name can be used as an external method name, an external field name, and an external initialization argument name. Overall, each internal

identifier must be distinct from all other internal identifiers, each external method name must be distinct from all other method names, each external field name must be distinct from all other field names, and each initialization argument name must be distinct from all other initialization argument names.

By default, external names have no lexical scope, which means, for example, that an external method name matches the same syntactic symbol in all uses of `send`. The `define-local-member-name` and `define-member-name` forms introduce scoped external names.

When a `class` expression is compiled, identifiers used in place of external names must be symbolically distinct (when the corresponding external names are required to be distinct), otherwise a syntax error is reported. When no external name is bound by `define-member-name`, then the actual external names are guaranteed to be distinct when `class` expression is evaluated. When any external name is bound by `define-member-name`, the `exn:fail:object` exception is raised by `class` if the actual external names are not distinct.

```
(define-local-member-name id ...)
```

Unless it appears as the top-level definition, binds each *id* so that, within the scope of the definition, each use of each *id* as an external name is resolved to a hidden name generated by the `define-local-member-name` declaration. Thus, methods, fields, and initialization arguments declared with such external-name *id*s are accessible only in the scope of the `define-local-member-name` declaration. As a top-level definition, `define-local-member-name` binds *id* to its symbolic form.

The binding introduced by `define-local-member-name` is a syntax binding that can be exported and imported with `modules`. Each evaluation of a `define-local-member-name` declaration generates a distinct hidden name (except as a top-level definition). The `interface->method-names` procedure does not expose hidden names.

Examples:

```
(define-values (r o)
  (let ()
    (define-local-member-name m)
    (define c% (class object%
                 (define/public (m) 10)
                 (super-new)))
    (define o (new c%))

    (values (send o m)
            o)))

> r
10
> (send o m)
send: no such method
```

method name: m
class name: c%

```
(define-member-name id key-expr)
```

Maps a single external name to an external name that is determined by an expression. The value of *key-expr* must be the result of either a `member-name-key` expression or a `generate-member-key` call.

```
(member-name-key identifier)
```

Produces a representation of the external name for `id` in the environment of the `member-name-key` expression.

```
(generate-member-key) → member-name-key?
```

Produces a hidden name, just like the binding for `define-local-member-name`.

```
(member-name-key? v) → boolean?
  v : any/c
```

Returns `#t` for values produced by `member-name-key` and `generate-member-key`, `#f` otherwise.

```
(member-name-key=? a-key b-key) → boolean?
  a-key : member-name-key?
  b-key : member-name-key?
```

Produces `#t` if member-name keys *a-key* and *b-key* represent the same external name, `#f` otherwise.

```
(member-name-key-hash-code a-key) → integer?
  a-key : member-name-key?
```

Produces an integer hash code consistent with `member-name-key=?` comparisons, analogous to `equal-hash-code`.

Examples:

```
(define (make-c% key)
  (define-member-name m key)
  (class object%
    (define/public (m) 10)
    (super-new)))
```

```
> (send (new (make-c% (member-name-key m))) m)
10
> (send (new (make-c% (member-name-key p))) m)
send: no such method
  method name: m
  class name: eval:57:0
> (send (new (make-c% (member-name-key p))) p)
10

(define (fresh-c%)
  (let ([key (generate-member-key)])
    (values (make-c% key) key)))

(define-values (fc% key) (fresh-c%))

> (send (new fc%) m)
send: no such method
  method name: m
  class name: eval:57:0
> (let ()
    (define-member-name p key)
    (send (new fc%) p))
10
```

6.3 Creating Objects

The `make-object` procedure creates a new object with by-position initialization arguments, the `new` form creates a new object with by-name initialization arguments, and the `instantiate` form creates a new object with both by-position and by-name initialization arguments.

All fields in the newly created object are initially bound to the special `#<undefined>` value (see §4.18 “Void”). Initialization variables with default value expressions (and no provided value) are also initialized to `#<undefined>`. After argument values are assigned to initialization variables, expressions in `field` clauses, `init-field` clauses with no provided argument, `init` clauses with no provided argument, private field definitions, and other expressions are evaluated. Those expressions are evaluated as they appear in the class expression, from left to right.

Sometime during the evaluation of the expressions, superclass-declared initializations must be evaluated once by using the `super-make-object` procedure, `super-new` form, or `super-instantiate` form.

By-name initialization arguments to a class that have no matching initialization variable are implicitly added as by-name arguments to a `super-make-object`, `super-new`, or `super-instantiate` invocation, after the explicit arguments. If multiple initialization arguments

are provided for the same name, the first (if any) is used, and the unused arguments are propagated to the superclass. (Note that converted by-position arguments are always placed before explicit by-name arguments.) The initialization procedure for the `object%` class accepts zero initialization arguments; if it receives any by-name initialization arguments, then `exn:fail:object` exception is raised.

If the end of initialization is reached for any class in the hierarchy without invoking the superclass's initialization, the `exn:fail:object` exception is raised. Also, if superclass initialization is invoked more than once, the `exn:fail:object` exception is raised.

Fields inherited from a superclass are not initialized until the superclass's initialization procedure is invoked. In contrast, all methods are available for an object as soon as the object is created; the overriding of methods is not affected by initialization (unlike objects in C++).

```
(make-object class init-v ...) → object?
  class : class?
  init-v : any/c
```

Creates an instance of *class*. The *init-v*s are passed as initialization arguments, bound to the initialization variables of *class* for the newly created object as described in §6.2.1 "Initialization Variables". If *class* is not a class, the `exn:fail:contract` exception is raised.

```
(new class-expr (id by-name-expr) ...)
```

Creates an instance of the value of *class-expr* (which must be a class), and the value of each *by-name-expr* is provided as a by-name argument for the corresponding *id*.

```
(instantiate class-expr (by-pos-expr ...) (id by-name-expr) ...)
```

Creates an instance of the value of *class-expr* (which must be a class), and the values of the *by-pos-expr*s are provided as by-position initialization arguments. In addition, the value of each *by-name-expr* is provided as a by-name argument for the corresponding *id*.

```
super-make-object
```

Produces a procedure that takes by-position arguments an invokes superclass initialization. See §6.3 "Creating Objects" for more information.

```
(super-instantiate (by-pos-expr ...) (id by-expr ...) ...)
```

Invokes superclass initialization with the specified by-position and by-name arguments. See §6.3 "Creating Objects" for more information.

```
(super-new (id by-name-expr ...) ...)
```

Invokes superclass initialization with the specified by-name arguments. See §6.3 "Creating Objects" for more information.

6.4 Field and Method Access

In expressions within a class definition, the initialization variables, fields, and methods of the class are all part of the environment. Within a method body, only the fields and other methods of the class can be referenced; a reference to any other class-introduced identifier is a syntax error. Elsewhere within the class, all class-introduced identifiers are available, and fields and initialization variables can be mutated with `set!`.

6.4.1 Methods

Method names used within a class can only be used in the procedure position of an application expression; any other use is a syntax error.

To allow methods to be applied to lists of arguments, a method application can have the following form:

```
(method-id arg ... . arg-list-expr)
```

This form calls the method in a way analogous to `(apply method-id arg ... arg-list-expr)`. The `arg-list-expr` must not be a parenthesized expression.

Methods are called from outside a class with the `send`, `send/apply`, and `send/keyword-apply` forms.

```
(send obj-expr method-id arg ...)
(send obj-expr method-id arg ... . arg-list-expr)
```

Evaluates *obj-expr* to obtain an object, and calls the method with (external) name *method-id* on the object, providing the *arg* results as arguments. Each *arg* is as for `#%app`: either *arg-expr* or *keyword arg-expr*. In the second form, *arg-list-expr* cannot be a parenthesized expression.

If *obj-expr* does not produce an object, the `exn:fail:contract` exception is raised. If the object has no public method named *method-id*, the `exn:fail:object` exception is raised.

```
(send/apply obj-expr method-id arg ... arg-list-expr)
```

Like the dotted form of `send`, but *arg-list-expr* can be any expression.

```
(send/keyword-apply obj-expr method-id
                    keyword-list-expr value-list-expr
                    arg ... arg-list-expr)
```

Like `send/apply`, but with expressions for keyword and argument lists like `keyword-apply`.

```
(dynamic-send obj
              method-name
              v ...
              #:<kw> kw-arg ...) → any
  obj : object?
  method-name : symbol?
  v : any/c
  kw-arg : any/c
```

Calls the method on *obj* whose name matches *method-name*, passing along all given *v*s and *kw-arg*s.

```
(send* obj-expr msg ...+)

msg = (method-id arg ...)
    | (method-id arg ... . arg-list-expr)
```

Calls multiple methods (in order) of the same object. Each *msg* corresponds to a use of `send`.

For example,

```
(send* edit (begin-edit-sequence)
            (insert "Hello")
            (insert #\newline)
            (end-edit-sequence))
```

is the same as

```
(let ([o edit])
  (send o begin-edit-sequence)
  (send o insert "Hello")
  (send o insert #\newline)
  (send o end-edit-sequence))
```

```
(send+ obj-expr msg ...)

msg = (method-id arg ...)
    | (method-id arg ... . arg-list-expr)
```

Calls methods (in order) starting with the object produced by *obj-expr*. Each method call will be invoked on the result of the last method call, which is expected to be an object. Each *msg* corresponds to a use of `send`.

This is the functional analogue of `send*`.

Examples:

```
(define point%
  (class object%
    (super-new)
    (init-field [x 0] [y 0])
    (define/public (move-x dx)
      (new this% [x (+ x dx)]))
    (define/public (move-y dy)
      (new this% [y (+ y dy)]))))

> (send+ (new point%)
         (move-x 5)
         (move-y 7)
         (move-x 12))
(object:point% ...)
```

```
(with-method ((id (obj-expr method-id)) ...)
  body ...+)
```

Extracts methods from an object and binds a local name that can be applied directly (in the same way as declared methods within a class) for each method. Each *obj-expr* must produce an object, which must have a public method named by the corresponding *method-id*. The corresponding *id* is bound so that it can be applied directly (see §6.4.1 “Methods”).

Example:

```
(let ([s (new stack%)])
  (with-method ([push (s push!)]
                [pop (s pop!)])
    (push 10)
    (push 9)
    (pop)))
```

is the same as

```
(let ([s (new stack%)])
  (send s push! 10)
  (send s push! 9)
  (send s pop!))
```

6.4.2 Fields

```
(get-field id obj-expr)
```

Extracts the field with (external) name *id* from the value of *obj-expr*.

If *obj-expr* does not produce an object, the exn:fail:contract exception is raised. If the object has no *id* field, the exn:fail:object exception is raised.

```
(dynamic-get-field field-name obj) → any/c
  field-name : symbol?
  obj : object?
```

Extracts the field from *obj* with the (external) name that matches *field-name*. If the object has no field matching *field-name*, the exn:fail:object exception is raised.

```
(set-field! id obj-expr expr)
```

Sets the field with (external) name *id* from the value of *obj-expr* to the value of *expr*.

If *obj-expr* does not produce an object, the exn:fail:contract exception is raised. If the object has no *id* field, the exn:fail:object exception is raised.

```
(dynamic-set-field! field-name obj v) → void?
  field-name : symbol?
  obj : object?
  v : any/c
```

Sets the field from *obj* with the (external) name that matches *field-name* to *v*. If the object has no field matching *field-name*, the exn:fail:object exception is raised.

```
(field-bound? id obj-expr)
```

Produces #t if the object result of *obj-expr* has a field with (external) name *id*, #f otherwise.

If *obj-expr* does not produce an object, the exn:fail:contract exception is raised.

```
(class-field-accessor class-expr field-id)
```

Returns an accessor procedure that takes an instance of the class produced by *class-expr* and returns the value of the object's field with (external) name *field-id*.

If *class-expr* does not produce a class, the exn:fail:contract exception is raised. If the class has no *field-id* field, the exn:fail:object exception is raised.

```
(class-field-mutator class-expr field-id)
```

Returns a mutator procedure that takes an instance of the class produced by *class-expr* and a value, and sets the value of the object's field with (external) name *field-id* to the given value. The result is #<void>.

If *class-expr* does not produce a class, the exn:fail:contract exception is raised. If the class has no *field-id* field, the exn:fail:object exception is raised.

6.4.3 Generics

A *generic* can be used instead of a method name to avoid the cost of relocating a method by name within a class.

```
(generic class-or-interface-expr id)
```

Produces a generic that works on instances of the class or interface produced by *class-or-interface-expr* (or an instance of a class/interface derived from class-or-interface) to call the method with (external) name *id*.

If *class-or-interface-expr* does not produce a class or interface, the exn:fail:contract exception is raised. If the resulting class or interface does not contain a method named *id*, the exn:fail:object exception is raised.

```
(send-generic obj-expr generic-expr arg ...)
(send-generic obj-expr generic-expr arg ... . arg-list-expr)
```

Calls a method of the object produced by *obj-expr* as indicated by the generic produced by *generic-expr*. Each *arg* is as for #%app: either *arg-expr* or *keyword arg-expr*. The second form is analogous to calling a procedure with apply, where *arg-list-expr* is not a parenthesized expression.

If *obj-expr* does not produce an object, or if *generic-expr* does not produce a generic, the exn:fail:contract exception is raised. If the result of *obj-expr* is not an instance of the class or interface encapsulated by the result of *generic-expr*, the exn:fail:object exception is raised.

```
(make-generic type method-name) → generic?
  type : (or/c class? interface?)
  method-name : symbol?
```

Like the generic form, but as a procedure that accepts a symbolic method name.

6.5 Mixins

```
(mixin (interface-expr ...) (interface-expr ...)
  class-clause ...)
```

Produces a *mixin*, which is a procedure that encapsulates a class extension, leaving the superclass unspecified. Each time that a mixin is applied to a specific superclass, it produces a new derived class using the encapsulated extension.

The given class must implement interfaces produced by the first set of *interface-expr*s. The result of the procedure is a subclass of the given class that implements the interfaces produced by the second set of *interface-expr*s. The *class-clause*s are as for `class*`, to define the class extension encapsulated by the mixin.

Evaluation of a `mixin` form checks that the *class-clause*s are consistent with both sets of *interface-expr*s.

6.6 Traits

`(require racket/trait)` package: `base`

The bindings documented in this section are provided by the `racket/trait` library, not `racket/base` or `racket`.

A *trait* is a collection of methods that can be converted to a mixin and then applied to a class. Before a trait is converted to a mixin, the methods of a trait can be individually renamed, and multiple traits can be merged to form a new trait.

```
(trait trait-clause ...)

trait-clause = (public maybe-renamed ...)
             | (pubment maybe-renamed ...)
             | (public-final maybe-renamed ...)
             | (override maybe-renamed ...)
             | (overment maybe-renamed ...)
             | (override-final maybe-renamed ...)
             | (augment maybe-renamed ...)
             | (augride maybe-renamed ...)
             | (augment-final maybe-renamed ...)
             | (inherit maybe-renamed ...)
             | (inherit/super maybe-renamed ...)
             | (inherit/inner maybe-renamed ...)
             | method-definition
             | (field field-declaration ...)
             | (inherit-field maybe-renamed ...)
```

Creates a trait. The body of a `trait` form is similar to the body of a `class*` form, but restricted to non-private method definitions. In particular, the grammar of *maybe-renamed*, *method-definition*, and *field-declaration* are the same as for `class*`, and every *method-definition* must have a corresponding declaration (one of `public`, `over-`

ride, etc.). As in class, uses of method names in direct calls, super calls, and inner calls depend on bringing method names into scope via inherit, inherit/super, inherit/inner, and other method declarations in the same trait; an exception, compared to class is that overment binds a method name only in the corresponding method, and not in other methods of the same trait. Finally, macros such as public* and define/public work in trait as in class.

External identifiers in trait, trait-exclude, trait-exclude-field, trait-alias, trait-rename, and trait-rename-field forms are subject to binding via define-member-name and define-local-member-name. Although private methods or fields are not allowed in a trait form, they can be simulated by using a public or field declaration and a name whose scope is limited to the trait form.

```
(trait? v) → boolean?
  v : any/c
```

Returns #t if v is a trait, #f otherwise.

```
(trait->mixin tr) → (class? . -> . class?)
  tr : trait?
```

Converts a trait to a mixin, which can be applied to a class to produce a new class. An expression of the form

```
(trait->mixin
 (trait
   trait-clause ...))
```

is equivalent to

```
(lambda (%)
  (class %
    trait-clause ...
    (super-new)))
```

Normally, however, a trait's methods are changed and combined with other traits before converting to a mixin.

```
(trait-sum tr ...+) → trait?
  tr : trait?
```

Produces a trait that combines all of the methods of the given trs. For example,

```
(define t1
  (trait
```

```
    (define/public (m1) 1)))
(define t2
  (trait
    (define/public (m2) 2)))
(define t3 (trait-sum t1 t2))
```

creates a trait `t3` that is equivalent to

```
(trait
  (define/public (m1) 1)
  (define/public (m2) 2))
```

but `t1` and `t2` can still be used individually or combined with other traits.

When traits are combined with `trait-sum`, the combination drops `inherit`, `inherit/super`, `inherit/inner`, and `inherit-field` declarations when a definition is supplied for the same method or field name by another trait. The `trait-sum` operation fails (the `exn:fail:contract` exception is raised) if any of the traits to combine define a method or field with the same name, or if an `inherit/super` or `inherit/inner` declaration to be dropped is inconsistent with the supplied definition. In other words, declaring a method with `inherit`, `inherit/super`, or `inherit/inner`, does not count as defining the method; at the same time, for example, a trait that contains an `inherit/super` declaration for a method `m` cannot be combined with a trait that defines `m` as `augment`, since no class could satisfy the requirements of both `augment` and `inherit/super` when the trait is later converted to a mixin and applied to a class.

```
(trait-exclude trait-expr id)
```

Produces a new trait that is like the trait result of *`trait-expr`*, but with the definition of a method named by *`id`* removed; as the method definition is removed, either an `inherit`, `inherit/super`, or `inherit/inner` declaration is added:

- A method declared with `public`, `pubment`, or `public-final` is replaced with an `inherit` declaration.
- A method declared with `override` or `override-final` is replaced with an `inherit/super` declaration.
- A method declared with `augment`, `augride`, or `augment-final` is replaced with an `inherit/inner` declaration.
- A method declared with `overment` is not replaced with any `inherit` declaration.

If the trait produced by *`trait-expr`* has no method definition for *`id`*, the `exn:fail:contract` exception is raised.

```
(trait-exclude-field trait-expr id)
```

Produces a new trait that is like the trait result of *trait-expr*, but with the definition of a field named by *id* removed; as the field definition is removed, an `inherit-field` declaration is added.

```
(trait-alias trait-expr id new-id)
```

Produces a new trait that is like the trait result of *trait-expr*, but the definition and declaration of the method named by *id* is duplicated with the name *new-id*. The consistency requirements for the resulting trait are the same as for `trait-sum`, otherwise the `exn:fail:contract` exception is raised. This operation does not rename any other use of *id*, such as in method calls (even method calls to `identifier` in the cloned definition for *new-id*).

```
(trait-rename trait-expr id new-id)
```

Produces a new trait that is like the trait result of *trait-expr*, but all definitions and references to methods named *id* are replaced by definitions and references to methods named by *new-id*. The consistency requirements for the resulting trait are the same as for `trait-sum`, otherwise the `exn:fail:contract` exception is raised.

```
(trait-rename-field trait-expr id new-id)
```

Produces a new trait that is like the trait result of *trait-expr*, but all definitions and references to fields named *id* are replaced by definitions and references to fields named by *new-id*. The consistency requirements for the resulting trait are the same as for `trait-sum`, otherwise the `exn:fail:contract` exception is raised.

6.7 Object and Class Contracts

```
(class/c maybe-opaque member-spec ...)
```

```
maybe-opaque =
             | #:opaque

member-spec = method-spec
             | (field field-spec ...)
             | (init field-spec ...)
             | (init-field field-spec ...)
             | (inherit method-spec ...)
             | (inherit-field field-spec ...)
             | (super method-spec ...)
             | (inner method-spec ...)
             | (override method-spec ...)
             | (augment method-spec ...)
             | (augride method-spec ...)
             | (absent absent-spec ...)

method-spec = method-id
             | (method-id method-contract-expr)

 field-spec = field-id
             | (field-id contract-expr)

absent-spec = method-id
             | (field field-id ...)
```

Produces a contract for a class.

There are two major categories of contracts listed in a `class/c` form: external and internal contracts. External contracts govern behavior when an object is instantiated from a class or when methods or fields are accessed via an object of that class. Internal contracts govern behavior when method or fields are accessed within the class hierarchy. This separation allows for stronger contracts for class clients and weaker contracts for subclasses.

Method contracts must contain an additional initial argument which corresponds to the implicit `this` parameter of the method. This allows for contracts which discuss the state of the object when the method is called (or, for dependent contracts, in other parts of the contract). Alternative contract forms, such as `->m`, are provided as a shorthand for writing method contracts.

Methods and fields listed in an `absent` clause must *not* be present in the class.

A class contract can be specified to be *opaque* with the `#:opaque` keyword. An opaque class contract will only accept a class that defines exactly the external methods and fields specified by the contract. A contract error is raised if the contracted class contains any methods or fields that are not specified.

The external contracts are as follows:

- An external method contract without a tag describes the behavior of the implementation of *method-id* on method sends to an object of the contracted class. This contract will continue to be checked in subclasses until the contracted class's implementation is no longer the entry point for dynamic dispatch.

  If only the field name is present, this is equivalent to insisting only that the method is present in the class.

  Examples:

```
(define woody%
  (class object%
    (define/public (draw who)
      (format "reach for the sky, ~a" who))
    (super-new)))

> (define/contract woody+c%
    (class/c [draw (->m symbol? string?)])
    woody%)

> (send (new woody%) draw #f)
"reach for the sky, #f"
> (send (new woody+c%) draw 'zurg)
"reach for the sky, zurg"
> (send (new woody+c%) draw #f)
draw: contract violation
  expected: symbol?
  given: #f
  in: the 1st argument of
        the draw method in
        (class/c (draw (->m symbol? string?)))
  contract from: (definition woody+c%)
  contract on: woody+c%
  blaming: top-level
   (assuming the contract is correct)
  at: eval:68.0
```

- An external field contract, tagged with `field`, describes the behavior of the value contained in that field when accessed from outside the class. Since fields may be mutated, these contracts are checked on any external access (via `get-field`) and external mutations (via `set-field!`) of the field.

  If only the field name is present, this is equivalent to using the contract `any/c` (but it is checked more efficiently).

  Examples:

```
(define woody/hat%
  (class woody%
    (field [hat-location 'uninitialized])
    (define/public (lose-hat) (set! hat-location 'lost))
    (define/public (find-hat) (set! hat-location 'on-head))
    (super-new)))

> (define/contract woody/hat+c%
    (class/c [draw (->m symbol? string?)]
             [lose-hat (->m void?)]
             [find-hat (->m void?)]
             (field [hat-location (or/c 'on-head 'lost)]))
    woody/hat%)

> (get-field hat-location (new woody/hat%))
'uninitialized
> (let ([woody (new woody/hat+c%)])
    (send woody lose-hat)
    (get-field hat-location woody))
'lost
> (get-field hat-location (new woody/hat+c%))
woody/hat+c%: broke its contract
  promised: (or/c (quote on-head) (quote lost))
  produced: 'uninitialized
  in: the hat-location field in
      (class/c
        (draw (->m symbol? string?))
        (lose-hat (->m void?))
        (find-hat (->m void?))
        (field (hat-location
                (or/c 'on-head 'lost))))
  contract from: (definition woody/hat+c%)
  blaming: (definition woody/hat+c%)
   (assuming the contract is correct)
  at: eval:73.0
> (let ([woody (new woody/hat+c%)])
    (set-field! hat-location woody 'under-the-dresser))
woody/hat+c%: contract violation
  expected: (or/c (quote on-head) (quote lost))
  given: 'under-the-dresser
  in: the hat-location field in
      (class/c
        (draw (->m symbol? string?))
        (lose-hat (->m void?))
        (find-hat (->m void?))
        (field (hat-location
```

```
                    (or/c 'on-head 'lost))))
  contract from: (definition woody/hat+c%)
  blaming: top-level
   (assuming the contract is correct)
  at: eval:73.0
```

- An initialization argument contract, tagged with `init`, describes the expected behavior of the value paired with that name during class instantiation. The same name can be provided more than once, in which case the first such contract in the `class/c` form is applied to the first value tagged with that name in the list of initialization arguments, and so on.

  If only the initialization argument name is present, this is equivalent to using the contract `any/c` (but it is checked more efficiently).

  Examples:

```
(define woody/init-hat%
  (class woody%
    (init init-hat-location)
    (field [hat-location init-hat-location])
    (define/public (lose-hat) (set! hat-location 'lost))
    (define/public (find-hat) (set! hat-location 'on-head))
    (super-new)))

> (define/contract woody/init-hat+c%
    (class/c [draw (->m symbol? string?)]
             [lose-hat (->m void?)]
             [find-hat (->m void?)]
             (init [init-hat-location (or/c 'on-head 'lost)])
             (field [hat-location (or/c 'on-head 'lost)]))
    woody/init-hat%)

> (get-field hat-location
             (new woody/init-hat+c%
                  [init-hat-location 'lost]))
'lost
> (get-field hat-location
             (new woody/init-hat+c%
                  [init-hat-location 'slinkys-mouth]))
woody/init-hat+c%: contract violation
  expected: (or/c (quote on-head) (quote lost))
  given: 'slinkys-mouth
  in: the init-hat-location init argument in
      (class/c
       (draw (->m symbol? string?))
       (lose-hat (->m void?))
       (find-hat (->m void?))
```

```
        (init (init-hat-location
                (or/c 'on-head 'lost)))
        (field (hat-location
                 (or/c 'on-head 'lost))))
  contract from:
      (definition woody/init-hat+c%)
  blaming: top-level
   (assuming the contract is correct)
  at: eval:79.0
```

- The contracts listed in an `init-field` section are treated as if each contract appeared in an `init` section and a `field` section.

The internal contracts restrict the behavior of method calls made between classes and their subclasses; such calls are not controlled by the class contracts described above.

As with the external contracts, when a method or field name is specified but no contract appears, the contract is satisfied merely with the presence of the corresponding field or method.

- A method contract tagged with `inherit` describes the behavior of the method when invoked directly (i.e., via `inherit`) in any subclass of the contracted class. This contract, like external method contracts, applies until the contracted class's method implementation is no longer the entry point for dynamic dispatch.

  Examples:

```
> (new (class woody+c%
         (inherit draw)
         (super-new)
         (printf "woody sez: “~a”\n" (draw "evil dr pork-
chop"))))
woody sez: “reach for the sky, evil dr porkchop”
(object:eval:82:0 ...)
> (define/contract woody+c-inherit%
    (class/c (inherit [draw (->m symbol? string?)]))
    woody+c%)

> (new (class woody+c-inherit%
         (inherit draw)
         (printf "woody sez: ~a\n" (draw "evil dr pork-
chop"))))
draw: contract violation
  expected: symbol?
  given: "evil dr porkchop"
  in: the 1st argument of
      the draw method in
```

```
      (class/c
       (inherit (draw (->m symbol? string?))))
contract from: (definition woody+c-inherit%)
contract on: woody+c-inherit%
blaming: top-level
 (assuming the contract is correct)
at: eval:83.0
```

- A method contract tagged with `super` describes the behavior of *`method-id`* when called by the `super` form in a subclass. This contract only affects `super` calls in subclasses which call the contract class's implementation of *`method-id`*.

  This example shows how to extend the `draw` method so that if it is passed two arguments, it combines two calls to the original `draw` method, but with a contract the controls how the `super` methods must be invoked.

  Examples:

```
> (define/contract woody2+c%
    (class/c (super [draw (->m symbol? string?)]))
    (class woody%
      (define/override draw
        (case-lambda
          [(a) (super draw a)]
          [(a b) (string-append (super draw a)
                                " and "
                                (super draw b))]))
      (super-new)))

> (send (new woody2+c%) draw 'evil-dr-porkchop  'zurg)
"reach for the sky, evil-dr-porkchop and reach for the sky,
zurg"
> (send (new woody2+c%) draw "evil dr porkchop" "zurg")
"reach for the sky, evil dr porkchop and reach for the sky,
zurg"
```

  The last call signals an error blaming `woody2%` because there is no contract checking the initial `draw` call.

- A method contract tagged with `inner` describes the behavior the class expects of an augmenting method in a subclass. This contract affects any implementations of *`method-id`* in subclasses which can be called via `inner` from the contracted class. This means a subclass which implements *`method-id`* via `augment` or `overment` stop future subclasses from being affected by the contract, since further extension cannot be reached via the contracted class.

- A method contract tagged with `override` describes the behavior expected by the contracted class for *`method-id`* when called directly (i.e. by the application `(`*`method-id`* `...)`). This form can only be used if overriding the method in subclasses will change

the entry point to the dynamic dispatch chain (i.e., the method has never been augmentable).

This time, instead of overriding `draw` to support two arguments, we can make a new method, `draw2` that takes the two arguments and calls `draw`. We also add a contract to make sure that overriding `draw` doesn't break `draw2`.

Examples:

```
> (define/contract woody2+override/c%
    (class/c (override [draw (->m symbol? string?)]))
    (class woody+c%
      (inherit draw)
      (define/public (draw2 a b)
        (string-append (draw a)
                       " and "
                       (draw b)))
      (super-new)))

(define woody2+broken-draw
  (class woody2+override/c%
    (define/override (draw x)
      'not-a-string)
    (super-new)))

> (send (new woody2+broken-draw) draw2
        'evil-dr-porkchop
        'zurg)
draw: broke its contract
  promised: string?
  produced: 'not-a-string
  in: the range of
      the draw method in
      (class/c
       (override (draw (->m symbol? string?))))
  contract from:
      (definition woody2+override/c%)
  contract on: woody2+override/c%
  blaming: top-level
   (assuming the contract is correct)
  at: eval:88.0
```

- A method contract tagged with either `augment` or `augride` describes the behavior provided by the contracted class for *`method-id`* when called directly from subclasses. These forms can only be used if the method has previously been augmentable, which means that no augmenting or overriding implementation will change the entry point to the dynamic dispatch chain. `augment` is used when subclasses can augment the method, and `augride` is used when subclasses can override the current augmentation.

- A field contract tagged with `inherit-field` describes the behavior of the value contained in that field when accessed directly (i.e., via `inherit-field`) in any subclass of the contracted class. Since fields may be mutated, these contracts are checked on any access and/or mutation of the field that occurs in such subclasses.

```
(absent absent-spec ...)
```

See `class/c`; use outside of a `class/c` form is a syntax error.

```
(->m dom ... range)
```

Similar to `->`, except that the domain of the resulting contract contains one more element than the stated domain, where the first (implicit) argument is contracted with `any/c`. This contract is useful for writing simpler method contracts when no properties of `this` need to be checked.

```
(->*m (mandatory-dom ...) (optional-dom ...) rest range)
```

Similar to `->*`, except that the mandatory domain of the resulting contract contains one more element than the stated domain, where the first (implicit) argument is contracted with `any/c`. This contract is useful for writing simpler method contracts when no properties of `this` need to be checked.

```
(case->m (-> dom ... rest range) ...)
```

Similar to `case->`, except that the mandatory domain of each case of the resulting contract contains one more element than the stated domain, where the first (implicit) argument is contracted with `any/c`. This contract is useful for writing simpler method contracts when no properties of `this` need to be checked.

```
(->dm (mandatory-dependent-dom ...)
      (optional-dependent-dom ...)
      dependent-rest
      pre-cond
      dep-range)
```

Similar to `->d`, except that the mandatory domain of the resulting contract contains one more element than the stated domain, where the first (implicit) argument is contracted with `any/c`. In addition, `this` is appropriately bound in the body of the contract. This contract is useful for writing simpler method contracts when no properties of `this` need to be checked.

```
(object/c member-spec ...)
```

```
member-spec = method-spec
            | (field field-spec ...)

method-spec = method-id
            | (method-id method-contract)

 field-spec = field-id
            | (field-id contract-expr)
```

Produces a contract for an object.

Unlike the older form `object-contract`, but like `class/c`, arbitrary contract expressions are allowed. Also, method contracts for `object/c` follow those for `class/c`. An object wrapped with `object/c` behaves as if its class had been wrapped with the equivalent `class/c` contract.

```
(instanceof/c class-contract) → contract?
  class-contract : contract?
```

Produces a contract for an object, where the object is an instance of a class that conforms to *class-contract*.

```
(object-contract member-spec ...)
```

```
           member-spec = (method-id method-contract)
                       | (field field-id contract-expr)

        method-contract = (-> dom ... range)
                       | (->* (mandatory-dom ...)
                              (optional-dom ...)
                              rest
                              range)
                       | (->d (mandatory-dependent-dom ...)
                              (optional-dependent-dom ...)
                              dependent-rest
                              pre-cond
                              dep-range)

                   dom = dom-expr
                       | keyword dom-expr

                 range = range-expr
                       | (values range-expr ...)
                       | any

          mandatory-dom = dom-expr
                       | keyword dom-expr

           optional-dom = dom-expr
                       | keyword dom-expr

                  rest =
                       | #:rest rest-expr

mandatory-dependent-dom = [id dom-expr]
                       | keyword [id dom-expr]

 optional-dependent-dom = [id dom-expr]
                       | keyword [id dom-expr]

         dependent-rest =
                       | #:rest id rest-expr

              pre-cond =
                       | #:pre-cond boolean-expr

             dep-range = any
                       | [id range-expr] post-cond
                       | (values [id range-expr] ...) post-cond

             post-cond =
                       | #:post-cond boolean-expr
```

Produces a contract for an object.

Each of the contracts for a method has the same semantics as the corresponding function contract, but the syntax of the method contract must be written directly in the body of the object-contract—much like the way that methods in class definitions use the same syntax as regular function definitions, but cannot be arbitrary procedures. Unlike the method contracts for `class/c`, the implicit `this` argument is not part of the contract. To allow for the use of `this` in dependent contracts, `->d` contracts implicitly bind `this` to the object itself.

```
mixin-contract : contract?
```

A function contract that recognizes mixins. It guarantees that the input to the function is a class and the result of the function is a subclass of the input.

```
(make-mixin-contract type ...) → contract?
  type : (or/c class? interface?)
```

Produces a function contract that guarantees the input to the function is a class that implements/subclasses each *type*, and that the result of the function is a subclass of the input.

```
(is-a?/c type) → flat-contract?
  type : (or/c class? interface?)
```

Accepts a class or interface and returns a flat contract that recognizes objects that instantiate the class/interface.

```
(implementation?/c interface) → flat-contract?
  interface : interface?
```

Returns a flat contract that recognizes classes that implement *interface*.

```
(subclass?/c class) → flat-contract?
  class : class?
```

Returns a flat contract that recognizes classes that are subclasses of *class*.

6.8 Object Equality and Hashing

By default, objects that are instances of different classes or that are instances of a non-transparent class are `equal?` only if they are `eq?`. Like transparent structures, two objects that are instances of the same transparent class (i.e., every superclass of the class has `#f` as its inspector) are `equal?` when their field values are `equal?`.

To customize the way that a class instance is compared to other instances by equal?, implement the equal<%> interface.

equal<%> : interface?

The equal<%> interface includes three methods, which are analogous to the functions provided for a structure type with prop:equal+hash:

- equal-to? — Takes two arguments. The first argument is an object that is an instance of the same class (or a subclass that does not re-declare its implementation of equal<%>) and that is being compared to the target object. The second argument is an equal?-like procedure of two arguments that should be used for recursive equality testing. The result should be a true value if the object and the first argument of the method are equal, #f otherwise.
- equal-hash-code-of — Takes one argument, which is a procedure of one argument that should be used for recursive hash-code computation. The result should be an exact integer representing the target object's hash code.
- equal-secondary-hash-code-of — Takes one argument, which is a procedure of one argument that should be used for recursive hash-code computation. The result should be an exact integer representing the target object's secondary hash code.

The equal<%> interface is unusual in that declaring the implementation of the interface is different from inheriting the interface. Two objects can be equal only if they are instances of classes whose most specific ancestor to explicitly implement equal<%> is the same ancestor.

See prop:equal+hash for more information on equality comparisons and hash codes. The equal<%> interface is implemented with interface* and prop:equal+hash.

Example:

```
#lang racket

;; Case insensitive words:
(define ci-word%
  (class* object% (equal<%>)

    ;; Initialization
    (init-field word)
    (super-new)

    ;; We define equality to ignore case:
```

```
    (define/public (equal-to? other recur)
      (string-ci=? word (get-field word other)))

    ;; The hash codes need to be insensitive to casing as well.
    ;; We'll just downcase the word and get its hash code.
    (define/public (equal-hash-code-of hash-code)
      (hash-code (string-downcase word)))

    (define/public (equal-secondary-hash-code-of hash-code)
      (hash-code (string-downcase word)))))

;; We can create a hash with a single word:
(define h (make-hash))
(hash-set! h (new ci-word% [word "inconceivable!"]) 'value)

;; Lookup into the hash should be case-insensitive, so that
;; both of these should return 'value.
(hash-ref h (new ci-word% [word "inconceivable!"]))
(hash-ref h (new ci-word% [word "INCONCEIVABLE!"]))

;; Comparison fails if we use a non-ci-word%:
(hash-ref h "inconceivable!" 'i-dont-think-it-means-what-you-
think-it-means)
```

6.9 Object Serialization

```
(define-serializable-class* class-id superclass-expr
                            (interface-expr ...)
  class-clause ...)
```

Binds *class-id* to a class, where *superclass-expr*, the *interface-expr*s, and the *class-clause*s are as in `class*`.

This form can only be used at the top level, either within a module or outside. The *class-id* identifier is bound to the new class, and `deserialize-info:`*class-id* is also defined; if the definition is within a module, then the latter is provided from a `deserialize-info` submodule via `module+`.

Serialization for the class works in one of two ways:

- If the class implements the built-in interface `externalizable<%>`, then an object is serialized by calling its `externalize` method; the result can be anything that is serializable (but, obviously, should not be the object itself). Deserialization creates an

instance of the class with no initialization arguments, and then calls the object's `internalize` method with the result of `externalize` (or, more precisely, a deserialized version of the serialized result of a previous call).

To support this form of serialization, the class must be instantiable with no initialization arguments. Furthermore, cycles involving only instances of the class (and other such classes) cannot be serialized.

- If the class does not implement `externalizable<%>`, then every superclass of the class must be either serializable or transparent (i.e,. have `#f` as its inspector). Serialization and deserialization are fully automatic, and may involve cycles of instances.

  To support cycles of instances, deserialization may create an instance of the call with all fields as the undefined value, and then mutate the object to set the field values. Serialization support does not otherwise make an object's fields mutable.

In the second case, a serializable subclass can implement `externalizable<%>`, in which case the `externalize` method is responsible for all serialization (i.e., automatic serialization is lost for instances of the subclass). In the first case, all serializable subclasses implement `externalizable<%>`, since a subclass implements all of the interfaces of its parent class.

In either case, if an object is an immediate instance of a subclass (that is not itself serializable), the object is serialized as if it was an immediate instance of the serializable class. In particular, overriding declarations of the `externalize` method are ignored for instances of non-serializable subclasses.

```
(define-serializable-class class-id superclass-expr
  class-clause ...)
```

Like `define-serializable-class*`, but without interface expressions (analogous to `class`).

```
externalizable<%> : interface?
```

The `externalizable<%>` interface includes only the `externalize` and `internalize` methods. See `define-serializable-class*` for more information.

6.10 Object Printing

To customize the way that a class instance is printed by `print`, `write` and `display`, implement the `printable<%>` interface.

```
printable<%> : interface?
```

The printable<%> interface includes only the custom-print, custom-write, and custom-display methods. The custom-print method accepts two arguments: the destination port and the current quasiquote depth as an exact nonnegative integer. The custom-write and custom-display methods each accepts a single argument, which is the destination port to write or display the object.

Calls to the custom-print, custom-write, or custom-display methods are like calls to a procedure attached to a structure type through the prop:custom-write property. In particular, recursive printing can trigger an escape from the call.

See prop:custom-write for more information. The printable<%> interface is implemented with interface* and prop:custom-write.

```
writable<%> : interface?
```

Like printable<%>, but includes only the custom-write and custom-display methods. A print request is directed to custom-write.

6.11 Object, Class, and Interface Utilities

```
(object? v) → boolean?
  v : any/c
```

Returns #t if v is an object, #f otherwise.

Examples:

```
> (object? (new object%))
#t
> (object? object%)
#f
> (object? "clam chowder")
#f
```

```
(class? v) → boolean?
  v : any/c
```

Returns #t if v is a class, #f otherwise.

Examples:

```
> (class? object%)
#t
> (class? (class object% (super-new)))
```

```
#t
> (class? (new object%))
#f
> (class? "corn chowder")
#f
```

```
(interface? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is an interface, `#f` otherwise.

Examples:

```
> (interface? (interface () empty cons first rest))
#t
> (interface? object%)
#f
> (interface? "gazpacho")
#f
```

```
(generic? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a generic, `#f` otherwise.

Examples:

```
(define c%
  (class object%
    (super-new)
    (define/public (m x)
      (+ 3.14 x))))

> (generic? (generic c% m))
#t
> (generic? c%)
#f
> (generic? "borscht")
#f
```

```
(object=? a b) → boolean?
  a : object?
  b : object?
```

Determines whether `a` and `b` were returned from the same call to `new` or not. If the two objects have fields, this procedure determines whether mutating a field of one would change that field in the other.

This procedure is similar in spirit to eq? but also works properly with contracts (and has a stronger guarantee).

Examples:

```
(define obj-1 (new object%))

(define obj-2 (new object%))

> (define/contract obj-3 (object/c) obj-1)

> (object=? obj-1 obj-1)
#t
> (object=? obj-1 obj-2)
#f
> (object=? obj-1 obj-3)
#t
> (eq? obj-1 obj-1)
#t
> (eq? obj-1 obj-2)
#f
> (eq? obj-1 obj-3)
#f
```

```
(object->vector object [opaque-v]) → vector?
  object : object?
  opaque-v : any/c = #f
```

Returns a vector representing *object* that shows its inspectable fields, analogous to struct->vector.

Examples:

```
> (object->vector (new object%))
'#(object:object% ...)
> (object->vector (new (class object%
                         (super-new)
                         (field [x 5] [y 10]))))
'#(object:eval:106:0 ...)
```

```
(class->interface class) → interface?
  class : class?
```

Returns the interface implicitly defined by *class*.

Example:

```
> (class->interface object%)
#<interface:object%>
```

```
(object-interface object) → interface?
  object : object?
```

Returns the interface implicitly defined by the class of *object*.

Example:

```
> (object-interface (new object%))
#<interface:object%>
```

```
(is-a? v type) → boolean?
  v : any/c
  type : (or/c interface? class?)
```

Returns #t if *v* is an instance of a class *type* or a class that implements an interface *type*, #f otherwise.

Examples:

```
(define point<%> (interface () get-x get-y))

(define 2d-point%
  (class* object% (point<%>)
    (super-new)
    (field [x 0] [y 0])
    (define/public (get-x) x)
    (define/public (get-y) y)))

> (is-a? (new 2d-point%) 2d-point%)
#t
> (is-a? (new 2d-point%) point<%>)
#t
> (is-a? (new object%) 2d-point%)
#f
> (is-a? (new object%) point<%>)
#f
```

```
(subclass? v cls) → boolean?
  v : any/c
  cls : class?
```

Returns #t if *v* is a class derived from (or equal to) *cls*, #f otherwise.

Examples:

```
> (subclass? (class object% (super-new)) object%)
#t
> (subclass? object% (class object% (super-new)))
#f
> (subclass? object% object%)
#t
```

```
(implementation? v intf) → boolean?
  v : any/c
  intf : interface?
```

Returns `#t` if *v* is a class that implements *intf*, `#f` otherwise.

Examples:

```
(define i<%> (interface () go))

(define c%
  (class* object% (i<%>)
    (super-new)
    (define/public (go) 'go)))

> (implementation? c% i<%>)
#t
> (implementation? object% i<%>)
#f
```

```
(interface-extension? v intf) → boolean?
  v : any/c
  intf : interface?
```

Returns `#t` if *v* is an interface that extends *intf*, `#f` otherwise.

Examples:

```
(define point<%> (interface () get-x get-y))

(define colored-point<%> (interface (point<%>) color))

> (interface-extension? colored-point<%> point<%>)
#t
> (interface-extension? point<%> colored-point<%>)
#f
> (interface-extension? (interface () get-x get-y get-z) point<%>)
#f
```

```
(method-in-interface? sym intf) → boolean?
  sym : symbol?
  intf : interface?
```

Returns `#t` if *intf* (or any of its ancestor interfaces) includes a member with the name *sym*, `#f` otherwise.

Examples:

```
(define i<%> (interface () get-x get-y))

> (method-in-interface? 'get-x i<%>)
#t
> (method-in-interface? 'get-z i<%>)
#f
```

```
(interface->method-names intf) → (listof symbol?)
  intf : interface?
```

Returns a list of symbols for the method names in *intf*, including methods inherited from superinterfaces, but not including methods whose names are local (i.e., declared with `define-local-member-name`).

Examples:

```
(define i<%> (interface () get-x get-y))

> (interface->method-names i<%>)
'(get-x get-y)
```

```
(object-method-arity-includes? object
                               sym
                               cnt)    → boolean?
  object : object?
  sym : symbol?
  cnt : exact-nonnegative-integer?
```

Returns `#t` if *object* has a method named *sym* that accepts *cnt* arguments, `#f` otherwise.

Examples:

```
(define c%
  (class object%
    (super-new)
    (define/public (m x [y 0])
      (+ x y))))

> (object-method-arity-includes? (new c%) 'm 1)
#t
> (object-method-arity-includes? (new c%) 'm 2)
#t
```

```
> (object-method-arity-includes? (new c%) 'm 3)
#f
> (object-method-arity-includes? (new c%) 'n 1)
#f
```

```
(field-names object) → (listof symbol?)
  object : object?
```

Returns a list of all of the names of the fields bound in *object*, including fields inherited from superinterfaces, but not including fields whose names are local (i.e., declared with `define-local-member-name`).

Examples:

```
> (field-names (new object%))
'()
> (field-names (new (class object% (super-
new) (field [x 0] [y 0]))))
'(y x)
```

```
(object-info object) → (or/c class? #f) boolean?
  object : object?
```

Returns two values, analogous to the return values of `struct-info`:

- *class*: a class or `#f`; the result is `#f` if the current inspector does not control any class for which the *object* is an instance.
- *skipped?*: `#f` if the first result corresponds to the most specific class of *object*, `#t` otherwise.

```
(class-info class)
    symbol?
    exact-nonnegative-integer?
    (listof symbol?)
 → (any/c exact-nonnegative-integer? . -> . any/c)
    (any/c exact-nonnegative-integer? any/c . -> . any/c)
    (or/c class? #f)
    boolean?
  class : class?
```

Returns seven values, analogous to the return values of `struct-type-info`:

- *name*: the class's name as a symbol;

- *field-cnt*: the number of fields (public and private) defined by the class;
- *field-name-list*: a list of symbols corresponding to the class's public fields; this list can be larger than *field-cnt* because it includes inherited fields;
- *field-accessor*: an accessor procedure for obtaining field values in instances of the class; the accessor takes an instance and a field index between 0 (inclusive) and *field-cnt* (exclusive);
- *field-mutator*: a mutator procedure for modifying field values in instances of the class; the mutator takes an instance, a field index between 0 (inclusive) and *field-cnt* (exclusive), and a new field value;
- *super-class*: a class for the most specific ancestor of the given class that is controlled by the current inspector, or #f if no ancestor is controlled by the current inspector;
- *skipped?*: #f if the sixth result is the most specific ancestor class, #t otherwise.

```
(struct exn:fail:object exn:fail ()
    #:extra-constructor-name make-exn:fail:object)
```

Raised for `class`-related failures, such as attempting to call a method that is not supplied by an object.

6.12 Surrogates

`(require racket/surrogate)` package: `base`

The bindings documented in this section are provided by the `racket/surrogate` library, not `racket/base` or `racket`.

The `racket/surrogate` library provides an abstraction for building an instance of the *proxy design pattern*. The pattern consists of two objects, a *host* and a *surrogate* object. The host object delegates method calls to its surrogate object. Each host has a dynamically assigned surrogate, so an object can completely change its behavior merely by changing the surrogate.

```
(surrogate method-spec ...)

method-spec = (method-id arg-spec ...)
            | (override method-id arg-spec ...)
            | (override-final method-id (lambda () default-expr)
                             arg-spec ...)

  arg-spec = (id ...)
           | id
```

If neither `override` nor `override-final` is specified for a *`method-id`*, then `override` is assumed.

The `surrogate` form produces four values: a host mixin (a procedure that accepts and returns a class), a host interface, a surrogate class, and a surrogate interface.

The host mixin adds one additional field, `surrogate`, to its argument. It also adds a getter method, `get-surrogate`, and a setter method, `set-surrogate`, for changing the field. The `set-surrogate` method accepts instances of the class returned by the `surrogate` form or `#f`, and it updates the field with its argument; then, `set-surrogate` calls the `on-disable-surrogate` on the previous value of the field and `on-enable-surrogate` for the new value of the field. The `get-surrogate` method returns the current value of the field.

The host mixin has a single overriding method for each *`method-id`* in the `surrogate` form. Each of these methods is defined with a `case-lambda` with one arm for each *`arg-spec`*. Each arm has the variables as arguments in the *`arg-spec`*. The body of each method tests the `surrogate` field. If it is `#f`, the method just returns the result of invoking the super or inner method. If the `surrogate` field is not `#f`, the corresponding method of the object in the field is invoked. This method receives the same arguments as the original method, plus two extras. The extra arguments come at the beginning of the argument list. The first is the original object. The second is a procedure that calls the super or inner method (i.e., the method of the class that is passed to the mixin or an extension, or the method in an overriding class), with the arguments that the procedure receives.

For example, the host-mixin for this surrogate:

```
(surrogate (override m (x y z)))
```

will override the `m` method and call the surrogate like this:

```
(define/override (m x y z)
  (if surrogate
      (send surrogate m
            this
            (λ (x y z) (super m x y z))
            x y z)
      (super m x y z)))
```

where *`surrogate`* is bound to the value most recently passed to the host mixin's *`set-surrogate`* method.

The host interface has the names `set-surrogate`, `get-surrogate`, and each of the *`method-id`*s in the original form.

The surrogate class has a single public method for each *`method-id`* in the `surrogate` form. These methods are invoked by classes constructed by the mixin. Each has a corresponding

method signature, as described in the above paragraph. Each method just passes its argument along to the super procedure it receives.

In the example above, this is the *m* method in the surrogate class:

```
(define/public (m original-object original-super x y z)
  (original-super x y z))
```

Note: if you derive a class from the surrogate class, do not both call the `super` argument and the super method of the surrogate class itself. Only call one or the other, since the default methods call the `super` argument.

Finally, the interface contains all of the names specified in surrogate's argument, plus `on-enable-surrogate` and `on-disable-surrogate`. The class returned by `surrogate` implements this interface.

7 Units

§14 “Units” in ***The Racket Guide*** introduces units.

Units organize a program into separately compilable and reusable components. The imports and exports of a unit are grouped into a *signature*, which can include “static” information (such as macros) in addition to placeholders for run-time values. Units with suitably matching signatures can be *linked* together to form a larger unit, and a unit with no imports can be *invoked* to execute its body.

```
(require racket/unit)
```
package: base

The bindings documented in this section are provided by the `racket/unit` and `racket` libraries, but not `racket/base`. The `racket/unit` module name can be used as a language name with `#lang`; see §7.10 “Single-Unit Modules”.

7.1 Creating Units

```
(unit
  (import tagged-sig-spec ...)
  (export tagged-sig-spec ...)
  init-depends-decl
  unit-body-expr-or-defn
  ...)

  tagged-sig-spec = sig-spec
                  | (tag id sig-spec)

        sig-spec = sig-id
                 | (prefix id sig-spec)
                 | (rename sig-spec (id id) ...)
                 | (only sig-spec id ...)
                 | (except sig-spec id ...)

init-depends-decl =
                  | (init-depend tagged-sig-id ...)

    tagged-sig-id = sig-id
                  | (tag id sig-id)
```

Produces a unit that encapsulates its *unit-body-expr-or-defn*s. Expressions in the `unit` body can refer to identifiers bound by the *sig-spec*s of the `import` clause, and the body must include one definition for each identifier of a *sig-spec* in the `export` clause. An identifier that is exported cannot be `set!`ed in either the defining unit or in importing units, although the implicit assignment to initialize the variable may be visible as a mutation.

Each import or export *sig-spec* ultimately refers to a *sig-id*, which is an identifier that is bound to a signature by `define-signature`.

In a specific import or export position, the set of identifiers bound or required by a particular *sig-id* can be adjusted in a few ways:

- `(prefix` *id sig-spec*`)` as an import binds the same as *sig-spec*, except that each binding is prefixed with *id*. As an export, this form causes definitions using the *id* prefix to satisfy the exports required by *sig-spec*.
- `(rename` *sig-spec* `(`*id id*`) ...)` as an import binds the same as *sig-spec*, except that the first *id* is used for the binding instead of the second *id* (where *sig-spec* by itself must imply a binding that is `bound-identifier=?` to second *id*). As an export, this form causes a definition for the first *id* to satisfy the export named by the second *id* in *sig-spec*.
- `(only` *sig-spec id ...*`)` as an import binds the same as *sig-spec*, but restricted to just the listed *id*s (where *sig-spec* by itself must imply a binding that is `bound-identifier=?` to each *id*). This form is not allowed for an export.
- `(except` *sig-spec id ...*`)` as an import binds the same as *sig-spec*, but excluding all listed *id*s (where *sig-spec* by itself must imply a binding that is `bound-identifier=?` to each *id*). This form is not allowed for an export.

As suggested by the grammar, these adjustments to a signature can be nested arbitrarily.

A unit's declared imports are matched with actual supplied imports by signature. That is, the order in which imports are supplied to a unit when linking is irrelevant; all that matters is the signature implemented by each supplied import. One actual import must be provided for each declared import. Similarly, when a unit implements multiple signatures, the order of the export signatures does not matter.

To support multiple imports or exports for the same signature, an import or export can be tagged using the form `(tag` *id sig-spec*`)`. When an import declaration of a unit is tagged, then one actual import must be given the same tag (with the same signature) when the unit is linked. Similarly, when an export declaration is tagged for a unit, then references to that particular export must explicitly use the tag.

A unit is prohibited syntactically from importing two signatures that are not distinct, unless they have different tags; two signatures are *distinct* only if they share no ancestor through `extends`. The same syntactic constraint applies to exported signatures. In addition, a unit is prohibited syntactically from importing the same identifier twice (after renaming and other transformations on a *sig-spec*), exporting the same identifier twice (again, after renaming), or exporting an identifier that is imported.

When units are linked, the bodies of the linked units are executed in an order that is specified at the linking site. An optional `(init-depend` *tagged-sig-id ...*`)` declaration constrains the allowed orders of linking by specifying that the current unit must be initialized

after the unit that supplies the corresponding import. Each *tagged-sig-id* in an `init-depend` declaration must have a corresponding import in the `import` clause.

```
(define-signature id extension-decl
  (sig-elem ...))

extension-decl =
               | extends sig-id

     sig-elem = id
              | (define-syntaxes (id ...) expr)
              | (define-values (id ...) expr)
              | (define-values-for-export (id ...) expr)
              | (contracted [id contract] ...)
              | (open sig-spec)
              | (struct id (field ...) struct-option ...)
              | (sig-form-id . datum)

        field = id
              | [id #:mutable]

 struct-option = #:mutable
               | #:constructor-name constructor-id
               | #:extra-constructor-name constructor-id
               | #:omit-constructor
               | #:omit-define-syntaxes
               | #:omit-define-values
```

Binds an identifier to a signature that specifies a group of bindings for import or export:

- Each *id* in a signature declaration means that a unit implementing the signature must supply a variable definition for the *id*. That is, *id* is available for use in units importing the signature, and *id* must be defined by units exporting the signature.
- Each `define-syntaxes` form in a signature declaration introduces a macro that is available for use in any unit that imports the signature. Free variables in the definition's *expr* refer to other identifiers in the signature first, or the context of the `define-signature` form if the signature does not include the identifier.
- Each `define-values` form in a signature declaration introduces code that effectively prefixes every unit that imports the signature. Free variables in the definition's *expr* are treated the same as for `define-syntaxes`.
- Each `define-values-for-export` form in a signature declaration introduces code that effectively suffixes every unit that exports the signature. Free variables in the definition's *expr* are treated the same as for `define-syntaxes`.

- Each `contracted` form in a signature declaration means that a unit exporting the signature must supply a variable definition for each *`id`* in that form. If the signature is imported, then uses of *`id`* inside the unit are protected by the appropriate contracts using the unit as the negative blame. If the signature is exported, then the exported values are protected by the appropriate contracts which use the unit as the positive blame, but internal uses of the exported identifiers are not protected. Variables in the *`contract`* expressions are treated the same as for `define-syntaxes`.
- Each `(open` *`sig-spec`*`)` adds to the signature everything specified by *`sig-spec`*.
- Each `(struct` *`id`* `(`*`field`* `...)` *`struct-option`* `...)` adds all of the identifiers that would be bound by `(struct` *`id`* `(`*`field`* `...)` `field-option` `...)`, where the extra option `#:omit-constructor` omits the constructor identifier.
- Each `(`*`sig-form-id`* `.` *`datum`*`)` extends the signature in a way that is defined by *`sig-form-id`*, which must be bound by `define-signature-form`. One such binding is for `struct/ctc`.

When a `define-signature` form includes an `extends` clause, then the define signature automatically includes everything in the extended signature. Furthermore, any implementation of the new signature can be used as an implementation of the extended signature.

`(open` *`sig-spec`*`)`

Allowed only in a *`sig-elem`*; see `define-signature`.

`(define-values-for-export (`*`id`* `...)` *`expr`*`)`

Allowed only in a *`sig-elem`*; see `define-signature`.

`(contracted [`*`id`* *`contract`*`] ...)`

Allowed only in a *`sig-elem`*; see `define-signature`.

`(only` *`sig-spec`* *`id`* `...)`

Allowed only in a *`sig-spec`*; see `unit`.

`(except` *`sig-spec`* *`id`* `...)`

Allowed only in a *`sig-spec`*; see `unit`.

```
(rename sig-spec (id id) ...)
```

Allowed only in a *sig-spec*; see `unit`.

```
(prefix id sig-spec)
```

Allowed only in a *sig-spec*; see `unit`.

```
(import tagged-sig-spec ...)
```

Allowed only in certain forms; see, for example, `unit`.

```
(export tagged-sig-spec ...)
```

Allowed only in certain forms; see, for example, `unit`.

```
(link linkage-decl ...)
```

Allowed only in certain forms; see, for example, `compound-unit`.

```
(tag id sig-spec)
(tag id sig-id)
```

Allowed only in certain forms; see, for example, `unit`.

```
(init-depend tagged-sig-id ...)
```

Allowed only in a `init-depend-decl`; see `unit`.

```
extends
```

Allowed only within `define-signature`.

7.2 Invoking Units

```
(invoke-unit unit-expr)
(invoke-unit unit-expr (import tagged-sig-spec ...))
```

Invokes the unit produced by *unit-expr*. For each of the unit's imports, the `invoke-unit` expression must contain a *tagged-sig-spec* in the `import` clause; see `unit` for the grammar of *tagged-sig-spec*. If the unit has no imports, the `import` clause can be omitted.

When no *tagged-sig-spec*s are provided, *unit-expr* must produce a unit that expects no imports. To invoke the unit, all bindings are first initialized to the `#<undefined>` value. Next, the unit's body definitions and expressions are evaluated in order; in the case of a definition, evaluation sets the value of the corresponding variable(s). Finally, the result of the last expression in the unit is the result of the `invoke-unit` expression.

Each supplied *tagged-sig-spec* takes bindings from the surrounding context and turns them into imports for the invoked unit. The unit need not declare an import for every provided *tagged-sig-spec*, but one *tagged-sig-spec* must be provided for each declared import of the unit. For each variable identifier in each provided *tagged-sig-spec*, the value of the identifier's binding in the surrounding context is used for the corresponding import in the invoked unit.

```
(define-values/invoke-unit unit-expr
  (import tagged-sig-spec ...)
  (export tagged-sig-spec ...))
```

Like `invoke-unit`, but the values of the unit's exports are copied to new bindings.

The unit produced by *unit-expr* is linked and invoked as for `invoke-unit`. In addition, the `export` clause is treated as a kind of import into the local definition context. That is, for every binding that would be available in a unit that used the `export` clause's *tagged-sig-spec* as an import, a definition is generated for the context of the `define-values/invoke-unit` form.

7.3 Linking Units and Creating Compound Units

```
(compound-unit
  (import link-binding ...)
  (export tagged-link-id ...)
  (link linkage-decl ...))

  link-binding = (link-id : tagged-sig-id)

tagged-link-id = (tag id link-id)
               | link-id

  linkage-decl = ((link-binding ...) unit-expr tagged-link-id ...)
```

Links several units into one new compound unit without immediately invoking any of the

linked units. The *unit-expr*s in the `link` clause determine the units to be linked in creating the compound unit. The *unit-expr*s are evaluated when the `compound-unit` form is evaluated.

The `import` clause determines the imports of the compound unit. Outside the compound unit, these imports behave as for a plain unit; inside the compound unit, they are propagated to some of the linked units. The `export` clause determines the exports of the compound unit. Again, outside the compound unit, these exports are treated the same as for a plain unit; inside the compound unit, they are drawn from the exports of the linked units. Finally, the left-hand and right-hand parts of each declaration in the `link` clause specify how the compound unit's imports and exports are propagated to the linked units.

Individual elements of an imported or exported signature are not available within the compound unit. Instead, imports and exports are connected at the level of whole signatures. Each specific import or export (i.e., an instance of some signature, possibly tagged) is given a *link-id* name. Specifically, a *link-id* is bound by the `import` clause or the left-hand part of a declaration in the `link` clause. A bound *link-id* is referenced in the right-hand part of a declaration in the `link` clause or by the `export` clause.

The left-hand side of a `link` declaration gives names to each expected export of the unit produced by the corresponding *unit-expr*. The actual unit may export additional signatures, and it may export an extension of a specific signature instead of just the specified one. If the unit does not export one of the specified signatures (with the specified tag, if any), the `exn:fail:contract` exception is raised when the `compound-unit` form is evaluated.

The right-hand side of a `link` declaration specifies the imports to be supplied to the unit produced by the corresponding *unit-expr*. The actual unit may import fewer signatures, and it may import a signature that is extended by the specified one. If the unit imports a signature (with a particular tag) that is not included in the supplied imports, the `exn:fail:contract` exception is raised when the `compound-unit` form is evaluated. Each *link-id* supplied as an import must be bound either in the `import` clause or in some declaration within the `link` clause.

The order of declarations in the `link` clause determines the order of invocation of the linked units. When the compound unit is invoked, the unit produced by the first *unit-expr* is invoked first, then the second, and so on. If the order specified in the `link` clause is inconsistent with `init-depend` declarations of the actual units, then the `exn:fail:contract` exception is raised when the `compound-unit` form is evaluated.

7.4 Inferred Linking

```
(define-unit unit-id
  (import tagged-sig-spec ...)
  (export tagged-sig-spec ...)
  init-depends-decl
  unit-body-expr-or-defn
  ...)
```

Binds *unit-id* to both a unit and static information about the unit.

Evaluating a reference to a *unit-id* bound by `define-unit` produces a unit, just like evaluating an `id` bound by `(define id (unit ...))`. In addition, however, *unit-id* can be used in `compound-unit/infer`. See `unit` for information on *tagged-sig-spec*, *init-depends-decl*, and *unit-body-expr-or-defn*.

```
(compound-unit/infer
  (import tagged-infer-link-import ...)
  (export tagged-infer-link-export ...)
  (link infer-linkage-decl ...))

tagged-infer-link-import = tagged-sig-id
                         | (link-id : tagged-sig-id)

tagged-infer-link-export = (tag id infer-link-export)
                         | infer-link-export

       infer-link-export = link-id
                         | sig-id

      infer-linkage-decl = ((link-binding ...) unit-id
                                               tagged-link-id ...)
                         | unit-id
```

Like `compound-unit`. Syntactically, the difference between `compound-unit` and `compound-unit/infer` is that the `unit-expr` for a linked unit is replaced with a *unit-id*, where a *unit-id* is bound by `define-unit` (or one of the other unit-binding forms that we introduce later in this section). Furthermore, an import can name just a *sig-id* without locally binding a *link-id*, and an export can be based on a *sig-id* instead of a *link-id*, and a declaration in the `link` clause can be simply a *unit-id* with no specified exports or imports.

The `compound-unit/infer` form expands to `compound-unit` by adding `sig-ids` as needed to the `import` clause, by replacing *sig-id*s in the `export` clause by *link-id*s, and by completing the declarations of the `link` clause. This completion is based on static information associated with each *unit-id*. Links and exports can be inferred when all signatures exported by the linked units are distinct from each other and from all imported

signatures, and when all imported signatures are distinct. Two signatures are *distinct* only if they share no ancestor through `extends`.

The long form of a `link` declaration can be used to resolve ambiguity by giving names to some of a unit's exports and supplying specific bindings for some of a unit's imports. The long form need not name all of a unit's exports or supply all of a unit's imports if the remaining parts can be inferred.

Like `compound-unit`, the `compound-unit/infer` form produces a (compound) unit without statically binding information about the result unit's imports and exports. That is, `compound-unit/infer` consumes static information, but it does not generate it. Two additional forms, `define-compound-unit` and `define-compound-unit/infer`, generate static information (where the former does not consume static information).

```
(define-compound-unit id
  (import link-binding ...)
  (export tagged-link-id ...)
  (link linkage-decl ...))
```

Like `compound-unit`, but binds static information about the compound unit like `define-unit`.

```
(define-compound-unit/infer id
  (import link-binding ...)
  (export tagged-infer-link-export ...)
  (link infer-linkage-decl ...))
```

Like `compound-unit/infer`, but binds static information about the compound unit like `define-unit`.

```
(define-unit-binding unit-id
  unit-expr
  (import tagged-sig-spec ...+)
  (export tagged-sig-spec ...+)
  init-depends-decl)
```

Like `define-unit`, but the unit implementation is determined from an existing unit produced by *unit-expr*. The imports and exports of the unit produced by *unit-expr* must be consistent with the declared imports and exports, otherwise the `exn:fail:contract` exception is raised when the `define-unit-binding` form is evaluated.

```
(invoke-unit/infer unit-spec)

unit-spec = unit-id
          | (link link-unit-id ...)
```

Like `invoke-unit`, but uses static information associated with *`unit-id`* to infer which imports must be assembled from the current context. If given a link form containing multiple *`link-unit-id`*s, then the units are first linked via `define-compound-unit/infer`.

```
(define-values/invoke-unit/infer maybe-exports unit-spec)

maybe-exports =
              | (export tagged-sig-spec ...)

    unit-spec = unit-id
              | (link link-unit-id ...)
```

Like `define-values/invoke-unit`, but uses static information associated with *`unit-id`* to infer which imports must be assembled from the current context and which exports should be bound by the definition. If given a link form containing multiple *`link-unit-id`*s, then the units are first linked via `define-compound-unit/infer`.

7.5 Generating A Unit from Context

```
(unit-from-context tagged-sig-spec)
```

Creates a unit that implements an interface using bindings in the enclosing environment. The generated unit is essentially the same as

```
(unit
  (import)
  (export tagged-sig-spec)
  (define id expr) ...)
```

for each `id` that must be defined to satisfy the exports, and each corresponding `expr` produces the value of `id` in the environment of the `unit-from-context` expression. (The unit cannot be written as above, however, since each `id` definition within the unit shadows the binding outside the `unit` form.)

See `unit` for the grammar of *`tagged-sig-spec`*.

```
(define-unit-from-context id tagged-sig-spec)
```

Like `unit-from-context`, in that a unit is constructed from the enclosing environment, and like `define-unit`, in that *`id`* is bound to static information to be used later with inference.

7.6 Structural Matching

```
(unit/new-import-export
  (import tagged-sig-spec ...)
  (export tagged-sig-spec ...)
  init-depends-decl
  ((tagged-sig-spec ...) unit-expr tagged-sig-spec))
```

Similar to `unit`, except the body of the unit is determined by an existing unit produced by *unit-expr*. The result is a unit whose implementation is *unit-expr*, but whose imports, exports, and initialization dependencies are as in the `unit/new-import-export` form (instead of as in the unit produced by *unit-expr*).

The final clause of the `unit/new-import-export` form determines the connection between the old and new imports and exports. The connection is similar to the way that `compound-unit` propagates imports and exports; the difference is that the connection between `import` and the right-hand side of the link clause is based on the names of elements in signatures, rather than the names of the signatures. That is, a *tagged-sig-spec* on the right-hand side of the link clause need not appear as a *tagged-sig-spec* in the `import` clause, but each of the bindings implied by the linking *tagged-sig-spec* must be implied by some *tagged-sig-spec* in the `import` clause. Similarly, each of the bindings implied by an `export` *tagged-sig-spec* must be implied by some left-hand-side *tagged-sig-spec* in the linking clause.

```
(define-unit/new-import-export unit-id
  (import tagged-sig-spec ...)
  (export tagged-sig-spec ...)
  init-depends-decl
  ((tagged-sig-spec ...) unit-expr tagged-sig-spec))
```

Like `unit/new-import-export`, but binds static information to *unit-id* like `define-unit`.

```
(unit/s
  (import tagged-sig-spec ...)
  (export tagged-sig-spec ...)
  init-depends-decl
  unit-id)
```

Like `unit/new-import-export`, but the linking clause is inferred, so *unit-id* must have the appropriate static information.

```
(define-unit/s name-id
  (import tagged-sig-spec ...)
  (export tagged-sig-spec ...)
  init-depends-decl
  unit-id)
```

Like `unit/s`, but binds static information to *name-id* like `define-unit`.

7.7 Extending the Syntax of Signatures

```
(define-signature-form sig-form-id expr)
(define-signature-form (sig-form-id id) body ...+)
```

Binds *sig-form-id* for use within a `define-signature` form.

In the first form, the result of *expr* must be a transformer procedure. In the second form, *sig-form-id* is bound to a transformer procedure whose argument is *id* and whose body is the *body*s. The result of the transformer must be a list of syntax objects, which are substituted for a use of *sig-form-id* in a `define-signature` expansion. (The result is a list so that the transformer can produce multiple declarations; `define-signature` has no splicing `begin` form.)

```
(struct/ctc id ([field contract-expr] ...) struct-option ...)

        field = id
              | [id #:mutable]

struct-option = #:mutable
              | #:omit-constructor
              | #:omit-define-syntaxes
              | #:omit-define-values
```

For use with `define-signature`. The `struct/ctc` form works similarly to `struct`, but the constructor, predicate, field accessors, and field mutators are contracted appropriately.

7.8 Unit Utilities

```
(unit? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a unit, `#f` otherwise.

```
(provide-signature-elements sig-spec ...)
```

Expands to a `provide` of all identifiers implied by the *sig-spec*s. See `unit` for the grammar of *sig-spec*.

7.9 Unit Contracts

```
(unit/c (import sig-block ...) (export sig-block ...))

sig-block = (tagged-sig-id [id contract] ...)
          | tagged-sig-id
```

A *unit contract* wraps a unit and checks both its imported and exported identifiers to ensure that they match the appropriate contracts. This allows the programmer to add contract checks to a single unit value without adding contracts to the imported and exported signatures.

The unit value must import a subset of the import signatures and export a superset of the export signatures listed in the unit contract. Any identifier which is not listed for a given signature is left alone. Variables used in a given *contract* expression first refer to other variables in the same signature, and then to the context of the `unit/c` expression.

```
(define-unit/contract unit-id
  (import sig-spec-block ...)
  (export sig-spec-block ...)
  init-depends-decl
  unit-body-expr-or-defn
  ...)

sig-spec-block = (tagged-sig-spec [id contract] ...)
               | tagged-sig-spec
```

The `define-unit/contract` form defines a unit compatible with link inference whose imports and exports are contracted with a unit contract. The unit name is used for the positive blame of the contract.

7.10 Single-Unit Modules

When `racket/unit` is used as a language name with `#lang`, the module body is treated as a unit body. The body must match the following *module-body* grammar:

```
module-body = require-decl ...
              (import tagged-sig-expr ...)
              (export tagged-sig-expr ...)
              init-depends-decl
              unit-body-expr-or-defn
              ...

require-decl = (require require-spec ...)
             | (begin require-decl ...)
```

| *derived-require-form*

After any number of *require-decl*s, the content of the module is the same as a `unit` body.

The resulting unit is exported as *base*`@`, where *base* is derived from the enclosing module's name (i.e., its symbolic name, or its path without the directory and file suffix). If the module name ends in `-unit`, then *base* corresponds to the module name before `-unit`. Otherwise, the module name serves as *base*.

7.11 Single-Signature Modules

`#lang racket/signature` package: `base`

The `racket/signature` language treats a module body as a unit signature.

The body must match the following *module-body* grammar:

module-body = (require *require-spec* ...) ... *sig-spec* ...

See §7.1 "Creating Units" for the grammar of *sig-spec*. Unlike the body of a `racket/unit` module, a `require` in a `racket/signature` module must be a literal use of `require`.

The resulting signature is exported as *base*`^`, where *base* is derived from the enclosing module's name (i.e., its symbolic name, or its path without the directory and file suffix). If the module name ends in `-sig`, then *base* corresponds to the module name before `-sig`. Otherwise, the module name serves as *base*.

7.12 Transformer Helpers

`(require racket/unit-exptime)` package: `base`

The `racket/unit-exptime` library provides procedures that are intended for use by macro transformers. In particular, the library is typically imported using `for-syntax` into a module that defines macro with `define-syntax`.

```
(unit-static-signatures unit-identifier
                        err-syntax)
   (list/c (cons/c (or/c symbol? #f)
                   identifier?))
→  (list/c (cons/c (or/c symbol? #f)
                   identifier?))
 unit-identifier : identifier?
 err-syntax : syntax?
```

If *unit-identifier* is bound to static unit information via `define-unit` (or other such forms), the result is two values. The first value is for the unit's imports, and the second is for the unit's exports. Each result value is a list, where each list element pairs a symbol or `#f` with an identifier. The symbol or `#f` indicates the import's or export's tag (where `#f` indicates no tag), and the identifier indicates the binding of the corresponding signature.

If *unit-identifier* is not bound to static unit information, then the `exn:fail:syntax` exception is raised. In that case, the given *err-syntax* argument is used as the source of the error, where *unit-identifier* is used as the detail source location.

```
(signature-members sig-identifier
                                    (or/c identifier? #f)
                   err-syntax)   →  (listof identifier?)
                                    (listof identifier?)
                                    (listof identifier?)
  sig-identifier : identifier?
  err-syntax : syntax?
```

If *sig-identifier* is bound to static unit information via `define-signature` (or other such forms), the result is four values:

- an identifier or `#f` indicating the signature (of any) that is extended by the *sig-identifier* binding;
- a list of identifiers representing the variables supplied/required by the signature;
- a list of identifiers for variable definitions in the signature (i.e., variable bindings that are provided on import, but not defined by units that implement the signature); and
- a list of identifiers with syntax definitions in the signature.

If *sig-identifier* is not bound to a signature, then the `exn:fail:syntax` exception is raised. In that case, the given *err-syntax* argument is used as the source of the error, where *sig-identifier* is used as the detail source location.

8 Contracts

§7 “Contracts” in ***The Racket Guide*** introduces contracts.

The contract system guards one part of a program from another. Programmers specify the behavior of a module’s exports via `(provide (contract-out ....))`, and the contract system enforces those constraints.

`(require racket/contract)` package: `base`

The bindings documented in this section are provided by the `racket/contract` and `racket` libraries, but not `racket/base`.

Contracts come in two forms: those constructed by the various operations listed in this section of the manual, and various ordinary Racket values that double as contracts, including

- symbols, booleans, characters, keywords, and `null`, which are treated as contracts that recognize themselves, using `eq?`,
- strings and byte strings, which are treated as contracts that recognize themselves using `equal?`,
- numbers, which are treated as contracts that recognize themselves using `=`,
- regular expressions, which are treated as contracts that recognize byte strings and strings that match the regular expression, and
- predicates: any procedure of arity 1 is treated as a predicate. During contract checking, it is applied to the values that appear and should return `#f` to indicate that the contract failed, and anything else to indicate it passed.

Contract combinators are functions such as `->` and `listof` that take contracts and produce other contracts.

Contracts in Racket are subdivided into three different categories:

- *Flat contract*s can be fully checked immediately for a given value. These kinds of contracts are essentially predicate functions. Using `flat-contract-predicate`, you can extract the predicate from an arbitrary flat contract; some flat contracts can be applied like functions, in which case they accept a single argument and return `#t` or `#f` to indicate if the given value would be accepted by the contract. All of the flat contracts returned by functions in this library can be used directly as predicates, but ordinary Racket values that double as flat contracts (e.g., numbers or symbols) cannot.

  The function `flat-contract?` recognizes a flat contract.

- *Chaperone contracts* are not always immediately checkable, but are guaranteed to not change any properties of any values that they check. That is, they may wrap a value in such a way that it signals contract violations later, as the value is used (e.g., a

function contract checks the inputs and outputs to the function only when the function is called and returned), but any properties that the value had before being wrapped by the contract are preserved by the contract wrapper.

All flat contracts are also chaperone contracts (but not vice-versa).

- *Impersonator contracts* do not provide any guarantees about values they check. Impersonator contracts may hide properties of values, or even make them completely opaque (e.g, new-∀/c).

  All contracts are impersonator contracts.

For more about this hierarchy, see chaperones and a research paper on chaperones, impersonators, and how they can be used to implement contracts~[Strickland12].

8.1 Data-structure Contracts

```
(flat-named-contract name
                     flat-contract
                     [generator])    → flat-contract?
  name : any/c
  flat-contract : flat-contract?
  generator : (or/c #f (-> contract (-> int? any))) = #f
```

Produces a contract like *flat-contract*, but with the name *name*.

For example,

```
(define/contract i
  (flat-named-contract
   'odd-integer
   (lambda (x) (and (integer? x) (odd? x))))
  2)
```

The generator argument adds a generator for the flat-named-contract. See contract-generate for more information.

```
any/c : flat-contract?
```

A flat contract that accepts any value.

When using this contract as the result portion of a function contract, consider using any instead; using any leads to better memory performance, but it also allows multiple results.

```
none/c : flat-contract?
```

A flat contract that accepts no values.

```
(or/c contract ...) → contract?
  contract : contract?
```

Takes any number of contracts and returns a contract that accepts any value that any one of the contracts accepts individually.

The `or/c` result tests any value by applying the contracts in order, from left to right, with the exception that it always moves the non-flat contracts (if any) to the end, checking them last. Thus, a contract such as `(or/c (not/c real?) positive?)` is guaranteed to only invoke the `positive?` predicate on real numbers.

If all of the arguments are procedures or flat contracts, the result is a flat contract. If only one of the arguments is a higher-order contract, the result is a contract that just checks the flat contracts and, if they don't pass, applies the higher-order contract.

If there are multiple higher-order contracts, `or/c` uses `contract-first-order-passes?` to distinguish between them. More precisely, when an `or/c` is checked, it first checks all of the flat contracts. If none of them pass, it calls `contract-first-order-passes?` with each of the higher-order contracts. If only one returns true, `or/c` uses that contract. If none of them return true, it signals a contract violation. If more than one returns true, it also signals a contract violation. For example, this contract

```
(or/c (-> number? number?)
      (-> string? string? string?))
```

does not accept a function like this one: `(lambda args ...)` since it cannot tell which of the two arrow contracts should be used with the function.

If all of its arguments are `list-contract?`s, then `or/c` returns a `list-contract?`.

```
(and/c contract ...) → contract?
  contract : contract?
```

Takes any number of contracts and returns a contract that accepts any value that satisfies all of the contracts simultaneously.

If all of the arguments are procedures or flat contracts, the result is a flat contract.

The contract produced by `and/c` tests any value by applying the contracts in order, from left to right.

```
(not/c flat-contract) → flat-contract?
  flat-contract : flat-contract?
```

Accepts a flat contracts or a predicate and returns a flat contract that checks the inverse of the argument.

```
(=/c z) → flat-contract?
  z : real?
```

Returns a flat contract that requires the input to be a number and = to z.

```
(</c n) → flat-contract?
  n : real?
```

Returns a flat contract that requires the input to be a number and < than n.

```
(>/c n) → flat-contract?
  n : real?
```

Like </c, but for >.

```
(<=/c n) → flat-contract?
  n : real?
```

Like </c, but for <=.

```
(>=/c n) → flat-contract?
  n : real?
```

Like </c, but for >=.

```
(between/c n m) → flat-contract?
  n : real?
  m : real?
```

Returns a flat contract that requires the input to be a real number between n and m or equal to one of them.

```
(real-in n m) → flat-contract?
  n : real?
  m : real?
```

An alias for between/c.

```
(integer-in j k) → flat-contract?
  j : exact-integer?
  k : exact-integer?
```

Returns a flat contract that requires the input to be an exact integer between *j* and *k*, inclusive.

```
natural-number/c : flat-contract?
```

A flat contract that requires the input to be an exact non-negative integer.

```
(string-len/c len) → flat-contract?
  len : real?
```

Returns a flat contract that recognizes strings that have fewer than *len* characters.

```
false/c : flat-contract?
```

An alias `#f` for backwards compatibility.

```
printable/c : flat-contract?
```

A flat contract that recognizes values that can be written out and read back in with `write` and `read`.

```
(one-of/c v ...+) → flat-contract?
  v : any/c
```

Accepts any number of atomic values and returns a flat contract that recognizes those values, using `eqv?` as the comparison predicate. For the purposes of `one-of/c`, atomic values are defined to be: characters, symbols, booleans, `null`, keywords, numbers, `#<void>`, and `#<undefined>`.

This is a backwards compatibility contract constructor. If neither `#<void>` nor `#<undefined>` are arguments, it simply passes its arguments to `or/c`.

```
(symbols sym ...+) → flat-contract?
  sym : symbol?
```

Accepts any number of symbols and returns a flat contract that recognizes those symbols.

This is a backwards compatibility constructor; it merely passes its arguments to `or/c`.

```
(vectorof  c
           [#:immutable immutable
            #:flat? flat?])        → contract?
  c : contract?
  immutable : (or/c #t #f 'dont-care) = 'dont-care
  flat? : boolean? = #f
```

Returns a contract that recognizes vectors. The elements of the vector must match *c*.

If the *flat?* argument is #t, then the resulting contract is a flat contract, and the *c* argument must also be a flat contract. Such flat contracts will be unsound if applied to mutable vectors, as they will not check future operations on the vector.

If the *immutable* argument is #t and the *c* argument is a flat contract, the result will be a flat contract. If the *c* argument is a chaperone contract, then the result will be a chaperone contract.

When a higher-order vectorof contract is applied to a vector, the result is not eq? to the input. The result will be a copy for immutable vectors and a chaperone or impersonator of the input for mutable vectors.

```
(vector-immutableof c) → contract?
  c : contract?
```

Returns the same contract as (vectorof *c* #:immutable #t). This form exists for backwards compatibility.

```
(vector/c  c
           ...
           [#:immutable immutable
            #:flat? flat?])        → contract?
  c : contract?
  immutable : (or/c #t #f 'dont-care) = 'dont-care
  flat? : boolean? = #f
```

Returns a contract that recognizes vectors whose lengths match the number of contracts given. Each element of the vector must match its corresponding contract.

If the *flat?* argument is #t, then the resulting contract is a flat contract, and the *c* arguments must also be flat contracts. Such flat contracts will be unsound if applied to mutable vectors, as they will not check future operations on the vector.

If the *immutable* argument is #t and the *c* arguments are flat contracts, the result will be a flat contract. If the *c* arguments are chaperone contracts, then the result will be a chaperone contract.

When a higher-order vector/c contract is applied to a vector, the result is not eq? to the input. The result will be a copy for immutable vectors and a chaperone or impersonator of the input for mutable vectors.

```
(vector-immutable/c c ...) → contract?
  c : contract?
```

Returns the same contract as (vector/c c ... #:immutable #t). This form exists for reasons of backwards compatibility.

```
(box/c  c
       [#:immutable immutable
        #:flat? flat?])        → contract?
  c : contract?
  immutable : (or/c #t #f 'dont-care) = 'dont-care
  flat? : boolean? = #f
```

Returns a contract that recognizes boxes. The content of the box must match c.

If the flat? argument is #t, then the resulting contract is a flat contract, and the c argument must also be a flat contract. Such flat contracts will be unsound if applied to mutable boxes, as they will not check future operations on the box.

If the immutable argument is #t and the c argument is a flat contract, the result will be a flat contract. If the c argument is a chaperone contract, then the result will be a chaperone contract.

When a higher-order box/c contract is applied to a box, the result is not eq? to the input. The result will be a copy for immutable boxes and either a chaperone or impersonator of the input for mutable boxes.

```
(box-immutable/c c) → contract?
  c : contract?
```

Returns the same contract as (box/c c #:immutable #t). This form exists for reasons of backwards compatibility.

```
(listof c) → list-contract?
  c : contract?
```

Returns a contract that recognizes a list whose every element matches the contract c. Beware that when this contract is applied to a value, the result is not necessarily eq? to the input.

```
(non-empty-listof c) → list-contract?
  c : contract?
```

Returns a contract that recognizes non-empty lists whose elements match the contract *c*. Beware that when this contract is applied to a value, the result is not necessarily eq? to the input.

```
(cons/c car-c cdr-c) → contract?
  car-c : contract?
  cdr-c : contract?
```

Produces a contract that recognizes pairs whose first and second elements match *car-c* and *cdr-c*, respectively. Beware that when this contract is applied to a value, the result is not necessarily eq? to the input.

If the *cdr-c* contract is a list-contract?, then cons/c returns a list-contract?.

Changed in version 6.0.1.13 of package base: Added the list-contract? propagating behavior.

```
(list/c c ...) → list-contract?
  c : contract?
```

Produces a contract for a list. The number of elements in the list must match the number of arguments supplied to list/c, and each element of the list must match the corresponding contract. Beware that when this contract is applied to a value, the result is not necessarily eq? to the input.

```
(syntax/c c) → flat-contract?
  c : flat-contract?
```

Produces a flat contract that recognizes syntax objects whose syntax-e content matches *c*.

```
(struct/c struct-id contract-expr ...)
```

Produces a contract that recognizes instances of the structure type named by *struct-id*, and whose field values match the contracts produced by the *contract-expr*s.

Contracts for immutable fields must be either flat or chaperone contracts. Contracts for mutable fields may be impersonator contracts. If all fields are immutable and the *contract-expr*s evaluate to flat contracts, a flat contract is produced. If all the *contract-expr*s are chaperone contracts, a chaperone contract is produced. Otherwise, an impersonator contract is produced.

```
(struct/dc struct-id field-spec ... maybe-inv)
```

```
        field-spec = [field-name maybe-lazy contract-expr]
                   | [field-name (dep-field-name ...)
                                 maybe-lazy
                                 maybe-contract-type
                                 maybe-dep-state
                                 contract-expr]

        field-name = field-id
                   | (#:selector selector-id)
                   | (field-id #:parent struct-id)

        maybe-lazy =
                   | #:lazy

maybe-contract-type =
                   | #:flat
                   | #:chaperone
                   | #:impersonator

    maybe-dep-state =
                   | #:depends-on-state

         maybe-inv =
                   | #:inv (dep-field-name ...) invariant-expr
```

Produces a contract that recognizes instances of the structure type named by *struct-id*, and whose field values match the contracts produced by the *field-spec*s.

If the *field-spec* lists the names of other fields, then the contract depends on values in those fields, and the *contract-expr* expression is evaluated each time a selector is applied, building a new contract for the fields based on the values of the *dep-field-name* fields (the *dep-field-name* syntax is the same as the *field-name* syntax). If the field is a dependent field and no `contract-type` annotation appears, then it is assumed that the contract is a chaperone, but not always a flat contract (and thus the entire `struct/dc` contract is not a flat contract). If this is not the case, and the contract is always flat then the field must be annotated with the `#:flat`, or the field must be annotated with `#:impersonator` (in which case, it must be a mutable field).

A *field-name* is either an identifier naming a field in the first case, an identifier naming a selector in the second case indicated by the `#:selector` keyword, or a field id for a struct that is a parent of *struct-id*, indicated by the `#:parent` keyword.

If the `#:lazy` keyword appears, then the contract on the field is check lazily (only when a selector is applied); `#:lazy` contracts cannot be put on mutable fields.

If a dependent contract depends on some mutable state, then use the `#:depends-on-state`

keyword argument (if a field's dependent contract depends on a mutable field, this keyword is automatically inferred). The presence of this keyword means that the contract expression is evaluated each time the corresponding field is accessed (or mutated, if it is a mutable field). Otherwise, the contract expression for a dependent field contract is evaluated when the contract is applied to a value.

If the `#:inv` clause appears, then the invariant expression is evaluated (and must return a non-`#f` value) when the contract is applied to a struct.

Contracts for immutable fields must be either flat or chaperone contracts. Contracts for mutable fields may be impersonator contracts. If all fields are immutable and the *contract-expr*s evaluate to flat contracts, a flat contract is produced. If all the *contract-expr*s are chaperone contracts, a chaperone contract is produced. Otherwise, an impersonator contract is produced.

As an example, the function `bst/c` below returns a contract for binary search trees whose values are all between `lo` and `hi`. The lazy annotations ensure that this contract does not change the running time of operations that do not inspect the entire tree.

```
(struct bt (val left right))
(define (bst/c lo hi)
  (or/c #f
        (struct/dc bt
                   [val (between/c lo hi)]
                   [left (val) #:lazy (bst lo val)]
                   [right (val) #:lazy (bst val hi)])))
```

Changed in version 6.0.1.6 of package `base`: Added `#:inv`.

```
(parameter/c in [out]) → contract?
  in : contract?
  out : contract? = in
```

Produces a contract on parameters whose values must match *out*. When the value in the contracted parameter is set, it must match *in*.

Examples:

```
> (define/contract current-snack
    (parameter/c string?)
    (make-parameter "potato-chip"))

> (define baked/c
    (flat-named-contract 'baked/c (λ (s) (regexp-
match #rx"baked" s))))
```

```
> (define/contract current-dinner
    (parameter/c string? baked/c)
    (make-parameter "turkey" (λ (s) (string-append "roasted
" s))))

> (current-snack 'not-a-snack)
current-snack: contract violation
  expected: string?
  given: 'not-a-snack
  in: the parameter of
      (parameter/c string?)
  contract from: (definition current-snack)
  blaming: top-level
   (assuming the contract is correct)
  at: eval:2.0
> (parameterize ([current-dinner "tofurkey"])
    (current-dinner))
current-dinner: broke its contract
  promised: baked/c
  produced: "roasted tofurkey"
  in: the parameter of
      (parameter/c string? baked/c)
  contract from: (definition current-dinner)
  blaming: (definition current-dinner)
   (assuming the contract is correct)
  at: eval:4.0
```

```
(procedure-arity-includes/c n) → flat-contract?
  n : exact-nonnegative-integer?
```

Produces a contract for procedures that accept *n* argument (i.e,. the `procedure?` contract is implied).

```
(hash/c  key
         val
        [#:immutable immutable
         #:flat? flat?])        → contract?
  key : chaperone-contract?
  val : contract?
  immutable : (or/c #t #f 'dont-care) = 'dont-care
  flat? : boolean? = #f
```

Produces a contract that recognizes `hash` tables with keys and values as specified by the *key* and *val* arguments.

Examples:

```
> (define/contract good-hash
    (hash/c integer? boolean?)
    (hash 1 #t
          2 #f
          3 #t))

> (define/contract bad-hash
    (hash/c integer? boolean?)
    (hash 1 "elephant"
          2 "monkey"
          3 "manatee"))
bad-hash: broke its contract
  promised: boolean?
  produced: "elephant"
  in: the values of
      (hash/c integer? boolean?)
  contract from: (definition bad-hash)
  blaming: (definition bad-hash)
   (assuming the contract is correct)
  at: eval:3.0
```

There are a number of technicalities that control how `hash/c` contracts behave.

- If the *`flat?`* argument is `#t`, then the resulting contract is a flat contract, and the *`key`* and `val` arguments must also be flat contracts.

  Examples:

  ```
  > (flat-contract? (hash/c integer? boolean?))
  #f
  > (flat-contract? (hash/c integer? boolean? #:flat? #t))
  #t
  > (hash/c integer? (-> integer? integer?) #:flat? #t)
  hash/c: contract violation
    expected: flat-contract?
    given: #<chaperone-contract: (-> integer? integer?)>
  ```

  Such flat contracts will be unsound if applied to mutable hash tables, as they will not check future mutations to the hash table.

  Examples:

  ```
  > (define original-h (make-hasheq))

  > (define/contract ctc-h
      (hash/c integer? boolean? #:flat? #t)
      original-h)
  ```

```
> (hash-set! original-h 1 "not a boolean")

> (hash-ref ctc-h 1)
"not a boolean"
```

- If the *immutable* argument is `#t` and the *key* and *val* arguments are `flat-contract?`s, the result will be a `flat-contract?`.

  Example:

  ```
  > (flat-contract? (hash/c integer? boolean? #:immutable #t))
  #t
  ```

  If either the domain or the range is a `chaperone-contract?`, then the result will be a `chaperone-contract?`.

  Examples:

  ```
  > (flat-contract? (hash/c (-> integer? integer?) boolean?
                            #:immutable #t))
  #f
  > (chaperone-contract? (hash/c (-> integer? integer?) boolean?
                                 #:immutable #t))
  #t
  ```

- If the *key* argument is a `chaperone-contract?` but not a `flat-contract?`, then the resulting contract can be applied only to `equal?`-based hash tables.

  Example:

  ```
  > (define/contract h
      (hash/c (-> integer? integer?) any/c)
      (make-hasheq))
  h: broke its contract
    promised equal?-based hash table due to higher-order
  domain contract
    produced: '#hasheq()
    in: (hash/c (-> integer? integer?) any/c)
    contract from: (definition h)
    blaming: (definition h)
     (assuming the contract is correct)
    at: eval:2.0
  ```

  Also, when such a `hash/c` contract is applied to a hash table, the result is not `eq?` to the input. The result of applying the contract will be a copy for immutable hash tables, and either a chaperone or impersonator of the original hash table for mutable hash tables.

`(channel/c val)` → `contract?`
  val : `contract?`

Produces a contract that recognizes channels that communicate values as specified by the *val* argument.

If the *val* argument is a chaperone contract, then the resulting contract is a chaperone contract. Otherwise, the resulting contract is an impersonator contract. When a channel contract is applied to a channel, the resulting channel is not `eq?` to the input.

Examples:

```
> (define/contract chan
    (channel/c string?)
    (make-channel))

> (thread (λ () (channel-get chan)))
#<thread>
> (channel-put chan 'not-a-string)
chan: contract violation
  expected: string?
  given: 'not-a-string
  in: (channel/c string?)
  contract from: (definition chan)
  blaming: top-level
   (assuming the contract is correct)
  at: eval:2.0
```

```
(prompt-tag/c contract ... maybe-call/cc)

maybe-call/cc =
              | #:call/cc contract
              | #:call/cc (values contract ...)

  contract : contract?
```

Takes any number of contracts and returns a contract that recognizes continuation prompt tags and will check any aborts or prompt handlers that use the contracted prompt tag.

Each *contract* will check the corresponding value passed to an `abort-current-continuation` and handled by the handler of a call to `call-with-continuation-prompt`.

If all of the *contract*s are chaperone contracts, the resulting contract will also be a chaperone contract. Otherwise, the contract is an impersonator contract.

If *maybe-call/cc* is provided, then the provided contracts are used to check the return values from a continuation captured with `call-with-current-continuation`.

Examples:

```
> (define/contract tag
    (prompt-tag/c (-> number? string?))
    (make-continuation-prompt-tag))

> (call-with-continuation-prompt
    (lambda ()
      (number->string
        (call-with-composable-continuation
          (lambda (k)
            (abort-current-continuation tag k)))))
    tag
    (lambda (k) (k "not a number")))
tag: contract violation
  expected: number?
  given: "not a number"
  in: the 1st argument of
      (prompt-tag/c
       (-> number? string?)
       #:call/cc)
  contract from: (definition tag)
  blaming: top-level
   (assuming the contract is correct)
  at: eval:2.0
```

```
(continuation-mark-key/c contract) → contract?
  contract : contract?
```

Takes a single contract and returns a contract that recognizes continuation marks and will check any mappings of marks to values or any accesses of the mark value.

If the argument *contract* is a chaperone contract, the resulting contract will also be a chaperone contract. Otherwise, the contract is an impersonator contract.

Examples:

```
> (define/contract mark-key
    (continuation-mark-key/c (-> symbol? (listof symbol?)))
    (make-continuation-mark-key))

> (with-continuation-mark
    mark-key
    (lambda (s) (append s '(truffle fudge ganache)))
    (let ([mark-value (continuation-mark-set-first
                       (current-continuation-marks) mark-key)])
      (mark-value "chocolate-bar")))
mark-key: contract violation
  expected: symbol?
```

```
  given: "chocolate-bar"
  in: the 1st argument of
        (continuation-mark-key/c
         (-> symbol? (listof symbol?)))
  contract from: (definition mark-key)
  blaming: top-level
   (assuming the contract is correct)
  at: eval:2.0
```

```
(evt/c contract ...) → chaperone-contract?
  contract : chaperone-contract?
```

Returns a contract that recognizes synchronizable events whose synchronization results are checked by the given *contract*s.

The resulting contract is always a chaperone contract and its arguments must all be chaperone contracts.

Examples:

```
> (define/contract my-evt
    (evt/c evt?)
    always-evt)

> (define/contract failing-evt
    (evt/c number? number?)
    (alarm-evt (+ (current-inexact-milliseconds) 50)))

> (sync my-evt)
#<always-evt>
> (sync failing-evt)
failing-evt: broke its contract
  promised: event that produces 2 values
  produced: event that produces 1 values
  in: (evt/c number? number?)
  contract from: (definition failing-evt)
  blaming: (definition failing-evt)
   (assuming the contract is correct)
  at: eval:3.0
```

```
(flat-rec-contract id flat-contract-expr ...)
```

Constructs a recursive flat contract. A *flat-contract-expr* can refer to *id* to refer recursively to the generated contract.

For example, the contract

```
(flat-rec-contract sexp
  (cons/c sexp sexp)
  number?
  symbol?)
```

is a flat contract that checks for (a limited form of) S-expressions. It says that a `sexp` is either two `sexp`s combined with `cons`, or a number, or a symbol.

Note that if the contract is applied to a circular value, contract checking will not terminate.

```
(flat-murec-contract ([id flat-contract-expr ...] ...) body ...+)
```

A generalization of `flat-rec-contract` for defining several mutually recursive flat contracts simultaneously. Each *id* is visible in the entire `flat-murec-contract` form, and the result of the final *body* is the result of the entire form.

```
any
```

Represents a contract that is always satisfied. In particular, it can accept multiple values. It can only be used in a result position of contracts like `->`. Using `any` elsewhere is a syntax error.

```
(promise/c c) → contract?
  c : contract?
```

Constructs a contract on a promise. The contract does not force the promise, but when the promise is forced, the contract checks that the result value meets the contract *c*.

```
(flat-contract predicate) → flat-contract?
  predicate : (-> any/c any/c)
```

Constructs a flat contract from *predicate*. A value satisfies the contract if the predicate returns a true value.

This function is a holdover from before flat contracts could be used directly as predicates. It exists today for backwards compatibilty.

```
(flat-contract-predicate v) → (-> any/c any/c)
  v : flat-contract?
```

Extracts the predicate from a flat contract.

This function is a holdover from before flat contracts could be used directly as predicates. It exists today for backwards compatibility.

8.2 Function Contracts

A *function contract* wraps a procedure to delay checks for its arguments and results. There are three primary function contract combinators that have increasing amounts of expressiveness and increasing additional overheads. The first `->` is the cheapest. It generates wrapper functions that can call the original function directly. Contracts built with `->*` require packaging up arguments as lists in the wrapper function and then using either `keyword-apply` or `apply`. Finally, `->i` is the most expensive (along with `->d`), because it requires delaying the evaluation of the contract expressions for the domain and range until the function itself is called or returns.

The `case->` contract is a specialized contract, designed to match `case-lambda` and `unconstrained-domain->` allows range checking without requiring that the domain have any particular shape (see below for an example use).

```
(-> dom ... range)

  dom = dom-expr
      | keyword dom-expr

range = range-expr
      | (values range-expr ...)
      | any
```

Produces a contract for a function that accepts a fixed number of arguments and returns either a fixed number of results or completely unspecified results (the latter when `any` is specified).

Each *dom-expr* is a contract on an argument to a function, and each *range-expr* is a contract on a result of the function.

Using a `->` between two whitespace-delimited `.`s is the same as putting the `->` right after the enclosing opening parenthesis. See §2.4.3 “Lists and Racket Syntax” or §1.3.6 “Reading Pairs and Lists” for more information.

For example,

```
(integer? boolean? . -> . integer?)
```

produces a contract on functions of two arguments. The first argument must be an integer, and the second argument must be a boolean. The function must produce an integer.

A domain specification may include a keyword. If so, the function must accept corresponding (mandatory) keyword arguments, and the values for the keyword arguments must match the corresponding contracts. For example:

```
(integer? #:x boolean? . -> . integer?)
```

is a contract on a function that accepts a by-position argument that is an integer and a `#:x` argument that is a boolean.

If any is used as the last sub-form for ->, no contract checking is performed on the result of the function, and thus any number of values is legal (even different numbers on different invocations of the function).

If (values *range-expr* ...) is used as the last sub-form of ->, the function must produce a result for each contract, and each value must match its respective contract.

```
(->* (mandatory-dom ...) optional-doms rest pre range post)

mandatory-dom = dom-expr
              | keyword dom-expr

optional-doms =
              | (optional-dom ...)

 optional-dom = dom-expr
              | keyword dom-expr

         rest =
              | #:rest rest-expr

          pre =
              | #:pre pre-cond-expr

        range = range-expr
              | (values range-expr ...)
              | any

         post =
              | #:post post-cond-expr
```

The ->* contract combinator produces contracts for functions that accept optional arguments (either keyword or positional) and/or arbitrarily many arguments. The first clause of a ->* contract describes the mandatory arguments, and is similar to the argument description of a -> contract. The second clause describes the optional arguments. The range of description can either be any or a sequence of contracts, indicating that the function must return multiple values.

If present, the *rest-expr* contract governs the arguments in the rest parameter. Note that the *rest-expr* contract governs only the arguments in the rest parameter, not those in mandatory arguments. For example, this contract:

```
(->* () #:rest (cons/c integer? (listof integer?)) any)
```

does not match the function

```
(λ (x . rest) x)
```

because the contract insists that the function accept zero arguments (because there are no mandatory arguments listed in the contract). The `->i` contract does not know that the contract on the rest argument is going to end up disallowing empty argument lists.

The *pre-cond-expr* and *post-cond-expr* expressions are checked as the function is called and returns, respectively, and allow checking of the environment without an explicit connection to an argument (or a result).

As an example, the contract

```
(->* () (boolean? #:x integer?) #:rest (listof symbol?) symbol?)
```

matches functions that optionally accept a boolean, an integer keyword argument `#:x` and arbitrarily more symbols, and that return a symbol.

```
(->i (mandatory-dependent-dom ...)
     dependent-rest
     pre-condition
     dependent-range
     post-condition)
(->i (mandatory-dependent-dom ...)
     (optional-dependent-dom ...)
     dependent-rest
     pre-condition
     dependent-range
     post-condition)
```

```
mandatory-dependent-dom = id+ctc
                        | keyword id+ctc

 optional-dependent-dom = id+ctc
                        | keyword id+ctc

         dependent-rest =
                        | #:rest id+ctc

          pre-condition =
                        | #:pre (id ...)
                          boolean-expr pre-condition
                        | #:pre/name (id ...)
                          string boolean-expr pre-condition

        dependent-range = any
                        | id+ctc
                        | un+ctc
                        | (values id+ctc ...)
                        | (values un+ctc ...)

         post-condition =
                        | #:post (id ...)
                          boolean-expr post-condition
                        | #:post/name (id ...)
                          string boolean-expr post-condition

                id+ctc = [id contract-expr]
                       | [id (id ...) contract-expr]

                un+ctc = [_ contract-expr]
                       | [_ (id ...) contract-expr]
```

The `->i` contract combinator differs from the `->*` combinator in that the support pre- and post-condition clauses and in that each argument and result is named. These names can then be used in the subcontracts and in the pre-/post-condition clauses. In short, contracts now express dependencies among arguments and results.

The first sub-form of a `->i` contract covers the mandatory and the second sub-form covers the optional arguments. Following that is an optional rest-args contract, and an optional pre-condition. The pre-condition is introduced with the `#:pre` keyword followed by the list of names on which it depends. If the `#:pre/name` keyword is used, the string supplied is used as part of the error message; similarly with `#:post/name`.

The *dependent-range* non-terminal specifies the possible result contracts. If it is `any`, then any value is allowed. Otherwise, the result contract pairs a name and a contract or a multiple

values return with names and contracts. In the last two cases, the range contract may be optionally followed by a post-condition; the post-condition expression is not allowed if the range contract is `any`. Like the pre-condition, the post-condition must specify the variables on which it depends.

Consider this sample contract:

```
(->i ([x number?]
      [y (x) (>=/c x)])
     [result (x y) (and/c number? (>=/c (+ x y)))])
```

It specifies a function of two arguments, both numbers. The contract on the second argument (`y`) demands that it is greater than the first argument. The result contract promises a number that is greater than the sum of the two arguments. While the dependency specification for `y` signals that the argument contract depends on the value of the first argument, the dependency sequence for `result` indicates that the contract depends on both argument values. Since the contract for `x` does not depend on anything else, it does not come with any dependency sequence, not even `()`.

In general, an empty sequence is (nearly) equivalent to not adding a sequence at all except that the former is more expensive than the latter.

This example is like the previous one, except the `x` and `y` arguments are now optional keyword arguments, instead of mandatory, by-position arguments:

```
(->i ()
     (#:x [x number?]
      #:y [y (x) (>=/c x)])
     [result (x y) (and/c number? (>=/c (+ x y)))])
```

The contract expressions are not always evaluated in order. First, if there is no dependency for a given contract expression, the contract expression is evaluated at the time that the `->i` expression is evaluated rather than the time when the function is called or returns. These dependency-free contract expressions are evaluated in the order in which they are listed. Second, the dependent contract sub-expressions are evaluated when the contracted function is called or returns in some order that satisfies the dependencies. That is, if a contract for an argument depends on the value of some other contract, the former is evaluated first (so that the argument, with its contract checked, is available for the other). When there is no dependency between two arguments (or the result and an argument), then the contract that appears earlier in the source text is evaluated first.

Finally, if all of the identifier positions of the range contract are `_` s (underscores), then the range contract expressions are evaluated when the function is called and the underscore is not bound in the range, after the argument contracts are evaluated and checked. Otherwise, the range expressions are evaluated when the function returns.

If there are optional arguments that are not supplied, then the corresponding variables will be bound to a special value called `the-unsupplied-arg` value.

```
(->d (mandatory-dependent-dom ...)
     dependent-rest
     pre-condition
     dependent-range
     post-condition)
(->d (mandatory-dependent-dom ...)
     (optional-dependent-dom ...)
     dependent-rest
     pre-condition
     dependent-range
     post-condition)

mandatory-dependent-dom = [id dom-expr]
                        | keyword [id dom-expr]

 optional-dependent-dom = [id dom-expr]
                        | keyword [id dom-expr]

         dependent-rest =
                        | #:rest id rest-expr

          pre-condition =
                        | #:pre boolean-expr
                        | #:pre-cond boolean-expr

        dependent-range = any
                        | [_ range-expr]
                        | (values [_ range-expr] ...)
                        | [id range-expr]
                        | (values [id range-expr] ...)

         post-condition =
                        | #:post-cond boolean-expr
```

This contract is here for backwards compatibility; any new code should use `->i` instead.

This contract is similar to `->i`, but is "lax", meaning that it does not enforce contracts internally. For example, using this contract

```
(->d ([f (-> integer? integer?)])
     #:pre
     (zero? (f #f))
     any)
```

will allow `f` to be called with `#f`, trigger whatever bad behavior the author of `f` was trying to prohibit by insisting that `f`'s contract accept only integers.

The `#:pre-cond` and `#:post-cond` keywords are aliases for `#:pre` and `#:post` and are provided for backwards compatibility.

```
(case-> (-> dom-expr ... rest range) ...)

 rest =
      | #:rest rest-expr

range = range-expr
      | (values range-expr ...)
      | any
```

This contract form is designed to match `case-lambda`. Each argument to `case->` is a contract that governs a clause in the `case-lambda`. If the `#:rest` keyword is present, the corresponding clause must accept an arbitrary number of arguments. The *range* specification is just like that for `->` and `->*`.

For example, this contract matches a function with two cases, one that accepts an integer, returning void, and one that accepts no arguments and returns an integer.

```
(case-> (-> integer? void?)
        (-> integer?))
```

Such a contract could be used to guard a function that controls access to a single shared integer.

```
(unconstrained-domain-> range-expr ...)
```

Constructs a contract that accepts a function, but makes no constraint on the function's domain. The *range-expr*s determine the number of results and the contract for each result.

Generally, this contract must be combined with another contract to ensure that the domain is actually known to be able to safely call the function itself.

For example, the contract

```
(provide
 (contract-out
  [f (->d ([size natural-number/c]
           [proc (and/c (unconstrained-domain-> number?)
                        (lambda (p)
                          (procedure-arity-includes? p size)))])
          ()
          [_ number?])]))
```

says that the function `f` accepts a natural number and a function. The domain of the function that `f` accepts must include a case for `size` arguments, meaning that `f` can safely supply `size` arguments to its input.

For example, the following is a definition of `f` that cannot be blamed using the above contract:

```
(define (f i g)
  (apply g (build-list i add1)))
```

`predicate/c : contract?`

Use this contract to indicate that some function is a predicate. It is semantically equivalent to `(-> any/c boolean?)`.

This contract also includes an optimization so that functions returning `#t` from `struct-predicate-procedure?` are just returned directly, without being wrapped. This contract is used by `provide/contract`'s `struct` subform so that struct predicates end up not being wrapped.

`the-unsupplied-arg : unsupplied-arg?`

Used by `->i` (and `->d`) to bind optional arguments that are not supplied by a call site.

```
(unsupplied-arg? v) → boolean?
  v : any/c
```

A predicate to determine whether `v` is `the-unsupplied-arg`.

8.3 Parametric Contracts

`(require racket/contract/parametric)` package: `base`

The most convenient way to use parametric contract is to use `contract-out`'s `#:exists` keyword. The `racket/contract/parametric` provides a few more, general-purpose parametric contracts.

`(parametric->/c (x ...) c)`

Creates a contract for parametric polymorphic functions. Each function is protected by `c`, where each `x` is bound in `c` and refers to a polymorphic type that is instantiated each time the function is applied.

At each application of a function, the `parametric->/c` contract constructs a new opaque wrapper for each `x`; values flowing into the polymorphic function (i.e. values protected by some `x` in negative position with respect to `parametric->/c`) are wrapped in the corresponding opaque wrapper. Values flowing out of the polymorphic function (i.e. values protected by some `x` in positive position with respect to `parametric->/c`) are checked for the appropriate wrapper. If they have it, they are unwrapped; if they do not, a contract violation is signaled.

Examples:

```
> (define/contract (check x y)
    (parametric->/c [X] (boolean? X . -> . X))
    (if (or (not x) (equal? y 'surprise))
        'invalid
        y))

> (check #t 'ok)
'ok
> (check #f 'ignored)
check: broke its contract
  promised: X
  produced: 'invalid
  in: the range of
      (parametric->/c (X) (-> boolean? X X))
  contract from: (function check)
  blaming: (function check)
   (assuming the contract is correct)
  at: eval:2.0
> (check #t 'surprise)
'surprise
```

```
(new-∀/c [name]) → contract?
  name : (or/c symbol? #f) = #f
```

Constructs a new universal contract.

Universal contracts accept all values when in negative positions (e.g., function inputs) and wrap them in an opaque struct, hiding the precise value. In positive positions (e.g. function returns), a universal contract accepts only values that were previously accepted in negative positions (by checking for the wrappers).

The name is used to identify the contract in error messages and defaults to a name based on the lexical context of `new-∀/c`.

For example, this contract:

```
(let ([a (new-∀/c 'a)])
  (-> a a))
```

describes the identity function (or a non-terminating function). That is, the first use of the a appears in a negative position and thus inputs to that function are wrapped with an opaque struct. Then, when the function returns, it is checked to determine whether the result is wrapped, since the second a appears in a positive position.

The new-∀/c construct constructor is dual to new-∃/c.

```
(new-∃/c [name]) → contract?
  name : (or/c symbol? #f) = #f
```

Constructs a new existential contract.

Existential contracts accept all values when in positive positions (e.g., function returns) and wrap them in an opaque struct, hiding the precise value. In negative positions (e.g. function inputs), they accepts only values that were previously accepted in positive positions (by checking for the wrappers).

The name is used to identify the contract in error messages and defaults to a name based on the lexical context of new-∀/c.

For example, this contract:

```
(let ([a (new-∃/c 'a)])
  (-> (-> a a)
      any/c))
```

describes a function that accepts the identity function (or a non-terminating function) and returns an arbitrary value. That is, the first use of the a appears in a positive position and thus inputs to that function are wrapped with an opaque struct. Then, when the function returns, it is checked to see if the result is wrapped, since the second a appears in a negative position.

The new-∃/c construct constructor is dual to new-∀/c.

8.4 Lazy Data-structure Contracts

```
(contract-struct id (field-id ...))
```

> **NOTE:** This library is deprecated; use `struct`, instead. Lazy struct contracts no longer require a separate struct declaration; instead `struct/dc` and `struct/c` work directly with `struct` and `define-struct`.

Like `struct`, but with two differences: they do not define field mutators, and they define two contract constructors: *id*`/c` and *id*`/dc`. The first is a procedure that accepts as many arguments as there are fields and returns a contract for struct values whose fields match the arguments. The second is a syntactic form that also produces contracts on the structs, but the contracts on later fields may depend on the values of earlier fields.

The generated contract combinators are *lazy*: they only verify the contract holds for the portion of some data structure that is actually inspected. More precisely, a lazy data structure contract is not checked until a selector extracts a field of a struct.

```
(id/dc field-spec ...)

field-spec = [field-id contract-expr]
           | [field-id (field-id ...) contract-expr]
```

In each *field-spec* case, the first *field-id* specifies which field the contract applies to; the fields must be specified in the same order as the original `contract-struct`. The first case is for when the contract on the field does not depend on the value of any other field. The second case is for when the contract on the field does depend on some other fields, and the parenthesized *field-id*s indicate which fields it depends on; these dependencies can only be to earlier fields.

```
(define-contract-struct id (field-id ...))
```

NOTE: This library is deprecated; use `struct`, instead. Lazy struct contracts no longer require a separate struct declaration; instead `struct/dc` and `struct/c` work directly with `struct` and `define-struct`.

Like `contract-struct`, but where the constructor's name is `make-`*id*, much like `define-struct`.

8.5 Structure Type Property Contracts

```
(struct-type-property/c value-contract) → contract?
  value-contract : contract?
```

Produces a contract for a structure type property. When the contract is applied to a struct type property, it produces a wrapped struct type property that applies *value-contract* to the value associated with the property when it used to create a new struct type (via `struct`, `make-struct-type`, etc).

The struct type property's accessor function is not affected; if it is exported, it must be protected separately.

As an example, consider the following module. It creates a structure type property, `prop`, whose value should be a function mapping a structure instance to a numeric predicate. The module also exports `app-prop`, which extracts the predicate from a structure instance and applies it to a given value.

```
> (module propmod racket
    (require racket/contract)
    (define-values (prop prop? prop-ref)
      (make-struct-type-property 'prop))
    (define (app-prop x v)
      (((prop-ref x) x) v))
    (provide/contract
     [prop? (-> any/c boolean?)]
     [prop (struct-type-property/c
            (-> prop? (-> integer? boolean?)))]
     [app-prop (-> prop? integer? boolean?)])
    (provide prop-ref))
```

The `structmod` module creates a structure type named `s` with a single field; the value of `prop` is a function that extracts the field value from an instance. Thus the field ought to be an integer predicate, but notice that `structmod` places no contract on `s` enforcing that constraint.

```
> (module structmod racket
    (require 'propmod)
    (struct s (f) #:property prop (lambda (s) (s-f s)))
    (provide (struct-out s)))

> (require 'propmod 'structmod)
```

First we create an `s` instance with an integer predicate, so the constraint on `prop` is in fact satisfied. The first call to `app-prop` is correct; the second simply violates the contract of `app-prop`.

```
> (define s1 (s even?))

> (app-prop s1 5)
#f
> (app-prop s1 'apple)
app-prop: contract violation
  expected: integer?
```

```
  given: 'apple
  in: the 2nd argument of
      (-> prop? integer? boolean?)
  contract from: propmod
  blaming: top-level
   (assuming the contract is correct)
  at: eval:2.0
```

We are able to create `s` instances with values other than integer predicates, but applying `app-prop` on them blames `structmod`, because the function associated with `prop`—that is, `(lambda (s) (s-f s))`—does not always produce a value satisfying `(-> integer? boolean?)`.

```
> (define s2 (s "not a fun"))

> (app-prop s2 5)
prop: contract violation
  expected: a procedure
  given: "not a fun"
  in: the range of
      the struct property value of
      (struct-type-property/c
       (-> prop? (-> integer? boolean?)))
  contract from: propmod
  blaming: structmod
   (assuming the contract is correct)
  at: eval:2.0
> (define s3 (s list))

> (app-prop s3 5)
prop: contract violation
  expected: boolean?
  given: '(5)
  in: the range of
      the range of
      the struct property value of
      (struct-type-property/c
       (-> prop? (-> integer? boolean?)))
  contract from: propmod
  blaming: structmod
   (assuming the contract is correct)
  at: eval:2.0
```

The fix would be to propagate the obligation inherited from `prop` to `s`:

```
(provide (contract-out
```

```
            [struct s ([f (-> integer? boolean?)])]))
```

Finally, if we directly apply the property accessor, `prop-ref`, and then misuse the resulting function, the `propmod` module is blamed:

```
> ((prop-ref s3) 'apple)
prop: broke its contract
  promised: prop?
  produced: 'apple
  in: the 1st argument of
      the struct property value of
      (struct-type-property/c
       (-> prop? (-> integer? boolean?)))
  contract from: propmod
  blaming: propmod
   (assuming the contract is correct)
  at: eval:2.0
```

The `propmod` module has an obligation to ensure a function associated with `prop` is applied only to values satisfying `prop?`. By directly providing `prop-ref`, it enables that constraint to be violated (and thus it is blamed), even though the bad application actually occurs elsewhere.

Generally there is no need to provide a structure type property accessor at all; it is typically only used by other functions within the module. But if it must be provided, it should be protected thus:

```
(provide (contract-out
          [prop-ref (-> prop? (-> prop? (-> integer? boolean?)))]))
```

8.6 Attaching Contracts to Values

```
(contract-out p/c-item ...)
```

```
      p/c-item = (struct id/super ((id contract-expr) ...)
                   struct-option)
               | (rename orig-id id contract-expr)
               | (id contract-expr)
               | #:∃ poly-variables
               | #:exists poly-variables
               | #:∀ poly-variables
               | #:forall poly-variables

poly-variables = id
               | (id ...)

      id/super = id
               | (id super-id)

 struct-option =
               | #:omit-constructor
```

A *provide-spec* for use in `provide` (currently only for the same phase level as the `provide` form; for example, `contract-out` cannot be nested within `for-syntax`). Each *id* is provided from the module. In addition, clients of the module must live up to the contract specified by *contract-expr* for each export.

The `contract-out` form treats modules as units of blame. The module that defines the provided variable is expected to meet the positive (co-variant) positions of the contract. Each module that imports the provided variable must obey the negative (contra-variant) positions of the contract. Each *contract-expr* in a `contract-out` form is effectively moved to the end of the enclosing module, so a *contract-expr* can refer to variables that are defined later in the same module.

Only uses of the contracted variable outside the module are checked. Inside the module, no contract checking occurs.

The `rename` form of `contract-out` exports the first variable (the internal name) with the name specified by the second variable (the external name).

The `struct` form of `contract-out` provides a structure-type definition, and each field has a contract that dictates the contents of the fields. The structure-type definition must appear before the `provide` clause within the enclosing module. If the structure type has a parent, the second `struct` form (above) must be used, with the first name referring to the structure type to export and the second name referring to the parent structure type. Unlike a `struct` definition, however, all of the fields (and their contracts) must be listed. The contract on the fields that the sub-struct shares with its parent are only used in the contract for the sub-struct's constructor, and the selector or mutators for the super-struct are not provided. The exported structure-type name always doubles as a constructor, even if the original structure-type name does not act as a constructor. If the `#:omit-constructor` option is present, the

constructor is not provided.

The `#:∃`, `#:exists`, `#:∀`, and `#:forall` clauses define new abstract contracts. The variables are bound in the remainder of the `contract-out` form to new contracts that hide the values they accept and ensure that the exported functions are treated parametrically. See `new-∃/c` and `new-∀/c` for details on how the clauses hide the values.

The implementation of `contract-out` uses `syntax-property` to attach properties to the code it generates that records the syntax of the contracts in the fully expanded program. Specifically, the symbol `'provide/contract-original-contract` is bound to vectors of two elements, the exported identifier and a syntax object for the expression that produces the contract controlling the export.

```
(recontract-out id ...)
```

A *provide-spec* for use in `provide` (currently, just like `contract-out`, only for the same phase level as the `provide` form).

It re-exports *id*, but with positive blame associated to the module containing `recontract-out` instead of the location of the original site of *id*.

This can be useful when a public module wants to export an identifier from a private module but where any contract violations should be reported in terms of the public module instead of the private one.

Examples:

```
> (module private-implementation racket/base
    (require racket/contract)
    (define (recip x) (/ 1 x))
    (define (non-zero? x) (not (= x 0)))
    (provide/contract [recip (-> (and/c real? non-zero?)
                                 (between/c -1 1))]))

> (module public racket/base
    (require racket/contract
             'private-implementation)
    (provide (recontract-out recip)))

> (require 'public)

> (recip +nan.0)
recip: broke its contract
  promised: (between/c -1 1)
  produced: +nan.0
  in: the range of
      (->
```

```
        (and/c real? non-zero?)
        (between/c -1 1))
  contract from: public
  blaming: public
   (assuming the contract is correct)
  at: eval:3.0
```

Replacing the use of `contract-out` with just `recip` would result in a contract violation blaming the private module.

```
(provide/contract p/c-item ...)
```

A legacy shorthand for `(provide (contract-out p/c-item ...))`, except that a *contract-expr* within `provide/contract` is evaluated at the position of the `provide/contract` form instead of at the end of the enclosing module.

8.6.1 Nested Contract Boundaries

```
(require racket/contract/region)      package: base
```

```
(with-contract blame-id (wc-export ...) free-var-list ... body ...+)
(with-contract blame-id results-spec free-var-list ... body ...+)

    wc-export = (id contract-expr)

  result-spec = #:result contract-expr
              | #:results (contract-expr ...)

free-var-list =
              | #:freevar id contract-expr
              | #:freevars ([id contract-expr] ...)
```

Generates a local contract boundary.

The first `with-contract` form cannot appear in expression position. All names defined within the first `with-contract` form are visible externally, but those names listed in the *wc-export* list are protected with the corresponding contract. The *body* of the form allows definition/expression interleaving if its context does.

The second `with-contract` form must appear in expression position. The final *body* expression should return the same number of values as the number of contracts listed in the *result-spec*, and each returned value is contracted with its respective contract. The sequence of *body* forms is treated as for `let`.

The *blame-id* is used for the positive positions of contracts paired with exported *id*s.

Contracts broken within the `with-contract` *body* will use the *blame-id* for their negative position.

If a free-var-list is given, then any uses of the free variables inside the *body* will be protected with contracts that blame the context of the `with-contract` form for the positive positions and the `with-contract` form for the negative ones.

```
(define/contract id contract-expr free-var-list init-value-expr)
(define/contract (head args) contract-expr free-var-list body ...+)
```

Works like `define`, except that the contract *contract-expr* is attached to the bound value. For the definition of *head* and *args*, see `define`. For the definition of *free-var-list*, see `with-contract`.

Examples:

```
> (define/contract distance (>=/c 0) 43.52)

> (define/contract (furlongs->feet fr)
    (-> real? real?)
    (* 660 fr))

; a contract violation expected here:
> (furlongs->feet "not a furlong")
furlongs->feet: contract violation
  expected: real?
  given: "not a furlong"
  in: the 1st argument of
      (-> real? real?)
  contract from: (function furlongs->feet)
  blaming: top-level
   (assuming the contract is correct)
  at: eval:3.0
```

The `define/contract` form treats the individual definition as a contract region. The definition itself is responsible for positive (co-variant) positions of the contract, and references to *id* outside of the definition must meet the negative positions of the contract. Since the contract boundary is between the definition and the surrounding context, references to *id* inside the `define/contract` form are not checked.

Examples:

```
; an unsual predicate that prints when called
> (define (printing-int? x)
    (displayln "I was called")
    (exact-integer? x))
```

```
> (define/contract (fact n)
    (-> printing-int? printing-int?)
    (if (zero? n)
        1
        (* n (fact (sub1 n)))))

> (fact 5)
I was called
I was called
120
; only prints twice, not for each recursive call
```

If a free-var-list is given, then any uses of the free variables inside the *body* will be protected with contracts that blame the context of the `define/contract` form for the positive positions and the `define/contract` form for the negative ones.

Examples:

```
> (define (integer->binary-string n)
    (number->string n 2))

> (define/contract (numbers->strings lst)
    (-> (listof number?) (listof string?))
    #:freevar integer->binary-string (-> exact-integer? string?)
    ; mistake, lst might contain inexact numbers
    (map integer->binary-string lst))

> (numbers->strings '(4.0 3.3 5.8))
integer->binary-string: contract violation
  expected: exact-integer?
  given: 4.0
  in: the 1st argument of
      (-> exact-integer? string?)
  contract from: top-level
  blaming: (function numbers->strings)
   (assuming the contract is correct)
  at: eval:3.0
```

```
(define-struct/contract struct-id ([field contract-expr] ...)
                        struct-option ...)
(define-struct/contract (struct-id super-struct-id)
                        ([field contract-expr] ...)
                        struct-option ...)
```

Works like `define-struct`, except that the arguments to the constructor, accessors, and mutators are protected by contracts. For the definitions of *field* and *struct-option*, see `define-struct`.

The `define-struct/contract` form only allows a subset of the *`struct-option`* keywords: `#:mutable`, `#:transparent`, `#:auto-value`, `#:omit-define-syntaxes`, `#:property` and `#:omit-define-values`.

Examples:

```
> (define-struct/contract fish ([color number?]))

> (make-fish 5)
#<fish>
> (make-fish #f)
make-fish: contract violation
  expected: number?
  given: #f
  in: the 1st argument of
      (-> number? symbol? any)
  contract from: (struct fish)
  blaming: top-level
   (assuming the contract is correct)
> (define-struct/contract (salmon fish) ([ocean symbol?]))

> (make-salmon 5 'atlantic)
#<salmon>
> (make-salmon 5 #f)
make-salmon: contract violation
  expected: symbol?
  given: #f
  in: the 2nd argument of
      (-> any/c symbol? symbol? any)
  contract from: (struct salmon)
  blaming: top-level
   (assuming the contract is correct)
> (make-salmon #f 'pacific)
make-fish: contract violation
  expected: number?
  given: #f
  in: the 1st argument of
      (-> number? symbol? any)
  contract from: (struct fish)
  blaming: top-level
   (assuming the contract is correct)
```

```
(invariant-assertion invariant-expr expr)
```

Establishes an invariant of *`expr`*, determined by *`invariant-expr`*.

Unlike the specification of a contract, an `invariant-assertion` does not establish a

boundary between two parties. Instead, it simply attaches a logical assertion to the value. Because the form uses contract machinery to check the assertion, the surrounding module is treated as the party to be blamed for any violations of the assertion.

This means, for example, that the assertion is checked on recursive calls, when an invariant is used on the right-hand side of a definition:

Examples:

```
> (define furlongss->feets
    (invariant-assertion
     (-> (listof real?) (listof real?))
     (λ (l)
       (cond
         [(empty? l) empty]
         [else
          (if (= 327 (car l))
              (furlongss->feets (list "wha?"))
              (cons (furlongs->feet (first l))
                    (furlongss->feets (rest l))))]))))

> (furlongss->feets (list 1 2 3))
'(660 1320 1980)
> (furlongss->feets (list 1 327 3))
furlongss->feets: contract violation
  expected: real?
  given: "wha?"
  in: an element of
      the 1st argument of
      (-> (listof real?) (listof real?))
  contract from: top-level
  blaming: top-level
   (assuming the contract is correct)
  at: eval:5.0
```

Added in version 6.0.1.11 of package `base`.

`current-contract-region`

Bound by `define-syntax-parameter`, this contains information about the current contract region, used by the above forms to determine the candidates for blame assignment.

8.6.2 Low-level Contract Boundaries

```
(define-module-boundary-contract id
  orig-id
  contract-expr
  pos-blame-party
  source-loc)

pos-blame-party =
                | #:pos-source pos-source-expr

     source-loc =
                | #:srcloc srcloc-expr
```

Defines *id* to be *orig-id*, but with the contract *contract-expr*.

The identifier *id* is defined as a macro transformer that consults the context of its use to determine the name for negative blame assignment (using the entire module where a reference appears as the negative party).

The positive party defaults to the module containing the use of `define-module-boundary-contract`, but can be specified explicitly via the `#:pos-source` keyword.

The source location used in the blame error messages for the location of the place where the contract was put on the value defaults to the source location of the use of `define-module-boundary-contract`, but can be specified via the `#:srcloc` argument, in which case it can be any of the things that the third argument to `datum->syntax` can be.

Examples:

```
> (module server racket/base
    (require racket/contract/base)
    (define (f x) #f)
    (define-module-boundary-contract g f (-> integer? integer?))
    (provide g))

> (module client racket/base
    (require 'server)
    (define (clients-fault) (g #f))
    (define (servers-fault) (g 1))
    (provide servers-fault clients-fault))

> (require 'client)

> (clients-fault)
g: contract violation
  expected: integer?
  given: #f
  in: the 1st argument of
```

```
      (-> integer? integer?)
  contract from: 'server
  blaming: client
   (assuming the contract is correct)
  at: eval:2.0
> (servers-fault)
g: broke its contract
  promised: integer?
  produced: #f
  in: the range of
      (-> integer? integer?)
  contract from: 'server
  blaming: (quote server)
   (assuming the contract is correct)
  at: eval:2.0
```

```
(contract contract-expr to-protect-expr
          positive-blame-expr negative-blame-expr)
(contract contract-expr to-protect-expr
          positive-blame-expr negative-blame-expr
          value-name-expr source-location-expr)
```

The primitive mechanism for attaching a contract to a value. The purpose of `contract` is as a target for the expansion of some higher-level contract specifying form.

The `contract` expression adds the contract specified by *`contract-expr`* to the value produced by *`to-protect-expr`*. The result of a `contract` expression is the result of the *`to-protect-expr`* expression, but with the contract specified by *`contract-expr`* enforced on *`to-protect-expr`*.

The values of *`positive-blame-expr`* and *`negative-blame-expr`* indicate how to assign blame for positive and negative positions of the contract specified by *`contract-expr`*. They may be any value, and are formatted as by `display` for purposes of contract violation error messages.

If specified, *`value-name-expr`* indicates a name for the protected value to be used in error messages. If not supplied, or if *`value-name-expr`* produces `#f`, no name is printed. Otherwise, it is also formatted as by `display`. More precisely, the *`value-name-expr`* ends up in the `blame-name` field of the blame record, which is used as the first portion of the error message.

Examples:

```
> (contract integer? #f 'pos 'neg 'timothy #f)
timothy: broke its contract
  promised: integer?
  produced: #f
```

```
  in: integer?
  contract from: pos
  blaming: pos
   (assuming the contract is correct)
> (contract integer? #f 'pos 'neg #f #f)
broke its contract:
  promised: integer?
  produced: #f
  in: integer?
  contract from: pos
  blaming: pos
   (assuming the contract is correct)
```

If specified, *source-location-expr* indicates the source location reported by contract violations. The expression must produce a `srcloc` structure, syntax object, `#f`, or a list or vector in the format accepted by the third argument to `datum->syntax`.

8.7 Building New Contract Combinators

```
(require racket/contract/combinator)
```
package: `base`

Contracts are represented internally as functions that accept information about the contract (who is to blame, source locations, etc.) and produce projections (in the spirit of Dana Scott) that enforce the contract. A projection is a function that accepts an arbitrary value, and returns a value that satisfies the corresponding contract. For example, a projection that accepts only integers corresponds to the contract `(flat-contract integer?)`, and can be written like this:

```
(define int-proj
  (λ (x)
    (if (integer? x)
        x
        (signal-contract-violation))))
```

As a second example, a projection that accepts unary functions on integers looks like this:

```
(define int->int-proj
  (λ (f)
    (if (and (procedure? f)
             (procedure-arity-includes? f 1))
        (λ (x) (int-proj (f (int-proj x))))
        (signal-contract-violation))))
```

Although these projections have the right error behavior, they are not quite ready for use as contracts, because they do not accommodate blame and do not provide good error messages.

In order to accommodate these, contracts do not just use simple projections, but use functions that accept a *blame object* encapsulating the names of two parties that are the candidates for blame, as well as a record of the source location where the contract was established and the name of the contract. They can then, in turn, pass that information to `raise-blame-error` to signal a good error message.

Here is the first of those two projections, rewritten for use in the contract system:

```
(define (int-proj blame)
  (λ (x)
    (if (integer? x)
        x
        (raise-blame-error
         blame
         val
         '(expected: "<integer>" given: "~e")
         val))))
```

The new argument specifies who is to be blamed for positive and negative contract violations.

Contracts, in this system, are always established between two parties. One party provides some value according to the contract, and the other consumes the value, also according to the contract. The first is called the "positive" person and the second the "negative". So, in the case of just the integer contract, the only thing that can go wrong is that the value provided is not an integer. Thus, only the positive party can ever accrue blame. The `raise-blame-error` function always blames the positive party.

Compare that to the projection for our function contract:

```
(define (int->int-proj blame)
  (define dom (int-proj (blame-swap blame)))
  (define rng (int-proj blame))
  (λ (f)
    (if (and (procedure? f)
             (procedure-arity-includes? f 1))
        (λ (x) (rng (f (dom x))))
        (raise-blame-error
         blame
         val
         '(expected "a procedure of one argument" given: "~e")
         val))))
```

In this case, the only explicit blame covers the situation where either a non-procedure is supplied to the contract or the procedure does not accept one argument. As with the integer projection, the blame here also lies with the producer of the value, which is why `raise-blame-error` is passed `blame` unchanged.

The checking for the domain and range are delegated to the `int-proj` function, which is supplied its arguments in the first two lines of the `int->int-proj` function. The trick here is that, even though the `int->int-proj` function always blames what it sees as positive, we can swap the blame parties by calling `blame-swap` on the given blame object, replacing the positive party with the negative party and vice versa.

This technique is not merely a cheap trick to get the example to work, however. The reversal of the positive and the negative is a natural consequence of the way functions behave. That is, imagine the flow of values in a program between two modules. First, one module defines a function, and then that module is required by another. So, far the function itself has to go from the original, providing module to the requiring module. Now, imagine that the providing module invokes the function, supplying it an argument. At this point, the flow of values reverses. The argument is traveling back from the requiring module to the providing module! And finally, when the function produces a result, that result flows back in the original direction. Accordingly, the contract on the domain reverses the positive and the negative blame parties, just like the flow of values reverses.

We can use this insight to generalize the function contracts and build a function that accepts any two contracts and returns a contract for functions between them.

This projection also goes further and uses `blame-add-context` to improve the error messages when a contract violation is detected.

```
(define (make-simple-function-contract dom-proj range-proj)
  (λ (blame)
    (define dom (dom-proj (blame-add-context blame
                                             "the argument of"
                                             #:swap? #t)))
    (define rng (range-proj (blame-add-context blame
                                               "the range of")))
    (λ (f)
      (if (and (procedure? f)
               (procedure-arity-includes? f 1))
          (λ (x) (rng (f (dom x))))
          (raise-blame-error
           blame
           val
           '(expected "a procedure of one argument" given: "~e")
           val)))))
```

While these projections are supported by the contract library and can be used to build new contracts, the contract library also supports a different API for projections that can be more efficient. Specifically, a *val first projection* accepts a blame object without the negative blame information and then returns a function that accepts the value to be contracted, and then finally accepts the name of the negative party to the contract before returning the value with the contract. Rewriting `int->int-proj` to use this API looks like this:

```
(define (int->int-proj blame)
  (define dom-blame (blame-add-context blame
                                       "the argument of"
                                       #:swap? #t))
  (define rng-blame (blame-add-context blame "the range of"))
  (define (check-int v to-blame neg-party)
    (unless (integer? x)
      (raise-blame-error
       to-blame #:missing-party neg-party
       val
       '(expected "an integer" given: "~e")
       val)))
  (λ (f)
    (if (and (procedure? f)
             (procedure-arity-includes? f 1))
        (λ (neg-party)
          (λ (x)
            (check-int x dom-blame neg-party)
            (define ans (f x))
            (check-int ans rng-blame neg-party)
            ans))
        (λ (neg-party)
          (raise-blame-error
           blame #:missing-party neg-party
           val
           '(expected "a procedure of one argument" given: "~e")
           val)))))
```

The advantage of this style of contract is that the *blame* and *v* arguments can be supplied on the server side of the contract boundary and the result can be used for every different client. With the simpler situation, a new blame object has to be created for each client.

Projections like the ones described above, but suited to other, new kinds of value you might make, can be used with the contract library primitives below.

```
(make-contract [#:name name
                #:first-order test
                #:val-first-projection val-first-proj
                #:projection proj
                #:stronger stronger
                #:list-contract is-list-contract?])
 → contract?
 name : any/c = 'anonymous-contract
 test : (-> any/c any/c) = (λ (x) #t)
 val-first-proj : (or/c #f (-> blame? (-> any/c (-> any/c any/c))))
                  = #f
 proj : (-> blame? (-> any/c any/c))
        (λ (b)
          (λ (x)
            (if (test x)
              x
      =       (raise-blame-error
               b x
               '(expected: "~a" given: "~e")
               name x))))
 stronger : (or/c #f (-> contract? contract? boolean?)) = #f
 is-list-contract? : boolean? = #f
(make-chaperone-contract
 [#:name name
  #:first-order test
  #:val-first-projection val-first-proj
  #:projection proj
  #:stronger stronger
  #:list-contract is-list-contract?])
 → chaperone-contract?
 name : any/c = 'anonymous-chaperone-contract
 test : (-> any/c any/c) = (λ (x) #t)
 val-first-proj : (or/c #f (-> blame? (-> any/c (-> any/c any/c))))
                  = #f
 proj : (-> blame? (-> any/c any/c))
        (λ (b)
          (λ (x)
            (if (test x)
              x
      =       (raise-blame-error
               b x
               '(expected: "~a" given: "~e")
               name x))))
 stronger : (or/c #f (-> contract? contract? boolean?)) = #f
 is-list-contract? : boolean? = #f
```

```
(make-flat-contract [#:name name
                     #:first-order test
                     #:val-first-projection val-first-proj
                     #:projection proj
                     #:stronger stronger
                     #:list-contract is-list-contract?])
 → flat-contract?
  name : any/c = 'anonymous-flat-contract
  test : (-> any/c any/c) = (λ (x) #t)
  val-first-proj : (or/c #f (-> blame? (-> any/c (-> any/c any/c))))
                 = #f
  proj : (-> blame? (-> any/c any/c))
         (λ (b)
           (λ (x)
             (if (test x)
               x
       =       (raise-blame-error
                b x
                '(expected: "~a" given: "~e")
                name x))))
  stronger : (or/c #f (-> contract? contract? boolean?)) = #f
  is-list-contract? : boolean? = #f
```

The precise details of the `val-first-projection` argument are subject to change. (Probably also the default values of the `project` arguments will change.

These functions build simple higher-order contracts, chaperone contracts, and flat contracts, respectively. They both take the same set of three optional arguments: a name, a first-order predicate, and a blame-tracking projection.

The *name* argument is any value to be rendered using `display` to describe the contract when a violation occurs. The default name for simple higher-order contracts is `anonymous-contract`, for chaperone contracts is `anonymous-chaperone-contract`, and for flat contracts is `anonymous-flat-contract`.

The first-order predicate *test* can be used to determine which values the contract applies to; usually, this is the set of values for which the contract fails immediately without any higher-order wrapping. This test is used by `contract-first-order-passes?`, and indirectly by `or/c` to determine which of multiple higher-order contracts to wrap a value with. The default test accepts any value.

The projection *proj* defines the behavior of applying the contract. It is a curried function of two arguments: the first application accepts a blame object, and the second accepts a value to protect with the contract. The projection must either produce the value, suitably wrapped to enforce any higher-order aspects of the contract, or signal a contract violation using `raise-blame-error`. The default projection produces an error when the first-order test fails, and

produces the value unchanged otherwise.

Projections for chaperone contracts must produce a value that passes `chaperone-of?` when compared with the original, uncontracted value. Projections for flat contracts must fail precisely when the first-order test does, and must produce the input value unchanged otherwise. Applying a flat contract may result in either an application of the predicate, or the projection, or both; therefore, the two must be consistent. The existence of a separate projection only serves to provide more specific error messages. Most flat contracts do not need to supply an explicit projection.

The *`stronger`* argument is used to implement `contract-stronger?`. The first argument is always the contract itself and the second argument is whatever was passed as the second argument to `contract-stronger?`.

The *`is-list-contract?`* argument is used by the `list-contract?` predicate to determine if this is a contract that accepts only `list?` values.

Examples:

```
(define int/c
  (make-flat-contract #:name 'int/c #:first-order integer?))

> (contract int/c 1 'positive 'negative)
1
> (contract int/c "not one" 'positive 'negative)
eval:4:0: broke its contract
  promised: int/c
  produced: "not one"
  in: int/c
  contract from: positive
  blaming: positive
   (assuming the contract is correct)
> (int/c 1)
#t
> (int/c "not one")
#f
(define int->int/c
  (make-contract
   #:name 'int->int/c
   #:first-order
   (λ (x) (and (procedure? x) (procedure-arity-includes? x 1)))
   #:projection
   (λ (b)
     (let ([domain ((contract-projection int/c) (blame-swap b))]
           [range ((contract-projection int/c) b)])
       (λ (f)
         (if (and (procedure? f) (procedure-arity-includes? f 1))
```

```
            (λ (x) (range (f (domain x))))
            (raise-blame-error
             b f
             '(expected "a function of one argument" 'given: "~e")
             f)))))))

> (contract int->int/c "not fun" 'positive 'negative)
regexp-match: contract violation
  expected: (or/c string? bytes? path? input-port?)
  given: "given:
  argument position: 2nd
  other arguments...:
   #rx"^[\n ]"
(define halve
  (contract int->int/c (λ (x) (/ x 2)) 'positive 'negative))

> (halve 2)
1
> (halve 1/2)
halve: contract violation
  expected: int/c
  given: 1/2
  in: ...
      int->int/c
  contract from: positive
  blaming: negative
   (assuming the contract is correct)
> (halve 1)
halve: broke its contract
  promised: int/c
  produced: 1/2
  in: ...
      int->int/c
  contract from: positive
  blaming: positive
   (assuming the contract is correct)
```

Changed in version 6.0.1.13 of package `base`: Added the `#:list-contract?` argument.

```
(build-compound-type-name c/s ...) → any
  c/s : any/c
```

Produces an S-expression to be used as a name for a contract. The arguments should be either contracts or symbols. It wraps parentheses around its arguments and extracts the names from any contracts it is supplied with.

```
(coerce-contract id x) → contract?
  id : symbol?
  x : any/c
```

Converts a regular Racket value into an instance of a contract struct, converting it according to the description of contracts.

If *x* is not one of the coercible values, `coerce-contract` signals an error, using the first argument in the error message.

```
(coerce-contracts id xs) → (listof contract?)
  id : symbol?
  xs : (listof any/c)
```

Coerces all of the arguments in 'xs' into contracts (via `coerce-contract/f`) and signals an error if any of them are not contracts. The error messages assume that the function named by *id* got *xs* as its entire argument list.

```
(coerce-chaperone-contract id x) → chaperone-contract?
  id : symbol?
  x : any/c
```

Like `coerce-contract`, but requires the result to be a chaperone contract, not an arbitrary contract.

```
(coerce-chaperone-contracts id x) → (listof chaperone-contract?)
  id : symbol?
  x : (listof any/c)
```

Like `coerce-contracts`, but requires the results to be chaperone contracts, not arbitrary contracts.

```
(coerce-flat-contract id x) → flat-contract?
  id : symbol?
  x : any/c
```

Like `coerce-contract`, but requires the result to be a flat contract, not an arbitrary contract.

```
(coerce-flat-contracts id x) → (listof flat-contract?)
  id : symbol?
  x : (listof any/c)
```

Like `coerce-contracts`, but requires the results to be flat contracts, not arbitrary contracts.

```
(coerce-contract/f x) → (or/c contract? #f)
  x : any/c
```

Like `coerce-contract`, but returns `#f` if the value cannot be coerced to a contract.

8.7.1 Blame Objects

```
(blame? x) → boolean?
  x : any/c
```

This predicate recognizes blame objects.

```
(blame-add-context  blame
                    context
                   [#:important important
                    #:swap? swap?])        → blame?
  blame : blame?
  context : (or/c string? #f)
  important : (or/c string? #f) = #f
  swap? : boolean? = #f
```

Adds some context information to blame error messages that explicates which portion of the contract failed (and that gets rendered by `raise-blame-error`).

The *context* argument describes one layer of the portion of the contract, typically of the form `"the 1st argument of"` (in the case of a function contract) or `"a conjunct of"` (in the case of an `and/c` contract).

For example, consider this contract violation:

```
> (define/contract f
    (list/c (-> integer? integer?))
    (list (λ (x) x)))

> ((car f) #f)
f: contract violation
  expected: integer?
  given: #f
  in: the 1st argument of
        the 1st element of
        (list/c (-> integer? integer?))
  contract from: (definition f)
  blaming: top-level
   (assuming the contract is correct)
  at: eval:2.0
```

It shows that the portion of the contract being violated is the first occurrence of `integer?`, because the `->` and the `list/c` combinators each internally called `blame-add-context` to add the two lines following “in” in the error message.

The *important* argument is used to build the beginning part of the contract violation. The last *important* argument that gets added to a blame object is used. The `class/c` contract adds an important argument, as does the `->` contract (when `->` knows the name of the function getting the contract).

The *swap?* argument has the effect of calling `blame-swap` while adding the layer of context, but without creating an extra blame object.

The context information recorded in blame structs keeps track of combinators that do not add information, and add the string `"..."` for them, so programmers at least see that there was some context they are missing in the error messages. Accordingly, since there are combinators that should not add any context (e.g., `recursive-contract`), passing `#f` as the context string argument avoids adding the `"..."` string.

```
(blame-positive b) → any/c
  b : blame?
(blame-negative b) → any/c
  b : blame?
```

These functions produce printable descriptions of the current positive and negative parties of a blame object.

```
(blame-contract b) → any/c
  b : blame?
```

This function produces a description of the contract associated with a blame object (the result of `contract-name`).

```
(blame-value b) → any/c
  b : blame?
```

This function produces the name of the value to which the contract was applied, or `#f` if no name was provided.

```
(blame-source b) → srcloc?
  b : blame?
```

This function produces the source location associated with a contract. If no source location was provided, all fields of the structure will contain `#f`.

```
(blame-swap b) → blame?
  b : blame?
```

This function swaps the positive and negative parties of a blame object. (See also `blame-add-context`.)

```
(blame-original? b) → boolean?
  b : blame?
(blame-swapped? b) → boolean?
  b : blame?
```

These functions report whether the current blame of a given blame object is the same as in the original contract invocation (possibly of a compound contract containing the current one), or swapped, respectively. Each is the negation of the other; both are provided for convenience and clarity.

```
(blame-replace-negative b neg) → blame?
  b : blame?
  neg : any/c
```

Produces a `blame?` object just like `b` except that it uses `neg` instead of the negative position `b` has.

```
(blame-update b pos neg) → blame?
  b : blame?
  pos : any/c
  neg : any/c
```

Produces a `blame?` object just like `b` except that it adds `pos` and `neg` to the positive and negative parties of `b` respectively.

```
(raise-blame-error b x fmt v ...) → none/c
  b : blame?
  x : any/c
        (or/c string?
              (listof (or/c string?
  fmt :                        'given 'given:
                               'expected 'expected:)))
  v : any/c
```

Signals a contract violation. The first argument, `b`, records the current blame information, including positive and negative parties, the name of the contract, the name of the value, and the source location of the contract application. The second argument, `x`, is the value that failed to satisfy the contract.

The remaining arguments are a format string, `fmt`, and its arguments, `v ...`, specifying an error message specific to the precise violation.

If `fmt` is a list, then the elements are concatenated together (with spaces added, unless there are already spaces at the ends of the strings), after first replacing symbols with either their string counterparts, or replacing `'given` with `"produced"` and `'expected` with

"promised", depending on whether or not the *b* argument has been swapped or not (see blame-swap).

If *fmt* contains the symbols 'given: or 'expected:, they are replaced like 'given: and 'expected: are, but the replacements are prefixed with the string "\n " to conform to the error message guidelines in §10.2.1 "Error Message Conventions".

```
(struct exn:fail:contract:blame exn:fail:contract (object)
    #:extra-constructor-name make-exn:fail:contract:blame)
  object : blame?
```

This exception is raised to signal a contract error. The object field contains a blame object associated with a contract violation.

```
(current-blame-format) → (-> blame? any/c string? string?)
(current-blame-format proc) → void?
  proc : (-> blame? any/c string? string?)
```

A parameter that is used when constructing a contract violation error. Its value is procedure that accepts three arguments:

- the blame object for the violation,
- the value that the contract applies to, and
- a message indicating the kind of violation.

The procedure then returns a string that is put into the contract error message. Note that the value is often already included in the message that indicates the violation.

Examples:

```
(define (show-blame-error blame value message)
  (string-append
   "Contract Violation!\n"
   (format "Guilty Party: ~a\n" (blame-positive blame))
   (format "Innocent Party: ~a\n" (blame-negative blame))
   (format "Contracted Value Name: ~a\n" (blame-value blame))
   (format "Contract Location: ~s\n" (blame-source blame))
   (format "Contract Name: ~a\n" (blame-contract blame))
   (format "Offending Value: ~s\n" value)
   (format "Offense: ~a\n" message)))

> (current-blame-format show-blame-error)
```

```
> (define/contract (f x)
    (-> integer? integer?)
    (/ x 2))

> (f 2)
1
> (f 1)
Contract Violation!
Guilty Party: (function f)
Innocent Party: top-level
Contracted Value Name: f
Contract Location: #(struct:srcloc eval 4 0 4 1)
Contract Name: (-> integer? integer?)
Offending Value: 1/2
Offense: promised: integer?
  produced: 1/2

> (f 1/2)
Contract Violation!
Guilty Party: top-level
Innocent Party: (function f)
Contracted Value Name: f
Contract Location: #(struct:srcloc eval 4 0 4 1)
Contract Name: (-> integer? integer?)
Offending Value: 1/2
Offense: expected: integer?
  given: 1/2
```

8.7.2 Contracts as structs

The property `prop:contract` allows arbitrary structures to act as contracts. The property `prop:chaperone-contract` allows arbitrary structures to act as chaperone contracts; `prop:chaperone-contract` inherits `prop:contract`, so chaperone contract structures may also act as general contracts. The property `prop:flat-contract` allows arbitrary structures to act as flat contracts; `prop:flat-contract` inherits both `prop:chaperone-contract` and `prop:procedure`, so flat contract structures may also act as chaperone contracts, as general contracts, and as predicate procedures.

```
prop:contract : struct-type-property?
prop:chaperone-contract : struct-type-property?
prop:flat-contract : struct-type-property?
```

These properties declare structures to be contracts or flat contracts, respectively. The value for `prop:contract` must be a contract property constructed by `build-contract-`

property; likewise, the value for prop:chaperone-contract must be a chaperone contract property constructed by build-chaperone-contract-property and the value for prop:flat-contract must be a flat contract property constructed by build-flat-contract-property.

```
prop:contracted : struct-type-property?
impersonator-prop:contracted : impersonator-property?
```

These properties attach a contract value to the protected structure, chaperone, or impersonator value. The function has-contract? returns #t for values that have one of these properties, and value-contract extracts the value from the property (which is expected to be the contract on the value).

```
prop:blame : struct-type-property?
impersonator-prop:blame : impersonator-property?
```

These properties attach a blame information to the protected structure, chaperone, or impersonator value. The function blame-contract? returns #t for values that have one of these properties, and blame-contract extracts the value from the property (which is expected to be the blame record for the contract on the value).

```
(build-flat-contract-property
 [#:name get-name
  #:first-order get-first-order
  #:val-first-projection val-first-proj
  #:projection get-projection
  #:stronger stronger
  #:generate generate]
  #:exercise exercise
 [#:list-contract? is-list-contract?])
→ flat-contract-property?
 get-name : (-> contract? any/c)
          = (λ (c) 'anonymous-flat-contract)
 get-first-order : (-> contract? (-> any/c boolean?))
                 = (λ (c) (λ (x) #t))
 val-first-proj : (or/c #f (-> contract? blame? (-> any/c (-> any/c any/c))))
                = #f
 get-projection : (-> contract? (-> blame? (-> any/c any/c)))
                  (λ (c)
                    (λ (b)
                      (λ (x)
                =       (if ((get-first-order c) x)
                            x
                            (raise-blame-error
                             b x '(expected: "~a" given: "~e")
                             (get-name c) x)))))
 stronger : (or/c (-> contract? contract? boolean?) #f) = #f
            (->i ([c contract?])
                 ([generator
                   (c)
 generate :        (-> (and/c positive? real?)
                       (or/c #f
                             (-> c)))]))
            (->i ([c contract?])
                 ([result
                   (c)
 exercise :        (-> (and/c positive? real?)
                       (values
                        (-> c void?)
                        (listof contract?)))]))
 is-list-contract? : (-> contract? boolean?) = (λ (c) #f)
```

```
(build-chaperone-contract-property
 [#:name get-name
  #:first-order get-first-order
  #:val-first-projection val-first-proj
  #:projection get-projection
  #:stronger stronger
  #:generate generate]
  #:exercise exercise
 [#:list-contract? is-list-contract?])
→ chaperone-contract-property?
 get-name : (-> contract? any/c)
          = (λ (c) 'anonymous-chaperone-contract)
 get-first-order : (-> contract? (-> any/c boolean?))
                 = (λ (c) (λ (x) #t))
 val-first-proj : (or/c #f (-> contract? blame? (-> any/c (-> any/c any/c))))
                = #f
 get-projection : (-> contract? (-> blame? (-> any/c any/c)))
                  (λ (c)
                    (λ (b)
                      (λ (x)
                =       (if ((get-first-order c) x)
                            x
                            (raise-blame-error
                             b x '(expected: "~a" given: "~e")
                             (get-name c) x)))))
 stronger : (or/c (-> contract? contract? boolean?) #f) = #f
            (->i ([c contract?])
                 ([generator
                   (c)
 generate :        (-> (and/c positive? real?)
                       (or/c #f
                             (-> c)))]))
            (->i ([c contract?])
                 ([result
                   (c)
 exercise :        (-> (and/c positive? real?)
                       (values
                        (-> c void?)
                        (listof contract?)))]))
 is-list-contract? : (-> contract? boolean?) = (λ (c) #f)
```

```
(build-contract-property
 [#:name get-name
  #:first-order get-first-order
  #:val-first-projection val-first-proj
  #:projection get-projection
  #:stronger stronger
  #:generate generate]
  #:exercise exercise
 [#:list-contract? is-list-contract?])
→ contract-property?
 get-name : (-> contract? any/c) = (λ (c) 'anonymous-contract)
 get-first-order : (-> contract? (-> any/c boolean?))
                 = (λ (c) (λ (x) #t))
 val-first-proj : (or/c #f (-> contract? blame? (-> any/c (-> any/c any/c))))
                = #f
 get-projection : (-> contract? (-> blame? (-> any/c any/c)))
                  (λ (c)
                    (λ (b)
                      (λ (x)
                =       (if ((get-first-order c) x)
                            x
                            (raise-blame-error
                             b x '(expected: "~a" given: "~e")
                             (get-name c) x)))))
 stronger : (or/c (-> contract? contract? boolean?) #f) = #f
            (->i ([c contract?])
                 ([generator
                   (c)
 generate :        (-> (and/c positive? real?)          = #f
                       (or/c #f
                             (-> c)))]))
            (->i ([c contract?])
                 ([result
                   (c)
 exercise :        (-> (and/c positive? real?)
                       (values
                        (-> c void?)
                        (listof contract?)))]))
 is-list-contract? : (-> contract? boolean?) = (λ (c) #f)
```

The precise details of the `val-first-projection` argument are subject to change. (Probably also the default values of the `project` arguments will change.)

These functions build the arguments for `prop:contract`, `prop:chaperone-contract`, and `prop:flat-contract`, respectively.

A *contract property* specifies the behavior of a structure when used as a contract. It is specified in terms of five accessors: *get-name*, which produces a description to write as part of a contract violation; *get-first-order*, which produces a first-order predicate to be used by contract-first-order-passes?; *get-projection*, which produces a blame-tracking projection defining the behavior of the contract; *stronger*, which is a predicate that determines whether this contract (passed in the first argument) is stronger than some other contract (passed in the second argument); *generate*, which returns a thunk that generates random values matching the contract or #f, indicating that random generation for this contract isn't supported; *exercise*, which returns a function that exercises values matching the contract (e.g., if it is a function contract, it may call the function) and a list of contracts whose values will be generated by this process; and is-flat-contract?, which is used by flat-contract? to determine if this contract accepts only list?s.

These accessors are passed as (optional) keyword arguments to build-contract-property, and are applied to instances of the appropriate structure type by the contract system. Their results are used analogously to the arguments of make-contract.

A *chaperone contract property* specifies the behavior of a structure when used as a chaperone contract. It is specified using build-chaperone-contract-property, and accepts exactly the same set of arguments as build-contract-property. The only difference is that the projection accessor must return a value that passes chaperone-of? when compared with the original, uncontracted value.

A *flat contract property* specifies the behavior of a structure when used as a flat contract. It is specified using build-flat-contract-property, and accepts exactly the same set of arguments as build-contract-property. The only difference is that the projection accessor is expected not to wrap its argument in a higher-order fashion, analogous to the constraint on projections in make-flat-contract.

Changed in version 6.0.1.13 of package base: Added the #:list-contract? argument.

```
(contract-property? x) → boolean?
  x : any/c
(chaperone-contract-property? x) → boolean?
  x : any/c
(flat-contract-property? x) → boolean?
  x : any/c
```

These predicates detect whether a value is a contract property, chaperone contract property, or a flat contract property, respectively.

8.7.3 Obligation Information in Check Syntax

Check Syntax in DrRacket shows obligation information for contracts according to syntax-propertys that the contract combinators leave in the expanded form of the program. These

properties indicate where contracts appear in the source and where the positive and negative positions of the contracts appear.

To make Check Syntax show obligation information for your new contract combinators, use the following properties (some helper macros and functions are below):

- `'racket/contract:contract :`
  `(vector/c symbol? (listof syntax?) (listof syntax?))`

  This property should be attached to the result of a transformer that implements a contract combinator. It signals to Check Syntax that this is where a contract begins.

  The first element in the vector should be a unique (in the sense of `eq?`) value that Check Syntax can use a tag to match up this contract with its subpieces (specified by the two following syntax properties).

  The second and third elements of the vector are syntax objects from pieces of the contract, and Check Syntax will color them. The first list should contain subparts that are the responsibility of parties (typically modules) that provide implementations of the contract. The second list should contain subparts that are the responsibility of clients.

  For example, in `(->* () #:pre #t any/c #:post #t)`, the `->*` and the `#:post` should be in the first list and `#:pre` in the second list.

- `'racket/contract:negative-position : symbol?`

  This property should be attached to sub-expressions of a contract combinator that are expected to be other contracts. The value of the property should be the key (the first element from the vector for the `'racket/contract:contract` property) indicating which contract this is.

  This property should be used when the expression's value is a contract that clients are responsible for.

- `'racket/contract:positive-position : symbol?`

  This form is just like `'racket/contract:negative-position`, except that it should be used when the expression's value is a contract that the original party should be responsible for.

- `'racket/contract:contract-on-boundary : symbol?`

  The presence of this property tells Check Syntax that it should start coloring from this point. It expects the expression to be a contract (and, thus, to have the `'racket/contract:contract` property); this property indicates that this contract is on a (module) boundary.

  (The value of the property is not used.)

- `'racket/contract:internal-contract : symbol?`

  Like `'racket/contract:contract-on-boundary`, the presence of this property triggers coloring, but this is meant for use when the party (module) containing the

contract (regardless of whether or not this module exports anything matching the contract) can be blamed for violating the contract. This comes into play for `->i` contracts, since the contract itself has access to values under contract via the dependency.

```
(define/final-prop header body ...)

header = main-id
       | (main-id id ...)
       | (main-id id ... . id)
```

The same as `(define header body ...)`, except that uses of *main-id* in the header are annotated with the `'racket/contract:contract` property (as above).

```
(define/subexpression-pos-prop header body ...)

header = main-id
       | (main-id id ...)
       | (main-id id ... . id)
```

The same as `(define header body ...)`, except that uses of *main-id* in the header are annotated with the `'racket/contract:contract` property (as above) and arguments are annotated with the `'racket/contract:positive-position` property.

8.7.4 Utilities for Building New Combinators

```
(contract-stronger? x y) → boolean?
  x : contract?
  y : contract?
```

Returns `#t` if the contract *x* accepts either fewer or the same number of values as *y* does.

This function is conservative, so it may return `#f` when *x* does, in fact, accept fewer values.

Examples:

```
> (contract-stronger? integer? integer?)
#t
> (contract-stronger? (between/c 25 75) (between/c 0 100))
#t
> (contract-stronger? (between/c 0 100) (between/c 25 75))
#f
> (contract-stronger? (between/c -10 0) (between/c 0 10))
#f
```

```
> (contract-stronger? (λ (x) (and (real? x) (<= x (random 10))))
                      (λ (x) (and (real? x) (<= x (+ 100 (random 10))))))
#f
```

```
(contract-first-order-passes? contract v) → boolean?
  contract : contract?
  v : any/c
```

Returns a boolean indicating whether the first-order tests of *contract* pass for *v*.

If it returns #f, the contract is guaranteed not to hold for that value; if it returns #t, the contract may or may not hold. If the contract is a first-order contract, a result of #t guarantees that the contract holds.

```
(contract-first-order c) → (-> any/c boolean?)
  c : contract?
```

Produces the first-order test used by or/c to match values to higher-order contracts.

8.8 Contract Utilities

```
(contract? v) → boolean?
  v : any/c
```

Returns #t if its argument is a contract (i.e., constructed with one of the combinators described in this section or a value that can be used as a contract) and #f otherwise.

```
(chaperone-contract? v) → boolean?
  v : any/c
```

Returns #t if its argument is a contract that guarantees that it returns a value which passes chaperone-of? when compared to the original, uncontracted value.

```
(impersonator-contract? v) → boolean?
  v : any/c
```

Returns #t if its argument is a contract that is not a chaperone contract nor a flat contract.

```
(flat-contract? v) → boolean?
  v : any/c
```

Returns #t when its argument is a contract that can be checked immediately (unlike, say, a function contract).

For example, flat-contract constructs flat contracts from predicates, and symbols, booleans, numbers, and other ordinary Racket values (that are defined as contracts) are also flat contracts.

```
(list-contract? v) → boolean?
  v : any/c
```

Recognizes certain `contract?` values that accept `list?`s.

A list contract is one that insists that its argument is a `list?`, meaning that the value cannot be cyclic and must either be the empty list or a pair constructed with `cons` and another list.

Added in version 6.0.1.13 of package `base`.

```
(contract-name c) → any/c
  c : contract?
```

Produces the name used to describe the contract in error messages.

```
(value-contract v) → contract?
  v : has-contract?
```

Returns the contract attached to `v`, if recorded. Otherwise it returns `#f`.

To support `value-contract` and `value-contract` in your own contract combinators, use `prop:contracted` or `impersonator-prop:contracted`.

```
(has-contract? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a value that has a recorded contract attached to it.

```
(value-blame v) → (or/c blame? #f)
  v : has-blame?
```

Returns the blame object for the contract attached to `v`, if recorded. Otherwise it returns `#f`.

To support `value-contract` and `value-blame` in your own contract combinators, use `prop:blame` or `impersonator-prop:blame`.

Added in version 6.0.1.12 of package `base`.

```
(has-blame? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a value that has a contract with blame information attached to it.

Added in version 6.0.1.12 of package `base`.

```
(contract-projection c) → (-> blame? (-> any/c any/c))
  c : contract?
```

Produces the projection defining a contract's behavior on protected values.

```
(make-none/c sexp-name) → contract?
  sexp-name : any/c
```

Makes a contract that accepts no values, and reports the name *sexp-name* when signaling a contract violation.

```
(recursive-contract contract-expr)
(recursive-contract contract-expr #:list-contract?)
(recursive-contract contract-expr type)
(recursive-contract contract-expr type #:list-contract?)
```

Delays the evaluation of its argument until the contract is checked, making recursive contracts possible. If *type* is given, it describes the expected type of contract and must be one of the keywords `#:impersonator`, `#:chaperone`, or `#:flat`. If *type* is not given, an impersonator contract is created.

If `#:list-contract?` is returned, then the result is a `list-contract?` and the *contract-expr* must evaluate to a `list-contract?`.

Changed in version 6.0.1.13 of package `base`: Added the `#:list-contract?` argument.

```
(opt/c contract-expr maybe-name)

maybe-name =
           | #:error-name id
```

This optimizes its argument contract expression by traversing its syntax and, for known contract combinators, fuses them into a single contract combinator that avoids as much allocation overhead as possible. The result is a contract that should behave identically to its argument, except faster.

If the `#:error-name` argument is present, and *contract-expr* evaluates to a non-contract expression, then `opt/c` raises an error using *id* as the name of the primitive, instead of using the name `opt/c`.

Examples:

```
> (define/contract (f x)
    (opt/c '(not-a-contract))
    x)
```

```
opt/c: contract violation
  expected: contract?
  given: '(not-a-contract)
> (define/contract (f x)
    (opt/c '(not-a-contract) #:error-name define/contract)
    x)
define/contract: contract violation
  expected: contract?
  given: '(not-a-contract)
```

```
(define-opt/c (id id ...) expr)
```

This defines a recursive contract and simultaneously optimizes it. Semantically, it behaves just as if the `-opt/c` were not present, defining a function on contracts (except that the body expression must return a contract). But, it also optimizes that contract definition, avoiding extra allocation, much like `opt/c` does.

For example,

```
(define-contract-struct bt (val left right))

(define-opt/c (bst-between/c lo hi)
  (or/c null?
        (bt/c [val (real-in lo hi)]
              [left (val) (bst-between/c lo val)]
              [right (val) (bst-between/c val hi)])))

(define bst/c (bst-between/c -inf.0 +inf.0))
```

defines the `bst/c` contract that checks the binary search tree invariant. Removing the `-opt/c` also makes a binary search tree contract, but one that is (approximately) 20 times slower.

```
contract-continuation-mark-key : continuation-mark-key?
```

Key used by continuation marks that are present during contract checking. The value of these marks are the blame objects that correspond to the contract currently being checked.

8.9 racket/contract/base

```
(require racket/contract/base)
```
package: `base`

The `racket/contract/base` module provides a subset of the exports of `racket/contract` module. In particular, it contains everything in the

- §8.1 “Data-structure Contracts”
- §8.2 “Function Contracts”
- §8.6 “Attaching Contracts to Values” and
- §8.8 “Contract Utilities” sections.

Unfortunately, using `racket/contract/base` does not yield a significantly smaller memory footprint than `racket/contract`, but it can still be useful to add contracts to libraries that `racket/contract` uses to implement some of the more sophisticated parts of the contract system.

8.10 Legacy Contracts

```
(make-proj-contract name proj first-order) → contract?
  name : any/c
         (or/c (-> any/c
                   any/c
                   (list/c any/c any/c)
                   contact?
                   (-> any/c any/c))
  proj :       (-> any/c
                   any/c
                   (list/c any/c any/c)
                   contact?
                   boolean?
                   (-> any/c any/c)))
  first-order : (-> any/c boolean?)
```

Builds a contract using an old interface.

Modulo errors, it is equivalent to:

```
(make-contract
 #:name name
 #:first-order first-order
 #:projection
 (cond
   [(procedure-arity-includes? proj 5)
    (lambda (blame)
      (proj (blame-positive blame)
            (blame-negative blame)
            (list (blame-source blame) (blame-value blame))
```

```
            (blame-contract blame)
            (not (blame-swapped? blame)))))]
    [(procedure-arity-includes? proj 4)
     (lambda (blame)
       (proj (blame-positive blame)
             (blame-negative blame)
             (list (blame-source blame) (blame-value blame))
             (blame-contract blame)))]))
```

```
(raise-contract-error val
                      src
                      pos
                      name
                      fmt
                      arg ...) → any/c
  val : any/c
  src : any/c
  pos : any/c
  name : any/c
  fmt : string?
  arg : any/c
```

Calls `raise-blame-error` after building a `blame` struct from the *val*, *src*, *pos*, and *name* arguments. The *fmt* string and following arguments are passed to `format` and used as the string in the error message.

```
(contract-proc c)
    (->* (symbol? symbol? (or/c syntax? (list/c any/c any/c)))
 →       (boolean?)
         (-> any/c any))
  c : contract?
```

Constructs an old-style projection from a contract.

The resulting function accepts the information that is in a `blame` struct and returns a projection function that checks the contract.

8.11 Random generation

```
(contract-random-generate ctc [fuel fail]) → any/c
  ctc : contract?
  fuel : 5 = exact-nonnegative-integer?
  fail : (or/c #f (-> any)) = #f
```

Attempts to randomly generate a value which will match the contract. The fuel argument limits how hard the generator tries to generate a value matching the contract and is a rough limit of the size of the resulting value.

The generator may fail to generate a value, either because some contracts do not have corresponding generators (for example, not all predicates have generators) or because there is not enough fuel. In either case, the thunk *fail* is invoked.

```
(contract-exercise val ...+) → void?
  val : any/c
```

Attempts to get the *val*s to break their contracts (if any).

Uses value-contract to determine if any of the *val*s have a contract and, for those that do, uses information about the contract's shape to poke and prod at the value. For example, if the value is function, it will use the contract to tell it what arguments to supply to the value.

9 Pattern Matching

§12 “Pattern Matching” in ***The Racket Guide*** introduces pattern matching.

The `match` form and related forms support general pattern matching on Racket values. See also §4.7 “Regular Expressions” for information on regular-expression matching on strings, bytes, and streams.

```
(require racket/match)
```
package: `base`

The bindings documented in this section are provided by the `racket/match` and `racket` libraries, but not `racket/base`.

```
(match val-expr clause ...)

clause = [pat body ...+]
       | [pat (=> id) body ...+]
       | [pat #:when cond-expr body ...+]
```

Finds the first *pat* that matches the result of *val-expr*, and evaluates the corresponding *body*s with bindings introduced by *pat* (if any). The last *body* in the matching clause is evaluated in tail position with respect to the `match` expression.

To find a match, the *clause*s are tried in order. If no *clause* matches, then the `exn:misc:match?` exception is raised.

An optional `#:when` *cond-expr* specifies that the pattern should only match if *cond-expr* produces a true value. *cond-expr* is in the scope of all of the variables bound in *pat*. *cond-expr* must not mutate the object being matched before calling the failure procedure, otherwise the behavior of matching is unpredictable. See also `failure-cont`, which is a lower-level mechanism achieving the same ends.

An optional `(=> id)` between a *pat* and the *body*s is bound to a *failure procedure* of zero arguments. If this procedure is invoked, it escapes back to the pattern matching expression, and resumes the matching process as if the pattern had failed to match. The *body*s must not mutate the object being matched before calling the failure procedure, otherwise the behavior of matching is unpredictable.

The grammar of *pat* is as follows, where non-italicized identifiers are recognized symbolically (i.e., not by binding).

pat	`::= id`	match anything, bind identifier
	`| (var id)`	match anything, bind identifier
	`| _`	match anything
	`| literal`	match literal
	`| (quote datum)`	match `equal?` value
	`| (list lvp ...)`	match sequence of *lvp*s
	`| (list-rest lvp ... pat)`	match *lvp*s consed onto a *pat*
	`| (list-no-order pat ...)`	match *pat*s in any order

	| `(list-no-order pat ... lvp)`	match `pat`s in any order
	| `(vector lvp ...)`	match vector of `pat`s
	| `(hash-table (pat pat) ...)`	match hash table
	| `(hash-table (pat pat) ...+ ooo)`	match hash table
	| `(cons pat pat)`	match pair of `pat`s
	| `(mcons pat pat)`	match mutable pair of `pat`s
	| `(box pat)`	match boxed `pat`
	| `(struct-id pat ...)`	match `struct-id` instance
	| `(struct struct-id (pat ...))`	match `struct-id` instance
	| `(regexp rx-expr)`	match string
	| `(regexp rx-expr pat)`	match string, result with `pat`
	| `(pregexp px-expr)`	match string
	| `(pregexp px-expr pat)`	match string, result with `pat`
	| `(and pat ...)`	match when all `pat`s match
	| `(or pat ...)`	match when any `pat` match
	| `(not pat ...)`	match when no `pat` matches
	| `(app expr pats ...)`	match `(expr` value`)` output values to `pat`s
	| `(? expr pat ...)`	match if `(expr` value`)` and `pat`s
	| `(quasiquote qp)`	match a quasipattern
	| `derived-pattern`	match using extension
`literal`	`::= #t`	match true
	| `#f`	match false
	| `string`	match `equal?` string
	| `bytes`	match `equal?` byte string
	| `number`	match `equal?` number
	| `char`	match `equal?` character
	| `keyword`	match `equal?` keyword
	| `regexp`	match `equal?` regexp literal
	| `pregexp`	match `equal?` pregexp literal
`lvp`	`::= pat ooo`	greedily match `pat` instances
	| `pat`	match `pat`
`qp`	`::= literal`	match literal
	| `id`	match symbol
	| `(qp ...)`	match sequences of `qp`s
	| `(qp ... . qp)`	match `qp`s ending `qp`
	| `(qp ooo . qp)`	match `qp`s beginning with repeated `qp`
	| `#(qp ...)`	match vector of `qp`s
	| `#&qp`	match boxed `qp`
	| `,pat`	match `pat`
	| `,@(list lvp ...)`	match `lvp`s, spliced
	| `,@(list-rest lvp ... pat)`	match `lvp`s plus `pat`, spliced
	| `,@'qp`	match list-matching `qp`, spliced
`ooo`	`::= ...`	zero or more; `...` is literal
	| `___`	zero or more
	| `..k`	`k` or more
	| `__k`	`k` or more

In more detail, patterns match as follows:

- *id* (excluding the reserved names _, ..., .._, ..*k*, and ..*k* for non-negative integers *k*) or (var *id*) — matches anything, and binds *id* to the matching values. If an *id* is used multiple times within a pattern, the corresponding matches must be the same according to (match-equality-test), except that instances of an *id* in different or and not sub-patterns are independent.

  Examples:

```
> (match '(1 2 3)
    [(list a b a) (list a b)]
    [(list a b c) (list c b a)])
'(3 2 1)
> (match '(1 (x y z) 1)
    [(list a b a) (list a b)]
    [(list a b c) (list c b a)])
'(1 (x y z))
```

- _ — matches anything, without binding any identifiers.

  Example:

```
> (match '(1 2 3)
    [(list _ _ a) a])
3
```

- #t, #f, *string*, *bytes*, *number*, *char*, or (quote *datum*) — matches an equal? constant.

  Example:

```
> (match "yes"
    ["no" #f]
    ["yes" #t])
#t
```

- (list *lvp* ...) — matches a list of elements. In the case of (list *pat* ...), the pattern matches a list with as many element as *pat*s, and each element must match the corresponding *pat*. In the more general case, each *lvp* corresponds to a "spliced" list of greedy matches.

  For spliced lists, ... and ___ are aliases for zero or more matches. The ..*k* and __*k* forms are also aliases, specifying *k* or more matches. Pattern variables that precede these splicing operators are bound to lists of matching forms.

  Examples:

```
> (match '(1 2 3)
    [(list a b c) (list c b a)])
'(3 2 1)
```

```
> (match '(1 2 3)
    [(list 1 a ...) a])
'(2 3)
> (match '(1 2 3)
    [(list 1 a ..3) a]
    [_ 'else])
'else
> (match '(1 2 3 4)
    [(list 1 a ..3) a]
    [_ 'else])
'(2 3 4)
> (match '(1 2 3 4 5)
    [(list 1 a ..3 5) a]
    [_ 'else])
'(2 3 4)
> (match '(1 (2) (2) (2) 5)
    [(list 1 (list a) ..3 5) a]
    [_ 'else])
'(2 2 2)
```

- `(list-rest` *lvp* `...` *pat*`)` — similar to a `list` pattern, but the final *pat* matches the "rest" of the list after the last *lvp*. In fact, the matched value can be a non-list chain of pairs (i.e., an "improper list") if *pat* matches non-list values.

  Examples:

```
> (match '(1 2 3 . 4)
    [(list-rest a b c d) d])
4
> (match '(1 2 3 . 4)
    [(list-rest a ... d) (list a d)])
'((1 2 3) 4)
```

- `(list-no-order` *pat* `...)` — similar to a `list` pattern, but the elements to match each *pat* can appear in the list in any order.

  Example:

```
> (match '(1 2 3)
    [(list-no-order 3 2 x) x])
1
```

- `(list-no-order` *pat* `...` *lvp*`)` — generalizes `list-no-order` to allow a pattern that matches multiple list elements that are interspersed in any order with matches for the other patterns.

  Example:

```
> (match '(1 2 3 4 5 6)
    [(list-no-order 6 2 y ...) y])
'(1 3 4 5)
```

- `(vector lvp ...)` — like a `list` pattern, but matching a vector.

  Example:

```
> (match #(1 (2) (2) (2) 5)
    [(vector 1 (list a) ..3 5) a])
'(2 2 2)
```

- `(hash-table (pat pat) ...)` — similar to `list-no-order`, but matching against hash table's key–value pairs.

  Example:

```
> (match #hash(("a" . 1) ("b" . 2))
    [(hash-table ("b" b) ("a" a)) (list b a)])
'(2 1)
```

- `(hash-table (pat pat) ...+ ooo)` — Generalizes `hash-table` to support a final repeating pattern.

  Example:

```
> (match #hash(("a" . 1) ("b" . 2))
    [(hash-table (key val) ...) key])
'("a" "b")
```

- `(cons pat1 pat2)` — matches a pair value.

  Example:

```
> (match (cons 1 2)
    [(cons a b) (+ a b)])
3
```

- `(mcons pat1 pat2)` — matches a mutable pair value.

  Example:

```
> (match (mcons 1 2)
    [(cons a b) 'immutable]
    [(mcons a b) 'mutable])
'mutable
```

- `(box pat)` — matches a boxed value.

  Example:

```
> (match #&1
    [(box a) a])
1
```

- `(`*`struct-id pat ...`*`)` or `(struct` *`struct-id`* `(`*`pat ...`*`))` — matches an instance of a structure type named *`struct-id`*, where each field in the instance matches the corresponding *`pat`*. See also `struct*`.

  Usually, *`struct-id`* is defined with `struct`. More generally, *`struct-id`* must be bound to expansion-time information for a structure type (see §5.7 “Structure Type Transformer Binding”), where the information includes at least a predicate binding and field accessor bindings corresponding to the number of field *`pat`*s. In particular, a module import or a `unit` import with a signature containing a `struct` declaration can provide the structure type information.

  Examples:

```
(define-struct tree (val left right))

> (match (make-tree 0 (make-tree 1 #f #f) #f)
    [(tree a (tree b  _ _) _) (list a b)])
'(0 1)
```

- `(struct` *`struct-id`* `_)` — matches any instance of *`struct-id`*, without regard to contents of the fields of the instance.

- `(regexp` *`rx-expr`*`)` — matches a string that matches the regexp pattern produced by *`rx-expr`*; see §4.7 “Regular Expressions” for more information about regexps.

  Examples:

```
> (match "apple"
    [(regexp #rx"p+") 'yes]
    [_ 'no])
'yes
> (match "banana"
    [(regexp #rx"p+") 'yes]
    [_ 'no])
'no
```

- `(regexp` *`rx-expr pat`*`)` — extends the `regexp` form to further constrain the match where the result of `regexp-match` is matched against *`pat`*.

  Examples:

```
> (match "apple"
    [(regexp #rx"p+(.)" (list _ "l")) 'yes]
    [_ 'no])
'yes
> (match "append"
    [(regexp #rx"p+(.)" (list _ "l")) 'yes]
    [_ 'no])
'no
```

- `(pregexp rx-expr)` or `(regexp rx-expr pat)` — like the `regexp` patterns, but if *rx-expr* produces a string, it is converted to a pattern using `pregexp` instead of `regexp`.

- `(and pat ...)` — matches if all of the *pat*s match. This pattern is often used as `(and id pat)` to bind *id* to the entire value that matches *pat*.

  Example:

  ```
  > (match '(1 (2 3) 4)
     [(list _ (and a (list _ ...)) _) a])
  '(2 3)
  ```

- `(or pat ...)` — matches if any of the *pat*s match. **Beware**: the result expression can be duplicated once for each *pat*! Identifiers in *pat* are bound only in the corresponding copy of the result expression; in a module context, if the result expression refers to a binding, then all *pat*s must include the binding.

  Example:

  ```
  > (match '(1 2)
     [(or (list a 1) (list a 2)) a])
  1
  ```

- `(not pat ...)` — matches when none of the *pat*s match, and binds no identifiers.

  Examples:

  ```
  > (match '(1 2 3)
     [(list (not 4) ...) 'yes]
     [_ 'no])
  'yes
  > (match '(1 4 3)
     [(list (not 4) ...) 'yes]
     [_ 'no])
  'no
  ```

- `(app expr pats ...)` — applies *expr* to the value to be matched; the result of the application is matched against *pats*.

  Examples:

  ```
  > (match '(1 2)
     [(app length 2) 'yes])
  'yes
  > (match '(1 2)
     [(app (lambda (v) (split-at v 1)) '(1) '(2)) 'yes])
  'yes
  ```

- `(? expr pat ...)` — applies *expr* to the value to be matched, and checks whether the result is a true value; the additional *pat*s must also match; i.e., `?` combines a predicate application and an `and` pattern. However, `?`, unlike `and`, guarantees that *expr* is matched before any of the *pat*s.

Example:

```
> (match '(1 3 5)
    [(list (? odd?) ...) 'yes])
'yes
```

- `(quasiquote qp)` — introduces a quasipattern, in which identifiers match symbols. Like the `quasiquote` expression form, `unquote` and `unquote-splicing` escape back to normal patterns.

  Example:

  ```
  > (match '(1 2 3)
      [`(1 ,a ,(? odd? b)) (list a b)])
  '(2 3)
  ```

- *derived-pattern* — matches a pattern defined by a macro extension via `define-match-expander`.

Note that the matching process may destructure the input multiple times, and may evaluate expressions embedded in patterns such as `(app expr pat)` in arbitrary order, or multiple times. Therefore, such expressions must be safe to call multiple times, or in an order other than they appear in the original program.

9.1 Additional Matching Forms

```
(match* (val-expr ...+) clause* ...)

clause* = [(pat ...+) body ...+]
        | [(pat ...+) (=> id) body ...+]
        | [(pat ...+) #:when cond-expr body ...+]
```

Matches a sequence of values against each clause in order, matching only when all patterns in a clause match. Each clause must have the same number of patterns as the number of *val-expr*s.

Examples:

```
> (match* (1 2 3)
   [(_ (? number?) x) (add1 x)])
4
> (match* (15 17)
   [((? number? a) (? number? b))
    #:when (= (+ a 2) b)
    'diff-by-two])
```

```
'diff-by-two
```

```
(match/values expr clause clause ...)
```

If *expr* evaluates to `n` values, then match all `n` values against the patterns in *clause* .... Each clause must contain exactly `n` patterns. At least one clause is required to determine how many values to expect from *expr*.

```
(define/match (head args)
  match*-clause ...)

        head = id
             | (head args)

        args = arg ...
             | arg ... . rest-id

         arg = arg-id
             | [arg-id default-expr]
             | keyword arg-id
             | keyword [arg-id default-expr]

match*-clause = [(pat ...+) body ...+]
              | [(pat ...+) (=> id) body ...+]
              | [(pat ...+) #:when cond-expr body ...+]
```

Binds *id* to a procedure that is defined by pattern matching clauses using `match*`. Each clause takes a sequence of patterns that correspond to the arguments in the function header. The arguments are ordered as they appear in the function header for matching purposes.

Examples:

```
> (define/match (fact n)
    [(0) 1]
    [(n) (* n (fact (sub1 n)))])

> (fact 5)
120
```

The function header may also contain optional or keyword arguments, may have curried arguments, and may also contain a rest argument.

Examples:

```
> (define/match ((f x) #:y [y '(1 2 3)])
    [((regexp #rx"p+") `(,a 2 3)) a]
    [(_ _) #f])
```

```
> ((f "ape") #:y '(5 2 3))
5
> ((f "dog"))
#f
> (define/match (g x y . rst)
    [(0 0 '()) #t]
    [(5 5 '(5 5)) #t]
    [(_ _ _) #f])

> (g 0 0)
#t
> (g 5 5 5 5)
#t
> (g 1 2)
#f
```

(match-lambda *clause* ...)

Equivalent to (lambda (id) (match id *clause* ...)).

(match-lambda* *clause* ...)

Equivalent to (lambda lst (match lst *clause* ...)).

(match-lambda** *clause** ...)

Equivalent to (lambda (args ...) (match* (args ...) *clause** ...)), where the number of args ... is computed from the number of patterns appearing in each of the *clause**.

(match-let ([*pat expr*] ...) *body* ...+)

Generalizes let to support pattern bindings. Each *expr* is matched against its corresponding *pat* (the match must succeed), and the bindings that *pat* introduces are visible in the *body*s.

Example:

```
> (match-let ([(list a b) '(1 2)]
              [(vector x ...) #(1 2 3 4)])
    (list b a x))
'(2 1 (1 2 3 4))
```

(match-let* ([*pat expr*] ...) *body* ...+)

Like `match-let`, but generalizes `let*`, so that the bindings of each *pat* are available in each subsequent *expr*.

Example:

```
> (match-let* ([(list a b) '(#(1 2 3 4) 2)]
                [(vector x ...) a])
    x)
'(1 2 3 4)
```

```
(match-let-values ([(pat ...) expr] ...) body ...+)
```

Like `match-let`, but generalizes `let-values`.

```
(match-let*-values ([(pat ...) expr] ...) body ...+)
```

Like `match-let*`, but generalizes `let*-values`.

```
(match-letrec ([pat expr] ...) body ...+)
```

Like `match-let`, but generalizes `letrec`.

```
(match-define pat expr)
```

Defines the names bound by *pat* to the values produced by matching against the result of *expr*.

Examples:

```
> (match-define (list a b) '(1 2))

> b
2
```

```
(match-define-values (pat pats ...) expr)
```

Like `match-define` but for when expr produces multiple values. Like match/values, it requires at least one pattern to determine the number of values to expect.

Examples:

```
> (match-define-values (a b) (values 1 2))

> b
2
```

```
(exn:misc:match? v) → boolean?
  v : any/c
```

A predicate for the exception raised in the case of a match failure.

```
(failure-cont)
```

Continues matching as if the current pattern failed. Note that unlike use of the => form, this does *not* escape the current context, and thus should only be used in tail position with respect to the `match` form.

9.2 Extending `match`

```
(define-match-expander id proc-expr)
(define-match-expander id proc-expr proc-expr)
```

Binds *id* to a *match expander*.

The first *proc-expr* sub-expression must evaluate to a transformer that produces a *pat* for `match`. Whenever *id* appears as the beginning of a pattern, this transformer is given, at expansion time, a syntax object corresponding to the entire pattern (including *id*). The pattern is replaced with the result of the transformer.

A transformer produced by a second *proc-expr* sub-expression is used when *id* is used in an expression context. Using the second *proc-expr*, *id* can be given meaning both inside and outside patterns.

Match expanders are not invoked unless *id* appears in the first position in a sequence. Instead, identifiers bound by `define-match-expander` are used as binding identifiers (like any other identifier) when they appear anywhere except the first position in a sequence.

For example, to extend the pattern matcher and destructure syntax lists,

```
(define (syntax-list? x)
  (and (syntax? x)
       (list? (syntax->list x))))
(define-match-expander syntax-list
  (lambda (stx)
    (syntax-case stx ()
      [(_ elts ...)
       #'(? syntax-list?
            (app syntax->list (list elts ...)))])))
```

```
(define (make-keyword-predicate keyword)
  (lambda (stx)
    (and (identifier? stx)
         (free-identifier=? stx keyword))))
(define or-keyword? (make-keyword-predicate #'or))
(define and-keyword? (make-keyword-predicate #'and))

> (match #'(or 3 4)
    [(syntax-list (? or-keyword?) b c)
     (list "OOORRR!" b c)]
    [(syntax-list (? and-keyword?) b c)
     (list "AAANND!" b c)])
'("OOORRR!" #<syntax:59:0 3> #<syntax:59:0 4>)
> (match #'(and 5 6)
    [(syntax-list (? or-keyword?) b c)
     (list "OOORRR!" b c)]
    [(syntax-list (? and-keyword?) b c)
     (list "AAANND!" b c)])
'("AAANND!" #<syntax:60:0 5> #<syntax:60:0 6>)
```

And here is an example showing how `define-match-expander`-bound identifiers are not treated specially unless they appear in the first position of pattern sequence.

```
(define-match-expander nil
  (λ (stx) #''())
  (λ (stx) #''()))
(define (len l)
  (match l
    [nil 0]
    [(cons hd tl) (+ 1 (len tl))]))

> (len nil)
0
> (len (cons 1 nil))
0
> (len (cons 1 (cons 2 nil)))
0
```

`prop:match-expander` : `struct-type-property?`

A structure type property to identify structure types that act as match expanders like the ones created by `define-match-expander`.

The property value must be an exact non-negative integer or a procedure of one or two arguments. In the former case, the integer designates a field within the structure that should con-

tain a procedure; the integer must be between 0 (inclusive) and the number of non-automatic fields in the structure type (exclusive, not counting supertype fields), and the designated field must also be specified as immutable.

If the property value is a procedure of one argument, then the procedure serves as the transformer for match expansion. If the property value is a procedure of two arguments, then the first argument is the structure whose type has `prop:match-expander` property, and the second argument is a syntax object as for a match expander..

If the property value is a assignment transformer, then the wrapped procedure is extracted with `set!-transformer-procedure` before it is called.

This binding is provided `for-syntax`.

```
prop:legacy-match-expander : struct-type-property?
```

Like `prop:match-expander`, but for the legacy match syntax.

This binding is provided `for-syntax`.

```
(match-expander? v) → boolean?
  v : any/c
(legacy-match-expander? v) → boolean?
  v : any/c
```

Predicates for values which implement the appropriate match expander properties.

```
(syntax-local-match-introduce stx) → syntax?
  stx : syntax?
```

Like `syntax-local-introduce`, but for match expanders.

```
(match-equality-test) → (any/c any/c . -> . any)
(match-equality-test comp-proc) → void?
  comp-proc : (any/c any/c . -> . any)
```

A parameter that determines the comparison procedure used to check whether multiple uses of an identifier match the “same” value. The default is `equal?`.

```
(match/derived val-expr original-datum clause ...)
(match*/derived (val-expr ...) original-datum clause* ...)
```

Like `match` and `match*` respectively, but includes a sub-expression to be used as the source for all syntax errors within the form. For example, `match-lambda` expands to `match/derived` so that errors in the body of the form are reported in terms of `match-lambda` instead of `match`.

9.3 Library Extensions

```
(== val comparator)
(== val)
```

A match expander which checks if the matched value is the same as *val* when compared by *comparator*. If *comparator* is not provided, it defaults to equal?.

Examples:

```
> (match (list 1 2 3)
    [(== (list 1 2 3)) 'yes]
    [_ 'no])
'yes
> (match (list 1 2 3)
    [(== (list 1 2 3) eq?) 'yes]
    [_ 'no])
'no
> (match (list 1 2 3)
    [(list 1 2 (== 3 =)) 'yes]
    [_ 'no])
'yes
```

```
(struct* struct-id ([field pat] ...))
```

A match pattern form that matches an instance of a structure type named *struct-id*, where the field *field* in the instance matches the corresponding *pat*.

Any field of *struct-id* may be omitted, and such fields can occur in any order.

Examples:

```
(define-struct tree (val left right))

> (match (make-tree 0 (make-tree 1 #f #f) #f)
    [(struct* tree ([val a]
                    [left (struct* tree ([right #f] [val b]))]))
     (list a b)])
'(0 1)
```

10 Control Flow

10.1 Multiple Values

See §1.1.3 “Multiple Return Values” for general information about multiple result values. In addition to `call-with-values` (described in this section), the `let-values`, `let*-values`, `letrec-values`, and `define-values` forms (among others) create continuations that receive multiple values.

```
(values v ...) → any
  v : any/c
```

Returns the given `v`s. That is, `values` returns its provided arguments.

Examples:

```
> (values 1)
1
> (values 1 2 3)
1
2
3
> (values)
```

```
(call-with-values generator receiver) → any
  generator : (-> any)
  receiver : procedure?
```

Calls *generator*, and passes the values that *generator* produces as arguments to *receiver*. Thus, `call-with-values` creates a continuation that accepts any number of values that *receiver* can accept. The *receiver* procedure is called in tail position with respect to the `call-with-values` call.

Examples:

```
> (call-with-values (lambda () (values 1 2)) +)
3
> (call-with-values (lambda () 1) (lambda (x y) (+ x y)))
#<procedure>: arity mismatch;
 the expected number of arguments does not match the given
number
  expected: 2
  given: 1
  arguments...:
   1
```

10.2 Exceptions

§10.1 “Exceptions” in ***The Racket Guide*** introduces exceptions.

See §1.1.15 “Exceptions” for information on the Racket exception model. It is based on a proposal by Friedman, Haynes, and Dybvig [Friedman95].

Whenever a primitive error occurs in Racket, an exception is raised. The value that is passed to the current exception handler for a primitive error is always an instance of the `exn` structure type. Every `exn` structure value has a `message` field that is a string, the primitive error message. The default exception handler recognizes exception values with the `exn?` predicate and passes the error message to the current error display handler (see `error-display-handler`).

Primitive procedures that accept a procedure argument with a particular required arity (e.g., `call-with-input-file`, `call/cc`) check the argument's arity immediately, raising `exn:fail:contract` if the arity is incorrect.

10.2.1 Error Message Conventions

Racket's *error message convention* is to produce error messages with the following shape:

⟨*srcloc*⟩`:` ⟨*name*⟩`:` ⟨*message*⟩`;`
⟨*continued-message*⟩ ...
⟨*field*⟩`:` ⟨*detail*⟩
...

The message starts with an optional source location, ⟨*srcloc*⟩, which is followed by a colon and space when present. The message continues with an optional ⟨*name*⟩ that usually identifies the complaining function, syntactic form, or other entity, but may also refer to an entity being complained about; the ⟨*name*⟩ is also followed by a colon and space when present.

The ⟨*message*⟩ should be relatively short, and it should be largely independent of specific values that triggered the error. More detailed explanation that requires multiple lines should continue with each line indented by a single space, in which case ⟨*message*⟩ should end in a semi-colon (but the semi-colon should be omitted if ⟨*continued-message*⟩ is not present). Message text should be lowercase—using semi-colons to separate sentences if needed, although long explanations may be better deferred to extra fields.

Specific values that triggered the error or other helpful information should appear in separate ⟨*field*⟩ lines, each of which is indented by two spaces. If a ⟨*detail*⟩ is especially long or takes multiple lines, it should start on its own line after the ⟨*field*⟩ label, and each of its lines should be indented by three spaces. Field names should be all lowercase.

A ⟨*field*⟩ name should end with `...` if the field provides relatively detailed information that might be distracting in common cases but useful in others. For example, when a contract failure is reported for a particular argument of a function, other arguments to the function

might be shown in an “other arguments...” field. The intent is that fields whose names end in ... might be hidden by default in an environment such as DrRacket.

Make ⟨*field*⟩ names as short as possible, relying on ⟨*message*⟩ or ⟨*continued message*⟩ text to clarify the meaning for a field. For example, prefer “given” to “given turtle” as a field name, where ⟨*message*⟩ is something like “given turtle is too sleepy” to clarify that “given” refers to a turtle.

10.2.2 Raising Exceptions

```
(raise v [barrier?]) → any
  v : any/c
  barrier? : any/c = #t
```

Raises an exception, where *v* represents the exception being raised. The *v* argument can be anything; it is passed to the current exception handler.

If *barrier?* is true, then the call to the exception handler is protected by a continuation barrier, so that multiple returns/escapes are impossible. All exceptions raised by `racket` functions effectively use `raise` with a `#t` value for *barrier?*.

Breaks are disabled from the time the exception is raised until the exception handler obtains control, and the handler itself is `parameterize-breaked` to disable breaks initially; see §10.6 “Breaks” for more information on breaks.

Examples:

```
> (with-handlers ([number? (lambda (n)
                             (+ n 5))])
    (raise 18 #t))
23
> (define-struct (my-exception exn:fail:user) ())

> (with-handlers ([my-exception? (lambda (e)
                                   #f)])
    (+ 5 (raise (make-my-exception
                 "failed"
                 (current-continuation-marks)))))
#f
> (raise 'failed #t)
uncaught exception: failed
```

```
(error sym) → any
  sym : symbol?
(error msg v ...) → any
```

```
  msg : string?
  v : any/c
(error src format v ...) → any
  src : symbol?
  format : string?
  v : any/c
```

Raises the exception `exn:fail`, which contains an error string. The different forms produce the error string in different ways:

- `(error sym)` creates a message string by concatenating `"error: "` with the string form of *sym*. Use this form sparingly.
- `(error msg v ...)` creates a message string by concatenating *msg* with string versions of the *v*s (as produced by the current error value conversion handler; see `error-value->string-handler`). A space is inserted before each *v*. Use this form sparingly, because it does not conform well to Racket's error message conventions; consider `raise-arguments-error`, instead.
- `(error src frmat v ...)` creates a message string equivalent to the string created by

  ```
  (format (string-append "~s: " frmat) src v ...)
  ```

  When possible, use functions such as `raise-argument-error`, instead, which construct messages that follow Racket's error message conventions.

In all cases, the constructed message string is passed to `make-exn:fail`, and the resulting exception is raised.

Examples:

```
> (error 'failed)
error: failed
> (error "failed" 23 'pizza (list 1 2 3))
failed 23 pizza (1 2 3)
> (error 'method-a "failed because ~a" "no argument supplied")
method-a: failed because no argument supplied
```

```
(raise-user-error sym) → any
  sym : symbol?
(raise-user-error msg v ...) → any
  msg : string?
  v : any/c
(raise-user-error src format v ...) → any
  src : symbol?
  format : string?
  v : any/c
```

Like error, but constructs an exception with make-exn:fail:user instead of make-exn:fail. The default error display handler does not show a "stack trace" for exn:fail:user exceptions (see §10.5 "Continuation Marks"), so raise-user-error should be used for errors that are intended for end users.

```
(raise-argument-error name expected v) → any
  name : symbol?
  expected : string?
  v : any/c
(raise-argument-error name
                      expected
                      bad-pos
                      v ...)   → any
  name : symbol?
  expected : string?
  bad-pos : exact-nonnegative-integer?
  v : any/c
```

Creates an exn:fail:contract value and raises it as an exception. The *name* argument is used as the source procedure's name in the error message. The *expected* argument is used as a description of the expected contract (i.e., as a string, but the string is intended to contain a contract expression).

In the first form, *v* is the value received by the procedure that does not have the expected type.

In the second form, the bad argument is indicated by an index *bad-pos* (counting from 0), and all of the original arguments *v* are provided (in order). The resulting error message names the bad argument and also lists the other arguments. If *bad-pos* is not less than the number of *v*s, the exn:fail:contract exception is raised.

Examples:

```
> (define (feed-machine bits)
    (if (not (integer? bits))
      (raise-argument-error 'feed-machine "integer?" bits)
      "fed the machine"))

> (feed-machine 'turkey)
feed-machine: contract violation
  expected: integer?
  given: 'turkey
> (define (feed-cow animal)
    (if (not (eq? animal 'cow))
      (raise-argument-error 'feed-cow "'cow" animal)
      "fed the cow"))
```

```
> (feed-cow 'turkey)
feed-cow: contract violation
  expected: 'cow
  given: 'turkey
> (define (feed-animals cow sheep goose cat)
    (if (not (eq? goose 'goose))
      (raise-argument-error 'feed-animals "'goose" 2 cow sheep goose cat)
      "fed the animals"))

> (feed-animals 'cow 'sheep 'dog 'cat)
feed-animals: contract violation
  expected: 'goose
  given: 'dog
  argument position: 3rd
  other arguments...:
   'cow
   'sheep
   'cat
```

```
(raise-result-error name expected v) → any
  name : symbol?
  expected : string?
  v : any/c
(raise-result-error name
                    expected
                    bad-pos
                    v ...)   → any
  name : symbol?
  expected : string?
  bad-pos : exact-nonnegative-integer?
  v : any/c
```

Like raise-argument-error, but the error message describe *v* as a "result" instead of an "argument."

```
(raise-arguments-error name
                       message
                       field
                       v ...
                       ...)     → any
  name : symbol?
  message : string?
  field : string?
  v : any/c
```

Creates an exn:fail:contract value and raises it as an exception. The *name* is used as the source procedure's name in the error message. The *message* is the error message;

if *message* contains newline characters, each extra line should be suitably indented (with one extra space at the start of each line), but it should not end with a newline character. Each *field* must have a corresponding *v*, and the two are rendered on their own line in the error message, with each *v* formatted using the error value conversion handler (see error-value->string-handler).

Example:

```
> (raise-arguments-error 'eat
                         "fish is smaller than its given meal"
                         "fish" 12
                         "meal" 13)
eat: fish is smaller than its given meal
  fish: 12
  meal: 13
```

```
(raise-range-error name
                   type-description
                   index-prefix
                   index
                   in-value
                   lower-bound
                   upper-bound
                   alt-lower-bound) → any
  name : symbol?
  type-description : string?
  index-prefix : string?
  index : exact-integer?
  in-value : any/c
  lower-bound : exact-integer?
  upper-bound : exact-integer?
  alt-lower-bound : (or/c #f exact-integer?)
```

Creates an exn:fail:contract value and raises it as an exception to report an out-of-range error. The *type-description* string describes the value for which the index is meant to select an element, and *index-prefix* is a prefix for the word "index." The *index* argument is the rejected index. The *in-value* argument is the value for which the index was meant. The *lower-bound* and *upper-bound* arguments specify the valid range of indices, inclusive; if *upper-bound* is below *lower-bound*, the value is characterized as "empty." If *alt-lower-bound* is not #f, and if *index* is between *alt-lower-bound* and *upper-bound*, then the error is report as *index* being less than the "starting" index *lower-bound*.

Since *upper-bound* is inclusive, a typical value is *one less than* the size of a collection—for example, (sub1 (vector-length *vec*)), (sub1 (length *lst*)), and so on.

Examples:

```
> (raise-range-error 'vector-ref "vector" "starting
" 5 #(1 2 3 4) 0 3)
vector-ref: starting index is out of range
  starting index: 5
  valid range: [0, 3]
  vector: '#(1 2 3 4)
> (raise-range-error 'vector-ref "vector" "ending
" 5 #(1 2 3 4) 0 3)
vector-ref: ending index is out of range
  ending index: 5
  valid range: [0, 3]
  vector: '#(1 2 3 4)
> (raise-range-error 'vector-ref "vector" "" 3 #() 0 -1)
vector-ref: index is out of range for empty vector
  index: 3
> (raise-range-error 'vector-ref "vector" "ending
" 1 #(1 2 3 4) 2 3 0)
vector-ref: ending index is smaller than starting index
  ending index: 1
  starting index: 2
  valid range: [0, 3]
  vector: '#(1 2 3 4)
```

```
(raise-type-error name expected v) → any
  name : symbol?
  expected : string?
  v : any/c
(raise-type-error name expected bad-pos v ...) → any
  name : symbol?
  expected : string?
  bad-pos : exact-nonnegative-integer?
  v : any/c
```

Like `raise-argument-error`, but with Racket's old formatting conventions, and where *expected* is used as a "type" description instead of a contract expression. Use `raise-argument-error` or `raise-result-error`, instead.

```
(raise-mismatch-error name
                      message
                      v ...+
                      ...+)    → any
  name : symbol?
  message : string?
  v : any/c
```

Similar to `raise-arguments-error`, but using Racket's old formatting conventions, with a required `v` immediately after the first *message* string, and with further *message* strings that

are spliced into the message without line breaks or space. Use raise-arguments-error, instead.

```
(raise-arity-error name arity-v [arg-v ...]) → any
  name : (or/c symbol? procedure?)
  arity-v :       (or/c exact-nonnegative-integer?
                        arity-at-least?
                        (listof
                         (or/c exact-nonnegative-integer?
                               arity-at-least?)))
  arg-v : any/c = #f
```

Creates an exn:fail:contract:arity value and raises it as an exception. The *name* is used for the source procedure's name in the error message.

The *arity-v* value must be a possible result from procedure-arity, except that it does not have to be normalized (see procedure-arity? for the details of normalized arities); raise-arity-error will normalize the arity and use the normalized form in the error message. If *name* is a procedure, its actual arity is ignored.

The *arg-v* arguments are the actual supplied arguments, which are shown in the error message (using the error value conversion handler; see error-value->string-handler); also, the number of supplied *arg-v*s is explicitly mentioned in the message.

```
(raise-syntax-error  name
                     message
                    [expr
                     sub-expr
                     extra-sources]) → any
  name : (or/c symbol? #f)
  message : string?
  expr : any/c = #f
  sub-expr : any/c = #f
  extra-sources : (listof syntax?) = null
```

Creates an exn:fail:syntax value and raises it as an exception. Macros use this procedure to report syntax errors.

The *name* argument is usually #f when *expr* is provided; it is described in more detail below. The *message* is used as the main body of the error message; if *message* contains newline characters, each new line should be suitably indented (with one space at the start), and it should not end with a newline character.

The optional *expr* argument is the erroneous source syntax object or S-expression (but the expression #f cannot be represented by itself; it must be wrapped as a syntax object). The optional *sub-expr* argument is a syntax object or S-expression (again, #f cannot represent

itself) within *expr* that more precisely locates the error. Both may appear in the generated error-message text if `error-print-source-location` is `#t`. Source location information in the error-message text is similarly extracted from *sub-expr* or *expr* when at least one is a syntax object and `error-print-source-location` is `#t`.

If *sub-expr* is provided and not `#f`, it is used (in syntax form) for the `exprs` field of the generated exception record, else the *expr* is used if provided and not `#f`. In either case, the syntax object is `cons`ed onto *extra-sources* to produce the `exprs` field, or *extra-sources* is used directly for `exprs` if neither *expr* nor *sub-expr* is provided and not `#f`.

The form name used in the generated error message is determined through a combination of the *name*, *expr*, and *sub-expr* arguments:

- When *name* is `#f`, and when *expr* is either an identifier or a syntax pair containing an identifier as its first element, then the form name from the error message is the identifier's symbol.
- When *name* is `#f` and when *expr* is not an identifier or a syntax pair containing an identifier as its first element, then the form name in the error message is `"?"`.
- When *name* is a symbol, then the symbol is used as the form name in the generated error message.

10.2.3 Handling Exceptions

```
(call-with-exception-handler f thunk) → any
  f : (any/c . -> . any)
  thunk : (-> any)
```

Installs *f* as the exception handler for the dynamic extent of the call to *thunk*. If an exception is raised during the evaluation of *thunk* (in an extension of the current continuation that does not have its own exception handler), then *f* is applied to the `raise`d value in the continuation of the `raise` call (but the continuation is normally extended with a continuation barrier; see §1.1.12 "Prompts, Delimited Continuations, and Barriers" and `raise`).

Any procedure that takes one argument can be an exception handler. Normally, an exception handler escapes from the context of the `raise` call via `abort-current-continuation` or some other escape mechanism. To propagate an exception to the "previous" exception handler—that is, the exception handler associated with the rest of the continuation after the point where the called exception handler was associated with the continuation—an exception handler can simply return a result instead of escaping, in which case the `raise` call propagates the value to the previous exception handler (still in the dynamic extent of the call to `raise`, and under the same barrier, if any). If an exception handler returns a result and no previous handler is available, the uncaught-exception handler is used.

A call to an exception handler is `parameterize-breaked` to disable breaks, and it is wrapped with `call-with-exception-handler` to install an exception handler that reports both the original and newly raised exceptions via the error display handler and then escapes via the error escape handler.

```
(uncaught-exception-handler) → (any/c . -> . any)
(uncaught-exception-handler f) → void?
  f : (any/c . -> . any)
```

A parameter that determines an *uncaught-exception handler* used by `raise` when the relevant continuation has no exception handler installed with `call-with-exception-handler` or `with-handlers`. Unlike exception handlers installed with `call-with-exception-handler`, the uncaught-exception handler must not return a value when called by `raise`; if it returns, an exception is raised (to be handled by an exception handler that reports both the original and newly raised exception).

The default uncaught-exception handler prints an error message using the current error display handler (see `error-display-handler`), unless the argument to the handler is an instance of `exn:break:hang-up`. If the argument to the handler is an instance of `exn:break:hang-up` or `exn:break:terminate`, the default uncaught-exception handler then calls the exit handler with `1`, which normally exits or escapes. For any argument, the default uncaught-exception handler then escapes by calling the current error escape handler (see `error-escape-handler`). The call to each handler is `parameterized` to set `error-display-handler` to the default error display handler, and it is `parameterize-breaked` to disable breaks. The call to the error escape handler is further parameterized to set `error-escape-handler` to the default error escape handler; if the error escape handler returns, then the default error escape handler is called.

When the current error display handler is the default handler, then the error-display call is parameterized to install an emergency error display handler that logs an error (see `log-error`) and never fails.

```
(with-handlers ([pred-expr handler-expr] ...)
  body ...+)
```

Evaluates each *`pred-expr`* and *`handler-expr`* in the order that they are specified, and then evaluates the *`body`*s with a new exception handler during its dynamic extent.

The new exception handler processes an exception only if one of the *`pred-expr`* procedures returns a true value when applied to the exception, otherwise the exception handler is invoked from the continuation of the `with-handlers` expression (by raising the exception again). If an exception is handled by one of the *`handler-expr`* procedures, the result of the entire `with-handlers` expression is the return value of the handler.

When an exception is raised during the evaluation of *`body`*s, each predicate procedure *`pred-expr`* is applied to the exception value; if a predicate returns a true value, the corresponding

handler-expr procedure is invoked with the exception as an argument. The predicates are tried in the order that they are specified.

Before any predicate or handler procedure is invoked, the continuation of the entire `with-handlers` expression is restored, but also `parameterize-breaked` to disable breaks. Thus, breaks are disabled by default during the predicate and handler procedures (see §10.6 "Breaks"), and the exception handler is the one from the continuation of the `with-handlers` expression.

The `exn:fail?` procedure is useful as a handler predicate to catch all error exceptions. Avoid using `(lambda (x) #t)` as a predicate, because the `exn:break` exception typically should not be caught (unless it will be re-raised to cooperatively break). Beware, also, of catching and discarding exceptions, because discarding an error message can make debugging unnecessarily difficult; instead of discarding an error message, consider logging it via `log-error` or a logging form created by `define-logger`.

Examples:

```
> (with-handlers ([exn:fail:syntax?
                   (λ (e) (displayln "got a syntax error"))])
    (raise-syntax-error #f "a syntax error"))
got a syntax error

> (with-handlers ([exn:fail:syntax?
                   (λ (e) (displayln "got a syntax error"))]
                  [exn:fail?
                   (λ (e) (displayln "fallback clause"))])
    (raise-syntax-error #f "a syntax error"))
got a syntax error
```

```
(with-handlers* ([pred-expr handler-expr] ...)
  body ...+)
```

Like `with-handlers`, but if a *handler-expr* procedure is called, breaks are not explicitly disabled, and the handler call is in tail position with respect to the `with-handlers*` form.

10.2.4 Configuring Default Handling

```
(error-escape-handler) → (-> any)
(error-escape-handler proc) → void?
  proc : (-> any)
```

A parameter for the *error escape handler*, which takes no arguments and escapes from the dynamic context of an exception. The default error escape handler escapes using `(abort-current-continuation (default-continuation-prompt-tag) void)`.

The error escape handler is normally called directly by an exception handler, in a parameterization that sets the error display handler and error escape handler to the default handlers, and it is normally `parameterize-break`ed to disable breaks. To escape from a run-time error in a different context, use `raise` or `error`.

Due to a continuation barrier around exception-handling calls, an error escape handler cannot invoke a full continuation that was created prior to the exception, but it can abort to a prompt (see `call-with-continuation-prompt`) or invoke an escape continuation (see `call-with-escape-continuation`).

```
(error-display-handler) → (string? any/c . -> . any)
(error-display-handler proc) → void?
  proc : (string? any/c . -> . any)
```

A parameter for the *error display handler*, which is called by the default exception handler with an error message and the exception value. More generally, the handler's first argument is a string to print as an error message, and the second is a value representing a raised exception.

The default error display handler `display`s its first argument to the current error port (determined by the `current-error-port` parameter) and extracts a stack trace (see `continuation-mark-set->context`) to display from the second argument if it is an `exn` value but not an `exn:fail:user` value.

The default error display handler in DrRacket also uses the second argument to highlight source locations.

To report a run-time error, use `raise` or procedures like `error`, instead of calling the error display handler directly.

```
(error-print-width) → (and/c exact-integer? (>=/c 3))
(error-print-width width) → void?
  width : (and/c exact-integer? (>=/c 3))
```

A parameter whose value is used as the maximum number of characters used to print a Racket value that is embedded in a primitive error message.

```
(error-print-context-length) → exact-nonnegative-integer?
(error-print-context-length cnt) → void?
  cnt : exact-nonnegative-integer?
```

A parameter whose value is used by the default error display handler as the maximum number of lines of context (or "stack trace") to print; a single "..." line is printed if more lines are available after the first *cnt* lines. A `0` value for *cnt* disables context printing entirely.

```
(error-value->string-handler)
    (any/c exact-nonnegative-integer?
 →          . -> .
            string?)
```

```
(error-value->string-handler proc) → void?
         (any/c exact-nonnegative-integer?
  proc :        . -> .
                string?)
```

A parameter that determines the *error value conversion handler*, which is used to print a Racket value that is embedded in a primitive error message.

The integer argument to the handler specifies the maximum number of characters that should be used to represent the value in the resulting string. The default error value conversion handler `print`s the value into a string (using the current global port print handler; see `global-port-print-handler`). If the printed form is too long, the printed form is truncated and the last three characters of the return string are set to "...".

If the string returned by an error value conversion handler is longer than requested, the string is destructively "truncated" by setting the first extra position in the string to the null character. If a non-string is returned, then the string `"..."` is used. If a primitive error string needs to be generated before the handler has returned, the default error value conversion handler is used.

Calls to an error value conversion handler are `parameterized` to re-install the default error value conversion handler, and to enable printing of unreadable values (see `print-unreadable`).

```
(error-print-source-location) → boolean?
(error-print-source-location include?) → void?
  include? : any/c
```

A parameter that controls whether read and syntax error messages include source information, such as the source line and column or the expression. This parameter also controls the error message when a module-defined variable is accessed before its definition is executed; the parameter determines whether the message includes a module name. Only the message field of an `exn:fail:read`, `exn:fail:syntax`, or `exn:fail:contract:variable` structure is affected by the parameter. The default is `#t`.

10.2.5 Built-in Exception Types

```
(struct exn (message continuation-marks)
    #:extra-constructor-name make-exn
    #:transparent)
  message : string?
  continuation-marks : continuation-mark-set?
```

The base structure type for exceptions. The `message` field contains an error message, and the `continuation-marks` field contains the value produced by `(current-continuation-`

marks) immediately before the exception was raised.

Exceptions raised by Racket form a hierarchy under exn:

```
exn
  exn:fail
    exn:fail:contract
      exn:fail:contract:arity
      exn:fail:contract:divide-by-zero
      exn:fail:contract:non-fixnum-result
      exn:fail:contract:continuation
      exn:fail:contract:variable
    exn:fail:syntax
      exn:fail:syntax:unbound
      exn:fail:syntax:missing-module
    exn:fail:read
      exn:fail:read:eof
      exn:fail:read:non-char
    exn:fail:filesystem
      exn:fail:filesystem:exists
      exn:fail:filesystem:version
      exn:fail:filesystem:errno
      exn:fail:filesystem:missing-module
    exn:fail:network
      exn:fail:network:errno
    exn:fail:out-of-memory
    exn:fail:unsupported
    exn:fail:user
  exn:break
    exn:break:hang-up
    exn:break:terminate
```

```
(struct exn:fail exn ()
    #:extra-constructor-name make-exn:fail
    #:transparent)
```

Raised for exceptions that represent errors, as opposed to exn:break.

```
(struct exn:fail:contract exn:fail ()
    #:extra-constructor-name make-exn:fail:contract
    #:transparent)
```

Raised for errors from the inappropriate run-time use of a function or syntactic form.

```
(struct exn:fail:contract:arity exn:fail:contract ()
    #:extra-constructor-name make-exn:fail:contract:arity
    #:transparent)
```

Raised when a procedure is applied to the wrong number of arguments.

```
(struct exn:fail:contract:divide-by-zero exn:fail:contract ()
    #:extra-constructor-name
    make-exn:fail:contract:divide-by-zero
    #:transparent)
```

Raised for division by exact zero.

```
(struct exn:fail:contract:non-fixnum-result exn:fail:contract ()
    #:extra-constructor-name
    make-exn:fail:contract:non-fixnum-result
    #:transparent)
```

Raised by functions like `fx+` when the result would not be a fixnum.

```
(struct exn:fail:contract:continuation exn:fail:contract ()
    #:extra-constructor-name make-exn:fail:contract:continuation
    #:transparent)
```

Raised when a continuation is applied where the jump would cross a continuation barrier.

```
(struct exn:fail:contract:variable exn:fail:contract (id)
    #:extra-constructor-name make-exn:fail:contract:variable
    #:transparent)
  id : symbol?
```

Raised for a reference to a not-yet-defined top-level variable or module-level variable.

```
(struct exn:fail:syntax exn:fail (exprs)
    #:extra-constructor-name make-exn:fail:syntax
    #:transparent)
  exprs : (listof syntax?)
```

Raised for a syntax error that is not a `read` error. The `exprs` indicate the relevant source expressions, least-specific to most-specific.

This structure type implements the `prop:exn:srclocs` property.

```
(struct exn:fail:syntax:unbound exn:fail:syntax ()
    #:extra-constructor-name make-exn:fail:syntax:unbound
    #:transparent)
```

Raised by `#%top` or `set!` for an unbound identifier within a module.

```
(struct exn:fail:syntax:missing-module exn:fail:syntax (path)
    #:extra-constructor-name make-exn:fail:syntax:missing-module
    #:transparent)
  path : module-path?
```

Raised by the default module name resolver or default load handler to report a module path—a reported in the `path` field—whose implementation file cannot be found.

The default module name resolver raises this exception only when it is given a syntax object as its second argument, and the default load handler raises this exception only when the value of `current-module-path-for-load` is a syntax object (in which case both the `exprs` field and the `path` field are determined by the syntax object.

This structure type implements the `prop:exn:missing-module` property.

```
(struct exn:fail:read exn:fail (srclocs)
    #:extra-constructor-name make-exn:fail:read
    #:transparent)
  srclocs : (listof srcloc?)
```

Raised for a `read` error. The `srclocs` indicate the relevant source expressions.

```
(struct exn:fail:read:eof exn:fail:read ()
    #:extra-constructor-name make-exn:fail:read:eof
    #:transparent)
```

Raised for a `read` error, specifically when the error is due to an unexpected end-of-file.

```
(struct exn:fail:read:non-char exn:fail:read ()
    #:extra-constructor-name make-exn:fail:read:non-char
    #:transparent)
```

Raised for a `read` error, specifically when the error is due to an unexpected non-character (i.e., “special”) element in the input stream.

```
(struct exn:fail:filesystem exn:fail ()
    #:extra-constructor-name make-exn:fail:filesystem
    #:transparent)
```

Raised for an error related to the filesystem (such as a file not found).

```
(struct exn:fail:filesystem:exists exn:fail:filesystem ()
    #:extra-constructor-name make-exn:fail:filesystem:exists
    #:transparent)
```

Raised for an error when attempting to create a file that exists already.

```
(struct exn:fail:filesystem:version exn:fail:filesystem ()
    #:extra-constructor-name make-exn:fail:filesystem:version
    #:transparent)
```

Raised for a version-mismatch error when loading an extension.

```
(struct exn:fail:filesystem:errno exn:fail:filesystem (errno)
    #:extra-constructor-name make-exn:fail:filesystem:errno
    #:transparent)
  errno : (cons/c exact-integer? (or/c 'posix 'windows 'gai))
```

Raised for a filesystem error for which a system error code is available. The symbol part of an `errno` field indicates the category of the error code: `'posix` indicates a C/Posix `errno` value, `'windows` indicates a Windows system error code (under Windows, only), and `'gai` indicates a `getaddrinfo` error code (which shows up only in `exn:fail:network:errno` exceptions for operations that resolve hostnames, but is allowed in `exn:fail:filesystem:errno` instances for consistency).

```
(struct exn:fail:filesystem:missing-module exn:fail:filesystem
        (path)
    #:extra-constructor-name
    make-exn:fail:filesystem:missing-module
    #:transparent)
  path : module-path?
```

Raised by the default module name resolver or default load handler to report a module path—a reported in the `path` field—whose implementation file cannot be found.

The default module name resolver raises this exception only when it is *not* given a syntax object as its second argument, and the default load handler raises this exception only when the value of `current-module-path-for-load` is *not* a syntax object.

This structure type implements the `prop:exn:missing-module` property.

```
(struct exn:fail:network exn:fail ()
    #:extra-constructor-name make-exn:fail:network
    #:transparent)
```

Raised for TCP and UDP errors.

```
(struct exn:fail:network:errno exn:fail:network (errno)
    #:extra-constructor-name make-exn:fail:network:errno
    #:transparent)
  errno : (cons/c exact-integer? (or/c 'posix 'windows 'gai))
```

Raised for a TCP or UDP error for which a system error code is available, where the `errno` field is as for `exn:fail:filesystem:errno`.

```
(struct exn:fail:out-of-memory exn:fail ()
    #:extra-constructor-name make-exn:fail:out-of-memory
    #:transparent)
```

Raised for an error due to insufficient memory, in cases where sufficient memory is at least available for raising the exception.

```
(struct exn:fail:unsupported exn:fail ()
    #:extra-constructor-name make-exn:fail:unsupported
    #:transparent)
```

Raised for an error due to an unsupported feature on the current platform or configuration.

```
(struct exn:fail:user exn:fail ()
    #:extra-constructor-name make-exn:fail:user
    #:transparent)
```

Raised for errors that are intended to be seen by end users. In particular, the default error printer does not show the program context when printing the error message.

```
(struct exn:break exn (continuation)
    #:extra-constructor-name make-exn:break
    #:transparent)
  continuation : continuation?
```

Raised asynchronously (when enabled) in response to a break request. The `continuation` field can be used by a handler to resume the interrupted computation.

```
(struct exn:break:hang-up exn:break ()
    #:extra-constructor-name make-exn:break:hang-up
    #:transparent)
```

Raised asynchronously for hang-up breaks. The default uncaught-exception handler reacts to this exception type by calling the exit handler.

```
(struct exn:break:terminate exn:break ()
    #:extra-constructor-name make-exn:break:terminate
    #:transparent)
```

Raised asynchronously for termination-request breaks. The default uncaught-exception handler reacts to this exception type by calling the exit handler.

prop:exn:srclocs : struct-type-property?

A property that identifies structure types that provide a list of srcloc values. The property is normally attached to structure types used to represent exception information.

The property value must be a procedure that accepts a single value—the structure type instance from which to extract source locations—and returns a list of srclocs. Some error display handlers use only the first returned location.

As an example,

```
#lang racket

;; We create a structure that supports the
;; prop:exn:srcloc protocol.  It carries
;; with it the location of the syntax that
;; is guilty.
(define-struct (exn:fail:he-who-shall-not-be-named
                exn:fail)
  (a-srcloc)
  #:property prop:exn:srclocs
  (lambda (a-struct)
    (match a-struct
      [(struct exn:fail:he-who-shall-not-be-named
         (msg marks a-srcloc))
       (list a-srcloc)])))

;; We can play with this by creating a form that
;; looks at identifiers, and only flags specific ones.
(define-syntax (skeeterize stx)
  (syntax-case stx ()
    [(_ expr)
     (cond
       [(and (identifier? #'expr)
             (eq? (syntax-e #'expr) 'voldemort))
        (quasisyntax/loc stx
          (raise (make-exn:fail:he-who-shall-not-be-named
                  "oh dear don't say his name"
                  (current-continuation-marks)
                  (srcloc '#,(syntax-source #'expr)
                          '#,(syntax-line #'expr)
                          '#,(syntax-column #'expr)
                          '#,(syntax-position #'expr)
                          '#,(syntax-span #'expr)))))]
```

```
       [else
        ;; Otherwise, leave the expression alone.
        #'expr])]))

(define (f x)
  (* (skeeterize x) x))

(define (g voldemort)
  (* (skeeterize voldemort) voldemort))

;; Examples:
(f 7)
(g 7)
;; The error should highlight the use
;; of the one-who-shall-not-be-named
;; in g.
```

```
(exn:srclocs? v) → boolean?
  v : any/c
```

Returns `#t` if `v` has the `prop:exn:srclocs` property, `#f` otherwise.

```
(exn:srclocs-accessor v)
 → (exn:srclocs? . -> . (listof srcloc))
  v : exn:srclocs?
```

Returns the `srcloc`-getting procedure associated with `v`.

```
(struct srcloc (source line column position span)
    #:extra-constructor-name make-srcloc
    #:transparent)
  source : any/c
  line : (or/c exact-positive-integer? #f)
  column : (or/c exact-nonnegative-integer? #f)
  position : (or/c exact-positive-integer? #f)
  span : (or/c exact-nonnegative-integer? #f)
```

The fields of a `srcloc` instance are as follows:

- `source` — An arbitrary value identifying the source, often a path (see §15.1 “Paths”).
- `line` — The line number (counts from 1) or `#f` (unknown).
- `column` — The column number (counts from 0) or `#f` (unknown).

- `position` — The starting position (counts from 1) or `#f` (unknown).
- `span` — The number of covered positions (counts from 0) or `#f` (unknown).

```
(srcloc->string srcloc) → (or/c string? #f)
  srcloc : srcloc?
```

Formats *srcloc* as a string suitable for error reporting. A path source in *srcloc* is shown relative to the value of `current-directory-for-user`. The result is `#f` if *srcloc* does not contain enough information to format a string.

```
prop:exn:missing-module : struct-type-property?
```

A property that identifies structure types that provide a module path for a load that fails because a module is not found.

The property value must be a procedure that accepts a single value—the structure type instance from which to extract source locations—and returns a module path.

```
(exn:missing-module? v) → boolean?
  v : any/c
```

Returns `#t` if *v* has the `prop:exn:missing-module` property, `#f` otherwise.

```
(exn:missing-module-accessor v)
 → (exn:missing-module? . -> . module-path?)
  v : exn:srclocs?
```

Returns the module path-getting procedure associated with *v*.

10.3 Delayed Evaluation

```
(require racket/promise)
```
package: `base`

The bindings documented in this section are provided by the `racket/promise` and `racket` libraries, but not `racket/base`.

A *promise* encapsulates an expression to be evaluated on demand via `force`. After a promise has been `force`d, every later `force` of the promise produces the same result.

```
(promise? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a promise, `#f` otherwise.

```
(delay body ...+)
```

Creates a promise that, when `force`d, evaluates the *body*s to produce its value. The result is then cached, so further uses of `force` produce the cached value immediately. This includes multiple values and exceptions.

```
(lazy body ...+)
```

Like `delay`, if the last *body* produces a promise when forced, then this promise is `force`d, too, to obtain a value. In other words, this form creates a composable promise, where the computation of its body is "attached" to the computation of the following promise, and a single `force` iterates through the whole chain, tail-calling each step.

Note that the last *body* of this form must produce a single value, but the value can itself be a `delay` promise that returns multiple values.

The `lazy` form is useful for implementing lazy libraries and languages, where tail calls can be wrapped in a promise.

```
(force v) → any
  v : any/c
```

If *v* is a promise, then the promise is forced to obtain a value. If the promise has not been forced before, then the result is recorded in the promise so that future `force`s on the promise produce the same value (or values). If forcing the promise raises an exception, then the exception is similarly recorded so that forcing the promise will raise the same exception every time.

If *v* is `force`d again before the original call to `force` returns, then the `exn:fail` exception is raised.

If *v* is not a promise, then it is returned as the result.

```
(promise-forced? promise) → boolean?
  promise : promise?
```

Returns `#t` if *promise* has been forced.

```
(promise-running? promise) → boolean?
  promise : promise?
```

Returns `#t` if *promise* is currently being forced. (Note that a promise can be either running or forced but not both.)

10.3.1 Additional Promise Kinds

```
(delay/name body ...+)
```

Creates a “call-by-name” promise that is similar to `delay`-promises, except that the resulting value is not cached. This kind of promise is essentially a thunk that is wrapped in a way that `force` recognizes.

If a `delay/name` promise forces itself, no exception is raised, the promise is never considered “running” or “forced” in the sense of `promise-running?` and `promise-forced?`.

```
(delay/strict body ...+)
```

Creates a “strict” promise: it is evaluated immediately, and the result is wrapped in a promise value. Note that the body can evaluate to multiple values, and forcing the resulting promise will return these values.

```
(delay/sync body ...+)
```

Produces a promise where an attempt to `force` the promise by a thread other than one currently running the promise causes the `force` to block until a result is available. This kind of promise is also a synchronizable event for use with `sync`; `sync`ing on the promise does not `force` it, but merely waits until a value is forced by another thread.

If a promise created by `delay/sync` is forced on a thread that is already running the promise, an exception is raised in the same way as for promises created with `delay`.

```
(delay/thread body/option ...+)

body/option = body
            | #:group thread-group-expr
```

Like `delay/sync`, but begins the computation immediately on a newly created thread. The thread is created under the thread group specified by *`thread-group-expr`*, which defaults to `(make-thread-group)`. A `#:group` specification can appear at most once.

Exceptions raised by the *`body`*s are caught as usual and raised only when the promise is `force`d. Unlike `delay/sync`, if the thread running *`body`* terminates without producing a result or exception, `force` of the promise raises an exception (instead of blocking).

```
(delay/idle body/option ...+)

body/option = body
            | #:wait-for wait-evt-expr
            | #:work-while while-evt-expr
            | #:tick tick-secs-expr
            | #:use use-ratio-expr
```

Like `delay/thread`, but with the following differences:

- the computation does not start until the event produced by *`wait-evt-expr`* is ready, where the default is `(system-idle-evt)`;
- the computation thread gets to work only when the process is otherwise idle as determined by *`while-evt-expr`*, which also defaults to `(system-idle-evt)`;
- the thread is allowed to run only periodically: out of every *`tick-secs-expr`* (defaults to `0.2`) seconds, the thread is allowed to run *`use-ratio-expr`* (defaults to `0.12`) of the time proportionally; i.e., the thread runs for `(* tick-secs-expr use-ratio-expr)` seconds.

If the promise is `forced` before the computation is done, it runs the rest of the computation immediately without waiting on events or periodically restricting evaluation.

A `#:wait-for`, `#:work-while`, `#:tick`, or `#:use` specification can appear at most once.

10.4 Continuations

§10.3 “Continuations” in ***The Racket Guide*** introduces continuations.

See §1.1.1 “Sub-expression Evaluation and Continuations” and §1.1.12 “Prompts, Delimited Continuations, and Barriers” for general information about continuations. Racket’s support for prompts and composable continuations most closely resembles Dorai Sitaram’s `%` and `fcontrol` operator [Sitaram93].

Racket installs a continuation barrier around evaluation in the following contexts, preventing full-continuation jumps into the evaluation context protected by the barrier:

- applying an exception handler, an error escape handler, or an error display handler (see §10.2 “Exceptions”);
- applying a macro transformer (see §12.4 “Syntax Transformers”), evaluating a compile-time expression, or applying a module name resolver (see §14.4.1 “Resolving Module Names”);
- applying a custom-port procedure (see §13.1.9 “Custom Ports”), an event guard procedure (see §11.2.1 “Events”), or a parameter guard procedure (see §11.3.2 “Parameters”);
- applying a security-guard procedure (see §14.6 “Security Guards”);
- applying a will procedure (see §16.3 “Wills and Executors”); or
- evaluating or loading code from the stand-alone Racket command line (see §18.1 “Running Racket or GRacket”).

In addition, extensions of Racket may install barriers in additional contexts. Finally, call-with-continuation-barrier applies a thunk barrier between the application and the current continuation.

```
(call-with-continuation-prompt proc
                               [prompt-tag
                                handler]
                               arg ...)   → any
  proc : procedure?
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
  handler : (or/c procedure? #f) = #f
  arg : any/c
```

Applies *proc* to the given *arg*s with the current continuation extended by a prompt. The prompt is tagged by *prompt-tag*, which must be a result from either default-continuation-prompt-tag (the default) or make-continuation-prompt-tag. The result of *proc* is the result of the call-with-continuation-prompt call.

The *handler* argument specifies a handler procedure to be called in tail position with respect to the call-with-continuation-prompt call when the installed prompt is the target of an abort-current-continuation call with *prompt-tag*; the remaining arguments of abort-current-continuation are supplied to the handler procedure. If *handler* is #f, the default handler accepts a single *abort-thunk* argument and calls (call-with-continuation-prompt *abort-thunk* *prompt-tag* #f); that is, the default handler reinstalls the prompt and continues with a given thunk.

```
(abort-current-continuation prompt-tag
                            v ...)     → any
  prompt-tag : any/c
  v : any/c
```

Resets the current continuation to that of the nearest prompt tagged by *prompt-tag* in the current continuation; if no such prompt exists, the exn:fail:contract:continuation exception is raised. The *v*s are delivered as arguments to the target prompt's handler procedure.

The protocol for *v*s supplied to an abort is specific to the *prompt-tag*. When abort-current-continuation is used with (default-continuation-prompt-tag), generally, a single thunk should be supplied that is suitable for use with the default prompt handler. Similarly, when call-with-continuation-prompt is used with (default-continuation-prompt-tag), the associated handler should generally accept a single thunk argument.

Each thread's continuation starts with a prompt for (default-continuation-prompt-tag) that uses the default handler, which accepts a single thunk to apply (with the prompt intact).

```
(make-continuation-prompt-tag) → continuation-prompt-tag?
(make-continuation-prompt-tag sym) → continuation-prompt-tag?
  sym : symbol?
```

Creates a prompt tag that is not `equal?` to the result of any other value (including prior or future results from `make-continuation-prompt-tag`). The optional *sym* argument, if supplied, is used when printing the prompt tag.

```
(default-continuation-prompt-tag) → continuation-prompt-tag?
```

Returns a constant prompt tag for which a prompt is installed at the start of every thread's continuation; the handler for each thread's initial prompt accepts any number of values and returns. The result of `default-continuation-prompt-tag` is the default tag for any procedure that accepts a prompt tag.

```
(call-with-current-continuation proc
                                [prompt-tag]) → any
  proc : (continuation? . -> . any)
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
```

Captures the current continuation up to the nearest prompt tagged by *prompt-tag*; if no such prompt exists, the `exn:fail:contract:continuation` exception is raised. The truncated continuation includes only continuation marks and `dynamic-wind` frames installed since the prompt.

The capture continuation is delivered to *proc*, which is called in tail position with respect to the `call-with-current-continuation` call.

If the continuation argument to *proc* is ever applied, then it removes the portion of the current continuation up to the nearest prompt tagged by *prompt-tag* (not including the prompt; if no such prompt exists, the `exn:fail:contract:continuation` exception is raised), or up to the nearest continuation frame (if any) shared by the current and captured continuations—whichever is first. While removing continuation frames, `dynamic-wind` *post-thunk*s are executed. Finally, the (unshared portion of the) captured continuation is appended to the remaining continuation, applying `dynamic-wind` *pre-thunk*s.

The arguments supplied to an applied procedure become the result values for the restored continuation. In particular, if multiple arguments are supplied, then the continuation receives multiple results.

If, at application time, a continuation barrier would be introduced by replacing the current continuation with the applied one, then the `exn:fail:contract:continuation` exception is raised.

A continuation can be invoked from the thread (see §11.1 “Threads”) other than the one where it was captured.

```
(call/cc proc [prompt-tag]) → any
  proc : (continuation? . -> . any)
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
```

The `call/cc` binding is an alias for `call-with-current-continuation`.

```
(call-with-composable-continuation  proc
                                    [prompt-tag]) → any
  proc : (continuation? . -> . any)
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
```

Similar to `call-with-current-continuation`, but applying the resulting continuation procedure does not remove any portion of the current continuation. Instead, application always extends the current continuation with the captured continuation (without installing any prompts other than those captured in the continuation).

When `call-with-composable-continuation` is called, if a continuation barrier appears in the continuation before the closest prompt tagged by *prompt-tag*, the `exn:fail:contract:continuation` exception is raised (because attempting to apply the continuation would always fail).

```
(call-with-escape-continuation proc) → any
  proc : (continuation? . -> . any)
```

Like `call-with-current-continuation`, but *proc* is not called in tail position, and the continuation procedure supplied to *proc* can only be called during the dynamic extent of the `call-with-escape-continuation` call.

Due to the limited applicability of its continuation, `call-with-escape-continuation` can be implemented more efficiently than `call-with-current-continuation`.

A continuation obtained from `call-with-escape-continuation` is actually a kind of prompt. Escape continuations are provided mainly for backwards compatibility, since they pre-date general prompts in Racket, and because `call/ec` is often an easy replacement for `call/cc` to improve performance.

```
(call/ec proc) → any
  proc : (continuation? . -> . any)
```

The `call/ec` binding is an alias for `call-with-escape-continuation`.

```
(let/cc k body ...+)
```

Equivalent to `(call/cc (lambda (k) body ...))`.

```
(let/ec k body ...+)
```

Equivalent to `(call/ec (lambda (k) body ...))`.

```
(call-with-continuation-barrier thunk) → any
  thunk : (-> any)
```

Applies *thunk* with a continuation barrier between the application and the current continuation. The results of *thunk* are the results of the `call-with-continuation-barrier` call.

```
(continuation-prompt-available? prompt-tag
                                [cont])     → any
  prompt-tag : continuation-prompt-tag?
  cont : continuation? = (call/cc values)
```

Returns `#t` if *cont*, which must be a continuation, includes a prompt tagged by *prompt-tag*, `#f` otherwise.

```
(continuation? v) → boolean?
  v : any/c
```

Return `#t` if *v* is a continuation as produced by `call-with-current-continuation`, `call-with-composable-continuation`, or `call-with-escape-continuation`, `#f` otherwise.

```
(continuation-prompt-tag? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a continuation prompt tag as produced by `default-continuation-prompt-tag` or `make-continuation-prompt-tag`.

```
(dynamic-wind pre-thunk
              value-thunk
              post-thunk) → any
  pre-thunk : (-> any)
  value-thunk : (-> any)
  post-thunk : (-> any)
```

Applies its three thunk arguments in order. The value of a `dynamic-wind` expression is the value returned by *`value-thunk`*. The *`pre-thunk`* procedure is invoked before calling *`value-thunk`* and *`post-thunk`* is invoked after *`value-thunk`* returns. The special properties of `dynamic-wind` are manifest when control jumps into or out of the *`value-thunk`* application (either due to a prompt abort or a continuation invocation): every time control jumps into the *`value-thunk`* application, *`pre-thunk`* is invoked, and every time control jumps out of *`value-thunk`*, *`post-thunk`* is invoked. (No special handling is performed for jumps into or out of the *`pre-thunk`* and *`post-thunk`* applications.)

When `dynamic-wind` calls *`pre-thunk`* for normal evaluation of *`value-thunk`*, the continuation of the *`pre-thunk`* application calls *`value-thunk`* (with `dynamic-wind`'s special jump handling) and then *`post-thunk`*. Similarly, the continuation of the *`post-thunk`* application returns the value of the preceding *`value-thunk`* application to the continuation of the entire `dynamic-wind` application.

When *`pre-thunk`* is called due to a continuation jump, the continuation of *`pre-thunk`*

- jumps to a more deeply nested *`pre-thunk`*, if any, or jumps to the destination continuation; then
- continues with the context of the *`pre-thunk`*'s `dynamic-wind` call.

Normally, the second part of this continuation is never reached, due to a jump in the first part. However, the second part is relevant because it enables jumps to escape continuations that are contained in the context of the `dynamic-wind` call. Furthermore, it means that the continuation marks (see §10.5 "Continuation Marks") and parameterization (see §11.3.2 "Parameters") for *`pre-thunk`* correspond to those of the `dynamic-wind` call that installed *`pre-thunk`*. The *`pre-thunk`* call, however, is `parameterize-breaked` to disable breaks (see also §10.6 "Breaks").

Similarly, when *`post-thunk`* is called due to a continuation jump, the continuation of *`post-thunk`* jumps to a less deeply nested *`post-thunk`*, if any, or jumps to a *`pre-thunk`* protecting the destination, if any, or jumps to the destination continuation, then continues from the *`post-thunk`*'s `dynamic-wind` application. As for *`pre-thunk`*, the parameterization of the original `dynamic-wind` call is restored for the call, and the call is `parameterize-breaked` to disable breaks.

In both cases, the target for a jump is recomputed after each *`pre-thunk`* or *`post-thunk`* completes. When a prompt-delimited continuation (see §1.1.12 "Prompts, Delimited Continuations, and Barriers") is captured in a *`post-thunk`*, it might be delimited and instantiated in such a way that the target of a jump turns out to be different when the continuation is applied than when the continuation was captured. There may even be no appropriate target, if a relevant prompt or escape continuation is not in the continuation after the restore; in that case, the first step in a *`pre-thunk`* or *`post-thunk`*'s continuation can raise an exception.

Examples:

```
> (let ([v (let/ec out
             (dynamic-wind
              (lambda () (display "in "))
              (lambda ()
                (display "pre ")
                (display (call/cc out))
                #f)
              (lambda () (display "out "))))])
    (when v (v "post ")))
in pre out in post out

> (let/ec k0
    (let/ec k1
      (dynamic-wind
       void
       (lambda () (k0 'cancel))
       (lambda () (k1 'cancel-canceled)))))
'cancel-canceled
> (let* ([x (make-parameter 0)]
         [l null]
         [add (lambda (a b)
                (set! l (append l (list (cons a b)))))])
    (let ([k (parameterize ([x 5])
               (dynamic-wind
                   (lambda () (add 1 (x)))
                   (lambda () (parameterize ([x 6])
                                (let ([k+e (let/cc k (cons k void))])
                                  (add 2 (x))
                                  ((cdr k+e))
                                  (car k+e))))
                   (lambda () (add 3 (x)))))])
      (parameterize ([x 7])
        (let/cc esc
          (k (cons void esc))))))
    l)
'((1 . 5) (2 . 6) (3 . 5) (1 . 5) (2 . 6) (3 . 5))
```

10.4.1 Additional Control Operators

`(require racket/control)` package: `base`

The bindings documented in this section are provided by the `racket/control` library, not `racket/base` or `racket`.

The `racket/control` library provides various control operators from the research litera-

ture on higher-order control operators, plus a few extra convenience forms. These control operators are implemented in terms of call-with-continuation-prompt, call-with-composable-continuation, etc., and they generally work sensibly together. Many are redundant; for example, reset and prompt are aliases.

```
(call/prompt proc [prompt-tag handler] arg ...) → any
  proc : procedure?
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
  handler : (or/c procedure? #f) = #f
  arg : any/c
```

The call/prompt binding is an alias for call-with-continuation-prompt.

```
(abort/cc prompt-tag v ...) → any
  prompt-tag : any/c
  v : any/c
```

The abort/cc binding is an alias for abort-current-continuation.

```
(call/comp proc [prompt-tag]) → any
  proc : (continuation? . -> . any)
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
```

The call/comp binding is an alias for call-with-composable-continuation.

```
(abort v ...) → any
  v : any/c
```

Returns the vs to a prompt using the default continuation prompt tag and the default abort handler.

That is, (abort v ...) is equivalent to

```
(abort-current-continuation
 (default-continuation-prompt-tag)
 (lambda () (values v ...)))
```

Example:

```
> (prompt
    (printf "start here\n")
    (printf "answer is ~a\n" (+ 2 (abort 3))))
start here
3
```

```
(% expr)
(% expr handler-expr)
(% expr handler-expr #:tag tag-expr)
(fcontrol v #:tag prompt-tag) → any
  v : any/c
  prompt-tag : (default-continuation-prompt-tag)
```

Sitaram's operators [Sitaram93].

The essential reduction rules are:

```
(% val proc) => val
(% E[(fcontrol val)] proc) => (proc val (lambda (x) E[x]))
  ; where E has no %
```

When *handler-expr* is omitted, % is the same as `prompt`. If *prompt-tag* is provided, % uses specific prompt tags like `prompt-at`.

Examples:

```
> (% (+ 2 (fcontrol 5))
     (lambda (v k)
       (k v)))
7
> (% (+ 2 (fcontrol 5))
     (lambda (v k)
       v))
5
```

```
(prompt expr ...+)
(control id expr ...+)
```

Among the earliest operators for higher-order control [Felleisen88a, Felleisen88, Sitaram90].

The essential reduction rules are:

```
(prompt val) => val
(prompt E[(control k expr)]) => (prompt ((lambda (k) expr)
                                         (lambda (v) E[v])))
  ; where E has no prompt
```

Examples:

```
> (prompt
    (+ 2 (control k (k 5))))
7
> (prompt
    (+ 2 (control k 5)))
5
> (prompt
    (+ 2 (control k (+ 1 (control k1 (k1 6))))))
7
> (prompt
    (+ 2 (control k (+ 1 (control k1 (k 6))))))
8
> (prompt
    (+ 2 (control k (control k1 (control k2 (k2 6))))))
6
```

```
(prompt-at prompt-tag-expr expr ...+)
(control-at prompt-tag-expr id expr ...+)
```

Like `prompt` and `control`, but using specific prompt tags:

```
(prompt-at tag val) => val
(prompt-at tag E[(control-at tag k expr)]) => (prompt-at tag
                                                ((lambda (k) expr)
                                                 (lambda (v) E[v])))
  ; where E has no prompt-at for tag
```

```
(reset expr ...+)
(shift id expr ...+)
```

Danvy and Filinski's operators [Danvy90].

The essential reduction rules are:

```
(reset val) => val
(reset E[(shift k expr)]) => (reset ((lambda (k) expr)
                                     (lambda (v) (reset E[v]))))
  ; where E has no reset
```

The `reset` and `prompt` forms are interchangeable.

```
(reset-at prompt-tag-expr expr ...+)
(shift-at prompt-tag-expr identifier expr ...+)
```

Like `reset` and `shift`, but using the specified prompt tags.

```
(prompt0 expr ...+)
(reset0 expr ...+)
(control0 id expr ...+)
(shift0 id expr ...+)
```

Generalizations of `prompt`, etc. [Shan04].

The essential reduction rules are:

```
(prompt0 val) => val
(prompt0 E[(control0 k expr)]) => ((lambda (k) expr)
                                   (lambda (v) E[v]))
(reset0 val) => val
(reset0 E[(shift0 k expr)]) => ((lambda (k) expr)
                                (lambda (v) (reset0 E[v])))
```

The `reset0` and `prompt0` forms are interchangeable. Furthermore, the following reductions apply:

```
(prompt E[(control0 k expr)]) => (prompt ((lambda (k) expr)
                                          (lambda (v) E[v])))
(reset E[(shift0 k expr)]) => (reset ((lambda (k) expr)
                                      (lambda (v) (reset0 E[v]))))
(prompt0 E[(control k expr)]) => (prompt0 ((lambda (k) expr)
                                           (lambda (v) E[v])))
(reset0 E[(shift k expr)]) => (reset0 ((lambda (k) expr)
                                       (lambda (v) (reset E[v]))))
```

That is, both the `prompt`/`reset` and `control`/`shift` sites must agree for `0`-like behavior, otherwise the non-`0` behavior applies.

```
(prompt0-at prompt-tag-expr expr ...+)
(reset0-at prompt-tag-expr expr ...+)
(control0-at prompt-tag-expr id expr ...+)
(shift0-at prompt-tag-expr id expr ...+)
```

Variants of `prompt0`, etc., that accept a prompt tag.

```
(spawn proc) → any
  proc : ((any/c . -> . any) . -> . any)
```

The operators of Hieb and Dybvig [Hieb90].

The essential reduction rules are:

```
(prompt-at tag obj) => obj
(spawn proc) => (prompt tag (proc (lambda (x) (abort tag x))))
(prompt-at tag E[(abort tag proc)])
  => (proc (lambda (x) (prompt-at tag E[x])))
  ; where E has no prompt-at for tag
```

```
(splitter proc) → any
        (((-> any) . -> . any)
 proc :  ((continuation? . -> . any) . -> . any)
         . -> . any)
```

The operator of Queinnec and Serpette [Queinnec91].

The essential reduction rules are:

```
(splitter proc) => (prompt-at tag
                     (proc (lambda (thunk)
                             (abort tag thunk))
                           (lambda (proc)
                             (control0-at tag k (proc k)))))
(prompt-at tag E[(abort tag thunk)]) => (thunk)
  ; where E has no prompt-at for tag
(prompt-at tag E[(control0-at tag k expr)]) => ((lambda (k) expr)
                                                (lambda (x) E[x]))
  ; where E has no prompt-at for tag
```

```
(new-prompt) → any
(set prompt-expr expr ...+)
(cupto prompt-expr id expr ...+)
```

The operators of Gunter et al. [Gunter95].

In this library, `new-prompt` is an alias for `make-continuation-prompt-tag`, `set` is an alias for `prompt0-at`, and `cupto` is an alias for `control0-at`.

10.5 Continuation Marks

See §1.1.11 “Continuation Frames and Marks” and §1.1.12 “Prompts, Delimited Continuations, and Barriers” for general information about continuation marks.

The list of continuation marks for a key k and a continuation C that extends C_0 is defined as follows:

- If C is an empty continuation, then the mark list is `null`.
- If C's first frame contains a mark m for k, then the mark list for C is `(cons m lst)`, where `lst` is the mark list for k in C_0.
- If C's first frame does not contain a mark keyed by k, then the mark list for C is the mark list for C_0.

The `with-continuation-mark` form installs a mark on the first frame of the current continuation (see §3.19 "Continuation Marks: `with-continuation-mark`"). Procedures such as `current-continuation-marks` allow inspection of marks.

Whenever Racket creates an exception record for a primitive exception, it fills the `continuation-marks` field with the value of `(current-continuation-marks)`, thus providing a snapshot of the continuation marks at the time of the exception.

When a continuation procedure returned by `call-with-current-continuation` or `call-with-composable-continuation` is invoked, it restores the captured continuation, and also restores the marks in the continuation's frames to the marks that were present when `call-with-current-continuation` or `call-with-composable-continuation` was invoked.

```
(continuation-marks cont [prompt-tag]) → continuation-mark-set?
  cont : (or/c continuation? thread? #f)
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
```

Returns an opaque value containing the set of continuation marks for all keys in the continuation *cont* (or the current continuation of *cont* if it is a thread) up to the prompt tagged by *prompt-tag*. If *cont* is `#f`, the resulting set of continuation marks is empty. If *cont* is an escape continuation (see §1.1.12 "Prompts, Delimited Continuations, and Barriers"), then the current continuation must extend *cont*, or the `exn:fail:contract` exception is raised. If *cont* was not captured with respect to *prompt-tag* and does not include a prompt for *prompt-tag*, the `exn:fail:contract` exception is raised. If *cont* is a dead thread, the result is an empty set of continuation marks.

```
(current-continuation-marks [prompt-tag])
 → continuation-mark-set?
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
```

Returns an opaque value containing the set of continuation marks for all keys in the current continuation up to *prompt-tag*. In other words, it produces the same value as

```
(call-with-current-continuation
  (lambda (k)
    (continuation-marks k prompt-tag))
  prompt-tag)
```

```
(continuation-mark-set->list  mark-set
                              key-v
                             [prompt-tag]) → list?
  mark-set : continuation-mark-set?
  key-v : any/c
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
```

Returns a newly-created list containing the marks for *key-v* in *mark-set*, which is a set of marks returned by `current-continuation-marks`. The result list is truncated at the first point, if any, where continuation frames were originally separated by a prompt tagged with *prompt-tag*.

```
(make-continuation-mark-key) → continuation-mark-key?
(make-continuation-mark-key sym) → continuation-mark-key?
  sym : symbol?
```

Creates a continuation mark key that is not `equal?` to the result of any other value (including prior and future results from `make-continuation-mark-key`). The continuation mark key can be used as the key argument for `with-continuation-mark` or accessor procedures like `continuation-mark-set-first`. The mark key can be chaperoned or impersonated, unlike other values that are used as the mark key.

The optional *sym* argument, if provided, is used when printing the continuation mark.

```
(continuation-mark-set->list*  mark-set
                               key-list
                              [none-v
                               prompt-tag]) → (listof vector?)
  mark-set : continuation-mark-set?
  key-list : (listof any/c)
  none-v : any/c = #f
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
```

Returns a newly-created list containing vectors of marks in *mark-set* for the keys in *key-list*, up to *prompt-tag*. The length of each vector in the result list is the same as the length of *key-list*, and a value in a particular vector position is the value for the corresponding key in *key-list*. Values for multiple keys appear in a single vector only when the marks are for the same continuation frame in *mark-set*. The *none-v* argument is used for vector elements to indicate the lack of a value.

```
(continuation-mark-set-first  mark-set
                              key-v
                             [none-v
                              prompt-tag]) → any
```

```
  mark-set : (or/c continuation-mark-set? #f)
  key-v : any/c
  none-v : any/c = #f
  prompt-tag : continuation-prompt-tag?
             = (default-continuation-prompt-tag)
```

Returns the first element of the list that would be returned by `(continuation-mark-set->list (or mark-set (current-continuation-marks prompt-tag)) key-v prompt-tag)`, or `none-v` if the result would be the empty list. Typically, this result can be computed more quickly using `continuation-mark-set-first` than using `continuation-mark-set->list`.

```
(call-with-immediate-continuation-mark key-v
                                       proc
                                       [default-v]) → any
  key-v : any/c
  proc : (any/c . -> . any)
  default-v : any/c = #f
```

Calls *proc* with the value associated with *key-v* in the first frame of the current continuation (i.e., a value that would be replaced if the call to `call-with-immediate-continuation-mark` were replaced with a `with-continuation-mark` form using *key-v* as the key expression). If no such value exists in the first frame, *default-v* is passed to *proc*. The *proc* is called in tail position with respect to the `call-with-immediate-continuation-mark` call.

This function could be implemented with a combination of `with-continuation-mark`, `current-continuation-marks`, and `continuation-mark-set->list`, but `call-with-immediate-continuation-mark` is implemented more efficiently; it inspects only the first frame of the current continuation.

```
(continuation-mark-key? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a mark key created by `make-continuation-mark-key`, `#f` otherwise.

```
(continuation-mark-set? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a mark set created by `continuation-marks` or `current-continuation-marks`, `#f` otherwise.

```
(continuation-mark-set->context mark-set) → list?
  mark-set : continuation-mark-set?
```

Returns a list representing an approximate "stack trace" for *mark-set*'s continuation. The list contains pairs, where the car of each pair contains either #f or a symbol for a procedure name, and the cdr of each pair contains either #f or a srcloc value for the procedure's source location (see §13.1.4 "Counting Positions, Lines, and Columns"); the car and cdr are never both #f.

Conceptually, the stack-trace list is the result of continuation-mark-set->list with *mark-set* and Racket's private key for procedure-call marks. The implementation may be different, however, and the results may merely approximate the correct answer. Thus, while the result may contain useful hints to humans about the context of an expression, it is not reliable enough for programmatic use.

A stack trace is extracted from an exception and displayed by the default error display handler (see error-display-handler) for exceptions other than exn:fail:user (see raise-user-error in §10.2.2 "Raising Exceptions").

Examples:

```
> (define (extract-current-continuation-marks key)
    (continuation-mark-set->list
     (current-continuation-marks)
     key))

> (with-continuation-mark 'key 'mark
    (extract-current-continuation-marks 'key))
'(mark)
> (with-continuation-mark 'key1 'mark1
    (with-continuation-mark 'key2 'mark2
      (list
       (extract-current-continuation-marks 'key1)
       (extract-current-continuation-marks 'key2))))
'((mark1) (mark2))
> (with-continuation-mark 'key 'mark1
    (with-continuation-mark 'key 'mark2 ; replaces previous mark
      (extract-current-continuation-marks 'key)))
'(mark2)
> (with-continuation-mark 'key 'mark1
    (list ; continuation extended to evaluate the argument
     (with-continuation-mark 'key 'mark2
        (extract-current-continuation-marks 'key))))
'((mark2 mark1))
> (let loop ([n 1000])
    (if (zero? n)
        (extract-current-continuation-marks 'key)
        (with-continuation-mark 'key n
          (loop (sub1 n)))))
'(1)
```

10.6 Breaks

A *break* is an asynchronous exception, usually triggered through an external source controlled by the user, or through the break-thread procedure. For example, the user may type Ctl-C in a terminal to trigger a break. On some platforms, the Racket process may receive SIGINT, SIGHUP, or SIGTERM; the latter two correspond to hang-up and terminate breaks as reflected by exn:break:hang-up and exn:break:terminate, respectively. Multiple breaks may be collapsed into a single exception, and multiple breaks of different kinds may be collapsed to a single "strongest" break, where a hang-up break is stronger than an interrupt break, and a terminate break is stronger than a hang-up break.

A break exception can only occur in a thread while breaks are enabled. When a break is detected and enabled, the exn:break (or exn:break:hang-up or exn:break:terminate) exception is raised in the thread sometime afterward; if breaking is disabled when break-thread is called, the break is suspended until breaking is again enabled for the thread. While a thread has a suspended break, additional breaks are ignored.

Breaks are enabled through the break-enabled parameter-like procedure and through the parameterize-break form, which is analogous to parameterize. The break-enabled procedure does not represent a parameter to be used with parameterize, because changing the break-enabled state of a thread requires an explicit check for breaks, and this check is incompatible with the tail evaluation of a parameterize expression's body.

Certain procedures, such as semaphore-wait/enable-break, enable breaks temporarily while performing a blocking action. If breaks are enabled for a thread, and if a break is triggered for the thread but not yet delivered as an exn:break exception, then the break is guaranteed to be delivered before breaks can be disabled in the thread. The timing of exn:break exceptions is not guaranteed in any other way.

Before calling a with-handlers predicate or handler, an exception handler, an error display handler, an error escape handler, an error value conversion handler, or a pre-thunk or post-thunk for a dynamic-wind, the call is parameterize-breaked to disable breaks. Furthermore, breaks are disabled during the transitions among handlers related to exceptions, during the transitions between pre-thunks and post-thunks for dynamic-wind, and during other transitions for a continuation jump. For example, if breaks are disabled when a continuation is invoked, and if breaks are also disabled in the target continuation, then breaks will remain disabled from the time of the invocation until the target continuation executes unless a relevant dynamic-wind pre-thunk or post-thunk explicitly enables breaks.

If a break is triggered for a thread that is blocked on a nested thread (see call-in-nested-thread), and if breaks are enabled in the blocked thread, the break is implicitly handled by transferring it to the nested thread.

When breaks are enabled, they can occur at any point within execution, which makes certain implementation tasks subtle. For example, assuming breaks are enabled when the following

code is executed,

```
(with-handlers ([exn:break? (lambda (x) (void))])
  (semaphore-wait s))
```

then it is *not* the case that a `#<void>` result means the semaphore was decremented or a break was received, exclusively. It is possible that *both* occur: the break may occur after the semaphore is successfully decremented but before a `#<void>` result is returned by `semaphore-wait`. A break exception will never damage a semaphore, or any other built-in construct, but many built-in procedures (including `semaphore-wait`) contain internal sub-expressions that can be interrupted by a break.

In general, it is impossible using only `semaphore-wait` to implement the guarantee that either the semaphore is decremented or an exception is raised, but not both. Racket therefore supplies `semaphore-wait/enable-break` (see §11.2.3 "Semaphores"), which does permit the implementation of such an exclusive guarantee:

```
(parameterize-break #f
  (with-handlers ([exn:break? (lambda (x) (void))])
    (semaphore-wait/enable-break s)))
```

In the above expression, a break can occur at any point until breaks are disabled, in which case a break exception is propagated to the enclosing exception handler. Otherwise, the break can only occur within `semaphore-wait/enable-break`, which guarantees that if a break exception is raised, the semaphore will not have been decremented.

To allow similar implementation patterns over blocking port operations, Racket provides `read-bytes-avail!/enable-break`, `write-bytes-avail/enable-break`, and other procedures.

```
(break-enabled) → boolean?
(break-enabled on?) → void?
  on? : any/c
```

Gets or sets the break enabled state of the current thread. If *on?* is not supplied, the result is `#t` if breaks are currently enabled, `#f` otherwise. If *on?* is supplied as `#f`, breaks are disabled, and if *on?* is a true value, breaks are enabled.

```
(parameterize-break boolean-expr body ...+)
```

Evaluates *boolean-expr* to determine whether breaks are initially enabled while evaluating the *body*s in sequence. The result of the `parameter-break` expression is the result of the last `expr`.

Like `parameterize` (see §11.3.2 "Parameters"), a fresh thread cell (see §11.3.1 "Thread Cells") is allocated to hold the break-enabled state of the continuation, and calls to `break-enabled` within the continuation access or modify the new cell. Unlike parameters, the break setting is not inherited by new threads.

```
(current-break-parameterization) → break-parameterization?
```

Analogous to (current-parameterization) (see §11.3.2 “Parameters”); it returns a break-parameterization (effectively, a thread cell) that holds the current continuation’s break-enable state.

```
(call-with-break-parameterization break-param
                                  thunk)      → any
  break-param : break-parameterization?
  thunk : (-> any)
```

Analogous to (call-with-parameterization parameterization *thunk*) (see §11.3.2 “Parameters”), calls *thunk* in a continuation whose break-enabled state is in *break-param*. The *thunk* is *not* called in tail position with respect to the call-with-break-parameterization call.

10.7 Exiting

```
(exit [v]) → any
  v : any/c = #t
```

Passes *v* to the current exit handler. If the exit handler does not escape or terminate the thread, #<void> is returned.

```
(exit-handler) → (any/c . -> . any)
(exit-handler proc) → void?
  proc : (any/c . -> . any)
```

A parameter that determines the current *exit handler*. The exit handler is called by exit.

The default exit handler in the Racket executable takes any argument, calls plumber-flush-all on the original plumber, and shuts down the OS-level Racket process. The argument is used as the OS-level exit code if it is an exact integer between 1 and 255 (which normally means “failure”); otherwise, the exit code is 0, (which normally means “success”).

```
(executable-yield-handler) → (byte? . -> . any)
(executable-yield-handler proc) → void?
  proc : (byte? . -> . any)
```

A parameter that determines a procedure to be called as the Racket process is about to exit normally. The procedure associated with this parameter is not called when exit (or, more precisely, the default exit handler) is used to exit early. The argument to the handler is the

status code that is returned to the system on exit. The default executable-yield handler simply returns `#<void>`.

The `scheme/gui/base` library sets this parameter to wait until all frames are closed, timers stopped, and queued events handled in the main eventspace. See `scheme/gui/base` for more information.

11 Concurrency and Parallelism

Racket supports multiple threads of control within a program, thread-local storage, some primitive synchronization mechanisms, and a framework for composing synchronization abstractions. In addition, the `racket/future` and `racket/place` libraries provide support for parallelism to improve performance.

11.1 Threads

§18 "Concurrency and Synchronization" in ***The Racket Guide*** introduces threads.

See §1.1.13 "Threads" for basic information on the Racket thread model. See also §11.4 "Futures".

When a thread is created, it is placed into the management of the current custodian and added to the current thread group. A thread can have any number of custodian managers added through `thread-resume`.

A thread that has not terminated can be garbage collected (see §1.1.7 "Garbage Collection") if it is unreachable and suspended or if it is unreachable and blocked on only unreachable events through functions such as `semaphore-wait`, `semaphore-wait/enable-break`, `channel-put`, `channel-get`, `sync`, `sync/enable-break`, or `thread-wait`. Beware, however, of a limitation on place-channel blocking; see the caveat in §11.5 "Places".

In GRacket, a handler thread for an eventspace is blocked on an internal semaphore when its event queue is empty. Thus, the handler thread is collectible when the eventspace is unreachable and contains no visible windows or running timers.

A thread can be used as a synchronizable event (see §11.2.1 "Events"). A thread is ready for synchronization when `thread-wait` would not block; the synchronization result of a thread is the thread itself.

All constant-time procedures and operations provided by Racket are thread-safe because they are *atomic*. For example, `set!` assigns to a variable as an atomic action with respect to all threads, so that no thread can see a "half-assigned" variable. Similarly, `vector-set!` assigns to a vector atomically. The `hash-set!` procedure is not atomic, but the table is protected by a lock; see §4.13 "Hash Tables" for more information. Port operations are generally not atomic, but they are thread-safe in the sense that a byte consumed by one thread from an input port will not be returned also to another thread, and procedures like `port-commit-peeked` and `write-bytes-avail` offer specific concurrency guarantees.

11.1.1 Creating Threads

```
(thread thunk) → thread?
  thunk : (-> any)
```

Calls *thunk* with no arguments in a new thread of control. The `thread` procedure returns immediately with a *thread descriptor* value. When the invocation of *thunk* returns, the thread created to invoke *thunk* terminates.

```
(thread? v) → thread?
  v : any/c
```

Returns #t if *v* is a thread descriptor, #f otherwise.

```
(current-thread) → thread?
```

Returns the thread descriptor for the currently executing thread.

```
(thread/suspend-to-kill thunk) → thread?
  thunk : (-> any)
```

Like thread, except that "killing" the thread through kill-thread or custodian-shutdown-all merely suspends the thread instead of terminating it.

```
(call-in-nested-thread thunk [cust]) → any
  thunk : (-> any)
  cust : custodian? = (current-custodian)
```

Creates a nested thread managed by *cust* to execute *thunk*. (The nested thread's current custodian is inherited from the creating thread, independent of the *cust* argument.) The current thread blocks until *thunk* returns, and the result of the call-in-nested-thread call is the result returned by *thunk*.

The nested thread's exception handler is initialized to a procedure that jumps to the beginning of the thread and transfers the exception to the original thread. The handler thus terminates the nested thread and re-raises the exception in the original thread.

If the thread created by call-in-nested-thread dies before *thunk* returns, the exn:fail exception is raised in the original thread. If the original thread is killed before *thunk* returns, a break is queued for the nested thread.

If a break is queued for the original thread (with break-thread) while the nested thread is running, the break is redirected to the nested thread. If a break is already queued on the original thread when the nested thread is created, the break is moved to the nested thread. If a break remains queued on the nested thread when it completes, the break is moved to the original thread.

11.1.2 Suspending, Resuming, and Killing Threads

```
(thread-suspend thd) → void?
  thd : thread?
```

Immediately suspends the execution of *thd* if it is running. If the thread has terminated or is already suspended, `thread-suspend` has no effect. The thread remains suspended (i.e., it does not execute) until it is resumed with `thread-resume`. If the current custodian does not solely manage *thd* (i.e., some custodian of *thd* is not the current custodian or a subordinate), the `exn:fail:contract` exception is raised, and the thread is not suspended.

```
(thread-resume thd [benefactor]) → void?
  thd : thread?
  benefactor : (or/c thread? custodian? #f) = #f
```

Resumes the execution of *thd* if it is suspended and has at least one custodian (possibly added through *benefactor*, as described below). If the thread has terminated, or if the thread is already running and *benefactor* is not supplied, or if the thread has no custodian and *benefactor* is not supplied, then `thread-resume` has no effect. Otherwise, if *benefactor* is supplied, it triggers up to three additional actions:

- If *benefactor* is a thread, whenever it is resumed from a suspended state in the future, then *thd* is also resumed. (Resuming *thd* may trigger the resumption of other threads that were previously attached to *thd* through `thread-resume`.)
- New custodians may be added to *thd*'s set of managers. If *benefactor* is a thread, then all of the thread's custodians are added to *thd*. Otherwise, *benefactor* is a custodian, and it is added to *thd* (unless the custodian is already shut down). If *thd* becomes managed by both a custodian and one or more of its subordinates, the redundant subordinates are removed from *thd*. If *thd* is suspended and a custodian is added, then *thd* is resumed only after the addition.
- If *benefactor* is a thread, whenever it receives a new managing custodian in the future, then *thd* also receives the custodian. (Adding custodians to *thd* may trigger adding the custodians to other threads that were previously attached to *thd* through `thread-resume`.)

```
(kill-thread thd) → void?
  thd : thread?
```

Terminates the specified thread immediately, or suspends the thread if *thd* was created with `thread/suspend-to-kill`. Terminating the main thread exits the application. If *thd* has already terminated, `kill-thread` does nothing. If the current custodian does not manage *thd* (and none of its subordinates manages *thd*), the `exn:fail:contract` exception is raised, and the thread is not killed or suspended.

Unless otherwise noted, procedures provided by Racket (and GRacket) are kill-safe and suspend-safe; that is, killing or suspending a thread never interferes with the application of procedures in other threads. For example, if a thread is killed while extracting a character from an input port, the character is either completely consumed or not consumed, and other threads can safely use the port.

```
(break-thread thd [kind]) → void?
  thd : thread?
  kind : (or/c #f 'hang-up 'terminate) = #f
```

Registers a break with the specified thread, where *kind* optionally indicates the kind of break to register. If breaking is disabled in *thd*, the break will be ignored until breaks are re-enabled (see §10.6 “Breaks”).

```
(sleep [secs]) → void?
  secs : (>=/c 0) = 0
```

Causes the current thread to sleep until at least *secs* seconds have passed after it starts sleeping. A zero value for *secs* simply acts as a hint to allow other threads to execute. The value of *secs* can be a non-integer to request a sleep duration to any precision; the precision of the actual sleep time is unspecified.

```
(thread-running? thd) → any
  thd : thread?
```

Returns `#t` if *thd* has not terminated and is not suspended, `#f` otherwise.

```
(thread-dead? thd) → any
  thd : thread?
```

Returns `#t` if *thd* has terminated, `#f` otherwise.

11.1.3 Synchronizing Thread State

```
(thread-wait thd) → void?
  thd : thread?
```

Blocks execution of the current thread until *thd* has terminated. Note that `(thread-wait (current-thread))` deadlocks the current thread, but a break can end the deadlock (if breaking is enabled; see §10.6 “Breaks”).

```
(thread-dead-evt thd) → evt?
  thd : thread?
```

Returns a synchronizable event (see §11.2.1 “Events”) that is ready for synchronization if and only if *thd* has terminated. Unlike using *thd* directly, however, a reference to the event does not prevent *thd* from being garbage collected (see §1.1.7 “Garbage Collection”). For a given *thd*, `thread-dead-evt` always returns the same (i.e., `eq?`) result. The synchronization result of a thread-dead event is the thread-dead event itself.

```
(thread-resume-evt thd) → evt?
  thd : thread?
```

Returns a synchronizable event (see §11.2.1 “Events”) that becomes ready for synchronization when *thd* is running. (If *thd* has terminated, the event never becomes ready.) If *thd* runs and is then suspended after a call to `thread-resume-evt`, the result event remains ready; after each suspend of *thd* a fresh event is generated to be returned by `thread-resume-evt`. The result of the event is *thd*, but if *thd* is never resumed, then reference to the event does not prevent *thd* from being garbage collected (see §1.1.7 “Garbage Collection”). The synchronization result of a thread-result event is the thread-result event itself.

```
(thread-suspend-evt thd) → evt?
  thd : thread?
```

Returns a synchronizable event (see §11.2.1 “Events”) that becomes ready for synchronization when *thd* is suspended. (If *thd* has terminated, the event will never unblock.) If *thd* is suspended and then resumes after a call to `thread-suspend-evt`, the result event remains ready; after each resume of *thd* created a fresh event to be returned by `thread-suspend-evt`. The synchronization result of a thread-suspend event is the thread-suspend event itself.

11.1.4 Thread Mailboxes

Each thread has a *mailbox* through which it can receive arbitrary messages. In other words, each thread has a built-in asynchronous channel.

See also §11.2.4 “Buffered Asynchronous Channels”.

```
(thread-send thd v [fail-thunk]) → any
  thd : thread?
  v : any/c
  fail-thunk : (or/c (-> any) #f)
             = (lambda () (raise-mismatch-error ....))
```

Queues *v* as a message to *thd* without blocking. If the message is queued, the result is `#<void>`. If *thd* stops running—as in `thread-running?`—before the message is queued, then *fail-thunk* is called (through a tail call) if it is a procedure to produce the result, or `#f` is returned if *fail-thunk* is `#f`.

```
(thread-receive) → any/c
```

Receives and dequeues a message queued for the current thread, if any. If no message is available, `thread-receive` blocks until one is available.

```
(thread-try-receive) → any/c
```

Receives and dequeues a message queued for the current thread, if any, or returns `#f` immediately if no message is available.

```
(thread-receive-evt) → evt?
```

Returns a constant synchronizable event (see §11.2.1 “Events”) that becomes ready for synchronization when the synchronizing thread has a message to receive. The synchronization result of a thread-receive event is the thread-receive event itself.

```
(thread-rewind-receive lst) → void?
  lst : list?
```

Pushes the elements of *lst* back onto the front of the current thread’s queue. The elements are pushed one by one, so that the first available message is the last element of *lst*.

11.2 Synchronization

Racket’s synchronization toolbox spans three layers:

- synchronizable events — a general framework for synchronization;
- channels — a primitive that can be used, in principle, to build most other kinds of synchronizable events (except the ones that compose events); and
- semaphores — a simple and especially cheap primitive for synchronization.

11.2.1 Events

A *synchronizable event* (or just *event* for short) works with the `sync` procedure to coordinate synchronization among threads. Certain kinds of objects double as events, including ports and threads. Other kinds of objects exist only for their use as events.

At any point in time, an event is either *ready for synchronization*, or it is not; depending on the kind of event and how it is used by other threads, an event can switch from not ready to ready (or back), at any time. If a thread synchronizes on an event when it is ready, then the event produces a particular *synchronization result*.

Synchronizing an event may affect the state of the event. For example, when synchronizing a semaphore, then the semaphore’s internal count is decremented, just as with `semaphore-wait`. For most kinds of events, however (such as a port), synchronizing does not modify the event’s state.

Racket values that act as synchronizable events include semaphores, channels, asynchronous channels, ports, TCP listeners, log receivers, threads, subprocesses, will executors, and custodian boxes. Libraries can define new synchronizable events, especially though `prop:evt`.

```
(evt? v) → boolean?
  v : any/c
```

Returns #t if v is a synchronizable event, #f otherwise.

Examples:

```
> (evt? never-evt)
#t
> (evt? (make-channel))
#t
> (evt? 5)
#f
```

```
(sync evt ...) → any
  evt : evt?
```

Blocks as long as none of the synchronizable events evts are ready, as defined above.

When at least one evt is ready, its synchronization result (often evt itself) is returned. If multiple evts are ready, one of the evts is chosen pseudo-randomly for the result; the current-evt-pseudo-random-generator parameter sets the random-number generator that controls this choice.

Examples:

```
> (define ch (make-channel))

> (thread (λ () (displayln (sync ch))))
#<thread>
> (channel-put ch 'hellooooo)
hellooooo
```

Changed in version 6.1.0.3 of package base: Allow 0 arguments instead of 1 or more.

```
(sync/timeout timeout evt ...) → any
  timeout : (or/c #f (and/c real? (not/c negative?)) (-> any))
  evt : evt?
```

Like sync if timeout is #f. If timeout is a real number, then the result is #f if timeout seconds pass without a successful synchronization. If timeout is a procedure, then it is called in tail position if polling the evts discovers no ready events.

A zero value for timeout is equivalent to (lambda () #f). In either case, each evt is checked at least once before returning #f or calling timeout.

See also alarm-evt for an alternative timeout mechanism.

Examples:

```
; times out before waking up
> (sync/timeout
   0.5
   (thread (λ () (sleep 1) (displayln "woke up!"))))
#f
> (sync/timeout
   (λ () (displayln "no ready events"))
   never-evt)
no ready events
```

Changed in version 6.1.0.3 of package `base`: Allow 1 argument instead of 2 or more.

```
(sync/enable-break evt ...) → any
  evt : evt?
```

Like `sync`, but breaking is enabled (see §10.6 “Breaks”) while waiting on the *evt*s. If breaking is disabled when `sync/enable-break` is called, then either all *evt*s remain unchosen or the `exn:break` exception is raised, but not both.

```
(sync/timeout/enable-break timeout evt ...) → any
  timeout : (or/c #f (and/c real? (not/c negative?)) (-> any))
  evt : evt?
```

Like `sync/enable-break`, but with a timeout as for `sync/timeout`.

```
(choice-evt evt ...) → evt?
  evt : evt?
```

Creates and returns a single event that combines the *evt*s. Supplying the result to `sync` is the same as supplying each *evt* to the same call.

That is, an event returned by `choice-evt` is ready for synchronization when one or more of the *evt*s supplied to `choice-evt` are ready for synchronization. If the choice event is chosen, one of its ready *evt*s is chosen pseudo-randomly, and the synchronization result is the chosen *evt*’s synchronization result.

Examples:

```
> (define ch1 (make-channel))

> (define ch2 (make-channel))

> (define either-channel (choice-evt ch1 ch2))
```

```
> (thread (λ () (displayln (sync either-channel))))
#<thread>
> (channel-put
   (if (> (random) 0.5) ch1 ch2)
   'tuturuu)
tuturuu
```

```
(wrap-evt evt wrap) → evt?
  evt : evt?
  wrap : (any/c ... . -> . any)
```

Creates an event that is ready for synchronization when *evt* is ready for synchronization, but whose synchronization result is determined by applying *wrap* to the synchronization result of *evt*. The number of arguments accepted by *wrap* must match the number of values for the synchronization result of *evt*.

The call to *wrap* is `parameterize-break`ed to disable breaks initially.

Examples:

```
> (define ch (make-channel))

> (define evt (wrap-evt ch (λ (v) (format "you've got mail:
~a" v))))

> (thread (λ () (displayln (sync evt))))
#<thread>
> (channel-put ch "Dear Alice ...")
you've got mail: Dear Alice ...
```

```
(handle-evt evt handle) → handle-evt?
  evt : evt?
  handle : (any/c ... . -> . any)
```

Like `wrap-evt`, except that *handle* is called in tail position with respect to the synchronization request—and without breaks explicitly disabled—when it is not wrapped by `wrap-evt`, `chaperone-evt`, or another `handle-evt`.

Examples:

```
> (define msg-ch (make-channel))

> (define exit-ch (make-channel))
```

```
> (thread
   (λ ()
     (let loop ([val 0])
       (printf "val = ~a~n" val)
       (sync (handle-evt
              msg-ch
              (λ (val) (loop val)))
             (handle-evt
              exit-ch
              (λ (val) (displayln val)))))))
val = 0
#<thread>
> (channel-put msg-ch 5)
val = 5

> (channel-put msg-ch 7)
val = 7

> (channel-put exit-ch 'done)
done
```

```
(guard-evt maker) → evt?
  maker : (-> (or/c evt? any/c))
```

Creates a value that behaves as an event, but that is actually an event maker.

An event *guard* returned by guard-evt generates an event when *guard* is used with sync (or whenever it is part of a choice event used with sync, etc.), where the generated event is the result of calling *maker*. The *maker* procedure may be called by sync at most once for a given call to sync, but *maker* may not be called if a ready event is chosen before *guard* is even considered.

If *maker* returns a non-event, then *maker*'s result is replaced with an event that is ready for synchronization and whose synchronization result is *guard*.

```
(nack-guard-evt maker) → evt?
  maker : (evt? . -> . (or/c evt? any/c))
```

Like guard-evt, but when *maker* is called, it is given a NACK ("negative acknowledgment") event. After starting the call to *maker*, if the event from *maker* is not ultimately chosen as the ready event, then the NACK event supplied to *maker* becomes ready for synchronization with a #<void> value.

The NACK event becomes ready for synchronization when the event is abandoned when either some other event is chosen, the synchronizing thread is dead, or control escapes from

the call to sync (even if nack-guard's maker has not yet returned a value). If the event returned by maker is chosen, then the NACK event never becomes ready for synchronization.

```
(poll-guard-evt maker) → evt?
  maker : (boolean? . -> . (or/c evt? any/c))
```

Like guard-evt, but when maker is called, it is provided a boolean value that indicates whether the event will be used for a poll, #t, or for a blocking synchronization, #f.

If #t is supplied to maker, if breaks are disabled, if the polling thread is not terminated, and if polling the resulting event produces a synchronization result, then the event will certainly be chosen for its result.

```
(replace-evt evt maker) → evt?
  evt : evt?
  maker : (any/c ... . -> . (or/c evt? any/c))
```

Like guard-evt, but maker is called only after evt becomes ready for synchronization, and the synchronization result of evt is passed to maker.

The attempt to synchronize on evt proceeds concurrently as the attempt to synchronize on the result guard from replace-evt; despite that concurrency, if maker is called, it is called in the thread that is synchronizing on guard. Synchronization can succeed for both evt and another synchronized with guard at the same time; the single-choice guarantee of synchronization applies only to the result of maker and other events synchronized with guard.

If maker returns a non-event, then maker's result is replaced with an event that is ready for synchronization and whose synchronization result is guard.

Added in version 6.1.0.3 of package base.

```
always-evt : evt?
```

A constant event that is always ready for synchronization, with itself as its synchronization result.

Example:

```
> (sync always-evt)
#<always-evt>
```

```
never-evt : evt?
```

A constant event that is never ready for synchronization.

Example:

```
> (sync/timeout 0.1 never-evt)
#f
```

(system-idle-evt) → evt?

Returns an event that is ready for synchronization when the system is otherwise idle: if the result event were replaced by never-evt, no thread in the system would be available to run. In other words, all threads must be suspended or blocked on events with timeouts that have not yet expired. The system-idle event's synchronization result is #<void>. The result of the system-idle-evt procedure is always the same event.

Examples:

```
> (define th (thread (λ () (let loop () (loop)))))

> (sync/timeout 0.1 (system-idle-evt))
#f
> (kill-thread th)

> (sync (system-idle-evt))
```

(alarm-evt *msecs*) → evt?
  msecs : real?

Returns a synchronizable event that is not ready for synchronization when (current-inexact-milliseconds) would return a value that is less than *msecs*, and it is ready for synchronization when (current-inexact-milliseconds) would return a value that is more than *msecs*. The synchronization result of a alarm event is the alarm event itself.

Examples:

```
> (define alarm (alarm-evt (+ (current-inexact-
milliseconds) 100)))

> (sync alarm)
#<alarm-evt>
```

(handle-evt? *evt*) → boolean?
  evt : evt?

Returns #t if *evt* was created by handle-evt or by choice-evt applied to another event for which handle-evt? produces #t. For any other event, handle-evt? produces #f.

Examples:

```
> (handle-evt? never-evt)
#f
> (handle-evt? (handle-evt always-evt values))
#t
```

```
prop:evt : struct-type-property?
```

A structure type property that identifies structure types whose instances can serve as synchronizable events. The property value can be any of the following:

- An event *evt*: In this case, using the structure as an event is equivalent to using *evt*.
- A procedure *proc* of one argument: In this case, the structure is similar to an event generated by guard-evt, except that the would-be guard procedure *proc* receives the structure as an argument, instead of no arguments; also, a non-event result from *proc* is replaced with an event that is already ready for synchronization and whose synchronization result is the structure.
- An exact, non-negative integer between 0 (inclusive) and the number of non-automatic fields in the structure type (exclusive, not counting supertype fields): The integer identifies a field in the structure, and the field must be designated as immutable. If the field contains an object or an event-generating procedure of one argument, the event or procedure is used as above. Otherwise, the structure acts as an event that is never ready.

Instances of a structure type with the prop:input-port or prop:output-port property are also synchronizable events by virtue of being a port. If the structure type has more than one of prop:evt, prop:input-port, and prop:output-port, then the prop:evt value (if any) takes precedence for determining the instance's behavior as an event, and the prop:input-port property takes precedence over prop:output-port for synchronization.

Examples:

```
> (define-struct wt (base val)
                 #:property prop:evt (struct-field-index base))

> (define sema (make-semaphore))

> (sync/timeout 0 (make-wt sema #f))
#f
> (semaphore-post sema)

> (sync/timeout 0 (make-wt sema #f))
#<semaphore>
> (semaphore-post sema)
```

```
> (sync/timeout 0 (make-wt (lambda (self) (wt-val self)) sema))
#<semaphore>
> (semaphore-post sema)

> (define my-wt (make-wt (lambda (self) (wrap-evt
                                         (wt-val self)
                                         (lambda (x) self)))
                         sema))

> (sync/timeout 0 my-wt)
#<wt>
> (sync/timeout 0 my-wt)
#f
```

```
(current-evt-pseudo-random-generator)
 → pseudo-random-generator?
(current-evt-pseudo-random-generator generator) → void?
  generator : pseudo-random-generator?
```

A parameter that determines the pseudo-random number generator used by `sync` for events created by `choice-evt`.

11.2.2 Channels

A *channel* both synchronizes a pair of threads and passes a value from one to the other. Channels are synchronous; both the sender and the receiver must block until the (atomic) transaction is complete. Multiple senders and receivers can access a channel at once, but a single sender and receiver is selected for each transaction.

Channel synchronization is *fair*: if a thread is blocked on a channel and transaction opportunities for the channel occur infinitely often, then the thread eventually participates in a transaction.

In addition to its use with channel-specific procedures, a channel can be used as a synchronizable event (see §11.2.1 “Events”). A channel is ready for synchronization when `make-channel` is ready when `channel-get` would not block; the channel’s synchronization result is the same as the `channel-get` result.

For buffered asynchronous channels, see §11.2.4 “Buffered Asynchronous Channels”.

```
(channel? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a channel, `#f` otherwise.

```
(make-channel) → channel?
```

Creates and returns a new channel. The channel can be used with `channel-get`, with `channel-try-get`, or as a synchronizable event (see §11.2.1 "Events") to receive a value through the channel. The channel can be used with `channel-put` or through the result of `channel-put-evt` to send a value through the channel.

```
(channel-get ch) → any
  ch : channel?
```

Blocks until a sender is ready to provide a value through *ch*. The result is the sent value.

```
(channel-try-get ch) → any
  ch : channel?
```

Receives and returns a value from *ch* if a sender is immediately ready, otherwise returns `#f`.

```
(channel-put ch v) → void?
  ch : channel?
  v : any/c
```

Blocks until a receiver is ready to accept the value *v* through *ch*.

```
(channel-put-evt ch v) → channel-put-evt?
  ch : channel?
  v : any/c
```

Returns a fresh synchronizable event for use with `sync`. The event is ready for synchronization when `(channel-put ch v)` would not block, and the event's synchronization result is the event itself.

```
(channel-put-evt? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a channel-put event produced by `channel-put-evt`, `#f` otherwise.

11.2.3 Semaphores

A *semaphore* has an internal counter; when this counter is zero, the semaphore can block a thread's execution (through `semaphore-wait`) until another thread increments the counter (using `semaphore-post`). The maximum value for a semaphore's internal counter is platform-specific, but always at least `10000`.

A semaphore's counter is updated in a single-threaded manner, so that semaphores can be used for reliable synchronization. Semaphore waiting is *fair*: if a thread is blocked on a semaphore and the semaphore's internal value is non-zero infinitely often, then the thread is eventually unblocked.

In addition to its use with semaphore-specific procedures, a semaphore can be used as a synchronizable event (see §11.2.1 "Events"). A semaphore is ready for synchronization when `semaphore-wait` would not block; the synchronization result of a semaphore is the semaphore itself.

```
(semaphore? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a semaphore, `#f` otherwise.

```
(make-semaphore [init]) → semaphore?
  init : exact-nonnegative-integer? = 0
```

Creates and returns a new semaphore with the counter initially set to *init*. If *init* is larger than a semaphore's maximum internal counter value, the `exn:fail` exception is raised.

```
(semaphore-post sema) → void?
  sema : semaphore?
```

Increments the semaphore's internal counter and returns `#<void>`. If the semaphore's internal counter has already reached its maximum value, the `exn:fail` exception is raised.

```
(semaphore-wait sema) → void?
  sema : semaphore?
```

Blocks until the internal counter for semaphore *sema* is non-zero. When the counter is non-zero, it is decremented and `semaphore-wait` returns `#<void>`.

```
(semaphore-try-wait? sema) → boolean?
  sema : semaphore?
```

Like `semaphore-wait`, but `semaphore-try-wait?` never blocks execution. If *sema*'s internal counter is zero, `semaphore-try-wait?` returns `#f` immediately without decrementing the counter. If *sema*'s counter is positive, it is decremented and `#t` is returned.

```
(semaphore-wait/enable-break sema) → void?
  sema : semaphore?
```

Like `semaphore-wait`, but breaking is enabled (see §10.6 "Breaks") while waiting on *sema*. If breaking is disabled when `semaphore-wait/enable-break` is called, then either the semaphore's counter is decremented or the `exn:break` exception is raised, but not both.

```
(semaphore-peek-evt sema) → semaphore-peek-evt?
  sema : semaphore?
```

Creates and returns a new synchronizable event (for use with `sync`, for example) that is ready for synchronization when *sema* is ready, but synchronizing the event does not decrement *sema*'s internal count. The synchronization result of a semaphore-peek event is the semaphore-peek event itself.

```
(semaphore-peek-evt? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a semaphore wrapper produced by `semaphore-peek-evt`, `#f` otherwise.

```
(call-with-semaphore  sema
                      proc
                     [try-fail-thunk]
                      arg ...)        → any
  sema : semaphore?
  proc : procedure?
  try-fail-thunk : (or/c (-> any) #f) = #f
  arg : any/c
```

Waits on *sema* using `semaphore-wait`, calls *proc* with all *arg*s, and then posts to *sema*. A continuation barrier blocks full continuation jumps into or out of *proc* (see §1.1.12 "Prompts, Delimited Continuations, and Barriers"), but escape jumps are allowed, and *sema* is posted on escape. If *try-fail-thunk* is provided and is not `#f`, then `semaphore-try-wait?` is called on *sema* instead of `semaphore-wait`, and *try-fail-thunk* is called if the wait fails.

```
(call-with-semaphore/enable-break  sema
                                   proc
                                  [try-fail-thunk]
                                   arg ...)        → any
  sema : semaphore?
  proc : procedure?
  try-fail-thunk : (or/c (-> any) #f) = #f
  arg : any/c
```

Like `call-with-semaphore`, except that `semaphore-wait/enable-break` is used with *sema* in non-try mode. When *try-fail-thunk* is provided and not `#f`, then breaks are enabled around the use of `semaphore-try-wait?` on *sema*.

11.2.4 Buffered Asynchronous Channels

```
(require racket/async-channel)      package: base
```

The bindings documented in this section are provided by the `racket/async-channel` library, not `racket/base` or `racket`.

See also §11.1.4 “Thread Mailboxes”.

An *asynchronous channel* is like a channel, but it buffers values so that a send operation does not wait on a receive operation.

In addition to its use with procedures that are specific to asynchronous channels, an asynchronous channel can be used as a synchronizable event (see §11.2.1 “Events”). An asynchronous channel is ready for synchronization when `async-channel-get` would not block; the asynchronous channel’s synchronization result is the same as the `async-channel-get` result.

```
(async-channel? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is an asynchronous channel, `#f` otherwise.

```
(make-async-channel [limit]) → async-channel?
  limit : (or/c exact-positive-integer? #f) = #f
```

Returns an asynchronous channel with a buffer limit of *limit* items. A get operation blocks when the channel is empty, and a put operation blocks when the channel has *limit* items already. If *limit* is `#f`, the channel buffer has no limit (so a put never blocks).

```
(async-channel-get ach) → any/c
  ach : async-channel?
```

Blocks until at least one value is available in *ach*, and then returns the first of the values that were put into `async-channel`.

```
(async-channel-try-get ach) → any/c
  ach : async-channel?
```

If at least one value is immediately available in *ach*, returns the first of the values that were put into *ach*. If `async-channel` is empty, the result is `#f`.

```
(async-channel-put ach v) → void?
  ach : async-channel?
  v : any/c
```

Puts *v* into *ach*, blocking if *ach*’s buffer is full until space is available.

```
(async-channel-put-evt ach v) → evt?
  ach : async-channel?
  v : any/c
```

Returns a synchronizable event that is ready for synchronization when `(async-channel-put ach v)` would return a value (i.e., when the channel holds fewer values already than its limit); the synchronization result of a asychronous channel-put event is the asychronous channel-put event itself.

Examples:

```
(define (server input-channel output-channel)
  (thread (lambda ()
            (define (get)
              (async-channel-get input-channel))
            (define (put x)
              (async-channel-put output-channel x))
            (define (do-large-computation)
              (sqrt 9))
            (let loop ([data (get)])
              (case data
                [(quit) (void)]
                [(add) (begin
                         (put (+ 1 (get)))
                         (loop (get)))]
                [(long) (begin
                          (put (do-large-computation))
                          (loop (get)))])))))

(define to-server (make-async-channel))

(define from-server (make-async-channel))

> (server to-server from-server)
#<thread>
> (async-channel? to-server)
#t
> (printf "Adding 1 to 4\n")
Adding 1 to 4

> (async-channel-put to-server 'add)

> (async-channel-put to-server 4)

> (printf "Result is ~a\n" (async-channel-get from-server))
Result is 5

> (printf "Ask server to do a long computation\n")
Ask server to do a long computation
```

```
> (async-channel-put to-server 'long)

> (printf "I can do other stuff\n")
I can do other stuff

> (printf "Ok, computation from server is ~a\n"
          (async-channel-get from-server))
Ok, computation from server is 3

> (async-channel-put to-server 'quit)
```

11.3 Thread-Local Storage

Thread cells provides primitive support for thread-local storage. Parameters combine thread cells and continuation marks to support thread-specific, continuation-specific binding.

11.3.1 Thread Cells

A *thread cell* contains a thread-specific value; that is, it contains a specific value for each thread, but it may contain different values for different threads. A thread cell is created with a default value that is used for all existing threads. When the cell's content is changed with `thread-cell-set!`, the cell's value changes only for the current thread. Similarly, `thread-cell-ref` obtains the value of the cell that is specific to the current thread.

A thread cell's value can be *preserved*, which means that when a new thread is created, the cell's initial value for the new thread is the same as the creating thread's current value. If a thread cell is non-preserved, then the cell's initial value for a newly created thread is the default value (which was supplied when the cell was created).

Within the current thread, the current values of all preserved threads cells can be captured through `current-preserved-thread-cell-values`. The captured set of values can be imperatively installed into the current thread through another call to `current-preserved-thread-cell-values`. The capturing and restoring threads can be different.

```
(thread-cell? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a thread cell, `#f` otherwise.

```
(make-thread-cell v [preserved?]) → thread-cell?
  v : any/c
  preserved? : any/c = #f
```

Creates and returns a new thread cell. Initially, *v* is the cell's value for all threads. If *preserved?* is true, then the cell's initial value for a newly created threads is the creating thread's value for the cell, otherwise the cell's value is initially *v* in all future threads.

```
(thread-cell-ref cell) → any
  cell : thread-cell?
```

Returns the current value of *cell* for the current thread.

```
(thread-cell-set! cell v) → any
  cell : thread-cell?
  v : any/c
```

Sets the value in *cell* to *v* for the current thread.

Examples:

```
> (define cnp (make-thread-cell '(nerve) #f))

> (define cp (make-thread-cell '(cancer) #t))

> (thread-cell-ref cnp)
'(nerve)
> (thread-cell-ref cp)
'(cancer)
> (thread-cell-set! cnp '(nerve nerve))

> (thread-cell-set! cp '(cancer cancer))

> (thread-cell-ref cnp)
'(nerve nerve)
> (thread-cell-ref cp)
'(cancer cancer)
> (define ch (make-channel))

> (thread (lambda ()
            (channel-put ch (thread-cell-ref cnp))
            (channel-put ch (thread-cell-ref cp))
            (channel-get ch)
            (channel-put ch (thread-cell-ref cp))))
#<thread>
> (channel-get ch)
'(nerve)
> (channel-get ch)
'(cancer cancer)
> (thread-cell-set! cp '(cancer cancer cancer))
```

```
> (thread-cell-ref cp)
'(cancer cancer cancer)
> (channel-put ch 'ok)

> (channel-get ch)
'(cancer cancer)
```

```
(current-preserved-thread-cell-values) → thread-cell-values?
(current-preserved-thread-cell-values thread-cell-vals) → void?
  thread-cell-vals : thread-cell-values?
```

When called with no arguments, this procedure produces a *thread-cell-vals* that represents the current values (in the current thread) for all preserved thread cells.

When called with a *thread-cell-vals* generated by a previous call to current-preserved-thread-cell-values, the values of all preserved thread cells (in the current thread) are set to the values captured in *thread-cell-vals*; if a preserved thread cell was created after *thread-cell-vals* was generated, then the thread cell's value for the current thread reverts to its initial value.

```
(thread-cell-values? v) → boolean?
  v : any/c
```

Returns #t if *v* is a set of thread cell values produced by current-preserved-thread-cell-values, #f otherwise.

11.3.2 Parameters

§4.13 "Dynamic Binding: `parameterize`" in ***The Racket Guide*** introduces parameters.

See §1.1.14 "Parameters" for basic information on the parameter model. Parameters correspond to *preserved thread fluids* in Scsh [Gasbichler02].

To parameterize code in a thread- and continuation-friendly manner, use `parameterize`. The `parameterize` form introduces a fresh thread cell for the dynamic extent of its body expressions.

When a new thread is created, the parameterization for the new thread's initial continuation is the parameterization of the creator thread. Since each parameter's thread cell is preserved, the new thread "inherits" the parameter values of its creating thread. When a continuation is moved from one thread to another, settings introduced with `parameterize` effectively move with the continuation.

In contrast, direct assignment to a parameter (by calling the parameter procedure with a value) changes the value in a thread cell, and therefore changes the setting only for the current thread. Consequently, as far as the memory manager is concerned, the value originally

associated with a parameter through `parameterize` remains reachable as long the continuation is reachable, even if the parameter is mutated.

```
(make-parameter v [guard]) → parameter?
  v : any/c
  guard : (or/c (any/c . -> . any) #f) = #f
```

Returns a new parameter procedure. The value of the parameter is initialized to *v* in all threads. If *guard* is supplied, it is used as the parameter's guard procedure. A guard procedure takes one argument. Whenever the parameter procedure is applied to an argument, the argument is passed on to the guard procedure. The result returned by the guard procedure is used as the new parameter value. A guard procedure can raise an exception to reject a change to the parameter's value. The *guard* is not applied to the initial *v*.

```
(parameterize ([parameter-expr value-expr] ...)
  body ...+)

  parameter-expr : parameter?
```

§4.13 "Dynamic Binding: `parameterize`" in ***The Racket Guide*** introduces `parameterize`.

The result of a `parameterize` expression is the result of the last *body*. The *parameter-expr*s determine the parameters to set, and the *value-expr*s determine the corresponding values to install while evaluating the `body-expr`s. All of the *parameter-expr*s are evaluated first (and checked with `parameter?`), then all *value-expr*s are evaluated, and then the parameters are bound in the continuation to preserved thread cells that contain the values of the *value-expr*s. The last `body-expr` is in tail position with respect to the entire `parameterize` form.

Outside the dynamic extent of a `parameterize` expression, parameters remain bound to other thread cells. Effectively, therefore, old parameters settings are restored as control exits the `parameterize` expression.

If a continuation is captured during the evaluation of `parameterize`, invoking the continuation effectively re-introduces the parameterization, since a parameterization is associated to a continuation via a continuation mark (see §10.5 "Continuation Marks") using a private key.

Examples:

```
> (parameterize ([exit-handler (lambda (x) 'no-exit)])
    (exit))

> (define p1 (make-parameter 1))

> (define p2 (make-parameter 2))

> (parameterize ([p1 3]
```

```
                [p2 (p1)])
    (cons (p1) (p2)))
'(3 . 1)
> (let ([k (let/cc out
             (parameterize ([p1 2])
               (p1 3)
               (cons (let/cc k
                       (out k))
                     (p1))))])
    (if (procedure? k)
        (k (p1))
        k))
'(1 . 3)
> (define ch (make-channel))

> (parameterize ([p1 0])
    (thread (lambda ()
              (channel-put ch (cons (p1) (p2))))))
#<thread>
> (channel-get ch)
'(0 . 2)
> (define k-ch (make-channel))

> (define (send-k)
    (parameterize ([p1 0])
      (thread (lambda ()
                (let/ec esc
                  (channel-put ch
                               ((let/cc k
                                  (channel-put k-ch k)
                                  (esc)))))))))

> (send-k)
#<thread>
> (thread (lambda () ((channel-get k-ch)
                      (let ([v (p1)])
                        (lambda () v)))))
#<thread>
> (channel-get ch)
1
> (send-k)
#<thread>
> (thread (lambda () ((channel-get k-ch) p1)))
#<thread>
> (channel-get ch)
0
```

```
(parameterize* ((parameter-expr value-expr) ...)
  body ...+)
```

Analogous to `let*` compared to `let`, `parameterize*` is the same as a nested series of single-parameter `parameterize` forms.

```
(make-derived-parameter parameter
                        guard
                        wrap)      → parameter?
  parameter : parameter?
  guard : (any/c . -> . any)
  wrap : (any/c . -> . any)
```

Returns a parameter procedure that sets or retrieves the same value as *parameter*, but with:

- *guard* applied when setting the parameter (before any guard associated with *parameter*), and
- *wrap* applied when obtaining the parameter's value.

See also `chaperone-procedure`, which can also be used to guard parameter procedures.

```
(parameter? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a parameter procedure, `#f` otherwise.

```
(parameter-procedure=? a b) → boolean?
  a : parameter?
  b : parameter?
```

Returns `#t` if the parameter procedures `a` and `b` always modify the same parameter with the same guards (although possibly with different chaperones), `#f` otherwise.

```
(current-parameterization) → parameterization?
```

Returns the current continuation's parameterization.

```
(call-with-parameterization parameterization
                            thunk)          → any
  parameterization : parameterization?
  thunk : (-> any)
```

Calls *thunk* (via a tail call) with *parameterization* as the current parameterization.

```
(parameterization? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a parameterization returned by `current-parameterization`, `#f` otherwise.

11.4 Futures

§20.1 "Parallelism with Futures" in ***The Racket Guide*** introduces futures.

```
(require racket/future)
```
package: `base`

The bindings documented in this section are provided by the `racket/future` and `racket` libraries, but not `racket/base`.

The `future` and `touch` functions from `racket/future` provide access to parallelism as supported by the hardware and operating system. In contrast to `thread`, which provides concurrency for arbitrary computations without parallelism, `future` provides parallelism for limited computations. A *future* executes its work in parallel (assuming that support for parallelism is available) until it detects an attempt to perform an operation that is too complex for the system to run safely in parallel. Similarly, work in a future is suspended if it depends in some way on the current continuation, such as raising an exception. A suspended computation for a future is resumed when `touch` is applied to the future.

Currently, parallel support for `future` is enabled by default for Windows, Linux x86/x86_64, and Mac OS X x86/x86_64. To enable support for other platforms, use `--enable-futures` with `configure` when building Racket.

"Safe" parallel execution of a future means that all operations provided by the system must be able to enforce contracts and produce results as documented. "Safe" does not preclude concurrent access to mutable data that is visible in the program. For example, a computation in a future might use `set!` to modify a shared variable, in which case concurrent assignment to the variable can be visible in other futures and threads. Furthermore, guarantees about the visibility of effects and ordering are determined by the operating system and hardware—which rarely support, for example, the guarantee of sequential consistency that is provided for `thread`-based concurrency. At the same time, operations that seem obviously safe may have a complex enough implementation internally that they cannot run in parallel. See also §20.1 "Parallelism with Futures" in ***The Racket Guide***.

A future never runs in parallel if all of the custodians that allow its creating thread to run are shut down. Such futures can execute through a call to `touch`, however.

11.4.1 Creating and Touching Futures

```
(future thunk) → future?
  thunk : (-> any)
(touch f) → any
  f : future?
```

The future procedure returns a future value that encapsulates thunk. The touch function forces the evaluation of the thunk inside the given future, returning the values produced by thunk. After touch forces the evaluation of a thunk, the resulting values are retained by the future in place of thunk, and additional touches of the future return those values.

Between a call to future and touch for a given future, the given thunk may run speculatively in parallel to other computations, as described above.

```
> (let ([f (future (lambda () (+ 1 2)))])
    (list (+ 3 4) (touch f)))
'(7 3)
```

(futures-enabled?) → boolean?

Returns whether parallel support for futures is enabled in the current Racket configuration.

(current-future) → (or/c #f future?)

Returns the descriptor of the future whose thunk execution is the current continuation. If a future thunk itself uses touch, future-thunk executions can be nested, in which case the descriptor of the most immediately executing future is returned. If the current continuation is not a future-thunk execution, the result is #f.

(future? v) → boolean?
  v : any/c

Returns #t if v is a future value, #f otherwise.

(would-be-future thunk) → future?
  thunk : (-> any)

Returns a future that never runs in parallel, but that consistently logs all potentially "unsafe" operations during the execution of the future's thunk (i.e., operations that interfere with parallel execution).

With a normal future, certain circumstances might prevent the logging of unsafe operations. For example, when executed with debug-level logging,

```
(touch (future (lambda ()
                 (printf "hello1")
                 (printf "hello2")
                 (printf "hello3"))))
```

might log three messages, one for each printf invocation. However, if the touch is performed before the future has a chance to start running in parallel, the future thunk evaluates in the same manner as any ordinary thunk, and no unsafe operations are logged. Replacing future with would-be-future ensures the logging of all three calls to printf.

```
(processor-count) → exact-positive-integer?
```

Returns the number of parallel computation units (e.g., processors or cores) that are available on the current machine.

11.4.2 Future Semaphores

```
(make-fsemaphore init) → fsemaphore?
  init : exact-nonnegative-integer?
```

Creates and returns a new *future semaphore* with the counter initially set to `init`.

A future semaphore is similar to a plain semaphore, but future-semaphore operations can be performed safely in parallel (to synchronize parallel computations). In contrast, operations on plain semaphores are not safe to perform in parallel, and they therefore prevent a computation from continuing in parallel.

```
(fsemaphore? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is an future semaphore value, `#f` otherwise.

```
(fsemaphore-post fsema) → void?
  fsema : fsemaphore?
```

Increments the future semaphore's internal counter and returns `#<void>`.

```
(fsemaphore-wait fsema) → void?
  fsema : fsemaphore?
```

Blocks until the internal counter for `fsema` is non-zero. When the counter is non-zero, it is decremented and `fsemaphore-wait` returns `#<void>`.

```
(fsemaphore-try-wait? fsema) → boolean?
  fsema : fsemaphore?
```

Like `fsemaphore-wait`, but `fsemaphore-try-wait?` never blocks execution. If `fsema`'s internal counter is zero, `fsemaphore-try-wait?` returns `#f` immediately without decrementing the counter. If `fsema`'s counter is positive, it is decremented and `#t` is returned.

```
(fsemaphore-count fsema) → exact-nonnegative-integer?
  fsema : fsemaphore?
```

Returns `fsema`'s current internal counter value.

11.4.3 Future Performance Logging

Racket traces use logging (see §15.5 “Logging”) extensively to report information about how futures are evaluated. Logging output is useful for debugging the performance of programs that use futures.

Though textual log output can be viewed directly (or retrieved in code via `trace-futures`), it is much easier to use the graphical profiler tool provided by `future-visualizer`.

Future events are reported to a logger named `'future`. In addition to its string message, each event logged for a future has a data value that is an instance of a `future-event` prefab structure:

```
(struct future-event (future-id proc-id action time prim-
name user-data)
  #:prefab)
```

The `future-id` field is an exact integer that identifies a future, or it is `#f` when `action` is `'missing`. The `future-id` field is particularly useful for correlating logged events.

The `proc-id` fields is an exact, non-negative integer that identifies a parallel process. Process 0 is the main Racket process, where all expressions other than future thunks evaluate.

The `time` field is an inexact number that represents time in the same way as `current-inexact-milliseconds`.

The `action` field is a symbol:

- `'create`: a future was created.
- `'complete`: a future’s thunk evaluated successfully, so that `touch` will produce a value for the future immediately.
- `'start-work` and `'end-work`: a particular process started and ended working on a particular future.
- `'start-0-work`: like `'start-work`, but for a future thunk that for some structural reason could not be started in a process other than 0 (e.g., the thunk requires too much local storage to start).
- `'start-overflow-work`: like `'start-work`, where the future thunk’s work was previously stopped due to an internal stack overflow.
- `'sync`: blocking (processes other than 0) or initiation of handing (process 0) for an “unsafe” operation in a future thunk’s evaluation; the operation must run in process 0.
- `'block`: like `'sync`, but for a part of evaluation that must be delayed until the future is `touch`ed, because the evaluation may depend on the current continuation.

- `'touch` (never in process 0): like `'sync` or `'block`, but for a `touch` operation within a future thunk.
- `'overflow` (never in process 0): like `'sync` or `'block`, but for the case that a process encountered an internal stack overflow while evaluating a future thunk.
- `'result` or `'abort`: waiting or handling for `'sync`, `'block`, or `'touch` ended with a value or an error, respectively.
- `'suspend` (never in process 0): a process blocked by `'sync`, `'block`, or `'touch` abandoned evaluation of a future; some other process may pick up the future later.
- `'touch-pause` and `'touch-resume` (in process 0, only): waiting in `touch` for a future whose thunk is being evaluated in another process.
- `'missing`: one or more events for the process were lost due to internal buffer limits before they could be reported, and the `time-id` field reports an upper limit on the time of the missing events; this kind of event is rare.

Assuming no `'missing` events, then `'start-work`, `'start-0-work`, `'start-overflow-work` is always paired with `'end-work`; `'sync`, `'block`, and `'touch` are always paired with `'result`, `'abort`, or `'suspend`; and `'touch-pause` is always paired with `'touch-resume`.

In process 0, some event pairs can be nested within other event pairs: `'sync`, `'block`, or `'touch` with `'result` or `'abort`; and `'touch-pause` with `'touch-resume`.

An `'block` in process 0 is generated when an unsafe operation is handled. This type of event will contain a symbol in the `unsafe-op-name` field that is the name of the operation. In all other cases, this field contains `#f`.

The `prim-name` field will always be `#f` unless the event occurred on process 0 and its `action` is either `'block` or `'sync`. If these conditions are met, `prim-name` will contain the name of the Racket primitive which required the future to synchronize with the runtime thread (represented as a symbol).

The `user-data` field may take on a number of different values depending on both the `action` and `prim-name` fields:

- `'touch` on process 0: contains the integer ID of the future being touched.
- `'sync` and `prim-name` = `|allocate memory|`: The size (in bytes) of the requested allocation.
- `'sync` and `prim-name` = `jit_on_demand`: The runtime thread is performing a JIT compilation on behalf of the future `future-id`. The field contains the name of the function being JIT compiled (as a symbol).
- `'create`: A new future was created. The field contains the integer ID of the newly created future.

11.5 Places

§20.2 “Parallelism with Places” in ***The Racket Guide*** introduces places.

```
(require racket/place)
```
package: base

The bindings documented in this section are provided by the `racket/place` and `racket` libraries, but not `racket/base`.

Places enable the development of parallel programs that take advantage of machines with multiple processors, cores, or hardware threads.

Currently, parallel support for places is enabled only for Racket 3m (which is the main variant of Racket), and only by default for Windows, Linux x86/x86_64, and Mac OS X x86/x86_64. To enable support for other platforms, use `--enable-places` with `configure` when building Racket. The `place-enabled?` function reports whether places run in parallel.

A *place* is a parallel task that is effectively a separate instance of the Racket virtual machine. Places communicate through *place channels*, which are endpoints for a two-way buffered communication.

To a first approximation, place channels support only immutable, transparent values as messages. In addition, place channels themselves can be sent across channels to establish new (possibly more direct) lines of communication in addition to any existing lines. Finally, mutable values produced by `shared-flvector`, `make-shared-flvector`, `shared-fxvector`, `make-shared-fxvector`, `shared-bytes`, and `make-shared-bytes` can be sent across place channels; mutation of such values is visible to all places that share the value, because they are allowed in a *shared memory space*. See `place-message-allowed?`.

A place channel can be used as a synchronizable event (see §11.2.1 “Events”) to receive a value through the channel. A place channel is ready for synchronization when a message is available on the channel, and the place channel’s synchronization result is the message (which is removed on synchronization). A place can also receive messages with `place-channel-get`, and messages can be sent with `place-channel-put`.

Two place channels are `equal?` if they are endpoints for the same underlying channels while both or neither is a place descriptor. Place channels can be `equal?` without being `eq?` after being sent messages through a place channel.

Constraints on messages across a place channel—and therefore on the kinds of data that places share—enable greater parallelism than `future`, even including separate garbage collection of separate places. At the same time, the setup and communication costs for places can be higher than for futures.

For example, the following expression launches two places, echoes a message to each, and then waits for the places to terminate:

```
(let ([pls (for/list ([i (in-range 2)])
             (dynamic-place "place-worker.rkt" 'place-main))])
  (for ([i (in-range 2)]
        [p pls])
    (place-channel-put p i)
    (printf "~a\n" (place-channel-get p)))
  (map place-wait pls))
```

The "place-worker.rkt" module must export the place-main function that each place executes, where place-main must accept a single place channel argument:

```
#lang racket
(provide place-main)

(define (place-main pch)
  (place-channel-put pch (format "Hello from place ~a"
                                 (place-channel-get pch))))
```

Place channels are subject to garbage collection, like other Racket values, and a thread that is blocked reading from a place channel can be garbage collected if place channel's writing end becomes unreachable. However, unlike normal channel blocking, if otherwise unreachable threads are mutually blocked on place channels that are reachable only from the same threads, the threads and place channels are all considered reachable, instead of unreachable.

When a place is created, its parameter values are generally set to the *initial* values of the parameters in the creating place, except that the *current* values of the following parameters are used: current-library-collection-paths, current-library-collection-links, and current-compiled-file-roots.

11.5.1 Using Places

```
(place-enabled?) → boolean?
```

Returns #t if Racket is configured so that dynamic-place and place create places that can run in parallel, #f if dynamic-place and place are simulated using thread.

```
(place? v) → boolean?
  v : any/c
```

Returns #t if v is a *place descriptor* value, #f otherwise. Every place descriptor is also a place channel.

```
(place-channel? v) → boolean?
  v : any/c
```

Returns #t if v is place channel, #f otherwise.

```
(dynamic-place  module-path
                start-name
               [#:at location
                #:named named]) → place?
```

```
module-path : (or/c module-path? path?)
start-name : symbol?
location : (or/c #f place-location?) = #f
named : any/c = #f
```

Creates a place to run the procedure that is identified by *module-path* and *start-name*. The result is a place descriptor value that represents the new parallel task; the place descriptor is returned immediately. The place descriptor value is also a place channel that permits communication with the place.

The module indicated by *module-path* must export a function with the name `start-proc`. The function must accept a single argument, which is a place channel that corresponds to the other end of communication for the place descriptor returned by `place`.

If *location* is provided, it must be a place location, such as a distributed places node produced by `create-place-node`.

When the place is created, the initial exit handler terminates the place, using the argument to the exit handler as the place's *completion value*. Use `(exit v)` to immediately terminate a place with the completion value *v*. Since a completion value is limited to an exact integer between `0` and `255`, any other value for `v` is converted to `0`.

If the function indicated by *module-path* and `start-proc` returns, then the place terminates with the completion value `0`.

In the created place, the `current-input-port` parameter is set to an empty input port, while the values of the `current-output-port` and `current-error-port` parameters are connected to the current ports in the creating place. If the output ports in the creating place are file-stream ports, then the connected ports in the created place share the underlying streams, otherwise a thread in the creating place pumps bytes from the created place's ports to the current ports in the creating place.

The *module-path* argument must not be a module path of the form `(quote sym)` unless the module is predefined (see `module-predefined?`).

The `dynamic-place` binding is protected in the sense of `protect-out`, so access to this operation can be prevented by adjusting the code inspector (see §14.10 "Code Inspectors").

```
(dynamic-place* module-path
                start-name
               [#:in in
                #:out out
                #:err err]) → place?
                              (or/c output-port? #f)
                              (or/c input-port? #f)
                              (or/c input-port? #f)
module-path : (or/c module-path? path?)
```

```
  start-name : symbol?
  in : (or/c input-port? #f) = #f
  out : (or/c output-port? #f) = (current-output-port)
  err : (or/c output-port? #f) = (current-error-port)
```

Like dynamic-place, but accepts specific ports to the new place's ports, and returns a created port when #f is supplied for a port. The *in*, *out*, and *err* ports are connected to the current-input-port, current-output-port, and current-error-port ports, respectively, for the place. Any of the ports can be #f, in which case a file-stream port (for an operating-system pipe) is created and returned by dynamic-place*. The *err* argument can be 'stdout, in which case the same file-stream port or that is supplied as standard output is also used for standard error. For each port or 'stdout that is provided, no pipe is created and the corresponding returned value is #f.

The caller of dynamic-place* is responsible for closing all returned ports; none are closed automatically.

The dynamic-place* procedure returns four values:

- a place descriptor value representing the created place;
- an output port piped to the place's standard input, or #f if *in* was a port;
- an input port piped from the place's standard output, or #f if *out* was a port;
- an input port piped from the place's standard error, or #f if *err* was a port or 'stdout.

The dynamic-place* binding is protected in the same way as dynamic-place.

```
(place id body ...+)
```

Creates a place that evaluates *body* expressions with *id* bound to a place channel. The *body*s close only over *id* plus the top-level bindings of the enclosing module, because the *body*s are lifted to a function that is exported by the module. The result of place is a place descriptor, like the result of dynamic-place.

The place binding is protected in the same way as dynamic-place.

```
(place* maybe-port ...
        id
        body ...+)

maybe-port =
           | #:in in-expr
           | #:out out-expr
           | #:err err-expr
```

Like `place`, but supports optional `#:in`, `#:out`, and `#:err` expressions (at most one of each) to specify ports in the same way and with the same defaults as `dynamic-place*`. The result of a `place*` form is also the same as for `dynamic-place*`.

The `place*` binding is protected in the same way as `dynamic-place`.

```
(place-wait p) → exact-integer?
  p : place?
```

Returns the completion value of the place indicated by *p*, blocking until the place has terminated.

If any pumping threads were created to connect a non-file-stream port to the ports in the place for *p* (see `dynamic-place`), `place-wait` returns only when the pumping threads have completed.

```
(place-dead-evt p) → evt?
  p : place?
```

Returns a synchronizable event (see §11.2.1 “Events”) that is ready for synchronization if and only if *p* has terminated. The synchronization result of a place-dead event is the place-dead event itself.

If any pumping threads were created to connect a non-file-stream port to the ports in the place for *p* (see `dynamic-place`), the event returned by `place-dead-evt` may become ready even if a pumping thread is still running.

```
(place-kill p) → void?
  p : place?
```

Immediately terminates the place, setting the place’s completion value to `1` if the place does not have a completion value already.

```
(place-break p [kind]) → void?
  p : place?
  kind : (or/c #f 'hang-up 'terminate) = #f
```

Sends the main thread of place *p* a break; see §10.6 “Breaks”.

```
(place-channel) → place-channel? place-channel?
```

Returns two place channels. Data sent through the first channel can be received through the second channel, and data sent through the second channel can be received from the first.

Typically, one place channel is used by the current place to send messages to a destination place; the other place channel is sent to the destination place (via an existing place channel).

```
(place-channel-put pch v) → void
  pch : place-channel?
  v : place-message-allowed?
```

Sends a message `v` on channel `pch`. Since place channels are asynchronous, `place-channel-put` calls are non-blocking.

See `place-message-allowed?` form information on automatic coercions in `v`, such as converting a mutable string to an immutable string.

```
(place-channel-get pch) → place-message-allowed?
  pch : place-channel?
```

Returns a message received on channel `pch`, blocking until a message is available.

```
(place-channel-put/get pch v) → any/c
  pch : place-channel?
  v : any/c
```

Sends an immutable message `v` on channel `pch` and then waits for a message (perhaps a reply) on the same channel.

```
(place-message-allowed? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is allowed as a message on a place channel, `#f` otherwise.

If `(place-enabled?)` returns `#f`, then the result is always `#t` and no conversions are performed on `v` as a message. Otherwise, the following kinds of data are allowed as messages:

- numbers, characters, booleans, keywords, and `#<void>`;
- symbols, where the `eq?`ness of uninterned symbols is preserved within a single message, but not across messages;
- strings and byte strings, where mutable strings and byte strings are automatically replaced by immutable variants;
- paths (for any platform);
- pairs, lists, vectors, and immutable prefab structures containing message-allowed values, where a mutable vector is automatically replaced by an immutable vector;
- hash tables where mutable hash tables are automatically replaced by immutable variants;

- place channels, where a place descriptor is automatically replaced by a plain place channel;
- file-stream ports and TCP ports, where the underlying representation (such as a file descriptor, socket, or handle) is duplicated and attached to a fresh port in the receiving place;
- C pointers as created or accessed via `ffi/unsafe`; and
- values produced by `shared-flvector`, `make-shared-flvector`, `shared-fxvector`, `make-shared-fxvector`, `shared-bytes`, and `make-shared-bytes`.

```
prop:place-location : struct-type-property?
(place-location? v) → boolean?
  v : any/c
```

A structure type property and associated predicate for implementations of *place locations*. The value of `prop:place-location` must be a procedure of four arguments: the place location itself, a module path, a symbol for the start function exported by the module, and a place name (which can be `#f` for an anonymous place).

A place location can be passed as the `#:at` argument to `dynamic-place`, which in turn simply calls the `prop:place-location` value of the place location.

A distributed places note created with `create-place-node` is an example of a place location.

11.5.2 Places Logging

Place events are reported to a logger named `'place`. In addition to its string message, each event logged for a place has a data value that is an instance of a `place-event` prefab structure:

```
(struct place-event (place-id action value time)
  #:prefab)
```

The `place-id` field is an exact integer that identifies a place.

The `time` field is an inexact number that represents time in the same way as `current-inexact-milliseconds`.

The `action` field is a symbol:

- `'create`: a place was created. This event is logged in the creating place, and the event's `value` field has the ID for the created place.

- `'reap`: a place that was previously created in the current place has exited (and that fact has been detected, possibly via `place-wait`). The event's `value` field has the ID for the exited place.
- `'enter`: a place has started, logged within the started place. The event's `value` field has `#f`.
- `'exit`: a place is exiting, logged within the exiting place. The event's `value` field has `#f`.
- `'put`: a place-channel message has been sent. The event's `value` field is a positive exact integer that approximates the message's size.
- `'get`: a place-channel message has been received. The event's `value` field is a positive exact integer that approximates the message's size.

Changed in version 6.0.0.2 of package `base`: Added logging via `'place` and `place-event`.

11.6 Engines

```
(require racket/engine)
```
package: `base`

The bindings documented in this section are provided by the `racket/engine` library, not `racket/base` or `racket`.

An *engine* is an abstraction that models processes that can be preempted by a timer or other external trigger. They are inspired by the work of Haynes and Friedman [Haynes84].

Engines log their behavior via a logger with the name `'racket/engine`. The logger is created when the module is instantiated and uses the result of `(current-logger)` as its parent. The library adds logs a `'debug` level message: when `engine-run` is called, when the engine timeout expires, and when the engine is stopped (either because it terminated or it reached a safe point to stop). Each log message holds a value of the struct:

```
(struct engine-info (msec name) #:prefab)
```

where the *msec* field holds the result of `(current-inexact-milliseconds)` at the moment of logging, and the *name* field holds the name of the procedure passed to `engine`.

```
(engine proc) → engine?
  proc : ((any/c . -> . void?) . -> . any/c)
```

Returns an engine object to encapsulate a thread that runs only when allowed. The *proc* procedure should accept one argument, and *proc* is run in the engine thread when `engine-run` is called. If `engine-run` returns due to a timeout, then the engine thread is suspended

until a future call to engine-run. Thus, *proc* only executes during the dynamic extent of a engine-run call.

The argument to *proc* is a procedure that takes a boolean, and it can be used to disable suspends (in case *proc* has critical regions where it should not be suspended). A true value passed to the procedure enables suspends, and #f disables suspends. Initially, suspends are allowed.

```
(engine? v) → any
  v : any/c
```

Returns #t if *v* is an engine produced by engine, #f otherwise.

```
(engine-run until engine) → boolean?
  until : (or/c evt? real?)
  engine : engine?
```

Allows the thread associated with *engine* to execute for up as long as *until* milliseconds (if *until* is a real number) or *until* is ready (if *until* is an event). If *engine*'s procedure disables suspends, then the engine can run arbitrarily long until it re-enables suspends.

The engine-run procedure returns #t if *engine*'s procedure completes (or if it completed earlier), and the result is available via engine-result. The engine-run procedure returns #f if *engine*'s procedure does not complete before it is suspended after timeout-secs. If *engine*'s procedure raises an exception, then it is re-raised by engine-run.

```
(engine-result engine) → any
  engine : engine?
```

Returns the result for *engine* if it has completed with a value (as opposed to an exception), #f otherwise.

```
(engine-kill engine) → void?
  engine : engine?
```

Forcibly terminates the thread associated with *engine* if it is still running, leaving the engine result unchanged.

12 Macros

§16 “Macros” in ***The Racket Guide*** introduces Macros.

See §1.2 “Syntax Model” for general information on how programs are parsed. In particular, the subsection §1.2.3.2 “Expansion Steps” describes how parsing triggers macros, and §1.2.3.5 “Transformer Bindings” describes how macro transformers are called.

12.1 Pattern-Based Syntax Matching

```
(syntax-case stx-expr (literal-id ...)
  clause ...)

     clause = [pattern result-expr]
            | [pattern fender-expr result-expr]

    pattern = _
            | id
            | (pattern ...)
            | (pattern ...+ . pattern)
            | (pattern ... pattern ellipsis pattern ...)
            | (pattern ... pattern ellipsis pattern ... . pattern)
            | #(pattern ...)
            | #(pattern ... pattern ellipsis pattern ...)
            | #&pattern
            | #s(key-datum pattern ...)
            | #s(key-datum pattern ... pattern ellipsis pattern ...)
            | (ellipsis stat-pattern)
            | const

stat-pattern = id
             | (stat-pattern ...)
             | (stat-pattern ...+ . stat-pattern)
             | #(stat-pattern ...)
             | const

   ellipsis = ...
```

Finds the first *pattern* that matches the syntax object produced by *stx-expr*, and for which the corresponding *fender-expr* (if any) produces a true value; the result is from the corresponding *result-expr*, which is in tail position for the `syntax-case` form. If no *clause* matches, then the `exn:fail:syntax` exception is raised; the exception is generated by calling `raise-syntax-error` with `#f` as the “name” argument, a string with a generic error message, and the result of *stx-expr*.

A syntax object matches a *pattern* as follows:

_

A _ pattern (i.e., an identifier with the same binding as _) matches any syntax object.

id

An *id* matches any syntax object when it is not bound to ... or _ and does not have the same binding as any *literal-id*. The *id* is further bound as *pattern variable* for the corresponding *fender-expr* (if any) and *result-expr*. A pattern-variable binding is a transformer binding; the pattern variable can be reference only through forms like `syntax`. The binding's value is the syntax object that matched the pattern with a *depth marker* of 0.

An *id* that has the same binding as a *literal-id* matches a syntax object that is an identifier with the same binding in the sense of `free-identifier=?`. The match does not introduce any pattern variables.

(*pattern* ...)

A (*pattern* ...) pattern matches a syntax object whose datum form (i.e., without lexical information) is a list with as many elements as sub-*pattern*s in the pattern, and where each syntax object that corresponds to an element of the list matches the corresponding sub-*pattern*.

Any pattern variables bound by the sub-*pattern*s are bound by the complete pattern; the bindings must all be distinct.

(*pattern* ...+ . *pattern*)

The last *pattern* must not be a (*pattern* ...), (*pattern* ...+ . *pattern*), (*pattern* ... *pattern ellipsis pattern* ...), or (*pattern* ... *pattern ellipsis pattern* ... . *pattern*) form.

Like the previous kind of pattern, but matches syntax objects that are not necessarily lists; for *n* sub-*pattern*s before the last sub-*pattern*, the syntax object's datum must be a pair such that *n*-1 `cdr`s produce pairs. The last sub-*pattern* is matched against the syntax object corresponding to the *n*th `cdr` (or the `datum->syntax` coercion of the datum using the nearest enclosing syntax object's lexical context and source location).

(*pattern* ... *pattern ellipsis pattern* ...)

Like the (*pattern* ...) kind of pattern, but matching a syntax object with any number (zero or more) elements that match the sub-*pattern* followed by *ellipsis* in the corresponding position relative to other sub-*pattern*s.

For each pattern variable bound by the sub-*pattern* followed by *ellipsis*, the larger pattern binds the same pattern variable to a list of values, one for each element of the syntax object matched to the sub-*pattern*, with an incremented depth marker. (The sub-*pattern* itself may contain *ellipsis*, leading to a pattern variables bound to lists of lists of syntax objects with a depth marker of 2, and so on.)

(*pattern* ... *pattern ellipsis pattern* ... . *pattern*)

Like the previous kind of pattern, but with a final sub-*pattern* as for (*pattern* ...+ . *pattern*). The final *pattern* never matches a syntax object whose datum is a pair.

#(*pattern* ...)

Like a (*pattern* ...) pattern, but matching a vector syntax object whose elements match the corresponding sub-*pattern*s.

#(*pattern* ... *pattern ellipsis pattern* ...)

Like a (*pattern* ... *pattern ellipsis pattern* ...) pattern, but matching a vector syntax object whose elements match the corresponding sub-*pattern*s.

#&*pattern*

Matches a box syntax object whose content matches the *pattern*.

#s(*key-datum pattern* ...)

Like a (*pattern* ...) pattern, but matching a prefab structure syntax object whose fields match the corresponding sub-*pattern*s. The *key-datum* must correspond to a valid first argument to make-prefab-struct.

#s(*key-datum pattern* ... *pattern ellipsis pattern* ...)

Like a (*pattern* ... *pattern ellipsis pattern* ...) pattern, but matching a prefab structure syntax object whose elements match the corresponding sub-*pattern*s.

(*ellipsis stat-pattern*)

Matches the same as *stat-pattern*, which is like a *pattern*, but identifiers with the binding `...` are treated the same as other *id*s.

const

A *const* is any datum that does not match one of the preceding forms; a syntax object matches a *const* pattern when its datum is `equal?` to the `quoted` *const*.

If *stx-expr* produces a syntax object that is tainted or armed, then any syntax object bound by a *pattern* are tainted—unless the binding corresponds to the whole syntax object produced by *stx-expr*, in which case it remains tainted or armed.

Examples:

```
> (require (for-syntax racket/base))

> (define-syntax (swap stx)
    (syntax-case stx ()
      [(_ a b) #'(let ([t a])
                   (set! a b)
                   (set! b t))]))

> (let ([x 5] [y 10])
    (swap x y)
    (list x y))
'(10 5)
> (syntax-case #'(ops 1 2 3 => +) (=>)
    [(_ x ... => op) #'(op x ...)])
#<syntax:564:0 (+ 1 2 3)>
> (syntax-case #'(let ([x 5] [y 9] [z 12])
                   (+ x y z))
                 (let)
    [(let ([var expr] ...) body ...)
     (list #'(var ...)
           #'(expr ...))])
'(#<syntax:565:0 (x y z)> #<syntax:565:0 (5 9 12)>)
```

(`syntax-case*` *stx-expr* (*literal-id* ...) *id-compare-expr*
  clause ...)

Like `syntax-case`, but *id-compare-expr* must produce a procedure that accepts two arguments. A *literal-id* in a *pattern* matches an identifier for which the procedure

returns true when given the identifier to match (as the first argument) and the identifier in the *pattern* (as the second argument).

In other words, `syntax-case` is like `syntax-case*` with an *id-compare-expr* that produces `free-identifier=?`.

```
(with-syntax ([pattern stx-expr] ...)
  body ...+)
```

Similar to `syntax-case`, in that it matches a *pattern* to a syntax object. Unlike `syntax-case`, all *pattern*s are matched, each to the result of a corresponding *stx-expr*, and the pattern variables from all matches (which must be distinct) are bound with a single *body* sequence. The result of the `with-syntax` form is the result of the last *body*, which is in tail position with respect to the `with-syntax` form.

If any *pattern* fails to match the corresponding *stx-expr*, the `exn:fail:syntax` exception is raised.

A `with-syntax` form is roughly equivalent to the following `syntax-case` form:

```
(syntax-case (list stx-expr ...) ()
  [(pattern ...) (let () body ...+)])
```

However, if any individual *stx-expr* produces a non-syntax object, then it is converted to one using `datum->syntax` and the lexical context and source location of the individual *stx-expr*.

Examples:

```
> (define-syntax (hello stx)
    (syntax-case stx ()
      [(_ name place)
       (with-syntax ([print-name #'(printf "~a\n" 'name)]
                     [print-place #'(printf "~a\n" 'place)])
         #'(begin
             (define (name times)
               (printf "Hello\n")
               (for ([i (in-range 0 times)])
                    print-name))
             (define (place times)
               (printf "From\n")
               (for ([i (in-range 0 times)])
                    print-place))))]))

> (hello jon utah)
```

```
> (jon 2)
Hello
jon
jon

> (utah 2)
From
utah
utah

> (define-syntax (math stx)
    (define (make+1 expression)
      (with-syntax ([e expression])
        #'(+ e 1)))
    (syntax-case stx ()
      [(_ numbers ...)
       (with-syntax ([(added ...)
                      (map make+1
                           (syntax->list #'(numbers ...)))])
         #'(begin
             (printf "got ~a\n" added)
             ...))]))

> (math 3 1 4 1 5 9)
got 4
got 2
got 5
got 2
got 6
got 10
```

(syntax *template*)

```
     template = id
              | (template-elem ...)
              | (template-elem ...+ . template)
              | #(template-elem ...)
              | #&template
              | #s(key-datum template-elem ...)
              | (ellipsis stat-template)
              | const

template-elem = template ellipsis ...

stat-template = id
              | (stat-template ...)
              | (stat-template ... . stat-template)
              | #(stat-template ...)
              | #&stat-template
              | #s(key-datum stat-template ...)
              | const

     ellipsis = ...
```

Constructs a syntax object based on a *template*, which can include pattern variables bound by `syntax-case` or `with-syntax`.

Template forms produce a syntax object as follows:

id

If *id* is bound as a pattern variable, then *id* as a template produces the pattern variable's match result. Unless the *id* is a sub-*template* that is replicated by *ellipsis* in a larger *template*, the pattern variable's value must be a syntax object with a depth marker of 0 (as opposed to a list of matches).

More generally, if the pattern variable's value has a depth marker *n*, then it can only appear within a template where it is replicated by at least *n* *ellipsis*es. In that case, the template will be replicated enough times to use each match result at least once.

If *id* is not bound as a pattern variable, then *id* as a template produces `(quote-syntax` *id*`)`.

`(`*template-elem* `...)`

Produces a syntax object whose datum is a list, and where the elements of the list correspond to syntax objects produced by the *template-elem*s.

A *template-elem* is a sub-*template* replicated by any number of *ellipsis*es:

- If the sub-*template* is replicated by no *ellipsis*es, then it generates a single syntax object to incorporate into the result syntax object.
- If the sub-*template* is replicated by one *ellipsis*, then it generates a sequence of syntax objects that is "inlined" into the resulting syntax object.

  The number of generated elements depends on the values of pattern variables referenced within the sub-*template*. There must be at least one pattern variable whose value has a depth marker less than the number of *ellipsis*es after the pattern variable within the sub-*template*.

  If a pattern variable is replicated by more *ellipsis*es in a *template* than the depth marker of its binding, then the pattern variable's result is determined normally for inner *ellipsis*es (up to the binding's depth marker), and then the result is replicated as necessary to satisfy outer *ellipsis*es.
- For each *ellipsis* after the first one, the preceding element (with earlier replicating *ellipsis*es) is conceptually wrapped with parentheses for generating output, and then the wrapping parentheses are removed in the resulting syntax object.

```
(template-elem ... . template)
```

Like the previous form, but the result is not necessarily a list; instead, the place of the empty list in the resulting syntax object's datum is taken by the syntax object produced by *template*.

```
#(template-elem ...)
```

Like the `(`*template-elem* `...)` form, but producing a syntax object whose datum is a vector instead of a list.

```
#&template
```

Produces a syntax object whose datum is a box holding the syntax object produced by *template*.

```
#s(key-datum template-elem ...)
```

Like the `(`*template-elem* `...)` form, but producing a syntax object whose datum is a prefab structure instead of a list. The *key-datum* must correspond to a valid first argument of `make-prefab-struct`.

(*ellipsis stat-template*)

> Produces the same result as *stat-template*, which is like a *template*, but `...` is treated like an *id* (with no pattern binding).

const

> A *const* template is any form that does not match the preceding cases, and it produces the result (`quote-syntax` *const*).

A (`syntax` *template*) form is normally abbreviated as `#'`*template*; see also §1.3.8 “Reading Quotes”. If *template* contains no pattern variables, then `#'`*template* is equivalent to (`quote-syntax` *template*).

(`quasisyntax` *template*)

Like `syntax`, but (`unsyntax` *expr*) and (`unsyntax-splicing` *expr*) escape to an expression within the *template*.

The *expr* must produce a syntax object (or syntax list) to be substituted in place of the `unsyntax` or `unsyntax-splicing` form within the quasiquoting template, just like `unquote` and `unquote-splicing` within `quasiquote`. (If the escaped expression does not generate a syntax object, it is converted to one in the same way as for the right-hand side of `with-syntax`.) Nested `quasisyntax`es introduce quasiquoting layers in the same way as nested `quasiquote`s.

Also analogous to `quasiquote`, the reader converts `#`` to `quasisyntax`, `#,` to `unsyntax`, and `#,@` to `unsyntax-splicing`. See also §1.3.8 “Reading Quotes”.

(`unsyntax` *expr*)

Illegal as an expression form. The `unsyntax` form is for use only with a `quasisyntax` template.

(`unsyntax-splicing` *expr*)

Illegal as an expression form. The `unsyntax-splicing` form is for use only with a `quasisyntax` template.

(`syntax/loc` *stx-expr template*)

Like `syntax`, except that the immediate resulting syntax object takes its source-location information from the result of *stx-expr* (which must produce a syntax object), unless the *template* is just a pattern variable, or both the source and position of *stx-expr* are `#f`.

```
(quasisyntax/loc stx-expr template)
```

Like `quasisyntax`, but with source-location assignment like `syntax/loc`.

```
(quote-syntax/prune id)
```

Like `quote-syntax`, but the lexical context of *id* is pruned via `identifier-prune-lexical-context` to including binding only for the symbolic name of *id* and for `'#%top`. Use this form to quote an identifier when its lexical information will not be transferred to other syntax objects (except maybe to `'#%top` for a top-level binding).

```
(syntax-rules (literal-id ...)
  [(id . pattern) template] ...)
```

Equivalent to

```
(lambda (stx)
  (syntax-case stx (literal-id ...)
    [(generated-id . pattern) (syntax-protect #'template)]  ...))
```

where each *generated-id* binds no identifier in the corresponding *template*.

```
(syntax-id-rules (literal-id ...)
  [pattern template] ...)
```

Equivalent to

```
(make-set!-transformer
 (lambda (stx)
   (syntax-case stx (literal-id ...)
     [pattern (syntax-protect #'template)]  ...))))
```

```
(define-syntax-rule (id . pattern) template)
```

Equivalent to

```
(define-syntax id
  (syntax-rules ()
   [(id . pattern) template]))
```

but with syntax errors potentially phrased in terms of *pattern*.

```
...
```

The ... transformer binding prohibits ... from being used as an expression. This binding is useful only in syntax patterns and templates, where it indicates repetitions of a pattern or template. See syntax-case and syntax.

```
_
```

The _ transformer binding prohibits _ from being used as an expression. This binding is useful only in syntax patterns, where it indicates a pattern that matches any syntax object. See syntax-case.

```
(syntax-pattern-variable? v) → boolean?
  v : any/c
```

Returns #t if v is a value that, as a transformer-binding value, makes the bound variable as pattern variable in syntax and other forms. To check whether an identifier is a pattern variable, use syntax-local-value to get the identifier's transformer value, and then test the value with syntax-pattern-variable?.

The syntax-pattern-variable? procedure is provided for-syntax by racket/base.

12.2 Syntax Object Content

```
(syntax? v) → boolean?
  v : any/c
```

Returns #t if v is a syntax object, #f otherwise. See also §1.2.2 "Syntax Objects".

Examples:

```
> (syntax? #'quinoa)
#t
> (syntax? #'(spelt triticale buckwheat))
#t
> (syntax? (datum->syntax #f 'millet))
#t
> (syntax? "barley")
#f
```

```
(identifier? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a syntax object and `(syntax-e stx)` produces a symbol.

Examples:

```
> (identifier? #'linguine)
#t
> (identifier? #'(if wheat? udon soba))
#f
> (identifier? 'ramen)
#f
> (identifier? 15)
#f
```

```
(syntax-source stx) → any
  stx : syntax?
```

Returns the source for the syntax object *stx*, or `#f` if none is known. The source is represented by an arbitrary value (e.g., one passed to `read-syntax`), but it is typically a file path string. Source-location information is dropped for a syntax object that is marshaled as part of compiled code; see also `current-compile`.

```
(syntax-line stx) → (or/c exact-positive-integer? #f)
  stx : syntax?
```

Returns the line number (positive exact integer) for the start of the syntax object in its source, or `#f` if the line number or source is unknown. The result is `#f` if and only if `(syntax-column stx)` produces `#f`. See also §13.1.4 “Counting Positions, Lines, and Columns”, and see `syntax-source` for information about marshaling compiled syntax objects.

```
(syntax-column stx) → (or/c exact-nonnegative-integer? #f)
  stx : syntax?
```

Returns the column number (non-negative exact integer) for the start of the syntax object in its source, or `#f` if the source column is unknown. The result is `#f` if and only if `(syntax-line stx)` produces `#f`. See also §13.1.4 “Counting Positions, Lines, and Columns”, and see `syntax-source` for information about marshaling compiled syntax objects.

```
(syntax-position stx) → (or/c exact-positive-integer? #f)
  stx : syntax?
```

Returns the character position (positive exact integer) for the start of the syntax object in its source, or `#f` if the source position is unknown. See also §13.1.4 “Counting Positions, Lines, and Columns”, and see `syntax-source` for information about marshaling compiled syntax objects.

```
(syntax-span stx) → (or/c exact-nonnegative-integer? #f)
  stx : syntax?
```

Returns the span (non-negative exact integer) in characters of the syntax object in its source, or `#f` if the span is unknown. See also `syntax-source` for information about marshaling compiled syntax objects.

```
(syntax-original? stx) → boolean?
  stx : syntax?
```

Returns `#t` if *`stx`* has the property that `read-syntax` attaches to the syntax objects that they generate (see §12.7 "Syntax Object Properties"), and if *`stx`*'s lexical information does not indicate that the object was introduced by a syntax transformer (see §1.2.2 "Syntax Objects"). The result is `#f` otherwise. This predicate can be used to distinguish syntax objects in an expanded expression that were directly present in the original expression, as opposed to syntax objects inserted by macros.

```
(syntax-source-module stx [source?])
 → (or/c module-path-index? symbol? path? resolved-module-path? #f)
  stx : syntax?
  source? : any/c = #f
```

Returns an indication of the module whose source contains *`stx`*, or `#f` if *`stx`* has no source module. If *`source?`* is `#f`, then result is a module path index or symbol (see §14.4.2 "Compiled Modules and References") or a resolved module path; if *`source?`* is true, the result is a path or symbol corresponding to the loaded module's source in the sense of `current-module-declare-source`.

```
(syntax-e stx) → any
  stx : syntax?
```

Unwraps the immediate datum structure from a syntax object, leaving nested syntax structure (if any) in place. The result of `(syntax-e stx)` is one of the following:

- a symbol
- a syntax pair (described below)
- the empty list
- an immutable vector containing syntax objects
- an immutable box containing syntax objects
- an immutable hash table containing syntax object values (but not necessarily syntax object keys)
- an immutable prefab structure containing syntax objects
- some other kind of datum—usually a number, boolean, or string—that is interned when `datum-intern-literal` would convert the value

Examples:

```
> (syntax-e #'a)
'a
> (syntax-e #'(x . y))
'(#<syntax:11:0 x> . #<syntax:11:0 y>)
> (syntax-e #'#(1 2 (+ 3 4)))
'#(#<syntax:12:0 1> #<syntax:12:0 2> #<syntax:12:0 (+ 3 4)>)
> (syntax-e #'#&"hello world")
'#&#<syntax:13:0 "hello world">
> (syntax-e #'#hash((imperial . "yellow") (festival . "green")))
'#hash((festival . #<syntax:14:0 "green">)
       (imperial . #<syntax:14:0 "yellow">))
> (syntax-e #'#(point 3 4))
'#(#<syntax:15:0 point> #<syntax:15:0 3> #<syntax:15:0 4>)
> (syntax-e #'3)
3
> (syntax-e #'"three")
"three"
> (syntax-e #'#t)
#t
```

A *syntax pair* is a pair containing a syntax object as its first element, and either the empty list, a syntax pair, or a syntax object as its second element.

A syntax object that is the result of `read-syntax` reflects the use of delimited `.` in the input by creating a syntax object for every pair of parentheses in the source, and by creating a pair-valued syntax object *only* for parentheses in the source. See §1.3.6 “Reading Pairs and Lists” for more information.

If *stx* is tainted or armed, then any syntax object in the result of `(syntax-e` *stx*`)` is tainted, and multiple calls to `syntax-e` may return values that are not `eq?`. For a *stx* that is not armed, the results from multiple calls to `syntax-e` of *stx* are `eq?`.

`(syntax->list` *stx*`)` → `(or/c list? #f)`
  stx : `syntax?`

Returns a list of syntax objects or `#f`. The result is a list of syntax objects when `(syntax->datum` *stx*`)` would produce a list. In other words, syntax pairs in `(syntax-e` *stx*`)` are flattened.

If *stx* is tainted or armed, then any syntax object in the result of `(syntax->list` *stx*`)` is tainted.

Examples:

```
> (syntax->list #'())
```

```
'()
> (syntax->list #'(1 (+ 3 4) 5 6))
'(#<syntax:20:0 1> #<syntax:20:0 (+ 3 4)> #<syntax:20:0 5>
#<syntax:20:0 6>)
> (syntax->list #'a)
#f
```

(syntax->datum *stx*) → any
  stx : syntax?

Returns a datum by stripping the lexical information, source-location information, properties, and tamper status from *stx*. Inside of pairs, (immutable) vectors, (immutable) boxes, immutable hash table values (not keys), and immutable prefab structures, syntax objects are recursively stripped.

The stripping operation does not mutate *stx*; it creates new pairs, vectors, boxes, hash tables, and prefab structures as needed to strip lexical and source-location information recursively.

Examples:

```
> (syntax->datum #'a)
'a
> (syntax->datum #'(x . y))
'(x . y)
> (syntax->datum #'#(1 2 (+ 3 4)))
'#(1 2 (+ 3 4))
> (syntax->datum #'#&"hello world")
'#&"hello world"
> (syntax->datum #'#hash((imperial . "yellow") (festival .
"green")))
'#hash((festival . "green") (imperial . "yellow"))
> (syntax->datum #'#(point 3 4))
'#(point 3 4)
> (syntax->datum #'3)
3
> (syntax->datum #'"three")
"three"
> (syntax->datum #'#t)
#t
```

(datum->syntax *ctxt v* [*srcloc prop ignored*]) → syntax?
  ctxt : (or/c syntax? #f)
  v : any/c

```
          (or/c syntax? #f
                (list/c any/c
                        (or/c exact-positive-integer? #f)
                        (or/c exact-nonnegative-integer? #f)
                        (or/c exact-positive-integer? #f)
srcloc :                (or/c exact-nonnegative-integer? #f))
                (vector/c any/c
                         (or/c exact-positive-integer? #f)
                         (or/c exact-nonnegative-integer? #f)
                         (or/c exact-positive-integer? #f)
                         (or/c exact-nonnegative-integer? #f)))
       = #f
prop : (or/c syntax? #f) = #f
ignored : (or/c syntax? #f) = #f
```

Converts the datum *v* to a syntax object. The contents of pairs, vectors, and boxes, the fields of prefab structures, and the values of immutable hash tables are recursively converted. The keys of prefab structures and the keys of immutable hash tables are not converted. Mutable vectors and boxes are replaced by immutable vectors and boxes. For any kind of value other than a pair, vector, box, immutable hash table, immutable prefab structure, or syntax object, conversion means wrapping the value with lexical information, source-location information, and properties after the value is interned via `datum-intern-literal`.

Converted objects in *v* are given the lexical context information of *ctxt* and the source-location information of *srcloc*. If *v* is not already a syntax object, then the resulting immediate syntax object is given the properties (see §12.7 “Syntax Object Properties”) of *prop* (even the hidden ones that would not be visible via `syntax-property-symbol-keys`); if *v* is a pair, vector, box, immutable hash table, or immutable prefab structure, recursively converted values are not given properties. If *ctxt* is tainted or armed, then the resulting syntax object from `datum->syntax` is tainted.

Any of *ctxt*, *srcloc*, or *prop* can be `#f`, in which case the resulting syntax has no lexical context, source information, and/or new properties.

If *srcloc* is not `#f` or a syntax object, it must be a list or vector of five elements:

```
(list source-name line column position span)
or (vector source-name line column position span)
```

where `source-name` is an arbitrary value for the source name; `line` is an integer for the source line, or `#f`; `column` is an integer for the source column, or `#f`; `position` is an integer for the source position, or `#f`; and `span` is an integer for the source span, or `#f`. The `line` and `column` values must both be numbers or both be `#f`, otherwise the `exn:fail:contract` exception is raised.

Graph structure is not preserved by the conversion of *v* to a syntax object. Instead, *v* is essentially unfolded into a tree. If *v* has a cycle through pairs, vectors, boxes, immutable hash tables, and immutable prefab structures, then the `exn:fail:contract` exception is raised.

The *ignored* argument is allowed for backward compatibility and has no effect on the returned syntax object.

```
(datum-intern-literal v) → any/c
  v : any/c
```

Converts some values to be consistent with an interned result produced by the default reader in `read-syntax` mode.

If *v* is a number, character, string, byte string, or regular expression, then the result is a value that is `equal?` to *v* and `eq?` to a potential result of the default reader. (Note that mutable strings and byte strings are interned as immutable strings and byte strings.)

If *v* is an uninterned or an unreadable symbol, the result is still *v*, since an interned symbol would not be `equal?` to *v*.

The conversion process does not traverse compound values. For example, if *v* is a pair containing strings, then the strings within *v* are not interned.

If *v1* and *v2* are `equal?` but not `eq?`, then it is possible that `(datum-intern-literal v1)` will return *v1* and—sometime after *v1* becomes unreachable as determined by the garbage collector (see §1.1.7 “Garbage Collection”)—`(datum-intern-literal v2)` can still return *v2*. In other words, `datum-intern-literal` may adopt a given value as an interned representative, but if a former representative becomes otherwise unreachable, then `datum-intern-literal` may adopt a new representative.

```
(syntax-shift-phase-level stx shift) → syntax?
  stx : syntax?
  shift : (or/c exact-integer? #f)
```

Returns a syntax object that is like *stx*, but with all of its top-level and module bindings shifted by *shift* phase levels. If *shift* is `#f`, then only bindings at phase level 0 are shifted to the label phase level. If *shift* is `0`, then the result is *stx*.

```
(generate-temporaries stx-pair) → (listof identifier?)
  stx-pair : (or syntax? list?)
```

Returns a list of identifiers that are distinct from all other identifiers. The list contains as many identifiers as *stx-pair* contains elements. The *stx-pair* argument must be a syntax pair that can be flattened into a list. The elements of *stx-pair* can be anything, but string, symbol, keyword (possibly wrapped as syntax), and identifier elements will be embedded in the corresponding generated name, which is useful for debugging purposes.

The generated identifiers are built with interned symbols (not gensyms); see also §1.4.16 “Printing Compiled Code”.

Examples:

```
> (generate-temporaries '(a b c d))
'(#<syntax a1> #<syntax b2> #<syntax c3> #<syntax d4>)
> (generate-temporaries #'(1 2 3 4))
'(#<syntax temp5> #<syntax temp6> #<syntax temp7> #<syntax temp8>)
> (define-syntax (set!-values stx)
    (syntax-case stx ()
      [(_ (id ...) expr)
       (with-syntax ([(temp ...) (generate-
temporaries #'(id ...))])
         #'(let-values ([(temp ...) expr])
             (set! id temp) ... (void)))]))
```

```
(identifier-prune-lexical-context  id-stx
                                   [syms]) → identifier?
  id-stx : identifier?
  syms : (listof symbol?) = (list (syntax-e id-stx))
```

Returns an identifier with the same binding as *id-stx*, but without lexical information from *id-stx* that does not apply to the symbols in *syms*, where even further extension of the lexical information drops information for other symbols. In particular, transferring the lexical context via datum->syntax from the result of this function to a symbol other than one in *syms* produces an identifier with no binding.

See also quote-syntax/prune.

```
(identifier-prune-to-source-module id-stx) → identifier?
  id-stx : identifier?
```

Returns an identifier with its lexical context minimized to that needed for syntax-source-module. The minimized lexical context does not include any bindings.

```
(syntax-recertify new-stx
                  old-stx
                  inspector
                  key)      → syntax?
  new-stx : syntax?
  old-stx : syntax?
  inspector : inspector?
  key : any/c
```

For backward compatibility only; returns *new-stx*.

12.3 Syntax Object Bindings

```
(bound-identifier=? a-id b-id [phase-level]) → boolean?
  a-id : syntax?
  b-id : syntax?
  phase-level : (or/c exact-integer? #f)
              = (syntax-local-phase-level)
```

Returns `#t` if the identifier `a-id` would bind `b-id` (or vice versa) if the identifiers were substituted in a suitable expression context at the phase level indicated by `phase-level`, `#f` otherwise. A `#f` value for `phase-level` corresponds to the label phase level.

Examples:

```
> (define-syntax (check stx)
    (syntax-case stx ()
      [(_ x y)
       (if (bound-identifier=? #'x #'y)
           #'(let ([y 'wrong]) (let ([x 'binds]) y))
           #'(let ([y 'no-binds]) (let ([x 'wrong]) y)))]))

> (check a a)
'binds
> (check a b)
'no-binds
> (define-syntax-rule (check-a x) (check a x))

> (check-a a)
'no-binds
```

```
(free-identifier=?  a-id
                    b-id
                   [a-phase-level
                    b-phase-level]) → boolean?
  a-id : identifier?
  b-id : identifier?
  a-phase-level : (or/c exact-integer? #f)
                = (syntax-local-phase-level)
  b-phase-level : (or/c exact-integer? #f) = a-phase-level
```

Returns `#t` if `a-id` and `b-id` access the same local binding, module binding, or top-level binding—perhaps via rename transformers—at the phase levels indicated by `a-phase-level` and `b-phase-level`, respectively. A `#f` value for `a-phase-level` or `b-phase-level` corresponds to the label phase level.

“Same module binding” means that the identifiers refer to the same original definition site, and not necessarily to the same `require` or `provide` site. Due to renaming in `require` and

`provide`, or due to a transformer binding to a rename transformer, the identifiers may return distinct results with `syntax-e`.

Examples:

```
> (define-syntax (check stx)
    (syntax-case stx ()
      [(_ x)
       (if (free-identifier=? #'car #'x)
           #'(list 'same: x)
           #'(list 'different: x))]))

> (check car)
'(same: #<procedure:car>)
> (check mcar)
'(different: #<procedure:mcar>)
> (let ([car list])
    (check car))
'(different: #<procedure:list>)
> (require (rename-in racket/base [car kar]))

> (check kar)
'(same: #<procedure:car>)
```

```
(free-transformer-identifier=? a-id b-id) → boolean?
  a-id : identifier?
  b-id : identifier?
```

Same as `(free-identifier=? a-id b-id (add1 (syntax-local-phase-level)))`.

```
(free-template-identifier=? a-id b-id) → boolean?
  a-id : identifier?
  b-id : identifier?
```

Same as `(free-identifier=? a-id b-id (sub1 (syntax-local-phase-level)))`.

```
(free-label-identifier=? a-id b-id) → boolean?
  a-id : identifier?
  b-id : identifier?
```

Same as `(free-identifier=? a-id b-id #f)`.

```
(check-duplicate-identifier ids) → (or/c identifier? #f)
  ids : (listof identifier?)
```

Compares each identifier in *ids* with every other identifier in the list with bound-identifier=?. If any comparison returns #t, one of the duplicate identifiers is returned (the first one in *ids* that is a duplicate), otherwise the result is #f.

```
(identifier-binding id-stx [phase-level])
    (or/c 'lexical
          #f
          (listof module-path-index?
                  symbol?
→                 module-path-index?
                  symbol?
                  exact-nonnegative-integer?
                  (or/c exact-integer? #f)
                  (or/c exact-integer? #f)))
  id-stx : identifier?
  phase-level : (or/c exact-integer? #f)
              = (syntax-local-phase-level)
```

Returns one of three kinds of values, depending on the binding of *id-stx* at the phase level indicated by *phase-level* (where a #f value for *phase-level* corresponds to the label phase level):

- The result is 'lexical if *id-stx* has a local binding. If 'lexical is produced for any *phase-level* value, then it is produced for all *phase-level* values.
- The result is a list of seven items when *id-stx* has a module binding: (list *source-mod source-id nominal-source-mod nominal-source-id source-phase import-phase nominal-export-phase*).
  - *source-mod* is a module path index (see §14.4.2 “Compiled Modules and References”) that indicates the defining module.
  - *source-id* is a symbol for the identifier’s name at its definition site in the source module. This can be different from the local name returned by syntax->datum for several reasons: the identifier is renamed on import, it is renamed on export, or it is implicitly renamed because the identifier (or its import) was generated by a macro invocation.
  - *nominal-source-mod* is a module path index (see §14.4.2 “Compiled Modules and References”) that indicates the module `required` into the context of *id-stx* to provide its binding. It can be different from *source-mod* due to a re-export in *nominal-source-mod* of some imported identifier. If the same binding is imported in multiple ways, an arbitrary representative is chosen.
  - *nominal-source-id* is a symbol for the identifier’s name as exported by *nominal-source-mod*. It can be different from *source-id* due to a renaming `provide`, even if *source-mod* and *nominal-source-mod* are the same.

  - *source-phase* is an exact non-negative integer representing the source phase. For example, it is 1 if the source definition is for-syntax.
  - *import-phase* is 0 if the binding import of *nominal-source-mode* is a plain require, 1 if it is from a for-syntax import, etc.
  - *nominal-export-phase* is the phase level of the export from *nominal-source-mod*.

- The result is #f if *id-stx* has a top-level binding (or, equivalently, if it is unbound).

If *id-stx* is bound to a rename-transformer, the result from identifier-binding is for the identifier in the transformer, so that identifier-binding is consistent with free-identifier=?.

```
(identifier-transformer-binding id-stx)
    (or/c 'lexical
          #f
          (listof module-path-index?
                  symbol?
 →                module-path-index?
                  symbol?
                  exact-nonnegative-integer?
                  (or/c exact-integer? #f)
                  (or/c exact-integer? #f)))
  id-stx : identifier?
```

Same as (identifier-binding *id-stx* (add1 (syntax-local-phase-level))).

```
(identifier-template-binding id-stx)
    (or/c 'lexical
          #f
          (listof module-path-index?
                  symbol?
 →                module-path-index?
                  symbol?
                  exact-nonnegative-integer?
                  (or/c exact-integer? #f)
                  (or/c exact-integer? #f)))
  id-stx : identifier?
```

Same as (identifier-binding *id-stx* (sub1 (syntax-local-phase-level))).

```
(identifier-label-binding id-stx)
```

```
   (or/c 'lexical
         #f
         (listof module-path-index?
                 symbol?
→                module-path-index?
                 symbol?
                 exact-nonnegative-integer?
                 (or/c exact-integer? #f)
                 (or/c exact-integer? #f)))
  id-stx : identifier?
```

Same as `(identifier-binding id-stx #f)`.

```
(identifier-binding-symbol  id-stx
                            [phase-level]) → symbol?
  id-stx : identifier?
  phase-level : (or/c exact-integer? #f)
              = (syntax-local-phase-level)
```

Like `identifier-binding`, but produces a symbol that corresponds to the binding. The symbol result is the same for any identifiers that are `free-identifier=?`, but the result may also be the same for identifiers that are not `free-identifier=?` (i.e., different symbols imply different bindings, but the same symbol does not imply the same binding).

When `identifier-binding` would produce a list, then the second element of that list is the result that `identifier-binding-symbol` produces.

12.4 Syntax Transformers

```
(set!-transformer? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a value created by `make-set!-transformer` or an instance of a structure type with the `prop:set!-transformer` property, `#f` otherwise.

```
(make-set!-transformer proc) → set!-transformer?
  proc : (syntax? . -> . syntax?)
```

Creates an assignment transformer that cooperates with `set!`. If the result of `make-set!-transformer` is bound to *id* as a transformer binding, then *proc* is applied as a transformer when *id* is used in an expression position, or when it is used as the target of a `set!` assignment as `(set! id expr)`. When the identifier appears as a `set!` target, the entire `set!` expression is provided to the transformer.

Example:

```
> (let ([x 1]
        [y 2])
    (let-syntax ([x (make-set!-transformer
                     (lambda (stx)
                       (syntax-case stx (set!)
                         ; Redirect mutation of x to y
                         [(set! id v) #'(set! y v)]
                         ; Normal use of x really gets x
                         [id (identifier? #'id)  #'x]))))])
      (begin
        (set! x 3)
        (list x y))))
'(1 3)
```

```
(set!-transformer-procedure transformer)
 → (syntax? . -> . syntax?)
  transformer : set!-transformer?
```

Returns the procedure that was passed to `make-set!-transformer` to create *transformer* or that is identified by the `prop:set!-transformer` property of *transformer*.

```
prop:set!-transformer : struct-type-property?
```

A structure type property to identify structure types that act as assignment transformers like the ones created by `make-set!-transformer`.

The property value must be an exact integer or procedure of one or two arguments. In the former case, the integer designates a field within the structure that should contain a procedure; the integer must be between `0` (inclusive) and the number of non-automatic fields in the structure type (exclusive, not counting supertype fields), and the designated field must also be specified as immutable.

If the property value is a procedure of one argument, then the procedure serves as a syntax transformer and for `set!` transformations. If the property value is a procedure of two arguments, then the first argument is the structure whose type has `prop:set!-transformer` property, and the second argument is a syntax object as for a syntax transformer and for `set!` transformations; `set!-transformer-procedure` applied to the structure produces a new function that accepts just the syntax object and calls the procedure associated through the property. Finally, if the property value is an integer, the target identifier is extracted from the structure instance; if the field value is not a procedure of one argument, then a procedure that always calls `raise-syntax-error` is used, instead.

If a value has both the `prop:set!-transformer` and `prop:rename-transformer` properties, then the latter takes precedence. If a structure type has the `prop:set!-transformer`

and prop:procedure properties, then the former takes precedence for the purposes of macro expansion.

```
(rename-transformer? v) → boolean?
  v : any/c
```

Returns #t if *v* is a value created by make-rename-transformer or an instance of a structure type with the prop:rename-transformer property, #f otherwise.

Examples:

```
> (rename-transformer? (make-rename-transformer #'values))
#t
> (rename-transformer? 'not-a-rename-transformer)
#f
```

```
(make-rename-transformer id-stx
                         [delta-introduce]) → rename-transformer?
  id-stx : syntax?
  delta-introduce : (identifier? . -> . identifier?)
                  = (lambda (id) id)
```

Creates a rename transformer that, when used as a transformer binding, acts as a transformer that inserts the identifier *id-stx* in place of whatever identifier binds the transformer, including in non-application positions, in set! expressions.

Such a transformer could be written manually, but the one created by make-rename-transformer triggers special cooperation with the parser and other syntactic forms when *id* is bound to the rename transformer:

- The parser installs a free-identifier=? and identifier-binding equivalence between *id* and *id-stx*, as long as *id-stx* does not have a true value for the 'not-free-identifier=? syntax property.
- A provide of *id* provides the binding indicated by *id-stx* instead of *id*, as long as *id-stx* does not have a true value for the 'not-free-identifier=? syntax property and as long as *id-stx* has a binding.
- If provide exports *id*, it uses a symbol-valued 'nominal-id property of *id-stx* to specify the "nominal source identifier" of the binding as reported by identifier-binding.
- If *id-stx* has a true value for the 'not-provide-all-defined syntax property, then *id* (or its target) is not exported by all-defined-out.
- The syntax-local-value and syntax-local-make-delta-introducer functions recognize rename-transformer bindings and consult their targets.

Examples:

```
> (define-syntax my-or (make-rename-transformer #'or))

> (my-or #f #t)
#t
> (free-identifier=? #'my-or #'or)
#t
```

```
(rename-transformer-target transformer) → identifier?
  transformer : rename-transformer?
```

Returns the identifier passed to `make-rename-transformer` to create *transformer* or as indicated by a `prop:rename-transformer` property on *transformer*.

Example:

```
> (rename-transformer-target (make-rename-transformer #'or))
#<syntax:8:0 or>
```

```
prop:rename-transformer : struct-type-property?
```

A structure type property to identify structure types that act as rename transformers like the ones created by `make-rename-transformer`.

The property value must be an exact integer or an identifier syntax object. In the former case, the integer designates a field within the structure that should contain an identifier; the integer must be between `0` (inclusive) and the number of non-automatic fields in the structure type (exclusive, not counting supertype fields), and the designated field must also be specified as immutable.

If the property value is an identifier, the identifier serves as the target for renaming, just like the first argument to `make-rename-transformer`. If the property value is an integer, the target identifier is extracted from the structure instance; if the field value is not an identifier, then an identifier `?` with an empty context is used, instead.

```
(local-expand  stx
               context-v
               stop-ids
              [intdef-ctx]) → syntax?
  stx : syntax?
  context-v : (or/c 'expression 'top-level 'module 'module-begin list?)
  stop-ids : (or/c (listof identifier?) #f)
               (or/c internal-definition-context?
                     (and/c pair?
                            (listof internal-definition-context?))
                     #f)
  intdef-ctx :
             = #f
```

Expands *stx* in the lexical context of the expression currently being expanded. The *context-v* argument is used as the result of `syntax-local-context` for immediate expansions; a list indicates an internal-definition context, and more information on the form of the list is below.

When an identifier in *stop-ids* is encountered by the expander in a sub-expression, expansions stops for the sub-expression. If *stop-ids* is a non-empty list and does not contain just `module*`, then `begin`, `quote`, `set!`, `lambda`, `case-lambda`, `let-values`, `letrec-values`, `if`, `begin0`, `with-continuation-mark`, `letrec-syntaxes+values`, `#%app`, `#%expression`, `#%top`, and `#%variable-reference` are added to *stop-ids*. If `#%app` or `#%datum` appears in *stop-ids*, then application and literal data expressions without the respective explicit form are not wrapped with the explicit form, and `#%top` wrappers are never added (even with an empty *stop-ids* list). If *stop-ids* is `#f` instead of a list, then *stx* is expanded only as long as the outermost form of *stx* is a macro (i.e., expansion does not proceed to sub-expressions). A fully expanded form can include the bindings listed in §1.2.3.1 "Fully Expanded Programs" plus the `letrec-syntaxes+values` form and `#%expression` in any expression position.

When `#%plain-module-begin` is not itself in *stop-ids* and `module*` is in *stop-ids*, then the `#%plain-module-begin` transformer refrains from expanding `module*` sub-forms. Otherwise, the `#%plain-module-begin` transformer detects and expands sub-forms (such as `define-values`) independent of the corresponding identifier's presence in *stop-ids*.

When *context-v* is `'module-begin`, and the result of expansion is a `#%plain-module-begin` form, then a `'submodule` syntax property is added to each enclosed `module` form (but not `module*` forms) in the same way as by `module` expansion.

The optional *intdef-ctx* argument must be either `#f`, the result of `syntax-local-make-definition-context`, or a list of such results. In the latter two cases, lexical information for internal definitions is added to *stx* before it is expanded (in reverse order relative to the list). The lexical information is also added to the expansion result (because the expansion might introduce bindings or references to internal-definition bindings).

For a particular internal-definition context, generate a unique value and put it into a list for *context-v*. To allow liberal expansion of `define` forms, the generated value should be an instance of a structure with a true value for `prop:liberal-define-context`. If the internal-definition context is meant to be self-contained, the list for *context-v* should contain only the generated value; if the internal-definition context is meant to splice into an immediately enclosing context, then when `syntax-local-context` produces a list, `cons` the generated value onto that list.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

Examples:

```
> (define-syntax-rule (do-print x ...)
    (printf x ...))

> (define-syntax-rule (hello x)
    (do-print "hello ~a" x))

> (define-syntax (show stx)
    (syntax-case stx ()
      [(_ x)
       (let ([partly (local-expand #'(hello x)
                                   'expression
                                   (list #'do-print))]
             [fully (local-expand #'(hello x)
                                  'expression
                                  #f)])
         (printf "partly expanded: ~s\n" (syntax->datum partly))
         (printf "fully expanded: ~s\n" (syntax->datum fully))
         fully)]))

> (show 1)
partly expanded: (do-print "hello ~a" 1)
fully expanded: (printf "hello ~a" 1)
hello 1
```

Changed in version 6.0.1.3 of package `base`: Changed treatment of `#%top` so that it is never introduced as an explicit wrapper.

```
(syntax-local-expand-expression stx) → syntax? syntax?
  stx : syntax?
```

Like `local-expand` given `'expression` and an empty stop list, but with two results: a syntax object for the fully expanded expression, and a syntax object whose content is opaque. The latter can be used in place of the former (perhaps in a larger expression produced by a macro transformer), and when the macro expander encounters the opaque object, it substitutes the fully expanded expression without re-expanding it; the `exn:fail:syntax` exception is raised if the expansion context includes bindings or marks that were not present for the original expansion, in which case re-expansion might produce different results. Consistent use of `syntax-local-expand-expression` and the opaque object thus avoids quadratic expansion times when local expansions are nested.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(local-transformer-expand  stx
                           context-v
                           stop-ids
                          [intdef-ctx]) → syntax?
  stx : syntax?
  context-v : (or/c 'expression 'top-level list?)
  stop-ids : (or/c (listof identifier?) #f)
  intdef-ctx : (or/c internal-definition-context? #f) = #f
```

Like `local-expand`, but *stx* is expanded as a transformer expression instead of a run-time expression. For `'expression` expansion, any lifted expressions—from calls to `syntax-local-lift-expression` during the expansion of *stx*—are captured into a `let-values` form in the result.

```
(local-expand/capture-lifts  stx
                             context-v
                             stop-ids
                            [intdef-ctx
                             lift-ctx]) → syntax?
  stx : syntax?
  context-v : (or/c 'expression 'top-level 'module 'module-begin list?)
  stop-ids : (or/c (listof identifier?) #f)
  intdef-ctx : (or/c internal-definition-context? #f) = #f
  lift-ctx : any/c = (gensym 'lifts)
```

Like `local-expand`, but the result is a syntax object that represents a `begin` expression. Lifted expressions—from calls to `syntax-local-lift-expression` during the expansion of *stx*—appear with their identifiers in `define-values` forms, and the expansion of *stx* is the last expression in the `begin`. The *lift-ctx* value is reported by `syntax-local-lift-context` during local expansion. The lifted expressions are not expanded, but instead left as provided in the `begin` form.

```
(local-transformer-expand/capture-lifts  stx
                                         context-v
                                         stop-ids
                                        [intdef-ctx
                                         lift-ctx]) → syntax?
  stx : syntax?
  context-v : (or/c 'expression 'top-level list?)
  stop-ids : (or/c (listof identifier?) #f)
  intdef-ctx : (or/c internal-definition-context? #f) = #f
  lift-ctx : any/c = (gensym 'lifts)
```

Like `local-expand/capture-lifts`, but *stx* is expanded as a transformer expression instead of a run-time expression. Lifted expressions are reported as `define-values` forms (in the transformer environment).

```
(internal-definition-context? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is an internal-definition context, `#f` otherwise.

```
(syntax-local-make-definition-context [intdef-ctx])
 → internal-definition-context?
  intdef-ctx : (or/c internal-definition-context? #f) = #f
```

Creates an opaque internal-definition context value to be used with `local-expand` and other functions. A transformer should create one context for each set of internal definitions to be expanded, and use it when expanding any form whose lexical context should include the definitions. After discovering an internal `define-values` or `define-syntaxes` form, use `syntax-local-bind-syntaxes` to add bindings to the context. Finally, the transformer must call `internal-definition-context-seal` after all bindings have been added; if an unsealed internal-definition context is detected in a fully expanded expression, the `exn:fail:contract` exception is raised.

If *intdef-ctx* is not `#f`, then the new internal-definition context extends the given one. That is, expanding in the new internal-definition context can use bindings previously introduced into *intdef-ctx*.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(syntax-local-bind-syntaxes id-list
                            expr
                            intdef-ctx) → void?
  id-list : (listof identifier?)
  expr : (or/c syntax? #f)
  intdef-ctx : internal-definition-context?
```

Binds each identifier in *id-list* within the internal-definition context represented by *intdef-ctx*, where *intdef-ctx* is the result of `syntax-local-make-definition-context`. Supply `#f` for *expr* when the identifiers correspond to `define-values` bindings, and supply a compile-time expression when the identifiers correspond to `define-syntaxes` bindings; in the latter case, the number of values produced by the expression should match the number of identifiers, otherwise the `exn:fail:contract:arity` exception is raised.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(internal-definition-context-seal intdef-ctx) → void?
  intdef-ctx : internal-definition-context?
```

Indicates that no further bindings will be added to *intdef-ctx*, which must not be sealed already. See also syntax-local-make-definition-context.

```
(identifier-remove-from-definition-context id-stx
                                           intdef-ctx)
 → identifier?
  id-stx : identifier?
  intdef-ctx : (or/c internal-definition-context?
                     (listof internal-definition-context?))
```

Removes *intdef-ctx* (or each identifier in the list) from the lexical information of *id-stx*. This operation is useful for correlating an identifier that is bound in an internal-definition context with its binding before the internal-definition context was created.

If simply removing the contexts produces a different binding than completely ignoring the contexts (due to nested internal definition contexts, for example), then the resulting identifier is given a syntax mark to simulate a non-existent lexical context. The *intdef-ctx* argument can be a list because removing internal-definition contexts one at a time can produce a different intermediate binding than removing them all at once.

```
(syntax-local-value  id-stx
                    [failure-thunk
                     intdef-ctx])  → any
  id-stx : syntax?
  failure-thunk : (or/c (-> any) #f) = #f
  intdef-ctx : (or/c internal-definition-context?
                     #f)                          = #f
```

Returns the transformer binding value of *id-stx* in either the context associated with *intdef-ctx* (if not #f) or the context of the expression being expanded (if *intdef-ctx* is #f). If *intdef-ctx* is provided, it must be an extension of the context of the expression being expanded.

If *id-stx* is bound to a rename transformer created with make-rename-transformer, syntax-local-value effectively calls itself with the target of the rename and returns that result, instead of the rename transformer.

If *id-stx* has no transformer binding (via define-syntax, let-syntax, etc.) in that environment, the result is obtained by applying *failure-thunk* if not #f. If *failure-thunk* is false, the exn:fail:contract exception is raised.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see syntax-transforming?), otherwise the exn:fail:contract exception is raised.

Examples:

```
> (define-syntax swiss-cheeses? #t)

> (define-syntax (transformer stx)
    (if (syntax-local-value #'swiss-cheeses?)
        #''(gruyère emmental raclette)
        #''(roquefort camembert boursin)))

> (transformer)
'(gruyère emmental raclette)
```

Examples:

```
> (define-syntax (transformer-2 stx)
    (syntax-local-value #'something-else (λ () (error "no bind-
ing"))))

> (transformer-2)
no binding
```

Examples:

```
> (define-syntax nachos #'(printf "nachos~n"))

> (define-syntax chips (make-rename-transformer #'nachos))

> (define-syntax (transformer-3 stx)
    (syntax-local-value #'chips))

> (transformer-3)
nachos
```

```
(syntax-local-value/immediate id-stx
                              [failure-thunk
                               intdef-ctx])  → any
  id-stx : syntax?
  failure-thunk : (or/c (-> any) #f) = #f
  intdef-ctx : (or/c internal-definition-context?
                     #f)                      = #f
```

Like `syntax-local-value`, but the result is normally two values. If *id-stx* is bound to a rename transformer, the results are the rename transformer and the identifier in the transformer. If *id-stx* is not bound to a rename transformer, then the results are the value that `syntax-local-value` would produce and `#f`.

If *id-stx* has no transformer binding, then *failure-thunk* is called (and it can return any number of values), or an exception is raised if *failure-thunk* is `#f`.

Beware that `provide` on an *id* bound to a rename transformer may export the target of the rename instead of *id*. See `make-rename-transformer` for more information.

```
(syntax-local-lift-expression stx) → identifier?
  stx : syntax?
```

Returns a fresh identifier, and cooperates with the `module`, `letrec-syntaxes+values`, `define-syntaxes`, `begin-for-syntax`, and top-level expanders to bind the generated identifier to the expression *stx*.

A run-time expression within a module is lifted to the module's top level, just before the expression whose expansion requests the lift. Similarly, a run-time expression outside of a module is lifted to a top-level definition. A compile-time expression in a `letrec-syntaxes+values` or `define-syntaxes` binding is lifted to a `let` wrapper around the corresponding right-hand side of the binding. A compile-time expression within `begin-for-syntax` is lifted to a `define` declaration just before the requesting expression within the `begin-for-syntax`.

Other syntactic forms can capture lifts by using `local-expand/capture-lifts` or `local-transformer-expand/capture-lifts`.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(syntax-local-lift-values-expression n stx)
 → (listof identifier?)
  n : exact-nonnegative-integer?
  stx : syntax?
```

Like `syntax-local-lift-expression`, but binds the result to *n* identifiers, and returns a list of the *n* identifiers.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(syntax-local-lift-context) → any/c
```

Returns a value that represents the target for expressions lifted via `syntax-local-lift-expression`. That is, for different transformer calls for which this procedure returns the same value (as determined by `eq?`), lifted expressions for the two transformer are moved to the same place. Thus, the result is useful for caching lift information to avoid redundant lifts.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(syntax-local-lift-module-end-declaration stx) → void?
  stx : syntax?
```

Cooperates with the `module` form to insert *stx* as a top-level declaration at the end of the module currently being expanded. If the current expression being transformed is in phase level 0 and not in the module top-level, then *stx* is eventually expanded in an expression context. If the current expression being transformed is in a higher phase level (i.e., nested within some number of `begin-for-syntax`es within a module top-level), then the lifted declaration is placed at the very end of the module (under a suitable number of `begin-for-syntax`es), instead of merely the end of the enclosing `begin-for-syntax`.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

If the current expression being transformed is not within a `module` form (see `syntax-transforming-module-expression?`), then the `exn:fail:contract` exception is raised.

```
(syntax-local-lift-require raw-require-spec
                           stx)             → syntax?
  raw-require-spec : any/c
  stx : syntax?
```

Lifts a `#%require` form corresponding to *raw-require-spec* (either as a syntax object or datum) to the top-level or to the top of the module currently being expanded or to an enclosing `begin-for-syntax`..

The resulting syntax object is the same as *stx*, except that a fresh syntax mark is added. The same syntax mark is added to the lifted `#%require` form, so that the `#%require` form can bind uses of imported identifiers in the resulting syntax object (assuming that the lexical information of *stx* includes the binding environment into which the `#%require` is lifted).

If *raw-require-spec* and *stx* are part of the input to a transformer, then typically `syntax-local-introduce` should be applied to each before passing them to `syntax-local-lift-require`, and then `syntax-local-introduce` should be applied to the result of `syntax-local-lift-require`. Otherwise, marks added by the macro expander can prevent access to the new imports.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(syntax-local-lift-provide raw-provide-spec-stx) → void?
  raw-provide-spec-stx : syntax?
```

Lifts a `#%provide` form corresponding to *raw-provide-spec-stx* to the top of the module currently being expanded or to an enclosing `begin-for-syntax`.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

If the current expression being transformed is not within a `module` form (see `syntax-transforming-module-expression?`), then the `exn:fail:contract` exception is raised.

```
(syntax-local-name) → any/c
```

Returns an inferred name for the expression position being transformed, or `#f` if no such name is available. A name is normally a symbol or an identifier. See also §1.2.6 “Inferred Value Names”.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(syntax-local-context)
 → (or/c 'expression 'top-level 'module 'module-begin list?)
```

Returns an indication of the context for expansion that triggered a syntax transformer call. See §1.2.3.3 “Expansion Context” for more information on contexts.

The symbol results indicate that the expression is being expanded for an expression context, a top-level context, a module context, or a module-begin context.

A list result indicates expansion in an internal-definition context. The identity of the list’s first element (i.e., its `eq?`ness) reflects the identity of the internal-definition context; in particular two transformer expansions receive the same first value if and only if they are invoked for the same internal-definition context. Later values in the list similarly identify internal-definition contexts that are still being expanded, and that required the expansion of nested internal-definition contexts.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(syntax-local-phase-level) → exact-integer?
```

During the dynamic extent of a syntax transformer application by the expander, the result is the phase level of the form being expanded. Otherwise, the result is `0`.

Examples:

```
; a macro bound at phase 0
> (define-syntax (print-phase-level stx)
    (printf "phase level: ~a~n" (syntax-local-phase-level))
    #'(void))

> (require (for-meta 2 racket/base))

> (begin-for-syntax
    ; a macro bound at phase 1
    (define-syntax (print-phase-level stx)
      (printf "phase level: ~a~n" (syntax-local-phase-level))
      #'(void)))

> (print-phase-level)
phase level: 0

> (begin-for-syntax (print-phase-level))
phase level: 1
```

```
(syntax-local-module-exports mod-path)
 → (listof (cons/c (or/c exact-integer? #f) (listof symbol?)))
             (or/c module-path?
                   (and/c syntax?
  mod-path :                (lambda (stx)
                              (module-path? (syntax->datum stx)))))
```

Returns an association list from phase-level numbers (or `#f` for the label phase level) to lists of symbols, where the symbols are the names of `provided` bindings from *mod-path* at the corresponding phase level.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(syntax-local-submodules) → (listof symbol?)
```

Returns a list of submodule names that are declared via `module` (as opposed to `module*`) in the current expansion context.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see `syntax-transforming?`), otherwise the `exn:fail:contract` exception is raised.

```
(syntax-local-get-shadower id-stx) → identifier?
  id-stx : identifier?
```

Returns *id-stx* if no binding in the current expansion context shadows *id-stx* (ignoring unsealed internal-definition contexts and identifiers that had the 'unshadowable syntax property), if *id-stx* has no module bindings in its lexical information, and if the current expansion context is not a module context.

If a binding of inner-identifier shadows *id-stx*, the result is the same as (syntax-local-get-shadower inner-identifier), except that it has the location and properties of *id-stx*. When searching for a shadowing binding, bindings from unsealed internal-definition contexts are ignored.

Otherwise, the result is the same as *id-stx* with its module bindings (if any) removed from its lexical information, and the lexical information of the current module context (if any) added.

Thus, the result is an identifier corresponding to the innermost shadowing of *id-stx* in the current context if it is shadowed, and a module-contextless version of *id-stx* otherwise.

If *id-stx* is tainted or armed, then the resulting identifier is tainted.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see syntax-transforming?), otherwise the exn:fail:contract exception is raised.

```
(syntax-local-certifier [active?])
  → ((syntax?) (any/c (or/c procedure? #f))
     . ->* . syntax?)
  active? : boolean? = #f
```

For backward compatibility only; returns a procedure that returns its first argument.

```
(syntax-transforming?) → boolean?
```

Returns #t during the dynamic extent of a syntax transformer application by the expander and while a module is being visited, #f otherwise.

```
(syntax-transforming-module-expression?) → boolean?
```

Returns #t during the dynamic extent of a syntax transformer application by the expander for an expression within a module form, #f otherwise.

```
(syntax-local-introduce stx) → syntax?
  stx : syntax?
```

Produces a syntax object that is like *stx*, except that a syntax mark for the current expansion is added (possibly canceling an existing mark in parts of *stx*). See §1.2.3.5 “Transformer Bindings” for information on syntax marks.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see syntax-transforming?), otherwise the exn:fail:contract exception is raised.

```
(make-syntax-introducer) → (syntax? . -> . syntax?)
```

Produces a procedure that behaves like syntax-local-introduce, but using a fresh syntax mark. Multiple applications of the same make-syntax-introducer result procedure use the same mark, and different result procedures use distinct marks.

```
(make-syntax-delta-introducer  ext-stx
                               base-stx
                              [phase-level])
 → (syntax? . -> . syntax?)
  ext-stx : syntax?
  base-stx : (or/c syntax? #f)
  phase-level : (or/c #f exact-integer?)
              = (syntax-local-phase-level)
```

Produces a procedure that behaves like syntax-local-introduce, but using the syntax marks of *ext-stx* that are not shared with *base-stx*. If *ext-stx* does not extend the set of marks in *base-stx* or if *base-stx* is #f, and if *ext-stx* has a module binding in the phase level indicated by *phase-level*, then any marks of *ext-stx* that would be needed to preserve its binding are not transferred in an introduction.

This procedure is potentially useful when *m-id* has a transformer binding that records some *orig-id*, and a use of *m-id* introduces a binding of *orig-id*. In that case, the syntax marks in the use of *m-id* since the binding of *m-id* should be transferred to the binding instance of *orig-id*, so that it captures uses with the same lexical context as the use of *m-id*.

More typically, however, syntax-local-make-delta-introducer should be used, since it cooperates with rename transformers.

If *ext-stx* is tainted or armed, then an identifier result from the created procedure is tainted.

```
(syntax-local-make-delta-introducer id)
 → (identifier? . -> . identifier?)
  id : identifier?
```

Determines the binding of *id*. If the binding is not a rename transformer, the result is an introducer as created by make-syntax-delta-introducer using *id* and the binding of *id* in the environment of expansion. If the binding is a rename transformer, then the introducer is one composed with the target of the rename transformer and its binding. Furthermore, the *delta-introduce* functions associated with the rename transformers (supplied as the

second argument to make-rename-transformer) are composed (in first-to-last order) before the introducers created with make-syntax-delta-introducer (which are composed last-to-first).

The exn:fail:contract exception is raised if *id* or any identifier in its rename-transformer chain has no binding.

This procedure must be called during the dynamic extent of a syntax transformer application by the expander or while a module is visited (see syntax-transforming?), otherwise the exn:fail:contract exception is raised.

```
(syntax-local-transforming-module-provides?) → boolean?
```

Returns #t while a provide transformer is running (see make-provide-transformer) or while an expand sub-form of #%provide is expanded, #f otherwise.

```
(syntax-local-module-defined-identifiers)
 → (and/c hash? immutable?)
```

Can be called only while syntax-local-transforming-module-provides? returns #t.

It returns a hash table mapping a phase-level number (such as 0) to a list of all definitions at that phase level within the module being expanded. This information is used for implementing provide sub-forms like all-defined-out.

Beware that the phase-level keys are absolute relative to the enclosing module, and not relative to the current transformer phase level as reported by syntax-local-phase-level.

```
(syntax-local-module-required-identifiers mod-path
                                          phase-level)
  (listof (cons/c (or/c exact-integer? #f)
→                 (listof identifier?)))
 mod-path : (or/c module-path? #f)
 phase-level : (or/c exact-integer? #f #t)
```

Can be called only while syntax-local-transforming-module-provides? returns #t.

It returns an association list mapping phase levels to lists of identifiers. Each list of identifiers includes all bindings imported (into the module being expanded) using the module path *mod-path*, or all modules if *mod-path* is #f. The association list includes all identifiers imported with a *phase-level* shift, of all shifts if *phase-level* is #t.

When an identifier is renamed on import, the result association list includes the identifier by its internal name. Use identifier-binding to obtain more information about the identifier.

Beware that the phase-level keys are absolute relative to the enclosing module, and not relative to the current transformer phase level as reported by syntax-local-phase-level.

```
prop:liberal-define-context : struct-type-property?
(liberal-define-context? v) → boolean?
  v : any/c
```

An instance of a structure type with a true value for the `prop:liberal-define-context` property can be used as an element of an internal-definition context representation in the result of `syntax-local-context` or the second argument of `local-expand`. Such a value indicates that the context supports *liberal expansion* of `define` forms into potentially multiple `define-values` and `define-syntaxes` forms. The `'module` and `'module-body` contexts implicitly allow liberal expansion.

The `liberal-define-context?` predicate returns `#t` if *v* is an instance of a structure with a true value for the `prop:liberal-define-context` property, `#f` otherwise.

12.4.1 `require` Transformers

`(require racket/require-transform)` package: `base`

The bindings documented in this section are provided by the `racket/require-transform` library, not `racket/base` or `racket`.

A transformer binding whose value is a structure with the `prop:require-transformer` property implements a derived *require-spec* for `require` as a *require transformer*.

A require transformer is called with the syntax object representing its use as a *require-spec* within a `require` form, and the result must be two lists: a list of `import`s and a list of `import-source`s.

If the derived form contains a sub-form that is a *require-spec*, then it can call `expand-import` to transform the sub-*require-spec* to lists of imports and import sources.

See also `define-require-syntax`, which supports macro-style `require` transformers.

```
(expand-import stx) → (listof import?) (listof import-source?)
  stx : syntax?
```

Expands the given *require-spec* to lists of imports and import sources. The latter specifies modules to be instantiated or visited, so the modules that it represents should be a superset of the modules represented in the former list (so that a module will be instantiated or visited even if all of imports are eventually filtered from the former list).

```
(make-require-transformer proc) → require-transformer?
        (syntax? . -> . (values
  proc :                  (listof import?)
                          (listof import-source?)))
```

Creates a require transformer using the given procedure as the transformer.

```
prop:require-transformer : struct-type-property?
```

A property to identify require transformers. The property value must be a procedure that takes the structure and returns a transformer procedure; the returned transformer procedure takes a syntax object and returns import and import-source lists.

```
(require-transformer? v) → boolean?
  v : any/c
```

Returns `#t` if `v` has the `prop:require-transformer` property, `#f` otherwise.

```
(struct import (local-id
                src-sym
                src-mod-path
                mode
                req-mode
                orig-mode
                orig-stx)
  #:extra-constructor-name make-import)
 local-id : identifier?
 src-sym : symbol?
 src-mod-path : (or/c module-path?
                      (and/c syntax?
                             (lambda (stx)
                               (module-path? (syntax->datum stx)))))
 mode : (or/c exact-integer? #f)
 req-mode : (or/c exact-integer? #f)
 orig-mode : (or/c exact-integer? #f)
 orig-stx : syntax?
```

A structure representing a single imported identifier:

- `local-id` — the identifier to be bound within the importing module.
- `src-sym` — the external name of the binding as exported from its source module.
- `src-mod-path` — a module path (relative to the importing module) for the source of the imported binding.
- `orig-stx` — a syntax object for the source of the import, used for error reporting.
- `mode` — the phase level of the binding in the importing module.
- `req-mode` — the phase level shift of the import relative to the exporting module.

- `orig-mode` — the phase level of the binding as exported by the exporting module.

```
(struct import-source (mod-path-stx mode)
    #:extra-constructor-name make-import-source)
                (and/c syntax?
  mod-path-stx :        (lambda (x)
                          (module-path? (syntax->datum x))))
  mode : (or/c exact-integer? #f)
```

A structure representing an imported module, which must be instantiated or visited even if no binding is imported into a module.

- `mod-path-stx` — a module path (relative to the importing module) for the source of the imported binding.
- `mode` — the phase level shift of the import.

```
(current-require-module-path) → (or/c #f module-path-index?)
(current-require-module-path module-path) → void?
  module-path : (or/c #f module-path-index?)
```

A parameter that determines how relative `require`-level module paths are expanded to `#%require`-level module paths by `convert-relative-module-path` (which is used implicitly by all built-in `require` sub-forms).

When the value of `current-require-module-path` is `#f`, relative module paths are left as-is, which means that the `require` context determines the resolution of the module path.

The `require` form parameterizes `current-require-module-path` as `#f` while invoking sub-form transformers, while `relative-in` parameterizes to a given module path.

```
(convert-relative-module-path module-path)
    (or/c module-path?
          (and/c syntax?
 →                (lambda (stx)
                    (module-path? (syntax-e stx)))))
                (or/c module-path?
                      (and/c syntax?
  module-path :               (lambda (stx)
                                (module-path? (syntax-e stx)))))
```

Converts *module-path* according to `current-require-module-path`.

If *module-path* is not relative or if the value of `current-require-module-path` is `#f`, then *module-path* is returned. Otherwise, *module-path* is converted to an absolute module path that is equivalent to *module-path* relative to the value of `current-require-module-path`.

```
(syntax-local-require-certifier)
  → ((syntax?) (or/c #f (syntax? . -> . syntax?))
     . ->* . syntax?)
```

For backward compatibility only; returns a procedure that returns its first argument.

12.4.2 `provide` Transformers

```
(require racket/provide-transform)
```
package: `base`

The bindings documented in this section are provided by the `racket/provide-transform` library, not `racket/base` or `racket`.

A transformer binding whose value is a structure with the `prop:provide-transformer` property implements a derived *provide-spec* for `provide` as a *provide transformer*. A provide transformer is applied as part of the last phase of a module's expansion, after all other declarations and expressions within the module are expanded.

A transformer binding whose value is a structure with the `prop:provide-pre-transformer` property implements a derived *provide-spec* for `provide` as a *provide pre-transformer*. A provide pre-transformer is applied as part of the first phase of a module's expansion. Since it is used in the first phase, a provide pre-transformer can use functions such as `syntax-local-lift-expression` to introduce expressions and definitions in the enclosing module.

An identifier can have a transformer binding to a value that acts both as a provide transformer and provide pre-transformer. The result of a provide pre-transformer is *not* automatically re-expanded, so a provide pre-transformer can usefully expand to itself in that case.

A transformer is called with the syntax object representing its use as a *provide-spec* within a `provide` form and a list of symbols representing the export modes specified by enclosing *provide-spec*s. The result of a provide transformer must be a list of `export`s, while the result of a provide pre-transformer is a syntax object to be used as a *provide-spec* in the last phase of module expansion.

If a derived form contains a sub-form that is a *provide-spec*, then it can call `expand-export` or `pre-expand-export` to transform the sub-*provide-spec* sub-form.

See also `define-provide-syntax`, which supports macro-style provide transformers.

```
(expand-export stx modes) → (listof export?)
  stx : syntax?
  modes : (listof (or/c exact-integer? #f))
```

Expands the given *provide-spec* to a list of exports. The *modes* list controls the expansion

of sub-*provide-specs*; for example, an identifier refers to a binding in the phase level of the enclosing `provide` form, unless the *modes* list specifies otherwise. Normally, *modes* is either empty or contains a single element.

```
(pre-expand-export stx modes) → syntax?
  stx : syntax?
  modes : (listof (or/c exact-integer? #f))
```

Expands the given *provide-spec* at the level of provide pre-transformers. The *modes* argument is the same as for `expand-export`.

```
(make-provide-transformer proc) → provide-transformer?
          (syntax? (listof (or/c exact-integer? #f))
  proc :   . -> . (listof export?))
(make-provide-transformer proc pre-proc)
 → (and/c provide-transformer? provide-pre-transformer?)
          (syntax? (listof (or/c exact-integer? #f))
  proc :   . -> . (listof export?))
              (syntax? (listof (or/c exact-integer? #f))
  pre-proc :   . -> . syntax?)
```

Creates a provide transformer (i.e., a structure with the `prop:provide-transformer` property) using the given procedure as the transformer. If a *pre-proc* is provided, then the result is also a provide pre-transformer.

```
(make-provide-pre-transformer pre-proc)
 → provide-pre-transformer?
              (syntax? (listof (or/c exact-integer? #f))
  pre-proc :   . -> . syntax?)
```

Like `make-provide-transformer`, but for a value that is a provide pre-transformer, only.

```
prop:provide-transformer : struct-type-property?
```

A property to identify provide transformers. The property value must be a procedure that takes the structure and returns a transformer procedure; the returned transformer procedure takes a syntax object and mode list and returns an export list.

```
prop:provide-pre-transformer : struct-type-property?
```

A property to identify provide pre-transformers. The property value must be a procedure that takes the structure and returns a transformer procedure; the returned transformer procedure takes a syntax object and mode list and returns a syntax object.

```
(provide-transformer? v) → boolean?
  v : any/c
```

Returns `#t` if *v* has the `prop:provide-transformer` property, `#f` otherwise.

```
(provide-pre-transformer? v) → boolean?
  v : any/c
```

Returns `#t` if *v* has the `prop:provide-pre-transformer` property, `#f` otherwise.

```
(struct export (local-id out-sym mode protect? orig-stx)
    #:extra-constructor-name make-export)
  local-id : identifier?
  out-sym : symbol?
  mode : (or/c exact-integer? #f)
  protect? : any/c
  orig-stx : syntax?
```

A structure representing a single imported identifier:

- `local-id` — the identifier that is bound within the exporting module.
- `out-sym` — the external name of the binding.
- `orig-stx` — a syntax object for the source of the export, used for error reporting.
- `protect?` — indicates whether the identifier should be protected (see §14.10 “Code Inspectors”).
- `mode` — the phase level of the binding in the exporting module.

```
(syntax-local-provide-certifier)
 → ((syntax?) (or/c #f (syntax? . -> . syntax?))
    . ->* . syntax?)
```

For backward compatibility only; returns a procedure that returns its first argument.

12.4.3 Keyword-Argument Conversion Introspection

`(require racket/keyword-transform)` package: `base`

The bindings documented in this section are provided by the `racket/keyword-transform` library, not `racket/base` or `racket`.

```
(syntax-procedure-alias-property stx)
    (or/c #f
          (letrec ([val? (recursive-contract
 →                        (or/c (cons/c identifier? identifier?)
                                (cons/c val? val?)))])
            val?))
  stx : syntax?
```

```
(syntax-procedure-converted-arguments-property stx)
   (or/c #f
         (letrec ([val? (recursive-contract
→                        (or/c (cons/c identifier? identifier?)
                              (cons/c val? val?)))])
           val?))
  stx : syntax?
```

Reports the value of a syntax property that can be attached to an identifier by the expansion of a keyword-application form. See `lambda` for more information about the property.

The property value is normally a pair consisting of the original identifier and an identifier that appears in the expansion. Property-value merging via `syntax-track-origin` can make the value a pair of such values, and so on.

12.5 Syntax Parameters

```
(require racket/stxparam)
```
package: `base`

The bindings documented in this section are provided by the `racket/stxparam` library, not `racket/base` or `racket`.

```
(define-syntax-parameter id expr)
```

Binds *id* as syntax to a *syntax parameter*. The *expr* is an expression in the transformer environment that serves as the default value for the syntax parameter. The value is typically obtained by a transformer using `syntax-parameter-value`.

The *id* can be used with `syntax-parameterize` or `syntax-parameter-value` (in a transformer). If *expr* produces a procedure of one argument or a `make-set!-transformer` result, then *id* can be used as a macro. If *expr* produces a `make-rename-transformer` result, then *id* can be used as a macro that expands to a use of the target identifier, but `syntax-local-value` of *id* does not produce the target's value.

Examples:

```
> (define-syntax-parameter current-class #f)

> (define-syntax-parameter yield (make-rename-
transformer #'abort))

> (define-syntax-parameter define/public
    (λ (stx)
      (raise-syntax-error #f "use of a class keyword not in a
class" stx)))
```

```
> (begin-for-syntax (displayln (syntax-parameter-value #'current-
class)))
#f

> (yield 5)
5
```

(syntax-parameterize ([*id expr*] ...) *body-expr* ...+)

See also splicing-syntax-parameteriz

Each *id* must be bound to a syntax parameter using define-syntax-parameter. Each *expr* is an expression in the transformer environment. During the expansion of the *body-expr*s, the value of each *expr* is bound to the corresponding *id*.

If an *expr* produces a procedure of one argument or a make-set!-transformer result, then its *id* can be used as a macro during the expansion of the *body-expr*s. If *expr* produces a make-rename-transformer result, then *id* can be used as a macro that expands to a use of the target identifier, but syntax-local-value of *id* does not produce the target's value.

Examples:

```
> (define-syntax-parameter abort (syntax-rules ()))

> (define-syntax forever
    (syntax-rules ()
      [(forever body ...)
       (call/cc (lambda (abort-k)
         (syntax-parameterize
             ([abort (syntax-rules () [(_) (abort-k)])])
           (let loop () body ... (loop)))))]))

> (define-syntax-parameter it (syntax-rules ()))

> (define-syntax aif
    (syntax-rules ()
      [(aif test then else)
       (let ([t test])
         (syntax-parameterize ([it (syntax-id-rules () [_ t])])
           (if t then else)))]))
```

12.5.1 Syntax Parameter Inspection

(require racket/stxparam-exptime) package: base

```
(syntax-parameter-value id-stx) → any
  id-stx : syntax?
```

This procedure is intended for use in a transformer environment, where *id-stx* is an identifier bound in the normal environment to a syntax parameter. The result is the current value of the syntax parameter, as adjusted by `syntax-parameterize` form.

This binding is provided `for-syntax` by `racket/stxparam`, since it is normally used in a transformer. It is provided normally by `racket/stxparam-exptime`.

```
(make-parameter-rename-transformer id-stx) → any
  id-stx : syntax?
```

This procedure is intended for use in a transformer, where *id-stx* is an identifier bound to a syntax parameter. The result is a transformer that behaves as *id-stx*, but that cannot be used with `syntax-parameterize` or `syntax-parameter-value`.

Using `make-parameter-rename-transformer` is analogous to defining a procedure that calls a parameter. Such a procedure can be exported to others to allow access to the parameter value, but not to change the parameter value. Similarly, `make-parameter-rename-transformer` allows a syntax parameter to be used as a macro, but not changed.

The result of `make-parameter-rename-transformer` is not treated specially by `syntax-local-value`, unlike the result of `make-rename-transformer`.

This binding is provided `for-syntax` by `racket/stxparam`, since it is normally used in a transformer. It is provided normally by `racket/stxparam-exptime`.

12.6 Local Binding with Splicing Body

`(require racket/splicing)` package: `base`

The bindings documented in this section are provided by the `racket/splicing` library, not `racket/base` or `racket`.

```
splicing-let
splicing-letrec
splicing-let-values
splicing-letrec-values
splicing-let-syntax
splicing-letrec-syntax
splicing-let-syntaxes
splicing-letrec-syntaxes
splicing-letrec-syntaxes+values
splicing-local
```

Like `let`, `letrec`, `let-values`, `letrec-values`, `let-syntax`, `letrec-syntax`, `let-syntaxes`, `letrec-syntaxes`, `letrec-syntaxes+values`, and `local`, except that in a definition context, the body forms are spliced into the enclosing definition context (in the same way as for `begin`).

Examples:

```
> (splicing-let-syntax ([one (lambda (stx) #'1)])
    (define o one))

> o
1
> one
one: undefined;
 cannot reference undefined identifier
```

When a splicing binding form occurs in a top-level context or module context, its local bindings are treated similarly to definitions. In particular, if a reference to one of the splicing form's bound variables is evaluated before the variable is initialized, an unbound variable error is raised, instead of the variable evaluating to the undefined value. Also, syntax bindings are evaluated every time the module is visited, instead of only once during compilation as in `let-syntax`, etc.

Example:

```
> (splicing-letrec ([x bad]
                    [bad 1])
    x)
bad.3: undefined;
 cannot reference undefined identifier
```

`splicing-syntax-parameterize`

Like `syntax-parameterize`, except that in a definition context, the body forms are spliced into the enclosing definition context (in the same way as for `begin`). In a definition context, the body of `splicing-syntax-parameterize` can be empty.

Note that require transformers and provide transformers are not affected by syntax parameterization. While all uses of `require` and `provide` will be spliced into the enclosing context, derived import or export specifications will expand as if they had not been inside of the `splicing-syntax-parameterize`.

Examples:

```
> (define-syntax-parameter place (lambda (stx) #'"Kansas"))

> (define-syntax-rule (where) `(at ,(place)))
```

```
> (where)
'(at "Kansas")
> (splicing-syntax-parameterize ([place (lambda (stx) #'"Oz")])
    (define here (where)))

> here
'(at "Oz")
```

12.7 Syntax Object Properties

Every syntax object has an associated *syntax property* list, which can be queried or extended with `syntax-property`. Properties are not preserved for a `syntax-quoted` syntax object in a compiled form that is marshaled to a byte string.

In `read-syntax`, the reader attaches a `'paren-shape` property to any pair or vector syntax object generated from parsing a pair `[` and `]` or `{` and `}`; the property value is `#\[` in the former case, and `#\{` in the latter case. The `syntax` form copies any `'paren-shape` property from the source of a template to corresponding generated syntax.

Both the syntax input to a transformer and the syntax result of a transformer may have associated properties. The two sets of properties are merged by the syntax expander: each property in the original and not present in the result is copied to the result, and the values of properties present in both are combined with `cons` (result value first, original value second).

Before performing the merge, however, the syntax expander automatically adds a property to the original syntax object using the key `'origin`. If the source syntax has no `'origin` property, it is set to the empty list. Then, still before the merge, the identifier that triggered the macro expansion (as syntax) is `cons`ed onto the `'origin` property so far. The `'origin` property thus records (in reverse order) the sequence of macro expansions that produced an expanded expression. Usually, the `'origin` value is an immutable list of identifiers. However, a transformer might return syntax that has already been expanded, in which case an `'origin` list can contain other lists after a merge. The `syntax-track-origin` procedure implements this tracking.

Besides `'origin` tracking for general macro expansion, Racket adds properties to expanded syntax (often using `syntax-track-origin`) to record additional expansion details:

- When a `begin` form is spliced into a sequence with internal definitions (see §1.2.3.7 “Internal Definitions”), `syntax-track-origin` is applied to every spliced element from the `begin` body. The second argument to `syntax-track-origin` is the `begin` form, and the third argument is the `begin` keyword (extracted from the spliced form).
- When an internal `define-values` or `define-syntaxes` form is converted into a `letrec-syntaxes+values` form (see §1.2.3.7 “Internal Definitions”), `syntax-`

`track-origin` is applied to each generated binding clause. The second argument to `syntax-track-origin` is the converted form, and the third argument is the `define-values` or `define-syntaxes` keyword form the converted form.

- When a `letrec-syntaxes+values` expression is fully expanded, syntax bindings disappear, and the result is either a `letrec-values` form (if the unexpanded form contained non-syntax bindings), or only the body of the `letrec-syntaxes+values` form (wrapped with `begin` if the body contained multiple expressions). To record the disappeared syntax bindings, a property is added to the expansion result: an immutable list of identifiers from the disappeared bindings, as a `'disappeared-binding` property.

- When a subtyping `struct` form is expanded, the identifier used to reference the base type does not appear in the expansion. Therefore, the `struct` transformer adds the identifier to the expansion result as a `'disappeared-use` property.

- When a reference to an unexported or protected identifier from a module is discovered, the `'protected` property is added to the identifier with a `#t` value.

- When `read-syntax` generates a syntax object, it attaches a property to the object (using a private key) to mark the object as originating from a read. The `syntax-original?` predicate looks for the property to recognize such syntax objects. (See §12.2 “Syntax Object Content” for more information.)

See §12.9.1 “Information on Expanded Modules” for information about properties generated by the expansion of a module declaration. See `lambda` and §1.2.6 “Inferred Value Names” for information about properties recognized when compiling a procedure. See `current-compile` for information on properties and byte codes.

```
(syntax-property stx key v) → syntax?
  stx : syntax?
  key : any/c
  v : any/c
(syntax-property stx key) → any
  stx : syntax?
  key : any/c
```

The three-argument form extends *stx* by associating an arbitrary property value *v* with the key *key*; the result is a new syntax object with the association (while *stx* itself is unchanged).

The two-argument form returns an arbitrary property value associated to *stx* with the key *key*, or `#f` if no value is associated to *stx* for *key*.

```
(syntax-property-symbol-keys stx) → list?
  stx : syntax?
```

Returns a list of all symbols that as keys have associated properties in *stx*. Uninterned symbols (see §4.6 “Symbols”) are not included in the result list.

```
(syntax-track-origin new-stx
                     orig-stx
                     id-stx)  → any
  new-stx : syntax?
  orig-stx : syntax?
  id-stx : syntax?
```

Adds properties to *new-stx* in the same way that macro expansion adds properties to a transformer result. In particular, it merges the properties of *orig-stx* into *new-stx*, first adding *id-stx* as an 'origin property, and it returns the property-extended syntax object. Use the syntax-track-origin procedure in a macro transformer that discards syntax (corresponding to *orig-stx* with a keyword *id-stx*) leaving some other syntax in its place (corresponding to *new-stx*).

For example, the expression

```
(or x y)
```

expands to

```
(let ([or-part x]) (if or-part or-part (or y)))
```

which, in turn, expands to

```
(let-values ([(or-part) x]) (if or-part or-part y))
```

The syntax object for the final expression will have an 'origin property whose value is (list (quote-syntax let) (quote-syntax or)).

12.8 Syntax Taints

§16.2.7 “Syntax Taints” in ***The Racket Guide*** introduces syntax taints.

The *tamper status* of a syntax object is either tainted, armed, or clean:

- A *tainted* identifier is rejected by the macro expander for use as either a binding or expression. If a syntax object is tainted, then any syntax object in the result of (syntax-e *stx*) is tainted, and datum->syntax with *stx* as its first argument produces a tainted syntax object.

  Other derived operations, such as pattern matching in syntax-case, also taint syntax objects when extracting them from a tainted syntax object.

- An *armed* syntax object has a set of *dye packs*, which creates taints if the armed syntax object is used without first disarming the dye packs. In particular, if a syntax object is armed, `syntax-e`, `datum->syntax`, `quote-syntax`, and derived operations effectively treat the syntax object as tainted. The macro expander, in contrast, disarms dye packs before pulling apart syntax objects.

  Each dye pack, which is added to a syntax object with the `syntax-arm` function, is keyed by an inspector. A dye pack can be *disarm*ed using `syntax-disarm` with an inspector that is the same as or a superior of the dye pack's inspector.

- A *clean* syntax object has no immediate taints or dye packs, although it may contain syntax objects that are tainted or armed.

Taints cannot be removed, and attempting to arm a syntax object that is already tainted has no effect on the resulting syntax object.

The macro expander disarms any syntax object that it encounters in an expression position or as a module body. A syntax object is therefore disarmed when it is provided to a syntax transformer. The transformer's result, however, is *rearm*ed by copying to it any dye packs that were originally attached to the transformer's input. The rearming process obeys the following rules:

- If the result has a `'taint-mode` property (see §12.7 "Syntax Object Properties") that is `'opaque`, then dye packs are attached to the immediate syntax object.

- If the result has a `'taint-mode` property that is `'none`, then no dye pack is attached to the syntax object. The `'none` mode is rarely appropriate.

- If the result has a `'taint-mode` property that is `'transparent`, then the dye packs are propagated recursively to syntax object that corresponds to elements of the syntax object's datum as a list (or, more precisely, to the `car`s of the datum as reached by any number of `cdr`s), and the immediate syntax object loses its lexical context; If the immediate syntax object is already armed, then recursive propagation taints the elements. Recursive propagation uses syntax properties and shapes, as for the immediate rearming.

- If the result has a `'taint-mode` property that is `'transparent-binding`, then dye packs are attached in a way similar to `'transparent`, but further treating the syntax object corresponding to the second list element as having a `'transparent` value for the `'taint-mode` property if it does not already have a `'taint-mode` property value.

- If the result has no `'taint-mode` property value, but its datum is a pair, and if the syntax object corresponding to the `car` of the pair is an identifier bound to `begin`, `module`, or `#%plain-module-begin`, then dye packs are propagated as if the syntax object had the `'transparent` property value.

- If the result has no `'taint-mode` property value, but its datum is a pair, and if the syntax object corresponding to the `car` of the pair is an identifier bound to `define-`

`values` or `define-syntaxes`, then dye packs are propagated as if the syntax object had the `'transparent-binding` property value.

For backward compatibility, a `'certify-mode` property is treated the same as a `'taint-mode` property if the former is not attached. To avoid accidental transfer of a `'taint-mode` or `'certify-mode` property value, the expander always removes any `'taint-mode` and `'certify-mode` property on a syntax object that is passed to a syntax transformer.

```
(syntax-tainted? stx) → boolean?
  stx : syntax?
```

Returns `#t` if *stx* is tainted, `#f` otherwise.

```
(syntax-arm stx [inspector use-mode?]) → syntax?
  stx : syntax?
  inspector : (or/c inspector? #f) = #f
  use-mode? : any/c = #f
```

Produces a syntax object like *stx*, but armed with a dye pack that is keyed by *inspector*.

A `#f` value for *inspector* is equivalent to an inspector that depends on the current dynamic context:

- when applying a syntax transformer is being applied, the declaration-time code inspector of the module in which a syntax transformer was bound;
- when a module is being visited, the module's declaration-time code inspector;
- `(current-code-inspector)`, otherwise.

If *use-mode?* is `#f`, then if *stx* is tainted or already armed with the key *inspector*, the result is *stx*.

If *use-mode?* is a true value, then a dye pack is not necessarily added directly to *stx*. Instead, the dye pack is pushed to interior syntax objects in the same way that the expander pushes armings into a syntax transformer's results when rearming (based on a `'taint-mode` syntax property or identifier bindings); see the expander's rearming rules for more information. To the degree that pushing dye packs into a syntax object must destructure *stx*, existing taints or dye packs can lead to tainted results rather than armed results.

```
(syntax-protect stx) → syntax?
  stx : syntax?
```

Equivalent to `(syntax-arm stx #f #t)`.

```
(syntax-disarm stx inspector) → syntax?
  stx : syntax?
  inspector : (or/c inspector? #f)
```

Produces a disarmed version of *stx*, removing any immediate dye packs that match *inspector*. An inspector matches when it is either the same as or a super-inspector of the dye pack's inspector. A `#f` value for *inspector* is replaced by a specific inspector in the same way as for `syntax-arm`.

```
(syntax-rearm stx from-stx [use-mode?]) → syntax?
  stx : syntax?
  from-stx : syntax?
  use-mode? : any/c = #f
```

Produces a rearmed or tainted version of *stx* by adding all immediate taints and dye packs of *from-stx*.

If *use-mode?* is a true value, *stx* is not necessarily tainted or armed directly. Instead, taints or dye packs are pushed to interior syntax objects in the same way as for `syntax-arm` or rearming by the expander.

```
(syntax-taint stx) → syntax?
  stx : syntax?
```

Returns tainted version of *stx*—equivalent to `(datum->syntax (syntax-arm stx) (syntax-e stx) stx stx)`—or *stx* if it is already tainted.

12.9 Expanding Top-Level Forms

```
(expand top-level-form) → syntax?
  top-level-form : any/c
```

Expands all non-primitive syntax in *top-level-form*, and returns a syntax object for the expanded form that contains only core forms, matching the grammar specified by §1.2.3.1 "Fully Expanded Programs".

Before *top-level-form* is expanded, its lexical context is enriched with `namespace-syntax-introduce`, just as for `eval`. Use `syntax->datum` to convert the returned syntax object into a printable datum.

Here's an example of using `expand` on a module:

```
(parameterize ([current-namespace (make-base-namespace)])
 (expand
  (datum->syntax
   #f
   '(module foo scheme
      (define a 3)
      (+ a 4)))))
```

Here's an example of using `expand` on a non-top-level form:

```
(define-namespace-anchor anchor)
(parameterize ([current-namespace
                (namespace-anchor->namespace anchor)])
 (expand
  (datum->syntax
   #f
   '(delay (+ 1 2)))))
```

(expand-syntax *stx*) → syntax?
  stx : syntax?

Like (expand *stx*), except that the argument must be a syntax object, and its lexical context is not enriched before expansion.

(expand-once *top-level-form*) → syntax?
  top-level-form : any/c

Partially expands *top-level-form* and returns a syntax object for the partially-expanded expression. Due to limitations in the expansion mechanism, some context information may be lost. In particular, calling `expand-once` on the result may produce a result that is different from expansion via `expand`.

Before *top-level-form* is expanded, its lexical context is enriched with `namespace-syntax-introduce`, as for `eval`.

(expand-syntax-once *stx*) → syntax?
  stx : syntax?

Like (expand-once *stx*), except that the argument must be a syntax object, and its lexical context is not enriched before expansion.

(expand-to-top-form *top-level-form*) → syntax?
  top-level-form : any/c

Partially expands *top-level-form* to reveal the outermost syntactic form. This partial expansion is mainly useful for detecting top-level uses of `begin`. Unlike the result of `expand-once`, expanding the result of `expand-to-top-form` with `expand` produces the same result as using `expand` on the original syntax.

Before `stx-or-sexpr` is expanded, its lexical context is enriched with `namespace-syntax-introduce`, as for `eval`.

```
(expand-syntax-to-top-form stx) → syntax?
  stx : syntax?
```

Like `(expand-to-top-form stx)`, except that the argument must be a syntax object, and its lexical context is not enriched before expansion.

12.9.1 Information on Expanded Modules

Information for an expanded `module` declaration is stored in a set of syntax properties (see §12.7 “Syntax Object Properties”) attached to the syntax object:

- `'module-direct-requires` — a list of module path indexes (or symbols) representing the modules explicitly imported into the module.
- `'module-direct-for-syntax-requires` — a list of module path indexes (or symbols) representing the modules explicitly for-syntax imported into the module.
- `'module-direct-for-template-requires` — a list of module path indexes (or symbols) representing the modules explicitly for-template imported into the module.
- `'module-variable-provides` — a list of provided items, where each item is one of the following:
  - `symbol` — represents a locally defined variable that is provided with its defined name.
  - `(cons provided-sym defined-sym)` — represents a locally defined variable that is provided with renaming; the first symbol is the exported name, and the second symbol is the defined name.
  - `(list* module-path-index provided-sym defined-sym)` — represents a re-exported and possibly re-named variable from the specified module; `module-path-index` is either a module path index or symbol (see §14.4.2 “Compiled Modules and References”), indicating the source module for the binding. The *provided-sym* is the external name for the re-export, and *defined-sym* is the originally defined name in the module specified by `module-path-index`.
- `'module-syntax-provides` — like `'module-variable-provides`, but for syntax exports instead of variable exports.

- `'module-indirect-provides` — a list of symbols for variables that are defined in the module but not exported; they may be exported indirectly through macro expansions. Definitions of macro-generated identifiers create uninterned symbols in this list.

12.10 File Inclusion

```
(require racket/include)
```
package: base

The bindings documented in this section are provided by the `racket/include` and `racket` libraries, but not `racket/base`.

```
(include path-spec)

path-spec = string
          | (file string)
          | (lib string ...+)
```

Inlines the syntax in the file designated by *path-spec* in place of the `include` expression.

A *path-spec* resembles a subset of the *mod-path* forms for `require`, but it specifies a file whose content need not be a module. That is, *string* refers to a file using a platform-independent relative path, `(file` *string*`)` refers to a file using platform-specific notation, and `(lib` *string* `...)` refers to a file within a collection.

If *path-spec* specifies a relative path, the path is resolved relative to the source for the `include` expression, if that source is a complete path string. If the source is not a complete path string, then *path-spec* is resolved relative to `(current-load-relative-directory)` if it is not `#f`, or relative to `(current-directory)` otherwise.

The included syntax is given the lexical context of the `include` expression, while the included syntax's source location refers to its actual source.

```
(include-at/relative-to context source path-spec)
```

Like `include`, except that the lexical context of *context* is used for the included syntax, and a relative *path-spec* is resolved with respect to the source of *source*. The *context* and *source* elements are otherwise discarded by expansion.

```
(include/reader path-spec reader-expr)
```

Like `include`, except that the procedure produced by the expression *reader-expr* is used to read the included file, instead of `read-syntax`.

The *reader-expr* is evaluated at expansion time in the transformer environment. Since it serves as a replacement for `read-syntax`, the expression's value should be a procedure that

consumes two inputs—a string representing the source and an input port—and produces a syntax object or `eof`. The procedure will be called repeatedly until it produces `eof`.

The syntax objects returned by the procedure should have source location information, but usually no lexical context; any lexical context in the syntax objects will be ignored.

```
(include-at/relative-to/reader context source path-spec reader-expr)
```

Combines `include-at/relative-to` and `include/reader`.

12.11 Syntax Utilities

```
(require racket/syntax)
```
package: base

The bindings documented in this section are provided by the `racket/syntax` library, not `racket/base` or `racket`.

12.11.1 Creating formatted identifiers

```
(format-id  lctx
            fmt
            v ...
           [#:source src
            #:props props
            #:cert ignored]) → identifier?
  lctx : (or/c syntax? #f)
  fmt : string?
  v : (or/c string? symbol? identifier? keyword? char? number?)
  src : (or/c syntax? #f) = #f
  props : (or/c syntax? #f) = #f
  ignored : (or/c syntax? #f) = #f
```

Like `format`, but produces an identifier using *lctx* for the lexical context, *src* for the source location, and *props* for the properties. An argument supplied with `#:cert` is ignored. (See `datum->syntax`.)

The format string must use only `~a` placeholders. Identifiers in the argument list are automatically converted to symbols.

Examples:

```
> (define-syntax (make-pred stx)
    (syntax-case stx ()
```

```
      [(make-pred name)
       (format-id #'name "~a?" (syntax-e #'name))]))

> (make-pred pair)
#<procedure:pair?>
> (make-pred none-such)
none-such?: undefined;
 cannot reference undefined identifier
> (define-syntax (better-make-pred stx)
    (syntax-case stx ()
      [(better-make-pred name)
       (format-id #'name #:source #'name
                  "~a?" (syntax-e #'name))]))

> (better-make-pred none-such)
none-such?: undefined;
 cannot reference undefined identifier
```

(Scribble doesn't show it, but the DrRacket pinpoints the location of the second error but not of the first.)

```
(format-symbol fmt v ...) → symbol?
  fmt : string?
  v : (or/c string? symbol? identifier? keyword? char? number?)
```

Like `format`, but produces a symbol. The format string must use only `~a` placeholders. Identifiers in the argument list are automatically converted to symbols.

Example:

```
> (format-symbol "make-~a" 'triple)
'make-triple
```

12.11.2 Pattern variables

```
(define/with-syntax pattern stx-expr)

  stx-expr : syntax?
```

Definition form of `with-syntax`. That is, it matches the syntax object result of `expr` against *`pattern`* and creates pattern variable definitions for the pattern variables of *`pattern`*.

Examples:

```
> (define/with-syntax (px ...) #'(a b c))
```

```
> (define/with-syntax (tmp ...) (generate-temporaries #'(px ...)))

> #'([tmp px] ...)
#<syntax:11:0 ((a9 a) (b10 b) (c11 c))>
> (define/with-syntax name #'Alice)

> #'(hello name)
#<syntax:13:0 (hello Alice)>
```

12.11.3 Error reporting

```
(current-syntax-context) → (or/c syntax? false/c)
(current-syntax-context stx) → void?
  stx : (or/c syntax? false/c)
```

The current contextual syntax object, defaulting to `#f`. It determines the special form name that prefixes syntax errors created by `wrong-syntax`.

```
(wrong-syntax stx format-string v ...) → any
  stx : syntax?
  format-string : string?
  v : any/c
```

Raises a syntax error using the result of `(current-syntax-context)` as the "major" syntax object and the provided *stx* as the specific syntax object. (The latter, *stx*, is usually the one highlighted by DrRacket.) The error message is constructed using the format string and arguments, and it is prefixed with the special form name as described under `current-syntax-context`.

Examples:

```
> (wrong-syntax #'here "expected ~s" 'there)
?: expected there
> (parameterize ([current-syntax-context #'(look over here)])
    (wrong-syntax #'here "expected ~s" 'there))
eval:15:0: look: expected there
  at: here
  in: (look over here)
```

A macro using `wrong-syntax` might set the syntax context at the very beginning of its transformation as follows:

```
(define-syntax (my-macro stx)
  (parameterize ([current-syntax-context stx])
    (syntax-case stx ()
      __)))
```

Then any calls to `wrong-syntax` during the macro's transformation will refer to `my-macro` (more precisely, the name that referred to `my-macro` where the macro was used, which may be different due to renaming, prefixing, etc).

12.11.4 Recording disappeared uses

```
(current-recorded-disappeared-uses)
 → (or/c (listof identifier?) false/c)
(current-recorded-disappeared-uses ids) → void?
  ids : (or/c (listof identifier?) false/c)
```

Parameter for tracking disappeared uses. Tracking is "enabled" when the parameter has a non-false value. This is done automatically by forms like `with-disappeared-uses`.

```
(with-disappeared-uses stx-expr)

  stx-expr : syntax?
```

Evaluates the *stx-expr*, catching identifiers looked up using `syntax-local-value/record`. Adds the caught identifiers to the `'disappeared-uses` syntax property of the resulting syntax object.

```
(syntax-local-value/record id predicate) → any/c
  id : identifier?
  predicate : (-> any/c boolean?)
```

Looks up *id* in the syntactic environment (as `syntax-local-value`). If the lookup succeeds and returns a value satisfying the predicate, the value is returned and *id* is recorded as a disappeared use by calling `record-disappeared-uses`. If the lookup fails or if the value does not satisfy the predicate, `#f` is returned and the identifier is not recorded as a disappeared use.

```
(record-disappeared-uses ids) → void?
  ids : (listof identifier?)
```

Add *ids* to `(current-recorded-disappeared-uses)` after calling `syntax-local-introduce` on each of the identifiers.

If not used within the extent of a `with-disappeared-uses` form or similar, has no effect.

12.11.5 Miscellaneous utilities

```
(generate-temporary [name-base]) → identifier?
  name-base : any/c = 'g
```

Generates one fresh identifier. Singular form of `generate-temporaries`. If *name-base* is supplied, it is used as the basis for the identifier's name.

```
(internal-definition-context-apply intdef-ctx
                                   stx)       → syntax?
  intdef-ctx : internal-definition-context?
  stx : syntax?
```

Applies the renamings of *intdef-ctx* to *stx*.

```
(syntax-local-eval stx [intdef-ctx]) → any
  stx : syntax?
  intdef-ctx : (or/c internal-definition-context? #f) = #f
```

Evaluates *stx* as an expression in the current transformer environment (that is, at phase level 1), optionally extended with *intdef-ctx*.

Examples:

```
> (define-syntax (show-me stx)
    (syntax-case stx ()
      [(show-me expr)
       (begin
         (printf "at compile time produces ~s\n"
                 (syntax-local-eval #'expr))
         #'(printf "at run time produces ~s\n"
                   expr))]))

> (show-me (+ 2 5))
at compile time produces 7
at run time produces 7

> (define-for-syntax fruit 'apple)

> (define fruit 'pear)

> (show-me fruit)
at compile time produces apple
at run time produces pear
```

```
(with-syntax* ([pattern stx-expr] ...)
  body ...+)

  stx-expr : syntax?
```

Similar to `with-syntax`, but the pattern variables of each *pattern* are bound in the *stx-expr*s of subsequent clauses as well as the *body*s, and the *pattern*s need not bind distinct pattern variables; later bindings shadow earlier bindings.

Example:

```
> (with-syntax* ([(x y) (list #'val1 #'val2)]
                 [nest #'((x) (y))])
    #'nest)
#<syntax:21:0 ((val1) (val2))>
```

13 Input and Output

13.1 Ports

§8 “Input and Output” in ***The Racket Guide*** introduces Ports and I/O.

Ports produce and/or consume bytes. An *input port* produces bytes, while an *output port* consumes bytes (and some ports are both input ports and output ports). When an input port is provided to a character-based operation, the bytes are decoded to a character, and character-based output operations similarly encode the character to bytes; see §13.1.1 “Encodings and Locales”. In addition to bytes and characters encoded as bytes, some ports can produce and/or consume arbitrary values as *special* results.

When a port corresponds to a file, network connection, or some other system resource, it must be explicitly closed via `close-input-port` or `close-output-port` (or indirectly via `custodian-shutdown-all`) to release low-level resources associated with the port. For any kind of port, after it is closed, attempting to read from or write to the port raises `exn:fail`.

Data produced by a input port can be read or *peeked*. When data is read, it is considered consumed and removed from the port’s stream. When data is peeked, it remains in the port’s stream to be returned again by the next read or peek. Previously peeked data can be *committed*, which causes the data to be removed from the port as for a read in a way that can be synchronized with other attempts to peek or read through a synchronizable event. Both read and peek operations are normally blocking, in the sense that the read or peek operation does not complete until data is available from the port; non-blocking variants of read and peek operations are also available.

The global variable `eof` is bound to the end-of-file value, and `eof-object?` returns `#t` only when applied to this value. Reading from a port produces an end-of-file result when the port has no more data, but some ports may also return end-of-file mid-stream. For example, a port connected to a Unix terminal returns an end-of-file when the user types control-D; if the user provides more input, the port returns additional bytes after the end-of-file.

Every port has a name, as reported by `object-name`. The name can be any value, and it is used mostly for error-reporting purposes. The `read-syntax` procedure uses the name of an input port as the default source location for the syntax objects that it produces.

A port can be used as a synchronizable event. An input port is ready for synchronization when `read-byte` would not block, and an output port is ready for synchronization when `write-bytes-avail` would not block or when the port contains buffered characters and `write-bytes-avail*` can flush part of the buffer (although `write-bytes-avail` might block). A value that can act as both an input port and an output port acts as an input port for a synchronizable event. The synchronization result of a port is the port itself.

13.1.1 Encodings and Locales

When a port is provided to a character-based operation, such as `read-char` or `read`, the port's bytes are read and interpreted as a UTF-8 encoding of characters. Thus, reading a single character may require reading multiple bytes, and a procedure like `char-ready?` may need to peek several bytes into the stream to determine whether a character is available. In the case of a byte stream that does not correspond to a valid UTF-8 encoding, functions such as `read-char` may need to peek one byte ahead in the stream to discover that the stream is not a valid encoding.

When an input port produces a sequence of bytes that is not a valid UTF-8 encoding in a character-reading context, then bytes that constitute an invalid sequence are converted to the character `#\uFFFD`. Specifically, bytes 255 and 254 are always converted to `#\uFFFD`, bytes in the range 192 to 253 produce `#\uFFFD` when they are not followed by bytes that form a valid UTF-8 encoding, and bytes in the range 128 to 191 are converted to `#\uFFFD` when they are not part of a valid encoding that was started by a preceding byte in the range 192 to 253. To put it another way, when reading a sequence of bytes as characters, a minimal set of bytes are changed to the encoding of `#\uFFFD` so that the entire sequence of bytes is a valid UTF-8 encoding.

See §4.4 "Byte Strings" for procedures that facilitate conversions using UTF-8 or other encodings. See also `reencode-input-port` and `reencode-output-port` for obtaining a UTF-8-based port from one that uses a different encoding of characters.

A *locale* captures information about a user's language-specific interpretation of character sequences. In particular, a locale determines how strings are "alphabetized," how a lowercase character is converted to an uppercase character, and how strings are compared without regard to case. String operations such as `string-ci=?` are *not* sensitive to the current locale, but operations such as `string-locale-ci=?` (see §4.3 "Strings") produce results consistent with the current locale.

A locale also designates a particular encoding of code-point sequences into byte sequences. Racket generally ignores this aspect of the locale, with a few notable exceptions: command-line arguments passed to Racket as byte strings are converted to character strings using the locale's encoding; command-line strings passed as byte strings to other processes (through `subprocess`) are converted to byte strings using the locale's encoding; environment variables are converted to and from strings using the locale's encoding; filesystem paths are converted to and from strings (for display purposes) using the locale's encoding; and, finally, Racket provides functions such as `string->bytes/locale` to specifically invoke a locale-specific encoding.

A Unix user selects a locale by setting environment variables, such as `LC_ALL`. On Windows and Mac OS X, the operating system provides other mechanisms for setting the locale. Within Racket, the current locale can be changed by setting the `current-locale` parameter. The locale name within Racket is a string, and the available locale names depend on the platform and its configuration, but the `""` locale means the current user's default locale;

on Windows and Mac OS X, the encoding for "" is always UTF-8, and locale-sensitive operations use the operating system's native interface. (In particular, setting the LC_ALL and LC_CTYPE environment variables does not affect the locale "" on Mac OS X. Use getenv and current-locale to explicitly install the environment-specified locale, if desired.) Setting the current locale to #f makes locale-sensitive operations locale-insensitive, which means using the Unicode mapping for case operations and using UTF-8 for encoding.

```
(current-locale) → (or/c string? #f)
(current-locale locale) → void?
  locale : (or/c string? #f)
```

A parameter that determines the current locale for procedures such as string-locale-ci=?.

When locale sensitivity is disabled by setting the parameter to #f, strings are compared, etc., in a fully portable manner, which is the same as the standard procedures. Otherwise, strings are interpreted according to a locale setting (in the sense of the C library's setlocale). The "" locale is always an alias for the current machine's default locale, and it is the default. The "C" locale is also always available; setting the locale to "C" is the same as disabling locale sensitivity with #f only when string operations are restricted to the first 128 characters. Other locale names are platform-specific.

String or character printing with write is not affected by the parameter, and neither are symbol case or regular expressions (see §4.7 "Regular Expressions").

13.1.2 Managing Ports

```
(input-port? v) → boolean?
  v : any/c
```

Returns #t if v is an input port, #f otherwise.

```
(output-port? v) → boolean?
  v : any/c
```

Returns #t if v is an output port, #f otherwise.

```
(port? v) → boolean?
  v : any/c
```

Returns #t if either (input-port? v) or (output-port? v) is #t, #f otherwise.

```
(close-input-port in) → void?
  in : input-port?
```

Closes the input port *in*. For some kinds of ports, closing the port releases lower-level resources, such as a file handle. If the port is already closed, close-input-port has no effect.

```
(close-output-port out) → void?
  out : output-port?
```

Closes the output port *out*. For some kinds of ports, closing the port releases lower-level resources, such as a file handle. Also, if the port is buffered, closing may first flush the port before closing it, and this flushing process can block. If the port is already closed, close-output-port has no effect.

```
(port-closed? port) → boolean?
  port : port?
```

Returns #t if the input or output port *port* is closed, #f otherwise.

```
(port-closed-evt port) → evt?
  port : port?
```

Return a synchronizable event that becomes ready for synchronization when *port* is closed. The synchronization result of a port-closed event is the port-closed event itself.

```
(current-input-port) → input-port?
(current-input-port in) → void?
  in : input-port?
```

A parameter that determines a default input port for many operations, such as read.

```
(current-output-port) → output-port?
(current-output-port out) → void?
  out : output-port?
```

A parameter that determines a default output port for many operations, such as write.

```
(current-error-port) → output-port?
(current-error-port out) → void?
  out : output-port?
```

A parameter that determines an output port that is typically used for errors and logging. For example, the default error display handler writes to this port.

```
(file-stream-port? port) → boolean?
  port : port?
```

Returns `#t` if the given port is a file-stream port (see §13.1.5 “File Ports”), `#f` otherwise.

```
(terminal-port? port) → boolean?
  port : port?
```

Returns `#t` if the given port is attached to an interactive terminal, `#f` otherwise.

```
eof : eof-object?
```

A value (distinct from all other values) that represents an end-of-file.

```
(eof-object? a) → boolean?
  a : any/c
```

Returns `#t` if `v` is `eof`, `#f` otherwise.

13.1.3 Port Buffers and Positions

Some ports—especially those that read from and write to files—are internally buffered:

- An input port is typically block-buffered by default, which means that on any read, the buffer is filled with immediately-available bytes to speed up future reads. Thus, if a file is modified between a pair of reads to the file, the second read can produce stale data. Calling `file-position` to set an input port’s file position flushes its buffer.
- An output port is typically block-buffered by default, though a terminal output port is line-buffered, and the initial error output port is unbuffered. An output buffer is filled with a sequence of written bytes to be committed as a group, either when the buffer is full (in block mode), when a newline is written (in line mode), when the port is closed via `close-output-port`, or when a flush is explicitly requested via a procedure like `flush-output`.

If a port supports buffering, its buffer mode can be changed via `file-stream-buffer-mode` (even if the port is not a file-stream port).

For an input port, peeking always places peeked bytes into the port’s buffer, even when the port’s buffer mode is `'none`; furthermore, on some platforms, testing the port for input (via `char-ready?` or `sync`) may be implemented with a peek. If an input port’s buffer mode is `'none`, then at most one byte is read for `read-bytes-avail!*`, `read-bytes-avail!`, `peek-bytes-avail!*`, or `peek-bytes-avail!`; if any bytes are buffered in the port (e.g., to satisfy a previous peek), the procedures may access multiple buffered bytes, but no further bytes are read.

In addition, the initial current output and error ports are automatically flushed when `read`, `read-line`, `read-bytes`, `read-string`, etc., are performed on the initial standard input port; more precisely, flushing is performed by the default port read handler (see `port-read-handler`).

```
(flush-output [out]) → void?
  out : output-port? = (current-output-port)
```

Forces all buffered data in the given output port to be physically written. Only file-stream ports, TCP ports, and custom ports (see §13.1.9 "Custom Ports") use buffers; when called on a port without a buffer, `flush-output` has no effect.

```
(file-stream-buffer-mode port) → (or/c 'none 'line 'block #f)
  port : port?
(file-stream-buffer-mode port mode) → void?
  port : port?
  mode : (or/c 'none 'line 'block)
```

Gets or sets the buffer mode for *port*, if possible. File-stream ports support setting the buffer mode, TCP ports (see §15.3 "Networking") support setting and getting the buffer mode, and custom ports (see §13.1.9 "Custom Ports") may support getting and setting buffer modes.

If *mode* is provided, it must be one of `'none`, `'line` (output only), or `'block`, and the port's buffering is set accordingly. If the port does not support setting the mode, the `exn:fail` exception is raised.

If *mode* is not provided, the current mode is returned, or `#f` is returned if the mode cannot be determined. If *port* is an input port and *mode* is `'line`, the `exn:fail:contract` exception is raised.

```
(file-position port) → exact-nonnegative-integer?
  port : port?
(file-position port pos) → void?
  port : port?
  pos : (or/c exact-nonnegative-integer? eof-object?)
```

Returns or sets the current read/write position of *port*.

Calling `file-position` without a position on a port other than a file-stream port or string port returns the number of bytes that have been read from that port if the position is known (see §13.1.4 "Counting Positions, Lines, and Columns"), otherwise the `exn:fail:filesystem` exception is raised.

For file-stream ports and string ports, the position-setting variant sets the read/write position to *pos* relative to the beginning of the file or (byte) string if *pos* is a number, or to the current end of the file or (byte) string if *pos* is `eof`. In position-setting mode, `file-position` raises

the exn:fail:contract exception for port kinds other than file-stream ports and string ports. Furthermore, not all file-stream ports support setting the position; if file-position is called with a position argument on such a file-stream port, the exn:fail:filesystem exception is raised.

When file-position sets the position *pos* beyond the current size of an output file or (byte) string, the file/string is enlarged to size *pos* and the new region is filled with 0 bytes. If *pos* is beyond the end of an input file or (byte) string, then reading thereafter returns eof without changing the port's position.

When changing the file position for an output port, the port is first flushed if its buffer is not empty. Similarly, setting the position for an input port clears the port's buffer (even if the new position is the same as the old position). However, although input and output ports produced by open-input-output-file share the file position, setting the position via one port does not flush the other port's buffer.

```
(file-position* port) → (or/c exact-nonnegative-integer? #f)
  port : port?
```

Like file-position on a single argument, but returns #f if the position is not known.

```
(file-truncate port size) → void?
  port : (and/c output-port? file-stream-port?)
  size : exact-nonnegative-integer?
```

Sets the size of the file written by *port* to *size*, assuming that the port is associated to a file whose size can be set.

The new file size can be either larger or smaller than its current size, but "truncate" in this function's name reflects that it is normally used to decrease the size of a file, since writing to a file or using file-position can extend a file's size.

13.1.4 Counting Positions, Lines, and Columns

By default, Racket keeps track of the *position* in a port as the number of bytes that have been read from or written to any port (independent of the read/write position, which is accessed or changed with file-position). Optionally, however, Racket can track the position in terms of characters (after UTF-8 decoding), instead of bytes, and it can track *line locations* and *column locations*; this optional tracking must be specifically enabled for a port via port-count-lines! or the port-count-lines-enabled parameter. Position, line, and column locations for a port are used by read-syntax. Position and line locations are numbered from 1; column locations are numbered from 0.

When counting lines, Racket treats linefeed, return, and return-linefeed combinations as a line terminator and as a single position (on all platforms). Each tab advances the column

count to one before the next multiple of 8. When a sequence of bytes in the range 128 to 253 forms a UTF-8 encoding of a character, the position/column is incremented once for each byte, and then decremented appropriately when a complete encoding sequence is discovered. See also §13.1 “Ports” for more information on UTF-8 decoding for ports.

A position is known for any port as long as its value can be expressed as a fixnum (which is more than enough tracking for realistic applications in, say, syntax-error reporting). If the position for a port exceeds the value of the largest fixnum, then the position for the port becomes unknown, and line and column tacking is disabled. Return-linefeed combinations are treated as a single character position only when line and column counting is enabled.

Custom ports can define their own counting functions, which are not subject to the rules above, except that the counting functions are invoked only when tracking is specifically enabled with `port-count-lines!`.

```
(port-count-lines! port) → void?
  port : port?
```

Turns on line location and column location counting for a port. Counting can be turned on at any time, though generally it is turned on before any data is read from or written to a port. At the point that line counting is turned on, `port-next-location` typically starts reporting as its last result (one more than) the number of characters read since line counting was enabled, instead of (one more than) bytes read since the port was opened.

When a port is created, if the value of the `port-count-lines-enabled` parameter is true, then line counting is automatically enabled for the port. Line counting cannot be disabled for a port after it is enabled.

```
(port-counts-lines? port) → boolean?
  port : port?
```

Returns `#t` if line location and column location counting has been enabled for *port*, `#f` otherwise.

```
(port-next-location port)
    (or/c exact-positive-integer? #f)
 → (or/c exact-nonnegative-integer? #f)
    (or/c exact-positive-integer? #f)
  port : port?
```

Returns three values: an integer or `#f` for the line number of the next read/written item, an integer or `#f` for the next item’s column, and an integer or `#f` for the next item’s position. The next column and position normally increase as bytes are read from or written to the port, but if line/character counting is enabled for *port*, the column and position results can decrease after reading or writing a byte that ends a UTF-8 encoding sequence.

If line counting is not enabled for a port, than the first two results are `#f`, and the last result is one more than the number of bytes read so far. At the point when line counting is enabled, the first two results typically become non-`#f`, and last result starts reporting characters instead of bytes, typically starting from the point when line counting is enabled.

Even with line counting enabled, a port may return `#f` values if it somehow cannot keep track of lines, columns, or positions.

```
(set-port-next-location! port
                         line
                         column
                         position) → void?
  port : port?
  line : (or/c exact-positive-integer? #f)
  column : (or/c exact-nonnegative-integer? #f)
  position : (or/c exact-positive-integer? #f)
```

Sets the next line, column, and position for *port*. If line counting has not been enabled for *port* or if *port* is a custom port that defines its own counting function, then `set-port-next-location!` has no effect.

```
(port-count-lines-enabled) → boolean?
(port-count-lines-enabled on?) → void?
  on? : any/c
```

A parameter that determines whether line counting is enabled automatically for newly created ports. The default value is `#f`.

13.1.5 File Ports

A port created by `open-input-file`, `open-output-file`, `subprocess`, and related functions is a *file-stream port*. The initial input, output, and error ports in `racket` are also file-stream ports. The `file-stream-port?` predicate recognizes file-stream ports.

When an input or output file-stream port is created, it is placed into the management of the current custodian (see §14.7 “Custodians”). In the case of an output port, a flush callback is registered with the current plumber to flush the port.

```
(open-input-file  path
                  [#:mode mode-flag
                   #:for-module? for-module?]) → input-port?
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
  for-module? : any/c = #f
```

Opens the file specified by *path* for input. The *mode-flag* argument specifies how the file's bytes are translated on input:

- `'binary` — bytes are returned from the port exactly as they are read from the file.
- `'text` — return and linefeed bytes (10 and 13) as read from the file are filtered by the port in a platform specific manner:
  - Unix and Mac OS X: no filtering occurs.
  - Windows: a return-linefeed combination from a file is returned by the port as a single linefeed; no filtering occurs for return bytes that are not followed by a linefeed, or for a linefeed that is not preceded by a return.

On Windows, `'text` mode works only with regular files; attempting to use `'text` with other kinds of files triggers an `exn:fail:filesystem` exception.

Otherwise, the file specified by *path* need not be a regular file. It might be a device that is connected through the filesystem, such as `"aux"` on Windows or `"/dev/null"` on Unix. In all cases, the port is buffered by default.

The port produced by `open-input-file` should be explicitly closed, either though `close-input-port` or indirectly via `custodian-shutdown-all`, to release the OS-level file handle. The input port will not be closed automatically if it is otherwise available for garbage collection (see §1.1.7 "Garbage Collection"); a will could be associated with an input port to close it more automatically (see §16.3 "Wills and Executors").

A path value that is the cleansed version of *path* is used as the name of the opened port.

If opening the file fails, if *for-module?* is true, and `current-module-path-for-load` has a non-`#f` value, then the raised exception is either `exn:fail:syntax:missing-module` (if the value of `current-module-path-for-load` is a syntax object) or `exn:fail:filesystem:missing-module` (otherwise).

Changed in version 6.0.1.6 of package `base`: Added `#:for-module?`.

Examples:

```
> (with-output-to-file some-file
    (lambda () (printf "hello world")))

(define in (open-input-file some-file))

> (read-string 11 in)
"hello world"
> (close-input-port in)
```

```
(open-output-file  path
                   [#:mode mode-flag
                    #:exists exists-flag]) → output-port?
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
                 (or/c 'error 'append 'update 'can-update
  exists-flag :        'replace 'truncate
                       'must-truncate 'truncate/replace)
              = 'error
```

Opens the file specified by *path* for output. The *mode-flag* argument specifies how bytes written to the port are translated when written to the file:

- `'binary` — bytes are written to the file exactly as written to the port.
- `'text` — on Windows, a linefeed byte (10) written to the port is translated to a return-linefeed combination in the file; no filtering occurs for returns.

On Windows, `'text` mode works only with regular files; attempting to use `'text` with other kinds of files triggers an `exn:fail:filesystem` exception.

The *exists-flag* argument specifies how to handle/require files that already exist:

- `'error` — raise `exn:fail:filesystem` if the file exists.
- `'replace` — remove the old file, if it exists, and write a new one.
- `'truncate` — remove all old data, if the file exists.
- `'must-truncate` — remove all old data in an existing file; if the file does not exist, the `exn:fail:filesystem` exception is raised.
- `'truncate/replace` — try `'truncate`; if it fails (perhaps due to file permissions), try `'replace`.
- `'update` — open an existing file without truncating it; if the file does not exist, the `exn:fail:filesystem` exception is raised. Use `file-position` to change the current read/write position.
- `'can-update` — open an existing file without truncating it, or create the file if it does not exist.
- `'append` — append to the end of the file, whether it already exists or not; on Windows, `'append` is equivalent to `'update`, except that the file is not required to exist, and the file position is immediately set to the end of the file after opening it.

The file specified by *path* need not be a regular file. It might be a device that is connected through the filesystem, such as `"aux"` on Windows or `"/dev/null"` on Unix. The output port is block-buffered by default, unless the file corresponds to a terminal, in which case it is line-buffered by default.

The port produced by `open-output-file` should be explicitly closed, either though `close-output-port` or indirectly via `custodian-shutdown-all`, to release the OS-level file handle. The output port will not be closed automatically if it is otherwise available for garbage collection (see §1.1.7 "Garbage Collection"); a will could be associated with an output port to close it more automatically (see §16.3 "Wills and Executors").

A path value that is the cleansed version of *path* is used as the name of the opened port.

Examples:

```
(define out (open-output-file some-file))

> (write "hello world" out)

> (close-output-port out)
```

```
(open-input-output-file path
                        [#:mode mode-flag
                         #:exists exists-flag])
 → input-port? output-port?
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
  exists-flag : (or/c 'error 'append 'update 'can-update
                      'replace 'truncate 'truncate/replace)
              = 'error
```

Like `open-output-file`, but producing two values: an input port and an output port. The two ports are connected in that they share the underlying file descriptor. This procedure is intended for use with special devices that can be opened by only one process, such as `"COM1"` in Windows. For regular files, sharing the file descriptor can be confusing. For example, using one port does not automatically flush the other port's buffer, and reading or writing in one port moves the file position (if any) for the other port. For regular files, use separate `open-input-file` and `open-output-file` calls to avoid confusion.

```
(call-with-input-file path
                      proc
                     [#:mode mode-flag]) → any
  path : path-string?
  proc : (input-port? . -> . any)
  mode-flag : (or/c 'binary 'text) = 'binary
```

Calls open-input-file with the *path* and *mode-flag* arguments, and passes the resulting port to *proc*. The result of *proc* is the result of the call-with-input-file call, but the newly opened port is closed when *proc* returns.

Examples:

```
> (with-output-to-file some-file
    (lambda () (printf "text in a file")))

> (call-with-input-file some-file
    (lambda (in) (read-string 14 in)))
"text in a file"
```

```
(call-with-output-file  path
                        proc
                       [#:mode mode-flag
                        #:exists exists-flag]) → any
  path : path-string?
  proc : (output-port? . -> . any)
  mode-flag : (or/c 'binary 'text) = 'binary
  exists-flag : (or/c 'error 'append 'update
                      'replace 'truncate 'truncate/replace)
              = 'error
```

Analogous to call-with-input-file, but passing *path*, *mode-flag* and *exists-flag* to open-output-file.

Examples:

```
> (call-with-output-file some-file
    (lambda (out)
      (write 'hello out)))

> (call-with-input-file some-file
    (lambda (in)
      (read-string 5 in)))
"hello"
```

```
(call-with-input-file*  path
                        proc
                       [#:mode mode-flag]) → any
  path : path-string?
  proc : (input-port? . -> . any)
  mode-flag : (or/c 'binary 'text) = 'binary
```

Like call-with-input-file, but the newly opened port is closed whenever control escapes the dynamic extent of the call-with-input-file* call, whether through *proc*'s return, a continuation application, or a prompt-based abort.

```
(call-with-output-file* path
                        proc
                       [#:mode mode-flag
                        #:exists exists-flag]) → any
  path : path-string?
  proc : (output-port? . -> . any)
  mode-flag : (or/c 'binary 'text) = 'binary
  exists-flag :
                (or/c 'error 'append 'update
                      'replace 'truncate 'truncate/replace)
              = 'error
```

Like `call-with-output-file`, but the newly opened port is closed whenever control escapes the dynamic extent of the `call-with-output-file*` call, whether through *proc*'s return, a continuation application, or a prompt-based abort.

```
(with-input-from-file path
                      thunk
                     [#:mode mode-flag]) → any
  path : path-string?
  thunk : (-> any)
  mode-flag : (or/c 'binary 'text) = 'binary
```

Like `call-with-input-file*`, but instead of passing the newly opened port to the given procedure argument, the port is installed as the current input port (see `current-input-port`) using `parameterize` around the call to *thunk*.

Examples:

```
> (with-output-to-file some-file
    (lambda () (printf "hello")))

> (with-input-from-file some-file
    (lambda () (read-string 5)))
"hello"
```

```
(with-output-to-file path
                     thunk
                    [#:mode mode-flag
                     #:exists exists-flag]) → any
  path : path-string?
  thunk : (-> any)
  mode-flag : (or/c 'binary 'text) = 'binary
  exists-flag :
                (or/c 'error 'append 'update
                      'replace 'truncate 'truncate/replace)
              = 'error
```

Like call-with-output-file*, but instead of passing the newly opened port to the given procedure argument, the port is installed as the current output port (see current-output-port) using parameterize around the call to *thunk*.

Examples:

```
> (with-output-to-file some-file
    (lambda () (printf "hello")))

> (with-input-from-file some-file
    (lambda () (read-string 5)))
"hello"
```

```
(port-try-file-lock? port mode) → boolean?
  port : file-stream-port?
  mode : (or/c 'shared 'exclusive)
```

Attempts to acquire a lock on the file using the current platform's facilities for file locking. Multiple processes can acquire a 'shared lock on a file, but at most one process can hold an 'exclusive lock, and 'shared and 'exclusive locks are mutually exclusive. When *mode* is 'shared, then *port* must be an input port; when *mode* is 'exclusive, then *port* must be an output port.

The result is #t if the requested lock is acquired, #f otherwise. When a lock is acquired, it is held until either it is released with port-file-unlock or the port is closed (perhaps because the process terminates).

Depending on the platform, locks may be merely advisory (i.e., locks affect only the ability of processes to acquire locks) or they may correspond to mandatory locks that prevent reads and writes to the locked file. Specifically, locks are mandatory on Windows and advisory on other platforms. Multiple tries for a 'shared lock on a single port can succeed; on Unix and Mac OS X, a single port-file-unlock release the lock, while on other Windows, a port-file-unlock is needed for each successful port-try-file-lock?. On Unix and Mac OS X, multiple tries for a 'exclusive lock can succeed and a single port-file-unlock releases the lock, while on Windows, a try for an 'exclusive lock fails for a given port if the port already holds the lock.

A lock acquired for an input port from open-input-output-file can be released through port-file-unlock on the corresponding output port, and vice versa. If the output port from open-input-output-file holds an 'exclusive lock, the corresponding input port can still acquire a 'shared lock, even multiple times; on Windows, a port-file-unlock is needed for each successful lock try, while a single port-file-unlock balances the lock tries on Unix and Mac OS X. A 'shared lock on an input port can be upgraded to an 'exclusive lock through the corresponding output port on Unix and Mac OS X, in which case a single port-file-unlock (on either port) releases the lock, while such upgrades are not allowed on Windows.

Locking is normally supported only for file ports, and attempting to acquire a lock with other kinds of file-stream ports raises an `exn:fail:filesystem` exception.

```
(port-file-unlock port) → void?
  port : file-stream-port?
```

Releases a lock held by the current process on the file of *port*.

```
(port-file-identity port) → exact-positive-integer?
  port : file-stream-port?
```

Returns a number that represents the identity of the device and file read or written by *port*. For two ports whose open times overlap, the result of `port-file-identity` is the same for both ports if and only if the ports access the same device and file. For ports whose open times do not overlap, no guarantee can be provided for the port identities (even if the ports actually access the same file)—except as can be inferred through relationships with other ports. If *port* is closed, the `exn:fail` exception is raised. On Windows 95, 98, and Me, if *port* is connected to a pipe instead of a file, the `exn:fail:filesystem` exception is raised.

Examples:

```
(define file1 (open-output-file some-file))

(define file2 (open-output-file some-other-file))

> (port-file-identity file1)
2214605061
> (port-file-identity file2)
2214990853
> (close-output-port file1)

> (close-output-port file2)
```

13.1.6 String Ports

A *string port* reads or writes from a byte string. An input string port can be created from either a byte string or a string; in the latter case, the string is effectively converted to a byte string using `string->bytes/utf-8`. An output string port collects output into a byte string, but `get-output-string` conveniently converts the accumulated bytes to a string.

Input and output string ports do not need to be explicitly closed. The `file-position` procedure works for string ports in position-setting mode.

§4.4 “Byte Strings” also provides information on bytestrings.

```
(string-port? p) → boolean?
  p : port?
```

Returns `#t` if *p* is a string port, `#f` otherwise.

Added in version 6.0.1.6 of package `base`.

```
(open-input-bytes bstr [name]) → (and/c input-port? string-port?)
  bstr : bytes?
  name : any/c = 'string
```

Creates an input string port that reads characters from *bstr* (see §4.4 “Byte Strings”). Modifying *bstr* afterward does not affect the byte stream produced by the port. The optional *name* argument is used as the name for the returned port.

Examples:

```
> (define sp (open-input-bytes #"(apples 42 day)"))

> (define sexp1 (read sp))

> (first sexp1)
'apples
> (rest sexp1)
'(42 day)
> (read-line (open-input-bytes
              #"the cow jumped over the moon\nthe little dog\n"))
"the cow jumped over the moon"
```

§4.3 “Strings” also provides information on strings.

```
(open-input-string str [name]) → (and/c input-port? string-port?)
  str : string?
  name : any/c = 'string
```

Creates an input string port that reads bytes from the UTF-8 encoding (see §13.1.1 “Encodings and Locales”) of *str*. The optional *name* argument is used as the name for the returned port.

Examples:

```
> (define sp (open-input-string "(λ (x) x)"))

> (read sp)
'(λ (x) x)
> (define names (open-input-string "Günter Harder\nFrédéric
Paulin\n"))
```

```
> (read-line names)
"Günter Harder"
> (read-line names)
"Frédéric Paulin"
```

```
(open-output-bytes [name]) → (and/c output-port? string-port?)
  name : any/c = 'string
```

Creates an output string port that accumulates the output into a byte string. The optional *name* argument is used as the name for the returned port.

Examples:

```
> (define op1 (open-output-bytes))

> (write '((1 2 3) ("Tom" "Dick") ('a 'b 'c)) op1)

> (get-output-bytes op1)
#"((1 2 3) (\"Tom\" \"Dick\") ((quote a) (quote b) (quote c)))"
> (define op2 (open-output-bytes))

> (write "Hi " op2)

> (write "there" op2)

> (get-output-bytes op2)
#"\"Hi \"\"there\""
> (define op3 (open-output-bytes))

> (write-bytes #"Hi " op3)
3
> (write-bytes #"there" op3)
5
> (get-output-bytes op3)
#"Hi there"
```

```
(open-output-string [name]) → (and/c output-port? string-port?)
  name : any/c = 'string
```

The same as `open-output-bytes`.

Examples:

```
> (define op1 (open-output-string))

> (write '((1 2 3) ("Tom" "Dick") ('a 'b 'c)) op1)
```

```
> (get-output-string op1)
"((1 2 3) (\"Tom\" \"Dick\") ((quote a) (quote b) (quote c)))"
> (define op2 (open-output-string))

> (write "Hi " op2)

> (write "there" op2)

> (get-output-string op2)
"\"Hi \"\"there\""
> (define op3 (open-output-string))

> (write-string "Hi " op3)
3
> (write-string "there" op3)
5
> (get-output-string op3)
"Hi there"
```

```
(get-output-bytes  out
                   [reset?
                    start-pos
                    end-pos]) → bytes?
  out : (and/c output-port? string-port?)
  reset? : any/c = #f
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = #f
```

Returns the bytes accumulated in the string port *out* so far in a freshly allocated byte string (including any bytes written after the port's current position, if any). The *out* port must be an output string port produced by `open-output-bytes` (or `open-output-string`) or a structure whose `prop:output-port` property refers to such an output port (transitively).

If *reset?* is true, then all bytes are removed from the port, and the port's position is reset to `0`; if *reset?* is `#f`, then all bytes remain in the port for further accumulation (so they are returned for later calls to `get-output-bytes` or `get-output-string`), and the port's position is unchanged.

The *start-pos* and *end-pos* arguments specify the range of bytes in the port to return; supplying *start-pos* and *end-pos* is the same as using `subbytes` on the result of `get-output-bytes`, but supplying them to `get-output-bytes` can avoid an allocation. The *end-pos* argument can be `#f`, which corresponds to not passing a second argument to `subbytes`.

Examples:

```
> (define op (open-output-bytes))

> (write '((1 2 3) ("Tom" "Dick") ('a 'b 'c)) op)

> (get-output-bytes op)
#"((1 2 3) (\"Tom\" \"Dick\") ((quote a) (quote b) (quote c)))"
> (get-output-bytes op #f 3 16)
#" 2 3) (\"Tom\" "
> (get-output-bytes op #t)
#"((1 2 3) (\"Tom\" \"Dick\") ((quote a) (quote b) (quote c)))"
> (get-output-bytes op)
#""
```

```
(get-output-string out) → string?
  out : (and/c output-port? string-port?)
```

Returns `(bytes->string/utf-8 (get-output-bytes out) #\?)`.

Examples:

```
> (define i (open-input-string "hello world"))

> (define o (open-output-string))

> (write (read i) o)

> (get-output-string o)
"hello"
```

13.1.7 Pipes

A Racket *pipe* is internal to Racket, and not related to OS-level pipes for communicating between different processes.

```
(make-pipe [limit input-name output-name])
 → input-port? output-port?
  limit : exact-positive-integer? = #f
  input-name : any/c = 'pipe
  output-name : any/c = 'pipe
```

Returns two port values: the first port is an input port and the second is an output port. Data written to the output port is read from the input port, with no intermediate buffering. Unlike some other kinds of ports, pipe ports do not need to be explicitly closed to be reclaimed by garbage collection.

OS-level pipes may be created by `subprocess`, opening an existing named file on a Unix filesystem, or starting Racket with pipes for its original input, output, or error port. Such pipes are file-stream ports, unlike the pipes produced by `make-pipe`.

If *limit* is #f, the new pipe holds an unlimited number of unread bytes (i.e., limited only by the available memory). If *limit* is a positive number, then the pipe will hold at most *limit* unread/unpeeked bytes; writing to the pipe's output port thereafter will block until a read or peek from the input port makes more space available. (Peeks effectively extend the port's capacity until the peeked bytes are read.)

The optional *input-name* and *output-name* are used as the names for the returned input and output ports, respectively.

```
(pipe-content-length pipe-port) → exact-nonnegative-integer?
  pipe-port : port?
```

Returns the number of bytes contained in a pipe, where *pipe-port* is either of the pipe's ports produced by make-pipe. The pipe's content length counts all bytes that have been written to the pipe and not yet read (though possibly peeked).

13.1.8 Structures as Ports

```
prop:input-port : struct-type-property?
```

```
prop:output-port : struct-type-property?
```

The prop:input-port and prop:output-port structure type properties identify structure types whose instances can serve as input and output ports, respectively.

Each property value can be either of the following:

- An input port (for prop:input-port) or output port (for prop:output-port): In this case, using the structure as port is equivalent to using the given input or output port.
- An exact, non-negative integer between 0 (inclusive) and the number of non-automatic fields in the structure type (exclusive, not counting supertype fields): The integer identifies a field in the structure, and the field must be designated as immutable. If the field contains an input port (for prop:input-port) or output port (for prop:output-port), the port is used. Otherwise, an empty string input port is used for prop:input-port, and a port that discards all data is used for prop:output-port.

Some procedures, such as file-position, work on both input and output ports. When given an instance of a structure type with both the prop:input-port and prop:output-port properties, the instance is used as an input port.

13.1.9 Custom Ports

The `make-input-port` and `make-output-port` procedures create *custom ports* with arbitrary control procedures (much like implementing a device driver). Custom ports are mainly useful to obtain fine control over the action of committing bytes as read or written.

```
(make-input-port  name
                  read-in
                  peek
                  close
                 [get-progress-evt
                  commit
                  get-location
                  count-lines!
                  init-position
                  buffer-mode])     → input-port?
  name : any/c
            (or/c
             (bytes?
              . -> . (or/c exact-nonnegative-integer?
  read-in :                 eof-object?
                            procedure?
                            evt?))
             input-port?)
         (or/c
          (bytes? exact-nonnegative-integer? (or/c evt? #f)
                  . -> . (or/c exact-nonnegative-integer?
                               eof-object?
  peek :                       procedure?
                               evt?
                               #f))
          input-port?)
  close : (-> any)
  get-progress-evt : (or/c (-> evt?) #f) = #f
           (or/c (exact-positive-integer? evt? evt? . -> . any)
  commit :       #f)
         = #f
                  (or/c
                   (->
                    (values (or/c exact-positive-integer? #f)
  get-location :            (or/c exact-nonnegative-integer? #f)
                            (or/c exact-positive-integer? #f)))
                   #f)
                = #f
  count-lines! : (-> any) = void
```

```
                (or/c exact-positive-integer?
                      port?
init-position :       #f
                      (-> (or/c exact-positive-integer? #f)))
              = 1
              (or/c (case-> ((or/c 'block 'none) . -> . any)
buffer-mode :                (-> (or/c 'block 'none #f)))
                    #f)
            = #f
```

Creates an input port, which is immediately open for reading. If *close* procedure has no side effects, then the port need not be explicitly closed. See also `make-input-port/peek-to-read`.

The arguments implement the port as follows:

- *name* — the name for the input port.
- *read-in* — either an input port, in which case reads are redirected to the given port, or a procedure that takes a single argument: a mutable byte string to receive read bytes. The procedure's result is one of the following:
  - the number of bytes read, as an exact, non-negative integer;
  - `eof`;
  - a procedure of arity four (representing a "special" result, as discussed further below) and optionally of arity zero, but a procedure result is allowed only when *peek* is not `#f`;
  - a pipe input port that supplies bytes to be used as long as the pipe has content (see `pipe-content-length`) or until *read-in* or *peek* is called again; or
  - a synchronizable event (see §11.2.1 "Events") other than a pipe input port or procedure of arity four; the event becomes ready when the read is complete (roughly): the event's value can be one of the above three results or another event like itself; in the last case, a reading process loops with `sync` until it gets a non-event result.

  The *read-in* procedure must not block indefinitely. If no bytes are immediately available for reading, the *read-in* must return `0` or an event, and preferably an event (to avoid busy waits). The *read-in* should not return `0` (or an event whose value is `0`) when data is available in the port, otherwise polling the port will behave incorrectly. An event result from an event can also break polling.

  If the result of a *read-in* call is not one of the above values, the `exn:fail:contract` exception is raised. If a returned integer is larger than the supplied byte string's length, the `exn:fail:contract` exception is raised. If *peek* is `#f` and a procedure for a special result is returned, the `exn:fail:contract` exception is raised.

The *read-in* procedure can report an error by raising an exception, but only if no bytes are read. Similarly, no bytes should be read if eof, an event, or a procedure is returned. In other words, no bytes should be lost due to spurious exceptions or non-byte data.

A port's reading procedure may be called in multiple threads simultaneously (if the port is accessible in multiple threads), and the port is responsible for its own internal synchronization. Note that improper implementation of such synchronization mechanisms might cause a non-blocking read procedure to block indefinitely.

If the result is a pipe input port, then previous *get-progress-evt* calls whose event is not yet ready must have been the pipe input port itself. Furthermore, *get-progress-evt* must continue to return the pipe as long as it contains data, or until the *read-in* or peek-in procedure is called again (instead of using the pipe, for whatever reason). If *read-in* or peek-in is called, any previously associated pipe (as returned by a previous call) is disassociated from the port and is not in use by any other thread as a result of the previous association.

If *peek*, *get-progress-evt*, and *commit* are all provided and non-#f, then the following is an acceptable implementation of *read-in*:

```
(lambda (bstr)
  (let* ([progress-evt (get-progress-evt)]
         [v (peek bstr 0 progress-evt)])
    (cond
     [(sync/timeout 0 progress-evt) 0] ; try again
     [(evt? v) (wrap-evt v (lambda (x) 0))] ; sync, try again
     [(and (number? v) (zero? v)) 0] ; try again
     [else
      (if (commit (if (number? v) v 1)
                  progress-evt
                  always-evt)
          v      ; got a result
          0)])))) ; try again
```

An implementor may choose not to implement the *peek*, *get-progress-evt*, and *commit* procedures, however, and even an implementor who does supply them may provide a different *read-in* that uses a fast path for non-blocking reads.

In an input port is provided for *read-in*, then an input port must also be provided for *peek*.

- *peek* — either #f, an input port (in which case peeks are redirected to the given port), or a procedure that takes three arguments:
  - a mutable byte string to receive peeked bytes;
  - a non-negative number of bytes (or specials) to skip before peeking; and
  - either #f or a progress event produced by *get-progress-evt*.

The results and conventions for *peek* are mostly the same as for *read-in*. The main difference is in the handling of the progress event, if it is not #f. If the given progress event becomes ready, the *peek* must abort any skip attempts and not peek any values. In particular, *peek* must not peek any values if the progress event is initially ready. If the port has been closed, the progress event should be ready, in which case *peek* should complete (instead of failing because the port is closed).

Unlike *read-in*, *peek* should produce #f (or an event whose value is #f) if no bytes were peeked because the progress event became ready. Like *read-in*, a 0 result indicates that another attempt is likely to succeed, so 0 is inappropriate when the progress event is ready. Also like *read-in*, *peek* must not block indefinitely.

The skip count provided to *peek* is a number of bytes (or specials) that must remain present in the port—in addition to the peek results—when the peek results are reported. If a progress event is supplied, then the peek is effectively canceled when another process reads data before the given number can be skipped. If a progress event is not supplied and data is read, then the peek must effectively restart with the original skip count.

The system does not check that multiple peeks return consistent results, or that peeking and reading produce consistent results.

If *peek* is #f, then peeking for the port is implemented automatically in terms of reads, but with several limitations. First, the automatic implementation is not thread-safe. Second, the automatic implementation cannot handle special results (non-byte and non-eof), so *read-in* cannot return a procedure for a special when *peek* is #f. Finally, the automatic peek implementation is incompatible with progress events, so if *peek* is #f, then progress-evt and *commit* must be #f. See also make-input-port/peek-to-read, which implements peeking in terms of *read-in* without these constraints.

In an input port is provided for *peek*, then an input port must also be provided for *read-in*.

- *close* — a procedure of zero arguments that is called to close the port. The port is not considered closed until the closing procedure returns. The port's procedures will never be used again via the port after it is closed. However, the closing procedure can be called simultaneously in multiple threads (if the port is accessible in multiple threads), and it may be called during a call to the other procedures in another thread; in the latter case, any outstanding reads and peeks should be terminated with an error.

- *get-progress-evt* — either #f (the default), or a procedure that takes no arguments and returns an event. The event must become ready only after data is next read from the port or the port is closed. If the port is already closed, the event must be ready. After the event becomes ready, it must remain so. See the description of *read-in* for information about the allowed results of this function when *read-in* returns a pipe input port. See also semaphore-peek-evt, which is sometimes useful for implementing *get-progress-evt*.

  If *get-progress-evt* is #f, then port-provides-progress-evts? applied to the

port will produce `#f`, and the port will not be a valid argument to `port-progress-evt`.

The result event will not be exposed directly by `port-progress-evt`. Instead, it will be wrapped in an event for which `progress-evt?` returns true.

- *`commit`* — either `#f` (the default), or a procedure that takes three arguments:
  - an exact, positive integer k_r;
  - a progress event produced by *`get-progress-evt`*;
  - an event, *`done`*, that is either a channel-put event, channel, semaphore, semaphore-peek event, always event, or never event.

  A *commit* corresponds to removing data from the stream that was previously peeked, but only if no other process removed data first. (The removed data does not need to be reported, because it has been peeked already.) More precisely, assuming that k_p bytes, specials, and mid-stream `eof`s have been previously peeked or skipped at the start of the port's stream, *`commit`* must satisfy the following constraints:
  - It must return only when the commit is complete or when the given progress event becomes ready.
  - It must commit only if k_p is positive.
  - If it commits, then it must do so with either k_r items or k_p items, whichever is smaller, and only if k_p is positive.
  - It must never choose *`done`* in a synchronization after the given progress event is ready, or after *`done`* has been synchronized once.
  - It must not treat any data as read from the port unless *`done`* is chosen in a synchronization.
  - It must not block indefinitely if *`done`* is ready; it must return soon after the read completes or soon after the given progress event is ready, whichever is first.
  - It can report an error by raising an exception, but only if no data has been committed. In other words, no data should be lost due to an exception, including a break exception.
  - It must return a true value if data has been committed, `#f` otherwise. When it returns a value, the given progress event must be ready (perhaps because data has just been committed).
  - It should return a byte string as a true result when line counting is enabled and *`get-location`* is `#f` (so that line counting is implemented the default way); the result byte string represents the data that was committed for the purposes of character and line counting. If any other true result is returned when a byte string is expected, it is treated like a byte string where each byte corresponds to a non-newline character.
  - It must raise an exception if no data (including `eof`) has been peeked from the beginning of the port's stream, or if it would have to block indefinitely to wait for the given progress event to become ready.

A call to *commit* is parameterize-breaked to disable breaks.

- *get-location* — either #f (the default), or a procedure that takes no arguments and returns three values: the line number for the next item in the port's stream (a positive number or #f), the column number for the next item in the port's stream (a non-negative number or #f), and the position for the next item in the port's stream (a positive number or #f). See also §13.1.4 "Counting Positions, Lines, and Columns".

  This procedure is called to implement port-next-location, but only if line counting is enabled for the port via port-count-lines! (in which case *count-lines!* is called). The read and read-syntax procedures assume that reading a non-whitespace character increments the column and position by one.

- *count-lines!* — a procedure of no arguments that is called if and when line counting is enabled for the port. The default procedure is void.

- *init-position* — normally an exact, positive integer that determines the position of the port's first item, which is used by file-position or when line counting is *not* enabled for the port. The default is 1. If *init-position* is #f, the port is treated as having an unknown position. If *init-position* is a port, then the given port's position is always used for the new port's position. If *init-position* is a procedure, it is called as needed to obtain the port's position.

- *buffer-mode* — either #f (the default) or a procedure that accepts zero or one arguments. If *buffer-mode* is #f, then the resulting port does not support a buffer-mode setting. Otherwise, the procedure is called with one symbol argument ('block or 'none) to set the buffer mode, and it is called with zero arguments to get the current buffer mode. In the latter case, the result must be 'block, 'none, or #f (unknown). See §13.1.3 "Port Buffers and Positions" for more information on buffer modes.

"Special" results: When *read-in* or *peek* (or an event produced by one of these) returns a procedure, the procedure is used to obtain a non-byte result. (This non-byte result is *not* intended to return a character or eof; in particular, read-char raises an exception if it encounters a special-result procedure, even if the procedure produces a byte.) A special-result procedure must accept four arguments, and it can optionally accept zero arguments:

- When the special read is triggered by read-syntax or read-syntax/recursive, the procedure is passed four arguments that represent a source location.

- When the special read is triggered by read, read-byte-or-special, read-char-or-special, peek-byte-or-special, or peek-char-or-special, the procedure is passed no arguments if it accepts zero arguments, otherwise it is passed four arguments that are all #f.

The special-value procedure can return an arbitrary value, and it will be called zero or one times (not necessarily before further reads or peeks from the port). See §13.7.2 "Reader-Extension Procedures" for more details on the procedure's result.

If *read-in* or *peek* returns a special procedure when called by any reading procedure other than read, read-syntax, read-char-or-special, peek-char-or-special, read-byte-or-special, or peek-byte-or-special, then the exn:fail:contract exception is raised.

Examples:

```
; A port with no input...
; Easy: (open-input-bytes #"")
; Hard:
> (define /dev/null-in
    (make-input-port 'null
                     (lambda (s) eof)
                     (lambda (skip s progress-evt) eof)
                     void
                     (lambda () never-evt)
                     (lambda (k progress-evt done-evt)
                       (error "no successful peeks!"))))

> (read-char /dev/null-in)
#<eof>
> (peek-char /dev/null-in)
#<eof>
> (read-byte-or-special /dev/null-in)
#<eof>
> (peek-byte-or-special /dev/null-in 100)
#<eof>
; A port that produces a stream of 1s:
> (define infinite-ones
    (make-input-port
     'ones
     (lambda (s)
       (bytes-set! s 0 (char->integer #\1)) 1)
     #f
     void))

> (read-string 5 infinite-ones)
"11111"
; But we can't peek ahead arbitrarily far, because the
; automatic peek must record the skipped bytes, so
; we'd run out of memory.
; An infinite stream of 1s with a specific peek procedure:
> (define infinite-ones
    (let ([one! (lambda (s)
                  (bytes-set! s 0 (char->integer #\1)) 1)])
      (make-input-port
```

```
       'ones
       one!
       (lambda (s skip progress-evt) (one! s))
       void)))

> (read-string 5 infinite-ones)
"11111"
; Now we can peek ahead arbitrarily far:
> (peek-string 5 (expt 2 5000) infinite-ones)
"11111"
; The port doesn't supply procedures to implement progress events:
> (port-provides-progress-evts? infinite-ones)
#f
> (port-progress-evt infinite-ones)
port-progress-evt: port does not provide progress evts
  port: #<input-port:ones>
; Non-byte port results:
> (define infinite-voids
    (make-input-port
     'voids
     (lambda (s) (lambda args 'void))
     (lambda (skip s evt) (lambda args 'void))
     void))

> (read-char infinite-voids)
read-char: non-character in an unsupported context
  port: #<input-port:voids>
> (read-char-or-special infinite-voids)
'void
; This port produces 0, 1, 2, 0, 1, 2, etc., but it is not
; thread-safe, because multiple threads might read and change n.
> (define mod3-cycle/one-thread
    (let* ([n 2]
           [mod! (lambda (s delta)
                   (bytes-set! s 0 (+ 48 (modulo (+ n delta) 3)))
                   1)])
      (make-input-port
       'mod3-cycle/not-thread-safe
       (lambda (s)
         (set! n (modulo (add1 n) 3))
         (mod! s 0))
       (lambda (s skip evt)
         (mod! s skip))
       void)))

> (read-string 5 mod3-cycle/one-thread)
```

```
"01201"
> (peek-string 5 (expt 2 5000) mod3-cycle/one-thread)
"20120"
; Same thing, but thread-safe and kill-safe, and with progress
; events. Only the server thread touches the stateful part
; directly. (See the output port examples for a simpler thread-
safe
; example, but this one is more general.)
> (define (make-mod3-cycle)
    (define read-req-ch (make-channel))
    (define peek-req-ch (make-channel))
    (define progress-req-ch (make-channel))
    (define commit-req-ch (make-channel))
    (define close-req-ch (make-channel))
    (define closed? #f)
    (define n 0)
    (define progress-sema #f)
    (define (mod! s delta)
      (bytes-set! s 0 (+ 48 (modulo (+ n delta) 3)))
      1)
    ; ----------------------------
    ; The server has a list of outstanding commit requests,
    ;  and it also must service each port operation (read,
    ;  progress-evt, etc.)
    (define (serve commit-reqs response-evts)
      (apply
       sync
       (handle-evt read-req-ch
                   (handle-read commit-reqs response-evts))
       (handle-evt progress-req-ch
                   (handle-progress commit-reqs response-evts))
       (handle-evt commit-req-ch
                   (add-commit commit-reqs response-evts))
       (handle-evt close-req-ch
                   (handle-close commit-reqs response-evts))
       (append
        (map (make-handle-response commit-reqs response-evts)
             response-evts)
        (map (make-handle-commit commit-reqs response-evts)
             commit-reqs))))
    ; Read/peek request: fill in the string and commit
    (define ((handle-read commit-reqs response-evts) r)
      (let ([s (car r)]
            [skip (cadr r)]
            [ch (caddr r)]
            [nack (cadddr r)]
```

```
        [evt (car (cddddr r))]
        [peek? (cdr (cddddr r))])
    (let ([fail? (and evt
                      (sync/timeout 0 evt))])
      (unless (or closed? fail?)
        (mod! s skip)
        (unless peek?
          (commit! 1)))
      ; Add an event to respond:
      (serve commit-reqs
             (cons (choice-evt
                    nack
                    (channel-put-evt ch (if closed?
                                            0
                                            (if fail? #f 1))))
                   response-evts)))))
; Progress request: send a peek evt for the current
;  progress-sema
(define ((handle-progress commit-reqs response-evts) r)
  (let ([ch (car r)]
        [nack (cdr r)])
    (unless progress-sema
      (set! progress-sema (make-semaphore (if closed? 1 0))))
    ; Add an event to respond:
    (serve commit-reqs
           (cons (choice-evt
                  nack
                  (channel-put-evt
                   ch
                   (semaphore-peek-evt progress-sema)))
                 response-evts))))
; Commit request: add the request to the list
(define ((add-commit commit-reqs response-evts) r)
  (serve (cons r commit-reqs) response-evts))
; Commit handling: watch out for progress, in which case
;  the response is a commit failure; otherwise, try
;  to sync for a commit. In either event, remove the
;  request from the list
(define ((make-handle-commit commit-reqs response-evts) r)
  (let ([k (car r)]
        [progress-evt (cadr r)]
        [done-evt (caddr r)]
        [ch (cadddr r)]
        [nack (cddddr r)])
    ; Note: we don't check that k is <= the sum of
    ;  previous peeks, because the entire stream is actually
```

```
        ;  known, but we could send an exception in that case.
        (choice-evt
         (handle-evt progress-evt
                     (lambda (x)
                       (sync nack (channel-put-evt ch #f))
                       (serve (remq r commit-reqs) response-
evts)))
         ; Only create an event to satisfy done-evt if progress-
evt
         ;  isn't already ready.
         ; Afterward, if progress-evt becomes ready, then this
         ;  event-making function will be called again, because
         ;  the server controls all posts to progress-evt.
         (if (sync/timeout 0 progress-evt)
             never-evt
             (handle-evt done-evt
                         (lambda (v)
                           (commit! k)
                           (sync nack (channel-put-evt ch #t))
                           (serve (remq r commit-reqs)
                                  response-evts)))))))
    ; Response handling: as soon as the respondee listens,
    ;  remove the response
    (define ((make-handle-response commit-reqs response-evts) evt)
      (handle-evt evt
                  (lambda (x)
                    (serve commit-reqs
                           (remq evt response-evts)))))
    ; Close handling: post the progress sema, if any, and set
    ;   the closed? flag
    (define ((handle-close commit-reqs response-evts) r)
      (let ([ch (car r)]
            [nack (cdr r)])
        (set! closed? #t)
        (when progress-sema
          (semaphore-post progress-sema))
        (serve commit-reqs
               (cons (choice-evt nack
                                 (channel-put-evt ch (void)))
                     response-evts))))
    ; Helper for reads and post-peek commits:
    (define (commit! k)
      (when progress-sema
        (semaphore-post progress-sema)
        (set! progress-sema #f))
      (set! n (+ n k)))
```

```
    ; Start the server thread:
    (define server-thread (thread (lambda () (serve null null))))
    ; ---------------------------
    ; Client-side helpers:
    (define (req-evt f)
      (nack-guard-evt
       (lambda (nack)
         ; Be sure that the server thread is running:
         (thread-resume server-thread (current-thread))
         ; Create a channel to hold the reply:
         (let ([ch (make-channel)])
           (f ch nack)
           ch))))
    (define (read-or-peek-evt s skip evt peek?)
      (req-evt (lambda (ch nack)
                 (channel-put read-req-ch
                              (list* s skip ch nack evt peek?)))))
    ; Make the port:
    (make-input-port 'mod3-cycle
                     ; Each handler for the port just sends
                     ;  a request to the server
                     (lambda (s) (read-or-peek-evt s 0 #f #f))
                     (lambda (s skip evt)
                       (read-or-peek-evt s skip evt #t))
                     (lambda () ; close
                       (sync (req-evt
                              (lambda (ch nack)
                                (channel-put progress-req-ch
                                             (list* ch nack))))))
                     (lambda () ; progress-evt
                       (sync (req-evt
                              (lambda (ch nack)
                                (channel-put progress-req-ch
                                             (list* ch nack))))))
                     (lambda (k progress-evt done-evt)  ; commit
                       (sync (req-evt
                              (lambda (ch nack)
                                (channel-put
                                 commit-req-ch
                                 (list* k progress-evt done-evt ch
                                        nack))))))))

> (define mod3-cycle (make-mod3-cycle))

> (let ([result1 #f]
        [result2 #f])
```

```
    (let ([t1 (thread
               (lambda ()
                 (set! result1 (read-string 5 mod3-cycle))))]
          [t2 (thread
               (lambda ()
                 (set! result2 (read-string 5 mod3-cycle))))])
      (thread-wait t1)
      (thread-wait t2)
      (string-append result1 "," result2)))
"12012,01200"
> (define s (make-bytes 1))

> (define progress-evt (port-progress-evt mod3-cycle))

> (peek-bytes-avail! s 0 progress-evt mod3-cycle)
1
> s
#"1"
> (port-commit-peeked 1 progress-evt (make-semaphore 1)
                      mod3-cycle)
#t
> (sync/timeout 0 progress-evt)
#<progress-evt>
> (peek-bytes-avail! s 0 progress-evt mod3-cycle)
0
> (port-commit-peeked 1 progress-evt (make-semaphore 1)
                      mod3-cycle)
#f
> (close-input-port mod3-cycle)
```

```
(make-output-port  name
                   evt
                   write-out
                   close
                  [write-out-special
                   get-write-evt
                   get-write-special-evt
                   get-location
                   count-lines!
                   init-position
                   buffer-mode])          → output-port?
  name : any/c
  evt : evt?
```

```
            (or/c
             (bytes? exact-nonnegative-integer?
                     exact-nonnegative-integer?
                     boolean?
                     boolean?
write-out :          . -> .
                     (or/c exact-nonnegative-integer?
                           #f
                           evt?))
             output-port?)
close : (-> any)
                    (or/c (any/c boolean? boolean?
                                 . -> .
                                 (or/c any/c
write-out-special :                    #f          = #f
                                       evt?))
                          output-port?
                          #f)
                (or/c
                 (bytes? exact-nonnegative-integer?
                         exact-nonnegative-integer?
get-write-evt :          . -> .                      = #f
                         evt?)
                 #f)
                        (or/c
get-write-special-evt :  (any/c . -> . evt?) = #f
                         #f)
               (or/c
                (->
                 (values (or/c exact-positive-integer? #f)
get-location :           (or/c exact-nonnegative-integer? #f)
                         (or/c exact-positive-integer? #f)))
                #f)
              = #f
count-lines! : (-> any) = void
               (or/c exact-positive-integer?
                     port?
init-position :      #f
                     (-> (or/c exact-positive-integer? #f)))
              = 1
              (or/c (case->
                     ((or/c 'block 'line 'none) . -> . any)
buffer-mode :        (-> (or/c 'block 'line 'none #f)))
                    #f)
            = #f
```

Creates an output port, which is immediately open for writing. If *close* procedure has no side effects, then the port need not be explicitly closed. The port can buffer data within its *write-out* and *write-out-special* procedures.

- *name* — the name for the output port.
- *evt* — a synchronization event (see §11.2.1 "Events"; e.g., a semaphore or another port). The event is used in place of the output port when the port is supplied to synchronization procedures like `sync`. Thus, the event should be unblocked when the port is ready for writing at least one byte without blocking, or ready to make progress in flushing an internal buffer without blocking. The event must not unblock unless the port is ready for writing; otherwise, the guarantees of `sync` will be broken for the output port. Use `always-evt` if writes to the port always succeed without blocking.
- *write-out* — either an output port, which indicates that writes should be redirected to the given port, or a procedure of five arguments:
  - an immutable byte string containing bytes to write;
  - a non-negative exact integer for a starting offset (inclusive) into the byte string;
  - a non-negative exact integer for an ending offset (exclusive) into the byte string;
  - a boolean; `#f` indicates that the port is allowed to keep the written bytes in a buffer, and that it is allowed to block indefinitely; `#t` indicates that the write should not block, and that the port should attempt to flush its buffer and completely write new bytes instead of buffering them;
  - a boolean; `#t` indicates that if the port blocks for a write, then it should enable breaks while blocking (e.g., using `sync/enable-break`); this argument is always `#f` if the fourth argument is `#t`.

  The procedure returns one of the following:

  - a non-negative exact integer representing the number of bytes written or buffered;
  - `#f` if no bytes could be written, perhaps because the internal buffer could not be completely flushed;
  - a pipe output port (when buffering is allowed and not when flushing) for buffering bytes as long as the pipe is not full and until *write-out* or *write-out-special* is called; or
  - a synchronizable event (see §11.2.1 "Events") other than a pipe output port that acts like the result of `write-bytes-avail-evt` to complete the write.

  Since *write-out* can produce an event, an acceptable implementation of *write-out* is to pass its first three arguments to the port's *get-write-evt*. Some port implementors, however, may choose not to provide *get-write-evt* (perhaps because writes cannot be made atomic), or may implement `write-proc` to enable a fast path for non-blocking writes or to enable buffering.

From a user's perspective, the difference between buffered and completely written data is (1) buffered data can be lost in the future due to a failed write, and (2) `flush-output` forces all buffered data to be completely written. Under no circumstances is buffering required.

If the start and end indices are the same, then the fourth argument to *write-out* will be `#f`, and the write request is actually a flush request for the port's buffer (if any), and the result should be `0` for a successful flush (or if there is no buffer).

The result should never be `0` if the start and end indices are different, otherwise the `exn:fail:contract` exception is raised. Similarly, the `exn:fail:contract` exception is raised if *write-out* returns a pipe output port when buffering is disallowed or when it is called for flushing. If a returned integer is larger than the supplied byte-string range, the `exn:fail:contract` exception is raised.

The `#f` result should be avoided, unless the next write attempt is likely to work. Otherwise, if data cannot be written, return an event instead.

An event returned by *write-out* can return `#f` or another event like itself, in contrast to events produced by `write-bytes-avail-evt` or *get-write-evt*. A writing process loops with `sync` until it obtains a non-event result.

The *write-out* procedure is always called with breaks disabled, independent of whether breaks were enabled when the write was requested by a client of the port. If breaks were enabled for a blocking operation, then the fifth argument to *write-out* will be `#t`, which indicates that *write-out* should re-enable breaks while blocking.

If the writing procedure raises an exception, due to write or commit operations, it must not have committed any bytes (though it may have committed previously buffered bytes).

A port's writing procedure may be called in multiple threads simultaneously (if the port is accessible in multiple threads). The port is responsible for its own internal synchronization. Note that improper implementation of such synchronization mechanisms might cause a non-blocking write procedure to block.

- *close* — a procedure of zero arguments that is called to close the port. The port is not considered closed until the closing procedure returns. The port's procedures will never be used again via the port after it is closed. However, the closing procedure can be called simultaneously in multiple threads (if the port is accessible in multiple threads), and it may be called during a call to the other procedures in another thread; in the latter case, any outstanding writes or flushes should be terminated immediately with an error.

- *write-out-special* — either `#f` (the default), an output port (which indicates that special writes should be redirected to the given port), or a procedure to handle `write-special` calls for the port. If `#f`, then the port does not support special output, and `port-writes-special?` will return `#f` when applied to the port.

  If a procedure is supplied, it takes three arguments: the special value to write, a boolean that is `#f` if the procedure can buffer the special value and block indefinitely,

and a boolean that is `#t` if the procedure should enable breaks while blocking. The result is one of the following:

  - a non-event true value, which indicates that the special is written;
  - `#f` if the special could not be written, perhaps because an internal buffer could not be completely flushed;
  - a synchronizable event (see §11.2.1 “Events”) that acts like the result of *`get-write-special-evt`* to complete the write.

Since *`write-out-special`* can return an event, passing the first argument to an implementation of *`get-write-special-evt`* is acceptable as a *`write-out-special`*.

As for *`write-out`*, the `#f` result is discouraged, since it can lead to busy waiting. Also as for *`write-out`*, an event produced by *`write-out-special`* is allowed to produce `#f` or another event like itself. The *`write-out-special`* procedure is always called with breaks disabled, independent of whether breaks were enabled when the write was requested by a client of the port.

- *`get-write-evt`* — either `#f` (the default) or a procedure of three arguments:

  - an immutable byte string containing bytes to write;
  - a non-negative exact integer for a starting offset (inclusive) into the byte string; and
  - a non-negative exact integer for an ending offset (exclusive) into the byte string.

  The result is a synchronizable event (see §11.2.1 “Events”) to act as the result of `write-bytes-avail-evt` for the port (i.e., to complete a write or flush), which becomes available only as data is committed to the port’s underlying device, and whose result is the number of bytes written.

  If *`get-write-evt`* is `#f`, then `port-writes-atomic?` will produce `#f` when applied to the port, and the port will not be a valid argument to procedures such as `write-bytes-avail-evt`. Otherwise, an event returned by *`get-write-evt`* must not cause data to be written to the port unless the event is chosen in a synchronization, and it must write to the port if the event is chosen (i.e., the write must appear atomic with respect to the synchronization).

  If the event’s result integer is larger than the supplied byte-string range, the `exn:fail:contract` exception is raised by a wrapper on the event. If the start and end indices are the same (i.e., no bytes are to be written), then the event should produce `0` when the buffer is completely flushed. (If the port has no buffer, then it is effectively always flushed.)

  If the event raises an exception, due to write or commit operations, it must not have committed any new bytes (though it may have committed previously buffered bytes).

  Naturally, a port’s events may be used in multiple threads simultaneously (if the port is accessible in multiple threads). The port is responsible for its own internal synchronization.

- *get-write-special-evt* — either #f (the default), or a procedure to handle write-special-evt calls for the port. This argument must be #f if either *write-out-special* or *get-write-evt* is #f, and it must be a procedure if both of those arguments are procedures.

  If it is a procedure, it takes one argument: the special value to write. The resulting event (with its constraints) is analogous to the result of *get-write-evt*.

  If the event raises an exception, due to write or commit operations, it must not have committed the special value (though it may have committed previously buffered bytes and values).

- *get-location* — either #f (the default), or a procedure that takes no arguments and returns three values: the line number for the next item written to the port's stream (a positive number or #f), the column number for the next item written to port's stream (a non-negative number or #f), and the position for the next item written to port's stream (a positive number or #f). See also §13.1.4 "Counting Positions, Lines, and Columns".

  This procedure is called to implement port-next-location for the port, but only if line counting is enabled for the port via port-count-lines! (in which case *count-lines!* is called).

- *count-lines!* — a procedure of no arguments that is called if and when line counting is enabled for the port. The default procedure is void.

- *init-position* — normally an exact, positive integer that determines the position of the port's first item, which is used by file-position or when line counting is *not* enabled for the port. The default is 1. If *init-position* is #f, the port is treated as having an unknown position. If *init-position* is a port, then the given port's position is always used for the new port's position. If *init-position* is a procedure, it is called as needed to obtain the port's position.

- *buffer-mode* — either #f (the default) or a procedure that accepts zero or one arguments. If *buffer-mode* is #f, then the resulting port does not support a buffer-mode setting. Otherwise, the procedure is called with one symbol argument ('block, 'line, or 'none) to set the buffer mode, and it is called with zero arguments to get the current buffer mode. In the latter case, the result must be 'block, 'line, 'none, or #f (unknown). See §13.1.3 "Port Buffers and Positions" for more information on buffer modes.

Examples:

```
; A port that writes anything to nowhere:
> (define /dev/null-out
    (make-output-port
     'null
     always-evt
     (lambda (s start end non-block? breakable?) (- end start))
     void
```

```
     (lambda (special non-block? breakable?) #t)
     (lambda (s start end) (wrap-evt
                            always-evt
                            (lambda (x)
                              (- end start))))
     (lambda (special) always-evt)))

> (display "hello" /dev/null-out)

> (write-bytes-avail #"hello" /dev/null-out)
5
> (write-special 'hello /dev/null-out)
#t
> (sync (write-bytes-avail-evt #"hello" /dev/null-out))
5
; A port that accumulates bytes as characters in a list,
;  but not in a thread-safe way:
> (define accum-list null)

> (define accumulator/not-thread-safe
    (make-output-port
     'accum/not-thread-safe
     always-evt
     (lambda (s start end non-block? breakable?)
       (set! accum-list
             (append accum-list
                     (map integer->char
                          (bytes->list (subbytes s start end)))))
       (- end start))
     void))

> (display "hello" accumulator/not-thread-safe)

> accum-list
'(#\h #\e #\l #\l #\o)
; Same as before, but with simple thread-safety:
> (define accum-list null)

> (define accumulator
    (let* ([lock (make-semaphore 1)]
           [lock-peek-evt (semaphore-peek-evt lock)])
      (make-output-port
       'accum
       lock-peek-evt
       (lambda (s start end non-block? breakable?)
         (if (semaphore-try-wait? lock)
```

```
            (begin
              (set! accum-list
                    (append accum-list
                            (map integer->char
                                 (bytes->list
                                  (subbytes s start end)))))
              (semaphore-post lock)
              (- end start))
            ; Cheap strategy: block until the list is unlocked,
            ;   then return 0, so we get called again
            (wrap-evt
             lock-peek
             (lambda (x) 0))))
       void)))

> (display "hello" accumulator)

> accum-list
'(#\h #\e #\l #\l #\o)
; A port that transforms data before sending it on
;  to another port. Atomic writes exploit the
;  underlying port's ability for atomic writes.
> (define (make-latin-1-capitalize port)
    (define (byte-upcase s start end)
      (list->bytes
       (map (lambda (b) (char->integer
                         (char-upcase
                          (integer->char b))))
            (bytes->list (subbytes s start end)))))
    (make-output-port
     'byte-upcase
     ; This port is ready when the original is ready:
     port
     ; Writing procedure:
     (lambda (s start end non-block? breakable?)
       (let ([s (byte-upcase s start end)])
         (if non-block?
             (write-bytes-avail* s port)
             (begin
               (display s port)
               (bytes-length s)))))
     ; Close procedure -- close original port:
     (lambda () (close-output-port port))
     #f
     ; Write event:
     (and (port-writes-atomic? port)
```

```
          (lambda (s start end)
            (write-bytes-avail-evt
             (byte-upcase s start end)
             port)))))

> (define orig-port (open-output-string))

> (define cap-port (make-latin-1-capitalize orig-port))

> (display "Hello" cap-port)

> (get-output-string orig-port)
"HELLO"
> (sync (write-bytes-avail-evt #"Bye" cap-port))
3
> (get-output-string orig-port)
"HELLOBYE"
```

13.1.10 More Port Constructors, Procedures, and Events

```
(require racket/port)
```
package: base

The bindings documented in this section are provided by the `racket/port` and `racket` libraries, but not `racket/base`.

Port String and List Conversions

```
(port->list [r in]) → (listof any/c)
  r : (input-port? . -> . any/c) = read
  in : input-port? = (current-input-port)
```

Returns a list whose elements are produced by calling *r* on *in* until it produces `eof`.

Examples:

```
> (define (read-number input-port)
    (define char (read-char input-port))
    (if (eof-object? char)
     char
     (string->number (string char))))

> (port->list read-number (open-input-string "12345"))
'(1 2 3 4 5)
```

```
(port->string [in]) → string?
  in : input-port? = (current-input-port)
```

Reads all characters from *in* and returns them as a string.

Example:

```
> (port->string (open-input-string "hello world"))
"hello world"
```

```
(port->bytes [in]) → bytes?
  in : input-port? = (current-input-port)
```

Reads all bytes from *in* and returns them as a byte string.

Example:

```
> (port->bytes (open-input-string "hello world"))
#"hello world"
```

```
(port->lines [in #:line-mode line-mode]) → (listof string?)
  in : input-port? = (current-input-port)
  line-mode : (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
            = 'any
```

Read all characters from *in*, breaking them into lines. The *line-mode* argument is the same as the second argument to `read-line`, but the default is `'any` instead of `'linefeed`.

Example:

```
> (port->lines
   (open-input-string "line 1\nline 2\n  line 3\nline 4"))
'("line 1" "line 2" "  line 3" "line 4")
```

```
(port->bytes-lines [in
                    #:line-mode line-mode]) → (listof bytes?)
  in : input-port? = (current-input-port)
  line-mode : (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
            = 'any
```

Like `port->lines`, but reading bytes and collecting them into lines like `read-bytes-line`.

Example:

```
> (port->bytes-lines
   (open-input-string "line 1\nline 2\n  line 3\nline 4"))
'(#"line 1" #"line 2" #"  line 3" #"line 4")
```

```
(display-lines  lst
                [out
                 #:separator separator]) → void?
  lst : list?
  out : output-port? = (current-output-port)
  separator : any/c = #"\n"
```

Uses `display` on each element of *lst* to *out*, adding *separator* after each element.

```
(call-with-output-string proc) → string?
  proc : (output-port? . -> . any)
```

Calls *proc* with an output port that accumulates all output into a string, and returns the string.

The port passed to *proc* is like the one created by `open-output-string`, except that it is wrapped via `dup-output-port`, so that *proc* cannot access the port's content using `get-output-string`. If control jumps back into *proc*, the port continues to accumulate new data, and `call-with-output-string` returns both the old data and newly accumulated data.

```
(call-with-output-bytes proc) → bytes?
  proc : (output-port? . -> . any)
```

Like `call-with-output-string`, but returns the accumulated result in a byte string instead of a string. Furthermore, the port's content is emptied when `call-with-output-bytes` returns, so that if control jumps back into *proc* and returns a second time, only the newly accumulated bytes are returned.

```
(with-output-to-string proc) → string?
  proc : (-> any)
```

Equivalent to

```
(call-with-output-string
 (lambda (p) (parameterize ([current-output-port p])
               (proc))))
```

```
(with-output-to-bytes proc) → bytes?
  proc : (-> any)
```

Equivalent to

```
(call-with-output-bytes
 (lambda (p) (parameterize ([current-output-port p])
               (proc))))
```

```
(call-with-input-string str proc) → any
  str : string?
  proc : (input-port? . -> . any)
```

Equivalent to `(proc (open-input-string str))`.

```
(call-with-input-bytes bstr proc) → any
  bstr : bytes?
  proc : (input-port? . -> . any)
```

Equivalent to `(proc (open-input-bytes bstr))`.

```
(with-input-from-string str proc) → any
  str : string?
  proc : (-> any)
```

Equivalent to

```
(parameterize ([current-input-port (open-input-string str)])
  (proc))
```

```
(with-input-from-bytes bstr proc) → any
  bstr : bytes?
  proc : (-> any)
```

Equivalent to

```
(parameterize ([current-input-port (open-input-bytes str)])
  (proc))
```

Creating Ports

```
(input-port-append close-at-eof? in ...) → input-port?
  close-at-eof? : any/c
  in : input-port?
```

Takes any number of input ports and returns an input port. Reading from the input port draws bytes (and special non-byte values) from the given input ports in order. If *close-at-eof?* is true, then each port is closed when an end-of-file is encountered from the port, or when

the result input port is closed. Otherwise, data not read from the returned input port remains available for reading in its original input port.

See also merge-input, which interleaves data from multiple input ports as it becomes available.

```
(make-input-port/read-to-peek  name
                               read-in
                               fast-peek
                               close
                              [get-location
                               count-lines!
                               init-position
                               buffer-mode
                               buffering?
                               on-consumed]) → input-port?
  name : any/c
           (bytes?
            . -> . (or/c exact-nonnegative-integer?
  read-in :                eof-object?
                           procedure?
                           evt?))
             (or/c #f
                   (bytes? exact-nonnegative-integer?
                    (bytes? exact-nonnegative-integer?
                     . -> . (or/c exact-nonnegative-integer?
                                  eof-object?
                                  procedure?
  fast-peek :                     evt?
                                  #f))
                    . -> . (or/c exact-nonnegative-integer?
                                 eof-object?
                                 procedure?
                                 evt?
                                 #f)))
  close : (-> any)
                 (or/c
                  (->
                   (values
  get-location :    (or/c exact-positive-integer? #f)    = #f
                    (or/c exact-nonnegative-integer? #f)
                    (or/c exact-positive-integer? #f)))
                  #f)
  count-lines! : (-> any) = void
  init-position : exact-positive-integer? = 1
```

```
                (or/c (case-> ((or/c 'block 'none) . -> . any)
buffer-mode :                 (-> (or/c 'block 'none #f)))
                      #f)
              = #f
buffering? : any/c = #f
                (or/c ((or/c exact-nonnegative-integer? eof-object?
                             procedure? evt?)
on-consumed :
                       . -> . any)
                      #f)
              = #f
```

Similar to `make-input-port`, but if the given *read-in* returns an event, the event's value must be `0`. The resulting port's peek operation is implemented automatically (in terms of *read-in*) in a way that can handle special non-byte values. The progress-event and commit operations are also implemented automatically. The resulting port is thread-safe, but not kill-safe (i.e., if a thread is terminated or suspended while using the port, the port may become damaged).

The *read-in*, *close*, *get-location*, *count-lines!*, *init-position*, and *buffer-mode* procedures are the same as for `make-input-port`.

The *fast-peek* argument can be either `#f` or a procedure of three arguments: a byte string to receive a peek, a skip count, and a procedure of two arguments. The *fast-peek* procedure can either implement the requested peek, or it can dispatch to its third argument to implement the peek. The *fast-peek* is not used when a peek request has an associated progress event.

The *buffering?* argument determines whether *read-in* can be called to read more characters than are immediately demanded by the user of the new port. If *buffer-mode* is not `#f`, then *buffering?* determines the initial buffer mode, and *buffering?* is enabled after a buffering change only if the new mode is `'block`.

If *on-consumed* is not `#f`, it is called when data is read (or committed) from the port, as opposed to merely peeked. The argument to *on-consumed* is the result value of the port's reading procedure, so it can be an integer or any result from *read-in*.

```
(make-limited-input-port  in
                          limit
                         [close-orig?]) → input-port?
  in : input-port?
  limit : exact-nonnegative-integer?
  close-orig? : any/c = #t
```

Returns a port whose content is drawn from *in*, but where an end-of-file is reported after *limit* bytes (and non-byte special values) have been read. If *close-orig?* is true, then the original port is closed if the returned port is closed.

Bytes are consumed from *in* only when they are consumed from the returned port. In particular, peeking into the returned port peeks into the original port.

If *in* is used directly while the resulting port is also used, then the *limit* bytes provided by the port need not be contiguous parts of the original port's stream.

```
(make-pipe-with-specials [limit
                          in-name
                          out-name]) → input-port? output-port?
  limit : exact-nonnegative-integer? = #f
  in-name : any/c = 'pipe
  out-name : any/c = 'pipe
```

Returns two ports: an input port and an output port. The ports behave like those returned by `make-pipe`, except that the ports support non-byte values written with procedures such as `write-special` and read with procedures such as `get-byte-or-special`.

The *limit* argument determines the maximum capacity of the pipe in bytes, but this limit is disabled if special values are written to the pipe before *limit* is reached. The limit is re-enabled after the special value is read from the pipe.

The optional *in-name* and *out-name* arguments determine the names of the result ports.

```
(merge-input a-in b-in [buffer-limit]) → input-port?
  a-in : input-port?
  b-in : input-port?
  buffer-limit : (or/c exact-nonnegative-integer? #f) = 4096
```

Accepts two input ports and returns a new input port. The new port merges the data from two original ports, so data can be read from the new port whenever it is available from either of the two original ports. The data from the original ports are interleaved. When an end-of-file has been read from an original port, it no longer contributes characters to the new port. After an end-of-file has been read from both original ports, the new port returns end-of-file. Closing the merged port does not close the original ports.

The optional *buffer-limit* argument limits the number of bytes to be buffered from *a-in* and *b-in*, so that the merge process does not advance arbitrarily beyond the rate of consumption of the merged data. A `#f` value disables the limit. As for `make-pipe-with-specials`, *buffer-limit* does not apply when a special value is produced by one of the input ports before the limit is reached.

See also `input-port-append`, which concatenates input streams instead of interleaving them.

```
(open-output-nowhere [name special-ok?]) → output-port?
  name : any/c = 'nowhere
  special-ok? : any/c = #t
```

Creates and returns an output port that discards all output sent to it (without blocking). The *name* argument is used as the port's name. If the *special-ok?* argument is true, then the resulting port supports `write-special`, otherwise it does not.

```
(peeking-input-port  in
                     [name
                      skip
                      #:init-position init-position]) → input-port
  in : input-port?
  name : any/c = (object-name in)
  skip : exact-nonnegative-integer? = 0
  init-position : exact-positive-integer? = 1
```

Returns an input port whose content is determined by peeking into *in*. In other words, the resulting port contains an internal skip count, and each read of the port peeks into *in* with the internal skip count, and then increments the skip count according to the amount of data successfully peeked.

The optional *name* argument is the name of the resulting port. The *skip* argument is the port initial skip count, and it defaults to `0`.

The resulting port's initial position (as reported by `file-position`) is `(- init-position 1)`, no matter the position of *in*.

The resulting port supports buffering, and a `'block` buffer mode allows the port to peek further into *in* than requested. The resulting port's initial buffer mode is `'block`, unless *in* supports buffer mode and its mode is initially `'none` (i.e., the initial buffer mode is taken from *in* when it supports buffering). If *in* supports buffering, adjusting the resulting port's buffer mode via `file-stream-buffer-mode` adjusts *in*'s buffer mode.

For example, when you read from a peeking port, you see the same answers as when you read from the original port:

Examples:

```
> (define an-original-port (open-input-string "123456789"))

> (define a-peeking-port (peeking-input-port an-original-port))

> (file-stream-buffer-mode a-peeking-port 'none)

> (read-string 3 a-peeking-port)
"123"
> (read-string 3 an-original-port)
"123"
```

Beware that the read from the original port is invisible to the peeking port, which keeps its own separate internal counter, and thus interleaving reads on the two ports can produce

confusing results. Continuing the example before, if we read three more characters from the peeking port, we end up skipping over the 456 in the port (but only because we disabled buffering above):

Example:

```
> (read-string 3 a-peeking-port)
"789"
```

If we had left the buffer mode of a-peeking-port alone, that last read-string would have likely produced "456" as a result of buffering bytes from an-original-port earlier.

Changed in version 6.1.0.3 of package base: Enabled buffering and buffer-mode adjustments via file-stream-buffer-mode, and set the port's initial buffer mode to that of *in*.

```
(reencode-input-port in
                     encoding
                     [error-bytes
                      close?
                      name
                      convert-newlines?
                      enc-error])        → input-port?
  in : input-port?
  encoding : string?
  error-bytes : (or/c #f bytes?) = #f
  close? : any/c = #f
  name : any/c = (object-name in)
  convert-newlines? : any/c = #f
  enc-error : (string? input-port? . -> . any)
            = (lambda (msg port) (error ...))
```

Produces an input port that draws bytes from *in*, but converts the byte stream using (bytes-open-converter encoding-str "UTF-8"). In addition, if *convert-newlines?* is true, then decoded sequences that correspond to UTF-8 encodings of "\r\n", "\r\u0085", "\r", "\u0085", and "\u2028" are all converted to the UTF-8 encoding of "\n".

If *error-bytes* is provided and not #f, then the given byte sequence is used in place of bytes from *in* that trigger conversion errors. Otherwise, if a conversion is encountered, *enc-error* is called, which must raise an exception.

If *close?* is true, then closing the result input port also closes *in*. The *name* argument is used as the name of the result input port.

In non-buffered mode, the resulting input port attempts to draw bytes from *in* only as needed to satisfy requests. Toward that end, the input port assumes that at least *n* bytes must be read to satisfy a request for *n* bytes. (This is true even if the port has already drawn some bytes, as long as those bytes form an incomplete encoding sequence.)

```
(reencode-output-port  out
                       encoding
                      [error-bytes
                       close?
                       name
                       newline-bytes
                       enc-error])    → output-port?
  out : output-port?
  encoding : string?
  error-bytes : (or/c #f bytes?) = #f
  close? : any/c = #f
  name : any/c = (object-name out)
  newline-bytes : (or/c #f bytes?) = #f
  enc-error : (string? output-port? . -> . any)
            = (lambda (msg port) (error ...))
```

Produces an output port that directs bytes to *out*, but converts its byte stream using `(bytes-open-converter "UTF-8" encoding-str)`. In addition, if *newline-bytes* is not `#f`, then bytes written to the port that are the UTF-8 encoding of `"\n"` are first converted to *newline-bytes* (before applying the convert from UTF-8 to `encoding-str`).

If *error-bytes* is provided and not `#f`, then the given byte sequence is used in place of bytes that have been sent to the output port and that trigger conversion errors. Otherwise, *enc-error* is called, which must raise an exception.

If *close?* is true, then closing the result output port also closes *out*. The *name* argument is used as the name of the result output port.

The resulting port supports buffering, and the initial buffer mode is `(or (file-stream-buffer-mode out) 'block)`. In `'block` mode, the port's buffer is flushed only when it is full or a flush is requested explicitly. In `'line` mode, the buffer is flushed whenever a newline or carriage-return byte is written to the port. In `'none` mode, the port's buffer is flushed after every write. Implicit flushes for `'line` or `'none` leave bytes in the buffer when they are part of an incomplete encoding sequence.

The resulting output port does not support atomic writes. An explicit flush or special-write to the output port can hang if the most recently written bytes form an incomplete encoding sequence.

When the port is buffered, a flush callback is registered with the current plumber to flush the buffer.

```
(dup-input-port in [close?]) → input-port?
  in : input-port?
  close? : any/c = #f
```

Returns an input port that draws directly from *in*. Closing the resulting port closes *in* only if *close?* is #t.

The new port is initialized with the port read handler of *in*, but setting the handler on the result port does not affect reading directly from *in*.

```
(dup-output-port out [close?]) → output-port?
  out : output-port?
  close? : any/c = #f
```

Returns an output port that propagates data directly to *out*. Closing the resulting port closes *out* only if *close?* is #t.

The new port is initialized with the port display handler and port write handler of *out*, but setting the handlers on the result port does not affect writing directly to *out*.

```
(relocate-input-port  in
                      line
                      column
                      position
                     [close?])  → input-port?
  in : input-port?
  line : (or/c exact-positive-integer? #f)
  column : (or/c exact-nonnegative-integer? #f)
  position : exact-positive-integer?
  close? : any/c = #t
```

Produces an input port that is equivalent to *in* except in how it reports location information. The resulting port's content starts with the remaining content of *in*, and it starts at the given line, column, and position. A #f for the line or column means that the line and column will always be reported as #f.

The *line* and *column* values are used only if line counting is enabled for *in* and for the resulting port, typically through port-count-lines!. The *column* value determines the column for the first line (i.e., the one numbered *line*), and later lines start at column 0. The given *position* is used even if line counting is not enabled.

When line counting is on for the resulting port, reading from *in* instead of the resulting port increments location reports from the resulting port. Otherwise, the resulting port's position does not increment when data is read from *in*.

If *close?* is true, then closing the resulting port also closes *in*. If *close?* is #f, then closing the resulting port does not close *in*.

```
(relocate-output-port  out
                       line
                       column
                       position
                      [close?]) → output-port?
  out : output-port?
  line : (or/c exact-positive-integer? #f)
  column : (or/c exact-nonnegative-integer? #f)
  position : exact-positive-integer?
  close? : any/c = #t
```

Like relocate-input-port, but for output ports.

```
(transplant-input-port  in
                        get-location
                        init-pos
                       [close?
                        count-lines!]) → input-port?
  in : input-port?
                 (or/c
                  (->
                   (values
  get-location :    (or/c exact-positive-integer? #f)
                    (or/c exact-nonnegative-integer? #f)
                    (or/c exact-positive-integer? #f)))
                  #f)
  init-pos : exact-positive-integer?
  close? : any/c = #t
  count-lines! : (-> any) = void
```

Like relocate-input-port, except that arbitrary position information can be produced (when line counting is enabled) via *get-location*, which is used as for make-input-port. If *get-location* is #f, then the port counts lines in the usual way starting from *init-pos*, independent of locations reported by *in*.

If *count-lines!* is supplied, it is called when line counting is enabled for the resulting port. The default is void.

```
(transplant-output-port  out
                         get-location
                         init-pos
                        [close?
                         count-lines!]) → output-port?
  out : output-port?
```

```
                (or/c
                 (->
                  (values
  get-location :    (or/c exact-positive-integer? #f)
                    (or/c exact-nonnegative-integer? #f)
                    (or/c exact-positive-integer? #f)))
                 #f)
  init-pos : exact-positive-integer?
  close? : any/c = #t
  count-lines! : (-> any) = void
```

Like transplant-input-port, but for output ports.

```
(filter-read-input-port  in
                         read-wrap
                         peek-wrap
                        [close?])   → input-port?
  in : input-port?
             (bytes? (or/c exact-nonnegative-integer?
                           eof-object?
                           procedure?
                           evt?)
  read-wrap :          . -> .
                       (or/c exact-nonnegative-integer?
                             eof-object?
                             procedure?
                             evt?))
             (bytes? exact-nonnegative-integer? (or/c evt? #f)
                     (or/c exact-nonnegative-integer?
                      eof-object?
                      procedure?
                      evt?
  peek-wrap :         #f)
               . -> . (or/c exact-nonnegative-integer?
                      eof-object?
                      procedure?
                      evt?
                      #f))
  close? : any/c = #t
```

Creates a port that draws from *in*, but each result from the port's read and peek procedures (in the sense of make-input-port) is filtered by *read-wrap* and *peek-wrap*. The filtering procedures each receive both the arguments and results of the read and peek procedures on *in* for each call.

If *close?* is true, then closing the resulting port also closes *in*.

```
(special-filter-input-port in proc [close?]) → input-port?
  in : input-port?
         (procedure? bytes? . -> . (or/c exact-nonnegative-integer?
                                         eof-object?
  proc :                                 procedure?
                                         evt?))
  close? : any/c = #t
```

Produces an input port that is equivalent to *in*, except that when *in* produces a procedure to access a special value, *proc* is applied to the procedure to allow the special value to be replaced with an alternative. The *proc* is called with the special-value procedure and the byte string that was given to the port's read or peek function (see `make-input-port`), and the result is used as the read or peek function's result. The *proc* can modify the byte string to substitute a byte for the special value, but the byte string is guaranteed only to hold at least one byte.

If *close?* is true, then closing the resulting input port also closes *in*.

Port Events

```
(eof-evt in) → evt?
  in : input-port?
```

Returns a synchronizable event that is ready when *in* produces an `eof`. If *in* produces a mid-stream `eof`, the `eof` is consumed by the event only if the event is chosen in a synchronization.

```
(read-bytes-evt k in) → evt?
  k : exact-nonnegative-integer?
  in : input-port?
```

Returns a synchronizable event that is ready when *k* bytes can be read from *in*, or when an end-of-file is encountered in *in*. If *k* is `0`, then the event is ready immediately with `""`. For non-zero *k*, if no bytes are available before an end-of-file, the event's result is `eof`. Otherwise, the event's result is a byte string of up to *k* bytes, which contains as many bytes as are available (up to *k*) before an available end-of-file. (The result is a byte string on less than *k* bytes only when an end-of-file is encountered.)

Bytes are read from the port if and only if the event is chosen in a synchronization, and the returned bytes always represent contiguous bytes in the port's stream.

The event can be synchronized multiple times—event concurrently—and each synchronization corresponds to a distinct read request.

The *in* must support progress events, and it must not produce a special non-byte value during the read attempt.

```
(read-bytes!-evt bstr in progress-evt) → evt?
  bstr : (and/c bytes? (not/c immutable?))
  in : input-port?
  progress-evt : (or/c progress-evt? #f)
```

Like read-bytes-evt, except that the read bytes are placed into bstr, and the number of bytes to read corresponds to (bytes-length bstr). The event's result is either eof or the number of read bytes.

The bstr may be mutated any time after the first synchronization attempt on the event and until either the event is selected, a non-#f progress-evt is ready, or the current custodian (at the time of synchronization) is shut down. Note that there is no time bound otherwise on when bstr might be mutated if the event is not selected by a synchronzation; nevertheless, multiple synchronization attempts can use the same result from read-bytes!-evt as long as there is no intervening read on in until one of the synchronization attempts selects the event.

```
(read-bytes-avail!-evt bstr in) → evt?
  bstr : (and/c bytes? (not/c immutable?))
  in : input-port?
```

Like read-bytes!-evt, except that the event reads only as many bytes as are immediately available, after at least one byte or one eof becomes available.

```
(read-string-evt k in) → evt?
  k : exact-nonnegative-integer?
  in : input-port?
```

Like read-bytes-evt, but for character strings instead of byte strings.

```
(read-string!-evt str in) → evt?
  str : (and/c string? (not/c immutable?))
  in : input-port?
```

Like read-bytes!-evt, but for a character string instead of a byte string.

```
(read-line-evt in mode) → evt?
  in : input-port?
  mode : (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
```

Returns a synchronizable event that is ready when a line of characters or end-of-file can be read from in. The meaning of mode is the same as for read-line. The event result is the read line of characters (not including the line separator).

A line is read from the port if and only if the event is chosen in a synchronization, and the returned line always represents contiguous bytes in the port's stream.

```
(read-bytes-line-evt in mode) → evt?
  in : input-port?
  mode : (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
```

Like `read-line-evt`, but returns a byte string instead of a string.

```
(peek-bytes-evt k skip progress-evt in) → evt?
  k : exact-nonnegative-integer?
  skip : exact-nonnegative-integer?
  progress-evt : (or/c progress-evt? #f)
  in : input-port?
(peek-bytes!-evt bstr skip progress-evt in) → evt?
  bstr : (and/c bytes? (not/c immutable?))
  skip : exact-nonnegative-integer?
  progress-evt : (or/c progress-evt? #f)
  in : input-port?
(peek-bytes-avail!-evt bstr
                       skip
                       progress-evt
                       in)          → evt?
  bstr : (and/c bytes? (not/c immutable?))
  skip : exact-nonnegative-integer?
  progress-evt : (or/c progress-evt? #f)
  in : input-port?
(peek-string-evt k skip progress-evt in) → evt?
  k : exact-nonnegative-integer?
  skip : exact-nonnegative-integer?
  progress-evt : (or/c progress-evt? #f)
  in : input-port?
(peek-string!-evt str skip progress-evt in) → evt?
  str : (and/c string? (not/c immutable?))
  skip : exact-nonnegative-integer?
  progress-evt : (or/c progress-evt? #f)
  in : input-port?
```

Like the `read-...-evt` functions, but for peeking. The *skip* argument indicates the number of bytes to skip, and *progress-evt* indicates an event that effectively cancels the peek (so that the event never becomes ready). The *progress-evt* argument can be `#f`, in which case the event is never canceled.

```
(regexp-match-evt pattern in) → any
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  in : input-port?
```

Returns a synchronizable event that is ready when *pattern* matches the stream of

bytes/characters from *in*; see also regexp-match. The event's value is the result of the match, in the same form as the result of regexp-match.

If *pattern* does not require a start-of-stream match, then bytes skipped to complete the match are read and discarded when the event is chosen in a synchronization.

Bytes are read from the port if and only if the event is chosen in a synchronization, and the returned match always represents contiguous bytes in the port's stream. If not-yet-available bytes from the port might contribute to the match, the event is not ready. Similarly, if *pattern* begins with a start-of-stream ^ and the *pattern* does not initially match, then the event cannot become ready until bytes have been read from the port.

The event can be synchronized multiple times—even concurrently—and each synchronization corresponds to a distinct match request.

The *in* port must support progress events. If *in* returns a special non-byte value during the match attempt, it is treated like eof.

Copying Streams

```
(convert-stream from-encoding
                in
                to-encoding
                out)          → void?
  from-encoding : string?
  in : input-port?
  to-encoding : string?
  out : output-port?
```

Reads data from *in*, converts it using (bytes-open-converter *from-encoding* *to-encoding*) and writes the converted bytes to *out*. The convert-stream procedure returns after reaching eof in *in*.

If opening the converter fails, the exn:fail exception is raised. Similarly, if a conversion error occurs at any point while reading from *in*, then exn:fail exception is raised.

```
(copy-port in out ...+) → void?
  in : input-port?
  out : output-port?
```

Reads data from *in* and writes it back out to *out*, returning when *in* produces eof. The copy is efficient, and it is without significant buffer delays (i.e., a byte that becomes available on *in* is immediately transferred to *out*, even if future reads on *in* must block). If *in* produces a special non-byte value, it is transferred to *out* using write-special.

This function is often called from a "background" thread to continuously pump data from one stream to another.

If multiple *out*s are provided, case data from *in* is written to every *out*. The different *out*s block output to each other, because each block of data read from *in* is written completely to one *out* before moving to the next *out*. The *out*s are written in the provided order, so non-blocking ports (e.g., file output ports) should be placed first in the argument list.

13.2 Byte and String Input

```
(read-char [in]) → (or/c char? eof-object?)
  in : input-port? = (current-input-port)
```

Reads a single character from *in*—which may involve reading several bytes to UTF-8-decode them into a character (see §13.1 “Ports”); a minimal number of bytes are read/peeked to perform the decoding. If no bytes are available before an end-of-file, then `eof` is returned.

Examples:

```
> (let ([ip (open-input-string "S2")])
    (print (read-char ip))
    (newline)
    (print (read-char ip))
    (newline)
    (print (read-char ip)))
#\S
#\2
#<eof>

> (let ([ip (open-input-bytes #"\316\273")])
    ; The byte string contains UTF-8-encoded content:
    (print (read-char ip)))
#\λ
```

```
(read-byte [in]) → (or/c byte? eof-object?)
  in : input-port? = (current-input-port)
```

Reads a single byte from *in*. If no bytes are available before an end-of-file, then `eof` is returned.

Examples:

```
> (let ([ip (open-input-string "a")])
    ; The two values in the following list should be the same.
    (list (read-byte ip) (char->integer #\a)))
'(97 97)
```

```
> (let ([ip (open-input-string (string #\λ))])
    ; This string has a two byte-encoding.
    (list (read-byte ip) (read-byte ip) (read-byte ip)))
'(206 187 #<eof>)
```

```
(read-line [in mode]) → (or/c string? eof-object?)
  in : input-port? = (current-input-port)
  mode : (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
       = 'linefeed
```

Returns a string containing the next line of bytes from *in*.

Characters are read from *in* until a line separator or an end-of-file is read. The line separator is not included in the result string (but it is removed from the port's stream). If no characters are read before an end-of-file is encountered, `eof` is returned.

The *mode* argument determines the line separator(s). It must be one of the following symbols:

- `'linefeed` breaks lines on linefeed characters.
- `'return` breaks lines on return characters.
- `'return-linefeed` breaks lines on return-linefeed combinations. If a return character is not followed by a linefeed character, it is included in the result string; similarly, a linefeed that is not preceded by a return is included in the result string.
- `'any` breaks lines on any of a return character, linefeed character, or return-linefeed combination. If a return character is followed by a linefeed character, the two are treated as a combination.
- `'any-one` breaks lines on either a return or linefeed character, without recognizing return-linefeed combinations.

Return and linefeed characters are detected after the conversions that are automatically performed when reading a file in text mode. For example, reading a file in text mode on Windows automatically changes return-linefeed combinations to a linefeed. Thus, when a file is opened in text mode, `'linefeed` is usually the appropriate `read-line` mode.

Examples:

```
> (let ([ip (open-input-string "x\ny\n")])
    (read-line ip))
"x"
> (let ([ip (open-input-string "x\ny\n")])
    (read-line ip 'return))
"x\ny\n"
```

```
> (let ([ip (open-input-string "x\ry\r")])
    (read-line ip 'return))
"x"
> (let ([ip (open-input-string "x\r\ny\r\n")])
    (read-line ip 'return-linefeed))
"x"
> (let ([ip (open-input-string "x\r\ny\nz")])
    (list (read-line ip 'any) (read-line ip 'any)))
'("x" "y")
> (let ([ip (open-input-string "x\r\ny\nz")])
    (list (read-line ip 'any-one) (read-line ip 'any-one)))
'("x" "")
```

```
(read-bytes-line [in mode]) → (or/c bytes? eof-object?)
  in : input-port? = (current-input-port)
  mode : (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
       = 'linefeed
```

Like `read-line`, but reads bytes and produces a byte string.

```
(read-string amt [in]) → (or/c string? eof-object?)
  amt : exact-nonnegative-integer?
  in : input-port? = (current-input-port)
```

To read an entire port as a string, use `port->string`.

Returns a string containing the next *amt* characters from *in*.

If *amt* is `0`, then the empty string is returned. Otherwise, if fewer than *amt* characters are available before an end-of-file is encountered, then the returned string will contain only those characters before the end-of-file; that is, the returned string's length will be less than *amt*. (A temporary string of size *amt* is allocated while reading the input, even if the size of the result is less than *amt* characters.) If no characters are available before an end-of-file, then `eof` is returned.

If an error occurs during reading, some characters may be lost; that is, if `read-string` successfully reads some characters before encountering an error, the characters are dropped.

Example:

```
> (let ([ip (open-input-string "supercalifragilisticexpialidocious")])
    (read-string 5 ip))
"super"
```

```
(read-bytes amt [in]) → (or/c bytes? eof-object?)
  amt : exact-nonnegative-integer?
  in : input-port? = (current-input-port)
```

To read an entire port as bytes, use `port->bytes`.

Like `read-string`, but reads bytes and produces a byte string.

Example:

```
> (let ([ip (open-input-bytes
                    (bytes 6
                           115 101 99 114 101
                           116))])
    (define length (read-byte ip))
    (bytes->string/utf-8 (read-bytes length ip)))
"secret"
```

```
(read-string! str [in start-pos end-pos])
 → (or/c exact-positive-integer? eof-object?)
  str : (and/c string? (not/c immutable?))
  in : input-port? = (current-input-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (string-length str)
```

Reads characters from *in* like read-string, but puts them into *str* starting from index *start-pos* (inclusive) up to *end-pos* (exclusive). Like substring, the exn:fail:contract exception is raised if *start-pos* or *end-pos* is out-of-range for *str*.

If the difference between *start-pos* and *end-pos* is 0, then 0 is returned and *str* is not modified. If no bytes are available before an end-of-file, then eof is returned. Otherwise, the return value is the number of characters read. If *m* characters are read and $m < \textit{end-pos} - \textit{start-pos}$, then *str* is not modified at indices $\textit{start-pos} + m$ through *end-pos*.

Example:

```
> (let ([buffer (make-string 10 #\_)]
        [ip (open-input-string "cketRa")])
    (printf "~s\n" buffer)
    (read-string! buffer ip 2 6)
    (printf "~s\n" buffer)
    (read-string! buffer ip 0 2)
    (printf "~s\n" buffer))
"__________"
"__cket____"
"Racket____"
```

```
(read-bytes! bstr [in start-pos end-pos])
 → (or/c exact-positive-integer? eof-object?)
  bstr : bytes?
  in : input-port? = (current-input-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like read-string!, but reads bytes, puts them into a byte string, and returns the number of bytes read.

Example:

```
> (let ([buffer (make-bytes 10 (char->integer #\_))]
        [ip (open-input-string "cketRa")])
    (printf "~s\n" buffer)
    (read-bytes! buffer ip 2 6)
    (printf "~s\n" buffer)
    (read-bytes! buffer ip 0 2)
    (printf "~s\n" buffer))
#"__________"
#"__cket____"
#"Racket____"
```

```
(read-bytes-avail! bstr [in start-pos end-pos])
 → (or/c exact-positive-integer? eof-object? procedure?)
  bstr : bytes?
  in : input-port? = (current-input-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like read-bytes!, but returns without blocking after having read the immediately available bytes, and it may return a procedure for a "special" result. The read-bytes-avail! procedure blocks only if no bytes (or specials) are yet available. Also unlike read-bytes!, read-bytes-avail! never drops bytes; if read-bytes-avail! successfully reads some bytes and then encounters an error, it suppresses the error (treating it roughly like an end-of-file) and returns the read bytes. (The error will be triggered by future reads.) If an error is encountered before any bytes have been read, an exception is raised.

When *in* produces a special value, as described in §13.1.9 "Custom Ports", the result is a procedure of four arguments. The four arguments correspond to the location of the special value within the port, as described in §13.1.9 "Custom Ports". If the procedure is called more than once with valid arguments, the exn:fail:contract exception is raised. If read-bytes-avail returns a special-producing procedure, then it does not place characters in *bstr*. Similarly, read-bytes-avail places only as many bytes into *bstr* as are available before a special value in the port's stream.

```
(read-bytes-avail!* bstr
                    [in
                     start-pos
                     end-pos])
 → (or/c exact-nonnegative-integer? eof-object? procedure?)
  bstr : bytes?
  in : input-port? = (current-input-port)
```

```
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like read-bytes-avail!, but returns 0 immediately if no bytes (or specials) are available for reading and the end-of-file is not reached.

```
(read-bytes-avail!/enable-break  bstr
                                 [in
                                  start-pos
                                  end-pos])
 → (or/c exact-positive-integer? eof-object? procedure?)
  bstr : bytes?
  in : input-port? = (current-input-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like read-bytes-avail!, but breaks are enabled during the read (see also §10.6 "Breaks"). If breaking is disabled when read-bytes-avail!/enable-break is called, and if the exn:break exception is raised as a result of the call, then no bytes will have been read from *in*.

```
(peek-string amt skip-bytes-amt [in]) → (or/c string? eof-object?)
  amt : exact-nonnegative-integer?
  skip-bytes-amt : exact-nonnegative-integer?
  in : input-port? = (current-input-port)
```

Similar to read-string, except that the returned characters are peeked: preserved in the port for future reads and peeks. (More precisely, undecoded bytes are left for future reads and peeks.) The *skip-bytes-amt* argument indicates a number of bytes (*not* characters) in the input stream to skip before collecting characters to return; thus, in total, the next *skip-bytes-amt* bytes plus *amt* characters are inspected.

For most kinds of ports, inspecting *skip-bytes-amt* bytes and *amt* characters requires at least *skip-bytes-amt*+*amt* bytes of memory overhead associated with the port, at least until the bytes/characters are read. No such overhead is required when peeking into a string port (see §13.1.6 "String Ports"), a pipe port (see §13.1.7 "Pipes"), or a custom port with a specific peek procedure (depending on how the peek procedure is implemented; see §13.1.9 "Custom Ports").

If a port produces eof mid-stream, attempts to skip beyond the eof for a peek always produce eof until the eof is read.

```
(peek-bytes amt skip-bytes-amt [in]) → (or/c bytes? eof-object?)
  amt : exact-nonnegative-integer?
  skip-bytes-amt : exact-nonnegative-integer?
  in : input-port? = (current-input-port)
```

Like peek-string, but peeks bytes and produces a byte string.

```
(peek-string!  str
               skip-bytes-amt
              [in
               start-pos
               end-pos])
 → (or/c exact-positive-integer? eof-object?)
  str : (and/c string? (not/c immutable?))
  skip-bytes-amt : exact-nonnegative-integer?
  in : input-port? = (current-input-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (string-length str)
```

Like read-string!, but for peeking, and with a *skip-bytes-amt* argument like peek-string.

```
(peek-bytes!  bstr
              skip-bytes-amt
             [in
              start-pos
              end-pos])
 → (or/c exact-positive-integer? eof-object?)
  bstr : (and/c bytes? (not/c immutable?))
  skip-bytes-amt : exact-nonnegative-integer?
  in : input-port? = (current-input-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like peek-string!, but peeks bytes, puts them into a byte string, and returns the number of bytes read.

```
(peek-bytes-avail!  bstr
                    skip-bytes-amt
                   [progress
                    in
                    start-pos
                    end-pos])
 → (or/c exact-nonnegative-integer? eof-object? procedure?)
  bstr : (and/c bytes? (not/c immutable?))
  skip-bytes-amt : exact-nonnegative-integer?
  progress : (or/c progress-evt? #f) = #f
  in : input-port? = (current-input-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like read-bytes-avail!, but for peeking, and with two extra arguments. The *skip-bytes-amt* argument is as in peek-bytes. The *progress* argument must be either #f or an event produced by port-progress-evt for *in*.

To peek, peek-bytes-avail! blocks until finding an end-of-file, at least one byte (or special) past the skipped bytes, or until a non-#f *progress* becomes ready. Furthermore, if *progress* is ready before bytes are peeked, no bytes are peeked or skipped, and *progress* may cut short the skipping process if it becomes available during the peek attempt. Furthermore, *progress* is checked even before determining whether the port is still open.

The result of peek-bytes-avail! is 0 only in the case that *progress* becomes ready before bytes are peeked.

```
(peek-bytes-avail!* bstr
                    skip-bytes-amt
                   [progress
                    in
                    start-pos
                    end-pos])
 → (or/c exact-nonnegative-integer? eof-object? procedure?)
  bstr : (and/c bytes? (not/c immutable?))
  skip-bytes-amt : exact-nonnegative-integer?
  progress : (or/c progress-evt? #f) = #f
  in : input-port? = (current-input-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like read-bytes-avail!*, but for peeking, and with *skip-bytes-amt* and *progress* arguments like peek-bytes-avail!. Since this procedure never blocks, it may return before even skip-amt bytes are available from the port.

```
(peek-bytes-avail!/enable-break bstr
                                skip-bytes-amt
                               [progress
                                in
                                start-pos
                                end-pos])
 → (or/c exact-nonnegative-integer? eof-object? procedure?)
  bstr : (and/c bytes? (not/c immutable?))
  skip-bytes-amt : exact-nonnegative-integer?
  progress : (or/c progress-evt? #f) = #f
  in : input-port? = (current-input-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like read-bytes-avail!/enable-break, but for peeking, and with *skip-bytes-amt* and *progress* arguments like peek-bytes-avail!.

```
(read-char-or-special [in]) → (or/c char? eof-object? any/c)
  in : input-port? = (current-input-port)
```

Like read-char, but if the input port returns a special value (through a value-generating procedure in a custom port; see §13.1.9 “Custom Ports” and §13.7.3 “Special Comments” for details), then the special value is returned.

```
(read-byte-or-special [in]) → (or/c byte? eof-object? any/c)
  in : input-port? = (current-input-port)
```

Like read-char-or-special, but reads and returns a byte instead of a character.

```
(peek-char [in skip-bytes-amt]) → (or/c char? eof-object?)
  in : input-port? = (current-input-port)
  skip-bytes-amt : exact-nonnegative-integer? = 0
```

Like read-char, but peeks instead of reading, and skips *skip-bytes-amt* bytes (not characters) at the start of the port.

```
(peek-byte [in skip-bytes-amt]) → (or/c byte? eof-object?)
  in : input-port? = (current-input-port)
  skip-bytes-amt : exact-nonnegative-integer? = 0
```

Like peek-char, but peeks and returns a byte instead of a character.

```
(peek-char-or-special [in skip-bytes-amt])
 → (or/c char? eof-object? any/c)
  in : input-port? = (current-input-port)
  skip-bytes-amt : exact-nonnegative-integer? = 0
```

Like peek-char, but if the input port returns a non-byte value after *skip-bytes-amt* byte positions, then it is returned.

```
(peek-byte-or-special [in
                       skip-bytes-amt
                       progress])
 → (or/c byte? eof-object? any/c)
  in : input-port? = (current-input-port)
  skip-bytes-amt : exact-nonnegative-integer? = 0
  progress : (or/c progress-evt? #f) = #f
```

Like peek-char-or-special, but peeks and returns a byte instead of a character, and it supports a *progress* argument like peek-bytes-avail!.

```
(port-progress-evt [in]) → progress-evt?
  in : (and/c input-port? port-provides-progress-evts?)
     = (current-input-port)
```

Returns a synchronizable event (see §11.2.1 “Events”) that becomes ready for synchronization after any subsequent read from *in* or after *in* is closed. After the event becomes ready, it remains ready. The synchronization result of a progress event is the progress event itself.

```
(port-provides-progress-evts? in) → boolean
  in : input-port?
```

Returns `#t` if `port-progress-evt` can return an event for *in*. All built-in kinds of ports support progress events, but ports created with `make-input-port` (see §13.1.9 “Custom Ports”) may not.

```
(port-commit-peeked amt progress evt [in]) → boolean?
  amt : exact-nonnegative-integer?
  progress : progress-evt?
  evt : evt?
  in : input-port? = (current-input-port)
```

Attempts to commit as read the first *amt* previously peeked bytes, non-byte specials, and `eof`s from *in*, or the first `eof` or special value peeked from *in*. Mid-stream `eof`s can be committed, but an `eof` when the port is exhausted does not necessarily commit, since it does not correspond to data in the stream.

The read commits only if *progress* does not become ready first (i.e., if no other process reads from *in* first), and only if *evt* is chosen by a `sync` within `port-commit-peeked` (in which case the event result is ignored); the *evt* must be either a channel-put event, channel, semaphore, semaphore-peek event, always event, or never event. Suspending the thread that calls `port-commit-peeked` may or may not prevent the commit from proceeding.

The result from `port-commit-peeked` is `#t` if data has been committed, and `#f` otherwise.

If no data has been peeked from *in* and *progress* is not ready, then `exn:fail:contract` exception is raised. If fewer than *amt* items have been peeked at the current start of *in*’s stream, then only the peeked items are committed as read. If *in*’s stream currently starts at an `eof` or a non-byte special value, then only the `eof` or special value is committed as read.

If *progress* is not a result of `port-progress-evt` applied to *in*, then `exn:fail:contract` exception is raised.

```
(byte-ready? [in]) → boolean?
  in : input-port? = (current-input-port)
```

Returns `#t` if `(read-byte in)` would not block (at the time that `byte-ready?` was called, at least). Equivalent to `(and (sync/timeout 0 in) #t)`.

```
(char-ready? [in]) → boolean?
  in : input-port? = (current-input-port)
```

Returns `#t` if `(read-char in)` would not block (at the time that `char-ready?` was called, at least). Depending on the initial bytes of the stream, multiple bytes may be needed to form a UTF-8 encoding.

```
(progress-evt? v) → boolean?
  v : any/c
(progress-evt? evt in) → boolean?
  evt : progress-evt?
  in : input-port?
```

With one argument, returns `#t` is *v* is a progress evt for some input port, `#f` otherwise.

With two arguments, returns `#t` if *evt* is a progress event for *in*, `#f` otherwise.

13.3 Byte and String Output

```
(write-char char [out]) → void?
  char : char?
  out : output-port? = (current-output-port)
```

Writes a single character to *out*; more precisely, the bytes that are the UTF-8 encoding of *char* are written to *out*.

```
(write-byte byte [out]) → void?
  byte : byte?
  out : output-port? = (current-output-port)
```

Writes a single byte to *out*.

```
(newline [out]) → void?
  out : output-port? = (current-output-port)
```

The same as `(write-char #\newline out)`.

```
(write-string str [out start-pos end-pos])
 → exact-nonnegative-integer?
  str : string?
  out : output-port? = (current-output-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (string-length str)
```

Writes characters to *out* from *str* starting from index *start-pos* (inclusive) up to *end-pos* (exclusive). Like substring, the exn:fail:contract exception is raised if *start-pos* or *end-pos* is out-of-range for *str*.

The result is the number of characters written to *out*, which is always (- *end-pos start-pos*).

```
(write-bytes bstr [out start-pos end-pos])
 → exact-nonnegative-integer?
  bstr : bytes?
  out : output-port? = (current-output-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like write-string, but writes bytes instead of characters.

```
(write-bytes-avail  bstr
                   [out
                    start-pos
                    end-pos]) → exact-nonnegative-integer?
  bstr : bytes?
  out : output-port? = (current-output-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like write-bytes, but returns without blocking after writing as many bytes as it can immediately flush. It blocks only if no bytes can be flushed immediately. The result is the number of bytes written and flushed to *out*; if *start-pos* is the same as *end-pos*, then the result can be 0 (indicating a successful flush of any buffered data), otherwise the result is between 1 and (- *end-pos start-pos*), inclusive.

The write-bytes-avail procedure never drops bytes; if write-bytes-avail successfully writes some bytes and then encounters an error, it suppresses the error and returns the number of written bytes. (The error will be triggered by future writes.) If an error is encountered before any bytes have been written, an exception is raised.

```
(write-bytes-avail*  bstr
                    [out
                     start-pos
                     end-pos])
 → (or/c exact-nonnegative-integer? #f)
  bstr : bytes?
  out : output-port? = (current-output-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like write-bytes-avail, but never blocks, returns #f if the port contains buffered data that cannot be written immediately, and returns 0 if the port's internal buffer (if any) is flushed but no additional bytes can be written immediately.

```
(write-bytes-avail/enable-break  bstr
                                 [out
                                  start-pos
                                  end-pos])
 → exact-nonnegative-integer?
  bstr : bytes?
  out : output-port? = (current-output-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like write-bytes-avail, except that breaks are enabled during the write. The procedure provides a guarantee about the interaction of writing and breaks: if breaking is disabled when write-bytes-avail/enable-break is called, and if the exn:break exception is raised as a result of the call, then no bytes will have been written to *out*. See also §10.6 "Breaks".

```
(write-special v [out]) → boolean?
  v : any/c
  out : output-port? = (current-output-port)
```

Writes *v* directly to *out* if the port supports special writes, or raises exn:fail:contract if the port does not support special write. The result is always #t, indicating that the write succeeded.

```
(write-special-avail* v [out]) → boolean?
  v : any/c
  out : output-port? = (current-output-port)
```

Like write-special, but without blocking. If *v* cannot be written immediately, the result is #f without writing *v*, otherwise the result is #t and *v* is written.

```
(write-bytes-avail-evt  bstr
                        [out
                         start-pos
                         end-pos]) → evt?
  bstr : bytes?
  out : output-port? = (current-output-port)
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Similar to write-bytes-avail, but instead of writing bytes immediately, it returns a synchronizable event (see §11.2.1 "Events"). The *out* must support atomic writes, as indicated by port-writes-atomic?.

Synchronizing on the object starts a write from *bstr*, and the event becomes ready when bytes are written (unbuffered) to the port. If *start-pos* and *end-pos* are the same, then the synchronization result is 0 when the port's internal buffer (if any) is flushed, otherwise the result is a positive exact integer. If the event is not selected in a synchronization, then no bytes will have been written to *out*.

```
(write-special-evt v [out]) → evt?
  v : any/c
  out : output-port? = (current-output-port)
```

Similar to write-special, but instead of writing the special value immediately, it returns a synchronizable event (see §11.2.1 "Events"). The *out* must support atomic writes, as indicated by port-writes-atomic?.

Synchronizing on the object starts a write of the special value, and the event becomes ready when the value is written (unbuffered) to the port. If the event is not selected in a synchronization, then no value will have been written to *out*.

```
(port-writes-atomic? out) → boolean?
  out : output-port?
```

Returns #t if write-bytes-avail/enable-break can provide an exclusive-or guarantee (break or write, but not both) for *out*, and if the port can be used with procedures like write-bytes-avail-evt. Racket's file-stream ports, pipes, string ports, and TCP ports all support atomic writes; ports created with make-output-port (see §13.1.9 "Custom Ports") may support atomic writes.

```
(port-writes-special? out) → boolean?
  out : output-port?
```

Returns #t if procedures like write-special can write arbitrary values to the port. Racket's file-stream ports, pipes, string ports, and TCP ports all reject special values, but ports created with make-output-port (see §13.1.9 "Custom Ports") may support them.

13.4 Reading

```
(read [in]) → any
  in : input-port? = (current-input-port)
```

Reads and returns a single datum from *in*. If *in* has a handler associated to it via port-read-handler, then the handler is called. Otherwise, the default reader is used, as parameterized by the current-readtable parameter, as well as many other parameters.

See §1.3 "The Reader" for information on the default reader.

```
(read-syntax [source-name in]) → (or/c syntax? eof-object?)
  source-name : any/c = (object-name in)
  in : input-port? = (current-input-port)
```

Like read, but produces a syntax object with source-location information. The *source-name* is used as the source field of the syntax object; it can be an arbitrary value, but it should generally be a path for the source file.

See §1.3 “The Reader” for information on the default reader in read-syntax mode.

```
(read/recursive [in start readtable graph?]) → any
  in : input-port? = (current-input-port)
  start : (or/c char? #f) = #f
  readtable : (or/c readtable? #f) = (current-readtable)
  graph? : any/c = #t
```

Similar to calling read, but normally used during the dynamic extent of read within a reader-extension procedure (see §13.7.2 “Reader-Extension Procedures”). The main effect of using read/recursive instead of read is that graph-structure annotations (see §1.3.17 “Reading Graph Structure”) in the nested read are considered part of the overall read, at least when the *graph?* argument is true; since the result is wrapped in a placeholder, however, it is not directly inspectable.

If *start* is provided and not #f, it is effectively prefixed to the beginning of *in*’s stream for the read. (To prefix multiple characters, use input-port-append.)

The *readtable* argument is used for top-level parsing to satisfy the read request; recursive parsing within the read (e.g., to read the elements of a list) instead uses the current readtable as determined by the current-readtable parameter. A reader macro might call read/recursive with a character and readtable to effectively invoke the readtable’s behavior for the character. If *readtable* is #f, the default readtable is used for top-level parsing.

When *graph?* is #f, graph structure annotations in the read datum are local to the datum.

When called within the dynamic extent of read, the read/recursive procedure produces either an opaque placeholder value, a special-comment value, or an end-of-file. The result is a special-comment value (see §13.7.3 “Special Comments”) when the input stream’s first non-whitespace content parses as a comment. The result is end-of-file when read/recursive encounters an end-of-file. Otherwise, the result is a placeholder that protects graph references that are not yet resolved. When this placeholder is returned within an S-expression that is produced by any reader-extension procedure (see §13.7.2 “Reader-Extension Procedures”) for the same outermost read, it will be replaced with the actual read value before the outermost read returns.

See §13.7.1 “Readtables” for an extended example that uses read/recursive.

```
(read-syntax/recursive [source-name
                        in
                        start
                        readtable
                        graph?])     → any
  source-name : any/c = (object-name in)
  in : input-port? = (current-input-port)
  start : (or/c char? #f) = #f
  readtable : (or/c readtable? #f) = (current-readtable)
  graph? : any/c = #t
```

Analogous to calling `read/recursive`, but the resulting value encapsulates S-expression structure with source-location information. As with `read/recursive`, when `read-syntax/recursive` is used within the dynamic extent of `read-syntax`, the result from `read-syntax/recursive` is either a special-comment value, end-of-file, or opaque graph-structure placeholder (not a syntax object). The placeholder can be embedded in an S-expression or syntax object returned by a reader macro, etc., and it will be replaced with the actual syntax object before the outermost `read-syntax` returns.

Using `read/recursive` within the dynamic extent of `read-syntax` does not allow graph structure for reading to be included in the outer `read-syntax` parsing, and neither does using `read-syntax/recursive` within the dynamic extent of `read`. In those cases, `read/recursive` and `read-syntax/recursive` produce results like `read` and `read-syntax`, except that a special-comment value is returned when the input stream starts with a comment (after whitespace).

See §13.7.1 “Readtables” for an extended example that uses `read-syntax/recursive`.

```
(read-language [in fail-thunk])
 → (or/c (any/c any/c . -> . any) #f)
  in : input-port? = (current-input-port)
  fail-thunk : (-> any) = (lambda () (error ...))
```

Reads from *in* in the same way as `read`, but stopping as soon as a reader language (or its absence) is determined.

A *reader language* is specified by `#lang` or `#!` (see §1.3.18 “Reading via an Extension”) at the beginning of the input, though possibly after comment forms. The default readtable is used by `read-language` (instead of the value of `current-readtable`), and `#reader` forms (which might produce comments) are not allowed before `#lang` or `#!`.

See also §17.3.5 “Source-Handling Configuration” in ***The Racket Guide***.

When it finds a `#lang` or `#!` specification, instead of dispatching to a `read` or `read-syntax` function as `read` and `read-syntax` do, `read-language` dispatches to the `get-info` function (if any) exported by the same module. The result of the `get-info` function is the result of `read-language` if it is a function of two arguments; if `get-info` produces any other

kind of result, the exn:fail:contract exception is raised. If no get-info function is exported, read-language returns #f.

The function produced by get-info reflects information about the expected syntax of the input stream. The first argument to the function serves as a key on such information; acceptable keys and the interpretation of results is up to external tools, such as DrRacket (see the DrRacket documentation). If no information is available for a given key, the result should be the second argument.

Examples:

```
> ((read-language (open-input-string "#lang algol60")) 'color-
lexer #f)
#<procedure:default-lexer>
> ((read-language (open-input-string "#lang algol60")) 'something-
else #f)
#f
```

The get-info function itself is applied to five arguments: the input port being read, the module path from which the get-info function was extracted, and the source line (positive exact integer or #f), column (non-negative exact integer or #f), and position (positive exact integer or #f) of the start of the #lang or #! form. The get-info function may further read from the given input port to determine its result, but it should read no further than necessary. The get-info function should not read from the port after returning a function.

If *in* starts with a reader language specification but the relevant module does not export get-info (but perhaps does export read and read-syntax), then the result of read-language is #f.

If *in* has a #lang or #! specification, but parsing and resolving the specification raises an exception, the exception is propagated by read-language. Having at least #l or #! (after comments and whitespace) counts as starting a #lang or #! specification.

If *in* does not specify a reader language with #lang or #!, then *fail-thunk* is called. The default *fail-thunk* raises exn:fail:read or exn:fail:read:eof.

```
(read-case-sensitive) → boolean?
(read-case-sensitive on?) → void?
  on? : any/c
```

A parameter that controls parsing and printing of symbols. When this parameter's value is #f, the reader case-folds symbols (e.g., producing 'hi when the input is any one of hi, Hi, HI, or hI). The parameter also affects the way that write prints symbols containing uppercase characters; if the parameter's value is #f, then symbols are printed with uppercase characters quoted by a \ or |. The parameter's value is overridden by quoting \ or | vertical-bar quotes and the #cs and #ci prefixes; see §1.3.2 "Reading Symbols" for more information. While a module is loaded, the parameter is set to #t (see current-load).

```
(read-square-bracket-as-paren) → boolean?
(read-square-bracket-as-paren on?) → void?
  on? : any/c
```

A parameter that controls whether `[` and `]` are treated as parentheses. See §1.3.6 “Reading Pairs and Lists” for more information.

```
(read-curly-brace-as-paren) → boolean?
(read-curly-brace-as-paren on?) → void?
  on? : any/c
```

A parameter that controls whether `{` and `}` are treated as parentheses. See §1.3.6 “Reading Pairs and Lists” for more information.

```
(read-accept-box) → boolean?
(read-accept-box on?) → void?
  on? : any/c
```

A parameter that controls parsing `#&` input. See §1.3.13 “Reading Boxes” for more information.

```
(read-accept-compiled) → boolean?
(read-accept-compiled on?) → void?
  on? : any/c
```

A parameter that controls parsing `#~` compiled input. See §1.3 “The Reader” and `current-compile` for more information.

```
(read-accept-bar-quote) → boolean?
(read-accept-bar-quote on?) → void?
  on? : any/c
```

A parameter that controls parsing and printing of `|` in symbols. See §1.3.2 “Reading Symbols” and §1.4 “The Printer” for more information.

```
(read-accept-graph) → boolean?
(read-accept-graph on?) → void?
  on? : any/c
```

A parameter value that controls parsing input with sharing. See §1.3.17 “Reading Graph Structure” for more information.

```
(read-decimal-as-inexact) → boolean?
(read-decimal-as-inexact on?) → void?
  on? : any/c
```

A parameter that controls parsing input numbers with a decimal point or exponent (but no explicit exactness tag). See §1.3.3 “Reading Numbers” for more information.

```
(read-accept-dot) → boolean?
(read-accept-dot on?) → void?
  on? : any/c
```

A parameter that controls parsing input with a dot, which is normally used for literal cons cells. See §1.3.6 “Reading Pairs and Lists” for more information.

```
(read-accept-infix-dot) → boolean?
(read-accept-infix-dot on?) → void?
  on? : any/c
```

A parameter that controls parsing input with two dots to trigger infix conversion. See §1.3.6 “Reading Pairs and Lists” for more information.

```
(read-accept-quasiquote) → boolean?
(read-accept-quasiquote on?) → void?
  on? : any/c
```

A parameter that controls parsing input with `` ` `` or `,`, which is normally used for `quasiquote`, `unquote`, and `unquote-splicing` abbreviations. See §1.3.8 “Reading Quotes” for more information.

```
(read-accept-reader) → boolean?
(read-accept-reader on?) → void?
  on? : any/c
```

A parameter that controls whether `#reader`, `#lang`, or `#!` are allowed for selecting a parser. See §1.3.18 “Reading via an Extension” for more information.

```
(read-accept-lang) → boolean?
(read-accept-lang on?) → void?
  on? : any/c
```

A parameter that (along with `read-accept-reader` controls whether `#lang` and `#!` are allowed for selecting a parser. See §1.3.18 “Reading via an Extension” for more information.

```
(current-readtable) → (or/c readtable? #f)
(current-readtable readtable) → void?
  readtable : (or/c readtable? #f)
```

A parameter whose value determines a readtable that adjusts the parsing of S-expression input, where `#f` implies the default behavior. See §13.7.1 “Readtables” for more information.

```
(call-with-default-reading-parameterization thunk) → any
  thunk : (-> any)
```

Calls *thunk* in tail position of a `parameterize` to set all reader parameters above to their default values.

Using the default parameter values ensures consistency, and it also provides safety when reading from untrusted sources, since the default values disable evaluation of arbitrary code via `#lang` or `#reader`.

```
(current-reader-guard) → (any/c . -> . any)
(current-reader-guard proc) → void?
  proc : (any/c . -> . any)
```

A parameter whose value converts or rejects (by raising an exception) a module-path datum following `#reader`. See §1.3.18 “Reading via an Extension” for more information.

```
(read-on-demand-source)
 → (or/c #f #t (and/c path? complete-path?))
(read-on-demand-source mode) → void?
  mode : (or/c #f #t (and/c path? complete-path?))
```

A parameter that enables lazy parsing of compiled code, so that closure bodies and syntax objects are extracted (and validated) from marshaled compiled code on demand. Normally, this parameter is set by the default load handler when `load-on-demand-enabled` is `#t`.

A `#f` value for `read-on-demand-source` disables lazy parsing of compiled code. A `#t` value enables lazy parsing. A path value furthers enable lazy retrieval from disk—instead of keeping unparsed compiled code in memory—when the `PLT_DELAY_FROM_ZO` environment variable is set (to any value) on start-up.

If the file at *mode* as a path changes before the delayed code is parsed when lazy retrieval from disk is enabled, then the on-demand parse most likely will encounter garbage, leading to an exception.

```
                          (case->
(port-read-handler in) →   (input-port? . -> . any)
                           (input-port?  any/c . -> . any))
  in : input-port?
(port-read-handler in proc) → void?
  in : input-port?
       (case->
  proc :  (input-port? . -> . any)
          (input-port? any/c . -> . any))
```

Gets or sets the *port read handler* for *in*. The handler called to read from the port when the built-in read or read-syntax procedure is applied to the port. (The port read handler is not used for read/recursive or read-syntax/recursive.)

A port read handler is applied to either one argument or two arguments:

- A single argument is supplied when the port is used with read; the argument is the port being read. The return value is the value that was read from the port (or end-of-file).
- Two arguments are supplied when the port is used with read-syntax; the first argument is the port being read, and the second argument is a value indicating the source. The return value is a syntax object that was read from the port (or end-of-file).

The default port read handler reads standard Racket expressions with Racket's built-in parser (see §1.3 "The Reader"). It handles a special result from a custom input port (see make-custom-input-port) by treating it as a single expression, except that special-comment values (see §13.7.3 "Special Comments") are treated as whitespace.

The default port read handler itself can be customized through a readtable; see §13.7.1 "Readtables" for more information.

13.5 Writing

```
(write datum [out]) → void?
  datum : any/c
  out : output-port? = (current-output-port)
```

Writes *datum* to *out*, normally in such a way that instances of core datatypes can be read back in. If *out* has a handler associated to it via port-write-handler, then the handler is called. Otherwise, the default printer is used (in write mode), as configured by various parameters.

See §1.4 "The Printer" for more information about the default printer. In particular, note that write may require memory proportional to the depth of the value being printed, due to the initial cycle check.

Examples:

```
> (write 'hi)
hi

> (write (lambda (n) n))
#<procedure>
```

```
> (define o (open-output-string))

> (write "hello" o)

> (get-output-string o)
"\"hello\""
```

```
(display datum [out]) → void?
  datum : any/c
  out : output-port? = (current-output-port)
```

Displays *datum* to *out*, similar to `write`, but usually in such a way that byte- and character-based datatypes are written as raw bytes or characters. If *out* has a handler associated to it via `port-display-handler`, then the handler is called. Otherwise, the default printer is used (in `display` mode), as configured by various parameters.

See §1.4 “The Printer” for more information about the default printer. In particular, note that `display` may require memory proportional to the depth of the value being printed, due to the initial cycle check.

```
(print datum [out quote-depth]) → void?
  datum : any/c
  out : output-port? = (current-output-port)
  quote-depth : (or/c 0 1) = 0
```

Writes *datum* to *out*, normally the same way as `write`. If *out* has a handler associated to it via `port-print-handler`, then the handler is called. Otherwise, the handler specified by `global-port-print-handler` is called; the default handler uses the default printer in `write` mode.

The optional *quote-depth* argument adjusts printing when the `print-as-expression` parameter is set to `#t`. In that case, *quote-depth* specifies the starting quote depth for printing *datum*.

The rationale for providing `print` is that `display` and `write` both have specific output conventions, and those conventions restrict the ways that an environment can change the behavior of `display` and `write` procedures. No output conventions should be assumed for `print`, so that environments are free to modify the actual output generated by `print` in any way.

```
(displayln datum [out]) → void?
  datum : any/c
  out : output-port? = (current-output-port)
```

The same as `(display datum out)` followed by `(newline out)`, which is similar to `println` in Pascal or Java.

```
(fprintf out form v ...) → void?
  out : output-port?
  form : string?
  v : any/c
```

Prints formatted output to *out*, where *form* is a string that is printed directly, except for special formatting escapes:

- ~n or ~% prints a newline character (which is equivalent to \n in a literal format string)
- ~a or ~A displays the next argument among the *v*s
- ~s or ~S writes the next argument among the *v*s
- ~v or ~V prints the next argument among the *v*s
- ~.⟨*c*⟩ where ⟨*c*⟩ is a, A, s, S, v, or V: truncates display, write, or print output to (error-print-width) characters, using ... as the last three characters if the untruncated output would be longer
- ~e or ~E outputs the next argument among the *v*s using the current error value conversion handler (see error-value->string-handler) and current error printing width
- ~c or ~C write-chars the next argument in *v*s; if the next argument is not a character, the exn:fail:contract exception is raised
- ~b or ~B prints the next argument among the *v*s in binary; if the next argument is not an exact number, the exn:fail:contract exception is raised
- ~o or ~O prints the next argument among the *v*s in octal; if the next argument is not an exact number, the exn:fail:contract exception is raised
- ~x or ~X prints the next argument among the *v*s in hexadecimal; if the next argument is not an exact number, the exn:fail:contract exception is raised
- ~~ prints a tilde.
- ~⟨*w*⟩, where ⟨*w*⟩ is a whitespace character (see char-whitespace?), skips characters in *form* until a non-whitespace character is encountered or until a second end-of-line is encountered (whichever happens first). On all platforms, an end-of-line can be #\return, #\newline, or #\return followed immediately by #\newline.

The *form* string must not contain any ~ that is not one of the above escapes, otherwise the exn:fail:contract exception is raised. When the format string requires more *v*s than are supplied, the exn:fail:contract exception is raised. Similarly, when more *v*s are supplied than are used by the format string, the exn:fail:contract exception is raised.

Example:

```
> (fprintf (current-output-port)
           "~a as a string is ~s.\n"
           '(3 4)
           "(3 4)")
(3 4) as a string is "(3 4)".
```

```
(printf form v ...) → void?
  form : string?
  v : any/c
```

The same as `(fprintf (current-output-port) form v ...)`.

```
(eprintf form v ...) → void?
  form : string?
  v : any/c
```

The same as `(fprintf (current-error-port) form v ...)`.

```
(format form v ...) → string?
  form : string?
  v : any/c
```

Formats to a string. The result is the same as

```
(let ([o (open-output-string)])
  (fprintf o form v ...)
  (get-output-string o))
```

Example:

```
> (format "~a as a string is ~s.\n" '(3 4) "(3 4)")
"(3 4) as a string is \"(3 4)\".\n"
```

```
(print-pair-curly-braces) → boolean?
(print-pair-curly-braces on?) → void?
  on? : any/c
```

A parameter that controls pair printing. If the value is true, then pairs print using `{` and `}` instead of `(` and `)`. The default is `#f`.

```
(print-mpair-curly-braces) → boolean?
(print-mpair-curly-braces on?) → void?
  on? : any/c
```

A parameter that controls pair printing. If the value is true, then mutable pairs print using { and } instead of (and). The default is #t.

```
(print-unreadable) → boolean?
(print-unreadable on?) → void?
  on? : any/c
```

A parameter that enables or disables printing of values that have no readable form (using the default reader), including structures that have a custom-write procedure (see prop:custom-write), but not including uninterned symbols and unreadable symbols (which print the same as interned symbols). If the parameter value is #f, an attempt to print an unreadable value raises exn:fail. The parameter value defaults to #t. See §1.4 “The Printer” for more information.

```
(print-graph) → boolean?
(print-graph on?) → void?
  on? : any/c
```

A parameter that controls printing data with sharing; defaults to #f. See §1.4 “The Printer” for more information.

```
(print-struct) → boolean?
(print-struct on?) → void?
  on? : any/c
```

A parameter that controls printing structure values in vector or prefab form; defaults to #t. See §1.4 “The Printer” for more information. This parameter has no effect on the printing of structures that have a custom-write procedure (see prop:custom-write).

```
(print-box) → boolean?
(print-box on?) → void?
  on? : any/c
```

A parameter that controls printing box values; defaults to #t. See §1.4.10 “Printing Boxes” for more information.

```
(print-vector-length) → boolean?
(print-vector-length on?) → void?
  on? : any/c
```

A parameter that controls printing vectors; defaults to #f. See §1.4.7 “Printing Vectors” for more information.

```
(print-hash-table) → boolean?
(print-hash-table on?) → void?
  on? : any/c
```

A parameter that controls printing hash tables; defaults to `#f`. See §1.4.9 “Printing Hash Tables” for more information.

```
(print-boolean-long-form) → boolean?
(print-boolean-long-form on?) → void?
  on? : any/c
```

A parameter that controls printing of booleans. When the parameter’s value is true, `#t` and `#f` print as `#true` and `#false`, otherwise they print as `#t` and `#f`. The default is `#f`.

```
(print-reader-abbreviations) → boolean?
(print-reader-abbreviations on?) → void?
  on? : any/c
```

A parameter that controls printing of two-element lists that start with `quote`, `'quasiquote`, `'unquote`, `'unquote-splicing`, `'syntax`, `'quasisyntax`, `'unsyntax`, or `'unsyntax-splicing`; defaults to `#f`. See §1.4.5 “Printing Pairs and Lists” for more information.

```
(print-as-expression) → boolean?
(print-as-expression on?) → void?
  on? : any/c
```

A parameter that controls printing in `print` mode (as opposed to `write` or `display`); defaults to `#t`. See §1.4 “The Printer” for more information.

```
(print-syntax-width)
 → (or/c +inf.0 0 (and/c exact-integer? (>/c 3)))
(print-syntax-width width) → void?
  width : (or/c +inf.0 0 (and/c exact-integer? (>/c 3)))
```

A parameter that controls printing of syntax objects. Up to *width* characters are used to show the datum form of a syntax object within `#<syntax...>` (after the syntax object’s source location, if any).

```
(current-write-relative-directory)
    (or/c (and/c path? complete-path?)
          (cons/c (and/c path? complete-path?)
                  (and/c path? complete-path?))
 →        #f)
(current-write-relative-directory path) → void?
         (or/c (and/c path-string? complete-path?)
               (cons/c (and/c path-string? complete-path?)
                       (and/c path-string? complete-path?))
  path :       #f)
```

A parameter that is used when writing compiled code (see §1.4.16 “Printing Compiled Code”) that contains pathname literals, including source-location pathnames for procedure names. When the parameter’s value is a *path*, paths that syntactically extend *path* are converted to relative paths; when the resulting compiled code is read, relative paths are converted back to complete paths using the current-load-relative-directory parameter (if it is not #f; otherwise, the path is left relative). When the parameter’s value is (cons *rel-to-path base-path*), then paths that syntactically extend *base-path* are converted as relative to *rel-to-path*; the *rel-to-path* must extend *base-path*, in which case 'up path elements (in the sense of build-path) may be used to make a path relative to *rel-to-path*.

```
(port-write-handler out) → (any/c output-port? . -> . any)
  out : output-port?
(port-write-handler out proc) → void?
  out : output-port?
  proc : (any/c output-port? . -> . any)
(port-display-handler out) → (any/c output-port? . -> . any)
  out : output-port?
(port-display-handler out proc) → void?
  out : output-port?
  proc : (any/c output-port? . -> . any)
(port-print-handler out)
 → ((any/c output-port?) ((or/c 0 1)) . ->* . any)
  out : output-port?
(port-print-handler out proc) → void?
  out : output-port?
  proc : (any/c output-port? . -> . any)
```

Gets or sets the *port write handler*, *port display handler*, or *port print handler* for *out*. This handler is called to output to the port when write, display, or print (respectively) is applied to the port. Each handler must accept two arguments: the value to be printed and the destination port. The handler’s return value is ignored.

A port print handler optionally accepts a third argument, which corresponds to the optional third argument to print; if a procedure given to port-print-handler does not accept a third argument, it is wrapped with a procedure that discards the optional third argument.

The default port display and write handlers print Racket expressions with Racket’s built-in printer (see §1.4 “The Printer”). The default print handler calls the global port print handler (the value of the global-port-print-handler parameter); the default global port print handler is the same as the default write handler.

```
(global-port-print-handler)
 → (->* (any/c output-port?) ((or/c 0 1)) any)
(global-port-print-handler proc) → void?
  proc : (or/c (->* (any/c output-port?) ((or/c 0 1)) any)
               (any/c output-port? . -> . any))
```

A parameter that determines *global port print handler*, which is called by the default port print handler (see port-print-handler) to print values into a port. The default value uses the built-in printer (see §1.4 "The Printer") in print mode.

A global port print handler optionally accepts a third argument, which corresponds to the optional third argument to print. If a procedure given to global-port-print-handler does not accept a third argument, it is wrapped with a procedure that discards the optional third argument.

13.6 Pretty Printing

(require racket/pretty) package: base

The bindings documented in this section are provided by the racket/pretty and racket libraries, but not racket/base.

```
(pretty-print v [port quote-depth]) → void?
  v : any/c
  port : output-port? = (current-output-port)
  quote-depth : (or/c 0 1) = 0
```

Pretty-prints the value *v* using the same printed form as the default print mode, but with newlines and whitespace inserted to avoid lines longer than (pretty-print-columns), as controlled by (pretty-print-current-style-table). The printed form ends in a newline, unless the pretty-print-columns parameter is set to 'infinity. When *port* has line counting enabled (see §13.1.4 "Counting Positions, Lines, and Columns"), then printing is sensitive to the column when printing starts—both for determining an initial line break and indenting subsequent lines.

In addition to the parameters defined in this section, pretty-print conforms to the print-graph, print-struct, print-hash-table, print-vector-length, print-box, and print-as-expression parameters.

The pretty printer detects structures that have the prop:custom-write property and calls the corresponding custom-write procedure. The custom-write procedure can check the parameter pretty-printing to cooperate with the pretty-printer. Recursive printing to the port automatically uses pretty printing, but if the structure has multiple recursively printed sub-expressions, a custom-write procedure may need to cooperate more to insert explicit newlines. Use port-next-location to determine the current output column, use pretty-print-columns to determine the target printing width, and use pretty-print-newline to insert a newline (so that the function in the pretty-print-print-line parameter can be called appropriately). Use make-tentative-pretty-print-output-port to obtain a port for tentative recursive prints (e.g., to check the length of the output).

```
(pretty-write v [port]) → void?
```

```
  v : any/c
  port : output-port? = (current-output-port)
```

Same as pretty-print, but v is printed like write instead of like print.

```
(pretty-display v [port]) → void?
  v : any/c
  port : output-port? = (current-output-port)
```

Same as pretty-print, but v is printed like display instead of like print.

```
(pretty-format v [columns]) → string?
  v : any/c
  columns : exact-nonnegative-integer? = (pretty-print-columns)
```

Like pretty-print, except that it returns a string containing the pretty-printed value, rather than sending the output to a port.

The optional argument columns argument is used to parameterize pretty-print-columns.

```
(pretty-print-handler v) → void?
  v : any/c
```

Pretty-prints v if v is not #<void>, or prints nothing if v is #<void>. Pass this procedure to current-print to install the pretty printer into the REPL run by read-eval-print-loop.

13.6.1 Basic Pretty-Print Options

```
(pretty-print-columns)
 → (or/c exact-positive-integer? 'infinity)
(pretty-print-columns width) → void?
  width : (or/c exact-positive-integer? 'infinity)
```

A parameter that determines the default width for pretty printing.

If the display width is 'infinity, then pretty-printed output is never broken into lines, and a newline is not added to the end of the output.

```
(pretty-print-depth) → (or/c exact-nonnegative-integer? #f)
(pretty-print-depth depth) → void?
  depth : (or/c exact-nonnegative-integer? #f)
```

Parameter that controls the default depth for recursive pretty printing. Printing to *depth* means that elements nested more deeply than *depth* are replaced with "..."; in particular, a depth of 0 indicates that only simple values are printed. A depth of #f (the default) allows printing to arbitrary depths.

```
(pretty-print-exact-as-decimal) → boolean?
(pretty-print-exact-as-decimal as-decimal?) → void?
  as-decimal? : any/c
```

A parameter that determines how exact non-integers are printed. If the parameter's value is #t, then an exact non-integer with a decimal representation is printed as a decimal number instead of a fraction. The initial value is #f.

```
(pretty-print-.-symbol-without-bars) → boolean?
(pretty-print-.-symbol-without-bars on?) → void?
  on? : any/c
```

A parameter that controls the printing of the symbol whose print name is just a period. If set to a true value, then such a symbol is printed as only the period. If set to a false value, it is printed as a period with vertical bars surrounding it.

```
(pretty-print-show-inexactness) → boolean?
(pretty-print-show-inexactness show?) → void?
  show? : any/c
```

A parameter that determines how inexact numbers are printed. If the parameter's value is #t, then inexact numbers are always printed with a leading #i. The initial value is #f.

13.6.2 Per-Symbol Special Printing

```
(pretty-print-abbreviate-read-macros) → boolean?
(pretty-print-abbreviate-read-macros abbrev?) → void?
  abbrev? : any/c
```

A parameter that controls whether or not quote, unquote, unquote-splicing, etc., are abbreviated with ', ,, ,@, etc. By default, the abbreviations are enabled.

See also pretty-print-remap-stylable.

```
(pretty-print-style-table? v) → boolean?
  v : any/c
```

Returns #t if *v* is a style table for use with pretty-print-current-style-table, #f otherwise.

```
(pretty-print-current-style-table) → pretty-print-style-table?
(pretty-print-current-style-table style-table) → void?
  style-table : pretty-print-style-table?
```

A parameter that holds a table of style mappings. See pretty-print-extend-style-table.

```
(pretty-print-extend-style-table style-table
                                 symbol-list
                                 like-symbol-list)
 → pretty-print-style-table?
  style-table : pretty-print-style-table?
  symbol-list : (listof symbol?)
  like-symbol-list : (listof symbol?)
```

Creates a new style table by extending an existing *style-table*, so that the style mapping for each symbol of *like-symbol-list* in the original table is used for the corresponding symbol of *symbol-list* in the new table. The *symbol-list* and *like-symbol-list* lists must have the same length. The *style-table* argument can be #f, in which case the default mappings are used from the original table (see below).

The style mapping for a symbol controls the way that whitespace is inserted when printing a list that starts with the symbol. In the absence of any mapping, when a list is broken across multiple lines, each element of the list is printed on its own line, each with the same indentation.

The default style mapping includes mappings for the following symbols, so that the output follows popular code-formatting rules:

```
'lambda 'case-lambda
'define 'define-macro 'define-syntax
'let 'letrec 'let*
'let-syntax 'letrec-syntax
'let-values 'letrec-values 'let*-values
'let-syntaxes 'letrec-syntaxes
'begin 'begin0 'do
'if 'set! 'set!-values
'unless 'when
'cond 'case 'and 'or
'module
'syntax-rules 'syntax-case 'letrec-syntaxes+values
'import 'export 'link
'require 'require-for-syntax 'require-for-template 'provide
'public 'private 'override 'rename 'inherit 'field 'init
'shared 'send 'class 'instantiate 'make-object
```

```
(pretty-print-remap-stylable)
 → (any/c . -> . (or/c symbol? #f))
(pretty-print-remap-stylable proc) → void?
  proc : (any/c . -> . (or/c symbol? #f))
```

A parameter that controls remapping for styles and for the determination of the reader shorthands.

This procedure is called with each sub-expression that appears as the first element in a sequence. If it returns a symbol, the style table is used, as if that symbol were at the head of the sequence. If it returns `#f`, the style table is treated normally. Similarly, when determining whether to abbreviate reader macros, this parameter is consulted.

13.6.3 Line-Output Hook

```
(pretty-print-newline port width) → void?
  port : output-port?
  width : exact-nonnegative-integer?
```

Calls the procedure associated with the `pretty-print-print-line` parameter to print a newline to *port*, if *port* is the output port that is redirected to the original output port for printing, otherwise a plain newline is printed to *port*. The *width* argument should be the target column width, typically obtained from `pretty-print-columns`.

```
(pretty-print-print-line)
    ((or/c exact-nonnegative-integer? #f)
     output-port?
     exact-nonnegative-integer?
 →   (or/c exact-nonnegative-integer? 'infinity)
     . -> .
     exact-nonnegative-integer?)
(pretty-print-print-line proc) → void?
         ((or/c exact-nonnegative-integer? #f)
          output-port?
          exact-nonnegative-integer?
  proc :  (or/c exact-nonnegative-integer? 'infinity)
          . -> .
          exact-nonnegative-integer?)
```

A parameter that determines a procedure for printing the newline separator between lines of a pretty-printed value. The procedure is called with four arguments: a new line number, an output port, the old line's length, and the number of destination columns. The return value from *proc* is the number of extra characters it printed at the beginning of the new line.

The *proc* procedure is called before any characters are printed with 0 as the line number and 0 as the old line length; *proc* is called after the last character of a value has been printed with #f as the line number and with the length of the last line. Whenever the pretty-printer starts a new line, *proc* is called with the new line's number (where the first new line is numbered 1) and the just-finished line's length. The destination-columns argument to *proc* is always the total width of the destination printing area, or 'infinity if pretty-printed values are not broken into lines.

The default *proc* procedure prints a newline whenever the line number is not 0 and the column count is not 'infinity, always returning 0. A custom *proc* procedure can be used to print extra text before each line of pretty-printed output; the number of characters printed before each line should be returned by *proc* so that the next line break can be chosen correctly.

The destination port supplied to *proc* is generally not the port supplied to pretty-print or pretty-display (or the current output port), but output to this port is ultimately redirected to the port supplied to pretty-print or pretty-display.

13.6.4 Value Output Hook

```
(pretty-print-size-hook)
    (any/c boolean? output-port?
 →  . -> .
     (or/c #f exact-nonnegative-integer?))
(pretty-print-size-hook proc) → void?
         (any/c boolean? output-port?
  proc :  . -> .
          (or/c #f exact-nonnegative-integer?))
```

A parameter that determines a sizing hook for pretty-printing.

The sizing hook is applied to each value to be printed. If the hook returns #f, then printing is handled internally by the pretty-printer. Otherwise, the value should be an integer specifying the length of the printed value in characters; the print hook will be called to actually print the value (see pretty-print-print-hook).

The sizing hook receives three arguments. The first argument is the value to print. The second argument is a boolean: #t for printing like display and #f for printing like write. The third argument is the destination port; the port is the one supplied to pretty-print or pretty-display (or the current output port). The sizing hook may be applied to a single value multiple times during pretty-printing.

```
(pretty-print-print-hook)
 → (any/c boolean? output-port? . -> . void?)
(pretty-print-print-hook proc) → void?
  proc : (any/c boolean? output-port? . -> . void?)
```

A parameter that determines a print hook for pretty-printing. The print-hook procedure is applied to a value for printing when the sizing hook (see pretty-print-size-hook) returns an integer size for the value.

The print hook receives three arguments. The first argument is the value to print. The second argument is a boolean: #t for printing like display and #f for printing like write. The third argument is the destination port; this port is generally not the port supplied to pretty-print or pretty-display (or the current output port), but output to this port is ultimately redirected to the port supplied to pretty-print or pretty-display.

```
(pretty-print-pre-print-hook)
 → (any/c output-port? . -> . void)
(pretty-print-pre-print-hook proc) → void?
  proc : (any/c output-port? . -> . void)
```

A parameter that determines a hook procedure to be called just before an object is printed. The hook receives two arguments: the object and the output port. The port is the one supplied to pretty-print or pretty-display (or the current output port).

```
(pretty-print-post-print-hook)
 → (any/c output-port? . -> . void)
(pretty-print-post-print-hook proc) → void?
  proc : (any/c output-port? . -> . void)
```

A parameter that determines a hook procedure to be called just after an object is printed. The hook receives two arguments: the object and the output port. The port is the one supplied to pretty-print or pretty-display (or the current output port).

13.6.5 Additional Custom-Output Support

```
(pretty-printing) → boolean?
(pretty-printing on?) → void?
  on? : any/c
```

A parameter that is set to #t when the pretty printer calls a custom-write procedure (see prop:custom-write) for output in a mode that supports line breaks. When pretty printer calls a custom-write procedure merely to detect cycles or to try to print on a single line, it sets this parameter to #f.

```
(make-tentative-pretty-print-output-port out
                                        width
                                        overflow-thunk)
 → output-port?
  out : output-port?
  width : exact-nonnegative-integer?
  overflow-thunk : (-> any)
```

Produces an output port that is suitable for recursive pretty printing without actually producing output. Use such a port to tentatively print when proper output depends on the size of recursive prints. After printing, determine the size of the tentative output using `file-position`.

The *out* argument should be a pretty-printing port, such as the one supplied to a custom-write procedure when `pretty-printing` is set to true, or another tentative output port. The *width* argument should be a target column width, usually obtained from `pretty-print-columns`, possibly decremented to leave room for a terminator. The *overflow-thunk* procedure is called if more than *width* items are printed to the port or if a newline is printed to the port via `pretty-print-newline`; it can escape from the recursive print through a continuation as a shortcut, but *overflow-thunk* can also return, in which case it is called every time afterward that additional output is written to the port.

After tentative printing, either accept the result with `tentative-pretty-print-port-transfer` or reject it with `tentative-pretty-print-port-cancel`. Failure to accept or cancel properly interferes with graph-structure printing, calls to hook procedures, etc. Explicitly cancel the tentative print even when *overflow-thunk* escapes from a recursive print.

```
(tentative-pretty-print-port-transfer tentative-out
                                      orig-out)     → void?
  tentative-out : output-port?
  orig-out : output-port?
```

Causes the data written to *tentative-out* to be transferred as if written to *orig-out*. The *tentative-out* argument should be a port produced by `make-tentative-pretty-print-output-port`, and *orig-out* should be either a pretty-printing port (provided to a custom-write procedure) or another tentative output port.

```
(tentative-pretty-print-port-cancel tentative-out) → void?
  tentative-out : output-port?
```

Cancels the content of *tentative-out*, which was produced by `make-tentative-pretty-print-output-port`. The main effect of canceling is that graph-reference definitions are undone, so that a future print of a graph-referenced object includes the defining `#`⟨*n*⟩`=`.

13.7 Reader Extension

Racket's reader can be extended in three ways: through a reader-macro procedure in a readtable (see §13.7.1 "Readtables"), through a `#reader` form (see §1.3.18 "Reading via an Extension"), or through a custom-port byte reader that returns a "special" result procedure (see §13.1.9 "Custom Ports"). All three kinds of *reader extension procedures* accept

similar arguments, and their results are treated in the same way by `read` and `read-syntax` (or, more precisely, by the default read handler; see `port-read-handler`).

13.7.1 Readtables

The dispatch table in §1.3.1 “Delimiters and Dispatch” corresponds to the default *readtable*. By creating a new readtable and installing it via the `current-readtable` parameter, the reader’s behavior can be extended.

A readtable is consulted at specific times by the reader:

- when looking for the start of a datum;
- when determining how to parse a datum that starts with `#`;
- when looking for a delimiter to terminate a symbol or number;
- when looking for an opener (such as `(`), closer (such as `)`), or `.` after the first character parsed as a sequence for a pair, list, vector, or hash table; or
- when looking for an opener after `#`⟨*n*⟩ in a vector of specified length ⟨*n*⟩.

The readtable is ignored at other times. In particular, after parsing a character that is mapped to the default behavior of `;`, the readtable is ignored until the comment’s terminating newline is discovered. Similarly, the readtable does not affect string parsing until a closing double-quote is found. Meanwhile, if a character is mapped to the default behavior of `(`, then it starts sequence that is closed by any character that is mapped to a closing parenthesis `)`. An apparent exception is that the default parsing of `|` quotes a symbol until a matching character is found, but the parser is simply using the character that started the quote; it does not consult the readtable.

For many contexts, `#f` identifies the default readtable. In particular, `#f` is the initial value for the `current-readtable` parameter, which causes the reader to behave as described in §1.3 “The Reader”.

```
(readtable? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a readtable, `#f` otherwise.

```
(make-readtable readtable
                key
                mode
                action ...+) → readtable?
  readtable : readtable?
```

```
key : (or/c char? #f)
      (or/c (or/c 'terminating-macro
                  'non-terminating-macro
mode :            'dispatch-macro)
            char?)
        (or/c procedure?
action :      readtable?)
```

Creates a new readtable that is like *readtable* (which can be `#f`), except that the reader's behavior is modified for each *key* according to the given *mode* and *action*. The `...+` for `make-readtable` applies to all three of *key*, *mode*, and *action*; in other words, the total number of arguments to `make-readtable` must be 1 modulo 3.

The possible combinations for *key*, *mode*, and *action* are as follows:

- *char* `'terminating-macro` *proc* — causes *char* to be parsed as a delimiter, and an unquoted/uncommented *char* in the input string triggers a call to the *reader macro* *proc*; the activity of *proc* is described further below. Conceptually, characters like `;`, `(`, and `)` are mapped to terminating reader macros in the default readtable.

- *char* `'non-terminating-macro` *proc* — like the `'terminating-macro` variant, but *char* is not treated as a delimiter, so it can be used in the middle of an identifier or number. Conceptually, `#` is mapped to a non-terminating macro in the default readtable.

- *char* `'dispatch-macro` *proc* — like the `'non-terminating-macro` variant, but for *char* only when it follows a `#` (or, more precisely, when the character follows one that has been mapped to the behavior of `#` in the default readtable).

- *char like-char readtable* — causes *char* to be parsed in the same way that *like-char* is parsed in *readtable*, where *readtable* can be `#f` to indicate the default readtable. Mapping a character to the same actions as `|` in the default reader means that the character starts quoting for symbols, and the same character terminates the quote; in contrast, mapping a character to the same action as a `"` means that the character starts a string, but the string is still terminated with a closing `"`. Finally, mapping a character to an action in the default readtable means that the character's behavior is sensitive to parameters that affect the original character; for example, mapping a character to the same action as a curly brace `{` in the default readtable means that the character is disallowed when the `read-curly-brace-as-paren` parameter is set to `#f`.

- `#f` `'non-terminating-macro` *proc* — replaces the macro used to parse characters with no specific mapping: i.e., the characters (other than `#` or `|`) that can start a symbol or number with the default readtable.

If multiple `'dispatch-macro` mappings are provided for a single *char*, all but the last

one are ignored. Similarly, if multiple non-'dispatch-macro mappings are provided for a single *char*, all but the last one are ignored.

A reader macro *proc* must accept six arguments, and it can optionally accept two arguments. The first two arguments are always the character that triggered the reader macro and the input port for reading. When the reader macro is triggered by read-syntax (or read-syntax/recursive), the procedure is passed four additional arguments that represent a source location for already-consumed character(s): the source name, a line number or #f, a column number or #f, and a position or #f. When the reader macro is triggered by read (or read/recursive), the procedure is passed only two arguments if it accepts two arguments, otherwise it is passed six arguments where the last four are all #f. See §13.7.2 "Reader-Extension Procedures" for information on the procedure's results.

A reader macro normally reads characters from the given input port to produce a value to be used as the "reader macro-expansion" of the consumed characters. The reader macro might produce a special-comment value (see §13.7.3 "Special Comments") to cause the consumed character to be treated as whitespace, and it might use read/recursive or read-syntax/recursive.

```
(readtable-mapping readtable char)
    (or/c char?
          (or/c 'terminating-macro
 →              'non-terminating-macro))
    (or/c #f procedure?)
    (or/c #f procedure?)
  readtable : readtable?
  char : char?
```

Produces information about the mappings in *readtable* for *char*. The result is three values:

- either a character (mapping to the same behavior as the character in the default readtable), 'terminating-macro, or 'non-terminating-macro; this result reports the main (i.e., non-'dispatch-macro) mapping for key. When the result is a character, then key is mapped to the same behavior as the returned character in the default readtable.
- either #f or a reader-macro procedure; the result is a procedure when the first result is 'terminating-macro or 'non-terminating-macro.
- either #f or a reader-macro procedure; the result is a procedure when the character has a 'dispatch-macro mapping in *readtable* to override the default dispatch behavior.

Note that reader-macro procedures for the default readtable are not directly accessible. To invoke default behaviors, use read/recursive or read-syntax/recursive with a character and the #f readtable.

Examples:

```
; Provides raise-read-error and raise-read-eof-error
> (require syntax/readerr)

> (define (skip-whitespace port)
    ; Skips whitespace characters, sensitive to the current
    ; readtable's definition of whitespace
    (let ([ch (peek-char port)])
      (unless (eof-object? ch)
        ; Consult current readtable:
        (let-values ([(like-ch/sym proc dispatch-proc)
                      (readtable-mapping (current-readtable) ch)])
          ; If like-ch/sym is whitespace, then ch is whitespace
          (when (and (char? like-ch/sym)
                     (char-whitespace? like-ch/sym))
            (read-char port)
            (skip-whitespace port))))))

> (define (skip-comments read-one port src)
    ; Recursive read, but skip comments and detect EOF
    (let loop ()
      (let ([v (read-one)])
        (cond
         [(special-comment? v) (loop)]
         [(eof-object? v)
          (let-values ([(l c p) (port-next-location port)])
            (raise-read-eof-error
             "unexpected EOF in tuple" src l c p 1))]
         [else v]))))

> (define (parse port read-one src)
    ; First, check for empty tuple
    (skip-whitespace port)
    (if (eq? #\> (peek-char port))
        null
        (let ([elem (read-one)])
          (if (special-comment? elem)
              ; Found a comment, so look for > again
              (parse port read-one src)
              ; Non-empty tuple:
              (cons elem
                    (parse-nonempty port read-one src))))))

> (define (parse-nonempty port read-one src)
    ; Need a comma or closer
```

```
    (skip-whitespace port)
    (case (peek-char port)
      [(#\>) (read-char port)
       ; Done
       null]
      [(#\,) (read-char port)
       ; Read next element and recur
       (cons (skip-comments read-one port src)
             (parse-nonempty port read-one src))]
      [else
       ; Either a comment or an error; grab location (in case
       ; of error) and read recursively to detect comments
       (let-values ([(l c p) (port-next-location port)]
                    [(v) (read-one)])
         (cond
          [(special-comment? v)
           ; It was a comment, so try again
           (parse-nonempty port read-one src)]
          [else
           ; Wasn't a comment, comma, or closer; error
           ((if (eof-object? v)
                raise-read-eof-error
                raise-read-error)
            "expected `,' or `>'" src l c p 1)]))]))

> (define (make-delims-table)
    ; Table to use for recursive reads to disallow delimiters
    ;  (except those in sub-expressions)
    (letrec ([misplaced-delimiter
              (case-lambda
               [(ch port) (misplaced-delimiter ch port #f #f #f #f)]
               [(ch port src line col pos)
                (raise-read-error
                 (format "misplaced `~a' in tuple" ch)
                 src line col pos 1)])])
      (make-readtable (current-readtable)
                      #\, 'terminating-macro misplaced-delimiter
                      #\> 'terminating-macro misplaced-
delimiter)))

> (define (wrap l)
    `(make-tuple (list ,@l)))

> (define parse-open-tuple
    (case-lambda
     [(ch port)
```

```
      ; `read' mode
      (wrap (parse port
                   (lambda ()
                     (read/recursive port #f
                                     (make-delims-table)))
                   (object-name port)))]
     [(ch port src line col pos)
      ; `read-syntax' mode
      (datum->syntax
       #f
       (wrap (parse port
                    (lambda ()
                      (read-syntax/recursive src port #f
                                             (make-delims-table)))
                    src))
       (let-values ([(l c p) (port-next-location port)])
         (list src line col pos (and pos (- p pos))))))]))

> (define tuple-readtable
    (make-readtable #f #\< 'terminating-macro parse-open-tuple))

> (parameterize ([current-readtable tuple-readtable])
    (read (open-input-string "<1 , 2 , \"a\">")))
'(make-tuple (list 1 2 "a"))
> (parameterize ([current-readtable tuple-readtable])
    (read (open-input-string
           "< #|# 1 #|# , #|# 2 #|# , #|# \"a\" #|# >")))
'(make-tuple (list 1 2 "a"))
> (define tuple-readtable+
    (make-readtable tuple-readtable
                    #\* 'terminating-macro (lambda a
                                             (make-special-
comment #f))
                    #\_ #\space #f))

> (parameterize ([current-readtable tuple-readtable+])
    (read (open-input-string "< * 1 __,__  2 __,__ * \"a\" * >")))
'(make-tuple (list 1 2 "a"))
```

13.7.2 Reader-Extension Procedures

Calls to reader extension procedures can be triggered through `read`, `read/recursive`, or `read-syntax`. In addition, a special-read procedure can be triggered by calls to `read-char-or-special`, or by the context of `read-bytes-avail!`, `peek-bytes-avail!*`, `read-bytes-avail!`, and `peek-bytes-avail!*`.

Optional arities for reader-macro and special-result procedures allow them to distinguish reads via `read`, etc., from reads via `read-syntax`, etc. (where the source value is `#f` and no other location information is available).

When a reader-extension procedure is called in syntax-reading mode (via `read-syntax`, etc.), it should generally return a syntax object that has no lexical context (e.g., a syntax object created using `datum->syntax` with `#f` as the first argument and with the given location information as the third argument). Another possible result is a special-comment value (see §13.7.3 "Special Comments"). If the procedure's result is not a syntax object and not a special-comment value, it is converted to one using `datum->syntax`.

When a reader-extension procedure is called in non-syntax-reading modes, it should generally not return a syntax object. If a syntax object is returned, it is converted to a plain value using `syntax->datum`.

In either context, when the result from a reader-extension procedure is a special-comment value (see §13.7.3 "Special Comments"), then `read`, `read-syntax`, etc. treat the value as a delimiting comment and otherwise ignore it.

Also, in either context, the result may be copied to prevent mutation to vectors or boxes before the read result is completed, and to support the construction of graphs with cycles. Mutable boxes, vectors, and prefab structures are copied, along with any pairs, boxes, vectors, prefab structures that lead to such mutable values, to placeholders produced by a recursive read (see `read/recursive`), or to references of a shared value. Graph structure (including cycles) is preserved in the copy.

13.7.3 Special Comments

```
(make-special-comment v) → special-comment?
  v : any/c
```

Creates a special-comment value that encapsulates `v`. The `read`, `read-syntax`, etc., procedures treat values constructed with `make-special-comment` as delimiting whitespace when returned by a reader-extension procedure (see §13.7.2 "Reader-Extension Procedures").

```
(special-comment? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is the result of `make-special-comment`, `#f` otherwise.

```
(special-comment-value sc) → any
  sc : special-comment?
```

Returns the value encapsulated by the special-comment value `sc`. This value is never used directly by a reader, but it might be used by the context of a `read-char-or-special`, etc., call that detects a special comment.

13.8 Printer Extension

`gen:custom-write` : `any/c`

A generic interface (see §5.4 “Generic Interfaces”) that supplies a method, `write-proc` used by the default printer to `display`, `write`, or `print` instances of the structure type.

A `write-proc` method takes three arguments: the structure to be printed, the target port, and an argument that is `#t` for `write` mode, `#f` for `display` mode, or `0` or `1` indicating the current quoting depth for `print` mode. The procedure should print the value to the given port using `write`, `display`, `print`, `fprintf`, `write-special`, etc.

The port write handler, port display handler, and print handler are specially configured for a port given to a custom-write procedure. Printing to the port through `display`, `write`, or `print` prints a value recursively with sharing annotations. To avoid a recursive print (i.e., to print without regard to sharing with a value currently being printed), print instead to a string or pipe and transfer the result to the target port using `write-string` or `write-special`. To print recursively to a port other than the one given to the custom-write procedure, copy the given port’s write handler, display handler, and print handler to the other port.

The port given to a custom-write handler is not necessarily the actual target port. In particular, to detect cycles and sharing, the printer invokes a custom-write procedure with a port that records recursive prints, and does not retain any other output.

Recursive print operations may trigger an escape from the call to the custom-write procedure (e.g., for pretty-printing where a tentative print attempt overflows the line, or for printing error output of a limited width).

The following example definition of a `tuple` type includes custom-write procedures that print the tuple’s list content using angle brackets in `write` and `print` mode and no brackets in `display` mode. Elements of the tuple are printed recursively, so that graph and cycle structure can be represented.

Examples:

```
(define (tuple-print tuple port mode)
  (when mode (write-string "<" port))
  (let ([l (tuple-ref tuple)]
        [recur (case mode
                 [(#t) write]
                 [(#f) display]
                 [else (lambda (p port) (print p port mode))])])
    (unless (zero? (vector-length l))
      (recur (vector-ref l 0) port)
      (for-each (lambda (e)
                  (write-string ", " port)
                  (recur e port))
```

```
                (cdr (vector->list l)))))
  (when mode (write-string ">" port)))

> (struct tuple (ref)
          #:methods gen:custom-write
          [(define write-proc tuple-print)])

> (display (tuple #(1 2 "a")))
1, 2, a

> (print (tuple #(1 2 "a")))
<1, 2, "a">

> (let ([t (tuple (vector 1 2 "a"))])
    (vector-set! (tuple-ref t) 0 t)
    (write t))
#0=<#0#, 2, "a">
```

```
prop:custom-write : struct-type-property?
```

A deprecated structure type property (see §5.3 “Structure Type Properties”) that supplies a procedure that corresponds to `gen:custom-write`’s `write-proc`. Use `gen:custom-write`, instead.

```
(custom-write? v) → boolean?
  v : any/c
```

Returns `#t` if *v* has the `prop:custom-write` property, `#f` otherwise.

```
(custom-write-accessor v)
 → (custom-write? output-port? boolean? . -> . any)
  v : custom-write?
```

Returns the custom-write procedure associated with *v*.

```
prop:custom-print-quotable : struct-type-property?
custom-print-quotable? : struct-type-property?
custom-print-quotable-accessor : struct-type-property?
```

A property and associated predicate and accessor. The property value is one of `'self`, `'never`, `'maybe`, or `'always`. When a structure has this property in addition to a `prop:custom-write` property value, then the property value affects printing in `print` mode; see §1.4 “The Printer”. When a value does not have the `prop:custom-print-quotable`, it is equivalent to having the `'self` property value, which is suitable both for self-quoting forms and printed forms that are unreadable.

13.9 Serialization

```
(require racket/serialize)
```
package: base

The bindings documented in this section are provided by the racket/serialize library, not racket/base or racket.

```
(serializable? v) → boolean?
  v : any/c
```

Returns #t if v appears to be serializable, without checking the content of compound values, and #f otherwise. See serialize for an enumeration of serializable values.

```
(serialize v) → any
  v : serializable?
```

Returns a value that encapsulates the value v. This value includes only readable values, so it can be written to a stream with write or s-exp->fasl, later read from a stream using read or fasl->s-exp, and then converted to a value like the original using deserialize. Serialization followed by deserialization produces a value with the same graph structure and mutability as the original value, but the serialized value is a plain tree (i.e., no sharing).

The following kinds of values are serializable:

- structures created through serializable-struct or serializable-struct/versions, or more generally structures with the prop:serializable property (see prop:serializable for more information);
- prefab structures;
- instances of classes defined with define-serializable-class or define-serializable-class*;
- booleans, numbers, characters, interned symbols, unreadable symbols, strings, byte strings, paths (for a specific convention), #<void>, and the empty list;
- pairs, mutable pairs, vectors, flvectors, fxvectors, boxes, hash tables, and sets;
- date, date*, arity-at-least and srcloc structures; and
- module path index values.

Serialization succeeds for a compound value, such as a pair, only if all content of the value is serializable. If a value given to serialize is not completely serializable, the exn:fail:contract exception is raised.

If *v* contains a cycle (i.e., a collection of objects that are all reachable from each other), then *v* can be serialized only if the cycle includes a mutable value, where a prefab structure counts as mutable only if all of its fields are mutable.

The `serialize` and `deserialize` functions currently do not handle certain cyclic values that `read` and `write` can handle, such as `'#0=(#0#)`.

See `deserialize` for information on the format of serialized data.

```
(deserialize v) → any
  v : any/c
```

Given a value *v* that was produced by `serialize`, produces a value like the one given to `serialize`, including the same graph structure and mutability.

A serialized representation *v* is a list of six or seven elements:

- An optional list `'(1)`, `'(2)`, or `'(3)` that represents the version of the serialization format. If the first element of a representation is not a list, then the version is `0`. Version 1 adds support for mutable pairs, version 2 adds support for unreadable symbols, and version 3 adds support for `date*` structures.
- A non-negative exact integer *s-count* that represents the number of distinct structure types represented in the serialized data.
- A list *s-types* of length *s-count*, where each element represents a structure type. Each structure type is encoded as a pair. The `car` of the pair is `#f` for a structure whose deserialization information is defined at the top level, otherwise it is a quoted module path or a byte string (to be converted into a platform-specific path using `bytes->path`) for a module that exports the structure's deserialization information. The `cdr` of the pair is the name of a binding (at the top level or exported from a module) for deserialization information, either a symbol or a string representing an unreadable symbol. These two are used with either `namespace-variable-binding` or `dynamic-require` to obtain deserialization information. See `make-deserialize-info` for more information on the binding's value. See also `deserialize-module-guard`.
- A non-negative exact integer, *g-count* that represents the number of graph points contained in the following list.
- A list *graph* of length *g-count*, where each element represents a serialized value to be referenced during the construction of other serialized values. Each list element is either a box or not:
  - A box represents a value that is part of a cycle, and for deserialization, it must be allocated with `#f` for each of its fields. The content of the box indicates the shape of the value:
    - a non-negative exact integer *i* for an instance of a structure type that is represented by the *i* th element of the *s-types* list;
    - `'c` for a pair, which fails on deserialization (since pairs are immutable; this case does not appear in output generated by `serialize`);

* `'m` for a mutable pair;
* `'b` for a box;
* a pair whose `car` is `'v` and whose `cdr` is a non-negative exact integer *s* for a vector of length *s*;
* a list whose first element is `'h` and whose remaining elements are symbols that determine the hash-table type:
  - `'equal` — `(make-hash)`
  - `'equal 'weak` — `(make-weak-hash)`
  - `'weak` — `(make-weak-hasheq)`
  - no symbols — `(make-hasheq)`
* `'date*` for a `date*` structure, which fails on deserialization (since dates are immutable; this case does not appear in output generated by `serialize`);
* `'date` for a `date` structure, which fails on deserialization (since dates are immutable; this case does not appear in output generated by `serialize`);
* `'arity-at-least` for an `arity-at-least` structure, which fails on deserialization (since arity-at-least are immutable; this case does not appear in output generated by `serialize`); or
* `'mpi` for a module path index, which fails on deserialization (since a module path index is immutable; this case does not appear in output generated by `serialize`).
* `'srcloc` for a `srcloc` structure, which fails on deserialization (since srclocs are immutable; this case does not appear in output generated by `serialize`).

The `#f`-filled value will be updated with content specified by the fifth element of the serialization list *v*.

– A non-box represents a *serial* value to be constructed immediately, and it is one of the following:

* a boolean, number, character, interned symbol, or empty list, representing itself.
* a string, representing an immutable string.
* a byte string, representing an immutable byte string.
* a pair whose `car` is `'?` and whose `cdr` is a non-negative exact integer *i*; it represents the value constructed for the *i*th element of *graph*, where *i* is less than the position of this element within *graph*.
* a pair whose `car` is a number *i*; it represents an instance of a structure type that is described by the *i*th element of the *s-types* list. The `cdr` of the pair is a list of serials representing arguments to be provided to the structure type's deserializer.
* a pair whose `car` is `'q` and whose `cdr` is an immutable value; it represents the quoted value.
* a pair whose `car` is `'f`; it represents an instance of a prefab structure type. The `cadr` of the pair is a prefab structure type key, and the `cddr` is a list of serials representing the field values.

* a pair whose `car` is `'void`, representing `#<void>`.
* a pair whose `car` is `'su` and whose `cdr` is a character string; it represents an unreadable symbol.
* a pair whose `car` is `'u` and whose `cdr` is either a byte string or character string; it represents a mutable byte or character string.
* a pair whose `car` is `'p` and whose `cdr` is a byte string; it represents a path using the serializer's path convention (deprecated in favor of `'p+`).
* a pair whose `car` is `'p+`, whose `cadr` is a byte string, and whose `cddr` is one of the possible symbol results of `system-path-convention-type`; it represents a path using the specified convention.
* a pair whose `car` is `'c` and whose `cdr` is a pair of serials; it represents an immutable pair.
* a pair whose `car` is `'c!` and whose `cdr` is a pair of serials; it represents a pair (but formerly represented a mutable pair), and does not appear in output generated by `serialize`.
* a pair whose `car` is `'m` and whose `cdr` is a pair of serials; it represents a mutable pair.
* a pair whose `car` is `'v` and whose `cdr` is a list of serials; it represents an immutable vector.
* a pair whose `car` is `'v!` and whose `cdr` is a list of serials; it represents a mutable vector.
* a pair whose `car` is `'vl` and whose `cdr` is a list of serials; it represents a flvector.
* a pair whose `car` is `'vx` and whose `cdr` is a list of serials; it represents a fxvector.
* a pair whose `car` is `'b` and whose `cdr` is a serial; it represents an immutable box.
* a pair whose `car` is `'b!` and whose `cdr` is a serial; it represents a mutable box.
* a pair whose `car` is `'h`, whose `cadr` is either `'!` or `'-` (mutable or immutable, respectively), whose `caddr` is a list of symbols (containing `'equal`, `'weak`, both, or neither) that determines the hash table type, and whose `cdddr` is a list of pairs, where the `car` of each pair is a serial for a hash-table key and the `cdr` is a serial for the corresponding value.
* a pair whose `car` is `'date*` and whose `cdr` is a list of serials; it represents a `date*` structure.
* a pair whose `car` is `'date` and whose `cdr` is a list of serials; it represents a `date` structure.
* a pair whose `car` is `'arity-at-least` and whose `cdr` is a serial; it represents an `arity-at-least` structure.
* a pair whose `car` is `'mpi` and whose `cdr` is a pair; it represents a module path index that joins the paired values.

    * a pair whose `car` is `'srcloc` and whose `cdr` is a list of serials; it represents a `srcloc` structure.

- A list of pairs, where the `car` of each pair is a non-negative exact integer *i* and the `cdr` is a serial (as defined in the previous bullet). Each element represents an update to an *i*th element of *graph* that was specified as a box, and the serial describes how to construct a new value with the same shape as specified by the box. The content of this new value must be transferred into the value created for the box in *graph*.

- A final serial (as defined in the two bullets back) representing the result of `deserialize`.

The result of `deserialize` shares no mutable values with the argument to `deserialize`.

If a value provided to `serialize` is a simple tree (i.e., no sharing), then the fourth and fifth elements in the serialized representation will be empty.

```
(serialized=? v1 v2) → boolean?
  v1 : any/c
  v2 : any/c
```

Returns `#t` if *v1* and *v2* represent the same serialization information.

More precisely, it returns the same value that `(equal? (deserialize v1) (deserialize v2))` would return if

- all structure types whose deserializers are accessed with distinct module paths are actually distinct types;

- all structure types are transparent; and

- all structure instances contain only the constituent values recorded in each of *v1* and *v2*.

```
(deserialize-module-guard)
 → (module-path? symbol? . -> . void?)
(deserialize-module-guard guard) → void?
  guard : (module-path? symbol? . -> . void?)
```

A parameter whose value is called by `deserialize` before dynamically loading a module via `dynamic-require`. The two arguments provided to the procedure are the same as the arguments to be passed to `dynamic-require`. The procedure can raise an exception to disallow the `dynamic-require`.

```
(serializable-struct id maybe-super (field ...)
                     struct-option ...)
```

Like `struct`, but instances of the structure type are serializable with `serialize`. This form is allowed only at the top level or in a module's top level (so that deserialization information can be found later).

Serialization only supports cycles involving the created structure type when all fields are mutable (or when the cycle can be broken through some other mutable value).

In addition to the bindings generated by `struct`, `serializable-struct` binds `deserialize-info:`*id*`-v0` to deserialization information. Furthermore, in a module context, it automatically `provide`s this binding in a `deserialize-info` submodule using `module+`.

The `serializable-struct` form enables the construction of structure instances from places where *id* is not accessible, since deserialization must construct instances. Furthermore, `serializable-struct` provides limited access to field mutation, but only for instances generated through the deserialization information bound to `deserialize-info:`*id*`-v0`. See `make-deserialize-info` for more information.

The `-v0` suffix on the deserialization enables future versioning on the structure type through `serializable-struct/version`.

When a supertype is supplied as *maybe-super*, compile-time information bound to the supertype identifier must include all of the supertype's field accessors. If any field mutator is missing, the structure type will be treated as immutable for the purposes of marshaling (so cycles involving only instances of the structure type cannot be handled by the deserializer).

Examples:

```
> (serializable-struct point (x y))

> (point-x (deserialize (serialize (point 1 2))))
1
```

```
(define-serializable-struct id-maybe-super (field ...)
                            struct-option ...)
```

Like `serializable-struct`, but with the supertype syntax and default constructor name of `define-struct`.

```
(serializable-struct/versions id maybe-super vers (field ...)
                              (other-version-clause ...)
                              struct-option ...)

other-version-clause = (other-vers make-proc-expr
                                   cycle-make-proc-expr)
```

Like `serializable-struct`, but the generated deserializer binding is `deserialize-info:`*id*`-v`*vers*. In addition, `deserialize-info:`*id*`-v`*other-vers* is bound for each *other-vers*. The *vers* and each *other-vers* must be a literal, exact, nonnegative integer.

Each *make-proc-expr* should produce a procedure, and the procedure should accept as many argument as fields in the corresponding version of the structure type, and it produce an instance of *id*. Each *cycle-make-proc-expr* should produce a procedure of no arguments; this procedure should return two values: an instance `x` of *id* (typically with `#f` for all fields) and a procedure that accepts another instance of *id* and copies its field values into `x`.

Examples:

```
> (serializable-struct point (x y) #:mutable #:transparent)

> (define ps (serialize (point 1 2)))

> (deserialize ps)
(point 1 2)
> (define x (point 1 10))

> (set-point-x! x x)

> (define xs (serialize x))

> (deserialize xs)
#0=(point #0# 10)
> (serializable-struct/versions point 1 (x y z)
    ([0
      ; Constructor for simple v0 instances:
      (lambda (x y) (point x y 0))
      ; Constructor for v0 instance in a cycle:
      (lambda ()
        (let ([p0 (point #f #f 0)])
          (values
            p0
            (lambda (p)
              (set-point-x! p0 (point-x p))
              (set-point-y! p0 (point-y p))))))])
    #:mutable #:transparent)

> (deserialize (serialize (point 4 5 6)))
(point 4 5 6)
> (deserialize ps)
(point 1 2 0)
> (deserialize xs)
#0=(point #0# 10 0)
```

```
(define-serializable-struct/versions id-maybe-super vers (field ...)
                                     (other-version-clause ...)
                                     struct-option ...)
```

Like `serializable-struct/versions`, but with the supertype syntax and default constructor name of `define-struct`.

```
(make-deserialize-info make cycle-make) → any
  make : procedure?
  cycle-make : (-> (values any/c procedure?))
```

Produces a deserialization information record to be used by `deserialize`. This information is normally tied to a particular structure because the structure has a `prop:serializable` property value that points to a top-level variable or module-exported variable that is bound to deserialization information.

The *make* procedure should accept as many arguments as the structure's serializer put into a vector; normally, this is the number of fields in the structure. It should return an instance of the structure.

The *cycle-make* procedure should accept no arguments, and it should return two values: a structure instance `x` (with dummy field values) and an update procedure. The update procedure takes another structure instance generated by the *make*, and it transfers the field values of this instance into `x`.

```
prop:serializable : property?
```

This property identifies structures and structure types that are serializable. The property value should be constructed with `make-serialize-info`.

```
(make-serialize-info to-vector
                     deserialize-id
                     can-cycle?
                     dir)            → any
  to-vector : (any/c . -> . vector?)
  deserialize-id :  (or identifier?
                        symbol?
                        (cons/c symbol?
                                module-path-index?))
  can-cycle? : any/c
  dir : path-string?
```

Produces a value to be associated with a structure type through the `prop:serializable` property. This value is used by `serialize`.

The *to-vector* procedure should accept a structure instance and produce a vector for the instance's content.

The *deserialize-id* value indicates a binding for deserialize information, to either a module export or a top-level definition. It must be one of the following:

- If *deserialize-id* is an identifier, and if (identifier-binding *deserialize-id*) produces a list, then the third element is used for the exporting module, otherwise the top-level is assumed. Before trying an exporting module directly, its deserialize-info submodule is tried; the module itself is tried if no deserialize-info submodule is available or if the export is not found. In either case, syntax-e is used to obtain the name of an exported identifier or top-level definition.
- If *deserialize-id* is a symbol, it indicates a top-level variable that is named by the symbol.
- If *deserialize-id* is a pair, the car must be a symbol to name an exported identifier, and the cdr must be a module path index to specify the exporting module.

See make-deserialize-info and deserialize for more information.

The *can-cycle?* argument should be false if instances should not be serialized in such a way that deserialization requires creating a structure instance with dummy field values and then updating the instance later.

The *dir* argument should be a directory path that is used to resolve a module reference for the binding of *deserialize-id*. This directory path is used as a last resort when *deserialize-id* indicates a module that was loaded through a relative path with respect to the top level. Usually, it should be (or (current-load-relative-directory) (current-directory)).

13.10 Fast-Load Serialization

(require racket/fasl) package: base

The bindings documented in this section are provided by the racket/fasl library, not racket/base or racket.

```
(s-exp->fasl v [out]) → (or/c (void) bytes?)
  v : any/c
  out : (or/c output-port? #f) = #f
(fasl->s-exp in) → any/c
  in : (or/c input-port? bytes?)
```

The s-exp->fasl function serializes *v* to a byte string, printing it directly to *out* if *out* is an output port or return the byte string otherwise. The fasl->s-exp function decodes a value from a byte string (supplied either directly or as an input port) that was encoded with s-exp->fasl.

The *v* argument must be a value that could be quoted as a literal, because s-exp->fasl essentially uses (compile `',v) to encode the value using Racket's built-in fast-load format

for bytecode.

The byte-string encoding produced by `s-exp->fasl` is specific to a version of Racket. That is, the resulting byte string can be decoded back to `v` only using the same version with which it was encoded.

Examples:

```
> (define fasl (s-exp->fasl (list #("speed") 'racer #\!)))

> fasl
#"#~\0056.1.1T\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\1\0\0\0\1\0\0\0\0\36\0\0\0\220$\24r\2
> (fasl->s-exp fasl)
'(#("speed") racer #\!)
```

14 Reflection and Security

14.1 Namespaces

See §1.2.5 “Namespaces” for basic information on the namespace model.

A new namespace is created with procedures like `make-empty-namespace`, and `make-base-namespace`, which return a first-class namespace value. A namespace is used by setting the `current-namespace` parameter value, or by providing the namespace to procedures such as `eval` and `eval-syntax`.

```
(namespace? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a namespace value, `#f` otherwise.

```
(make-empty-namespace) → namespace?
```

Creates a new namespace that is empty, and whose module registry contains no mappings. The namespace’s base phase is the same as the base phase of the current namespace. Attach modules from an existing namespace to the new one with `namespace-attach-module`.

```
(make-base-empty-namespace) → namespace?
```

Creates a new empty namespace, but with `racket/base` attached. The namespace’s base phase is the same as the phase in which the `make-base-empty-namespace` function was created.

```
(make-base-namespace) → namespace?
```

Creates a new namespace with `racket/base` attached and `required` into the top-level environment. The namespace’s base phase is the same as the phase in which the `make-base-namespace` function was created.

```
(define-namespace-anchor id)
```

Binds *id* to a namespace anchor that can be used with `namespace-anchor->empty-namespace` and `namespace-anchor->namespace`.

This form can be used only in a top-level context or in a module-context.

```
(namespace-anchor? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a namespace-anchor value, `#f` otherwise.

```
(namespace-anchor->empty-namespace a) → namespace?
  a : namespace-anchor?
```

Returns an empty namespace that shares a module registry with the source of the anchor, and whose base phase is the phase in which the anchor was created.

If the anchor is from a `define-namespace-anchor` form in a module context, then the source is the namespace in which the containing module is instantiated. If the anchor is from a `define-namespace-anchor` form in a top-level content, then the source is the namespace in which the anchor definition was evaluated.

```
(namespace-anchor->namespace a) → namespace?
  a : namespace-anchor?
```

Returns a namespace corresponding to the source of the anchor.

If the anchor is from a `define-namespace-anchor` form in a module context, then the result is a namespace for the module's body in the anchor's phase. The result is the same as a namespace obtained via `module->namespace`.

If the anchor is from a `define-namespace-anchor` form in a top-level content, then the result is the namespace in which the anchor definition was evaluated.

```
(current-namespace) → namespace?
(current-namespace n) → void?
  n : namespace?
```

A parameter that determines the current namespace.

```
(namespace-symbol->identifier sym) → identifier?
  sym : symbol?
```

Similar to `datum->syntax` restricted to symbols. The lexical information of the resulting identifier corresponds to the top-level environment of the current namespace; the identifier has no source location or properties.

```
(namespace-base-phase [namespace]) → exact-integer?
  namespace : namespace? = (current-namespace)
```

Returns the base phase of *namespace*.

```
(namespace-module-identifier [where]) → identifier?
  where : (or/c namespace? exact-integer? #f)
        = (current-namespace)
```

Returns an identifier whose binding is `module` in the base phase of *where* if it is a namespace, or in the *where* phase level otherwise.

The lexical information of the identifier includes bindings (in the same phase level) for all syntactic forms that appear in fully expanded code (see §1.2.3.1 "Fully Expanded Programs"), but using the name reported by the second element of `identifier-binding` for the binding; the lexical information may also include other bindings.

```
(namespace-variable-value  sym
                           [use-mapping?
                            failure-thunk
                            namespace])    → any
  sym : symbol?
  use-mapping? : any/c = #t
  failure-thunk : (or/c (-> any) #f) = #f
  namespace : namespace? = (current-namespace)
```

Returns a value for *sym* in *namespace*, using *namespace*'s base phase. The returned value depends on *use-mapping?*:

- If *use-mapping?* is true (the default), and if *sym* maps to a top-level variable or an imported variable (see §1.2.5 "Namespaces"), then the result is the same as evaluating *sym* as an expression. If *sym* maps to syntax or imported syntax, then *failure-thunk* is called or the `exn:fail:syntax` exception is raised. If *sym* is mapped to an undefined variable or an uninitialized module variable, then *failure-thunk* is called or the `exn:fail:contract:variable` exception is raised.
- If *use-mapping?* is `#f`, the namespace's syntax and import mappings are ignored. Instead, the value of the top-level variable named *sym* in namespace is returned. If the variable is undefined, then *failure-thunk* is called or the `exn:fail:contract:variable` exception is raised.

If *failure-thunk* is not `#f`, `namespace-variable-value` calls *failure-thunk* to produce the return value in place of raising an `exn:fail:contract:variable` or `exn:fail:syntax` exception.

```
(namespace-set-variable-value!  sym
                                v
                                [map?
                                 namespace]) → void?
  sym : symbol?
  v : any/c
  map? : any/c = #f
  namespace : namespace? = (current-namespace)
```

Sets the value of *sym* in the top-level environment of *namespace* in the base phase, defining *sym* if it is not already defined.

If *map?* is supplied as true, then the namespace's identifier mapping is also adjusted (see §1.2.5 "Namespaces") in the phase level corresponding to the base phase, so that *sym* maps to the variable.

```
(namespace-undefine-variable! sym
                              [namespace]) → void?
  sym : symbol?
  namespace : namespace? = (current-namespace)
```

Removes the *sym* variable, if any, in the top-level environment of *namespace* in its base phase. The namespace's identifier mapping (see §1.2.5 "Namespaces") is unaffected.

```
(namespace-mapped-symbols [namespace]) → (listof symbol?)
  namespace : namespace? = (current-namespace)
```

Returns a list of all symbols that are mapped to variables, syntax, and imports in *namespace* for the phase level corresponding to the namespace's base phase.

```
(namespace-require quoted-raw-require-spec) → void?
  quoted-raw-require-spec : any/c
```

Performs the import corresponding to *quoted-raw-require-spec* in the top-level environment of the current namespace, like a top-level `#%require`. The *quoted-raw-require-spec* argument must be a datum that corresponds to a quoted *raw-require-spec* for `#%require`, which includes module paths.

Module paths in *quoted-raw-require-spec* are resolved with respect to `current-load-relative-directory` or `current-directory` (if the former is `#f`), even if the current namespace corresponds to a module body.

```
(namespace-require/copy quoted-raw-require-spec) → void?
  quoted-raw-require-spec : any/c
```

Like `namespace-require` for syntax exported from the module, but exported variables at the namespace's base phase are treated differently: the export's current value is copied to a top-level variable in the current namespace.

```
(namespace-require/constant quoted-raw-require-spec) → void?
  quoted-raw-require-spec : any/c
```

Like `namespace-require`, but for each exported variable at the namespace's base phase, the export's value is copied to a corresponding top-level variable that is made immutable. Despite setting the top-level variable, the corresponding identifier is bound as imported.

```
(namespace-require/expansion-time quoted-raw-require-spec)
 → void?
  quoted-raw-require-spec : any/c
```

Like namespace-require, but only the transformer part of the module is executed relative to the namespace's base phase; that is, the module is merely visited, and not instantiated (see §1.2.3.8 "Module Phases and Visits"). If the required module has not been instantiated before, the module's variables remain undefined.

```
(namespace-attach-module  src-namespace
                          modname
                         [dest-namespace]) → void?
  src-namespace : namespace?
  modname : module-path?
  dest-namespace : namespace? = (current-namespace)
```

Attaches the instantiated module named by *modname* in *src-namespace* (at its base phase) to the module registry of *dest-namespace*.

In addition to *modname*, every module that it imports (directly or indirectly) is also recorded in the current namespace's module registry, and instances at the same phase or lower are also attached to *dest-namespace* (while visits at the module's phase and instances at higher phases are not attached, nor even made available for on-demand visits). The inspector of the module invocation in *dest-namespace* is the same as inspector of the invocation in *src-namespace*.

If *modname* is not a symbol, the current module name resolver is called to resolve the path, but no module is loaded; the resolved form of *modname* is used as the module name in *dest-namespace*.

If *modname* refers to a submodule or a module with submodules, unless the module was loaded from bytecode (i.e., a ".zo" file) independently from submodules within the same top-level module, then declarations for all submodules within the module's top-level module are also attached to *dest-namespace*.

If *modname* does not refer to an instantiated module in *src-namespace*, or if the name of any module to be attached already has a different declaration or same-phase instance in *dest-namespace*, then the exn:fail:contract exception is raised.

If *src-namespace* and *dest-namespace* do not have the same base phase, then the exn:fail:contract exception is raised.

```
(namespace-attach-module-declaration  src-namespace
                                      modname
                                     [dest-namespace]) → void?
  src-namespace : namespace?
  modname : module-path?
  dest-namespace : namespace? = (current-namespace)
```

Like namespace-attach-module, but the module specified by *modname* need only be de-

clared (and not necessarily instantiated) in *src-namespace*, and the module is merely declared in *dest-namespace*.

```
(namespace-unprotect-module  inspector
                             modname
                            [namespace]) → void?
  inspector : inspector?
  modname : module-path?
  namespace : namespace? = (current-namespace)
```

Changes the inspector for the instance of the module referenced by *modname* in *namespace*'s module registry so that it is controlled by the current code inspector. The given *inspector* must currently control the invocation of the module in *namespace*'s module registry, otherwise the `exn:fail:contract` exception is raised. See also §14.10 "Code Inspectors".

```
(namespace-module-registry namespace) → any
  namespace : namespace?
```

Returns the module registry of the given namespace. This value is useful only for identification via `eq?`.

```
(module->namespace mod) → namespace?
       (or/c module-path?
  mod :      resolved-module-path?
             module-path-index?)
```

Returns a namespace that corresponds to the body of an instantiated module in the current namespace's module registry and in the current namespace's base phase, making the module available for on-demand visits at the namespace's base phase. The returned namespace has the same module registry as the current namespace. Modifying a binding in the namespace changes the binding seen in modules that require the namespace's module.

Module paths in a top-level `require` expression are resolved with respect to the namespace's module. New `provide` declarations are not allowed.

If the current code inspector does not control the invocation of the module in the current namespace's module registry, the `exn:fail:contract` exception is raised; see also §14.10 "Code Inspectors".

Bindings in the namespace cannot be modified if the `compile-enforce-module-constants` parameter was true when the module was declared, unless the module declaration itself included assignments to the binding via `set!`.

```
(namespace-syntax-introduce stx) → syntax?
  stx : syntax?
```

Returns a syntax object like *stx*, except that the current namespace's bindings are included in the syntax object's lexical information (see §1.2.2 "Syntax Objects"). The additional context is overridden by any existing top-level bindings in the syntax object's lexical information, or by any existing or future module bindings in the lexical information.

```
(module-provide-protected? module-path-index
                           sym)              → boolean?
  module-path-index : (or/c symbol? module-path-index?)
  sym : symbol?
```

Returns `#f` if the module declaration for *module-path-index* defines *sym* and exports it unprotected, `#t` otherwise (which may mean that the symbol corresponds to an unexported definition, a protected export, or an identifier that is not defined at all within the module).

The *module-path-index* argument can be a symbol; see §14.4.2 "Compiled Modules and References" for more information on module path indices.

Typically, the arguments to `module-provide-protected?` correspond to the first two elements of a list produced by `identifier-binding`.

```
(variable-reference? v) → boolean?
  v : any/c
```

Return `#t` if *v* is a variable reference produced by `#%variable-reference`, `#f` otherwise.

```
(variable-reference-constant? varref) → boolean?
  varref : variable-reference?
```

Returns `#t` if the variable represented by *varref* will retain its current value (i.e., *varref* refers to a variable that cannot be further modified by `set!` or `define`), `#f` otherwise.

```
(variable-reference->empty-namespace varref) → namespace?
  varref : variable-reference?
```

Returns an empty namespace that shares module declarations and instances with the namespace in which *varref* is instantiated, and with the same phase as *varref*.

```
(variable-reference->namespace varref) → namespace?
  varref : variable-reference?
```

If *varref* refers to a module-level variable, then the result is a namespace for the module's body in the referenced variable's phase; the result is the same as a namespace obtained via `module->namespace`.

If *varref* refers to a top-level variable, then the result is the namespace in which the referenced variable is defined.

```
(variable-reference->resolved-module-path varref)
 → (or/c resolved-module-path? #f)
  varref : variable-reference?
```

If *varref* refers to a module-level variable, the result is a resolved module path naming the module.

If *varref* refers to a top-level variable, then the result is `#f`.

```
(variable-reference->module-path-index varref)
 → (or/c module-path-index? #f)
  varref : variable-reference?
```

If *varref* refers to a module-level variable, the result is a module path index naming the module.

If *varref* refers to a top-level variable, then the result is `#f`.

```
(variable-reference->module-source varref)
 → (or/c symbol? (and/c path? complete-path?) #f)
  varref : variable-reference?
```

If *varref* refers to a module-level variable, the result is a path or symbol naming the module's source (which is typically, but not always, the same as in the resolved module path). If the relevant module is a submodule, the result corresponds to the enclosing top-level module's source.

If *varref* refers to a top-level variable, then the result is `#f`.

```
(variable-reference->phase varref) → exact-nonnegative-integer?
  varref : variable-reference?
```

Returns the phase of the variable referenced by *varref*.

```
(variable-reference->module-base-phase varref) → exact-integer?
  varref : variable-reference?
```

Returns the phase in which the module is instantiated for the variable referenced by *varref*, or `0` if the variable for *varref* is not within a module.

For a variable with a module, the result is less than the result of `(variable-reference->phase varref)` by *n* when the variable is bound at phase level *n* within the module.

```
(variable-reference->module-declaration-inspector varref)
 → inspector?
  varref : variable-reference?
```

Returns the declaration inspector (see §14.10 “Code Inspectors”) for the module of *varref*, where *varref* must refer to an anonymous module variable as produced by `(#%variable-reference)`.

14.2 Evaluation and Compilation

```
(current-eval) → (any/c . -> . any)
(current-eval proc) → void?
  proc : (any/c . -> . any)
```

A parameter that determines the current *evaluation handler*. The evaluation handler is a procedure that takes a top-level form and evaluates it, returning the resulting values. The evaluation handler is called by `eval`, `eval-syntax`, the default load handler, and `read-eval-print-loop` to evaluate a top-level form. The handler should evaluate its argument in tail position.

The *top-level-form* provided to the handler can be a syntax object, a compiled form, a compiled form wrapped as a syntax object, or an arbitrary datum.

The default handler converts an arbitrary datum to a syntax object using `datum->syntax`, and then enriches its lexical information in the same way as `eval`. (If *top-level-form* is a syntax object, then its lexical information is not enriched.) The default evaluation handler partially expands the form to splice the body of top-level `begin` forms into the top level (see `expand-to-top-form`), and then individually compiles and evaluates each spliced form before continuing to expand, compile, and evaluate later forms.

```
(eval top-level-form [namespace]) → any
  top-level-form : any/c
  namespace : namespace? = (current-namespace)
```

See also §15.1.2 “Namespaces” in ***The Racket Guide***.

Calls the current evaluation handler to evaluate *top-level-form*. The evaluation handler is called in tail position with respect to the `eval` call, and `parameterized` to set `current-namespace` to *namespace*.

If *top-level-form* is a syntax object whose datum is not a compiled form, then its lexical information is enriched before it is sent to the evaluation handler:

- If *top-level-form* is a pair whose `car` is a symbol or identifier, and if applying `namespace-syntax-introduce` to the (`datum->syntax`-converted) identifier produces an identifier bound to `module` in a phase level that corresponds to *namespace*’s base phase, then only that identifier is enriched.
- For any other *top-level-form*, `namespace-syntax-introduce` is applied to the entire syntax object.

For interactive evaluation in the style of `read-eval-print-loop` and `load`, wrap each expression with `#%top-interaction`, which is normally bound to `#%top-interaction`, before passing it to `eval`.

```
(eval-syntax stx [namespace]) → any
  stx : syntax?
  namespace : namespace? = (current-namespace)
```

Like `eval`, except that *stx* must be a syntax object, and its lexical context is not enriched before it is passed to the evaluation handler.

```
(current-load)
    (path? (or/c #f
                 symbol?
                 (cons/c (or/c #f symbol?)
 →                       (non-empty-listof symbol?)))
           . -> .
           any)
(current-load proc) → void?
         (path? (or/c #f
                      symbol?
                      (cons/c (or/c #f symbol?)
  proc :                      (non-empty-listof symbol?)))
                . -> .
                any)
```

A parameter that determines the current *load handler* to load top-level forms from a file. The load handler is called by `load`, `load-relative`, `load/cd`, and the default compiled-load handler.

A load handler takes two arguments: a path (see §15.1 “Paths”) and an expected module name. The expected module name is a symbol or a list when the call is to load a module declaration in response to a `require` (in which case the file should contain a module declaration), or `#f` for any other load.

When loading a module from a stream that starts with a compiled module that contains submodules, the load handler should load only the requested module, where a symbol as the load handler’s indicates the root module and a list indicates a submodule whose path relative to the root module is given by the `cdr` of the list. The list starts with `#f` when a submodule should be loaded *only* if it can be loaded independently (i.e., from compiled form—never from source); if the submodule cannot be loaded independently, the load handler should return without loading from a file. When the expected module name is a list that starts with a symbol, the root module and any other submodules can be loaded from the given file, which might be from source, and the load handler still should not complain if the expected submodule is not found.

The default load handler reads forms from the file in `read-syntax` mode with line-counting enabled for the file port, unless the path has a `".zo"` suffix. It also `parameterizes` each read to set `read-accept-compiled`, `read-accept-reader`, and `read-accept-lang` to `#t`. In addition, if `load-on-demand-enabled` is `#t`, then `read-on-demand-source` is set to the cleansed, absolute form of `path` during the `read-syntax` call. After reading a single form, the form is passed to the current evaluation handler, wrapping the evaluation in a continuation prompt (see `call-with-continuation-prompt`) for the default continuation prompt tag with handler that propagates the abort to the continuation of the `load` call.

If the second argument to the load handler is a symbol, then:

- The `read-syntax` from the file is additionally `parameterized` as follows (to provide consistent reading of module source):

```
(current-readtable #f)
(read-case-sensitive #t)
(read-square-bracket-as-paren #t)
(read-curly-brace-as-paren #t)
(read-accept-box #t)
(read-accept-compiled #t)
(read-accept-bar-quote #t)
(read-accept-graph #t)
(read-decimal-as-inexact #t)
(read-accept-dot #t)
(read-accept-infix-dot #t)
(read-accept-quasiquote #t)
(read-accept-reader #t)
(read-accept-lang #t)
```

- If the read result is not a `module` form, or if a second `read-syntax` does not produce an end-of-file, then the `exn:fail` exception is raised without evaluating the form that was read from the file. (In previous versions, the module declaration was checked to match the name given as the second argument to the load handler, but this check is no longer performed.)
- The lexical information of the initial `module` identifier is enriched with a binding for `module`, so that the form corresponds to a module declaration independent of the current namespace's bindings.

If the second argument to the load handler is `#f`, then each expression read from the file is wrapped with `#%top-interaction`, which is normally bound to `#%top-interaction`, before passing it to the evaluation handler.

The return value from the default load handler is the value of the last form from the loaded file, or `#<void>` if the file contains no forms. If the given path is a relative path, then it is resolved using the value of `current-directory`.

```
(load file) → any
  file : path-string?
```

See also §15.1.2 “Namespaces” in ***The Racket Guide***.

Calls the current load handler in tail position. The call is `parameterized` to set `current-load-relative-directory` to the directory of *file*, which is resolved relative to the value of `current-directory`.

```
(load-relative file) → any
  file : path-string?
```

Like `load/use-compiled`, but when *file* is a relative path, it is resolved using the value of `current-load-relative-directory` instead of the value of `current-directory` if the former is not `#f`, otherwise `current-directory` is used.

```
(load/cd file) → any
  file : path-string?
```

Like `load`, but `load/cd` sets both `current-directory` and `current-load-relative-directory` before calling the load handler.

```
(current-load-extension)
 → (path? (or/c symbol? #f) . -> . any)
(current-load-extension proc) → void?
  proc : (path? (or/c symbol? #f) . -> . any)
```

A parameter that determines a *extension-load handler*, which is called by `load-extension` and the default compiled-load handler.

An extension-load handler takes the same arguments as a load handler, but the file should be a platform-specific *dynamic extension*, typically with the file suffix `".so"` (Unix), `".dll"` (Windows), or `".dylib"` (Mac OS X). The file is loaded using internal, OS-specific primitives. See *Inside: Racket C API* for more information on dynamic extensions.

```
(load-extension file) → any
  file : path-string?
```

Sets `current-load-relative-directory` like `load`, and calls the extension-load handler in tail position.

```
(load-relative-extension file) → any
  file : path-string?
```

Like `load-extension`, but resolves *file* using `current-load-relative-directory` like `load-relative`.

```
(current-load/use-compiled)
    (path? (or/c #f
                 symbol?
 →               (cons/c (or/c #f symbol?)
                         (non-empty-listof symbol?)))
           . -> . any)
(current-load/use-compiled proc) → void?
         (path? (or/c #f
                      symbol?
  proc :              (cons/c (or/c #f symbol?)
                              (non-empty-listof symbol?)))
                . -> . any)
```

A parameter that determines the current *compiled-load handler* to load from a file that may have a compiled form. The compiled-load handler is called by `load/use-compiled`.

The protocol for a compiled-load handler is the same as for the load handler (see `current-load`), except that a compiled-load handler is expected to set `current-load-relative-directory` itself. The default compiled-load handler, however, checks for a `".ss"` file when the given path ends with `".rkt"`, no `".rkt"` file exists, and when the handler's second argument is a symbol. In addition, the default compiled-load handler checks for `".zo"` (bytecode) files and `".so"` (native Unix), `".dll"` (native Windows), or `".dylib"` (native Mac OS X) files.

The check for a compiled file occurs whenever the given path *file* ends with any extension (e.g., `".rkt"` or `".scrbl"`), and the check consults the subdirectories indicated by the `current-compiled-file-roots` and `use-compiled-file-paths` parameters relative to *file*, where the former supplies "roots" for compiled files and the latter provides subdirectories. A "root" can be an absolute path, in which case *file*'s directory is combined with `reroot-path` and the root as the second argument; if the "root" is a relative path, then the relative path is instead suffixed onto the directory of *file*. The roots are tried in order, and the subdirectories are checked in order within each root. A `".zo"` version of the file (whose name is formed by passing *file* and `#".zo"` to `path-add-suffix`) is loaded if it exists directly in one of the indicated subdirectories, or a `".so"/".dll"/".dylib"` version of the file is loaded if it exists within a `"native"` subdirectory of a `use-compiled-file-paths` directory, in an even deeper subdirectory as named by `system-library-subpath`. A compiled file is loaded only if its modification date is not older than the date for *file*. If both `".zo"` and `".so"/".dll"/".dylib"` files are available, the `".so"/".dll"/".dylib"` file is used. If *file* ends with `".rkt"`, no such file exists, the handler's second argument is a symbol, and a `".ss"` file exists, then `".zo"` and `".so"/".dll"/".dylib"` files are used only with names based on *file* with its suffixed replaced by `".ss"`.

See also `compiler/compilation-path`.

While a `".zo"`, `".so"`, `".dll"`, or `".dylib"` file is loaded, the current `load-relative` directory is set to the directory of the original *file*. If the file to be loaded has the suffix `".ss"` while the requested file has the suffix `".rkt"`, then the `current-module-declare-`

`source` parameter is set to the full path of the loaded file, otherwise the `current-module-declare-source` parameter is set to `#f`.

If the original *`file`* is loaded or a `".zo"` variant is loaded, the load handler is called to load the file. If any other kind of file is loaded, the extension-load handler is called.

When the default compiled-load handler loads a module from a bytecode (i.e., `".zo"`) file, the handler records the bytecode file path in the current namespace's module registry. More specifically, the handler records the path for the top-level module of the loaded module, which is an enclosing module if the loaded module is a submodule. Thereafter, loads via the default compiled-load handler for modules within the same top-level module use the recorded file, independent of the file that otherwise would be selected by the compiled-load handler (e.g., even if the `use-compiled-file-paths` parameter value changes). The default module name resolver transfers bytecode-file information when a module declaration is attached to a new namespace. This protocol supports independent but consistent loading of submodules from bytecode files.

```
(load/use-compiled file) → any
  file : path-string?
```

Calls the current compiled-load handler in tail position.

```
(current-load-relative-directory)
 → (or/c (and/c path-string? complete-path?) #f)
(current-load-relative-directory path) → void?
  path : (or/c (and/c path-string? complete-path?) #f)
```

A parameter that is set by `load`, `load-relative`, `load-extension`, `load-relative-extension`, and the default compiled-load handler, and used by `load-relative`, `load-relative-extension`, and the default compiled-load handler.

When a new path or string is provided as the parameter's value, it is immediately expanded (see §15.1 "Paths") and converted to a path. (The directory need not exist.)

```
(use-compiled-file-paths)
 → (listof (and/c path? relative-path?))
(use-compiled-file-paths paths) → void?
  paths : (listof (and/c path-string? relative-path?))
```

A list of relative paths, which defaults to `(list (string->path "compiled"))`. It is used by the compiled-load handler (see `current-load/use-compiled`).

```
(current-compiled-file-roots) → (listof (or/c path? 'same))
(current-compiled-file-roots paths) → void?
  paths : (listof (or/c path-string? 'same))
```

A list of paths and `'same`s that is is used by the default compiled-load handler (see `current-load/use-compiled`).

The parameter is normally initialized to `(list 'same)`, but the parameter's initial value can be adjusted by the `PLTCOMPILEDROOTS` environment variable or the `--compiled` or `-R` command-line flag for `racket`. If the environment variable is defined and not overridden by a command-line flag, it is parsed by first replacing any `@(version)` with the result of `(version)`, then using `path-list-string->path-list` with a default path list `(list (build-path 'same))` to arrive at the parameter's initial value.

```
(read-eval-print-loop) → any
```

Starts a new *REPL* using the current input, output, and error ports. The REPL wraps each expression to evaluate with `#%top-interaction`, which is normally bound to `#%top-interaction`, and it wraps each evaluation with a continuation prompt using the default continuation prompt tag and prompt handler (see `call-with-continuation-prompt`). The REPL also wraps the read and print operations with a prompt for the default tag whose handler ignores abort arguments and continues the loop. The `read-eval-print-loop` procedure does not return until `eof` is read, at which point it returns `#<void>`.

The `read-eval-print-loop` procedure can be configured through the `current-prompt-read`, `current-eval`, and `current-print` parameters.

```
(current-prompt-read) → (-> any)
(current-prompt-read proc) → void?
  proc : (-> any)
```

A parameter that determines a *prompt read handler*, which is a procedure that takes no arguments, displays a prompt string, and returns a top-level form to evaluate. The prompt read handler is called by `read-eval-print-loop`, and after printing a prompt, the handler typically should call the read interaction handler (as determined by the `current-read-interaction` parameter) with the port produced by the interaction port handler (as determined by the `current-get-interaction-input-port` parameter).

The default prompt read handler prints `> ` and returns the result of

```
(let ([in ((current-get-interaction-input-port))])
  ((current-read-interaction) (object-name in) in))
```

If the input and output ports are both terminals (in the sense of `terminal-port?`) and if the output port appears to be counting lines (because `port-next-location` returns a non-`#f` line and column), then the output port's line is incremented and its column is reset to `0` via `set-port-next-location!` before returning the read result.

```
(current-get-interaction-input-port) → (-> input-port?)
(current-get-interaction-input-port proc) → void?
  proc : (-> input-port?)
```

A parameter that determines the *interaction port handler*, which returns a port to use for `read-eval-print-loop` inputs.

The default interaction port handler returns the current input port. In addition, if that port is the initial current input port, the initial current output and error ports are flushed.

The `racket/gui/base` library adjusts this parameter's value by extending the current value. The extension wraps the result port so that GUI events can be handled when reading from the port blocks.

```
(current-read-interaction) → (any/c input-port? -> any)
(current-read-interaction proc) → void?
  proc : (any/c input-port? -> any)
```

A parameter that determines the current *read interaction handler*, which is procedure that takes an arbitrary value and an input port and returns an expression read from the input port.

The default read interaction handler accepts *src* and *in* and returns

```
(parameterize ([read-accept-reader #t]
               [read-accept-lang #f])
  (read-syntax src in))
```

```
(current-print) → (any/c -> any)
(current-print proc) → void?
  proc : (any/c -> any)
```

A parameter that determines the *print handler* that is called by `read-eval-print-loop` to print the result of an evaluation (and the result is ignored).

The default print handler `print`s the value to the current output port (as determined by the `current-output-port` parameter) and then outputs a newline, except that it prints nothing when the value is `#<void>`.

```
(current-compile)
 → (any/c boolean? . -> . compiled-expression?)
(current-compile proc) → void?
  proc : (any/c boolean? . -> . compiled-expression?)
```

A parameter that determines the current *compilation handler*. The compilation handler is a procedure that takes a top-level form and returns a compiled form; see §1.2.4 "Compilation" for more information on compilation.

The compilation handler is called by `compile`, and indirectly by the default evaluation handler and the default load handler.

The handler's second argument is `#t` if the compiled form will be used only for immediate evaluation, or `#f` if the compiled form may be saved for later use; the default compilation handler is optimized for the special case of immediate evaluation.

When a compiled form is written to an output port, the written form starts with `#~`. See §1.4.16 "Printing Compiled Code" for more information.

For internal testing purposes, when the `PLT_VALIDATE_COMPILE` environment variable is set, the default compilation handler runs a bytecode validator on its own compilation results.

```
(compile top-level-form) → compiled-expression?
  top-level-form : any/c
```

Like `eval`, but calls the current compilation handler in tail position with *top-level-form*.

```
(compile-syntax stx) → compiled-expression?
  stx : syntax?
```

Like `eval-syntax`, but calls the current compilation handler in tail position with *stx*.

```
(compiled-expression? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a compiled form, `#f` otherwise.

```
(compile-enforce-module-constants) → boolean?
(compile-enforce-module-constants on?) → void?
  on? : any/c
```

A parameter that determines how a module declaration is compiled.

When constants are enforced, and when the macro-expanded body of a module contains no `set!` assignment to a particular variable defined within the module, then the variable is marked as constant when the definition is evaluated. Afterward, the variable's value cannot be assigned or undefined through `module->namespace`, and it cannot be defined by redeclaring the module.

Enforcing constants allows the compiler to inline some variable values, and it allows the native-code just-in-time compiler to generate code that skips certain run-time checks.

```
(compile-allow-set!-undefined) → boolean?
(compile-allow-set!-undefined allow?) → void?
  allow? : any/c
```

A parameter that determines how a `set!` expression is compiled when it mutates a global variable. If the value of this parameter is a true value, `set!` expressions for global

variables are compiled so that the global variable is set even if it was not previously defined. Otherwise, `set!` expressions for global variables are compiled to raise the `exn:fail:contract:variable` exception if the global variable is not defined at the time the `set!` is performed. Note that this parameter is used when an expression is *compiled*, not when it is *evaluated*.

```
(compile-context-preservation-enabled) → boolean?
(compile-context-preservation-enabled on?) → void?
  on? : any/c
```

A parameter that determines whether compilation should avoid function-call inlining and other optimizations that may cause information to be lost from stack traces (as reported by `continuation-mark-set->context`). The default is `#f`, which allows such optimizations.

```
(eval-jit-enabled) → boolean?
(eval-jit-enabled on?) → void?
  on? : any/c
```

See also §19.2 "The Bytecode and Just-in-Time (JIT) Compilers" in ***The Racket Guide***.

A parameter that determines whether the native-code just-in-time compiler (*JIT*) is enabled for code (compiled or not) that is passed to the default evaluation handler. A true parameter value is effective only on platforms for which the JIT is supported, and changing the value from its initial setting affects only forms that are outside of `module`.

The default is `#t`, unless the JIT is not supported by the current platform, unless it is disabled through the `-j`/`--no-jit` command-line flag to stand-alone Racket (or GRacket), and unless it is disabled through the `PLTNOMZJIT` environment variable (set to any value).

```
(load-on-demand-enabled) → boolean?
(load-on-demand-enabled on?) → void?
  on? : any/c
```

A parameter that determines whether the default load handler sets `read-on-demand-source`. See `current-load` for more information. The default is `#t`, unless it is disabled through the `-d`/`--no-delay` command-line flag.

14.3 The `racket/load` Language

`#lang racket/load` package: `base`

The `racket/load` language supports evaluation where each top-level form in the module body is separately passed to `eval` in the same way as for `load`.

The namespace for evaluation shares the module registry with the `racket/load` module instance, but it has a separate top-level environment, and it is initialized with the bindings of `racket`. A single namespace is created for each instance of the `racket/load`

module (i.e., multiple modules using the racket/load language share a namespace). The racket/load library exports only #%module-begin and #%top-interaction forms that effectively swap in the evaluation namespace and call eval.

For example, the body of a module using racket/load can include module forms, so that running the following module prints 5:

```
#lang racket/load

(module m racket/base
  (provide x)
  (define x 5))

(module n racket/base
  (require 'm)
  (display x))

(require 'n)
```

Definitions in a module using racket/load are evaluated in the current namespace, which means that load and eval can see the definitions. For example, running the following module prints 6:

```
#lang racket/load

(define x 6)
(display (eval 'x))
```

Since all forms within a racket/load module are evaluated in the top level, bindings cannot be exported from the module using provide. Similarly, since evaluation of the module-body forms is inherently dynamic, compilation of the module provides essentially no benefit. For these reasons, use racket/load for interactive exploration of top-level forms only, and not for constructing larger programs.

14.4 Module Names and Loading

14.4.1 Resolving Module Names

The name of a declared module is represented by a *resolved module path*, which encapsulates either a symbol or a complete filesystem path (see §15.1 “Paths”). A symbol normally refers to a predefined module or module declared through reflective evaluation (e.g., eval). A filesystem path normally refers to a module declaration that was loaded on demand via require or other forms.

The syntax/modresolve library provides additional operations for resolving and manipulating module names.

A *module path* is a datum that matches the grammar for *module-path* for `require`. A module path is relative to another module.

```
(resolved-module-path? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a resolved module path, `#f` otherwise.

```
(make-resolved-module-path path) → resolved-module-path?
        (or/c symbol?
              (and/c path? complete-path?)
  path :      (cons/c (or/c symbol?
                              (and/c path? complete-path?))
                      (non-empty-listof symbol?)))
```

Returns a resolved module path that encapsulates *path*, where a list *path* corresponds to a submodule path. If *path* is a path or starts with a path, the path normally should be cleansed (see `cleanse-path`) and simplified (see `simplify-path`).

A resolved module path is interned. That is, if two resolved module path values encapsulate paths that are `equal?`, then the resolved module path values are `eq?`.

```
(resolved-module-path-name module-path)
    (or/c symbol?
          (and/c path? complete-path?)
 →        (cons/c (or/c symbol?
                        (and/c path? complete-path?))
                 (non-empty-listof symbol?)))
  module-path : resolved-module-path?
```

Returns the path or symbol encapsulated by a resolved module path. A list result corresponds to a submodule path.

```
(module-path? v) → boolean?
  v : any/c
```

Returns `#t` if *v* corresponds to a datum that matches the grammar for *module-path* for `require`, `#f` otherwise. Note that a path (in the sense of `path?`) is a module path.

```
(current-module-name-resolver)
    (case->
     (resolved-module-path? (or/c #f namespace?) . -> . any)
     (module-path?
      (or/c #f resolved-module-path?)
 →    (or/c #f syntax?)
      boolean?
      . -> .
      resolved-module-path?))
```

```
(current-module-name-resolver proc) → void?
          (case->
           (resolved-module-path? (or/c #f namespace?) . -> . any)
           (module-path?
            (or/c #f resolved-module-path?)
 proc :     (or/c #f syntax?)
            boolean?
            . -> .
            resolved-module-path?))
```

A parameter that determines the current *module name resolver*, which manages the conversion from other kinds of module references to a resolved module path. For example, when the expander encounters `(require module-path)` where *module-path* is not an identifier, then the expander passes `'module-path` to the module name resolver to obtain a symbol or resolved module path. When such a `require` appears within a module, the *module path resolver* is also given the name of the enclosing module, so that a relative reference can be converted to an absolute symbol or resolved module path.

The default module name resolver uses `collection-file-path` to convert `lib` and symbolic-shorthand module paths to filesystem paths. The `collection-file-path` function, in turn, uses the `current-library-collection-links` and `current-library-collection-paths` parameters.

A module name resolver takes two and four arguments:

- When given two arguments, the first is a name for a module that is now declared in the current namespace, and the second is optionally a namespace from which the declaration was copied. The module name resolver's result in this case is ignored.

  The current module name resolver is called with two arguments by `namespace-attach-module` or `namespace-attach-module-declaration` to notify the resolver that a module declaration was attached to the current namespace (and should not be loaded in the future for the namespace's module registry). Evaluation of a module declaration also calls the current module name resolver with two arguments, where the first is the declared module and the second is `#f`. No other Racket operation invokes the module name resolver with two arguments, but other tools (such as DrRacket) might call this resolver in this mode to avoid redundant module loads.

- When given four arguments, the first is a module path, equivalent to a quoted *module-path* for `require`. The second is name for the source module, if any, to which the path is relative; if the second argument is `#f`, the module path is relative to `(or (current-load-relative-directory) (current-directory))`. The third argument is a syntax object that can be used for error reporting, if it is not `#f`. If the last argument is `#t`, then the module declaration should be loaded (if it is not already), otherwise the module path should be simply resolved to a name. The result is the resolved name.

For the second case, the standard module name resolver keeps a table per module registry containing loaded module name. If a resolved module path is not in the table, and `#f` is not provided as the fourth argument to the module name resolver, then the name is put into the table and the corresponding file is loaded with a variant of `load/use-compiled` that passes the expected module name to the compiled-load handler.

While loading a file, the default module name resolver sets the `current-module-declare-name` parameter to the resolved module name (while the compiled-load handler sets `current-module-declare-source`). Also, the default module name resolver records in a private continuation mark the module being loaded, and it checks whether such a mark already exists; if such a continuation mark does exist in the current continuation, then the `exn:fail` exception is raised with a message about a dependency cycle.

The default module name resolver cooperates with the default compiled-load handler: on a module-attach notification, bytecode-file information recorded by the compiled-load handler for the source namespace's module registry is transferred to the target namespace's module registry.

The default module name resolver also maintains a small, module registry-specific cache that maps `lib` and symbolic module paths to their resolutions. This cache is consulted before checking parameters such as `current-library-collection-links` and `current-library-collection-paths`, so results may "stick" even if those parameter values change. An entry is added to the cache only when the fourth argument to the module name resolver is true (indicating that a module should be loaded) and only when loading succeeds.

Module loading is suppressed (i.e., `#f` is supplied as a fourth argument to the module name resolver) when resolving module paths in syntax objects (see §1.2.2 "Syntax Objects"). When a syntax object is manipulated, the current namespace might not match the original namespace for the syntax object, and the module should not necessarily be loaded in the current namespace.

For historical reasons, the default module name resolver currently accepts three arguments, in addition to two and four. Three arguments are treated the same as four arguments with the fourth argument as `#t`, except that an error is also logged. Support for three arguments will be removed in a future version.

Changed in version 6.0.1.12 of package `base`: Added error logging to the default module name resolver when called with three arguments.

```
(current-module-declare-name)
 → (or/c resolved-module-path? #f)
(current-module-declare-name name) → void?
  name : (or/c resolved-module-path? #f)
```

A parameter that determines a module name that is used when evaluating a `module` declaration (when the parameter value is not `#f`). In that case, the *id* from the `module` declaration is ignored, and the parameter's value is used as the name of the declared module.

When declaring submodules, current-module-declare-name determines the name used for the submodule's root module, while its submodule path relative to the root module is unaffected.

```
(current-module-declare-source)
 → (or/c symbol? (and/c path? complete-path?) #f)
(current-module-declare-source src) → void?
  src : (or/c symbol? (and/c path? complete-path?) #f)
```

A parameter that determines source information to be associated with a module when evaluating a module declaration. Source information is used in error messages and reflected by variable-reference->module-source. When the parameter value is #f, the module's name (as determined by current-module-declare-name) is used as the source name instead of the parameter value.

```
(current-module-path-for-load)
    (or/c #f module-path?
          (and/c syntax?
 →                (lambda (stx)
                   (module-path? (syntax->datum s)))))
(current-module-path-for-load path) → void?
         (or/c #f module-path?
               (and/c syntax?
  path :               (lambda (stx)
                        (module-path? (syntax->datum s)))))
```

A parameter that determines a module path used for exn:fail:syntax:missing-module and exn:fail:filesystem:missing-module exceptions as raised by the default load handler. The parameter is normally set by a module name resolver.

14.4.2 Compiled Modules and References

While expanding a module declaration, the expander resolves module paths for imports to load module declarations as necessary and to determine imported bindings, but the compiled form of a module declaration preserves the original module path. Consequently, a compiled module can be moved to another filesystem, where the module name resolver can resolve inter-module references among compiled code.

When a module reference is extracted from compiled form (see module-compiled-imports) or from syntax objects in macro expansion (see §12.2 "Syntax Object Content"), the module reference is reported in the form of a *module path index*. A module path index is a semi-interned (multiple references to the same relative module tend to use the same module path index value, but not always) opaque value that encodes a module path (see module-path?) and either a resolved module path or another module path index to which it is relative.

A module path index that uses both #f for its path and base module path index represents "self"—i.e., the module declaration that was the source of the module path index—and such a module path index can be used as the root for a chain of module path indexes at compile time. For example, when extracting information about an identifier's binding within a module, if the identifier is bound by a definition within the same module, the identifier's source module is reported using the "self" module path index. If the identifier is instead defined in a module that is imported via a module path (as opposed to a literal module name), then the identifier's source module will be reported using a module path index that contains the `required` module path and the "self" module path index. A "self" module path index has a submodule path when the module that it refers to is a submodule.

A module path index has state. When it is *resolved* to a resolved module path, then the resolved module path is stored with the module path index. In particular, when a module is loaded, its root module path index is resolved to match the module's declaration-time name. This resolved path is forgotten, however, in identifiers that the module contributes to the compiled and marshaled form of other modules. The transient nature of resolved names allows the module code to be loaded with a different resolved name than the name when it was compiled.

Two module path index values are equal? when they have equal? path and base values (even if they have different resolved values).

```
(module-path-index? v) → boolean?
  v : any/c
```

Returns #t if *v* is a module path index, #f otherwise.

```
(module-path-index-resolve mpi) → resolved-module-path?
  mpi : module-path-index?
```

Returns a resolved module path for the resolved module name, computing the resolved name (and storing it in *mpi*) if it has not been computed before.

Resolving a module path index uses the current module name resolver (see current-module-name-resolver). Depending on the kind of module paths encapsulated by *mpi*, the computed resolved name can depend on the value of current-load-relative-directory or current-directory.

```
(module-path-index-split mpi)
  → (or/c module-path? #f)
    (or/c module-path-index? resolved-module-path? #f)
  mpi : module-path-index?
```

Returns two values: a module path, and a base path—either a module path index, resolved module path, or #f—to which the first path is relative.

A #f second result means that the path is relative to an unspecified directory (i.e., its resolution depends on the value of current-load-relative-directory and/or current-directory).

A #f for the first result implies a #f for the second result, and means that *mpi* represents "self" (see above). Such a module path index may have a non-#f submodule path as reported by module-path-index-submodule.

```
(module-path-index-submodule mpi)
 → (or/c #f (non-empty-listof symbol?))
  mpi : module-path-index?
```

Returns a non-empty list of symbols if *mpi* is a "self" (see above) module path index that refers to a submodule. The result is always #f if either result of (module-path-index-split *mpi*) is non-#f.

```
(module-path-index-join path base [submod]) → module-path-index?
  path : (or/c module-path? #f)
  base : (or/c module-path-index? resolved-module-path? #f)
  submod : (or/c #f (non-empty-listof symbol?)) = #f
```

Combines *path*, *base*, and *submod* to create a new module path index. The *path* argument can #f only if *base* is also #f. The *submod* argument can be a list only when *path* and *base* are both #f.

```
(compiled-module-expression? v) → boolean?
  v : any/c
```

Returns #t if *v* is a compiled module declaration, #f otherwise. See also current-compile.

```
(module-compiled-name compiled-module-code)
 → (or/c symbol? (cons/c symbol? (non-empty-listof symbol?)))
  compiled-module-code : compiled-module-expression?
(module-compiled-name compiled-module-code
                      name)
 → compiled-module-expression?
  compiled-module-code : compiled-module-expression?
  name : (or/c symbol? (cons/c symbol? (non-empty-listof symbol?)))
```

Takes a module declaration in compiled form and either gets the module's declared name (when *name* is not provided) or returns a revised module declaration with the given *name*.

The name is a symbol for a top-level module, and it list of symbols for a submodule, where a list reflects the submodule path to the module starting with the top-level module's declared name.

```
(module-compiled-submodules compiled-module-code
                            non-star?)
 → (listof compiled-module-expression?)
  compiled-module-code : compiled-module-expression?
  non-star? : any/c
(module-compiled-submodules compiled-module-code
                            non-star?
                            submodules)
 → compiled-module-expression?
  compiled-module-code : compiled-module-expression?
  non-star? : any/c
  submodules : (listof compiled-module-expression?)
```

Takes a module declaration in compiled form and either gets the module's submodules (when *submodules* is not provided) or returns a revised module declaration with the given *submodules*. The `pre-module?` argument determines whether the result or new submodule list corresponds to `module` declarations (when *non-star?* is true) or `module*` declarations (when *non-star?* is `#f`).

```
(module-compiled-imports compiled-module-code)
    (listof (cons/c (or/c exact-integer? #f)
 →                  (listof module-path-index?)))
  compiled-module-code : compiled-module-expression?
```

Takes a module declaration in compiled form and returns an association list mapping phase level shifts (where `#f` corresponds to a shift into the label phase level) to module references for the module's explicit imports.

```
(module-compiled-exports compiled-module-code)
    (listof (cons/c (or/c exact-integer? #f) list?))
 →  (listof (cons/c (or/c exact-integer? #f) list?))
  compiled-module-code : compiled-module-expression?
```

Returns two association lists mapping phase level values (where `#f` corresponds to the label phase level) to exports at the corresponding phase. The first association list is for exported variables, and the second is for exported syntax. Beware however, that value bindings re-exported though a rename transformer are in the syntax list instead of the value list.

Each associated list, which is represented by `list?` in the result contracts above, more precisely matches the contract

```
(listof (list/c symbol?
                (listof
                 (or/c module-path-index?
```

```
                    (list/c module-path-index?
                            (or/c exact-integer? #f)
                            symbol?
                            (or/c exact-integer? #f))))))
```

For each element of the list, the leading symbol is the name of the export.

The second part—the list of module path index values, etc.—describes the origin of the exported identifier. If the origin list is `null`, then the exported identifier is defined in the module. If the exported identifier is re-exported, instead, then the origin list provides information on the import that was re-exported. The origin list has more than one element if the binding was imported multiple times from (possibly) different sources.

For each origin, a module path index by itself means that the binding was imported with a phase level shift of `0` (i.e., a plain `require` without `for-meta`, `for-syntax`, etc.), and imported identifier has the same name as the re-exported name. An origin represented with a list indicates explicitly the import, the import phase level shift (where `#f` corresponds to a `for-label` import), the import name of the re-exported binding, and the phase level of the import.}

```
(module-compiled-language-info compiled-module-code)
 → (or/c #f (vector/c module-path? symbol? any/c))
  compiled-module-code : compiled-module-expression?
```

See also §17.3.6 "Module-Handling Configuration" in ***The Racket Guide***.

Returns information intended to reflect the "language" of the module's implementation as originally attached to the syntax of the module's declaration though the `'module-language` syntax property. See also `module`.

If no information is available for the module, the result is `#f`. Otherwise, the result is `(vector mp name val)` such that `((dynamic-require mp name) val)` should return function that takes two arguments. The function's arguments are a key for reflected information and a default value. Acceptable keys and the interpretation of results is up to external tools, such as DrRacket. If no information is available for a given key, the result should be the given default value.

See also `module->language-info` and `racket/language-info`.

```
(module-compiled-cross-phase-persistent? compiled-module-code)
 → boolean?
  compiled-module-code : compiled-module-expression?
```

Return `#t` if *compiled-module-code* represents a cross-phase persistent module, `#f` otherwise.

14.4.3 Dynamic Module Access

```
(dynamic-require mod provided [fail-thunk]) → any
        (or/c module-path?
  mod :       resolved-module-path?
              module-path-index?)
  provided : (or/c symbol? #f 0 void?)
  fail-thunk : (-> any) = (lambda () ....)
```

Dynamically instantiates the module specified by *mod* in the current namespace's registry at the namespace's base phase, if it is not yet instantiated. The current module name resolver may load a module declaration to resolve *mod* (see current-module-name-resolver); the path is resolved relative to current-load-relative-directory and/or current-directory.

If *provided* is #f, then the result is #<void>, and the module is not visited (see §1.2.3.8 "Module Phases and Visits") or even made available (for on-demand visits) in phases above the base phase.

When *provided* is a symbol, the value of the module's export with the given name is returned, and still the module is not visited or made available in higher phases. If the module exports *provided* as syntax, then a use of the binding is expanded and evaluated in a fresh namespace to which the module is attached, which means that the module is visited in the fresh namespace. If the module has no such exported variable or syntax, then *fail-thunk* is called; the default *fail-thunk* raises exn:fail:contract. If the variable named by *provided* is exported protected (see §14.10 "Code Inspectors"), then the exn:fail:contract exception is raised.

If *provided* is 0, then the module is instantiated but not visited, the same as when *provided* is #f. With 0, however, the module is made available in higher phases.

If *provided* is #<void>, then the module is visited but not instantiated (see §1.2.3.8 "Module Phases and Visits"), and the result is #<void>.

```
(dynamic-require-for-syntax  mod
                             provided
                            [fail-thunk]) → any
  mod : module-path?
  provided : (or/c symbol? #f)
  fail-thunk : (-> any) = (lambda () ....)
```

Like dynamic-require, but in a phase that is 1 more than the namespace's base phase.

```
(module-declared? mod [load?]) → boolean?
        (or/c module-path? module-path-index?
  mod :       resolved-module-path?)
  load? : any/c = #f
```

Returns `#t` if the module indicated by *mod* is declared (but not necessarily instantiated or visited) in the current namespace, `#f` otherwise.

If *load?* is `#t` and *mod* is not a resolved module path, the module is loaded in the process of resolving *mod* (as for `dynamic-require` and other functions). Checking for the declaration of a submodule does not trigger an exception if the submodule cannot be loaded because it does not exist, either within a root module that does exist or because the root module does not exist.

```
(module->language-info mod [load?])
 → (or/c #f (vector/c module-path? symbol? any/c))
  mod : (or/c module-path? module-path-index?
              resolved-module-path?)
  load? : any/c = #f
```

Returns information intended to reflect the "language" of the implementation of *mod*. If *mod* is a resolved module path or *load?* is `#f`, the module named by *mod* must be declared (but not necessarily instantiated or visited) in the current namespace; otherwise, *mod* may be loaded (as for `dynamic-require` and other functions). The information returned by `module->language-info` is the same as would have been returned by `module-compiled-language-info` applied to the module's implementation as compiled code.

```
(module->imports mod)
 → (listof (cons/c (or/c exact-integer? #f)
                   (listof module-path-index?)))
  mod : (or/c module-path? module-path-index?
              resolved-module-path?)
```

Like `module-compiled-imports`, but produces the imports of *mod*, which must be declared (but not necessarily instantiated or visited) in the current namespace.

```
(module->exports mod)
 → (listof (cons/c (or/c exact-integer? #f) list?))
   (listof (cons/c (or/c exact-integer? #f) list?))
  mod : (or/c module-path? resolved-module-path?)
```

Like `module-compiled-exports`, but produces the exports of *mod*, which must be declared (but not necessarily instantiated or visited) in the current namespace.

```
(module-predefined? mod) → boolean?
  mod : (or/c module-path? resolved-module-path?)
```

Reports whether *mod* refers to a module that is predefined for the running Racket instance. Predefined modules always have a symbolic resolved module path, and they may be predefined always or specifically within a particular executable (such as one created by `raco exe` or `create-embedding-executable`).

14.5 Impersonators and Chaperones

An *impersonator* is a wrapper for a value where the wrapper redirects some of the value's operations. Impersonators apply only to procedures, structures for which an accessor or mutator is available, structure types, hash tables, vectors, boxes, channels, and prompt tags. An impersonator is `equal?` to the original value, but not `eq?` to the original value.

A *chaperone* is a kind of impersonator whose refinement of a value's operation is restricted to side effects (including, in particular, raising an exception) or chaperoning values supplied to or produced by the operation. For example, a vector chaperone can redirect `vector-ref` to raise an exception if the accessed vector slot contains a string, or it can cause the result of `vector-ref` to be a chaperoned variant of the value that is in the accessed vector slot, but it cannot redirect `vector-ref` to produce a value that is arbitrarily different from the value in the vector slot.

A non-chaperone impersonator, in contrast, can refine an operation to swap one value for any other. An impersonator cannot be applied to an immutable value or refine the access to an immutable field in an instance of a structure type, since arbitrary redirection of an operation amounts to mutation of the impersonated value.

Beware that each of the following operations can be redirected to an arbitrary procedure through an impersonator on the operation's argument—assuming that the operation is available to the creator of the impersonator:

- a structure-field accessor
- a structure-field mutator
- a structure type property accessor
- application of a procedure
- `unbox`
- `set-box!`
- `vector-ref`
- `vector-set!`
- `hash-ref`
- `hash-set`
- `hash-set!`
- `hash-remove`
- `hash-remove!`

- channel-get
- channel-put
- call-with-continuation-prompt
- abort-current-continuation

Derived operations, such as printing a value, can be redirected through impersonators due to their use of accessor functions. The equal?, equal-hash-code, and equal-secondary-hash-code operations, in contrast, may bypass impersonators (but they are not obliged to).

In addition to redirecting operations that work on a value, a impersonator can include *impersonator properties* for an impersonated value. An impersonator property is similar to a structure type property, but it applies to impersonators instead of structure types and their instances.

```
(impersonator? v) → boolean?
  v : any/c
```

Returns #t if v is an impersonator, #f otherwise.

Programs and libraries generally should avoid impersonator? and treat impersonators the same as non-impersonator values. In rare cases, impersonator? may be needed to guard against redirection by an impersonator of an operation to an arbitrary procedure.

```
(chaperone? v) → boolean?
  v : any/c
```

Returns #t if v is a chaperone, #f otherwise.

Programs and libraries generally should avoid chaperone? for the same reason that they should avoid impersonator?.

```
(impersonator-of? v1 v2) → boolean?
  v1 : any/c
  v2 : any/c
```

Indicates whether v1 can be considered equivalent modulo impersonators to v2.

For values that include no impersonators, v1 and v2 can be considered impersonators of each other if they are equal?.

Otherwise, impersonators within v2 must be intact within v1:

- If a part of v2 is an impersonator created from one of the impersonator constructors (e.g., impersonate-procedure or chaperone-procedure), and if the impersonator is constructed with at least one redirection procedure (i.e., a value other than #f

was supplied for a redirection procedure), then the corresponding part of *v1* must be one of the following:

  - the same value that is a part of *v2*;
  - a value further derived from the part of *v2* value using an impersonator constructor; or
  - a value with the `prop:impersonator-of` property whose procedure produces an impersonator of the value that is a part of *v2*.

- If a part of *v2* is a structure or procedure impersonator that was created with no redirection procedures (i.e, `#f` in place of all redirection procedures for specified operations), then the impersonated value is considered in place of that part of *v2*. In other words, an impersonator construction that does not redirect any access or mutation (but that includes some impersonator properties) need not be preserved in *v1*.

```
(chaperone-of? v1 v2) → boolean?
  v1 : any/c
  v2 : any/c
```

Indicates whether *v1* can be considered equivalent modulo chaperones to *v2*.

For values that include no chaperones, *v1* and *v2* can be considered chaperones of each other if they are `equal?`, except that the mutability of vectors and boxes with *v1* and *v2* must be the same.

Otherwise, chaperones within *v2* must be intact within *v1* analogous to way that `impersonator-of?` requires that impersonators are preserved, except that `prop:impersonator-of` has no analog for `chaperone-of?`.

```
(impersonator-ephemeron v) → ephemeron?
  v : any/c
```

Produces an ephemeron that can be used to connect the reachability of *v* (in the sense of garbage collection; see §1.1.7 “Garbage Collection”) with the reachability of any value for which *v* is an impersonator. That is, the value *v* will be considered reachable as long as the result ephemeron is reachable in addition to any value that *v* impersonates (including itself).

14.5.1 Impersonator Constructors

```
(impersonate-procedure proc
                       wrapper-proc
                       prop
                       prop-val ...
                       ...)
```

```
→ (and/c procedure? impersonator?)
 proc : procedure?
 wrapper-proc : (or/c procedure? #f)
 prop : impersonator-property?
 prop-val : any
```

Returns an impersonator procedure that has the same arity, name, and other attributes as *proc*. When the impersonator procedure is applied, the arguments are first passed to *wrapper-proc* (when it is not `#f`), and then the results from *wrapper-proc* are passed to *proc*. The *wrapper-proc* can also supply a procedure that processes the results of *proc*.

The arity of *wrapper-proc* must include the arity of *proc*. The allowed keyword arguments of *wrapper-proc* must be a superset of the allowed keywords of *proc*. The required keyword arguments of *wrapper-proc* must be a subset of the required keywords of *proc*.

For applications without keywords, the result of *wrapper-proc* must be either the same number of values as supplied to it or one more than the number of supplied values, where an extra result is supplied before the others. The additional result, if any, must be a procedure that accepts as many results as produced by *proc*; it must return the same number of results. If *wrapper-proc* returns the same number of values as it is given (i.e., it does not return a procedure to impersonator *proc*'s result), then *proc* is called in tail position with respect to the call to the impersonator.

For applications that include keyword arguments, *wrapper-proc* must return an additional value before any other values but after the result-impersonating procedure (if any). The additional value must be a list of replacements for the keyword arguments that were supplied to the impersonator (i.e., not counting optional arguments that were not supplied). The arguments must be ordered according to the sorted order of the supplied arguments' keywords.

If `wrapper` is `#f`, then applying the resulting impersonator is the same as applying *proc*. If `wrapper` is `#f` and no *prop* is provided, then the result is *proc* unimpersonated.

Pairs of *prop* and *prop-val* (the number of arguments to `procedure-impersonator` must be even) add impersonator properties or override impersonator-property values of *proc*.

If any *prop* is `impersonator-prop:application-mark` and if the associated *prop-val* is a pair, then the call to *proc* is wrapped with `with-continuation-mark` using `(car` *prop-val*`)` as the mark key and `(cdr` *prop-val*`)` as the mark value. In addition, if `continuation-mark-set-first` with `(car` *prop-val*`)` produces a value for the immediate continuation frame of the call to the impersonated procedure, the value is also installed as an immediate value for `(car` *prop-val*`)` as a mark during the call to *wrapper-proc* (which allows tail-calls of impersonators with respect to wrapping impersonators to be detected within *wrapper-proc*).

```
(impersonate-struct v
                    orig-proc
                    redirect-proc ...
                    ...
                    prop
                    prop-val ...
                    ...)              → any/c
  v : any/c
  orig-proc :       (or/c struct-accessor-procedure?
                          struct-mutator-procedure?
                          struct-type-property-accessor-procedure?)
  redirect-proc : (or/c procedure? #f)
  prop : impersonator-property?
  prop-val : any
```

Returns an impersonator of *v*, which redirects certain operations on the impersonated value. The *orig-proc*s indicate the operations to redirect, and the corresponding *redirect-proc*s supply the redirections.

The protocol for a *redirect-proc* depends on the corresponding *orig-proc*:

- A structure-field accessor: *redirect-proc* must accept two arguments, *v* and the value *field-v* that *orig-proc* produces for *v*; it must return a replacement for *field-v*. The corresponding field must not be immutable, and either the field's structure type must be accessible via the current inspector or one of the other *orig-proc*s must be a structure-field mutator for the same field.
- A structure-field mutator: *redirect-proc* must accept two arguments, *v* and the value *field-v* supplied to the mutator; it must return a replacement for *field-v* to be propagated to *orig-proc* and *v*.
- A property accessor: *redirect-proc* uses the same protocol as for a structure-field accessor. The accessor's property must have been created with `'can-impersonate` as the second argument to `make-struct-type-property`.

When a *redirect-proc* is `#f`, the corresponding *orig-proc* is unaffected. Supplying `#f` for a *redirect-proc* is useful to allow its *orig-proc* to act as a "witness" of *v*'s representation and enable the addition of *prop*s.

Pairs of *prop* and *prop-val* (the number of arguments to `impersonate-struct` must be odd) add impersonator properties or override impersonator-property values of *v*.

Each *orig-proc* must indicate a distinct operation. If no *orig-proc*s are supplied, then no *prop*s must be supplied. If *orig-proc*s are supplied only with `#f` *redirect-proc*s and no *prop*s are supplied, then *v* is returned unimpersonated.

If any *orig-proc* is itself an impersonator, then a use of the accessor or mutator that *orig-proc* impersonates is redirected for the resulting impersonated structure to use *orig-proc* on *v* before *redirect-proc* (in the case of accessor) or after *redirect-proc* (in the case of a mutator).

```
(impersonate-vector vec
                    ref-proc
                    set-proc
                    prop
                    prop-val ...
                    ...)
 → (and/c vector? impersonator?)
  vec : (and/c vector? (not/c immutable?))
  ref-proc : (vector? exact-nonnegative-integer? any/c . -> . any/c)
  set-proc : (vector? exact-nonnegative-integer? any/c . -> . any/c)
  prop : impersonator-property?
  prop-val : any
```

Returns an impersonator of *vec*, which redirects the `vector-ref` and `vector-set!` operations.

The *ref-proc* must accept *vec*, an index passed to `vector-ref`, and the value that `vector-ref` on *vec* produces for the given index; it must produce a replacement for the value, which is the result of `vector-ref` on the impersonator.

The *set-proc* must accept *vec*, an index passed to `vector-set!`, and the value passed to `vector-set!`; it must produce a replacement for the value, which is used with `vector-set!` on the original *vec* to install the value.

Pairs of *prop* and *prop-val* (the number of arguments to `impersonate-vector` must be odd) add impersonator properties or override impersonator-property values of *vec*.

```
(impersonate-box box
                 unbox-proc
                 set-proc
                 prop
                 prop-val ...
                 ...)          → (and/c box? impersonator?)
  box : (and/c box? (not/c immutable?))
  unbox-proc : (box? any/c . -> . any/c)
  set-proc : (box? any/c . -> . any/c)
  prop : impersonator-property?
  prop-val : any
```

Returns an impersonator of *box*, which redirects the `unbox` and `set-box!` operations.

The *unbox-proc* must accept *box* and the value that unbox produces on *box*; it must produce a replacement value, which is the result of unbox on the impersonator.

The *set-proc* must accept *box* and the value passed to set-box!; it must produce a replacement value, which is used with set-box! on the original *box* to install the value.

Pairs of *prop* and *prop-val* (the number of arguments to impersonate-box must be odd) add impersonator properties or override impersonator-property values of *box*.

```
(impersonate-hash  hash
                   ref-proc
                   set-proc
                   remove-proc
                   key-proc
                  [clear-proc]
                   prop
                   prop-val ...
                   ...)          → (and/c hash? impersonator?)
  hash : (and/c hash? (not/c immutable?))
             (hash? any/c . -> . (values
  ref-proc :                      any/c
                                  (hash? any/c any/c . -> . any/c)))
  set-proc : (hash? any/c any/c . -> . (values any/c any/c))
  remove-proc : (hash? any/c . -> . any/c)
  key-proc : (hash? any/c . -> . any/c)
  clear-proc : (or/c #f (hash? . -> . any)) = #f
  prop : impersonator-property?
  prop-val : any
```

Returns an impersonator of *hash*, which redirects the hash-ref, hash-set! or hash-set (as applicable), hash-remove or hash-remove! (as applicable), hash-clear or hash-clear! (as applicable and if *clear-proc* is not #f) operations. When hash-set, hash-remove or hash-clear is used on an impersonator of a hash table, the result is an impersonator with the same redirecting procedures. In addition, operations like hash-iterate-key or hash-map, which extract keys from the table, use *key-proc* to filter keys extracted from the table. Operations like hash-iterate-value or hash-iterate-map implicitly use hash-ref and therefore redirect through *ref-proc*.

The *ref-proc* must accept *hash* and a key passed to hash-ref. It must return a replacement key as well as a procedure. The returned procedure is called only if the returned key is found in *hash* via hash-ref, in which case the procedure is called with *hash*, the previously returned key, and the found value. The returned procedure must itself return a replacement for the found value.

The *set-proc* must accept *hash*, a key passed to hash-set! or hash-set, and the value passed to hash-set! or hash-set; it must produce two values: a replacement for the key

and a replacement for the value. The returned key and value are used with hash-set! or hash-set on the original hash to install the value.

The remove-proc must accept hash and a key passed to hash-remove! or hash-remove; it must produce the a replacement for the key, which is used with hash-remove! or hash-remove on the original hash to remove any mapping using the (impersonator-replaced) key.

The key-proc must accept hash and a key that has been extracted from hash (by hash-iterate-key or other operations that use hash-iterate-key internally); it must produce a replacement for the key, which is then reported as a key extracted from the table.

If clear-proc is not #f, it must accept hash as an argument, and its result is ignored. The fact that clear-proc returns (as opposed to raising an exception or otherwise escaping) grants the capability to remove all keys from hash. If clear-proc is #f, then hash-clear or hash-clear! on the impersonator is implemented using hash-iterate-key and hash-remove or hash-remove!.

The hash-iterate-value, hash-map, or hash-for-each functions use a combination of hash-iterate-key and hash-ref. If a key produced by key-proc does not yield a value through hash-ref, then the exn:fail:contract exception is raised.

Pairs of prop and prop-val (the number of arguments to impersonate-hash must be odd) add impersonator properties or override impersonator-property values of hash.

```
(impersonate-channel channel
                     get-proc
                     put-proc
                     prop
                     prop-val ...
                     ...)
 → (and/c channel? impersonator?)
  channel : channel?
  get-proc : (channel? . -> . (values channel? (any/c . -> . any/c)))
  put-proc : (channel? any/c . -> . any/c)
  prop : impersonator-property?
  prop-val : any
```

Returns an impersonator of channel, which redirects the channel-get and channel-put operations.

The get-proc generator is called on channel-get or any other operation that fetches results from the channel (such as a sync on the channel). The get-proc must return two values: a channel that is an impersonator of channel, and a procedure that is used to check the channel's contents.

The put-proc must accept channel and the value passed to channel-put; it must produce a replacement value, which is used with channel-put on the original channel to

send the value over the channel.

Pairs of *prop* and *prop-val* (the number of arguments to impersonate-channel must be odd) add impersonator properties or override impersonator-property values of *channel*.

```
(impersonate-prompt-tag prompt-tag
                        handle-proc
                        abort-proc
                        [cc-guard-proc
                         callcc-impersonate-proc]
                        prop
                        prop-val ...
                        ...)
 → (and/c continuation-prompt-tag? impersonator?)
  prompt-tag : continuation-prompt-tag?
  handle-proc : procedure?
  abort-proc : procedure?
  cc-guard-proc : procedure? = values
  callcc-impersonate-proc : (procedure? . -> . procedure?)
                          = (lambda (p) p)
  prop : impersonator-property?
  prop-val : any
```

Returns an impersonator of *prompt-tag*, which redirects the call-with-continuation-prompt and abort-current-continuation operations.

The *handle-proc* must accept the values that the handler of a continuation prompt would take and it must produce replacement values, which will be passed to the handler.

The *abort-proc* must accept the values passed to abort-current-continuation; it must produce replacement values, which are aborted to the appropriate prompt.

The *cc-guard-proc* must accept the values produced by call-with-continuation-prompt in the case that a non-composable continuation is applied to replace the continuation that is delimited by the prompt, but only if abort-current-continuation is not later used to abort the continuation delimited by the prompt (in which case *abort-proc* is used).

The *callcc-impersonate-proc* must accept a procedure that guards the result of a continuation captured by call-with-current-continuation with the impersonated prompt tag. The *callcc-impersonate-proc* is applied (under a continuation barrier) when the captured continuation is applied to refine a guard function (initially values) that is specific to the delimiting prompt; this prompt-specific guard is ultimately composed with any *cc-guard-proc* that is in effect at the delimiting prompt, and it is not used in the same case that a *cc-guard-proc* is not used (i.e., when abort-current-continuation is used to abort to the prompt). In the special case where the delimiting prompt at application time is a thread's built-in initial prompt, *callcc-impersonate-proc* is ignored (partly on the grounds that the initial prompt's result is ignored).

Pairs of *prop* and *prop-val* (the number of arguments to impersonate-prompt-tag must be odd) add impersonator properties or override impersonator-property values of *prompt-tag*.

Examples:

```
> (define tag
    (impersonate-prompt-tag
     (make-continuation-prompt-tag)
     (lambda (n) (* n 2))
     (lambda (n) (+ n 1))))

> (call-with-continuation-prompt
    (lambda ()
      (abort-current-continuation tag 5))
    tag
    (lambda (n) n))
12
```

```
(impersonate-continuation-mark-key key
                                   get-proc
                                   set-proc
                                   prop
                                   prop-val ...
                                   ...)
 → (and/c continuation-mark? impersonator?)
  key : continuation-mark-key?
  get-proc : procedure?
  set-proc : procedure?
  prop : impersonator-property?
  prop-val : any
```

Returns an impersonator of *key*, which redirects with-continuation-mark and continuation mark accessors such as continuation-mark-set->list.

The *get-proc* must accept the value attached to a continuation mark and it must produce a replacement value, which will be returned by the continuation mark accessor.

The *set-proc* must accept a value passed to with-continuation-mark; it must produce a replacement value, which is attached to the continuation frame.

Pairs of *prop* and *prop-val* (the number of arguments to impersonate-prompt-tag must be odd) add impersonator properties or override impersonator-property values of *key*.

Examples:

```
> (define mark-key
```

```
    (impersonate-continuation-mark-key
     (make-continuation-mark-key)
     (lambda (l) (map char-upcase l))
     (lambda (s) (string->list s))))

> (with-continuation-mark mark-key "quiche"
    (continuation-mark-set-first
     (current-continuation-marks)
     mark-key))
'(#\Q #\U #\I #\C #\H #\E)
```

```
prop:impersonator-of : struct-type-property?
```

A structure type property (see §5.3 “Structure Type Properties”) that supplies a procedure for extracting an impersonated value from a structure that represents an impersonator. The property is used for `impersonator-of?` as well as `equal?`.

The property value must be a procedure of one argument, which is a structure whose structure type has the property. The result can be `#f` to indicate the structure does not represent an impersonator, otherwise the result is a value for which the original structure is an impersonator (so the original structure is an `impersonator-of?` and `equal?` to the result value). The result value must have the same `prop:impersonator-of` and `prop:equal+hash` property values as the original structure, and the property values must be inherited from the same structure type (which ensures some consistency between `impersonator-of?` and `equal?`).

14.5.2 Chaperone Constructors

```
(chaperone-procedure proc
                     wrapper-proc
                     prop
                     prop-val ...
                     ...)
 → (and/c procedure? chaperone?)
  proc : procedure?
  wrapper-proc : (or/c procedure? #f)
  prop : impersonator-property?
  prop-val : any
```

Like `impersonate-procedure`, but for each value supplied to *`wrapper-proc`*, the corresponding result must be the same or a chaperone of (in the sense of `chaperone-of?`) the supplied value. The additional result, if any, that precedes the chaperoned values must be a procedure that accepts as many results as produced by *`proc`*; it must return the same number of results, each of which is the same or a chaperone of the corresponding original result.

For applications that include keyword arguments, *wrapper-proc* must return an additional value before any other values but after the result-chaperoning procedure (if any). The additional value must be a list of chaperones of the keyword arguments that were supplied to the chaperone procedure (i.e., not counting optional arguments that were not supplied). The arguments must be ordered according to the sorted order of the supplied arguments' keywords.

```
(chaperone-struct v
                  orig-proc
                  redirect-proc ...
                  ...
                  prop
                  prop-val ...
                  ...)              → any/c
  v : any/c
              (or/c struct-accessor-procedure?
                    struct-mutator-procedure?
  orig-proc :       struct-type-property-accessor-procedure?
                    (one-of/c struct-info))
  redirect-proc : (or/c procedure? #f)
  prop : impersonator-property?
  prop-val : any
```

Like `impersonate-struct`, but with the following refinements:

- With a structure-field accessor as *orig-proc*, *redirect-proc* must accept two arguments, *v* and the value *field-v* that *orig-proc* produces for *v*; it must return a chaperone of *field-v*. The corresponding field may be immutable.
- With structure-field mutator as *orig-proc*, *redirect-proc* must accept two arguments, *v* and the value *field-v* supplied to the mutator; it must return a chaperone of *field-v* to be propagated to *orig-proc* and *v*.
- A property accessor can be supplied as *orig-proc*, and the property need not have been created with `'can-impersonate`. The corresponding *redirect-proc* uses the same protocol as for a structure-field accessor.
- With `struct-info` as *orig-proc*, the corresponding *redirect-proc* must accept two values, which are the results of `struct-info` on *v*; it must return each values or a chaperone of each value. The *redirect-proc* is not called if `struct-info` would return `#f` as its first argument. An *orig-proc* can be `struct-info` only if some other *orig-proc* is supplied.
- Any accessor or mutator *orig-proc* that is an impersonator must be specifically a chaperone.

```
(chaperone-vector vec
                  ref-proc
                  set-proc
                  prop
                  prop-val ...
                  ...)          → (and/c vector? chaperone?)
  vec : vector?
  ref-proc : (vector? exact-nonnegative-integer? any/c . -> . any/c)
  set-proc : (vector? exact-nonnegative-integer? any/c . -> . any/c)
  prop : impersonator-property?
  prop-val : any
```

Like `impersonate-vector`, but with support for immutable vectors. The *ref-proc* procedure must produce the same value or a chaperone of the original value, and *set-proc* must produce the value that is given or a chaperone of the value. The *set-proc* will not be used if *vec* is immutable.

```
(chaperone-box box
               unbox-proc
               set-proc
               prop
               prop-val ...
               ...)          → (and/c box? chaperone?)
  box : box?
  unbox-proc : (box? any/c . -> . any/c)
  set-proc : (box? any/c . -> . any/c)
  prop : impersonator-property?
  prop-val : any
```

Like `impersonate-box`, but with support for immutable boxes. The *unbox-proc* procedure must produce the same value or a chaperone of the original value, and *set-proc* must produce the same value or a chaperone of the value that it is given. The *set-proc* will not be used if *box* is immutable.

```
(chaperone-hash  hash
                 ref-proc
                 set-proc
                 remove-proc
                 key-proc
                [clear-proc]
                 prop
                 prop-val ...
                 ...)          → (and/c hash? chaperone?)
  hash : hash?
```

```
            (hash? any/c . -> . (values
 ref-proc :                      any/c
                                 (hash? any/c any/c . -> . any/c)))
 set-proc : (hash? any/c any/c . -> . (values any/c any/c))
 remove-proc : (hash? any/c . -> . any/c)
 key-proc : (hash? any/c . -> . any/c)
 clear-proc : (or/c #f (hash? . -> . any)) = #f
 prop : impersonator-property?
 prop-val : any
```

Like impersonate-hash, but with constraints on the given functions and support for immutable hashes. The *ref-proc* procedure must return a found value or a chaperone of the value. The *set-proc* procedure must produce two values: the key that it is given or a chaperone of the key and the value that it is given or a chaperone of the value. The *remove-proc* and *key-proc* procedures must produce the given key or a chaperone of the key.

```
(chaperone-struct-type struct-type
                       struct-info-proc
                       make-constructor-proc
                       guard-proc
                       prop
                       prop-val ...
                       ...)
 → (and/c struct-type? chaperone?)
  struct-type : struct-type?
  struct-info-proc : procedure?
  make-constructor-proc : (procedure? . -> . procedure?)
  guard-proc : procedure?
  prop : impersonator-property?
  prop-val : any
```

Returns a chaperoned value like *struct-type*, but with struct-type-info and struct-type-make-constructor operations on the chaperoned structure type redirected. In addition, when a new structure type is created as a subtype of the chaperoned structure type, *guard-proc* is interposed as an extra guard on creation of instances of the subtype.

The *struct-info-proc* must accept 8 arguments—the result of struct-type-info on *struct-type*. It must return 8 values, where each is the same or a chaperone of the corresponding argument. The 8 values are used as the results of struct-type-info for the chaperoned structure type.

The *make-constructor-proc* must accept a single procedure argument, which is a constructor produced by struct-type-make-constructor on *struct-type*. It must return the same or a chaperone of the procedure, which is used as the result of struct-type-make-constructor on the chaperoned structure type.

The *guard-proc* must accept as many argument as a constructor for *struct-type*; it must return the same number of arguments, each the same or a chaperone of the corresponding argument. The *guard-proc* is added as a constructor guard when a subtype is created of the chaperoned structure type.

Pairs of *prop* and *prop-val* (the number of arguments to chaperone-struct-type must be even) add impersonator properties or override impersonator-property values of *struct-type*.

```
(chaperone-evt evt proc prop prop-val ... ...)
 → (and/c evt? chaperone?)
  evt : evt?
  proc : (evt? . -> . (values evt? (any/c . -> . any/c)))
  prop : impersonator-property?
  prop-val : any
```

Returns a chaperoned value like *evt*, but with *proc* as an event generator when the result is synchronized with functions like sync.

The *proc* generator is called on synchronization, much like the procedure passed to guard-evt, except that *proc* is given *evt*. The *proc* must return two values: a synchronizable event that is a chaperone of *evt*, and a procedure that is used to check the event's result if it is chosen in a selection. The latter procedure accepts the result of *evt*, and it must return a chaperone of that value.

Pairs of *prop* and *prop-val* (the number of arguments to chaperone-evt must be even) add impersonator properties or override impersonator-property values of *evt*.

```
(chaperone-channel channel
                   get-proc
                   put-proc
                   prop
                   prop-val ...
                   ...)          → (and/c channel? chaperone?)
  channel : channel?
  get-proc : (channel? . -> . (values channel? (any/c . -> . any/c)))
  put-proc : (channel? any/c . -> . any/c)
  prop : impersonator-property?
  prop-val : any
```

Like impersonate-channel, but with restrictions on the *get-proc* and *put-proc* procedures.

The *get-proc* must return two values: a channel that is a chaperone of *channel*, and a procedure that is used to check the channel's contents. The latter procedure must return the original value or a chaperone of that value.

The *put-proc* must produce a replacement value that is either the original value communicated on the channel or a chaperone of that value.

Pairs of *prop* and *prop-val* (the number of arguments to chaperone-channel must be odd) add impersonator properties or override impersonator-property values of *channel*.

```
(chaperone-prompt-tag  prompt-tag
                       handle-proc
                       abort-proc
                      [cc-guard-proc
                       callcc-chaperone-proc]
                       prop
                       prop-val ...
                       ...)
 → (and/c continuation-prompt-tag? chaperone?)
  prompt-tag : continuation-prompt-tag?
  handle-proc : procedure?
  abort-proc : procedure?
  cc-guard-proc : procedure? = values
  callcc-chaperone-proc : (procedure? . -> . procedure?)
                        = (lambda (p) p)
  prop : impersonator-property?
  prop-val : any
```

Like impersonate-prompt-tag, but produces a chaperoned value. The *handle-proc* procedure must produce the same values or chaperones of the original values, *abort-proc* must produce the same values or chaperones of the values that it is given, and *cc-guard-proc* must produce the same values or chaperones of the original result values, and *callcc-chaperone-proc* must produce a procedure that is a chaperone or the same as the given procedure.

Examples:

```
> (define bad-chaperone
    (chaperone-prompt-tag
     (make-continuation-prompt-tag)
     (lambda (n) (* n 2))
     (lambda (n) (+ n 1))))

> (call-with-continuation-prompt
    (lambda ()
      (abort-current-continuation bad-chaperone 5))
    bad-chaperone
    (lambda (n) n))
abort-current-continuation: non-chaperone result;
 received a value that is not a chaperone of the original
```

```
value
  original: 5
  received: 6
> (define good-chaperone
    (chaperone-prompt-tag
     (make-continuation-prompt-tag)
     (lambda (n) (if (even? n) n (error "not even")))
     (lambda (n) (if (even? n) n (error "not even")))))

> (call-with-continuation-prompt
    (lambda ()
      (abort-current-continuation good-chaperone 2))
    good-chaperone
    (lambda (n) n))
2
```

```
(chaperone-continuation-mark-key key
                                 get-proc
                                 set-proc
                                 prop
                                 prop-val ...
                                 ...)
 → (and/c continuation-mark-key? chaperone?)
  key : continuation-mark-key?
  get-proc : procedure?
  set-proc : procedure?
  prop : impersonator-property?
  prop-val : any
```

Like `impersonate-continuation-mark-key`, but produces a chaperoned value. The *get-proc* procedure must produce the same value or a chaperone of the original value, and *set-proc* must produce the same value or a chaperone of the value that it is given.

Examples:

```
> (define bad-chaperone
    (chaperone-continuation-mark-key
     (make-continuation-mark-key)
     (lambda (l) (map char-upcase l))
     string->list))

> (with-continuation-mark bad-chaperone "timballo"
    (continuation-mark-set-first
     (current-continuation-marks)
     bad-chaperone))
with-continuation-mark: non-chaperone result;
```

```
  received a value that is not a chaperone of the original
value
  original: "timballo"
  received: '(#\t #\i #\m #\b #\a #\l #\l #\o)
> (define (checker s)
    (if (> (string-length s) 5)
        s
        (error "expected string of length at least 5")))

> (define good-chaperone
    (chaperone-continuation-mark-key
     (make-continuation-mark-key)
     checker
     checker))

> (with-continuation-mark good-chaperone "zabaione"
    (continuation-mark-set-first
     (current-continuation-marks)
     good-chaperone))
"zabaione"
```

14.5.3 Impersonator Properties

```
                                      impersonator-property?
(make-impersonator-property name) → (-> any/c boolean?)
                                      (-> impersonator? any)
  name : symbol?
```

Creates a new impersonator property and returns three values:

- an *impersonator property descriptor*, for use with `impersonate-procedure`, `chaperone-procedure`, and other impersonator constructors;
- an *impersonator property predicate* procedure, which takes an arbitrary value and returns `#t` if the value is an impersonator with a value for the property, `#f` otherwise;
- an *impersonator property accessor* procedure, which returns the value associated with an impersonator for the property; if a value given to the accessor is not an impersonator or does not have a value for the property (i.e. if the corresponding impersonator property predicate returns `#f`), the `exn:fail:contract` exception is raised.

```
(impersonator-property? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a impersonator property descriptor value, `#f` otherwise.

```
(impersonator-property-accessor-procedure? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is an accessor procedure produced by `make-impersonator-property`, `#f` otherwise.

```
impersonator-prop:application-mark : impersonator-property?
```

An impersonator property that is recognized by `impersonate-procedure` and `chaperone-procedure`.

14.6 Security Guards

```
(security-guard? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a security guard value as created by `make-security-guard`, `#f` otherwise.

A *security guard* provides a set of access-checking procedures to be called when a thread initiates access of a file, directory, or network connection through a primitive procedure. For example, when a thread calls `open-input-file`, the thread's current security guard is consulted to check whether the thread is allowed read access to the file. If access is granted, the thread receives a port that it may use indefinitely, regardless of changes to the security guard (although the port's custodian could shut down the port; see §14.7 "Custodians").

A thread's current security guard is determined by the `current-security-guard` parameter. Every security guard has a parent, and a parent's access procedures are called whenever a child's access procedures are called. Thus, a thread cannot increase its own access arbitrarily by installing a new guard. The initial security guard enforces no access restrictions other than those enforced by the host platform.

```
(make-security-guard parent
                     file-guard
                     network-guard
                     [link-guard])  → security-guard?
  parent : security-guard?
  file-guard : (symbol?
                (or/c path? #f)
                (listof symbol?)
                . -> . any)
```

```
                (symbol?
                 (or/c (and/c string? immutable?) #f)
network-guard :  (or/c (integer-in 1 65535) #f)
                 (or/c 'server 'client)
                 . -> . any)
link-guard : (or/c (symbol? path? path? . -> . any) #f) = #f
```

Creates a new security guard as child of *parent*.

The *file-guard* procedure must accept three arguments:

- a symbol for the primitive procedure that triggered the access check, which is useful for raising an exception to deny access.
- a path (see §15.1 “Paths”) or `#f` for pathless queries, such as `(current-directory)`, `(filesystem-root-list)`, and `(find-system-path` *symbol*`)`. A path provided to *file-guard* is not expanded or otherwise normalized before checking access; it may be a relative path, for example.
- a list containing one or more of the following symbols:
  - `'read` — read a file or directory
  - `'write` — modify or create a file or directory
  - `'execute` — execute a file
  - `'delete` — delete a file or directory
  - `'exists` — determine whether a file or directory exists, or that a path string is well-formed

  The `'exists` symbol is never combined with other symbols in the last argument to *file-guard*, but any other combination is possible. When the second argument to *file-guard* is `#f`, the last argument always contains only `'exists`.

The *network-guard* procedure must accept four arguments:

- a symbol for the primitive operation that triggered the access check, which is useful for raising an exception to deny access.
- an immutable string representing the target hostname for a client connection or the accepting hostname for a listening server; `#f` for a listening server or UDP socket that accepts connections at all of the host’s address; or `#f` an unbound UDP socket.
- an exact integer between `1` and `65535` (inclusive) representing the port number, or `#f` for an unbound UDP socket. In the case of a client connection, the port number is the target port on the server. For a listening server, the port number is the local port number.

- a symbol, either `'client` or `'server`, indicating whether the check is for the creation of a client connection or a listening server. The opening of an unbound UDP socket is identified as a `'client` connection; explicitly binding the socket is identified as a `'server` action.

The *link-guard* argument can be `#f` or a procedure of three arguments:

- a symbol for the primitive procedure that triggered the access check, which is useful for raising an exception to deny access.
- a complete path (see §15.1 "Paths") representing the file to create as link.
- a path representing the content of the link, which may be relative the second-argument path; this path is not expanded or otherwise normalized before checking access.

If *link-guard* is `#f`, then a default procedure is used that always raises `exn:fail`.

The return value of *file-guard*, *network-guard*, or *link-guard* is ignored. To deny access, the procedure must raise an exception or otherwise escape from the context of the primitive call. If the procedure returns, the parent's corresponding procedure is called on the same inputs, and so on up the chain of security guards.

The *file-guard*, *network-guard*, and *link-guard* procedures are invoked in the thread that called the access-checked primitive. Breaks may or may not be enabled (see §10.6 "Breaks"). Full continuation jumps are blocked going into or out of the *file-guard* or *network-guard* call (see §1.1.12 "Prompts, Delimited Continuations, and Barriers").

```
(current-security-guard) → security-guard?
(current-security-guard guard) → void?
  guard : security-guard?
```

A parameter that determines the current security guard that controls access to the filesystem and network.

14.7 Custodians

See §1.1.16 "Custodians" for basic information on the Racket custodian model.

```
(custodian? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a custodian value, `#f` otherwise.

```
(make-custodian [cust]) → custodian?
  cust : custodian? = (current-custodian)
```

Creates a new custodian that is subordinate to *cust*. When *cust* is directed (via `custodian-shutdown-all`) to shut down all of its managed values, the new subordinate custodian is automatically directed to shut down its managed values as well.

```
(custodian-shutdown-all cust) → void?
  cust : custodian?
```

In `racket/gui/base`, eventspaces managed by *cust* are also shut down.

Closes all file-stream ports, TCP ports, TCP listeners, and UDP sockets that are managed by *cust* (and its subordinates), and empties all custodian boxes associated with *cust* (and its subordinates). It also removes *cust* (and its subordinates) as managers of all threads; when a thread has no managers, it is killed (or suspended; see `thread/suspend-to-kill`) If the current thread is to be killed, all other shut-down actions take place before killing the thread.

```
(current-custodian) → custodian?
(current-custodian cust) → void?
  cust : custodian?
```

Custodians also manage eventspaces from `racket/gui/base`.

A parameter that determines a custodian that assumes responsibility for newly created threads, file-stream ports, TCP ports, TCP listeners, UDP sockets, and byte converters.

```
(custodian-managed-list cust super) → list?
  cust : custodian?
  super : custodian?
```

Returns a list of immediately managed objects (not including custodian boxes) and subordinate custodians for *cust*, where *cust* is itself subordinate to *super* (directly or indirectly). If *cust* is not strictly subordinate to *super*, the `exn:fail:contract` exception is raised.

```
(custodian-memory-accounting-available?) → boolean?
```

Memory accounting is normally available in Racket 3m, which is the main variant of Racket, and not normally available in Racket CGC.

Returns `#t` if Racket is compiled with support for per-custodian memory accounting, `#f` otherwise.

```
(custodian-require-memory limit-cust
                          need-amt
                          stop-cust) → void?
  limit-cust : custodian?
  need-amt : exact-nonnegative-integer?
  stop-cust : custodian?
```

Registers a required-memory check if Racket is compiled with support for per-custodian memory accounting, otherwise the `exn:fail:unsupported` exception is raised.

If a check is registered, and if Racket later reaches a state after garbage collection (see §1.1.7 “Garbage Collection”) where allocating *need-amt* bytes charged to *limit-cust* would fail or trigger some shutdown, then *stop-cust* is shut down.

```
(custodian-limit-memory limit-cust
                        limit-amt
                       [stop-cust]) → void?
  limit-cust : custodian?
  limit-amt : exact-nonnegative-integer?
  stop-cust : custodian? = limit-cust
```

Registers a limited-memory check if Racket is compiled with support for per-custodian memory accounting, otherwise the `exn:fail:unsupported` exception is raised.

If a check is registered, and if Racket later reaches a state after garbage collection (see §1.1.7 “Garbage Collection”) where *limit-cust* owns more than *limit-amt* bytes, then *stop-cust* is shut down.

A custodian’s limit is checked only after a garbage collection, except that it may also be checked during certain large allocations that are individually larger than the custodian’s limit. A single garbage collection may shut down multiple custodians, even if shutting down only one of the custodians would have reduced memory use for other custodians.

For reliable shutdown, *limit-amt* for `custodian-limit-memory` must be much lower than the total amount of memory available (minus the size of memory that is potentially used and not charged to *limit-cust*). Moreover, if individual allocations that are initially charged to *limit-cust* can be arbitrarily large, then *stop-cust* must be the same as *limit-cust*, so that excessively large immediate allocations can be rejected with an `exn:fail:out-of-memory` exception.

```
(make-custodian-box cust v) → custodian-box?
  cust : custodian?
  v : any/c
```

Returns a custodian box that contains *v* as long as *cust* has not been shut down.

A custodian box is a synchronizable event (see §11.2.1 “Events”). The custodian box becomes ready when its custodian is shut down; the synchronization result of a custodian box is the custodian box itself.

```
(custodian-box? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a custodian box produced by `make-custodian-box`, `#f` otherwise.

```
(custodian-box-value cb) → any
  cb : custodian-box?
```

Returns the value in the given custodian box, or `#f` if the value has been removed.

14.8 Thread Groups

A *thread group* is a collection of threads and other thread groups that have equal claim to the CPU. By nesting thread groups and by creating certain threads within certain groups,

a programmer can control the amount of CPU allocated to a set of threads. Every thread belongs to a thread group, which is determined by the current-thread-group parameter when the thread is created. Thread groups and custodians (see §14.7 "Custodians") are independent.

The root thread group receives all of the CPU that the operating system gives Racket. Every thread or nested group in a particular thread group receives equal allocation of the CPU (a portion of the group's access), although a thread may relinquish part of its allocation by sleeping or synchronizing with other processes.

```
(make-thread-group [group]) → thread-group?
  group : thread-group? = (current-thread-group)
```

Creates a new thread group that belongs to *group*.

```
(thread-group? v) → boolean?
  v : any/c
```

Returns #t if *v* is a thread group value, #f otherwise.

```
(current-thread-group) → thread-group?
(current-thread-group group) → void?
  group : thread-group?
```

A parameter that determines the thread group for newly created threads.

14.9 Structure Inspectors

An *inspector* provides access to structure fields and structure type information without the normal field accessors and mutators. (Inspectors are also used to control access to module bindings; see §14.10 "Code Inspectors".) Inspectors are primarily intended for use by debuggers.

When a structure type is created, an inspector can be supplied. The given inspector is not the one that will control the new structure type; instead, the given inspector's parent will control the type. By using the parent of the given inspector, the structure type remains opaque to "peer" code that cannot access the parent inspector.

The current-inspector parameter determines a default inspector argument for new structure types. An alternate inspector can be provided though the #:inspector option of the define-struct form (see §5.1 "Defining Structure Types: struct"), or through an optional inspector argument to make-struct-type.

```
(inspector? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is an inspector, `#f` otherwise.

```
(make-inspector [inspector]) → inspector?
  inspector : inspector? = (current-inspector)
```

Returns a new inspector that is a subinspector of *inspector*. Any structure type controlled by the new inspector is also controlled by its ancestor inspectors, but no other inspectors.

```
(make-sibling-inspector [inspector]) → inspector?
  inspector : inspector? = (current-inspector)
```

Returns a new inspector that is a subinspector of the same inspector as *inspector*. That is, *inspector* and the result inspector control mutually disjoint sets of structure types.

```
(current-inspector) → inspector?
(current-inspector insp) → void?
  insp : inspector?
```

A parameter that determines the default inspector for newly created structure types.

```
(struct-info v) → (or/c struct-type? #f) boolean?
  v : any/c
```

Returns two values:

- *struct-type*: a structure type descriptor or `#f`; the result is a structure type descriptor of the most specific type for which *v* is an instance, and for which the current inspector has control, or the result is `#f` if the current inspector does not control any structure type for which the `struct` is an instance.
- *skipped?*: `#f` if the first result corresponds to the most specific structure type of *v*, `#t` otherwise.

```
(struct-type-info struct-type)
    symbol?
    exact-nonnegative-integer?
    exact-nonnegative-integer?
    struct-accessor-procedure?
→
    struct-mutator-procedure?
    (listof exact-nonnegative-integer?)
    (or/c struct-type? #f)
    boolean?
  struct-type : struct-type?
```

Returns eight values that provide information about the structure type descriptor *struct-type*, assuming that the type is controlled by the current inspector:

- *name*: the structure type's name as a symbol;
- *init-field-cnt*: the number of fields defined by the structure type provided to the constructor procedure (not counting fields created by its ancestor types);
- *auto-field-cnt*: the number of fields defined by the structure type without a counterpart in the constructor procedure (not counting fields created by its ancestor types);
- *accessor-proc*: an accessor procedure for the structure type, like the one returned by `make-struct-type`;
- *mutator-proc*: a mutator procedure for the structure type, like the one returned by `make-struct-type`;
- *immutable-k-list*: an immutable list of exact non-negative integers that correspond to immutable fields for the structure type;
- *super-type*: a structure type descriptor for the most specific ancestor of the type that is controlled by the current inspector, or `#f` if no ancestor is controlled by the current inspector;
- *skipped?*: `#f` if the seventh result is the most specific ancestor type or if the type has no supertype, `#t` otherwise.

If the type for *struct-type* is not controlled by the current inspector, the `exn:fail:contract` exception is raised.

```
(struct-type-make-constructor struct-type)
 → struct-constructor-procedure?
  struct-type : struct-type?
```

Returns a constructor procedure to create instances of the type for *struct-type*. If the type for *struct-type* is not controlled by the current inspector, the `exn:fail:contract` exception is raised.

```
(struct-type-make-predicate struct-type) → any
  struct-type : any/c
```

Returns a predicate procedure to recognize instances of the type for *struct-type*. If the type for *struct-type* is not controlled by the current inspector, the `exn:fail:contract` exception is raised.

```
(object-name v) → any
  v : any/c
```

Returns a value for the name of `v` if `v` has a name, `#f` otherwise. The argument `v` can be any value, but only (some) procedures, structures, structure types, structure type properties, regexp values, ports, and loggers have names. See also §1.2.6 “Inferred Value Names”.

The name (if any) of a procedure is always a symbol. The `procedure-rename` function creates a procedure with a specific name.

The name of a structure, structure type, structure type property is always a symbol. If a structure is a procedure as implemented by one of its fields (i.e., the `prop:procedure` property value for the structure’s type is an integer), then its name is the implementing procedure’s name; otherwise, its name matches the name of the structure type that it instantiates.

The name of a regexp value is a string or byte string. Passing the string or byte string to `regexp`, `byte-regexp`, `pregexp`, or `byte-pregexp` (depending on the kind of regexp whose name was extracted) produces a value that matches the same inputs.

The name of a port can be any value, but many tools use a path or string name as the port’s for (to report source locations, for example).

The name of a logger is either a symbol or `#f`.

14.10 Code Inspectors

In the same way that inspectors control access to structure fields (see §14.9 “Structure Inspectors”), inspectors also control access to module bindings. The default inspector for module bindings is determined by the `current-code-inspector` parameter, instead of the `current-inspector` parameter.

When a `module` declaration is evaluated, the value of the `current-code-inspector` parameter is associated with the module declaration. When the module is invoked via `require` or `dynamic-require`, a sub-inspector of the module’s declaration-time inspector is created, and this sub-inspector is associated with the module invocation. Any inspector that controls the sub-inspector (including the declaration-time inspector and its superior) controls the module invocation. In particular, if the value of `current-code-inspector` never changes, then no control is lost for any module invocation, since the module’s invocation is associated with a sub-inspector of `current-code-inspector`.

When an inspector that controls a module invocation is installed `current-code-inspector`, it enables the following `module->namespace` on the module, and it enables access to the module’s protected exports (i.e., those identifiers exported from the module with `protect-out`) via `dynamic-require`.

When a `module` form is expanded or a namespace is created, the value of `current-code-inspector` is associated with the module or namespace’s top-level lexical information. Syntax objects with that lexical information gain access to the protected and unexported bindings

of any module that the inspector controls. In the case of a `module`, the inspector sticks with such syntax objects even the syntax object is used in the expansion of code in a less powerful context; furthermore, if the syntax object is an identifier that is compiled as a variable reference, the inspector sticks with the variable reference even if it appears in a module form that is evaluated (i.e., declared) with a weaker inspector. When a syntax object or variable reference is within compiled code that is printed (see §1.4.16 "Printing Compiled Code"), the associated inspector is not preserved.

When compiled code in printed form is `read` back in, no inspectors are associated with the code. When the code is `eval`uated, the instantiated syntax-object literals and module-variable references acquire value of `current-code-inspector` as their inspector.

When a module instantiation is attached to multiple namespaces, each with its own module registry, the inspector for the module invocation can be registry-specific. The invocation inspector in a particular module registry can be changed via `namespace-unprotect-module` (but changing the inspector requires control over the old one).

```
(current-code-inspector) → inspector?
(current-code-inspector insp) → void?
  insp : inspector?
```

A parameter that determines an inspector to control access to module bindings and redefinitions.

14.11 Plumbers

A *plumber* supports *flush callbacks*, which are normally triggered just before a Racket process or place exits. For example, a flush callback might flush an output port's buffer.

Flush callbacks are roughly analogous to the standard C library's `atexit`, but flush callback can also be used in other, similar scenarios.

There is no guarantee that a flush callback will be called before a process terminates—either because the plumber is not the original plumber that is flushed by the default exit handler, or because the process is terminated forcibly (e.g., through a custodian shutdown).

```
(plumber? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a plumber value, `#f` otherwise.

Added in version 6.0.1.8 of package `base`.

```
(make-plumber) → pluber?
```

Creates a new plumber.

Plumbers have no hierarchy (unlike custodians or inspectors), but a flush callback can be registered in one plumber to call `plumber-flush-all` with another plumber.

Added in version 6.0.1.8 of package `base`.

```
(current-plumber) → plumber?
(current-plumber plumber) → void?
  plumber : plumber?
```

A parameter that determines a *current plumber* for flush callbacks. For example, creating an output file stream port registers a flush callback with the current plumber to flush the port as long as the port is opened.

Added in version 6.0.1.8 of package `base`.

```
(plumber-flush-all plumber) → void?
  plumber : plumber?
```

Calls all flush callbacks that are registered with *plumber*.

The flush callbacks to call are collected from *plumber* before the first one is called. If a flush callback registers a new flush callback, the new one is *not* called. If a flush callback raises an exception or otherwise escapes, then the remaining flush callbacks are not called.

Added in version 6.0.1.8 of package `base`.

```
(plumber-flush-handle? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a *flush handle* represents the registration of a flush callback, `#f` otherwise.

Added in version 6.0.1.8 of package `base`.

```
(plumber-add-flush! plumber proc [weak?]) → plumber-flush-handle?
  plumber : plumber?
  proc : (plumber-flush-handle? . -> . any)
  weak? : any/c = #f
```

Registers *proc* as a flush callback with *plumber*, so that *proc* is called when `plumber-flush-all` is applied to *plumber*.

The result flush handle represents the registration of the callback, and it can be used with `plumber-flush-handle-remove!` to unregister the callback.

The given *proc* is reachable from the flush handle, but if *weak?* is true, then *plumber* retains only a weak reference to the result flush handle (and thus *proc*).

When *proc* is called as a flush callback, it is passed the same value that is returned by `plumber-add-flush!` so that *proc* can conveniently unregister itself. The call of *proc* is within a continuation barrier.

Added in version 6.0.1.8 of package `base`.

```
(plumber-flush-handle-remove! handle) → void?
  handle : plumber-flush-handle?
```

Unregisters the flush callback that was registered by the `plumber-add-flush!` call that produced *handle*.

If the registration represented by *handle* has been removed already, then `plumber-flush-handle-remove!` has no effect.

Added in version 6.0.1.8 of package `base`.

14.12 Sandboxed Evaluation

`(require racket/sandbox)` package: `sandbox-lib`

The bindings documented in this section are provided by the `racket/sandbox` library, not `racket/base` or `racket`.

The `racket/sandbox` module provides utilities for creating "sandboxed" evaluators, which are configured in a particular way and can have restricted resources (memory and time), filesystem and network access, and much more. Sandboxed evaluators can be configured through numerous parameters — and the defaults are set for the common use case where sandboxes are very limited.

```
(make-evaluator  language
                 input-program ...
                [#:requires requires
                 #:allow-for-require allow-for-require
                 #:allow-for-load allow-for-load
                 #:allow-read allow-read])
 → (any/c . -> . any)
             (or/c module-path?
  language :       (list/c 'special symbol?)
                   (cons/c 'begin list?))
  input-program : any/c
              (listof (or/c module-path? path-string?
  requires :                  (cons/c 'for-syntax (listof module-path?))))
           = null
  allow-for-require : (listof (or/c module-path? path?)) = null
  allow-for-load : (listof path-string?) = null
```

```
  allow-read : (listof (or/c module-path? path-string?)) = null
(make-module-evaluator module-decl
                       [#:language lang
                        #:allow-for-require allow-for-require
                        #:allow-for-load allow-for-load
                        #:allow-read allow-read])
 → (any/c . -> . any)
  module-decl : (or/c syntax? pair? path? input-port? string? bytes?)
  lang : (or/c #f module-path?) = #f
  allow-for-require : (listof (or/c module-path? path?)) = null
  allow-for-load : (listof path-string?) = null
  allow-read : (listof (or/c module-path? path-string?)) = null
```

The `make-evaluator` function creates an evaluator with a *`language`* and *`requires`* specification, and starts evaluating the given *`input-program`*s. The `make-module-evaluator` function creates an evaluator that works in the context of a given module. The result in either case is a function for further evaluation.

The returned evaluator operates in an isolated and limited environment. In particular, filesystem access is restricted, which may interfere with using modules from the filesystem. See below for information on the *`allow-for-require`*, *`allow-for-load`*, and *`allow-read`* arguments. When *`language`* is a module path or when *`requires`* is provided, the indicated modules are implicitly included in the *`allow-for-require`* list. (For backward compatibility, non-`module-path?` path strings are allowed in *`requires`*; they are implicitly converted to paths before addition to *`allow-for-require`*.)

Each *`input-program`* or *`module-decl`* argument provides a program in one of the following forms:

- an input port used to read the program;
- a string or a byte string holding the complete input;
- a path that names a file holding the input; or
- an S-expression or a syntax object, which is evaluated as with `eval` (see also `get-uncovered-expressions`).

In the first three cases above, the program is read using `sandbox-reader`, with line-counting enabled for sensible error messages, and with `'program` as the source (used for testing coverage). In the last case, the input is expected to be the complete program, and is converted to a syntax object (using `'program` as the source), unless it already is a syntax object.

The returned evaluator function accepts additional expressions (each time it is called) in essentially the same form: a string or byte string holding a sequence of expressions, a path

for a file holding expressions, an S-expression, or a syntax object. If the evaluator receives an `eof` value, it is terminated and raises errors thereafter. See also `kill-evaluator`, which terminates the evaluator without raising an exception.

For `make-evaluator`, multiple *`input-program`*s are effectively concatenated to form a single program. The way that the *`input-program`*s are evaluated depends on the *`language`* argument:

- The *`language`* argument can be a module path (i.e., a datum that matches the grammar for *`module-path`* of `require`).

  In this case, the *`input-program`*s are automatically wrapped in a `module`, and the resulting evaluator works within the resulting module's namespace.

- The *`language`* argument can be a list starting with `'special`, which indicates a built-in language with special input configuration. The possible values are `'(special r5rs)` or a value indicating a teaching language: `'(special beginner)`, `'(special beginner-abbr)`, `'(special intermediate)`, `'(special intermediate-lambda)`, or `'(special advanced)`.

  In this case, the *`input-program`*s are automatically wrapped in a `module`, and the resulting evaluator works within the resulting module's namespace. In addition, certain parameters (such as such as `read-accept-infix-dot`) are set to customize reading programs from strings and ports.

  This option is provided mainly for older test systems. Using `make-module-evaluator` with input starting with `#lang` is generally better.

- Finally, *`language`* can be a list whose first element is `'begin`.

  In this case, a new namespace is created using `sandbox-namespace-specs`, which by default creates a new namespace using `sandbox-make-namespace` (which, in turn, uses `make-base-namespace` or `make-gui-namespace` depending on `sandbox-gui-available` and `gui-available?`).

  In the new namespace, *`language`* is evaluated as an expression to further initialize the namespace.

The *`requires`* list adds additional imports to the module or namespace for the *`input-program`*s, even in the case that `require` is not made available through the *`language`*.

The following examples illustrate the difference between an evaluator that puts the program in a module and one that merely initializes a top-level namespace:

```
> (define base-module-eval
    ; a module cannot have free variables...
    (make-evaluator 'racket/base '(define (f) later)))
program:1:0: later: unbound identifier in module
  in: later
```

```
> (define base-module-eval
    (make-evaluator 'racket/base '(define (f) later)
                                 '(define later 5)))

> (base-module-eval '(f))
5
> (define base-top-eval
    ; non-module code can have free variables:
    (make-evaluator '(begin) '(define (f) later)))

> (base-top-eval '(+ 1 2))
3
> (base-top-eval '(define later 5))

> (base-top-eval '(f))
5
```

The make-module-evaluator function is essentially a restriction of make-evaluator, where the program must be a module, and all imports are part of the program. In some cases it is useful to restrict the program to be a module using a specific module in its language position — use the optional *lang* argument to specify such a restriction (the default, #f, means no restriction is enforced). When the program is specified as a path, then the path is implicitly added to the *allow-for-load* list.

```
(define base-module-eval2
  ; equivalent to base-module-eval:
  (make-module-evaluator '(module m racket/base
                            (define (f) later)
                            (define later 5))))
```

The make-module-evaluator function can be convenient for testing module files: pass in a path value for the file name, and you get back an evaluator in the module's context which you can use with your favorite test facility.

In all cases, the evaluator operates in an isolated and limited environment:

- It uses a new custodian and namespace. When gui-available? and sandbox-gui-available produce true, it is also runs in its own eventspace.
- The evaluator works under the sandbox-security-guard, which restricts file system and network access.
- The evaluator is contained in a memory-restricted environment, and each evaluation is wrapped in a call-with-limits (when memory accounting is available); see also sandbox-memory-limit, sandbox-eval-limits and set-eval-limits.

Note that these limits apply to the creation of the sandbox environment too — so, for example, if the memory that is required to create the sandbox is higher than the limit, then `make-evaluator` will fail with a memory limit exception.

The *`allow-for-require`* and *`allow-for-load`* arguments adjust filesystem permissions to extend the set of files that are usable by the evaluator. The *`allow-for-require`* argument lists modules that can be `required` along with their imports (transitively). The *`allow-for-load`* argument lists files that can be `load`ed. (The precise permissions needed for `require` versus `load` can differ.) The *`allow-read`* argument is for backward compatibility, only; each `module-path?` element of *`allow-read`* is effectively moved to *`allow-for-require`*, while other elements are moved to `all-for-load`.

The sandboxed environment is well isolated, and the evaluator function essentially sends it an expression and waits for a result. This form of communication makes it impossible to have nested (or concurrent) calls to a single evaluator. Usually this is not a problem, but in some cases you can get the evaluator function available inside the sandboxed code, for example:

```
> (let ([e (make-evaluator 'racket/base)])
    (e `(,e 1)))
evaluator: nested evaluator call with: 1
```

An error will be signaled in such cases.

If the value of `sandbox-propagate-exceptions` is true (the default) when the sandbox is created, then exceptions (both syntax and run-time) are propagated as usual to the caller of the evaluation function (i.e., catch them with `with-handlers`). If the value of `sandbox-propagate-exceptions` is `#f` when the sandbox is created, then uncaught exceptions in a sandbox evaluation cause the error to be printed to the sandbox's error port, and the caller of the evaluation receives `#<void>`.

Finally, the fact that a sandboxed evaluator accept syntax objects makes it usable as the value for `current-eval`, which means that you can easily start a sandboxed read-eval-print-loop. For example, here is a quick implementation of a networked REPL:

```
(define e (make-evaluator 'racket/base))
(let-values ([(i o) (tcp-accept (tcp-listen 9999))])
  (parameterize ([current-input-port  i]
                 [current-output-port o]
                 [current-error-port  o]
                 [current-eval e])
    (read-eval-print-loop)
    (fprintf o "\nBye...\n")
    (close-output-port o)))
```

Note that in this code it is only the REPL interactions that are going over the network con-

nection; using I/O operations inside the REPL will still use the usual sandbox parameters (defaulting to no I/O). In addition, the code works only from an existing toplevel REPL — specifically, `read-eval-print-loop` reads a syntax value and gives it the lexical context of the current namespace. Here is a variation that uses the networked ports for user I/O, and works when used from a module (by using a new namespace):

```
(let-values ([(i o) (tcp-accept (tcp-listen 9999))])
  (parameterize ([current-input-port   i]
                 [current-output-port  o]
                 [current-error-port   o]
                 [sandbox-input        i]
                 [sandbox-output       o]
                 [sandbox-error-output o]
                 [current-namespace (make-empty-namespace)])
    (parameterize ([current-eval
                    (make-evaluator 'racket/base)])
      (read-eval-print-loop))
    (fprintf o "\nBye...\n")
    (close-output-port o)))
```

```
(exn:fail:sandbox-terminated? v) → boolean?
  v : any/c
(exn:fail:sandbox-terminated-reason exn) → symbol?
  exn : exn:fail:sandbox-terminated?
```

A predicate and accessor for exceptions that are raised when a sandbox is terminated. Once a sandbox raises such an exception, it will continue to raise it on further evaluation attempts.

14.12.1 Customizing Evaluators

The sandboxed evaluators that `make-evaluator` creates can be customized via many parameters. Most of the configuration parameters affect newly created evaluators; changing them has no effect on already-running evaluators.

The default configuration options are set for a very restricted sandboxed environment — one that is safe to make publicly available. Further customizations might be needed in case more privileges are needed, or if you want tighter restrictions. Another useful approach for customizing an evaluator is to begin with a relatively unrestricted configuration and add the desired restrictions. This approach is made possible by the `call-with-trusted-sandbox-configuration` function.

The sandbox environment uses two notions of restricting the time that evaluations takes: shallow time and deep time. *Shallow time* refers to the immediate execution of an expression. For example, a shallow time limit of five seconds would restrict `(sleep 6)` and other

computations that take longer than five seconds. *Deep time* refers to the total execution of the expression and all threads and sub-processes that the expression creates. For example, a deep time limit of five seconds would restrict `(thread (λ () (sleep 6)))`, which shallow time would not, *as well as* all expressions that shallow time would restrict. By default, most sandboxes only restrict shallow time to facilitate expressions that use threads.

```
(call-with-trusted-sandbox-configuration thunk) → any
  thunk : (-> any)
```

Invokes the *thunk* in a context where sandbox configuration parameters are set for minimal restrictions. More specifically, there are no memory or time limits, and the existing existing inspectors, security guard, exit handler, logger, plumber, and environment variable set are used. (Note that the I/O ports settings are not included.)

```
(sandbox-init-hook) → (-> any)
(sandbox-init-hook thunk) → void?
  thunk : (-> any)
```

A parameter that determines a thunk to be called for initializing a new evaluator. The hook is called just before the program is evaluated in a newly-created evaluator context. It can be used to setup environment parameters related to reading, writing, evaluation, and so on. Certain languages (`'(special r5rs)` and the teaching languages) have initializations specific to the language; the hook is used after that initialization, so it can override settings.

```
(sandbox-reader) → (any/c . -> . any)
(sandbox-reader proc) → void?
  proc : (any/c . -> . any)
```

A parameter that specifies a function that reads all expressions from `(current-input-port)`. The function is used to read program source for an evaluator when a string, byte string, or port is supplied. The reader function receives a value to be used as input source (i.e., the first argument to `read-syntax`), and it should return a list of syntax objects. The default reader calls `read-syntax`, accumulating results in a list until it receives `eof`.

Note that the reader function is usually called as is, but when it is used to read the program input for `make-module-evaluator`, `read-accept-lang` and `read-accept-reader` are set to `#t`.

```
                    (or/c #f
                          string? bytes?
(sandbox-input) →         input-port?
                          'pipe
                          (-> input-port?))
(sandbox-input in) → void?
```

```
      (or/c #f
            string? bytes?
 in :       input-port?
            'pipe
            (-> input-port?))
```

A parameter that determines the initial `current-input-port` setting for a newly created evaluator. It defaults to `#f`, which creates an empty port. The following other values are allowed:

- a string or byte string, which is converted to a port using `open-input-string` or `open-input-bytes`;
- an input port;
- the symbol `'pipe`, which triggers the creation of a pipe, where `put-input` can return the output end of the pipe or write directly to it;
- a thunk, which is called to obtain a port (e.g., using `current-input-port` means that the evaluator input is the same as the calling context's input).

```
                    (or/c #f
                          output-port?
(sandbox-output) →        'pipe
                          'bytes
                          'string
                          (-> output-port?))
(sandbox-output in) → void?
      (or/c #f
            output-port?
 in :       'pipe
            'bytes
            'string
            (-> output-port?))
```

A parameter that determines the initial `current-output-port` setting for a newly created evaluator. It defaults to `#f`, which creates a port that discards all data. The following other values are allowed:

- an output port, which is used as-is;
- the symbol `'bytes`, which causes `get-output` to return the complete output as a byte string as long as the evaluator has not yet terminated (so that the size of the bytes can be charged to the evaluator);

- the symbol `'string`, which is similar to `'bytes`, but makes `get-output` produce a string;
- the symbol `'pipe`, which triggers the creation of a pipe, where `get-output` returns the input end of the pipe;
- a thunk, which is called to obtain a port (e.g., using `current-output-port` means that the evaluator output is not diverted).

```
                          (or/c #f
                                output-port?
(sandbox-error-output) →        'pipe
                                'bytes
                                'string
                                (-> output-port?))
(sandbox-error-output in) → void?
      (or/c #f
            output-port?
  in :      'pipe
            'bytes
            'string
            (-> output-port?))
```

Like `sandbox-output`, but for the initial `current-error-port` value. An evaluator's error output is set after its output, so using `current-output-port` (the parameter itself, not its value) for this parameter value means that the error port is the same as the evaluator's initial output port.

The default is `(lambda () (dup-output-port (current-error-port)))`, which means that the error output of the generated evaluator goes to the calling context's error port.

```
(sandbox-coverage-enabled) → boolean?
(sandbox-coverage-enabled enabled?) → void?
  enabled? : any/c
```

A parameter that controls whether syntactic coverage information is collected by sandbox evaluators. Use `get-uncovered-expressions` to retrieve coverage information.

```
(sandbox-propagate-breaks) → boolean?
(sandbox-propagate-breaks propagate?) → void?
  propagate? : any/c
```

When both this boolean parameter and `(break-enabled)` are true, breaking while an evaluator is running propagates the break signal to the sandboxed context. This makes the sandboxed evaluator break, typically, but beware that sandboxed evaluation can capture and avoid

the breaks (so if safe execution of code is your goal, make sure you use it with a time limit). Also, beware that a break may be received after the evaluator's result, in which case the evaluation result is lost. Finally, beware that a break may be propagated after an evaluator has produced a result, so that the break is visible on the next interaction with the evaluator (or the break is lost if the evaluator is not used further). The default is `#t`.

```
(sandbox-propagate-exceptions) → boolean?
(sandbox-propagate-exceptions propagate?) → void?
  propagate? : any/c
```

A parameter that controls how uncaught exceptions during a sandbox evaluation are treated. When the parameter value is `#t`, then the exception is propagated to the caller of sandbox. When the parameter value is `#f`, the exception message is printed to the sandbox's error port, and the caller of the sandbox receives `#<void>` for the evaluation. The default is `#t`.

```
(sandbox-namespace-specs) → (cons/c (-> namespace?)
                                    (listof module-path?))
(sandbox-namespace-specs spec) → void?
  spec : (cons/c (-> namespace?)
                 (listof module-path?))
```

A parameter that holds a list of values that specify how to create a namespace for evaluation in `make-evaluator` or `make-module-evaluator`. The first item in the list is a thunk that creates the namespace, and the rest are module paths for modules to be attached to the created namespace using `namespace-attach-module`.

The default is `(list sandbox-make-namespace)`.

The module paths are needed for sharing module instantiations between the sandbox and the caller. For example, sandbox code that returns `posn` values (from the `lang/posn` module) will not be recognized as such by your own code by default, since the sandbox will have its own instance of `lang/posn` and thus its own struct type for `posn`s. To be able to use such values, include `'lang/posn` in the list of module paths.

When testing code that uses a teaching language, the following piece of code can be helpful:

```
(sandbox-namespace-specs
 (let ([specs (sandbox-namespace-specs)])
   `(,(car specs)
     ,@(cdr specs)
     lang/posn
     ,@(if (gui-available?) '(mrlib/cache-image-snip) '()))))
```

```
(sandbox-make-namespace) → namespace?
```

Calls `make-gui-namespace` when `(sandbox-gui-available)` produces true, `make-base-namespace` otherwise.

```
(sandbox-gui-available) → boolean?
(sandbox-gui-available avail?) → void?
  avail? : any/c
```

Determines whether the `racket/gui` module can be used when a sandbox evaluator is created. If `gui-available?` produces `#f` during the creation of a sandbox evaluator, this parameter is forced to `#f` during initialization of the sandbox. The default value of the parameter is `#t`.

Various aspects of the library change when the GUI library is available, such as using a new eventspace for each evaluator.

```
(sandbox-override-collection-paths) → (listof path-string?)
(sandbox-override-collection-paths paths) → void?
  paths : (listof path-string?)
```

A parameter that determines a list of collection directories to prefix `current-library-collection-paths` in an evaluator. This parameter is useful for cases when you want to test code using an alternate, test-friendly version of a collection, for example, testing code that uses a GUI (like the `htdp/world` teachpack) can be done using a fake library that provides the same interface but no actual interaction. The default is `null`.

```
(sandbox-security-guard)
 → (or/c security-guard? (-> security-guard?))
(sandbox-security-guard guard) → void?
  guard : (or/c security-guard? (-> security-guard?))
```

A parameter that determines the initial `(current-security-guard)` for sandboxed evaluations. It can be either a security guard, or a function to construct one. The default is a function that restricts the access of the current security guard by forbidding all filesystem I/O except for specifications in `sandbox-path-permissions`, and it uses `sandbox-network-guard` for network connections.

```
(sandbox-path-permissions)
    (listof (list/c (or/c 'execute 'write 'delete
 →                        'read-bytecode 'read 'exists)
                    (or/c byte-regexp? bytes? string? path?)))
(sandbox-path-permissions perms) → void?
          (listof (list/c (or/c 'execute 'write 'delete
  perms :                       'read-bytecode 'read 'exists)
                          (or/c byte-regexp? bytes? string? path?)))
```

A parameter that configures the behavior of the default sandbox security guard by listing paths and access modes that are allowed for them. The contents of this parameter is a list

of specifications, each is an access mode and a byte-regexp for paths that are granted this access.

The access mode symbol is one of: `'execute`, `'write`, `'delete`, `'read`, or `'exists`. These symbols are in decreasing order: each implies access for the following modes too (e.g., `'read` allows reading or checking for existence).

The path regexp is used to identify paths that are granted access. It can also be given as a path (or a string or a byte string), which is (made into a complete path, cleansed, simplified, and then) converted to a regexp that allows the path and sub-directories; e.g., `"/foo/bar"` applies to `"/foo/bar/baz"`.

An additional mode symbol, `'read-bytecode`, is not part of the linear order of these modes. Specifying this mode is similar to specifying `'read`, but it is not implied by any other mode. (For example, even if you specify `'write` for a certain path, you need to also specify `'read-bytecode` to grant this permission.) The sandbox usually works in the context of a lower code inspector (see `sandbox-make-code-inspector`) which prevents loading of untrusted bytecode files — the sandbox is set-up to allow loading bytecode from files that are specified with `'read-bytecode`. This specification is given by default to the Racket collection hierarchy (including user-specific libraries) and to libraries that are explicitly specified in an `#:allow-read` argument. (Note that this applies for loading bytecode files only, under a lower code inspector it is still impossible to use protected module bindings (see §14.10 “Code Inspectors”).)

The default value is null, but when an evaluator is created, it is augmented by `'read-bytecode` permissions that make it possible to use collection libraries (including `sandbox-override-collection-paths`). See `make-evaluator` for more information.

```
(sandbox-network-guard)
    (symbol?
     (or/c (and/c string? immutable?) #f)
 →   (or/c (integer-in 1 65535) #f)
     (or/c 'server 'client)
     . -> . any)
(sandbox-network-guard proc) → void?
         (symbol?
          (or/c (and/c string? immutable?) #f)
  proc :  (or/c (integer-in 1 65535) #f)
          (or/c 'server 'client)
          . -> . any)
```

A parameter that specifies a procedure to be used (as is) by the default `sandbox-security-guard`. The default forbids all network connection.

```
(sandbox-exit-handler) → (any/c . -> . any)
(sandbox-exit-handler handler) → void?
  handler : (any/c . -> . any)
```

A parameter that determines the initial (exit-handler) for sandboxed evaluations. The default kills the evaluator with an appropriate error message (see exn:fail:sandbox-terminated-reason).

```
(sandbox-memory-limit) → (or/c (>=/c 0) #f)
(sandbox-memory-limit limit) → void?
  limit : (or/c (>=/c 0) #f)
```

A parameter that determines the total memory limit on the sandbox in megabytes (it can hold a rational or a floating point number). When this limit is exceeded, the sandbox is terminated. This value is used when the sandbox is created and the limit cannot be changed afterwards. It defaults to 30mb. See sandbox-eval-limits for per-evaluation limits and a description of how the two limits work together.

Note that (when memory accounting is enabled) memory is attributed to the highest custodian that refers to it. This means that if you inspect a value that sandboxed evaluation returns outside of the sandbox, your own custodian will be charged for it. To ensure that it is charged back to the sandbox, you should remove references to such values when the code is done inspecting it.

This policy has an impact on how the sandbox memory limit interacts with the per-expression limit specified by sandbox-eval-limits: values that are reachable from the sandbox, as well as from the interaction will count against the sandbox limit. For example, in the last interaction of this code,

```
(define e (make-evaluator 'racket/base))
(e '(define a 1))
(e '(for ([i (in-range 20)]) (set! a (cons (make-
bytes 500000) a))))
```

the memory blocks are allocated within the interaction limit, but since they're chained to the defined variable, they're also reachable from the sandbox — so they will count against the sandbox memory limit but not against the interaction limit (more precisely, no more than one block counts against the interaction limit).

```
                          (or/c (list/c (or/c (>=/c 0) #f)
(sandbox-eval-limits) →                 (or/c (>=/c 0) #f))
                                #f)
(sandbox-eval-limits limits) → void?
           (or/c (list/c (or/c (>=/c 0) #f)
  limits :               (or/c (>=/c 0) #f))
                 #f)
```

A parameter that determines the default limits on *each* use of a make-evaluator function, including the initial evaluation of the input program. Its value should be a list of two numbers; where the first is a shallow time value in seconds, and the second is a memory limit in

megabytes (note that they don't have to be integers). Either one can be `#f` for disabling the corresponding limit; alternately, the parameter can be set to `#f` to disable all per-evaluation limits (useful in case more limit kinds are available in future versions). The default is `(list 30 20)`.

Note that these limits apply to the creation of the sandbox environment too — even `(make-evaluator 'racket/base)` can fail if the limits are strict enough. For example,

```
(parameterize ([sandbox-eval-limits '(0.25 5)])
  (make-evaluator 'racket/base '(sleep 2)))
```

will throw an error instead of creating an evaluator. Therefore, to avoid surprises you need to catch errors that happen when the sandbox is created.

When limits are set, `call-with-limits` (see below) is wrapped around each use of the evaluator, so consuming too much time or memory results in an exception. Change the limits of a running evaluator using `set-eval-limits`.

A custodian's limit is checked only after a garbage collection, except that it may also be checked during certain large allocations that are individually larger than the custodian's limit.

The memory limit that is specified by this parameter applies to each individual evaluation, but not to the whole sandbox — that limit is specified via `sandbox-memory-limit`. When the global limit is exceeded, the sandbox is terminated, but when the per-evaluation limit is exceeded the `exn:fail:resource` exception is raised. For example, say that you evaluate an expression like

```
(for ([i (in-range 1000)])
  (set! a (cons (make-bytes 1000000) a))
  (collect-garbage))
```

then, assuming sufficiently small limits,

- if a global limit is set but no per-evaluation limit, the sandbox will eventually be terminated and no further evaluations possible;
- if there is a per-evaluation limit, but no global limit, the evaluation will abort with an error and it can be used again — specifically, `a` will still hold a number of blocks, and you can evaluate the same expression again which will add more blocks to it;
- if both limits are set, with the global one larger than the per-evaluation limit, then the evaluation will abort and you will be able to repeat it, but doing so several times will eventually terminate the sandbox (this will be indicated by the error message, and by the `evaluator-alive?` predicate).

`(sandbox-eval-handlers)`

```
  (list/c (or/c #f ((-> any) . -> . any))
→         (or/c #f ((-> any) . -> . any)))
(sandbox-eval-handlers handlers) → void?
  handlers : (list/c (or/c #f ((-> any) . -> . any))
                     (or/c #f ((-> any) . -> . any)))
```

A parameter that determines two (optional) handlers that wrap sandboxed evaluations. The first one is used when evaluating the initial program when the sandbox is being set-up, and the second is used for each interaction. Each of these handlers should expect a thunk as an argument, and they should execute these thunks — possibly imposing further restrictions. The default values are `#f` and `call-with-custodian-shutdown`, meaning no additional restrictions on initial sandbox code (e.g., it can start background threads), and a custodian-shutdown around each interaction that follows. Another useful function for this is `call-with-killing-threads` which kills all threads, but leaves other resources intact.

```
(sandbox-run-submodules) → (list/c symbol?)
(sandbox-run-submodules submod-syms) → void?
  submod-syms : (list/c symbol?)
```

A parameter that determines submodules to run when a sandbox is created by `make-module-evaluator`. The parameter's default value is the empty list.

```
(sandbox-make-inspector) → (-> inspector?)
(sandbox-make-inspector make) → void?
  make : (-> inspector?)
```

A parameter that determines the (nullary) procedure that is used to create the inspector for sandboxed evaluation. The procedure is called when initializing an evaluator. The default parameter value is `(lambda () (make-inspector (current-inspector)))`.

```
(sandbox-make-code-inspector) → (-> inspector?)
(sandbox-make-code-inspector make) → void?
  make : (-> inspector?)
```

A parameter that determines the (nullary) procedure that is used to create the code inspector for sandboxed evaluation. The procedure is called when initializing an evaluator. The default parameter value is `(lambda () (make-inspector (current-code-inspector)))`.

The `current-load/use-compiled` handler is setup to allow loading of bytecode files under the original code inspector when `sandbox-path-permissions` allows it through a `'read-bytecode` mode symbol, which makes loading libraries possible.

```
(sandbox-make-logger) → (-> logger?)
(sandbox-make-logger make) → void?
  make : (-> logger?)
```

A parameter that determines the procedure used to create the logger for sandboxed evaluation. The procedure is called when initializing an evaluator, and the default parameter value is `current-logger`. This means that it is not creating a new logger (this might change in the future).

```
(sandbox-make-plumber) → (or/c (-> plumber?) 'propagate)
(sandbox-make-plumber make) → void?
  make : (or/c (-> plumber?) 'propagate)
```

A parameter that determines the procedure used to create the plumber for sandboxed evaluation. The procedure is called when initializing an evaluator.

If the value is `'propagate` (the default), then a new plumber is created, and a flush callback is added to the current plumber to propagate the request to the new plumber within the created sandbox (if the sandbox has not already terminated).

Added in version 6.0.1.8 of package `sandbox-lib`.

```
(sandbox-make-environment-variables)
 → (-> environment-variables?)
(sandbox-make-environment-variables make) → void?
  make : (-> environment-variables?)
```

A parameter that determines the procedure used to create the environment variable set for sandboxed evaluation. The procedure is called when initializing an evaluator, and the default parameter value constructs a new environment variable set using `(environment-variables-copy (current-environment-variables))`.

14.12.2 Interacting with Evaluators

The following functions are used to interact with a sandboxed evaluator in addition to using it to evaluate code.

```
(evaluator-alive? evaluator) → boolean?
  evaluator : (any/c . -> . any)
```

Determines whether the evaluator is still alive.

```
(kill-evaluator evaluator) → void?
  evaluator : (any/c . -> . any)
```

Releases the resources that are held by *evaluator* by shutting down the evaluator's custodian. Attempting to use an evaluator after killing raises an exception, and attempts to kill a dead evaluator are ignored.

Killing an evaluator is similar to sending an `eof` value to the evaluator, except that an `eof` value will raise an error immediately.

```
(break-evaluator evaluator) → void?
  evaluator : (any/c . -> . any)
```

Sends a break to the running evaluator. The effect of this is as if Ctrl-C was typed when the evaluator is currently executing, which propagates the break to the evaluator's context.

```
(get-user-custodian evaluator) → void?
  evaluator : (any/c . -> . any)
```

Retrieves the *evaluator*'s toplevel custodian. This returns a value that is different from `(evaluator '(current-custodian))` or `call-in-sandbox-context evaluator current-custodian` — each sandbox interaction is wrapped in its own custodian, which is what these would return.

(One use for this custodian is with `current-memory-use`, where the per-interaction sub-custodians will not be charged with the memory for the whole sandbox.)

```
(set-eval-limits evaluator secs mb) → void?
  evaluator : (any/c . -> . any)
  secs : (or/c exact-nonnegative-integer? #f)
  mb : (or/c exact-nonnegative-integer? #f)
```

Changes the per-expression limits that *evaluator* uses to *secs* seconds of shallow time and *mb* megabytes (either one can be `#f`, indicating no limit).

This procedure should be used to modify an existing evaluator limits, because changing the `sandbox-eval-limits` parameter does not affect existing evaluators. See also `call-with-limits`.

```
(set-eval-handler evaluator handler) → void?
  evaluator : (any/c . -> . any)
  handler : (or/c #f ((-> any) . -> . any))
```

Changes the per-expression handler that the *evaluator* uses around each interaction. A `#f` value means no handler is used.

This procedure should be used to modify an existing evaluator handler, because changing the `sandbox-eval-handlers` parameter does not affect existing evaluators. See also `call-with-custodian-shutdown` and `call-with-killing-threads` for two useful handlers that are provided.

```
(call-with-custodian-shutdown thunk) → any
  thunk : (-> any)
(call-with-killing-threads thunk) → any
  thunk : (-> any)
```

These functions are useful for use as an evaluation handler. `call-with-custodian-shutdown` will execute the *thunk* in a fresh custodian, then shutdown that custodian, making sure that *thunk* could not have left behind any resources. `call-with-killing-threads` is similar, except that it kills threads that were left, but leaves other resources as is.

```
(put-input evaluator) → output-port?
  evaluator : (any/c . -> . any)
(put-input evaluator i/o) → void?
  evaluator : (any/c . -> . any)
  i/o : (or/c bytes? string? eof-object?)
```

If `(sandbox-input)` is `'pipe` when an evaluator is created, then this procedure can be used to retrieve the output port end of the pipe (when used with no arguments), or to add a string or a byte string into the pipe. It can also be used with `eof`, which closes the pipe.

```
(get-output evaluator) → (or/c #f input-port? bytes? string?)
  evaluator : (any/c . -> . any)
(get-error-output evaluator)
 → (or/c #f input-port? bytes? string?)
  evaluator : (any/c . -> . any)
```

Returns the output or error-output of the *evaluator*, in a way that depends on the setting of `(sandbox-output)` or `(sandbox-error-output)` when the evaluator was created:

- if it was `'pipe`, then `get-output` returns the input port end of the created pipe;
- if it was `'bytes` or `'string`, then the result is the accumulated output, and the output port is reset so each call returns a different piece of the evaluator's output (note that results are available only until the evaluator has terminated, and any allocations of the output are subject to the sandbox memory limit);
- otherwise, it returns `#f`.

```
(get-uncovered-expressions evaluator
                           [prog?
                            src])      → (listof syntax?)
  evaluator : (any/c . -> . any)
  prog? : any/c = #t
  src : any/c = default-src
```

Retrieves uncovered expression from an evaluator, as longs as the `sandbox-coverage-enabled` parameter had a true value when the evaluator was created. Otherwise, an exception is raised to indicate that no coverage information is available.

The *prog?* argument specifies whether to obtain expressions that were uncovered after only the original input program was evaluated (`#t`) or after all later uses of the evaluator (`#f`). Using `#t` retrieves a list that is saved after the input program is evaluated, and before the evaluator is used, so the result is always the same.

A `#t` value of *prog?* is useful for testing student programs to find out whether a submission has sufficient test coverage built in. A `#f` value is useful for writing test suites for a program to ensure that your tests cover the whole code.

The second optional argument, *src*, specifies that the result should be filtered to hold only syntax objects whose source matches *src*. The default is the source that was used in the program code, if there was one. Note that `'program` is used as the source value if the input program was given as S-expressions or as a string (and in these cases it will be the default for filtering). If given `#f`, the result is the unfiltered list of expressions.

The resulting list of syntax objects has at most one expression for each position and span. Thus, the contents may be unreliable, but the position information is reliable (i.e., it always indicates source code that would be painted red in DrRacket when coverage information is used).

Note that if the input program is a sequence of syntax values, either make sure that they have `'program` as the source field, or use the *src* argument. Using a sequence of S-expressions (not syntax objects) for an input program leads to unreliable coverage results, since each expression may be assigned a single source location.

```
(call-in-sandbox-context  evaluator
                          thunk
                         [unrestricted?]) → any
  evaluator : (any/c . -> . any)
  thunk : (-> any)
  unrestricted? : boolean? = #f
```

Calls the given *thunk* in the context of a sandboxed evaluator. The call is performed under the resource limits and evaluation handler that are used for evaluating expressions, unless *unrestricted?* is specified as true.

This process is usually similar to `(evaluator (list thunk))`, except that it does not rely on the common meaning of a sexpr-based syntax with list expressions as function application (which is not true in all languages). Note that this is more useful for meta-level operations such as namespace manipulation, it is not intended to be used as a safe-evaluation replacement (i.e., using the sandbox evaluator as usual).

In addition, you can avoid some of the sandboxed restrictions by using your own permissions, for example,

```
(let ([guard (current-security-guard)])
```

```
(call-in-sandbox-context
  ev
  (lambda ()
    (parameterize ([current-security-guard guard])
      ; can access anything you want here
      (delete-file "/some/file")))))
```

14.12.3 Miscellaneous

```
gui? : boolean?
```

For backward compatibility, only: the result of `gui-available?` at the time that `racket/sandbox` was instantiated.

The value of `gui?` is no longer used by `racket/sandbox` itself. Instead, `gui-available?` and `sandbox-gui-available` are checked at the time that a sandbox evaluator is created.

```
(call-with-limits secs mb thunk) → any
  secs : (or/c exact-nonnegative-integer? #f)
  mb : (or/c exact-nonnegative-integer? #f)
  thunk : (-> any)
```

Executes the given *thunk* with memory and time restrictions: if execution consumes more than *mb* megabytes or more than *secs* shallow time seconds, then the computation is aborted and the `exn:fail:resource` exception is raised. Otherwise the result of the thunk is returned as usual (a value, multiple values, or an exception). Each of the two limits can be `#f` to indicate the absence of a limit. See also `custodian-limit-memory` for information on memory limits.

Sandboxed evaluators use `call-with-limits`, according to the `sandbox-eval-limits` setting and uses of `set-eval-limits`: each expression evaluation is protected from timeouts and memory problems. Use `call-with-limits` directly only to limit a whole testing session, instead of each expression.

```
(with-limits sec-expr mb-expr body ...)
```

A macro version of `call-with-limits`.

```
(call-with-deep-time-limit secs thunk) → any
  secs : exact-nonnegative-integer?
  thunk : (-> any)
```

Executes the given *thunk* with deep time restrictions.

```
(with-deep-time-limit secs-expr body ...)
```

A macro version of `call-with-deep-time-limit`.

```
(exn:fail:resource? v) → boolean?
  v : any/c
(exn:fail:resource-resource exn)
 → (or/c 'time 'memory 'deep-time)
  exn : exn:fail:resource?
```

A predicate and accessor for exceptions that are raised by `call-with-limits`. The `resource` field holds a symbol, representing the resource that was expended. `'time` is used for shallow time and `'deep-time` is used for deep time.

15 Operating System

15.1 Paths

When a Racket procedure takes a filesystem path as an argument, the path can be provided either as a string or as an instance of the *path* datatype. If a string is provided, it is converted to a path using `string->path`. A Racket procedure that generates a filesystem path always generates a path value.

By default, paths are created and manipulated for the current platform, but procedures that merely manipulate paths (without using the filesystem) can manipulate paths using conventions for other supported platforms. The `bytes->path` procedure accepts an optional argument that indicates the platform for the path, either `'unix` or `'windows`. For other functions, such as `build-path` or `simplify-path`, the behavior is sensitive to the kind of path that is supplied. Unless otherwise specified, a procedure that requires a path accepts only paths for the current platform.

Two path values are `equal?` when they are use the same convention type and when their byte-string representations are `equal?`. A path string (or byte string) cannot be empty, and it cannot contain a nul character or byte. When an empty string or a string containing nul is provided as a path to any procedure except `absolute-path?`, `relative-path?`, or `complete-path?`, the `exn:fail:contract` exception is raised.

Most Racket primitives that accept paths first *cleanse* the path before using it. Procedures that build paths or merely check the form of a path do not cleanse paths, with the exceptions of `cleanse-path`, `expand-user-path`, and `simplify-path`. For more information about path cleansing and other platform-specific details, see §15.1.3 “Unix and Mac OS X Paths” for Unix and Mac OS X paths and §15.1.4 “Windows Path Conventions” for Windows paths.

15.1.1 Manipulating Paths

```
(path? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a path value for the current platform (not a string, and not a path for a different platform), `#f` otherwise.

```
(path-string? v) → boolean?
  v : any/c
```

Return `#t` if `v` is either a path value for the current platform or a non-empty string without nul characters, `#f` otherwise.

```
(path-for-some-system? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a path value for some platform (not a string), `#f` otherwise.

```
(string->path str) → path?
  str : string?
```

Produces a path whose byte-string name is `(string->bytes/locale string (char->integer #\?))`.

Beware that the current locale might not encode every string, in which case `string->path` can produce the same path for different *str*s. See also `string->path-element`, which should be used instead of `string->path` when a string represents a single path element.

See also `string->some-system-path`.

```
(bytes->path bstr [type]) → path?
  bstr : bytes?
  type : (or/c 'unix 'windows) = (system-path-convention-type)
```

Produces a path (for some platform) whose byte-string name is *bstr*. The optional *type* specifies the convention to use for the path.

For converting relative path elements from literals, use instead `bytes->path-element`, which applies a suitable encoding for individual elements.

```
(path->string path) → string?
  path : path?
```

Produces a string that represents *path* by decoding *path*'s byte-string name using the current locale's encoding; `?` is used in the result string where encoding fails, and if the encoding result is the empty string, then the result is `"?"`.

The resulting string is suitable for displaying to a user, string-ordering comparisons, etc., but it is not suitable for re-creating a path (possibly modified) via `string->path`, since decoding and re-encoding the path's byte string may lose information.

Furthermore, for display and sorting based on individual path elements (such as pathless file names), use `path-element->string`, instead, to avoid special encodings use to represent some relative paths. See §15.1.4 "Windows Path Conventions" for specific information about the conversion of Windows paths.

See also `some-system-path->string`.

```
(path->bytes path) → bytes?
  path : path-for-some-system?
```

Produces *path*'s byte string representation. No information is lost in this translation, so that (bytes->path (path->bytes *path*) (path-convention-type *path*)) always produces a path that is equal? to *path*. The *path* argument can be a path for any platform.

Conversion to and from byte values is useful for marshaling and unmarshaling paths, but manipulating the byte form of a path is generally a mistake. In particular, the byte string may start with a \\?\REL encoding for Windows paths. Instead of path->bytes, use split-path and path-element->bytes to manipulate individual path elements.

```
(string->path-element str) → path?
  str : string?
```

Like string->path, except that *str* corresponds to a single relative element in a path, and it is encoded as necessary to convert it to a path. See §15.1.3 "Unix and Mac OS X Paths" for more information on the conversion for Unix and Mac OS X paths, and see §15.1.4 "Windows Path Conventions" for more information on the conversion for Windows paths.

If *str* does not correspond to any path element (e.g., it is an absolute path, or it can be split), or if it corresponds to an up-directory or same-directory indicator on Unix and Mac OS X, then exn:fail:contract exception is raised.

As for path->string, information can be lost from *str* in the locale-specific conversion to a path.

```
(bytes->path-element bstr [type]) → path-for-some-system?
  bstr : bytes?
  type : (or/c 'unix 'windows) = (system-path-convention-type)
```

Like bytes->path, except that *bstr* corresponds to a single relative element in a path. In terms of conversions and restrictions on *bstr*, bytes->path-element is like string->path-element.

The bytes->path-element procedure is generally the best choice for reconstructing a path based on another path (where the other path is deconstructed with split-path and path-element->bytes) when ASCII-level manipulation of path elements is necessary.

```
(path-element->string path) → string?
  path : path-element?
```

Like path->string, except that trailing path separators are removed (as by split-path). On Windows, any \\?\REL encoding prefix is also removed; see §15.1.4 "Windows Path Conventions" for more information on Windows paths.

The *path* argument must be such that split-path applied to *path* would return 'relative as its first result and a path as its second result, otherwise the exn:fail:contract exception is raised.

The path-element->string procedure is generally the best choice for presenting a path-less file or directory name to a user.

```
(path-element->bytes path) → bytes?
  path : path-element?
```

Like path->bytes, except that any encoding prefix is removed, etc., as for path-element->string.

For any reasonable locale, consecutive ASCII characters in the printed form of *path* are mapped to consecutive byte values that match each character's code-point value, and a leading or trailing ASCII character is mapped to a leading or trailing byte, respectively. The *path* argument can be a path for any platform.

The path-element->bytes procedure is generally the right choice (in combination with split-path) for extracting the content of a path to manipulate it at the ASCII level (then reassembling the result with bytes->path-element and build-path).

```
(path<? a-path b-path ...) → boolean?
  a-path : path?
  b-path : path?
```

Returns #t if the arguments are sorted, where the comparison for each pair of paths is the same as using path->bytes and bytes<?.

```
(path-convention-type path) → (or/c 'unix 'windows)
  path : path-for-some-system?
```

Accepts a path value (not a string) and returns its convention type.

```
(system-path-convention-type) → (or/c 'unix 'windows)
```

Returns the path convention type of the current platform: 'unix for Unix and Mac OS X, 'windows for Windows.

```
(build-path base sub ...) → path-for-some-system?
  base : (or/c path-string? path-for-some-system? 'up 'same)
       (or/c (and/c (or/c path-string? path-for-some-system?)
  sub :              (not/c complete-path?))
             (or/c 'up 'same))
```

Creates a path given a base path and any number of sub-path extensions. If *base* is an absolute path, the result is an absolute path, otherwise the result is a relative path.

The *base* and each *sub* must be either a relative path, the symbol 'up (indicating the relative parent directory), or the symbol 'same (indicating the relative current directory). For Windows paths, if *base* is a drive specification (with or without a trailing slash) the first *sub* can be an absolute (driveless) path. For all platforms, the last *sub* can be a filename.

The *base* and *sub* arguments can be paths for any platform. The platform for the resulting path is inferred from the *base* and *sub* arguments, where string arguments imply a path for the current platform. If different arguments are for different platforms, the exn:fail:contract exception is raised. If no argument implies a platform (i.e., all are 'up or 'same), the generated path is for the current platform.

Each *sub* and *base* can optionally end in a directory separator. If the last *sub* ends in a separator, it is included in the resulting path.

If *base* or *sub* is an illegal path string (because it is empty or contains a nul character), the exn:fail:contract exception is raised.

The build-path procedure builds a path *without* checking the validity of the path or accessing the filesystem.

See §15.1.3 “Unix and Mac OS X Paths” for more information on the construction of Unix and Mac OS X paths, and see §15.1.4 “Windows Path Conventions” for more information on the construction of Windows paths.

The following examples assume that the current directory is "/home/joeuser" for Unix examples and "C:\Joe's Files" for Windows examples.

```
(define p1 (build-path (current-directory) "src" "racket"))
 ; Unix: p1 is "/home/joeuser/src/racket"
 ; Windows: p1 is "C:\\Joe's Files\\src\\racket"
(define p2 (build-path 'up 'up "docs" "Racket"))
 ; Unix: p2 is "../../docs/Racket"
 ; Windows: p2 is "..\\..\\docs\\Racket"
(build-path p2 p1)
 ; Unix and Windows: raises exn:fail:contract; p1 is absolute
(build-path p1 p2)
 ; Unix: is "/home/joeuser/src/racket/../../docs/Racket"
 ; Windows: is "C:\\Joe's Files\\src\\racket\\..\\..\\docs\\Racket"
```

```
(build-path/convention-type type
                            base
                            sub ...) → path-for-some-system?
  type : (or/c 'unix 'windows)
  base : (or/c path-string? path-for-some-system? 'up 'same)
       (or/c (and/c (or/c path-string? path-for-some-system?)
  sub :              (not/c complete-path?))
             (or/c 'up 'same))
```

Like build-path, except a path convention type is specified explicitly.

```
(absolute-path? path) → boolean?
  path : (or/c path-string? path-for-some-system?)
```

Returns #t if path is an absolute path, #f otherwise. The path argument can be a path for any platform. If path is not a legal path string (e.g., it contains a nul character), #f is returned. This procedure does not access the filesystem.

```
(relative-path? path) → boolean?
  path : (or/c path-string? path-for-some-system?)
```

Returns #t if path is a relative path, #f otherwise. The path argument can be a path for any platform. If path is not a legal path string (e.g., it contains a nul character), #f is returned. This procedure does not access the filesystem.

```
(complete-path? path) → boolean?
  path : (or/c path-string? path-for-some-system?)
```

Returns #t if path is a *complete*ly determined path (*not* relative to a directory or drive), #f otherwise. The path argument can be a path for any platform. Note that for Windows paths, an absolute path can omit the drive specification, in which case the path is neither relative nor complete. If path is not a legal path string (e.g., it contains a nul character), #f is returned.

This procedure does not access the filesystem.

```
(path->complete-path path [base]) → path-for-some-system?
  path : (or/c path-string? path-for-some-system?)
  base : (or/c path-string? path-for-some-system?)
       = (current-directory)
```

Returns path as a complete path. If path is already a complete path, it is returned as the result. Otherwise, path is resolved with respect to the complete path base. If base is not a complete path, the exn:fail:contract exception is raised.

The path and base arguments can be paths for any platform; if they are for different platforms, the exn:fail:contract exception is raised.

This procedure does not access the filesystem.

```
(path->directory-path path) → path-for-some-system?
  path : (or/c path-string? path-for-some-system?)
```

Returns path if path syntactically refers to a directory and ends in a separator, otherwise it returns an extended version of path that specifies a directory and ends with a separator.

For example, on Unix and Mac OS X, the path `"x/y/"` syntactically refers to a directory and ends in a separator, but `"x/y"` would be extended to `"x/y/"`, and `"x/.."` would be extended to `"x/../"`. The *path* argument can be a path for any platform, and the result will be for the same platform.

This procedure does not access the filesystem.

```
(resolve-path path) → path?
  path : path-string?
```

Cleanses *path* and returns a path that references the same file or directory as *path*. If *path* is a soft link to another path, then the referenced path is returned (this may be a relative path with respect to the directory owning *path*), otherwise *path* is returned (after expansion).

On Windows, the path for a link should be simplified syntactically, so that an up-directory indicator removes a preceding path element independent of whether the preceding element itself refers to a link. For relative-paths links, the path should be parsed specially; see §15.1.4 “Windows Path Conventions” for more information.

Changed in version 6.0.1.12 of package `base`: Added support for links on Windows.

```
(cleanse-path path) → path-for-some-system?
  path : (or/c path-string? path-for-some-system?)
```

Cleanses *path* (as described at the beginning of this chapter) without consulting the filesystem.

```
(expand-user-path path) → path?
  path : path-string?
```

Cleanses *path*. In addition, on Unix and Mac OS X, a leading ~ is treated as user’s home directory and expanded; the username follows the ~ (before a / or the end of the path), where ~ by itself indicates the home directory of the current user.

```
(simplify-path path [use-filesystem?]) → path-for-some-system?
  path : (or/c path-string? path-for-some-system?)
  use-filesystem? : boolean? = #t
```

Eliminates redundant path separators (except for a single trailing separator), up-directory `..`, and same-directory `.` indicators in *path*, and changes / separators to \ separators in Windows paths, such that the result accesses the same file or directory (if it exists) as *path*.

In general, the pathname is normalized as much as possible—without consulting the filesystem if *use-filesystem?* is `#f`, and (on Windows) without changing the case of letters within the path. If *path* syntactically refers to a directory, the result ends with a directory separator.

When *path* is simplified and *use-filesystem?* is true (the default), a complete path is returned. If *path* is relative, it is resolved with respect to the current directory. On Unix and Mac OS X, up-directory indicators are removed taking into account soft links (so that the resulting path refers to the same directory as before); on Windows, up-directory indicators are removed by by deleting a preceding path element.

When *use-filesystem?* is `#f`, up-directory indicators are removed by deleting a preceding path element, and the result can be a relative path with up-directory indicators remaining at the beginning of the path; up-directory indicators are dropped when they refer to the parent of a root directory. Similarly, the result can be the same as `(build-path 'same)` (but with a trailing separator) if eliminating up-directory indicators leaves only same-directory indicators.

The *path* argument can be a path for any platform when *use-filesystem?* is `#f`, and the resulting path is for the same platform.

The filesystem might be accessed when *use-filesystem?* is true, but the source or simplified path might be a non-existent path. If *path* cannot be simplified due to a cycle of links, the `exn:fail:filesystem` exception is raised (but a successfully simplified path may still involve a cycle of links if the cycle did not inhibit the simplification).

See §15.1.3 “Unix and Mac OS X Paths” for more information on simplifying Unix and Mac OS X paths, and see §15.1.4 “Windows Path Conventions” for more information on simplifying Windows paths.

```
(normal-case-path path) → path-for-some-system?
  path : (or/c path-string? path-for-some-system?)
```

Returns *path* with “normalized” case letters. For Unix and Mac OS X paths, this procedure always returns the input path, because filesystems for these platforms can be case-sensitive. For Windows paths, if *path* does not start `\\?\`, the resulting string uses only lowercase letters, based on the current locale. In addition, for Windows paths when the path does not start `\\?\`, all `/`s are converted to `\`s, and trailing spaces and `.`s are removed.

The *path* argument can be a path for any platform, but beware that local-sensitive decoding and conversion of the path may be different on the current platform than for the path’s platform.

This procedure does not access the filesystem.

```
                      (or/c path-for-some-system? 'relative #f)
(split-path path) → (or/c path-for-some-system? 'up 'same)
                      boolean?
  path : (or/c path-string? path-for-some-system?)
```

Deconstructs *path* into a smaller path and an immediate directory or file name. Three values are returned:

- `base` is either
  - a path,
  - `'relative` if *path* is an immediate relative directory or filename, or
  - `#f` if *path* is a root directory.
- `name` is either
  - a directory-name path,
  - a filename,
  - `'up` if the last part of *path* specifies the parent directory of the preceding path (e.g., `..` on Unix), or
  - `'same` if the last part of *path* specifies the same directory as the preceding path (e.g., `.` on Unix).
- `must-be-dir?` is `#t` if *path* explicitly specifies a directory (e.g., with a trailing separator), `#f` otherwise. Note that `must-be-dir?` does not specify whether `name` is actually a directory or not, but whether *path* syntactically specifies a directory.

Compared to *path*, redundant separators (if any) are removed in the result `base` and `name`. If `base` is `#f`, then `name` cannot be `'up` or `'same`. The *path* argument can be a path for any platform, and resulting paths for the same platform.

This procedure does not access the filesystem.

See §15.1.3 “Unix and Mac OS X Paths” for more information on splitting Unix and Mac OS X paths, and see §15.1.4 “Windows Path Conventions” for more information on splitting Windows paths.

```
(explode-path path)
 → (listof (or/c path-for-some-system? 'up 'same))
  path : (or/c path-string? path-for-some-system?)
```

Returns the list of path elements that constitute *path*. If *path* is simplified in the sense of `simple-form-path`, then the result is always a list of paths, and the first element of the list is a root.

The `explode-path` function computes its result in time proportional to the length of *path* (unlike a loop in that uses `split-path`, which must allocate intermediate paths).

```
(path-replace-suffix path suffix) → path-for-some-system?
  path : (or/c path-string? path-for-some-system?)
  suffix : (or/c string? bytes?)
```

Returns a path that is the same as *path*, except that the suffix for the last element of the path is changed to *suffix*. If the last element of *path* has no suffix, then *suffix* is added to

the path. A suffix is defined as a `.` followed by any number of non-`.` characters/bytes at the end of the path element, as long as the path element is not `".."` or `"."`. The *path* argument can be a path for any platform, and the result is for the same platform. If *path* represents a root, the `exn:fail:contract` exception is raised.

```
(path-add-suffix path suffix) → path-for-some-system?
  path : (or/c path-string? path-for-some-system?)
  suffix : (or/c string? bytes?)
```

Similar to `path-replace-suffix`, but any existing suffix on *path* is preserved by replacing every `.` in the last path element with `_`, and then the *suffix* is added to the end.

```
(reroot-path path root-path) → path-for-some-system?
  path : (or/c path-string? path-for-some-system?)
  root-path : (or/c path-string? path-for-some-system?)
```

Produces a path that extends *root-path* based on the complete form of *path*.

If *path* is not already complete, is it completed via `path->complete-path`, in which case *path* must be a path for the current platform. The *path* argument is also cleansed and case-normalized via `normal-case-path`. The path is then appended to *root-path*; in the case of Windows paths, a root letter drive becomes a letter path element, while a root UNC path is prefixed with `"UNC"` as a path element and the machine and volume names become path elements.

Examples:

```
> (reroot-path (bytes->path #"/home/caprica/baltar" 'unix)
               (bytes->path #"/earth" 'unix))
#<path:/earth/home/caprica/baltar>
> (reroot-path (bytes->path #"c:\\usr\\adama" 'windows)
               (bytes->path #"\\\\earth\\africa\\" 'windows))
#<windows-path:\\earth\africa\c\usr\adama>
> (reroot-path (bytes->path #"\\\\galactica\\cac\\adama" 'windows)
               (bytes->path #"s:\\earth\\africa\\" 'windows))
#<windows-path:s:\earth\africa\UNC\galactica\cac\adama>
```

15.1.2 More Path Utilities

`(require racket/path)` package: `base`

The bindings documented in this section are provided by the `racket/path` and `racket` libraries, but not `racket/base`.

```
(file-name-from-path path) → (or/c path-for-some-system? #f)
  path : (or/c path-string? path-for-some-system?)
```

Returns the last element of *path*. If *path* is syntactically a directory path (see split-path), then the result is #f.

```
(filename-extension path) → (or/c bytes? #f)
  path : (or/c path-string? path-for-some-system?)
```

Returns a byte string that is the extension part of the filename in *path* without the . separator. If *path* is syntactically a directory (see split-path) or if the path has no extension, #f is returned.

```
(find-relative-path  base
                     path
                    [#:more-than-root? more-than-root?])
 → path-for-some-system?
  base : (or/c path-string? path-for-some-system?)
  path : (or/c path-string?  path-for-some-system?)
  more-than-root? : any/c = #f
```

Finds a relative pathname with respect to *base* that names the same file or directory as *path*. Both *base* and *path* must be simplified in the sense of simple-form-path. If *path* shares no subpath in common with *base*, *path* is returned.

If *more-than-root?* is true, if *base* and *path* share only a Unix root in common, and if neither *base* nor *path* is just a root path, then *path* is returned.

```
(normalize-path path [wrt]) → path?
  path : path-string?
  wrt :  (and/c path-string? complete-path?)
      = (current-directory)
```

Returns a complete version of *path* by making the path complete, expanding the complete path, and resolving all soft links (which requires consulting the filesystem). If *path* is relative, then *wrt* is used as the base path.

For most purposes, simple-form-path is the preferred mechanism to normalize a path, because it works for paths that include non-existent directory components, and it avoids unnecessarily expanding soft links.

Letter case is *not* normalized by normalize-path. For this and other reasons, such as whether the path is syntactically a directory, the result of normalize-path is not suitable for comparisons that determine whether two paths refer to the same file or directory (i.e., the comparison may produce false negatives).

An error is signaled by normalize-path if the input path contains an embedded path for a non-existent directory, or if an infinite cycle of soft links is detected.

```
(path-element? path) → boolean?
  path : any/c
```

Returns #t if *path* is a *path element*: a path value for some platform (see path-for-some-system?) such that split-path applied to *path* would return 'relative as its first result and a path as its second result. Otherwise, the result is #f.

```
(path-only path) → (or/c #f path-for-some-system?)
  path : (or/c path-string? path-for-some-system?)
```

If *path* is a filename, the file's path is returned. If *path* is syntactically a directory, *path* is returned (as a path, if it was a string). If *path* has no directory part `#f` is returned.

```
(simple-form-path path) → path?
  path : path-string?
```

Returns `(simplify-path (path->complete-path path))`, which ensures that the result is a complete path containing no up- or same-directory indicators.

```
(some-system-path->string path) → string?
  path : path-for-some-system?
```

Converts *path* to a string using a UTF-8 encoding of the path's bytes.

Use this function when working with paths for a different system (whose encoding of pathnames might be unrelated to the current locale's encoding) and when starting and ending with strings.

```
(string->some-system-path str kind) → path-for-some-system?
  str : string?
  kind : (or/c 'unix 'windows)
```

Converts *str* to a *kind* path using a UTF-8 encoding of the path's bytes.

Use this function when working with paths for a different system (whose encoding of pathnames might be unrelated to the current locale's encoding) and when starting and ending with strings.

```
(shrink-path-wrt pth other-pths) → (or/c #f path?)
  pth : path?
  other-pths : (listof path?)
```

Returns a suffix of *pth* that shares nothing in common with the suffixes of *other-pths*, or *pth*, if not possible (e.g. when *other-pths* is empty or contains only paths with the same elements as *pth*).

Examples:

```
> (shrink-path-wrt (build-path "racket" "list.rkt")
                   (list (build-path "racket" "list.rkt")
                         (build-path "racket" "base.rkt")))
#<path:list.rkt>
```

```
> (shrink-path-wrt (build-path "racket" "list.rkt")
                   (list (build-path "racket" "list.rkt")
                         (build-path "racket" "private" "list.rkt")
                         (build-path "racket" "base.rkt")))
#<path:racket/list.rkt>
```

15.1.3 Unix and Mac OS X Paths

In Unix and Mac OS X paths, a `/` separates elements of the path, `.` as a path element always means the directory indicated by preceding path, and `..` as a path element always means the parent of the directory indicated by the preceding path. A leading `~` in a path is not treated specially, but `expand-user-path` can be used to convert a leading `~` element to a user-specific directory. No other character or byte has a special meaning within a path. Multiple adjacent `/` are equivalent to a single `/` (i.e., they act as a single path separator).

A path root is always `/`. A path starting with `/` is an absolute, complete path, and a path starting with any other character is a relative path.

Any pathname that ends with a `/` syntactically refers to a directory, as does any path whose last element is `.` or `..`.

A Unix and Mac OS X path is cleansed by replacing multiple adjacent `/`s with a single `/`.

For `(bytes->path-element` *bstr*`)`, `bstr` must not contain any `/`, otherwise the `exn:fail:contract` exception is raised. The result of `(path-element->bytes` *path*`)` or `(path-element->string` *path*`)` is always the same as the result of `(path->bytes` *path*`)` and `(path->string` *path*`)`. Since that is not the case for other platforms, however, `path-element->bytes` and `path-element->string` should be used when converting individual path elements.

On Mac OS X, Finder aliases are zero-length files.

15.1.4 Windows Path Conventions

In general, a Windows pathname consists of an optional drive specifier and a drive-specific path. A Windows path can be *absolute* but still relative to the current drive; such paths start with a `/` or `\` separator and are not UNC paths or paths that start with `\\?\`.

A path that starts with a drive specification is *complete*. Roughly, a drive specification is either a Latin letter followed by a colon, a UNC path of the form `\\`⟨*machine*⟩`\`⟨*volume*⟩, or a `\\?\` form followed by something other than `REL\`⟨*element*⟩, or `RED\`⟨*element*⟩. (Variants of `\\?\` paths are described further below.)

Racket fails to implement the usual Windows path syntax in one way. Outside of Racket,

a pathname "C:rant.txt" can be a drive-specific relative path. That is, it names a file "rant.txt" on drive "C:", but the complete path to the file is determined by the current working directory for drive "C:". Racket does not support drive-specific working directories (only a working directory across all drives, as reflected by the current-directory parameter). Consequently, Racket implicitly converts a path like "C:rant.txt" into "C:\rant.txt".

- *Racket-specific:* Whenever a path starts with a drive specifier ⟨*letter*⟩: that is not followed by a / or \, a \ is inserted as the path is cleansed.

Otherwise, Racket follows standard Windows path conventions, but also adds \\?\REL and \\?\RED conventions to deal with paths inexpressible in the standard convention, plus conventions to deal with excessive \s in \\?\ paths.

In the following, ⟨*letter*⟩ stands for a Latin letter (case does not matter), ⟨*machine*⟩ stands for any sequence of characters that does not include \ or / and is not ?, ⟨*volume*⟩ stands for any sequence of characters that does not include \ or / , and ⟨*element*⟩ stands for any sequence of characters that does not include \.

- Trailing spaces and . in a path element are ignored when the element is the last one in the path, unless the path starts with \\?\ or the element consists of only spaces and .s.
- The following special "files", which access devices, exist in all directories, case-insensitively, and with all possible endings after a period or colon, except in pathnames that start with \\?\: "NUL", "CON", "PRN", "AUX", "COM1", "COM2", "COM3", "COM4", "COM5", "COM6", "COM7", "COM8", "COM9", "LPT1", "LPT2", "LPT3", "LPT4", "LPT5", "LPT6", "LPT7", "LPT8", "LPT9".
- Except for \\?\ paths, /s are equivalent to \s. Except for \\?\ paths and the start of UNC paths, multiple adjacent /s and \s count as a single \. In a path that starts \\?\ paths, elements can be separated by either a single or double \.
- A directory can be accessed with or without a trailing separator. In the case of a non-\\?\ path, the trailing separator can be any number of /s and \s; in the case of a \\?\ path, a trailing separator must be a single \, except that two \s can follow \\?\⟨*letter*⟩:.
- Except for \\?\ paths, a single . as a path element means "the current directory," and a .. as a path element means "the parent directory." Up-directory path elements (i.e., ..) immediately after a drive are ignored.
- A pathname that starts \\⟨*machine*⟩\⟨*volume*⟩ (where a / can replace any \) is a UNC path, and the starting \\⟨*machine*⟩\⟨*volume*⟩ counts as the drive specifier.
- Normally, a path element cannot contain any of the following characters:

`< > : " / \ |`

Except for `\`, path elements containing these characters can be accessed using a `\\?\` path (assuming that the underlying filesystem allows the characters).

- In a pathname that starts `\\?\`⟨*letter*⟩`:\`, the `\\?\`⟨*letter*⟩`:\` prefix counts as the path's drive, as long as the path does not both contain non-drive elements and end with two consecutive `\`s, and as long as the path contains no sequence of three or more `\`s. Two `\`s can appear in place of the `\` before ⟨*letter*⟩. A `/` cannot be used in place of a `\` (but `/`s can be used in element names, though the result typically does not name an actual directory or file).

- In a pathname that starts `\\?\UNC\`⟨*machine*⟩`\`⟨*volume*⟩, the `\\?\UNC\`⟨*machine*⟩`\`⟨*volume*⟩ prefix counts as the path's drive, as long as the path does not end with two consecutive `\`s, and as long as the path contains no sequence of three or more `\`s. Two `\`s can appear in place of the `\` before `UNC`, the `\`s after `UNC`, and/or the `\`s after⟨*machine*⟩. The letters in the `UNC` part can be uppercase or lowercase, and `/` cannot be used in place of `\`s (but `/` can be used in element names).

- *Racket-specific:* A pathname that starts `\\?\REL\`⟨*element*⟩ or `\\?\REL\\`⟨*element*⟩ is a relative path, as long as the path does not end with two consecutive `\`s, and as long as the path contains no sequence of three or more `\`s. This Racket-specific path form supports relative paths with elements that are not normally expressible in Windows paths (e.g., a final element that ends in a space). The `REL` part must be exactly the three uppercase letters, and `/`s cannot be used in place of `\`s. If the path starts `\\?\REL\..` then for as long as the path continues with repetitions of `\..`, each element counts as an up-directory element; a single `\` must be used to separate the up-directory elements. As soon as a second `\` is used to separate the elements, or as soon as a non-`..` element is encountered, the remaining elements are all literals (never up-directory elements). When a `\\?\REL` path value is converted to a string (or when the path value is written or displayed), the string does not contain the starting `\\?\REL` or the immediately following `\`s; converting a path value to a byte string preserves the `\\?\REL` prefix.

- *Racket-specific:* A pathname that starts `\\?\RED\`⟨*element*⟩ or `\\?\RED\\`⟨*element*⟩ is a drive-relative path, as long as the path does not end with two consecutive `\`s, and as long as the path contains no sequence of three or more `\`s. This Racket-specific path form supports drive-relative paths (i.e., absolute given a drive) with elements that are not normally expressible in Windows paths. The `RED` part must be exactly the three uppercase letters, and `/`s cannot be used in place of `\`s. Unlike `\\?\REL` paths, a `..` element is always a literal path element. When a `\\?\RED` path value is converted to a string (or when the path value is written or displayed), the string does not contain the starting `\\?\RED` and it contains a single starting `\`; converting a path value to a byte string preserves the `\\?\RED` prefix.

Three additional Racket-specific rules provide meanings to character sequences that are otherwise ill-formed as Windows paths:

- *Racket-specific:* In a pathname of the form `\\?\`⟨*any*⟩`\\` where ⟨*any*⟩ is any non-empty sequence of characters other than ⟨*letter*⟩`:` or `\`⟨*letter*⟩`:`, the entire path counts as the path's (non-existent) drive.

- *Racket-specific:* In a pathname of the form `\\?\`⟨*any*⟩`\\\`⟨*elements*⟩, where ⟨*any*⟩ is any non-empty sequence of characters and ⟨*elements*⟩ is any sequence that does not start with a `\`, does not end with two `\`s, and does not contain a sequence of three `\`s, then `\\?\`⟨*any*⟩`\\` counts as the path's (non-existent) drive.

- *Racket-specific:* In a pathname that starts `\\?\` and does not match any of the patterns from the preceding bullets, `\\?\` counts as the path's (non-existent) drive.

Outside of Racket, except for `\\?\` paths, pathnames are typically limited to 259 characters. Racket internally converts pathnames to `\\?\` form as needed to avoid this limit. The operating system cannot access files through `\\?\` paths that are longer than 32,000 characters or so.

Where the above descriptions says "character," substitute "byte" for interpreting byte strings as paths. The encoding of Windows paths into bytes preserves ASCII characters, and all special characters mentioned above are ASCII, so all of the rules are the same.

Beware that the `\` path separator is an escape character in Racket strings. Thus, the path `\\?\REL\..\\..` as a string must be written `"\\\\?\\REL\\..\\\\.."`.

A path that ends with a directory separator syntactically refers to a directory. In addition, a path syntactically refers to a directory if its last element is a same-directory or up-directory indicator (not quoted by a `\\?\` form), or if it refers to a root.

Even on variants of Windows that support symbolic links, up-directory `..` indicators in a path are resolved syntactically, not sensitive to links. For example, if a path ends with `d\..\f` and `d` refers to a symbolic link that references a directory with a different parent than `d`, the path nevertheless refers to `f` in the same directory as `d`. A relative-path link is parsed as if prefixed with `\\?\REL` paths, except that `..` and `.` elements are allowed throughout the path, and any number of redundant `/` separators are allowed.

Windows paths are cleansed as follows: In paths that start `\\?\`, redundant `\`s are removed, an extra `\` is added in a `\\?\REL` if an extra one is not already present to separate up-directory indicators from literal path elements, and an extra `\` is similarly added after `\\?\RED` if an extra one is not already present. For other paths, multiple `/`s and `\`s are converted to single `/`s or `\` (except at the beginning of a shared folder name), and a `\` is inserted after the colon in a drive specification if it is missing.

For `(bytes->path-element` *`bstr`*`)`, `/`s, colons, trailing dots, trailing whitespace, and special device names (e.g., "aux") in *`bstr`* are encoded as a literal part of the path element by using a `\\?\REL` prefix. The `bstr` argument must not contain a `\`, otherwise the `exn:fail:contract` exception is raised.

For (path-element->bytes path) or (path-element->string path), if the byte-string form of path starts with a \\?\REL, the prefix is not included in the result.

For (build-path base-path sub-path ...), trailing spaces and periods are removed from the last element of base-path and all but the last sub-path (unless the element consists of only spaces and periods), except for those that start with \\?\. If base-path starts \\?\, then after each non-\\?\REL\ and non-\\?\RED\ sub-path is added, all /s in the addition are converted to \s, multiple consecutive \s are converted to a single \, added . elements are removed, and added .. elements are removed along with the preceding element; these conversions are not performed on the original base-path part of the result or on any \\?\REL\ or \\?\RED\ or sub-path. If a \\?\REL\ or \\?\RED\ sub-path is added to a non-\\?\ base-path, the base-path (with any additions up to the \\?\REL\ or \\?\RED\ sub-path) is simplified and converted to a \\?\ path. In other cases, a \ may be added or removed before combining paths to avoid changing the root meaning of the path (e.g., combining //x and y produces /x/y, because //x/y would be a UNC path instead of a drive-relative path).

For (simplify-path path use-filesystem?), path is expanded, and if path does not start with \\?\, trailing spaces and periods are removed, a / is inserted after the colon in a drive specification if it is missing, and a \ is inserted after \\?\ as a root if there are elements and no extra \ already. Otherwise, if no indicators or redundant separators are in path, then path is returned.

For (split-path path) producing base, name, and must-be-dir?, splitting a path that does not start with \\?\ can produce parts that start with \\?\. For example, splitting C:/x~/aux/ produces \\?\C:\x~\ and \\?\REL\\aux; the \\?\ is needed in these cases to preserve a trailing space after x and to avoid referring to the AUX device instead of an "aux" file.

15.2 Filesystem

15.2.1 Locating Paths

(find-system-path kind) → path?
  kind : symbol?

Returns a machine-specific path for a standard type of path specified by kind, which must be one of the following:

- 'home-dir — the current *user's home directory*.

  On all platforms, if the PLTUSERHOME environment variable is defined as a complete path, then the path is used as the user's home directory.

  On Unix and Mac OS X, when PLTUSERHOME does not apply, the user's home directory is determined by expanding the path "~", which is expanded by first checking for

a HOME environment variable. If none is defined, the USER and LOGNAME environment variables are consulted (in that order) to find a user name, and then system files are consulted to locate the user's home directory.

On Windows, when PLTUSERHOME does not apply, the user's home directory is the user-specific profile directory as determined by the Windows registry. If the registry cannot provide a directory for some reason, the value of the USERPROFILE environment variable is used instead, as long as it refers to a directory that exists. If USERPROFILE also fails, the directory is the one specified by the HOMEDRIVE and HOMEPATH environment variables. If those environment variables are not defined, or if the indicated directory still does not exist, the directory containing the current executable is used as the home directory.

- 'pref-dir — the standard directory for storing the current user's preferences. On Unix, the directory is ".racket" in the user's home directory. On Windows, it is "Racket" in the user's home directory if determined by PLTUSERHOME, otherwise in the user's application-data folder as specified by the Windows registry; the application-data folder is usually "Application Data" in the user's profile directory. On Mac OS X, the preferences directory is "Library/Preferences" in the user's home directory. The preferences directory might not exist.

- 'pref-file — a file that contains a symbol-keyed association list of preference values. The file's directory path always matches the result returned for 'pref-dir. The file name is "racket-prefs.rktd" on Unix and Windows, and it is "org.racket-lang.prefs.rktd" on Mac OS X. The file's directory might not exist. See also get-preference.

- 'temp-dir — the standard directory for storing temporary files. On Unix and Mac OS X, this is the directory specified by the TMPDIR environment variable, if it is defined, otherwise it is the first path that exists among "/var/tmp", "/usr/tmp", and "/tmp". On Windows, the result is the directory specified by the TMP or TEMP environment variable, if it is defined, otherwise it is the current directory.

- 'init-dir — the directory containing the initialization file used by the Racket executable. It is the same as the user's home directory.

- 'init-file — the file loaded at start-up by the Racket executable. The directory part of the path is the same path as returned for 'init-dir. The file name is platform-specific:

  - Unix and Mac OS X: ".racketrc"
  - Windows: "racketrc.rktl"

- 'config-dir — a directory for the installation's configuration. This directory is specified by the PLTCONFIGDIR environment variable, and it can be overridden by the --config or -G command-line flag. If no environment variable or flag is specified, or if the value is not a legal path name, then this directory defaults to an "etc" directory

relative to the current executable. If the result of (find-system-path 'config-dir) is a relative path, it is relative to the current executable. The directory might not exist.

- 'addon-dir — a directory for user-specific Racket configuration, packages, and extension. This directory is specified by the PLTADDONDIR environment variable, and it can be overridden by the --addon or -A command-line flag. If no environment variable or flag is specified, or if the value is not a legal path name, then this directory defaults to "Library/Racket" in the user's home directory on Mac OS X and 'pref-dir otherwise. The directory might not exist.

- 'doc-dir — the standard directory for storing the current user's documents. On Unix, it's the user's home directory. On Windows, it is the user's home directory if determined by PLTUSERHOME, otherwise it is the user's documents folder as specified by the Windows registry; the documents folder is usually "My Documents" in the user's home directory. On Mac OS X, it's the "Documents" directory in the user's home directory.

- 'desk-dir — the directory for the current user's desktop. On Unix, it's the user's home directory. On Windows, it is the user's home directory if determined by PLTUSERHOME, otherwise it is the user's desktop folder as specified by the Windows registry; the desktop folder is usually "Desktop" in the user's home directory. On Mac OS X, it is "Desktop" in the user's home directory

- 'sys-dir — the directory containing the operating system for Windows. On Unix and Mac OS X, the result is "/".

- 'exec-file — the path of the Racket executable as provided by the operating system for the current invocation. For some operating systems, the path can be relative.

For GRacket, the executable path is the name of a GRacket executable.

- 'run-file — the path of the current executable; this may be different from result for 'exec-file because an alternate path was provided through a --name or -N command-line flag to the Racket (or GRacket) executable, or because an embedding executable installed an alternate path. In particular a "launcher" script created by make-racket-launcher sets this path to the script's path.

- 'collects-dir — a path to the main collection of libraries (see §18.2 "Libraries and Collections"). If this path is relative, then it is relative to the executable as reported by (find-system-path 'exec-file)—though the latter could be a soft-link or relative to the user's executable search path, so that the two results should be combined with find-executable-path. The 'collects-dir path is normally embedded in the Racket executable, but it can be overridden by the --collects or -X command-line flag.

- 'orig-dir — the current directory at start-up, which can be useful in converting a relative-path result from (find-system-path 'exec-file) or (find-system-path 'run-file) to a complete path.

Changed in version 6.0.0.3 of package base: Added PLTUSERHOME.

```
(path-list-string->path-list str
                             default-path-list)
 → (listof path?)
  str : (or/c string? bytes?)
  default-path-list : (listof path?)
```

Parses a string or byte string containing a list of paths, and returns a list of path strings. On Unix and Mac OS X, paths in a path list are separated by a `:`; on Windows, paths are separated by a `;`, and all `"`s in the string are discarded. Whenever the path list contains an empty path, the list *default-path-list* is spliced into the returned list of paths. Parts of *str* that do not form a valid path are not included in the returned list.

```
(find-executable-path  program-sub
                       [related-sub
                        deepest?])   → (or/c path? #f)
  program-sub : path-string?
  related-sub : (or/c path-string? #f) = #f
  deepest? : any/c = #f
```

Finds a path for the executable *program-sub*, returning `#f` if the path cannot be found.

If *related-sub* is not `#f`, then it must be a relative path string, and the path found for *program-sub* must be such that the file or directory *related-sub* exists in the same directory as the executable. The result is then the full path for the found *related-sub*, instead of the path for the executable.

This procedure is used by the Racket executable to find the standard library collection directory (see §18.2 “Libraries and Collections”). In this case, `program` is the name used to start Racket and `related` is `"collects"`. The *related-sub* argument is used because, on Unix and Mac OS X, *program-sub* may involve a sequence of soft links; in this case, *related-sub* determines which link in the chain is relevant.

If *related-sub* is not `#f`, then when `find-executable-path` does not find a *program-sub* that is a link to another file path, the search can continue with the destination of the link. Further links are inspected until *related-sub* is found or the end of the chain of links is reached. If *deepest?* is `#f` (the default), then the result corresponds to the first path in a chain of links for which *related-sub* is found (and further links are not actually explored); otherwise, the result corresponds to the last link in the chain for which *related-sub* is found.

If *program-sub* is a pathless name, `find-executable-path` gets the value of the `PATH` environment variable; if this environment variable is defined, `find-executable-path` tries each path in `PATH` as a prefix for *program-sub* using the search algorithm described above for path-containing *program-sub*s. If the `PATH` environment variable is not defined, *program-sub* is prefixed with the current directory and used in the search algorithm above.

(On Windows, the current directory is always implicitly the first item in PATH, so find-executable-path checks the current directory first on Windows.)

15.2.2 Files

```
(file-exists? path) → boolean?
  path : path-string?
```

Returns #t if a file (not a directory) path exists, #f otherwise.

On Windows, file-exists? reports #t for all variations of the special filenames (e.g., "LPT1", "x:/baddir/LPT1").

```
(link-exists? path) → boolean?
  path : path-string?
```

Returns #t if a link path exists, #f otherwise.

The predicates file-exists? or directory-exists? work on the final destination of a link or series of links, while link-exists? only follows links to resolve the base part of path (i.e., everything except the last name in the path).

This procedure never raises the exn:fail:filesystem exception.

On Windows, link-exists? reports #t for both symbolic links and junctions.

Changed in version 6.0.1.12 of package base: Added support for links on Windows.

```
(delete-file path) → void?
  path : path-string?
```

Deletes the file with path path if it exists, otherwise the exn:fail:filesystem exception is raised. If path is a link, the link is deleted rather than the destination of the link.

On Windows, delete-file can delete a symbolic link, but not a junction. Use delete-directory to delete a junction.

```
(rename-file-or-directory old
                         new
                        [exists-ok?]) → void?
  old : path-string?
  new : path-string?
  exists-ok? : any/c = #f
```

Renames the file or directory with path old—if it exists—to the path new. If the file or directory is not renamed successfully, the exn:fail:filesystem exception is raised.

This procedure can be used to move a file/directory to a different directory (on the same disk) as well as rename a file/directory within a directory. Unless *exists-ok?* is provided as a true value, *new* cannot refer to an existing file or directory. Even if *exists-ok?* is true, *new* cannot refer to an existing file when *old* is a directory, and vice versa.

If *new* exists and is replaced, the replacement is atomic on Unix and Mac OS X, but it is not guaranteed to be atomic on Windows. Furthermore, if *new* exists and is opened by any process for reading or writing, then attempting to replace it will typically fail on Windows. See also `call-with-atomic-output-file`.

If *old* is a link, the link is renamed rather than the destination of the link, and it counts as a file for replacing any existing *new*.

```
(file-or-directory-modify-seconds  path
                                   [secs-n
                                    fail-thunk]) → any
  path : path-string?
  secs-n : (or/c exact-integer? #f) = #f
  fail-thunk :  (-> any)
             = (lambda () (raise (make-exn:fail:filesystem ....)))
```

Returns the file or directory's last modification date in seconds since midnight UTC, January 1, 1970 (see also §15.6 "Time") when *secs-n* is not provided or is `#f`.

For FAT filesystems on Windows, directories do not have modification dates. Therefore, the creation date is returned for a directory, but the modification date is returned for a file.

If *secs-n* is provided and not `#f`, the access and modification times of *path* are set to the given time.

On error (e.g., if no such file exists), *fail-thunk* is called, and the default *fail-thunk* raises `exn:fail:filesystem`.

```
(file-or-directory-permissions path [mode])
 → (listof (or/c 'read 'write 'execute))
  path : path-string?
  mode : #f = #f
(file-or-directory-permissions path mode) → (integer-in 0 65535)
  path : path-string?
  mode : 'bits
(file-or-directory-permissions path mode) → void
  path : path-string?
  mode : (integer-in 0 65535)
```

When given one argument or `#f` as the second argument, returns a list containing `'read`, `'write`, and/or `'execute` to indicate permission the given file or directory path by the

current user and group. On Unix and Mac OS X, permissions are checked for the current effective user instead of the real user.

If `'bits` is supplied as the second argument, the result is a platform-specific integer encoding of the file or directory properties (mostly permissions), and the result is independent of the current user and group. The lowest nine bits of the encoding are somewhat portable, reflecting permissions for the file or directory's owner, members of the file or directory's group, or other users:

- `#o100` : owner has read permission
- `#o200` : owner has write permission
- `#o400` : owner has execute permission
- `#o010` : group has read permission
- `#o020` : group has write permission
- `#o040` : group has execute permission
- `#o001` : others have read permission
- `#o002` : others have write permission
- `#o004` : others have execute permission

See also `user-read-bit`, etc. On Windows, permissions from all three (owner, group, and others) are always the same, and read and execute permission are always available. On Unix and Mac OS X, higher bits have a platform-specific meaning.

If an integer is supplied as the second argument, its is used as an encoding of properties (mostly permissions) to install for the file.

In all modes, the `exn:fail:filesystem` exception is raised on error (e.g., if no such file exists).

```
(file-or-directory-identity path [as-link?])
 → exact-positive-integer?
  path : path-string?
  as-link? : any/c = #f
```

Returns a number that represents the identity of *path* in terms of the device and file or directory that it accesses. This function can be used to check whether two paths correspond to the same filesystem entity under the assumption that the path's entity selection does not change.

If *as-link?* is a true value, then if *path* refers to a filesystem link, the identity of the link is returned instead of the identity of the referenced file or directory (if any).

```
(file-size path) → exact-nonnegative-integer?
  path : path-string?
```

Returns the (logical) size of the specified file in bytes. On Mac OS X, this size excludes the resource-fork size. On error (e.g., if no such file exists), the `exn:fail:filesystem` exception is raised.

```
(copy-file src dest [exists-ok?]) → void?
  src : path-string?
  dest : path-string?
  exists-ok? : any/c = #f
```

Creates the file *dest* as a copy of *src*, if *dest* does not already exist. If *dest* already exists and *exists-ok?* is `#f`, the copy fails with `exn:fail:filesystem:exists?` exception is raised; otherwise, if *dest* exists, its content is replaced with the content of *src*. File permissions are transferred from *src* to *dest*; on Windows, the modification time of *src* is also transferred to *dest*. If *src* refers to a link, the target of the link is copied, rather than the link itself; if *dest* refers to a link and *exists-ok?* is true, the target of the link is updated.

```
(make-file-or-directory-link to path) → void?
  to : path-string?
  path : path-string?
```

Creates a link *path* to *to*. The creation will fail if *path* already exists. The *to* need not refer to an existing file or directory, and *to* is not expanded before writing the link. If the link is not created successfully,the `exn:fail:filesystem` exception is raised.

On Windows XP and earlier, the `exn:fail:unsupported` exception is raised. On later versions of Windows, the creation of links tends to be disallowed by security policies. Furthermore, a relative-path link is parsed specially; see §15.1.4 “Windows Path Conventions” for more information. When `make-file-or-directory-link` succeeds, it creates a symbolic link as opposed to a junction.

Changed in version 6.0.1.12 of package `base`: Added support for links on Windows.

15.2.3 Directories

See also: `rename-file-or-directory`, `file-or-directory-modify-seconds`, `file-or-directory-permissions`.

```
(current-directory) → (and/c path? complete-path?)
(current-directory path) → void?
  path : path-string?
```

A parameter that determines the current directory for resolving relative paths.

When the parameter procedure is called to set the current directory, the path argument is cleansed using `cleanse-path`, simplified using `simplify-path`, and then converted to a directory path with `path->directory-path`; cleansing and simplification raise an exception if the path is ill-formed. Thus, the current value of `current-directory` is always a cleansed, simplified, complete, directory path.

The path is not checked for existence when the parameter is set.

On Unix and Mac OS X, the initial value of the parameter for a Racket process is taken from the `PWD` environment variable—if the value of the environment variable identifies the same directory as the operating system's report of the current directory.

```
(current-directory-for-user) → (and/c path? complete-path?)
(current-directory-for-user path) → void?
  path : path-string?
```

Like `current-directory`, but use only by `srcloc->string` for reporting paths relative to a directory.

Normally, `current-directory-for-user` should stay at its initial value, reflecting the directory where a user started a process. A tool such as DrRacket, however, implicitly lets a user select a directory (for the file being edited), in which case updating `current-directory-for-user` makes sense.

```
(current-drive) → path?
```

Returns the current drive name Windows. For other platforms, the `exn:fail:unsupported` exception is raised. The current drive is always the drive of the current directory.

```
(directory-exists? path) → boolean?
  path : path-string?
```

Returns `#t` if *path* refers to a directory, `#f` otherwise.

```
(make-directory path) → void?
  path : path-string?
```

Creates a new directory with the path *path*. If the directory is not created successfully, the `exn:fail:filesystem` exception is raised.

```
(delete-directory path) → void?
  path : path-string?
```

Deletes an existing directory with the path *path*. If the directory is not deleted successfully, the `exn:fail:filesystem` exception is raised.

```
(directory-list [path #:build? build?]) → (listof path?)
  path : path-string? = (current-directory)
  build? : any/c = #f
```

See also the `in-directory` sequence constructor.

Returns a list of all files and directories in the directory specified by *path*. If *build?* is `#f`, the resulting paths are all path elements; otherwise, the individual results are combined with *path* using `build-path`. On Windows, an element of the result list may start with `\\?\REL\\`.

The resulting paths are always sorted using `path<?`.

```
(filesystem-root-list) → (listof path?)
```

Returns a list of all current root directories. Obtaining this list can be particularly slow on Windows.

15.2.4 Detecting Filesystem Changes

Many operating systems provide notifications for filesystem changes, and those notifications are reflected in Racket by filesystem change events.

```
(filesystem-change-evt? v) → boolean?
  v : any/c
```

Returns `#f` if *v* is a filesystem change event, `#f` otherwise.

```
(filesystem-change-evt path) → filesystem-change-evt?
  path : path-string?
(filesystem-change-evt path failure-thunk) → any
  path : path-string?
  failure-thunk : (-> any)
```

Creates a *filesystem change event*, which is a synchronizable event that becomes ready for synchronization after a change to *path*:

- If *path* refers to a file, the event becomes ready for synchronization when the file's content or attributes change, or when the file is deleted.
- If *path* refers to a directory, the event becomes ready for synchronization if a file or subdirectory is added, renamed, or removed within the directory.

The event also becomes ready for synchronization if it is passed to `filesystem-change-evt-cancel`.

Finally, depending on the precision of information available from the operating system, the event may become ready for synchronization under other circumstances. For example, on Windows, an event for a file becomes ready when any file changes within in the same directory as the file.

After a filesystem change event becomes ready for synchronization, it stays ready for synchronization. The event's synchronization result is the event itself.

If the current platform does not support filesystem-change notifications, then the `exn:fail:unsupported` exception is raised if *failure-thunk* is not provided, or *failure-thunk* is called in tail position if provided. Similarly, if there is any operating-system error when creating the event (such as a non-existent file), then the `exn:fail:filesystem` exception is raised or *failure-thunk* is called.

Creation of a filesystem change event alloates resources at the operating-system level. The resources are released at latest when the event is sychronized and ready for synchronization or when the event is canceled with `filesystem-change-evt-cancel`. See also `system-type` in `'fs-change` mode.

A filesystem change event is placed under the management of the current custodian when it is created. If the custodian is shut down, `filesystem-change-evt-cancel` is applied to the event.

```
(filesystem-change-evt-cancel evt) → void?
  evt : filesystem-change-evt?
```

Causes *evt* to become immediately ready for synchronization, whether it was ready or before not, and releases and resources (at the operating-system level) for tracking filesystem changes.

15.2.5 Declaring Paths Needed at Run Time

```
(require racket/runtime-path)
```
package: `base`

The bindings documented in this section are provided by the `racket/runtime-path` library, not `racket/base` or `racket`.

The `racket/runtime-path` library provides forms for accessing files and directories at run time using a path that are usually relative to an enclosing source file. Unlike using `collection-path`, `define-runtime-path` exposes each run-time path to tools like the executable and distribution creators, so that files and directories needed at run time are carried along in a distribution.

In addition to the bindings described below, `racket/runtime-path` provides `#%datum` in phase level 1, since string constants are often used as compile-time expressions with `define-runtime-path`.

```
(define-runtime-path id expr)
```

Uses *expr* as both a compile-time (i.e., phase 1) expression and a run-time (i.e., phase 0) expression. In either context, *expr* should produce a path, a string that represents a path, a list of the form `(list 'lib str ...+)`, or a list of the form `(list 'so str)` or `(list 'so str vers)`.

For run time, *id* is bound to a path that is based on the result of *expr*. The path is normally computed by taking a relative path result from *expr* and adding it to a path for the enclosing file (which is computed as described below). However, tools like the executable creator can also arrange (by colluding with `racket/runtime-path`) to have a different base path substituted in a generated executable. If *expr* produces an absolute path, it is normally returned directly, but again may be replaced by an executable creator. In all cases, the executable creator preserves the relative locations of all paths within a given package (treating paths outside of any package as being together). When *expr* produces a relative or absolute path, then the path bound to *id* is always an absolute path.

If *expr* produces a list of the form `(list 'lib str ...+)`, the value bound to *id* is an absolute path. The path refers to a collection-based file similar to using the value as a module path.

If *expr* produces a list of the form `(list 'so str)` or `(list 'so str vers)`, the value bound to *id* can be either *str* or an absolute path; it is an absolute path when searching in the Racket-specific shared-object library directories (as determined by `get-lib-search-dirs`) locates the path. In this way, shared-object libraries that are installed specifically for Racket get carried along in distributions. The search tries each directory in order; within a directory, the search tries using *str* directly, then it tries adding each version specified by *vers*—which defaults to `'(#f)`—along with a platform-specific shared-library extension—as produced by `(system-type 'so-suffix)`. A *vers* can be a string, or it can be a list of strings and `#f`.

If *expr* produces a list of the form `(list 'module module-path var-ref)` or `(list 'so str (list str-or-false ...))`, the value bound to *id* is a module path index, where *module-path* is treated as relative (if it is relative) to the module that is the home of the variable reference *var-ref*, where *var-ref* can be `#f` if *module-path* is absolute. In an executable, the corresponding module is carried along, including all of its dependencies.

For compile-time, the *expr* result is used by an executable creator—but not the result when the containing module is compiled. Instead, *expr* is preserved in the module as a compile-time expression (in the sense of `begin-for-syntax`). Later, at the time that an executable is created, the compile-time portion of the module is executed (again), and the result of *expr* is the file to be included with the executable. The reason for the extra compile-time execution is that the result of *expr* might be platform-dependent, so the result should not be stored in the (platform-independent) bytecode form of the module; the platform at executable-creation time, however, is the same as at run time for the executable. Note that *expr* is still evaluated

at run-time; consequently, avoid procedures like `collection-path`, which depends on the source installation, and instead use relative paths and forms like `(list 'lib str ...+)`.

If a path is needed only on some platforms and not on others, use `define-runtime-path-list` with an *expr* that produces an empty list on platforms where the path is not needed.

Beware that `define-runtime-path` in a phase level other than 0 does not cooperate properly with an executable creator. To work around that limitation, put `define-runtime-path` in a separate module—perhaps a submodule created by `module`—then export the definition, and then the module containing the definition can be `required` into any phase level. Using `define-runtime-path` in a phase level other than 0 logs a warning at expansion time.

The enclosing path for a `define-runtime-path` is determined as follows from the `define-runtime-path` syntactic form:

- If the form has a source module according to `syntax-source-module`, then the source location is determined by preserving the original expression as a syntax object, extracting its source module path at run time (again using `syntax-source-module`), and then resolving the resulting module path index.
- If the expression has no source module, the `syntax-source` location associated with the form is used, if is a string or path.
- If no source module is available, and `syntax-source` produces no path, then `current-load-relative-directory` is used if it is not `#f`. Finally, `current-directory` is used if all else fails.

In the latter two cases, the path is normally preserved in (platform-specific) byte form, but if the enclosing path corresponds to a result of `collection-file-path`, then the path is record as relative to the corresponding module path.

Changed in version 6.0.1.6 of package `base`: Preserve relative paths only within a package.

Examples:

```
; Access a file "data.txt" at run-time that is originally
; located in the same directory as the module source file:
(define-runtime-path data-file "data.txt")
(define (read-data)
  (with-input-from-file data-file
    (lambda ()
      (read-bytes (file-size data-file)))))

; Load a platform-specific shared object (using ffi-lib)
; that is located in a platform-specific sub-directory of the
; module's source directory:
(define-runtime-path libfit-path
```

```
  (build-path "compiled" "native" (system-library-subpath #f)
              (path-replace-suffix "libfit"
                                   (system-type 'so-suffix))))
(define libfit (ffi-lib libfit-path))

; Load a platform-specific shared object that might be installed
; as part of the operating system, or might be installed
; specifically for Racket:
(define-runtime-path libssl-so
  (case (system-type)
    [(windows) '(so "ssleay32")]
    [else '(so "libssl")]))
(define libssl (ffi-lib libssl-so))
```

(define-runtime-paths (*id* ...) *expr*)

Like define-runtime-path, but declares and binds multiple paths at once. The *expr* should produce as many values as *id*s.

(define-runtime-path-list *id* *expr*)

Like define-runtime-path, but *expr* should produce a list of paths.

(define-runtime-module-path-index *id* *module-path-expr*)

Similar to define-runtime-path, but *id* is bound to a module path index that encapsulates the result of *module-path-expr* relative to the enclosing module.

Use define-runtime-module-path-index to bind a module path that is passed to a reflective function like dynamic-require while also creating a module dependency for building and distributing executables.

(runtime-require *module-path*)

Similar to define-runtime-module-path-index, but creates the distribution dependency without binding a module path index. When runtime-require is used multiple times within a module with the same *module-path*, all but the first use expands to an empty begin.

(define-runtime-module-path *id* *module-path*)

Similar to define-runtime-path, but *id* is bound to a resolved module path. The resolved module path for *id* corresponds to *module-path* (with the same syntax as a module path for require), which can be relative to the enclosing module.

The `define-runtime-module-path-index` form is usually preferred, because it creates a weaker link to the referenced module. Unlike `define-runtime-module-path-index`, the `define-runtime-module-path` form creates a `for-label` dependency from an enclosing module to *module-path*. Since the dependency is merely `for-label`, *module-path* is not instantiated or visited when the enclosing module is instantiated or visited (unless such a dependency is created by other `require`s), but the code for the referenced module is loaded when the enclosing module is loaded.

```
(runtime-paths module-path)
```

This form is mainly for use by tools such as executable builders. It expands to a quoted list containing the run-time paths declared by *module-path*, returning the compile-time results of the declaration `expr`s, except that paths are converted to byte strings. The enclosing module must require (directly or indirectly) the module specified by *module-path*, which is an unquoted module path. The resulting list does *not* include module paths bound through `define-runtime-module-path`.

15.2.6 More File and Directory Utilities

`(require racket/file)` package: `base`

The bindings documented in this section are provided by the `racket/file` and `racket` libraries, but not `racket/base`.

```
(file->string path [#:mode mode-flag]) → string?
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
```

Reads all characters from *path* and returns them as a string. The *mode-flag* argument is the same as for `open-input-file`.

```
(file->bytes path [#:mode mode-flag]) → bytes?
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
```

Reads all characters from *path* and returns them as a byte string. The *mode-flag* argument is the same as for `open-input-file`.

```
(file->value path [#:mode mode-flag]) → any
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
```

Reads a single S-expression from *path* using `read`. The *mode-flag* argument is the same as for `open-input-file`.

```
(file->list path [proc #:mode mode-flag]) → (listof any/c)
  path : path-string?
  proc : (input-port? . -> . any/c) = read
  mode-flag : (or/c 'binary 'text) = 'binary
```

Repeatedly calls *proc* to consume the contents of *path*, until `eof` is produced. The *mode-flag* argument is the same as for `open-input-file`.

```
(file->lines  path
              [#:mode mode-flag
               #:line-mode line-mode]) → (listof string?)
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
  line-mode :  (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
            = 'any
```

Read all characters from *path*, breaking them into lines. The *line-mode* argument is the same as the second argument to `read-line`, but the default is `'any` instead of `'linefeed`. The *mode-flag* argument is the same as for `open-input-file`.

```
(file->bytes-lines  path
                    [#:mode mode-flag
                     #:line-mode line-mode]) → (listof bytes?)
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
  line-mode :  (or/c 'linefeed 'return 'return-linefeed 'any 'any-one)
            = 'any
```

Like `file->lines`, but reading bytes and collecting them into lines like `read-bytes-line`.

```
(display-to-file  v
                  path
                 [#:mode mode-flag
                  #:exists exists-flag]) → void?
  v : any/c
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
  exists-flag :  (or/c 'error 'append 'update
                       'replace 'truncate 'truncate/replace)
              = 'error
```

Uses `display` to print *v* to *path*. The *mode-flag* and *exists-flag* arguments are the same as for `open-output-file`.

```
(write-to-file  v
                path
               [#:mode mode-flag
                #:exists exists-flag]) → void?
  v : any/c
  path : path-string?
  mode-flag : (or/c 'binary 'text) = 'binary
  exists-flag :  (or/c 'error 'append 'update
                       'replace 'truncate 'truncate/replace)
              = 'error
```

Like `display-to-file`, but using `write` instead of `display`.

```
(display-lines-to-file  lst
                        path
                       [#:separator separator
                        #:mode mode-flag
                        #:exists exists-flag]) → void?
  lst : list?
  path : path-string?
  separator : any/c = #"\n"
  mode-flag : (or/c 'binary 'text) = 'binary
  exists-flag :  (or/c 'error 'append 'update
                       'replace 'truncate 'truncate/replace)
              = 'error
```

Displays each element of *lst* to *path*, adding *separator* after each element. The *mode-flag* and *exists-flag* arguments are the same as for `open-output-file`.

```
 (copy-directory/files
  src
  dest
  #:keep-modify-seconds? keep-modify-seconds?)
 → void?
  src : path-string?
  dest : path-string?
  keep-modify-seconds? : #f
```

Copies the file or directory *src* to *dest*, raising `exn:fail:filesystem` if the file or directory cannot be copied, possibly because *dest* exists already. If *src* is a directory, the copy applies recursively to the directory's content. If a source is a link, the target of the link is copied rather than the link itself.

If *keep-modify-seconds?* is `#f`false, then file copies keep only the properties kept by `copy-file`, If *keep-modify-seconds?* is true, then each file copy also keeps the modification date of the original.

```
(delete-directory/files path
                        #:must-exist? must-exist?) → void?
  path : path-string?
  must-exist? : #t
```

Deletes the file or directory specified by *path*, raising `exn:fail:filesystem` if the file or directory cannot be deleted. If *path* is a directory, then `delete-directory/files` is first applied to each file and directory in *path* before the directory is deleted.

If *must-exist?* is true, then `exn:fail:filesystem` is raised if *path* does not exist. If *must-exist?* is false, then `delete-directory/files` succeeds if *path* does not exist (but a failure is possible if *path* initially exists and is removed by another thread or process before `delete-directory/files` deletes it).

```
(find-files  predicate
             [start-path]
              #:follow-links? follow-links?) → (listof path?)
  predicate : (path? . -> . any/c)
  start-path : (or/c path-string? #f) = #f
  follow-links? : #f
```

Traverses the filesystem starting at *start-path* and creates a list of all files and directories for which *predicate* returns true. If *start-path* is `#f`, then the traversal starts from `(current-directory)`. In the resulting list, each directory precedes its content.

The *predicate* procedure is called with a single argument for each file or directory. If *start-path* is `#f`, the argument is a pathname string that is relative to the current directory. Otherwise, it is a path building on *start-path*. Consequently, supplying `(current-directory)` for *start-path* is different from supplying `#f`, because *predicate* receives complete paths in the former case and relative paths in the latter. Another difference is that *predicate* is not called for the current directory when *start-path* is `#f`.

If *follow-links?* is true, the `find-files` traversal follows links, and links are not included in the result. If *follow-links?* is `#f`, then links are not followed, and links are included in the result.

If *start-path* does not refer to an existing file or directory, then *predicate* will be called exactly once with *start-path* as the argument.

The `find-files` procedure raises and exception if it encounters a directory for which `directory-list` fails.

```
(pathlist-closure path-list
                  #:follow-links? follow-links?)
 → (listof path?)
  path-list : (listof path-string?)
```

```
  follow-links? : #f
```

Given a list of paths, either absolute or relative to the current directory, returns a list such that

- if a nested path is given, all of its ancestors are also included in the result (but the same ancestor is not added twice);
- if a path refers to directory, all of its descendants are also included in the result;
- ancestor directories appear before their descendants in the result list.

If *follow-links?* is true, then the traversal of directories and files follows links, and the link paths are not included in the result. If *follow-links?* is `#f`, then he result list includes paths to link and the links are not followed.

```
(fold-files  proc
             init-val
            [start-path
             follow-links?]) → any
        (or/c (path? (or/c 'file 'dir 'link) any/c
                . -> . any/c)
  proc :      (path? (or/c 'file 'dir 'link) any/c
                . -> . (values any/c any/c)))
  init-val : any/c
  start-path : (or/c path-string? #f) = #f
  follow-links? : any/c = #t
```

Traverses the filesystem starting at *start-path*, calling *proc* on each discovered file, directory, and link. If *start-path* is `#f`, then the traversal starts from `(current-directory)`.

The *proc* procedure is called with three arguments for each file, directory, or link:

- If *start-path* is `#f`, the first argument is a pathname string that is relative to the current directory. Otherwise, the first argument is a pathname that starts with *start-path*. Consequently, supplying `(current-directory)` for *start-path* is different from supplying `#f`, because *proc* receives complete paths in the former case and relative paths in the latter. Another difference is that *proc* is not called for the current directory when *start-path* is `#f`.
- The second argument is a symbol, either `'file`, `'dir`, or `'link`. The second argument can be `'link` when *follow-links?* is `#f`, in which case the filesystem traversal does not follow links. If *follow-links?* is `#t`, then *proc* will only get a `'link` as a second argument when it encounters a dangling symbolic link (one that does not resolve to an existing file or directory).

- The third argument is the accumulated result. For the first call to *proc*, the third argument is *init-val*. For the second call to *proc* (if any), the third argument is the result from the first call, and so on. The result of the last call to *proc* is the result of fold-files.

The *proc* argument is used in an analogous way to the procedure argument of foldl, where its result is used as the new accumulated result. There is an exception for the case of a directory (when the second argument is 'dir): in this case the procedure may return two values, the second indicating whether the recursive scan should include the given directory or not. If it returns a single value, the directory is scanned. In the cases of files or links (when the second argument is 'file or 'link), a second value is permitted but ignored.

If the *start-path* is provided but no such path exists, or if paths disappear during the scan, then an exception is raised.

```
(make-directory* path) → void?
  path : path-string?
```

Creates directory specified by *path*, creating intermediate directories as necessary, and never failing if *path* exists already.

```
(make-temporary-file [template
                      copy-from-filename
                      directory])         → path?
  template : string? = "rkttmp~a"
  copy-from-filename : (or/c path-string? #f 'directory) = #f
  directory : (or/c path-string? #f) = #f
```

Creates a new temporary file and returns a pathname string for the file. Instead of merely generating a fresh file name, the file is actually created; this prevents other threads or processes from picking the same temporary name.

The *template* argument must be a format string suitable for use with format and one additional string argument (where the string contains only digits). If the resulting string is a relative path, it is combined with the result of (find-system-path 'temp-dir), unless *directory* is provided and non-#f, in which case the file name generated from *template* is combined with *directory* to obtain a full path.

The *template* argument's default is only the string "rkttmp~a" when there is no source location information for the callsite of make-temporary-file (or if make-temporary-file is used in a higher-order position). If there is such information, then the template string is based on the source location.

If *copy-from-filename* is provided as path, the temporary file is created as a copy of the named file (using copy-file). If *copy-from-filename* is #f, the temporary file is created

as empty. If *copy-from-filename* is 'directory, then the temporary "file" is created as a directory.

When a temporary file is created, it is not opened for reading or writing when the pathname is returned. The client program calling make-temporary-file is expected to open the file with the desired access and flags (probably using the 'truncate flag; see open-output-file) and to delete it when it is no longer needed.

```
(call-with-atomic-output-file  file
                               proc
                              [#:security-guard security-guard])
 → any
  file : path-string?
  proc : ([port input-port?] [tmp-path path?]  . -> . any)
  security-guard : (or/c #f security-guard?) = #f
```

Opens a temporary file for writing in the same directory as *file*, calls *proc* to write to the temporary file, and then atomically moves the temporary file in place of *proc*. The atomic move simply uses rename-file-or-directory on Unix and Mac OS X, but it uses an extra rename step (see below) on Windows to avoid problems due to concurrent readers of *file*.

The *proc* function is called with an output port for the temporary file, plus the path of the temporary file. The result of *proc* is the result of call-with-atomic-output.

The call-with-atomic-output function arranges to delete temporary files on exceptions.

Windows prevents programs from deleting or replacing files that are open, but it allows renaming of open files. Therefore, on Windows, call-with-atomic-output-file creates a second temporary file *extra-tmp-file*, renames *file* to *extra-tmp-file*, renames the temporary file written by *proc* to p, and finally deletes *extra-tmp-file*.

```
(get-preference  name
                [failure-thunk
                 flush-mode
                 filename
                 #:use-lock? use-lock?
                 #:timeout-lock-there timeout-lock-there
                 #:lock-there lock-there])                → any
  name : symbol?
  failure-thunk : (-> any) = (lambda () #f)
  flush-mode : any/c = 'timestamp
  filename : (or/c string-path? #f) = #f
  use-lock? : any/c = #t
  timeout-lock-there : (or/c (path? . -> . any) #f) = #f
```

```
lock-there : (or/c (path? . -> . any) #f)
             (make-handle-get-preference-locked
           =  0.01 name failure-thunk flush-mode filename
              #:lock-there timeout-lock-there)
```

Extracts a preference value from the file designated by `(find-system-path 'pref-file)`, or by *filename* if it is provided and is not `#f`. In the former case, if the preference file doesn't exist, `get-preferences` attempts to read an old preferences file, and then a `"racket-prefs.rktd"` file in the configuration directory (as reported by `find-config-dir`), instead. If none of those files exists, the preference set is empty.

The preference file should contain a list of symbol–value lists written with the default parameter settings. Keys starting with `racket:`, `mzscheme:`, `mred:`, and `plt:` in any letter case are reserved for use by Racket implementors. If the preference file does not contain a list of symbol–value lists, an error is logged via `log-error` and *failure-thunk* is called.

The result of `get-preference` is the value associated with *name* if it exists in the association list, or the result of calling *failure-thunk* otherwise.

Preference settings are cached (weakly) across calls to `get-preference`, using `(path->complete-path` *filename*`)` as a cache key. If *flush-mode* is provided as `#f`, the cache is used instead of re-consulting the preferences file. If *flush-mode* is provided as `'timestamp` (the default), then the cache is used only if the file has a timestamp that is the same as the last time the file was read. Otherwise, the file is re-consulted.

On platforms for which `preferences-lock-file-mode` returns `'file-lock` and when *use-lock?* is true, preference-file reading is guarded by a lock; multiple readers can share the lock, but writers take the lock exclusively. If the preferences file cannot be read because the lock is unavailable, *lock-there* is called on the path of the lock file; if *lock-there* is `#f`, an exception is raised. The default *lock-there* handler retries about 5 times (with increasing delays between each attempt) before trying *timeout-lock-there*, and the default *timeout-lock-there* triggers an exception.

See also `put-preferences`. For a more elaborate preference system, see `preferences:get`.

Old preferences files: When a *filename* is not provided and the file indicated by `(find-system-path 'pref-file)` does not exist, the following paths are checked for compatibility with old versions of Racket:

- Windows: `(build-path (find-system-path 'pref-dir) 'up "PLT Scheme" "plt-prefs.ss")`
- Mac OS X: `(build-path (find-system-path 'pref-dir) "org.plt-scheme.prefs.ss")`
- Unix: `(expand-user-path "~/.plt-scheme/plt-prefs.ss")`

```
(put-preferences  names
                  vals
                 [locked-proc
                  filename])   → void?
  names : (listof symbol?)
  vals : list?
  locked-proc : (or/c #f (path? . -> . any)) = #f
  filename : (or/c #f path-string?) = #f
```

Installs a set of preference values and writes all current values to the preference file designated by `(find-system-path 'pref-file)`, or *filename* if it is supplied and not `#f`.

The *names* argument supplies the preference names, and *vals* must have the same length as *names*. Each element of *vals* must be an instance of a built-in data type whose `write` output is `read`able (i.e., the `print-unreadable` parameter is set to `#f` while writing preferences).

Current preference values are read from the preference file before updating, and a write lock is held starting before the file read, and lasting until after the preferences file is updated. The lock is implemented by the existence of a file in the same directory as the preference file; see `preferences-lock-file-mode` for more information. If the directory of the preferences file does not already exist, it is created.

If the write lock is already held, then *locked-proc* is called with a single argument: the path of the lock file. The default *locked-proc* (used when the *locked-proc* argument is `#f`) reports an error; an alternative thunk might wait a while and try again, or give the user the choice to delete the lock file (in case a previous update attempt encountered disaster and locks are implemented by the presence of the lock file).

If *filename* is `#f` or not supplied, and the preference file does not already exist, then values read from the `"defaults"` collection (if any) are written for preferences that are not mentioned in *names*.

```
(preferences-lock-file-mode) → (or/c 'exists 'file-lock)
```

Reports the way that the lock file is used to implement preference-file locking on the current platform.

The `'exists` mode is currently used on all platforms except Windows. In `'exists` mode, the existence of the lock file indicates that a write lock is held, and readers need no lock (because the preferences file is atomically updated via `rename-file-or-directory`).

The `'file-lock` mode is currently used on Windows. In `'file-lock` mode, shared and exclusive locks (in the sense of `port-try-file-lock?`) on the lock file reflect reader and writer locks on the preference-file content. (The preference file itself is not locked, because a lock would interfere with replacing the file via `rename-file-or-directory`.)

```
(make-handle-get-preference-locked  delay
                                    name
                                   [failure-thunk
                                    flush-mode
                                    filename
                                    #:lock-there lock-there
                                    #:max-delay max-delay])
 → (path-string? . -> . any)
  delay : real?
  name : symbol?
  failure-thunk : (-> any) = (lambda () #f)
  flush-mode : any/c = 'timestamp
  filename : (or/c path-string? #f) = #f
  lock-there : (or/c (path? . -> . any) #f) = #f
  max-delay : real? = 0.2
```

Creates a procedure suitable for use as the `#:lock-there` argument to `get-preference`, where the *name*, *failure-thunk*, *flush-mode*, and *filename* are all passed on to `get-preference` by the result procedure to retry the preferences lookup.

Before calling `get-preference`, the result procedure uses `(sleep delay)` to pause. Then, if `(* 2 delay)` is less than *max-delay*, the result procedure calls `make-handle-get-preference-locked` to generate a new retry procedure to pass to `get-preference`, but with a *delay* of `(* 2 delay)`. If `(* 2 delay)` is not less than *max-delay*, then `get-preference` is called with the given *lock-there*, instead.

```
(call-with-file-lock/timeout  filename
                              kind
                              thunk
                              failure-thunk
                             [#:lock-file lock-file
                              #:delay delay
                              #:max-delay max-delay]) → any
  filename : (or/c path-string? #f)
  kind : (or/c 'shared 'exclusive)
  thunk : (-> any)
  failure-thunk : (-> any)
  lock-file : (or/c #f path-string?) = #f
  delay : (and/c real? (not/c negative?)) = 0.01
  max-delay : (and/c real? (not/c negative?)) = 0.2
```

Obtains a lock for the filename *lock-file* and then calls *thunk*. The *filename* argument specifies a file path prefix that is used only to generate the lock filename when *lock-file* is `#f`. Specifically, when *lock-file* is `#f`, then `call-with-file-lock/timeout` uses

`make-lock-file-name` to build the lock filename. If the lock file does not yet exist, it is created; beware that the lock file is *not* deleted by `call-with-file-lock/timeout`.

When *`thunk`* returns, `call-with-file-lock/timeout` releases the lock, returning the result of *`thunk`*. The `call-with-file-lock/timeout` function will retry after *`delay`* seconds and continue retrying with exponential backoff until delay reaches *`max-delay`*. If `call-with-file-lock/timeout` fails to obtain the lock, *`failure-thunk`* is called in tail position. The *`kind`* argument specifies whether the lock is `'shared` or `'exclusive` in the sense of `port-try-file-lock?`.

Examples:

```
> (call-with-file-lock/timeout filename 'exclusive
    (lambda () (printf "File is locked\n"))
    (lambda () (printf "Failed to obtain lock for file\n")))
File is locked

> (call-with-file-lock/timeout #f 'exclusive
    (lambda ()
      (call-with-file-lock/timeout filename 'shared
        (lambda () (printf "Shouldn't get here\n"))
        (lambda () (printf "Failed to obtain lock for file\n"))))
    (lambda () (printf "Shouldn't get here either\n"))
    #:lock-file (make-lock-file-name filename))
Failed to obtain lock for file
```

```
(make-lock-file-name path) → path?
  path : (or path-string? path-for-some-system?)
(make-lock-file-name dir name) → path?
  dir : (or path-string? path-for-some-system?)
  name : path-element?
```

Creates a lock filename by prepending `"_LOCK"` on Windows or `".LOCK"` on other platforms to the file portion of the path.

Example:

```
> (make-lock-file-name "/home/george/project/important-file")
#<path:/home/george/project/.LOCKimportant-file>
```

```
user-read-bit : #o400
user-write-bit : #o200
user-execute-bit : #o100
group-read-bit : #o040
group-write-bit : #o020
group-execute-bit : #o010
```

```
other-read-bit : #o004
other-write-bit : #o002
other-execute-bit : #o001
```

Constants that are useful with `file-or-directory-permissions` and bitwise operations such as `bitwise-ior`, and `bitwise-and`.

15.3 Networking

15.3.1 TCP

`(require racket/tcp)` package: `base`

The bindings documented in this section are provided by the `racket/tcp` and `racket` libraries, but not `racket/base`.

For information about TCP in general, see *TCP/IP Illustrated, Volume 1* by W. Richard Stevens.

```
(tcp-listen  port-no
             [max-allow-wait
              reuse?
              hostname])      → tcp-listener?
  port-no : (integer-in 0 65535)
  max-allow-wait : exact-nonnegative-integer? = 4
  reuse? : any/c = #f
  hostname : (or/c string? #f) = #f
```

Creates a "listening" server on the local machine at the port number specified by *port-no*. If *port-no* is 0 the socket binds to an ephemeral port, which can be determined by calling `tcp-addresses`. The *max-allow-wait* argument determines the maximum number of client connections that can be waiting for acceptance. (When *max-allow-wait* clients are waiting acceptance, no new client connections can be made.)

If the *reuse?* argument is true, then `tcp-listen` will create a listener even if the port is involved in a `TIME_WAIT` state. Such a use of *reuse?* defeats certain guarantees of the TCP protocol; see Stevens's book for details. Furthermore, on many modern platforms, a true value for *reuse?* overrides `TIME_WAIT` only if the listener was previously created with a true value for *reuse?*.

If *hostname* is `#f` (the default), then the listener accepts connections to all of the listening machine's addresses. Otherwise, the listener accepts connections only at the interface(s) associated with the given hostname. For example, providing `"127.0.0.1"` as *hostname* creates a listener that accepts only connections to `"127.0.0.1"` (the loopback interface) from the local machine.

(Racket implements a listener with multiple sockets, if necessary, to accommodate multiple addresses with different protocol families. On Linux, if *hostname* maps to both IPv4 and IPv6 addresses, then the behavior depends on whether IPv6 is supported and IPv6 sockets can be configured to listen to only IPv6 connections: if IPv6 is not supported or IPv6 sockets are not configurable, then the IPv6 addresses are ignored; otherwise, each IPv6 listener accepts only IPv6 connections.)

The return value of `tcp-listen` is a *TCP listener*. This value can be used in future calls to `tcp-accept`, `tcp-accept-ready?`, and `tcp-close`. Each new TCP listener value is placed into the management of the current custodian (see §14.7 "Custodians").

If the server cannot be started by `tcp-listen`, the `exn:fail:network` exception is raised.

A TCP listener can be used as a synchronizable event (see §11.2.1 "Events"). A TCP listener is ready for synchronization when `tcp-accept` would not block; the synchronization result of a TCP listener is the TCP listener itself.

```
(tcp-connect  hostname
              port-no
             [local-hostname
              local-port-no]) → input-port? output-port?
  hostname : string?
  port-no : (integer-in 1 65535)
  local-hostname : (or/c string? #f) = #f
  local-port-no : (or/c (integer-in 1 65535) #f) = #f
```

Attempts to connect as a client to a listening server. The *hostname* argument is the server host's Internet address name, and *port-no* is the port number where the server is listening.

(If *hostname* is associated with multiple addresses, they are tried one at a time until a connection succeeds. The name `"localhost"` generally specifies the local machine.)

The optional *local-hostname* and *local-port-no* specify the client's address and port. If both are `#f` (the default), the client's address and port are selected automatically. If *local-hostname* is not `#f`, then *local-port-no* must be non-`#f`. If *local-port-no* is non-`#f` and *local-hostname* is `#f`, then the given port is used but the address is selected automatically.

Two values are returned by `tcp-connect`: an input port and an output port. Data can be received from the server through the input port and sent to the server through the output port. If the server is a Racket program, it can obtain ports to communicate to the client with `tcp-accept`. These ports are placed into the management of the current custodian (see §14.7 "Custodians").

Initially, the returned input port is block-buffered, and the returned output port is block-buffered. Change the buffer mode using `file-stream-buffer-mode`.

Both of the returned ports must be closed to terminate the TCP connection. When both ports are still open, closing the output port with close-output-port sends a TCP close to the server (which is seen as an end-of-file if the server reads the connection through a port). In contrast, tcp-abandon-port (see below) closes the output port, but does not send a TCP close until the input port is also closed.

Note that the TCP protocol does not support a state where one end is willing to send but not read, nor does it include an automatic message when one end of a connection is fully closed. Instead, the other end of a connection discovers that one end is fully closed only as a response to sending data; in particular, some number of writes on the still-open end may appear to succeed, though writes will eventually produce an error.

If a connection cannot be established by tcp-connect, the exn:fail:network exception is raised.

```
(tcp-connect/enable-break  hostname
                           port-no
                          [local-hostname]
                           local-port-no)
 → input-port? output-port?
  hostname : string?
  port-no : (integer-in 1 65535)
  local-hostname : (or/c string? #f) = #f
  local-port-no : (or/c (integer-in 1 65535) #f)
```

Like tcp-connect, but breaking is enabled (see §10.6 “Breaks”) while trying to connect. If breaking is disabled when tcp-connect/enable-break is called, then either ports are returned or the exn:break exception is raised, but not both.

```
(tcp-accept listener) → input-port? output-port?
  listener : tcp-listener?
```

Accepts a client connection for the server associated with *listener*. If no client connection is waiting on the listening port, the call to tcp-accept will block. (See also tcp-accept-ready?.)

Two values are returned by tcp-accept: an input port and an output port. Data can be received from the client through the input port and sent to the client through the output port. These ports are placed into the management of the current custodian (see §14.7 “Custodians”).

In terms of buffering and connection states, the ports act the same as ports from tcp-connect.

If a connection cannot be accepted by tcp-accept, or if the listener has been closed, the exn:fail:network exception is raised.

```
(tcp-accept/enable-break listener) → input-port? output-port?
  listener : tcp-listener?
```

Like tcp-accept, but breaking is enabled (see §10.6 “Breaks”) while trying to accept a connection. If breaking is disabled when tcp-accept/enable-break is called, then either ports are returned or the exn:break exception is raised, but not both.

```
(tcp-accept-ready? listener) → boolean?
  listener : tcp-listener?
```

Tests whether an unaccepted client has connected to the server associated with *listener*. If a client is waiting, the return value is #t, otherwise it is #f. A client is accepted with the tcp-accept procedure, which returns ports for communicating with the client and removes the client from the list of unaccepted clients.

If the listener has been closed, the exn:fail:network exception is raised.

```
(tcp-close listener) → void?
  listener : tcp-listener?
```

Shuts down the server associated with *listener*. All unaccepted clients receive an end-of-file from the server; connections to accepted clients are unaffected.

If the listener has already been closed, the exn:fail:network exception is raised.

The listener’s port number may not become immediately available for new listeners (with the default *reuse?* argument of tcp-listen). For further information, see Stevens’s explanation of the TIME_WAIT TCP state.

```
(tcp-listener? v) → boolean?
  v : any/c
```

Returns #t if *v* is a TCP listener created by tcp-listen, #f otherwise.

```
(tcp-accept-evt listener) → evt?
  listener : tcp-listener?
```

Returns a synchronizable event (see §11.2.1 “Events”) that is ready for synchronization when tcp-accept on *listener* would not block. The synchronization result is a list of two items, which correspond to the two results of tcp-accept. (If the event is not chosen in a syntax, no connections are accepted.) The ports are placed into the management of the custodian that is the current custodian (see §14.7 “Custodians”) at the time that tcp-accept-evt is called.

```
(tcp-abandon-port tcp-port) → void?
  tcp-port : tcp-port?
```

Like close-output-port or close-input-port (depending on whether *tcp-port* is an input or output port), but if *tcp-port* is an output port and its associated input port is not yet closed, then the other end of the TCP connection does not receive a TCP close message until the input port is also closed.

The TCP protocol does not include a "no longer reading" state on connections, so tcp-abandon-port is equivalent to close-input-port on input TCP ports.

```
(tcp-addresses tcp-port [port-numbers?])
    (or/c (values string? string?)
 →        (values string? (integer-in 1 65535)
                  string? (integer-in 0 65535)))
  tcp-port : (or/c tcp-port? tcp-listener?)
  port-numbers? : any/c = #f
```

Returns two strings when *port-numbers?* is #f (the default). The first string is the Internet address for the local machine a viewed by the given TCP port's connection or for the TCP listener. (For most machines, the answer corresponds to the current machine's only Internet address, but when a machine serves multiple addresses, the result is connection-specific or listener-specific.) If a listener is given and it has no specific host, the first string result is "0.0.0.0". The second string is the Internet address for the other end of the connection, or always "0.0.0.0" for a listener.

If *port-numbers?* is true, then four results are returned: a string for the local machine's address, an exact integer between 1 and 65535 for the local machine's port number, a string for the remote machine's address, and an exact integer between 1 and 65535 for the remote machine's port number or 0 for a listener.

If the given port has been closed, the exn:fail:network exception is raised.

```
(tcp-port? v) → boolean?
  v : any/c
```

Returns #t if *v* is a *TCP port*—which is a port returned by tcp-accept, tcp-connect, tcp-accept/enable-break, or tcp-connect/enable-break—#f otherwise.

15.3.2 UDP

(require racket/udp) package: base

The bindings documented in this section are provided by the racket/udp and racket libraries, but not racket/base.

For information about UDP in general, see *TCP/IP Illustrated, Volume 1* by W. Richard Stevens.

```
(udp-open-socket [family-hostname
                  family-port-no]) → udp?
  family-hostname : (or/c string? #f) = #f
  family-port-no : (or/c (integer-in 1 65535) #f) = #f
```

Creates and returns a *UDP socket* to send and receive datagrams (broadcasting is allowed). Initially, the socket is not bound or connected to any address or port.

If *family-hostname* or *family-port-no* is not `#f`, then the socket's protocol family is determined from these arguments. The socket is *not* bound to the hostname or port number. For example, the arguments might be the hostname and port to which messages will be sent through the socket, which ensures that the socket's protocol family is consistent with the destination. Alternately, the arguments might be the same as for a future call to `udp-bind!`, which ensures that the socket's protocol family is consistent with the binding. If neither *family-hostname* nor *family-port-no* is non-`#f`, then the socket's protocol family is IPv4.

```
(udp-bind!  udp-socket
            hostname-string
            port-no
           [reuse?])          → void?
  udp-socket : udp?
  hostname-string : (or/c string? #f)
  port-no : (integer-in 0 65535)
  reuse? : any/c = #f
```

Binds an unbound *udp-socket* to the local port number *port-no*. If *port-no* is 0 the *udp-socket* is bound to an ephemeral port, which can be determined by calling `udp-addresses`.

If *hostname-string* is `#f`, then the socket accepts connections to all of the listening machine's IP addresses at *port-no*. Otherwise, the socket accepts connections only at the IP address associated with the given name. For example, providing `"127.0.0.1"` as *hostname-string* typically creates a listener that accepts only connections to `"127.0.0.1"` from the local machine.

A socket cannot receive datagrams until it is bound to a local address and port. If a socket is not bound before it is used with a sending procedure `udp-send`, `udp-send-to`, etc., the sending procedure binds the socket to a random local port. Similarly, if an event from `udp-send-evt` or `udp-send-to-evt` is chosen for a synchronization (see §11.2.1 "Events"), the socket is bound; if the event is not chosen, the socket may or may not become bound.

The binding of a bound socket cannot be changed, with one exception: on some systems, if the socket is bound automatically when sending, if the socket is disconnected via `udp-`

connect!, and if the socket is later used again in a send, then the later send may change the socket's automatic binding.

If udp-socket is already bound or closed, the exn:fail:network exception is raised.

If the reuse? argument is true, then udp-bind! will set the SO_REUSEADDR socket option before binding, permitting the sharing of access to a UDP port between many processes on a single machine when using UDP multicast.

```
(udp-connect! udp-socket
              hostname-string
              port-no)        → void?
  udp-socket : udp?
  hostname-string : (or/c string? #f)
  port-no : (or/c (integer-in 1 65535)
                  #f)
```

Connects the socket to the indicated remote address and port if hostname-string is a string and port-no is an exact integer.

If hostname-string is #f, then port-no also must be #f, and the port is disconnected (if connected). If one of hostname-string or port-no is #f and the other is not, the exn:fail:contract exception is raised.

A connected socket can be used with udp-send (not udp-send-to), and it accepts datagrams only from the connected address and port. A socket need not be connected to receive datagrams. A socket can be connected, re-connected, and disconnected any number of times.

If udp-socket is closed, the exn:fail:network exception is raised.

```
(udp-send-to  udp-socket
              hostname
              port-no
              bstr
             [start-pos
              end-pos])   → void
  udp-socket : udp?
  hostname : string?
  port-no : (integer-in 1 65535)
  bstr : bytes?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Sends (subbytes bytes start-pos end-pos) as a datagram from the unconnected udp-socket to the socket at the remote machine hostname-address on the port port-no. The udp-socket need not be bound or connected; if it is not bound, udp-send-to

binds it to a random local port. If the socket's outgoing datagram queue is too full to support the send, udp-send-to blocks until the datagram can be queued.

If *start-pos* is greater than the length of *bstr*, or if *end-pos* is less than *start-pos* or greater than the length of *bstr*, the exn:fail:contract exception is raised.

If *udp-socket* is closed or connected, the exn:fail:network exception is raised.

```
(udp-send udp-socket bstr [start-pos end-pos]) → void
  udp-socket : udp?
  bstr : bytes?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like udp-send-to, except that *udp-socket* must be connected, and the datagram goes to the connection target. If *udp-socket* is closed or unconnected, the exn:fail:network exception is raised.

```
(udp-send-to*  udp-socket
               hostname
               port-no
               bstr
              [start-pos
               end-pos])   → boolean?
  udp-socket : udp?
  hostname : string?
  port-no : (integer-in 1 65535)
  bstr : bytes?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like udp-send-to, but never blocks; if the socket's outgoing queue is too full to support the send, #f is returned, otherwise the datagram is queued and the result is #t.

```
(udp-send* udp-socket bstr [start-pos end-pos]) → boolean?
  udp-socket : udp?
  bstr : bytes?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like udp-send, except that (like udp-send-to) it never blocks and returns #f or #t.

```
(udp-send-to/enable-break  udp-socket
                           hostname
                           port-no
                           bstr
                          [start-pos
                           end-pos])   → void
```

```
  udp-socket : udp?
  hostname : string?
  port-no : (integer-in 1 65535)
  bstr : bytes?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like udp-send-to, but breaking is enabled (see §10.6 “Breaks”) while trying to send the datagram. If breaking is disabled when udp-send-to/enable-break is called, then either the datagram is sent or the exn:break exception is raised, but not both.

```
(udp-send/enable-break  udp-socket
                        bstr
                       [start-pos
                        end-pos])   → void
  udp-socket : udp?
  bstr : bytes?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like udp-send, except that breaks are enabled like udp-send-to/enable-break.

```
(udp-receive!  udp-socket
               bstr
              [start-pos
                           exact-nonnegative-integer?
               end-pos])   → string?
                           (integer-in 1 65535)
  udp-socket : udp?
  bstr : (and/c bytes? (not immutable?))
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Accepts up to *end-pos-start-pos* bytes of *udp-socket*’s next incoming datagram into *bstr*, writing the datagram bytes starting at position *start-pos* within *bstr*. The *udp-socket* must be bound to a local address and port (but need not be connected). If no incoming datagram is immediately available, udp-receive! blocks until one is available.

Three values are returned: the number of received bytes (between 0 and *end-pos-start-pos*, a hostname string indicating the source address of the datagram, and an integer indicating the source port of the datagram. If the received datagram is longer than *end-pos-start-pos* bytes, the remainder is discarded.

If *start-pos* is greater than the length of *bstr*, or if *end-pos* is less than *start-pos* or greater than the length of *bstr*, the exn:fail:contract exception is raised.

```
(udp-receive!* udp-socket
               bstr
              [start-pos
               end-pos])
   (or/c exact-nonnegative-integer? #f)
 → (or/c string? #f)
   (or/c (integer-in 1 65535) #f)
  udp-socket : udp?
  bstr : (and/c bytes? (not immutable?))
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like `udp-receive!`, except that it never blocks. If no datagram is available, the three result values are all `#f`.

```
(udp-receive!/enable-break udp-socket
                           bstr
                          [start-pos
                           end-pos])
   exact-nonnegative-integer?
 → string?
   (integer-in 1 65535)
  udp-socket : udp?
  bstr : (and/c bytes? (not immutable?))
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Like `udp-receive!`, but breaking is enabled (see §10.6 “Breaks”) while trying to receive the datagram. If breaking is disabled when `udp-receive!/enable-break` is called, then either a datagram is received or the `exn:break` exception is raised, but not both.

```
(udp-close udp-socket) → void?
  udp-socket : udp?
```

Closes *udp-socket*, discarding unreceived datagrams. If the socket is already closed, the `exn:fail:network` exception is raised.

```
(udp? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a socket created by `udp-open-socket`, `#f` otherwise.

```
(udp-bound? udp-socket) → boolean?
  udp-socket : udp?
```

Returns `#t` if *udp-socket* is bound to a local address and port, `#f` otherwise.

```
(udp-connected? udp-socket) → boolean?
  udp-socket : udp?
```

Returns `#t` if *udp-socket* is connected to a remote address and port, `#f` otherwise.

```
(udp-send-ready-evt udp-socket) → evt?
  udp-socket : udp?
```

Returns a synchronizable event (see §11.2.1 “Events”) that is in a blocking state when `udp-send-to` on *udp-socket* would block. The synchronization result is the event itself.

```
(udp-receive-ready-evt udp-socket) → evt?
  udp-socket : udp?
```

Returns a synchronizable event (see §11.2.1 “Events”) that is in a blocking state when `udp-receive!` on *udp-socket* would block. The synchronization result is the event itself.

```
(udp-send-to-evt  udp-socket
                  hostname
                  port-no
                  bstr
                 [start-pos
                  end-pos])   → evt?
  udp-socket : udp?
  hostname : string?
  port-no : (integer-in 1 65535)
  bstr : bytes?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Returns a synchronizable event. The event is in a blocking state when `udp-send-to` on *udp-socket* would block. Otherwise, if the event is chosen in a synchronization, data is sent as for `(udp-send-to` *udp-socket* `hostname-address` *port-no bstr start-pos end-pos*`)`, and the synchronization result is `#<void>`. (No bytes are sent if the event is not chosen.)

```
(udp-send-evt  udp-socket
               bstr
              [start-pos
               end-pos])   → evt?
  udp-socket : udp?
  bstr : bytes?
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Returns a synchronizable event. The event is ready for synchronization when `udp-send` on *udp-socket* would not block. Otherwise, if the event is chosen in a synchronization, data is sent as for `(udp-send-to` *udp-socket bstr start-pos end-pos*`)`, and the synchronization result is `#<void>`. (No bytes are sent if the event is not chosen.) If *udp-socket* is closed or unconnected, the `exn:fail:network` exception is raised during a synchronization attempt.

```
(udp-receive!-evt  udp-socket
                   bstr
                  [start-pos
                   end-pos])   → evt?
  udp-socket : udp?
  bstr : (and/c bytes? (not immutable?))
  start-pos : exact-nonnegative-integer? = 0
  end-pos : exact-nonnegative-integer? = (bytes-length bstr)
```

Returns a synchronizable event. The event is ready for synchronization when `udp-receive` on *udp-socket* would not block. Otherwise, if the event is chosen in a synchronization, data is received into *bstr* as for `(udp-receive!` *udp-socket bytes start-pos end-pos*`)`, and the synchronization result is a list of three values, corresponding to the three results from `udp-receive!`. (No bytes are received and the *bstr* content is not modified if the event is not chosen.)

```
(udp-addresses udp-port [port-numbers?])
    (or/c (values string? string?)
 →        (values string? (integer-in 0 65535)
                  string? (integer-in 0 65535)))
  udp-port : udp?
  port-numbers? : any/c = #f
```

Returns two strings when *port-numbers?* is `#f` (the default). The first string is the Internet address for the local machine a viewed by the given UDP socket's connection. (For most machines, the answer corresponds to the current machine's only Internet address, but when a machine serves multiple addresses, the result is connection-specific.) The second string is the Internet address for the other end of the connection.

If *port-numbers?* is true, then four results are returned: a string for the local machine's address, an exact integer between `1` and `65535` for the local machine's port number or `0` if the socket is unbound, a string for the remote machine's address, and an exact integer between `1` and `65535` for the remote machine's port number or `0` if the socket is unconnected.

If the given port has been closed, the `exn:fail:network` exception is raised.

```
(udp-multicast-join-group! udp-socket
                           multicast-addr
                           hostname)      → void?
  udp-socket : udp?
  multicast-addr : string?
  hostname : (or/c string? #f)
(udp-multicast-leave-group! udp-socket
                            multicast-addr
                            hostname)      → void?
  udp-socket : udp?
  multicast-addr : string?
  hostname : (or/c string? #f)
```

Adds or removes *udp-socket* to a named multicast group.

The *multicast-addr* argument must be a valid IPv4 multicast IP address; for example, `"224.0.0.251"` is the appropriate address for the mDNS protocol. The *hostname* argument selects the interface that the socket uses to receive (not send) multicast datagrams; if *hostname* is `#f` or `"0.0.0.0"`, the kernel selects an interface automatically.

Leaving a group requires the same *multicast-addr* and *hostname* arguments that were used to join the group.

```
(udp-multicast-interface udp-socket) → string?
  udp-socket : udp?
(udp-multicast-set-interface! udp-socket
                              hostname)  → void?
  udp-socket : udp?
  hostname : (or/c string? #f)
```

Retrieves or sets the interface that *udp-socket* uses to send (not receive) multicast datagrams. If the result or *hostname* is either `#f` or `"0.0.0.0"`, the kernel automatically selects an interface when a multicast datagram is sent.

```
(udp-multicast-set-loopback! udp-socket
                             loopback?) → void?
  udp-socket : udp?
  loopback? : any/c
(udp-multicast-loopback? udp-socket) → boolean?
  udp-socket : udp?
```

Loopback settings correspond to the `IP_MULTICAST_LOOP` setting of the socket.

Sets or checks whether *udp-socket* receives its own multicast datagrams: a `#t` result or a true value for *loopback?* indicates that self-receipt is enabled, and `#f` indicates that self-receipt is disabled.

```
(udp-multicast-set-ttl! udp-socket ttl) → void?
  udp-socket : udp?
  ttl : byte?
(udp-multicast-ttl udp-socket) → byte?
  udp-socket : udp?
```

Time-to-live settings correspond to the `IP_MULTICAST_TTL` setting of the socket.

Sets or retrieves the current time-to-live setting of *udp-socket*.

The time-to-live setting should almost always be 1, and it is important that this number is as low as possible. In fact, these functions seldom should be used at all. See the documentation for your platform's IP stack.

15.4 Processes

```
(subprocess stdout
            stdin
            stderr
            command
            arg ...)
  → subprocess?
    (or/c (and/c input-port? file-stream-port?) #f)
    (or/c (and/c output-port? file-stream-port?) #f)
    (or/c (and/c input-port? file-stream-port?) #f)
  stdout : (or/c (and/c output-port? file-stream-port?) #f)
  stdin : (or/c (and/c input-port? file-stream-port?) #f)
  stderr : (or/c (and/c output-port? file-stream-port?) #f 'stdout)
  command : path-string?
  arg : (or/c path? string-no-nuls? bytes-no-nuls?)
(subprocess stdout
            stdin
            stderr
            command
            exact
            arg)
  → subprocess?
    (or/c (and/c input-port? file-stream-port?) #f)
    (or/c (and/c output-port? file-stream-port?) #f)
    (or/c (and/c input-port? file-stream-port?) #f)
  stdout : (or/c (and/c output-port? file-stream-port?) #f)
  stdin : (or/c (and/c input-port? file-stream-port?) #f)
  stderr : (or/c (and/c output-port? file-stream-port?) #f)
  command : path-string?
  exact : 'exact
  arg : string?
```

Creates a new process in the underlying operating system to execute *command* asynchronously, providing the new process with environment variables current-environment-variables. See also system and process from racket/system.

The *command* argument is a path to a program executable, and the *arg*s are command-line arguments for the program. See find-executable-path for locating an executable based on the PATH environment variable. On Unix and Mac OS X, command-line arguments are passed as byte strings, and string *arg*s are converted using the current locale's encoding (see §13.1.1 "Encodings and Locales"). On Windows, command-line arguments are passed as strings, and bytes strings are converted using UTF-8.

On Unix and Mac OS X, subprocess creation is separate from starting the program indicated by *command*. In particular, if *command* refers to a non-existent or non-executable file, an error will be reported (via standard error and a non-0 exit code) in the subprocess, not in the creating process.

On Windows, the first *arg* can be replaced with 'exact, which triggers a Windows-specific behavior: the sole *arg* is used exactly as the command-line for the subprocess. Otherwise, on Windows, a command-line string is constructed from *command* and *arg* so that a typical Windows console application can parse it back to an array of arguments. If 'exact is provided on a non-Windows platform, the exn:fail:contract exception is raised.

For information on the Windows command-line conventions, search for "command line parsing" at http://msdn.microsoft.com/.

When provided as a port, *stdout* is used for the launched process's standard output, *stdin* is used for the process's standard input, and *stderr* is used for the process's standard error. All provided ports must be file-stream ports. Any of the ports can be #f, in which case a system pipe is created and returned by subprocess. The *stderr* argument can be 'stdout, in which case the same file-stream port or system pipe that is supplied as standard output is also used for standard error. For each port or 'stdout that is provided, no pipe is created and the corresponding returned value is #f.

The subprocess procedure returns four values:

- a *subprocess* value representing the created process;
- an input port piped from the process's standard output, or #f if stdout-output-port was a port;
- an output port piped to the process standard input, or #f if stdin-input-port was a port;
- an input port piped from the process's standard error, or #f if stderr-output-port was a port or 'stdout.

Important: All ports returned from subprocess must be explicitly closed, usually with close-input-port or close-output-port.

A file-stream port for communicating with a subprocess is normally a pipe with a limited capacity. Beware of creating deadlock by serializing a write to a subprocess followed by a read, while the subprocess does the same, so that both processes end up blocking on a write because the other end must first read to make room in the pipe. Beware also of waiting for a subprocess to finish without reading its

The returned ports are file-stream ports (see §13.1.5 "File Ports"), and they are placed into the management of the current custodian (see §14.7 "Custodians"). The exn:fail exception is raised when a low-level error prevents the spawning of a process or the creation of operating system pipes for process communication.

If the subprocess-group-enabled parameter's value is true, then the new process is created as a new OS-level process group. In that case, subprocess-kill attempts to termi-

nate all processes within the group, which may include additional processes created by the subprocess. See subprocess-kill for details, and see subprocess-group-enabled for additional caveats.

The current-subprocess-custodian-mode parameter determines whether the subprocess itself is registered with the current custodian so that a custodian shutdown calls subprocess-kill for the subprocess.

A subprocess can be used as a synchronizable event (see §11.2.1 "Events"). A subprocess value is ready for synchronization when subprocess-wait would not block; the synchronization result of a subprocess value is the subprocess value itself.

```
(subprocess-wait subproc) → void?
  subproc : subprocess?
```

Blocks until the process represented by *subproc* terminates. The *subproc* value also can be used with sync and sync/timeout.

```
(subprocess-status subproc) → (or/c 'running
                                    exact-nonnegative-integer?)
  subproc : subprocess?
```

Returns 'running if the process represented by *subproc* is still running, or its exit code otherwise. The exit code is an exact integer, and 0 typically indicates success. If the process terminated due to a fault or signal, the exit code is non-zero.

```
(subprocess-kill subproc force?) → void?
  subproc : subprocess?
  force? : any/c
```

Terminates the subprocess represented by *subproc*. The precise action depends on whether *force?* is true, whether the process was created in its own group by setting the subprocess-group-enabled parameter to a true value, and the current platform:

- *force?* is true, not a group, all platforms: Terminates the process if the process still running.
- *force?* is false, not a group, on Unix or Mac OS X: Sends the process an interrupt signal instead of a kill signal.
- *force?* is false, not a group, on Windows: No action is taken.
- *force?* is true, a group, on Unix or Mac OS X: Terminates all processes in the group, but only if subprocess-status has never produced a non-'running result for the subprocess and only if functions like subprocess-wait and sync have not detected the subprocess's completion. Otherwise, no action is taken (because the immediate process is known to have terminated while the continued existence of the group is unknown).

- *force?* is true, a group, on Windows: Terminates the process if the process still running.
- *force?* is false, a group, on Unix or Mac OS X: The same as when *force?* is #t, but when the group is sent a signal, it is an interrupt signal instead of a kill signal.
- *force?* is false, a group, on Windows: All processes in the group receive a CTRL-BREAK signal (independent of whether the immediate subprocess has terminated).

If an error occurs during termination, the exn:fail exception is raised.

```
(subprocess-pid subproc) → exact-nonnegative-integer?
  subproc : subprocess?
```

Returns the operating system's numerical ID (if any) for the process represented by *subproc*. The result is valid only as long as the process is running.

```
(subprocess? v) → boolean?
  v : any/c
```

Returns #t if *v* is a subprocess value, #f otherwise.

```
(current-subprocess-custodian-mode)
 → (or/c #f 'kill 'interrupt)
(current-subprocess-custodian-mode mode) → void?
  mode : (or/c #f 'kill 'interrupt)
```

A parameter that determines whether a subprocess (as created by subprocess or wrappers like process) is registered with the current custodian. If the parameter value is #f, then the subprocess is not registered with the custodian—although any created ports are registered. If the parameter value is 'kill or 'interrupt, then the subprocess is shut down through subprocess-kill, where 'kill supplies a #t value for the *force?* argument and 'interrupt supplies a #f value. The shutdown may occur either before or after ports created for the subprocess are closed.

Custodian-triggered shutdown is limited by details of process handling in the host system. For example, process and system may create an intermediate shell process to run a program, in which case custodian-based termination shuts down the shell process and probably not the process started by the shell. See also subprocess-kill. Process groups (see subprocess-group-enabled) can address some limitations, but not all of them.

```
(subprocess-group-enabled) → boolean?
(subprocess-group-enabled on?) → void?
  on? : any/c
```

A parameter that determines whether a subprocess is created as a new process group. See `subprocess-kill` for more information.

Beware that creating a group may interfere with the job control in an interactive shell, since job control is based on process groups.

```
(shell-execute verb
               target
               parameters
               dir
               show-mode) → #f
  verb : (or/c string? #f)
  target : string?
  parameters : string?
  dir : path-string?
  show-mode : symbol?
```

Performs the action specified by *verb* on *target* in Windows. For platforms other than Windows, the `exn:fail:unsupported` exception is raised.

For example,

```
(shell-execute #f "http://racket-lang.org" ""
               (current-directory) 'sw_shownormal)
```

Opens the Racket home page in a browser window.

The *verb* can be `#f`, in which case the operating system will use a default verb. Common verbs include `"open"`, `"edit"`, `"find"`, `"explore"`, and `"print"`.

The *target* is the target for the action, usually a filename path. The file could be executable, or it could be a file with a recognized extension that can be handled by an installed application.

The *parameters* argument is passed on to the system to perform the action. For example, in the case of opening an executable, the *parameters* is used as the command line (after the executable name).

The *dir* is used as the current directory when performing the action.

The *show-mode* sets the display mode for a Window affected by the action. It must be one of the following symbols; the description of each symbol's meaning is taken from the Windows API documentation.

- `'sw_hide` or `'SW_HIDE` — Hides the window and activates another window.
- `'sw_maximize` or `'SW_MAXIMIZE` — Maximizes the window.

- `'sw_minimize` or `'SW_MINIMIZE` — Minimizes the window and activates the next top-level window in the z-order.
- `'sw_restore` or `'SW_RESTORE` — Activates and displays the window. If the window is minimized or maximized, Windows restores it to its original size and position.
- `'sw_show` or `'SW_SHOW` — Activates the window and displays it in its current size and position.
- `'sw_showdefault` or `'SW_SHOWDEFAULT` — Uses a default.
- `'sw_showmaximized` or `'SW_SHOWMAXIMIZED` — Activates the window and displays it as a maximized window.
- `'sw_showminimized` or `'SW_SHOWMINIMIZED` — Activates the window and displays it as a minimized window.
- `'sw_showminnoactive` or `'SW_SHOWMINNOACTIVE` — Displays the window as a minimized window. The active window remains active.
- `'sw_showna` or `'SW_SHOWNA` — Displays the window in its current state. The active window remains active.
- `'sw_shownoactivate` or `'SW_SHOWNOACTIVATE` — Displays a window in its most recent size and position. The active window remains active.
- `'sw_shownormal` or `'SW_SHOWNORMAL` — Activates and displays a window. If the window is minimized or maximized, Windows restores it to its original size and position.

If the action fails, the `exn:fail` exception is raised. If the action succeeds, the result is `#f`.

In future versions of Racket, the result may be a subprocess value if the operating system did returns a process handle (but if a subprocess value is returned, its process ID will be `0` instead of the real process ID).

15.4.1 Simple Subprocesses

`(require racket/system)` package: `base`

The bindings documented in this section are provided by the `racket/system` and `racket` libraries, but not `racket/base`.

```
(system command [#:set-pwd? set-pwd?]) → boolean?
  command : (or/c string-no-nuls? bytes-no-nuls?)
  set-pwd? : any/c = (member (system-type) '(unix macosx))
```

Executes a Unix, Mac OS X, or Windows shell command synchronously (i.e., the call to `system` does not return until the subprocess has ended). The *command* argument is a string or byte string containing no nul characters. If the command succeeds, the return value is `#t`, `#f` otherwise.

See also `subprocess` for notes about error handling and the limited buffer capacity of subprocess pipes.

If *set-pwd?* is true, then the `PWD` environment variable is set to the value of `(current-directory)` when starting the shell process.

See also `current-subprocess-custodian-mode` and `subprocess-group-enabled`, which affect the subprocess used to implement `system`.

The resulting process writes to `(current-output-port)`, reads from `(current-input-port)`, and logs errors to `(current-error-port)`. To gather the process's non-error output to a string, for example, use `with-output-to-string`, which sets `current-output-port` while calling the given function:

```
(with-output-to-string (lambda () (system "date")))
```

```
(system* command arg ... [#:set-pwd? set-pwd?]) → boolean?
  command : path-string?
  arg : (or/c path? string-no-nuls? bytes-no-nuls?)
  set-pwd? : any/c = (member (system-type) '(unix macosx))
(system*  command
          exact
          arg
         [#:set-pwd? set-pwd?]) → boolean?
  command : path-string?
  exact : 'exact
  arg : string?
  set-pwd? : any/c = (member (system-type) '(unix macosx))
```

Like `system`, except that *command* is a filename that is executed directly (instead of through a shell command; see `find-executable-path` for locating an executable based on the `PATH` environment variable), and the *arg*s are the arguments. The executed file is passed the specified string arguments (which must contain no nul characters).

On Windows, the first argument after *command* can be `'exact`, and the final *arg* is a complete command line. See `subprocess` for details.

```
(system/exit-code  command
                   [#:set-pwd? set-pwd?]) → byte?
  command : (or/c string-no-nuls? bytes-no-nuls?)
  set-pwd? : any/c = (member (system-type) '(unix macosx))
```

Like `system`, except that the result is the exit code returned by the subprocess. A `0` result normally indicates success.

```
(system*/exit-code  command
                    arg ...
                   [#:set-pwd? set-pwd?]) → byte?
  command : path-string?
  arg : (or/c path? string-no-nuls? bytes-no-nuls?)
  set-pwd? : any/c = (member (system-type) '(unix macosx))
(system*/exit-code  command
                    exact
                    arg
                   [#:set-pwd? set-pwd?]) → byte?
  command : path-string?
  exact : 'exact
  arg : string?
  set-pwd? : any/c = (member (system-type) '(unix macosx))
```

Like `system*`, but returns the exit code like `system/exit-code`.

```
(process command [#:set-pwd? set-pwd?])
    (list input-port?
          output-port?
 →        exact-nonnegative-integer?
          input-port?
          ((or/c 'status 'wait 'interrupt 'kill) . -> . any))
  command : (or/c string-no-nuls? bytes-no-nuls?)
  set-pwd? : any/c = (member (system-type) '(unix macosx))
```

Executes a shell command asynchronously (using `sh` on Unix and Mac OS X, `cmd` on Windows). The result is a list of five values:

See also `subprocess` for notes about error handling and the limited buffer capacity of subprocess pipes.

- an input port piped from the subprocess's standard output,
- an output port piped to the subprocess standard input,
- the system process id of the subprocess,
- an input port piped from the subprocess's standard error, and
- a procedure of one argument, either `'status`, `'wait`, `'interrupt`, `'exit-code` or `'kill`:
  - `'status` returns the status of the subprocess as one of `'running`, `'done-ok`, or `'done-error`.
  - `'exit-code` returns the integer exit code of the subprocess or `#f` if it is still running.
  - `'wait` blocks execution in the current thread until the subprocess has completed.

– `'interrupt` sends the subprocess an interrupt signal on Unix and Mac OS X, and takes no action on Windows. The result is `#<void>`.

– `'kill` terminates the subprocess and returns `#<void>`. Note that the immediate process created by `process` is a shell process that may run another program; terminating the shell process may not terminate processes that the shell starts, particularly on Windows.

On Unix and Mac OS X, if *command* runs a single program, then `sh` typically runs the program in such a way that it replaces `sh` in the same process. For reliable and precise control over process creation, however, use `process*`.

Important: All three ports returned from `process` must be explicitly closed with `close-input-port` or `close-output-port`.

If *set-pwd?* is true, then `PWD` is set in the same way as `system`.

See also `current-subprocess-custodian-mode` and `subprocess-group-enabled`, which affect the subprocess used to implement `process`. In particular, the `'interrupt` and `'kill` process-control messages are implemented via `subprocess-kill`, so they can affect a process group instead of a single process.

```
(process*  command
           arg ...
          [#:set-pwd? set-pwd?]) → list?
  command : path-string?
  arg : (or/c path? string-no-nuls? bytes-no-nuls?)
  set-pwd? : any/c = (member (system-type) '(unix macosx))
(process*  command
           exact
           arg
          [#:set-pwd? set-pwd?]) → list?
  command : path-string?
  exact : 'exact
  arg : string?
  set-pwd? : any/c = (member (system-type) '(unix macosx))
```

Like `process`, except that *command* is a filename that is executed directly like `system*`, and the *arg*s are the arguments. On Windows, as for `system*`, the first *arg* can be replaced with `'exact`.

```
(process/ports  out
                in
                error-out
                command
               [#:set-pwd? set-pwd?]) → list?
  out : (or/c #f output-port?)
  in : (or/c #f input-port?)
  error-out : (or/c #f output-port? 'stdout)
  command : (or/c path? string-no-nuls? bytes-no-nuls?)
  set-pwd? : any/c = (member (system-type) '(unix macosx))
```

Like process, except that *out* is used for the process's standard output, *in* is used for the process's standard input, and *error-out* is used for the process's standard error. Any of the ports can be #f, in which case a system pipe is created and returned, as in process. If *error-out* is 'stdout, then standard error is redirected to standard output. For each port or 'stdout that is provided, no pipe is created, and the corresponding value in the returned list is #f.

```
(process*/ports  out
                 in
                 error-out
                 command
                 arg ...
                [#:set-pwd? set-pwd?]) → list?
  out : (or/c #f output-port?)
  in : (or/c #f input-port?)
  error-out : (or/c #f output-port? 'stdout)
  command : path-string?
  arg : (or/c path? string-no-nuls? bytes-no-nuls?)
  set-pwd? : any/c = (member (system-type) '(unix macosx))
(process*/ports  out
                 in
                 error-out
                 command
                 exact
                 arg
                [#:set-pwd? set-pwd?]) → list?
  out : (or/c #f output-port?)
  in : (or/c #f input-port?)
  error-out : (or/c #f output-port? 'stdout)
  command : path-string?
  exact : 'exact
  arg : string?
  set-pwd? : any/c = (member (system-type) '(unix macosx))
```

Like process*, but with the port handling of process/ports.

The contracts of system and related functions may signal a contract error with references to the following functions.

```
(string-no-nuls? x) → boolean?
  x : any/c
```

Ensures that *x* is a string and does not contain "\u0000".

```
(bytes-no-nuls? x) → boolean?
  x : any/c
```

Ensures that `x` is a byte-string and does not contain `#"\0"`.

15.5 Logging

A *logger* accepts events that contain information to be logged for interested parties. A *log receiver* represents an interested party that receives logged events asynchronously. Each event has a level of importance, and a log receiver subscribes to logging events at a certain level of importance and higher. The levels, in decreasing order of importance, are `'fatal`, `'error`, `'warning`, `'info`, and `'debug`.

To help organize logged events, loggers can be named and hierarchical. Every event reported to a logger is also propagated to its parent (if any), but the event message is prefixed with a name (if any) that is typically the name of the logger to which is was originally reported. A logger is not required to have a parent or name.

On start-up, Racket creates an initial logger that is used to record events from the core run-time system. For example, an `'debug` event is reported for each garbage collection (see §1.1.7 "Garbage Collection"). For this initial logger, two log receivers are also created: one that writes events to the process's original error output port, and one that writes events to the system log. The level of written events in each case is system-specific, and the default can be changed through command-line flags (see §18.1.4 "Command Line") or through environment variables:

- If the `PLTSTDERR` environment variable is defined and is not overridden by a command-line flag, it determines the level of the log receiver that propagates events to the original error port.

  The environment variable's value can be a ⟨*level*⟩: `none`, `fatal`, `error`, `warning`, `info`, or `debug`; all events the corresponding level of higher are printed. After an initial ⟨*level*⟩, the value can contain space-separated specifications of the form ⟨*level*⟩`@`⟨*name*⟩, which prints events whose names match ⟨*name*⟩ only at the given ⟨*level*⟩ or higher (where a ⟨*name*⟩ contains any character other than a space or `@`). For example, the value `"error debug@GC"` prints all events at the `'error` level and higher, but prints events named `'GC` at the `'debug` level and higher (which includes all levels).

  The default is `"error"`.

- If the `PLTSYSLOG` environment variable is defined and is not overridden by a command-line flag, it determines the level of the log receiver that propagates events to the system log. The possible values are the same as for `PLTSTDERR`.

  The default is `"none"` for Unix or `"error"` for Windows and Mac OS X.

The `current-logger` parameter determines the *current logger* that is used by forms such as `log-warning`. On start-up, the initial value of this parameter is the initial logger. The

run-time system sometimes uses the current logger to report events. For example, the byte-code compiler sometimes reports `'warning` events when it detects an expression that would produce a run-time error if evaluated.

15.5.1 Creating Loggers

```
(logger? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a logger, `#f` otherwise.

```
(make-logger [name parent notify-callback]) → logger?
  name : (or/c symbol? #f) = #f
  parent : (or/c logger? #f) = #f
  notify-callback : (vector? . -> . any/c) = #f
```

Creates a new logger with an optional name and parent.

If *notify-callback* is provided, then it is called (under a continuation barrier) whenever an event is logged to the result logger or one of its descendants, but only if some log receiver is interested in the event in the same sense as `log-level?`. The event is not propagated to any log receivers until *notify-callback* returns.

```
(logger-name logger) → (or/c symbol? #f)
  logger : logger?
```

Reports *logger*'s name, if any.

```
(current-logger) → logger?
(current-logger logger) → void?
  logger : logger?
```

A parameter that determines the current logger.

```
(define-logger id)
```

Defines `log-`*id*`-fatal`, `log-`*id*`-error`, `log-`*id*`-warning`, `log-`*id*`-info`, and `log-`*id*`-debug` as forms like `log-fatal`, `log-error`,`log-warning`, `log-info`, and `log-debug`. The `define-logger` form also defines *id*`-logger`, which is a logger named `'`*id* that is a child of `(current-logger)`; the `log-`*id*`-fatal`, etc. forms use this new logger. The new logger is created when `define-logger` is evaluated.

15.5.2 Logging Events

```
(log-message  logger
              level
             [name]
              message
              data
             [prefix-message?]) → void?
  logger : logger?
  level : (or/c 'fatal 'error 'warning 'info 'debug)
  name : (or/c symbol? #f) = (object-name logger)
  message : string?
  data : any/c
  prefix-message? : any/c = #t
```

Reports an event to *logger*, which in turn distributes the information to any log receivers attached to *logger* or its ancestors that are interested in events at *level* or higher.

Log receivers can filter events based on *name*. In addition, if *name* and *prefix-message?* are not `#f`, then *message* is prefixed with the name followed by `": "` before it is sent to receivers.

Changed in version 6.0.1.10 of package `base`: Added the *prefix-message?* argument.

```
(log-level? logger level) → boolean?
  logger : logger?
  level : (or/c 'fatal 'error 'warning 'info 'debug)
```

Reports whether any log receiver attached to *logger* or one of its ancestors is interested in *level* events (or potentially lower). Use this function to avoid work generating an event for `log-message` if no receiver is interested in the information; this shortcut is built into `log-fatal`, `log-error`, `log-warning`, `log-info`, `log-debug`, and forms bound by `define-logger`, however, so it should not be used with those forms.

The result of this function can change if a garbage collection determines that a log receiver is no longer accessible (and therefore that any event information it receives will never become accessible).

```
(log-max-level logger)
 → (or/c #f 'fatal 'error 'warning 'info 'debug)
  logger : logger?
```

Similar to `log-level?`, but reports the maximum level of logging for which `log-level?` on *logger* returns `#t`. The result is `#f` if `log-level?` with *logger* currently returns `#f` for all levels.

```
(log-fatal string-expr)
(log-fatal format-string-expr v ...)
(log-error string-expr)
(log-error format-string-expr v ...)
(log-warning string-expr)
(log-warning format-string-expr v ...)
(log-info string-expr)
(log-info format-string-expr v ...)
(log-debug string-expr)
(log-debug format-string-expr v ...)
```

Log an event with the current logger, evaluating *string-expr* or (format *format-string-expr v ...*) only if the logger has receivers that are interested in the event. In addition, the current continuation's continuation marks are sent to the logger with the message string.

These form are convenient for using the current logger, but libraries should generally use a specifically named logger—typically through similar convenience forms generated by `define-logger`.

For each `log-`*level*,

```
(log-level string-expr)
```

is equivalent to

```
(let ([l (current-logger)])
  (when (log-level? l 'level)
    (log-message l 'level string-expr
                 (current-continuation-marks))))
```

while

```
(log-level format-string-expr v ...)
```

is equivalent to

```
(log-level (format format-string-expr v ...))
```

15.5.3 Receiving Logged Events

```
(log-receiver? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a log receiver, `#f` otherwise.

```
(make-log-receiver logger level [name ...] ...) → log-receiver?
  logger : logger?
  level : (or/c 'none 'fatal 'error 'warning 'info 'debug)
  name : (or/c #f symbol?) = #f
```

Creates a log receiver to receive events of importance *level* and higher as reported to *logger* and its descendants, as long as either *name* is `#f` or the event's name matches *name*.

A log receiver is a synchronizable event. It becomes ready for synchronization when a logging event is received, so use `sync` to receive an logged event. The log receiver's synchronization result is an immutable vector containing four values: the level of the event as a symbol, an immutable string for the event message, an arbitrary value that was supplied as the last argument to `log-message` when the event was logged, and a symbol or `#f` for the event name (where a symbol is usually the name of the original logger for the event).

Multiple pairs of *level* and *name* can be provided to indicate different specific *level*s for different *name*s (where *name* defaults to `#f` only for the last given *level*). A *level* for a `#f` *name* applies only to events whose names do not match any other provided *name*. If the same *name* is provided multiple times, the *level* provided with the last instance in the argument list takes precedence.

15.6 Time

```
(current-seconds) → exact-integer?
```

Returns the current time in seconds since midnight UTC, January 1, 1970.

```
(current-inexact-milliseconds) → real?
```

Returns the current time in milliseconds since midnight UTC, January 1, 1970. The result may contain fractions of a millisecond.

Example:

```
> (current-inexact-milliseconds)
1289513737015.418
```

In this example, `1289513737015` is in milliseconds and `418` is in microseconds.

```
(seconds->date secs-n [local-time?]) → date*?
  secs-n : real?
  local-time? : any/c = #t
```

Takes *secs-n*, a platform-specific time in seconds returned by current-seconds, file-or-directory-modify-seconds, or 1/1000th of current-inexact-milliseconds, and returns an instance of the date* structure type. Note that *secs-n* can include fractions of a second. If *secs-n* is too small or large, the exn:fail exception is raised.

The resulting date* reflects the time according to the local time zone if *local-time?* is #t, otherwise it reflects a date in UTC.

```
(struct date (second
              minute
              hour
              day
              month
              year
              week-day
              year-day
              dst?
              time-zone-offset)
    #:extra-constructor-name make-date
    #:transparent)
  second : (integer-in 0 60)
  minute : (integer-in 0 59)
  hour : (integer-in 0 23)
  day : (integer-in 1 31)
  month : (integer-in 1 12)
  year : exact-integer?
  week-day : (integer-in 0 6)
  year-day : (integer-in 0 365)
  dst? : boolean?
  time-zone-offset : exact-integer?
```

Represents a date. The *second* field reaches 60 only for leap seconds. The week-day field is 0 for Sunday, 1 for Monday, etc. The year-day field is 0 for January 1, 1 for January 2, etc.; the year-day field reaches 365 only in leap years.

The dst? field is #t if the date reflects a daylight-saving adjustment. The time-zone-offset field reports the number of seconds east of UTC (GMT) for the current time zone (e.g., Pacific Standard Time is -28800), including any daylight-saving adjustment (e.g., Pacific Daylight Time is -25200). When a date record is generated by seconds->date with #f as the second argument, then the dst? and time-zone-offset fields are #f and 0, respectively.

The date constructor accepts any value for dst? and converts any non-#f value to #t.

The value produced for the time-zone-offset field tends to be sensitive to the value of the TZ environment variable, especially on Unix platforms; consult the system documentation (usually under tzset) for details.

See also the racket/date library.

```
(struct date* date (nanosecond time-zone-name)
    #:extra-constructor-name make-date*)
  nanosecond : (integer-in 0 999999999)
  time-zone-name : (and/c string? immutable?)
```

Extends date with nanoseconds and a time zone name, such as "MDT", "Mountain Daylight Time", or "UTC".

When a date* record is generated by seconds->date with #f as the second argument, then the time-zone-name field is "UTC".

The date* constructor accepts a mutable string for time-zone-name and converts it to an immutable one.

```
(current-milliseconds) → exact-integer?
```

Like current-inexact-milliseconds, but coerced to a fixnum (possibly negative). Since the result is a fixnum, the value increases only over a limited (though reasonably long) time on a 32-bit platform.

```
(current-process-milliseconds thread) → exact-integer?
  thread : (or/c thread? #f)
```

Returns an amount of processor time in fixnum milliseconds that has been consumed by the Racket process on the underlying operating system. (On Unix and Mac OS X, this includes both user and system time.) If *thread* is #f, the reported time is for all Racket threads, otherwise the result is specific to the time while *thread* ran. The precision of the result is platform-specific, and since the result is a fixnum, the value increases only over a limited (though reasonably long) time on a 32-bit platform.

```
(current-gc-milliseconds) → exact-integer?
```

Returns the amount of processor time in fixnum milliseconds that has been consumed by Racket's garbage collection so far. This time is a portion of the time reported by (current-process-milliseconds), and is similarly limited.

```
                          list?
(time-apply proc lst) →   exact-integer?
                          exact-integer?
                          exact-integer?
  proc : procedure?
  lst : list?
```

Collects timing information for a procedure application.

Four values are returned: a list containing the result(s) of applying *proc* to the arguments in *lst*, the number of milliseconds of CPU time required to obtain this result, the number of "real" milliseconds required for the result, and the number of milliseconds of CPU time (included in the first result) spent on garbage collection.

The reliability of the timing numbers depends on the platform. If multiple Racket threads are running, then the reported time may include work performed by other threads.

```
(time body ...+)
```

Reports `time-apply`-style timing information for the evaluation of `expr` directly to the current output port. The result is the result of the last *body*.

15.6.1 Date Utilities

```
(require racket/date)
```
package: base

The bindings documented in this section are provided by the `racket/date` library, not `racket/base` or `racket`.

```
(current-date) → date*?
```

An abbreviation for `(seconds->date (* 0.001 (current-inexact-milliseconds)))`.

```
(date->string date [time?]) → string?
  date : date?
  time? : any/c = #f
```

Converts a date to a string. The returned string contains the time of day only if *time?*. See also `date-display-format`.

```
                          (or/c 'american
                                'chinese
                                'german
(date-display-format) →         'indian
                                'irish
                                'iso-8601
                                'rfc2822
                                'julian)
(date-display-format format) → void?
```

```
format : (or/c 'american
               'chinese
               'german
               'indian
               'irish
               'iso-8601
               'rfc2822
               'julian)
```

Parameter that determines the date string format. The initial format is `'american`.

```
(date->seconds date [local-time?]) → exact-integer?
  date : date?
  local-time? : any/c = #t
```

Finds the representation of a date in platform-specific seconds. If the platform cannot represent the specified date, `exn:fail` exception is raised.

The `week-day`, `year-day` fields of *date* are ignored. The `dst?` and `time-zone-offset` fields of *date* are also ignored; the date is assumed to be in local time by default or in UTC if *local-time?* is `#f`.

```
(date*->seconds date [local-time?]) → real?
  date : date?
  local-time? : any/c = #t
```

Like `date->seconds`, but returns an exact number that can include a fraction of a second based on `(date*-nanosecond date)` if *date* is a `date*` instance.

```
(find-seconds second
              minute
              hour
              day
              month
              year
             [local-time?]) → exact-integer?
  second : (integer-in 0 61)
  minute : (integer-in 0 59)
  hour : (integer-in 0 23)
  day : (integer-in 1 31)
  month : (integer-in 1 12)
  year : exact-nonnegative-integer?
  local-time? : any/c = #t
```

Finds the representation of a date in platform-specific seconds. The arguments correspond to the fields of the `date` structure—in local time by default or UTC if *local-time?* is `#f`.

If the platform cannot represent the specified date, an error is signaled, otherwise an integer is returned.

```
(date->julian/scalinger date) → exact-integer?
  date : date?
```

Converts a date structure (up to 2099 BCE Gregorian) into a Julian date number. The returned value is not a strict Julian number, but rather Scalinger's version, which is off by one for easier calculations.

```
(julian/scalinger->string date-number) → string?
  date-number : exact-integer?
```

Converts a Julian number (Scalinger's off-by-one version) into a string.

15.7 Environment Variables

An *environment variable set* encapsulates a partial mapping from byte strings to bytes strings. A Racket process's initial environment variable set is connected to the operating system's environment variables: accesses or changes to the set read or change operating-system environment variables for the Racket process.

Since Windows environment variables are case-insensitive, environment variable set's key byte strings on Windows are case-folded. More precisely, key byte strings are coerced to a UTF-8 encoding of characters that are converted to lowercase via `string-locale-downcase`.

The current environment variable set, which is determined by the `current-environment-variables` parameter, is propagated to a subprocess when the subprocess is created.

```
(environment-variables? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is an environment variable set, `#f` otherwise.

```
(current-environment-variables) → environment-variables?
(current-environment-variables env) → void?
  env : environment-variables?
```

A parameter that determines the environment variable set that is propagated to a subprocess and that is used as the default set for `getenv` and `putenv`.

```
(bytes-environment-variable-name? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a byte string and if it is valid for an environment variable name. An environment variable name must contain no bytes with the value `0` or `61`, where `61` is `(char->integer #\=)`. On Windows, an environment variable name also must have a non-zero length.

```
(make-environment-variables name val ... ...)
 → environment-variables?
  name : bytes-environment-variable-name?
  val : bytes-no-nuls?
```

Creates a fresh environment variable set that is initialized with the given *name* to *val* mappings.

```
(environment-variables-ref env name)
 → (or/c #f (and/c bytes-no-nuls? immutable?))
  env : environment-variables?
  name : bytes-environment-variable-name?
```

Returns the mapping for *name* in *env*, returning `#f` if *name* has no mapping.

Normally, *name* should be a byte-string encoding of a string using the default encoding of the current locale. On Windows, *name* is coerced to a UTF-8 encoding and case-normalized.

```
(environment-variables-set! env
                            name
                            maybe-bstr
                           [fail])      → any
  env : environment-variables?
  name : bytes-environment-variable-name?
  maybe-bstr : (or/c bytes-no-nuls? #f)
  fail : (-> any) = (lambda ()
                      (raise (make-exn:fail ....)))
```

Changes the mapping for *name* in *env* to *maybe-bstr*. If *maybe-bstr* is `#f` and *env* is the initial environment variable set of the Racket process, then the operating system environment-variable mapping for *name* is removed.

Normally, *name* and *maybe-bstr* should be a byte-string encoding of a string using the default encoding of the current locale. On Windows, *name* is coerced to a UTF-8 encoding and case-normalized, and *maybe-bstr* is coerced to a UTF-8 encoding if *env* is the initial environment variable set of the Racket process.

On success, the result of `environment-variables-set!` is `#<void>`. If *env* is the initial environment variable set of the Racket process, then attempting to adjust the operating system environment-variable mapping might fail for some reason, in which case *fail* is called in tail position with respect to the `environment-variables-set!`. The default *fail* raises an exception.

```
(environment-variables-names env)
 → (listof (and/c bytes-environment-variable-name? immutable?))
  env : environment-variables?
```

Returns a list of byte strings that corresponds to names mapped by *env*.

```
(environment-variables-copy env) → environment-variables?
  env : environment-variables?
```

Returns an environment variable set that is initialized with the same mappings as *env*.

```
(getenv name) → (or/c string-no-nuls? #f)
  name : string-environment-variable-name?
(putenv name value) → boolean?
  name : string-environment-variable-name?
  value : string-no-nuls?
```

Convenience wrappers for `environment-variables-ref` and `environment-variables-set!` that convert between strings and byte strings using the current locale's default encoding (using `#\?` as the replacement character for encoding errors) and always using the current environment variable set from `current-environment-variables`. The `putenv` function returns `#t` for success and `#f` for failure.

```
(string-environment-variable-name? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is a string and if its encoding using the current locale's encoding is valid for an environment variable name according to `bytes-environment-variable-name?`.

15.8 Environment and Runtime Information

```
(system-type [mode])
 → (or/c symbol? string? bytes? exact-positive-integer? vector?)
  mode : (or/c 'os 'word 'gc 'link 'machine
               'so-suffix 'so-mode 'fs-change) = 'os
```

Returns information about the operating system, build mode, or machine for a running Racket.

In `'os` mode, the possible symbol results are:

- `'unix`

- 'windows
- 'macosx

In 'word mode, the result is either 32 or 64 to indicate whether Racket is running as a 32-bit program or 64-bit program.

In 'gc mode, the possible symbol results are:

- 'cgc
- '3m

In 'link mode, the possible symbol results are:

- 'static (Unix)
- 'shared (Unix)
- 'dll (Windows)
- 'framework (Mac OS X)

Future ports of Racket may expand the list of 'os, 'gc, and 'link results.

In 'machine mode, then the result is a string, which contains further details about the current machine in a platform-specific format.

In 'so-suffix mode, then the result is a byte string that represents the file extension used for shared objects on the current platform. The byte string starts with a period, so it is suitable as a second argument to path-replace-suffix.

In 'so-mode mode, then the result is 'local if foreign libraries should be opened in "local" mode by default (as on most platforms) or 'global if foreign libraries should be opened in "global" mode.

In 'fs-change mode, the result is an immutable vector of four elements. Each element is either #f or a symbol, where a symbol indicates the presence of a property and #f indicates the absence of a property. The possible symbols, in order, are:

- 'supported — filesystem-change-evt can produce a filesystem change event to monitor filesystem changes; if this symbol is not first in the vector, all other vector elements are #f
- 'scalable — resources consumed by a filesystem change event are effectively limited only by available memory, as opposed to file-descriptor limits; this property is #f on Mac OS X and BSD variants of Unix

- `'low-latency` — creation and checking of a filesystem change event is practically instantaneous; this property is `#f` on Linux
- `'file-level` — a filesystem change event can track changes at the level of a file, as opposed to the file's directory; this property is `#f` on Windows

```
(system-language+country) → string?
```

Returns a string to identify the current user's language and country.

On Unix and Mac OS X, the string is five characters: two lowercase ASCII letters for the language, an underscore, and two uppercase ASCII letters for the country. On Windows, the string can be arbitrarily long, but the language and country are in English (all ASCII letters or spaces) separated by an underscore.

On Unix, the result is determined by checking the `LC_ALL`, `LC_TYPE`, and `LANG` environment variables, in that order (and the result is used if the environment variable's value starts with two lowercase ASCII letters, an underscore, and two uppercase ASCII letters, followed by either nothing or a period). On Windows and Mac OS X, the result is determined by system calls.

```
(system-library-subpath [mode]) → path?
  mode : (or/c 'cgc '3m #f) = (system-type 'gc)
```

Returns a relative directory path. This string can be used to build paths to system-specific files. For example, when Racket is running on Solaris on a Sparc architecture, the subpath starts `"sparc-solaris"`, while the subpath for Windows on an i386 architecture starts `"win32\\i386"`.

The optional *mode* argument specifies the relevant garbage-collection variant, which one of the possible results of `(system-type 'gc)`: `'cgc` or `'3m`. It can also be `#f`, in which case the result is independent of the garbage-collection variant.

```
(version) → (and/c string? immutable?)
```

Returns an string indicating the currently executing version of Racket.

```
(banner) → (and/c string? immutable?)
```

Returns an immutable string for Racket's start-up banner text (or the banner text for an embedding program, such as GRacket). The banner string ends with a newline.

```
(current-command-line-arguments) → (vectorof string?)
(current-command-line-arguments argv) → void?
  argv : (vectorof (and/c string? immutable?))
```

A parameter that is initialized with command-line arguments when Racket starts (not including any command-line arguments that were treated as flags for the system).

```
(current-thread-initial-stack-size) → exact-positive-integer?
(current-thread-initial-stack-size size) → void?
  size : exact-positive-integer?
```

A parameter that provides a hint about how much space to reserve for a newly created thread's local variables. The actual space used by a computation is affected by JIT compilation, but it is otherwise platform-independent.

```
(vector-set-performance-stats! results [thd]) → void?
  results : (and/c vector?
                   (not/c immutable?))
  thd : (or/c thread? #f) = #f
```

Sets elements in *results* to report current performance statistics. If *thd* is not `#f`, a particular set of thread-specific statistics are reported, otherwise a different set of global statics are reported.

For global statistics, up to 11 elements are set in the vector, starting from the beginning. If *results* has n elements where $n < 11$, then the n elements are set to the first n performance-statistics values. The reported statistics values are as follows, in the order that they are set within *results*:

- `0`: The same value as returned by `current-process-milliseconds`.
- `1`: The same value as returned by `current-milliseconds`.
- `2`: The same value as returned by `current-gc-milliseconds`.
- `3`: The number of garbage collections performed since start-up.
- `4`: The number of thread context switches performed since start-up.
- `5`: The number of internal stack overflows handled since start-up.
- `6`: The number of threads currently scheduled for execution (i.e., threads that are running, not suspended, and not unscheduled due to a synchronization).
- `7`: The number of syntax objects read from compiled code since start-up.
- `8`: The number of hash-table searches performed. When this counter reaches the maximum value of a fixnum, it overflows to the most negative fixnum.
- `9`: The number of additional hash slots searched to complete hash searches (using double hashing). When this counter reaches the maximum value of a fixnum, it overflows to the most negative fixnum.

- `10`: The number of bytes allocated for machine code that is not reported by `current-memory-use`.

For thread-specific statistics, up to 4 elements are set in the vector:

- `0`: `#t` if the thread is running, `#f` otherwise (same result as `thread-running?`).
- `1`: `#t` if the thread has terminated, `#f` otherwise (same result as `thread-dead?`).
- `2`: `#t` if the thread is currently blocked on a synchronizable event (or sleeping for some number of milliseconds), `#f` otherwise.
- `3`: The number of bytes currently in use for the thread's continuation.

15.9 Command-Line Parsing

`(require racket/cmdline)` package: `base`

The bindings documented in this section are provided by the `racket/cmdline` and `racket` libraries, but not `racket/base`.

```
(command-line optional-name-expr optional-argv-expr
              flag-clause ...
              finish-clause)
```

```
optional-name-expr =
                   | #:program name-expr

optional-argv-expr =
                   | #:argv argv-expr

      flag-clause = #:multi flag-spec ...
                  | #:once-each flag-spec ...
                  | #:once-any flag-spec ...
                  | #:final flag-spec ...
                  | #:usage-help string ...
                  | #:help-labels string ...
                  | #:ps string ...

        flag-spec = (flags id ... help-spec body ...+)
                  | (flags => handler-expr help-expr)

            flags = flag-string
                  | (flag-string ...+)

        help-spec = string
                  | (string-expr ...+)

    finish-clause =
                  | #:args arg-formals body ...+
                  | #:handlers handlers-exprs

      arg-formals = rest-id
                  | (arg ...)
                  | (arg ...+ . rest-id)

              arg = id
                  | [id default-expr]

   handlers-exprs = finish-expr arg-strings-expr
                  | finish-expr arg-strings-expr help-expr
                  | finish-expr arg-strings-expr help-expr
                    unknown-expr
```

Parses a command line according to the specification in the *flag-clause*s.

The *name-expr*, if provided, should produce a path or string to be used as the program name for reporting errors when the command-line is ill-formed. It defaults to `(find-system-path 'run-file)`. When a path is provided, only the last element of the path is used to report an error.

The *argv-expr*, if provided, must evaluate to a list or a vector of strings. It defaults to `(current-command-line-arguments)`.

The command-line is disassembled into flags, each possibly with flag-specific arguments, followed by (non-flag) arguments. Command-line strings starting with `-` or `+` are parsed as flags, but arguments to flags are never parsed as flags, and integers and decimal numbers that start with `-` or `+` are not treated as flags. Non-flag arguments in the command-line must appear after all flags and the flags' arguments. No command-line string past the first non-flag argument is parsed as a flag. The built-in `--` flag signals the end of command-line flags; any command-line string past the `--` flag is parsed as a non-flag argument.

A `#:multi`, `#:once-each`, `#:once-any`, or `#:final` clause introduces a set of command-line flag specifications. The clause tag indicates how many times the flag can appear on the command line:

- `#:multi` — Each flag specified in the set can be represented any number of times on the command line; i.e., the flags in the set are independent and each flag can be used multiple times.
- `#:once-each` — Each flag specified in the set can be represented once on the command line; i.e., the flags in the set are independent, but each flag should be specified at most once. If a flag specification is represented in the command line more than once, the `exn:fail` exception is raised.
- `#:once-any` — Only one flag specified in the set can be represented on the command line; i.e., the flags in the set are mutually exclusive. If the set is represented in the command line more than once, the `exn:fail` exception is raised.
- `#:final` — Like `#:multi`, except that no argument after the flag is treated as a flag. Note that multiple `#:final` flags can be specified if they have short names; for example, if `-a` is a `#:final` flag, then `-aa` combines two instances of `-a` in a single command-line argument.

A normal flag specification has four parts:

- *flags* — a flag string, or a set of flag strings. If a set of flags is provided, all of the flags are equivalent. Each flag string must be of the form `"-x"` or `"+x"` for some character *x*, or `"-x"` or `"++x"` for some sequence of characters *x*. An *x* cannot contain only digits or digits plus a single decimal point, since simple (signed) numbers are not treated as flags. In addition, the flags `"--"`, `"-h"`, and `"--help"` are predefined and cannot be changed.
- *id*s — identifier that are bound to the flag's arguments. The number of identifiers determines how many arguments can be provided on the command line with the flag, and the names of these identifiers will appear in the help message describing the flag. The *id*s are bound to string values in the *body*s for handling the flag.

- *help-spec* — a string or sequence of strings that describes the flag. This string is used in the help message generated by the handler for the built-in `-h` (or `--help`) flag. A single literal string can be provided, or any number of expressions that produce strings; in the latter case, strings after the first one are displayed on subsequent lines.

- *body*s — expressions that are evaluated when one of the *flags* appears on the command line. The flags are parsed left-to-right, and each sequence of *body*s is evaluated as the corresponding flag is encountered. When the *body*s are evaluated, the preceding *id*s are bound to the arguments provided for the flag on the command line.

A flag specification using `=>` escapes to a more general method of specifying the handler and help strings. In this case, the handler procedure and help string list returned by *handler-expr* and *help-expr* are used as in the *table* argument of `parse-command-line`.

A `#:usage-help` clause inserts text lines immediately after the usage line. Each string in the clause provides a separate line of text.

A `#:help-labels` clause inserts text lines into the help table of command-line flags. Each string in the clause provides a separate line of text.

A `#:ps` clause inserts text lines at the end of the help output. Each string in the clause provides a separate line of text.

After the flag clauses, a final clause handles command-line arguments that are not parsed as flags:

- Supplying no finish clause is the same as supplying `#:args () (void)`.

- For an `#:args` finish clause, identifiers in *arg-formals* are bound to the leftover command-line strings in the same way that identifiers are bound for a `lambda` expression. Thus, specifying a single *id* (without parentheses) collects all of the leftover arguments into a list. The effective arity of the *arg-formals* specification determines the number of extra command-line arguments that the user can provide, and the names of the identifiers in *arg-formals* are used in the help string. When the command-line is parsed, if the number of provided arguments cannot be matched to identifiers in *arg-formals*, the `exn:fail` exception is raised. Otherwise, `args` clause's *body*s are evaluated to handle the leftover arguments, and the result of the last *body* is the result of the `command-line` expression.

- A `#:handlers` finish clause escapes to a more general method of handling the leftover arguments. In this case, the values of the expressions are used like the last two to four arguments `parse-command-line`.

Example:

```
(define verbose-mode (make-parameter #f))
```

```
(define profiling-on (make-parameter #f))
(define optimize-level (make-parameter 0))
(define link-flags (make-parameter null))

(define file-to-compile
  (command-line
   #:program "compiler"
   #:once-each
   [("-v" "--verbose") "Compile with verbose messages"
                       (verbose-mode #t)]
   [("-p" "--profile") "Compile with profiling"
                       (profiling-on #t)]
   #:once-any
   [("-o" "--optimize-1") "Compile with optimization level 1"
                          (optimize-level 1)]
   ["--optimize-2"        (; show help on separate lines
                           "Compile with optimization level 2,"
                           "which includes all of level 1")
                          (optimize-level 2)]
   #:multi
   [("-l" "--link-flags") lf ; flag takes one argument
                          "Add a flag <lf> for the linker"
                          (link-flags (cons lf (link-flags)))]
   #:args (filename) ; expect one command-line argument: <file-
name>
   ; return the argument as a filename to compile
   filename))
```

```
(parse-command-line  name
                     argv
                     table
                     finish-proc
                     arg-help-strs
                    [help-proc
                     unknown-proc]) → any
  name : (or/c string? path?)
  argv : (or/c (listof string?) (vectorof string?))
  table : (listof (cons/c symbol? list?))
  finish-proc : ((list?) () #:rest list? . ->* . any)
  arg-help-strs : (listof string?)
  help-proc : (string? . -> . any) = (lambda (str) ....)
  unknown-proc : (string? . -> . any) = (lambda (str) ...)
```

Parses a command-line using the specification in *table*. For an overview of command-line parsing, see the `command-line` form, which provides a more convenient notation for most purposes.

The *table* argument to this procedural form encodes the information in `command-line`'s clauses, except for the `args` clause. Instead, arguments are handled by the *finish-proc* procedure, and help information about non-flag arguments is provided in *arg-help-strs*. In addition, the *finish-proc* procedure receives information accumulated while parsing flags. The *help-proc* and *unknown-proc* arguments allow customization that is not possible with `command-line`.

When there are no more flags, *finish-proc* is called with a list of information accumulated for command-line flags (see below) and the remaining non-flag arguments from the command-line. The arity of *finish-proc* determines the number of non-flag arguments accepted and required from the command-line. For example, if *finish-proc* accepts either two or three arguments, then either one or two non-flag arguments must be provided on the command-line. The *finish-proc* procedure can have any arity (see `procedure-arity`) except `0` or a list of `0`s (i.e., the procedure must at least accept one or more arguments).

The *arg-help-strs* argument is a list of strings identifying the expected (non-flag) command-line arguments, one for each argument. If an arbitrary number of arguments are allowed, the last string in *arg-help-strs* represents all of them.

The *help-proc* procedure is called with a help string if the `-h` or `--help` flag is included on the command line. If an unknown flag is encountered, the *unknown-proc* procedure is called just like a flag-handling procedure (as described below); it must at least accept one argument (the unknown flag), but it may also accept more arguments. The default *help-proc* displays the string and exits and the default *unknown-proc* raises the `exn:fail` exception.

A *table* is a list of flag specification sets. Each set is represented as a pair of two items: a mode symbol and a list of either help strings or flag specifications. A mode symbol is one of `'once-each`, `'once-any`, `'multi`, `'final`, `'help-labels`, `'usage-help`, or `'ps` with the same meanings as the corresponding clause tags in `command-line`. For the `'help-labels`, `'usage-help` or `'ps` mode, a list of help strings is provided. For the other modes, a list of flag specifications is provided, where each specification maps a number of flags to a single handler procedure. A specification is a list of three items:

- A list of strings for the flags defined by the spec. See `command-line` for information about the format of flag strings.
- A procedure to handle the flag and its arguments when one of the flags is found on the command line. The arity of this handler procedure determines the number of arguments consumed by the flag: the handler procedure is called with a flag string plus the next few arguments from the command line to match the arity of the handler procedure. The handler procedure must accept at least one argument to receive the flag. If the handler accepts arbitrarily many arguments, all of the remaining arguments are passed to the handler. A handler procedure's arity must either be a number or an `arity-at-least` value.

  The return value from the handler is added to a list that is eventually passed to *finish-proc*. If the handler returns `#<void>`, no value is added onto this list. For

all non-#<void> values returned by handlers, the order of the values in the list is the same as the order of the arguments on the command-line.

- A non-empty list for constructing help information for the spec. The first element of the list describes the flag; it can be a string or a non-empty list of strings, and in the latter case, each string is shown on its own line. Additional elements of the main list must be strings to name the expected arguments for the flag. The number of extra help strings provided for a spec must match the number of arguments accepted by the spec's handler procedure.

The following example is the same as the core example for `command-line`, translated to the procedural form:

```
(parse-command-line "compile" (current-command-line-arguments)
  `((once-each
     [("-v" "--verbose")
      ,(lambda (flag) (verbose-mode #t))
      ("Compile with verbose messages")]
     [("-p" "--profile")
      ,(lambda (flag) (profiling-on #t))
      ("Compile with profiling")])
    (once-any
     [("-o" "--optimize-1")
      ,(lambda (flag) (optimize-level 1))
      ("Compile with optimization level 1")]
     [("--optimize-2")
      ,(lambda (flag) (optimize-level 2))
      (("Compile with optimization level 2,"
        "which implies all optimizations of level 1"))])
    (multi
     [("-l" "--link-flags")
      ,(lambda (flag lf) (link-flags (cons lf (link-flags))))
      ("Add a flag <lf> for the linker" "lf")]))
   (lambda (flag-accum file) file)
   '("filename"))
```

16 Memory Management

16.1 Weak Boxes

A *weak box* is similar to a normal box (see §4.12 “Boxes”), but when the garbage collector (see §1.1.7 “Garbage Collection”) can prove that the content value of a weak box is only reachable via weak references, the content of the weak box is replaced with `#f`. A *weak reference* is a reference through a weak box, through a key reference in a weak hash table (see §4.13 “Hash Tables”), through a value in an ephemeron where the value can be replaced by `#f` (see §16.2 “Ephemerons”), or through a custodian (see §14.7 “Custodians”).

```
(make-weak-box v) → weak-box?
  v : any/c
```

Returns a new weak box that initially contains `v`.

```
(weak-box-value weak-box [gced-v]) → any/c
  weak-box : weak-box?
  gced-v : any/c = #f
```

Returns the value contained in `weak-box`. If the garbage collector has proven that the previous content value of `weak-box` was reachable only through a weak reference, then `gced-v` (which defaults to `#f`) is returned.

```
(weak-box? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a weak box, `#f` otherwise.

16.2 Ephemerons

An *ephemeron* [Hayes97] is a generalization of a weak box (see §16.1 “Weak Boxes”). Instead of just containing one value, an ephemeron holds two values: one that is considered the value of the ephemeron and another that is the ephemeron’s key. Like the value in a weak box, the value in an ephemeron may be replaced by `#f`, but when the *key* is no longer reachable (except possibly via weak references) instead of when the value is no longer reachable.

As long as an ephemeron’s value is retained, the reference is considered a non-weak reference. References to the key via the value are treated specially, however, in that the reference does not necessarily count toward the key’s reachability. A weak box can be seen as a specialization of an ephemeron where the key and value are the same.

One particularly common use of ephemerons is to combine them with a weak hash table (see §4.13 “Hash Tables”) to produce a mapping where the memory manager can reclaim key–value pairs even when the value refers to the key. A related use is to retain a reference to a value as long as any value for which it is an impersonator is reachable; see `impersonator-ephemeron`.

More precisely,

- the value in an ephemeron is replaced by `#f` when the automatic memory manager can prove that either the ephemeron or the key is reachable only through weak references (see §16.1 “Weak Boxes”); and
- nothing reachable from the value in an ephemeron counts toward the reachability of an ephemeron key (whether for the same ephemeron or another), unless the same value is reachable through a non-weak reference, or unless the value’s ephemeron key is reachable through a non-weak reference (see §16.1 “Weak Boxes” for information on weak references).

```
(make-ephemeron key v) → ephemeron?
  key : any/c
  v : any/c
```

Returns a new ephemeron whose key is *key* and whose value is initially *v*.

```
(ephemeron-value ephemeron [gced-v]) → any/c
  ephemeron : ephemeron?
  gced-v : any/c = #f
```

Returns the value contained in *ephemeron*. If the garbage collector has proven that the key for *ephemeron* is only weakly reachable, then the result is *gced-v* (which defaults to `#f`).

```
(ephemeron? v) → boolean?
  v : any/c
```

Returns `#t` if *v* is an ephemeron, `#f` otherwise.

16.3 Wills and Executors

A *will executor* manages a collection of values and associated *will* procedures (a.k.a. *finalizers*). The will procedure for each value is ready to be executed when the value has been proven (by the garbage collector) to be unreachable, except through weak references (see §16.1 “Weak Boxes”) or as the registrant for other will executors. A will is useful for triggering clean-up actions on data associated with an unreachable value, such as closing a port embedded in an object when the object is no longer used.

Calling the `will-execute` or `will-try-execute` procedure executes a will that is ready in the specified will executor. A will executor is also a synchronizable event, so `sync` or `sync/timeout` can be used to detect when a will executor has ready wills. Wills are not executed automatically, because certain programs need control to avoid race conditions. However, a program can create a thread whose sole job is to execute wills for a particular executor.

If a value is registered with multiple wills (in one or multiple executors), the wills are readied in the reverse order of registration. Since readying a will procedure makes the value reachable again, the will must be executed and the value must be proven again unreachable through only weak references before another of the wills is readied or executed. However, wills for distinct unreachable values are readied at the same time, regardless of whether the values are reachable from each other.

A will executor's register is held non-weakly until after the corresponding will procedure is executed. Thus, if the content value of a weak box (see §16.1 "Weak Boxes") is registered with a will executor, the weak box's content is not changed to `#f` until all wills have been executed for the value and the value has been proven again reachable through only weak references.

A will executor can be used as a synchronizable event (see §11.2.1 "Events"). A will executor is ready for synchronization when `will-execute` would not block; the synchronization result of a will executor is the will executor itself.

These examples show how to run cleanup actions when no synchronization is necessary. It simply runs the registered executors as they become ready in another thread.

Examples:

```
> (define an-executor (make-will-executor))

> (void
   (thread
    (λ ()
      (let loop ()
        (will-execute an-executor)
        (loop)))))

> (define (executor-proc v) (printf "a-box is now garbage\n"))

> (define a-box-to-track (box #f))

> (will-register an-executor a-box-to-track executor-proc)

> (collect-garbage)

> (set! a-box-to-track #f)
```

```
> (collect-garbage)
executor-proc: arity mismatch;
 the expected number of arguments does not match the given
number
  expected: 1
  given: 0
```

```
(make-will-executor) → will-executor?
```

Returns a new will executor with no managed values.

```
(will-executor? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a will executor, `#f` otherwise.

```
(will-register executor v proc) → void?
  executor : will-executor?
  v : any/c
  proc : (any/c . -> . any)
```

Registers the value `v` with the will procedure `proc` in the will executor `executor`. When `v` is proven unreachable, then the procedure `proc` is ready to be called with `v` as its argument via `will-execute` or `will-try-execute`. The `proc` argument is strongly referenced until the will procedure is executed.

```
(will-execute executor) → any
  executor : will-executor?
```

Invokes the will procedure for a single "unreachable" value registered with the executor `executor`. The values returned by the will procedure are the result of the `will-execute` call. If no will is ready for immediate execution, `will-execute` blocks until one is ready.

```
(will-try-execute executor) → any
  executor : any/c
```

Like `will-execute` if a will is ready for immediate execution. Otherwise, `#f` is returned.

16.4 Garbage Collection

Set the `PLTDISABLEGC` environment variable (to any value) before Racket starts to disable garbage collection.

In Racket 3m (the main variant of Racket), each garbage collection logs a message (see §15.5 “Logging”) at the `'debug` level to a logger named `'GC`. The data portion of the message is an instance of a `gc-info` prefab structure type with 10 fields as follows, but future versions of Racket may use a `gc-info` prefab structure with additional fields:

```
(struct gc-info (major? pre-amount pre-admin-amount code-amount
                        post-amount post-admin-amount
                        start-process-time end-process-time
                        start-time end-time)
  #:prefab)
```

The `major?` field indicates whether the collection was a “major” collection that inspects all memory or a “minor” collection that mostly inspects just recent allocations. The `pre-amount` field reports place-local memory use (i.e., not counting the memory use of child places) in bytes at the time that the garbage collection started. The `pre-admin-amount` is a larger number that includes memory use for the garbage collector’s overhead (such as space on memory pages that is not yet used). The `code-amount` field reports additional memory use for generated native code (which is the same just before and after a garbage collection, since it is released via finalization). The `post-amount` and `post-admin-amount` fields correspond to `pre-amount` and `pre-admin-amount`, but after garbage collection. The `start-process-time` and `end-process-time` fields report processor time (in the sense of `current-process-milliseconds`) at the start and end of garbage collection; the difference is the processor time consumed by collection. The `start-time` and `end-time` fields report real time (in the sense of `current-inexact-milliseconds`) at the start and end of garbage collection; the difference is the real time consumed by garbage collection.

```
(collect-garbage) → void?
```

Forces an immediate garbage collection (unless garbage collection is disabled by setting `PLTDISABLEGC`). Some effectively unreachable data may remain uncollected, because the collector cannot prove that it is unreachable.

The `collect-garbage` procedure provides some control over the timing of collections, but garbage will obviously be collected even if this procedure is never called (unless garbage collection is disabled).

```
(current-memory-use [cust]) → exact-nonnegative-integer?
  cust : custodian? = #f
```

Returns an estimate of the number of bytes of memory occupied by reachable data from *`cust`*. This estimate is calculated by the last garbage collection, and can be 0 if none occurred (or if none occurred since the given custodian was created). The `current-memory-use` function does *not* perform a collection by itself; doing one before the call will generally decrease the result (or increase it from 0 if no collections happened yet).

If *`cust`* is not provided, the estimate is a total reachable from any custodians.

When Racket is compiled without support for memory accounting, the estimate is the same (i.e., all memory) for any individual custodian; see also `custodian-memory-accounting-available?`.

```
(dump-memory-stats v ...) → any
  v : any/c
```

Dumps information about memory usage to the low-level error port or console.

Various combinations of `v` arguments can control the information in a dump. The information that is available depends on your Racket build; check the end of a dump from a particular build to see if it offers additional information; otherwise, all `v`s are ignored.

16.5 Phantom Byte Strings

A *phantom byte string* is a small Racket value that is treated by the Racket memory manager as having an arbitrary size, which is specified when the phantom byte string is created or when it is changed via `set-phantom-bytes!`.

A phantom byte string acts as a hint to Racket's memory manager that memory is allocated within the process but through a separate allocator, such as through a foreign library that is accessed via `ffi/unsafe`. This hint is used to trigger garbage collections or to compute the result of `current-memory-use`. Currently, the hint is used only in Racket 3m (the main variant of Racket).

```
(phantom-bytes? v) → boolean?
  v : any/c
```

Returns `#t` if `v` is a phantom byte string, `#f` otherwise.

```
(make-phantom-bytes k) → phantom-bytes?
  k : exact-nonnegative-integer?
```

Creates a phantom byte string that is treated by the Racket memory manager as being `k` bytes in size. For a large enough `k`, the `exn:fail:out-of-memory` exception is raised—either because the size is implausibly large, or because a memory limit has been installed with `custodian-limit-memory`.

```
(set-phantom-bytes! phantom-bstr k) → phantom-bytes?
  phantom-bstr : phantom-bytes?
  k : exact-nonnegative-integer?
```

Adjusts the size of a phantom byte string as it is treated by the Racket memory manager.

For example, if the memory that *phantom-bstr* represents is released through a foreign library, then `(set-phantom-bytes! phantom-bstr 0)` can reflect the change in memory use.

When *k* is larger than the current size of *phantom-bstr*, then this function can raise `exn:fail:out-of-memory`, like `make-phantom-bytes`.

17 Unsafe Operations

```
(require racket/unsafe/ops)
```
package: base

All functions and forms provided by `racket/base` and `racket` check their arguments to ensure that the arguments conform to contracts and other constraints. For example, `vector-ref` checks its arguments to ensure that the first argument is a vector, that the second argument is an exact integer, and that the second argument is between `0` and one less than the vector's length, inclusive.

Functions provided by `racket/unsafe/ops` are *unsafe*. They have certain constraints, but the constraints are not checked, which allows the system to generate and execute faster code. If arguments violate an unsafe function's constraints, the function's behavior and result is unpredictable, and the entire system can crash or become corrupted.

All of the exported bindings of `racket/unsafe/ops` are protected in the sense of `protect-out`, so access to unsafe operations can be prevented by adjusting the code inspector (see §14.10 "Code Inspectors").

17.1 Unsafe Numeric Operations

```
(unsafe-fx+ a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fx- a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fx* a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fxquotient a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fxremainder a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fxmodulo a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fxabs a) → fixnum?
  a : fixnum?
```

For fixnums: Like `+`, `-`, `*`, `quotient`, `remainder`, `modulo`, and `abs`, but constrained to consume fixnums and produce a fixnum result. The mathematical operation on `a` and `b` must

be representable as a fixnum. In the case of unsafe-fxquotient, unsafe-fxremainder, and unsafe-fxmodulo, b must not be 0.

```
(unsafe-fxand a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fxior a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fxxor a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fxnot a) → fixnum?
  a : fixnum?
(unsafe-fxlshift a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fxrshift a b) → fixnum?
  a : fixnum?
  b : fixnum?
```

For fixnums: Like bitwise-and, bitwise-ior, bitwise-xor, bitwise-not, and arithmetic-shift, but constrained to consume fixnums; the result is always a fixnum. The unsafe-fxlshift and unsafe-fxrshift operations correspond to arithmetic-shift, but require non-negative arguments; unsafe-fxlshift is a positive (i.e., left) shift, and unsafe-fxrshift is a negative (i.e., right) shift, where the number of bits to shift must be less than the number of bits used to represent a fixnum. In the case of unsafe-fxlshift, bits in the result beyond the number of bits used to represent a fixnum are effectively replaced with a copy of the high bit.

```
(unsafe-fx= a b) → boolean?
  a : fixnum?
  b : fixnum?
(unsafe-fx< a b) → boolean?
  a : fixnum?
  b : fixnum?
(unsafe-fx> a b) → boolean?
  a : fixnum?
  b : fixnum?
(unsafe-fx<= a b) → boolean?
  a : fixnum?
  b : fixnum?
(unsafe-fx>= a b) → boolean?
  a : fixnum?
  b : fixnum?
```

```
(unsafe-fxmin a b) → fixnum?
  a : fixnum?
  b : fixnum?
(unsafe-fxmax a b) → fixnum?
  a : fixnum?
  b : fixnum?
```

For fixnums: Like `=`, `<`, `>`, `<=`, `>=`, `min`, and `max`, but constrained to consume fixnums.

```
(unsafe-fl+ a b) → flonum?
  a : flonum?
  b : flonum?
(unsafe-fl- a b) → flonum?
  a : flonum?
  b : flonum?
(unsafe-fl* a b) → flonum?
  a : flonum?
  b : flonum?
(unsafe-fl/ a b) → flonum?
  a : flonum?
  b : flonum?
(unsafe-flabs a) → flonum?
  a : flonum?
```

For flonums: Unchecked versions of `fl+`, `fl-`, `fl*`, `fl/`, and `flabs`.

```
(unsafe-fl= a b) → boolean?
  a : flonum?
  b : flonum?
(unsafe-fl< a b) → boolean?
  a : flonum?
  b : flonum?
(unsafe-fl> a b) → boolean?
  a : flonum?
  b : flonum?
(unsafe-fl<= a b) → boolean?
  a : flonum?
  b : flonum?
(unsafe-fl>= a b) → boolean?
  a : flonum?
  b : flonum?
(unsafe-flmin a b) → flonum?
  a : flonum?
  b : flonum?
(unsafe-flmax a b) → flonum?
  a : flonum?
  b : flonum?
```

For flonums: Unchecked versions of `fl=`, `fl<`, `fl>`, `fl<=`, `fl>=`, `flmin`, and `flmax`.

```
(unsafe-flround a) → flonum?
  a : flonum?
(unsafe-flfloor a) → flonum?
  a : flonum?
(unsafe-flceiling a) → flonum?
  a : flonum?
(unsafe-fltruncate a) → flonum?
  a : flonum?
```

For flonums: Unchecked (potentially) versions of `flround`, `flfloor`, `flceiling`, and `fltruncate`. Currently, these bindings are simply aliases for the corresponding safe bindings.

```
(unsafe-flsin a) → flonum?
  a : flonum?
(unsafe-flcos a) → flonum?
  a : flonum?
(unsafe-fltan a) → flonum?
  a : flonum?
(unsafe-flasin a) → flonum?
  a : flonum?
(unsafe-flacos a) → flonum?
  a : flonum?
(unsafe-flatan a) → flonum?
  a : flonum?
(unsafe-fllog a) → flonum?
  a : flonum?
(unsafe-flexp a) → flonum?
  a : flonum?
(unsafe-flsqrt a) → flonum?
  a : flonum?
(unsafe-flexpt a b) → flonum?
  a : flonum?
  b : flonum?
```

For flonums: Unchecked (potentially) versions of `flsin`, `flcos`, `fltan`, `flasin`, `flacos`, `flatan`, `fllog`, `flexp`, `flsqrt`, and `flexpt`. Currently, some of these bindings are simply aliases for the corresponding safe bindings.

```
(unsafe-make-flrectangular a b)
    (and/c complex?
 →         (lambda (c) (flonum? (real-part c)))
           (lambda (c) (flonum? (imag-part c))))
  a : flonum?
  b : flonum?
```

```
(unsafe-flreal-part a) → flonum?
      (and/c complex?
  a :         (lambda (c) (flonum? (real-part c)))
              (lambda (c) (flonum? (imag-part c))))
(unsafe-flimag-part a) → flonum?
      (and/c complex?
  a :         (lambda (c) (flonum? (real-part c)))
              (lambda (c) (flonum? (imag-part c))))
```

For flonums: Unchecked versions of `make-flrectangular`, `flreal-part`, and `flimag-part`.

```
(unsafe-fx->fl a) → flonum?
  a : fixnum?
(unsafe-fl->fx a) → fixnum?
  a : flonum?
```

Unchecked conversion of a fixnum to an integer flonum and vice versa. These are similar to the safe bindings `->fl` and `fl->exact-integer`, but further constrained to consume or produce a fixnum.

```
(unsafe-flrandom rand-gen) → (and flonum? (>/c 0) (</c 1))
  rand-gen : pseudo-random-generator?
```

Unchecked version of `flrandom`.

17.2 Unsafe Data Extraction

```
(unsafe-car p) → any/c
  p : pair?
(unsafe-cdr p) → any/c
  p : pair?
(unsafe-mcar p) → any/c
  p : mpair?
(unsafe-mcdr p) → any/c
  p : mpair?
(unsafe-set-mcar! p v) → void?
  p : mpair?
  v : any/c
(unsafe-set-mcdr! p v) → void?
  p : mpair?
  v : any/c
```

Unsafe variants of `car`, `cdr`, `mcar`, `mcdr`, `set-mcar!`, and `set-mcdr!`.

```
(unsafe-cons-list v rest) → (and/c pair? list?)
  v : any/c
  rest : list?
```

Unsafe variant of `cons` that produces a pair that claims to be a list—without checking whether *rest* is a list.

```
(unsafe-list-ref lst pos) → any/c
  lst : pair?
  pos : (and/c exact-nonnegative-integer? fixnum?)
(unsafe-list-tail lst pos) → any/c
  lst : any/c
  pos : (and/c exact-nonnegative-integer? fixnum?)
```

Unsafe variants of `list-ref` and `list-tail`, where *pos* must be a fixnum, and *lst* must start with at least `(add1 pos)` (for `unsafe-list-ref`) or *pos* (for `unsafe-list-tail`) pairs.

```
(unsafe-unbox b) → fixnum?
  b : box?
(unsafe-set-box! b k) → void?
  b : box?
  k : fixnum?
(unsafe-unbox* v) → any/c
  v : (and/c box? (not/c impersonator?))
(unsafe-set-box*! v val) → void?
  v : (and/c box? (not/c impersonator?))
  val : any/c
```

Unsafe versions of `unbox` and `set-box!`, where the `box*` variants can be faster but do not work on impersonators.

```
(unsafe-box*-cas! loc old new) → boolean?
  loc : box?
  old : any/c
  new : any/c
```

Unsafe version of `box-cas!`. Like `unsafe-set-box*!`, it does not work on impersonators.

```
(unsafe-vector-length v) → fixnum?
  v : vector?
(unsafe-vector-ref v k) → any/c
  v : vector?
  k : fixnum?
```

```
(unsafe-vector-set! v k val) → void?
  v : vector?
  k : fixnum?
  val : any/c
(unsafe-vector*-length v) → fixnum?
  v : (and/c vector? (not/c impersonator?))
(unsafe-vector*-ref v k) → any/c
  v : (and/c vector? (not/c impersonator?))
  k : fixnum?
(unsafe-vector*-set! v k val) → void?
  v : (and/c vector? (not/c impersonator?))
  k : fixnum?
  val : any/c
```

Unsafe versions of `vector-length`, `vector-ref`, and `vector-set!`, where the `vector*` variants can be faster but do not work on impersonators.

A vector's size can never be larger than a fixnum, so even `vector-length` always returns a fixnum.

```
(unsafe-string-length str) → fixnum?
  str : string?
(unsafe-string-ref str k)
 → (and/c char? (lambda (ch) (<= 0 (char->integer ch) 255)))
  str : string?
  k : fixnum?
(unsafe-string-set! str k ch) → void?
  str : (and/c string? (not/c immutable?))
  k : fixnum?
  ch : char?
```

Unsafe versions of `string-length`, `string-ref`, and `string-set!`. The `unsafe-string-ref` procedure can be used only when the result will be a Latin-1 character. A string's size can never be larger than a fixnum (so even `string-length` always returns a fixnum).

```
(unsafe-bytes-length bstr) → fixnum?
  bstr : bytes?
(unsafe-bytes-ref bstr k) → byte?
  bstr : bytes?
  k : fixnum?
(unsafe-bytes-set! bstr k b) → void?
  bstr : (and/c bytes? (not/c immutable?))
  k : fixnum?
  b : byte?
```

Unsafe versions of `bytes-length`, `bytes-ref`, and `bytes-set!`. A bytes's size can never be larger than a fixnum (so even `bytes-length` always returns a fixnum).

```
(unsafe-flvector-length v) → fixnum?
  v : flvector?
(unsafe-flvector-ref v k) → any/c
  v : flvector?
  k : fixnum?
(unsafe-flvector-set! v k x) → void?
  v : flvector?
  k : fixnum?
  x : flonum?
```

Unsafe versions of `flvector-length`, `flvector-ref`, and `flvector-set!`. A flvector's size can never be larger than a fixnum (so even `flvector-length` always returns a fixnum).

```
(unsafe-f64vector-ref vec k) → flonum?
  vec : f64vector?
  k : fixnum?
(unsafe-f64vector-set! vec k n) → void?
  vec : f64vector?
  k : fixnum?
  n : flonum?
```

Unsafe versions of `f64vector-ref` and `f64vector-set!`.

```
(unsafe-s16vector-ref vec k) → (integer-in -32768 32767)
  vec : s16vector?
  k : fixnum?
(unsafe-s16vector-set! vec k n) → void?
  vec : s16vector?
  k : fixnum?
  n : (integer-in -32768 32767)
```

Unsafe versions of `s16vector-ref` and `s16vector-set!`.

```
(unsafe-u16vector-ref vec k) → (integer-in 0 65535)
  vec : u16vector?
  k : fixnum?
(unsafe-u16vector-set! vec k n) → void?
  vec : u16vector?
  k : fixnum?
  n : (integer-in 0 65535)
```

Unsafe versions of `u16vector-ref` and `u16vector-set!`.

```
(unsafe-struct-ref v k) → any/c
  v : any/c
  k : fixnum?
(unsafe-struct-set! v k val) → void?
  v : any/c
  k : fixnum?
  val : any/c
(unsafe-struct*-ref v k) → any/c
  v : (not/c impersonator?)
  k : fixnum?
(unsafe-struct*-set! v k val) → void?
  v : (not/c impersonator?)
  k : fixnum?
  val : any/c
```

Unsafe field access and update for an instance of a structure type, where the `struct*` variants can be faster but do not work on impersonators. The index *k* must be between `0` (inclusive) and the number of fields in the structure (exclusive). In the case of `unsafe-struct-set!`, the field must be mutable.

17.3 Unsafe Extflonum Operations

```
(unsafe-extfl+ a b) → extflonum?
  a : extflonum?
  b : extflonum?
(unsafe-extfl- a b) → extflonum?
  a : extflonum?
  b : extflonum?
(unsafe-extfl* a b) → extflonum?
  a : extflonum?
  b : extflonum?
(unsafe-extfl/ a b) → extflonum?
  a : extflonum?
  b : extflonum?
(unsafe-extflabs a) → extflonum?
  a : extflonum?
```

Unchecked versions of `extfl+`, `extfl-`, `extfl*`, `extfl/`, and `extflabs`.

```
(unsafe-extfl= a b) → boolean?
  a : extflonum?
  b : extflonum?
```

```
(unsafe-extfl< a b) → boolean?
  a : extflonum?
  b : extflonum?
(unsafe-extfl> a b) → boolean?
  a : extflonum?
  b : extflonum?
(unsafe-extfl<= a b) → boolean?
  a : extflonum?
  b : extflonum?
(unsafe-extfl>= a b) → boolean?
  a : extflonum?
  b : extflonum?
(unsafe-extflmin a b) → extflonum?
  a : extflonum?
  b : extflonum?
(unsafe-extflmax a b) → extflonum?
  a : extflonum?
  b : extflonum?
```

Unchecked versions of `extfl=`, `extfl<`, `extfl>`, `extfl<=`, `extfl>=`, `extflmin`, and `extflmax`.

```
(unsafe-extflround a) → extflonum?
  a : extflonum?
(unsafe-extflfloor a) → extflonum?
  a : extflonum?
(unsafe-extflceiling a) → extflonum?
  a : extflonum?
(unsafe-extfltruncate a) → extflonum?
  a : extflonum?
```

Unchecked (potentially) versions of `extflround`, `extflfloor`, `extflceiling`, and `extfltruncate`. Currently, these bindings are simply aliases for the corresponding safe bindings.

```
(unsafe-extflsin a) → extflonum?
  a : extflonum?
(unsafe-extflcos a) → extflonum?
  a : extflonum?
(unsafe-extfltan a) → extflonum?
  a : extflonum?
(unsafe-extflasin a) → extflonum?
  a : extflonum?
(unsafe-extflacos a) → extflonum?
  a : extflonum?
(unsafe-extflatan a) → extflonum?
  a : extflonum?
```

```
(unsafe-extfllog a) → extflonum?
  a : extflonum?
(unsafe-extflexp a) → extflonum?
  a : extflonum?
(unsafe-extflsqrt a) → extflonum?
  a : extflonum?
(unsafe-extflexpt a b) → extflonum?
  a : extflonum?
  b : extflonum?
```

Unchecked (potentially) versions of `extflsin`, `extflcos`, `extfltan`, `extflasin`, `extflacos`, `extflatan`, `extfllog`, `extflexp`, `extflsqrt`, and `extflexpt`. Currently, some of these bindings are simply aliases for the corresponding safe bindings.

```
(unsafe-fx->extfl a) → extflonum?
  a : fixnum?
(unsafe-extfl->fx a) → fixnum?
  a : extflonum?
```

Unchecked conversion of a fixnum to an integer extflonum and vice versa. These are similar to the safe bindings `->extfl` and `extfl->exact-integer`, but further constrained to consume or produce a fixnum.

```
(unsafe-extflvector-length v) → fixnum?
  v : extflvector?
(unsafe-extflvector-ref v k) → any/c
  v : extflvector?
  k : fixnum?
(unsafe-extflvector-set! v k x) → void?
  v : extflvector?
  k : fixnum?
  x : extflonum?
```

Unchecked versions of `extflvector-length`, `extflvector-ref`, and `extflvector-set!`. A extflvector's size can never be larger than a fixnum (so even `extflvector-length` always returns a fixnum).

17.4 Unsafe Undefined

`(require racket/unsafe/undefined)` package: `base`

The bindings documented in this section are provided by the `racket/unsafe/undefined` library, not `racket/base` or `racket`.

The constant unsafe-undefined is used internally as a placeholder value. For example, it is used by letrec as a value for a variable that has not yet been assigned a value. Unlike the undefined value exported by racket/undefined, however, the unsafe-undefined value should not leak as the result of a safe expression. Expression results that potentially produce unsafe-undefined can be guarded by check-not-unsafe-undefined, so that an exception can be raised instead of producing an undefined value.

The unsafe-undefined value is always eq? to itself.

Added in version 6.0.1.2 of package base.

```
unsafe-undefined : any/c
```

The unsafe "undefined" constant.

```
(check-not-unsafe-undefined v sym)
 → (and/c any/c (not/c (one-of/c unsafe-undefined)))
  v : any/c
  sym : symbol?
```

Checks whether v is unsafe-undefined, and raises exn:fail:contract:variable in that case with an error message along the lines of "sym: undefined; use before initialization." If v is not unsafe-undefined, then v is returned.

```
(check-not-unsafe-undefined/assign v sym)
 → (and/c any/c (not/c (one-of/c unsafe-undefined)))
  v : any/c
  sym : symbol?
```

The same as check-not-unsafe-undefined, except that the error message (if any) is along the lines of "sym: undefined; assignment before initialization."

```
(chaperone-struct-unsafe-undefined v) → any/c
  v : any/c
```

Chaperones v if it is a structure (as viewed through some inspector). Every access of a field in the structure is checked to prevent returning unsafe-undefined. Similarly, every assignment to a field in the structure is checked (unless the check disabled as described below) to prevent assignment of a field whose current value is unsafe-undefined.

When a field access would otherwise produce unsafe-undefined or when a field assignment would replace unsafe-undefined, the exn:fail:contract exception is raised.

The chaperone's field-assignment check is disabled whenever (continuation-mark-set-first #f prop:chaperone-unsafe-undefined) returns unsafe-undefined. Thus, a field-initializing assignment—one that is intended to replace the unsafe-undefined value

of a field—should be wrapped with (with-continuation-mark prop:chaperone-unsafe-undefined unsafe-undefined ....).

prop:chaperone-unsafe-undefined : struct-type-property?

A structure type property that causes a structure type's constructor to produce a chaperone of an instance in the same way as chaperone-struct-unsafe-undefined.

The property value should be a list of symbols used as field names, but the list should be in reverse order of the structure's fields. When a field access or assignment would produce or replace unsafe-undefined, the exn:fail:contract:variable exception is raised if a field name is provided by the structure property's value, otherwise the exn:fail:contract exception is raised.

18 Running Racket

18.1 Running Racket or GRacket

The core Racket run-time system is available in two main variants:

- Racket, which provides the primitives libraries on which `racket/base` is implemented. On Unix and Mac OS X, the executable is called `racket`. On Windows, the executable is called `Racket.exe`.
- GRacket, which is a GUI variant of `racket` to the degree that the system distinguishes them. On Unix, the executable is called `gracket`, and single-instance flags and X11-related flags are handled and communicated specially to the `racket/gui/base` library. On Windows, the executable is called `GRacket.exe`, and it is a GUI application (as opposed to a console application) that implements single-instance support. On Mac OS X, the `gracket` script launches `GRacket.app`.

18.1.1 Initialization

On start-up, the top-level environment contains no bindings—not even `#%app` for function application. Primitive modules with names that start with `#%` are defined, but they are not meant for direct use, and the set of such modules can change. For example, the `'#%kernel` module is eventually used to bootstrap the implementation of `racket/base`.

The first action of Racket or GRacket is to initialize `current-library-collection-paths` to the result of `(find-library-collection-paths` *`pre-extras extras`*`)`, where *`pre-extras`* is normally `null` and *`extras`* are extra directory paths provided in order in the command line with `-S`/`--search`. An executable created from the Racket or GRacket executable can embed paths used as *`pre-extras`*.

Racket and GRacket next `require` `racket/init` and `racket/gui/init`, respectively, but only if the command line does not specify a `require` flag (`-t`/`--require`, `-l`/`--lib`, or `-u`/`--require-script`) before any `eval`, `load`, or read-eval-print-loop flag (`-e`/`--eval`, `-f`/`--load`, `-r`/`--script`, `-m`/`--main`, or `-i`/`--repl`). The initialization library can be changed with the `-I` configuration option. The `configure-runtime` submodule of the initialization library or the `'configure-runtime` property of the initialization library's language is used before the library is instantiated; see §18.1.5 "Language Run-Time Configuration".

After potentially loading the initialization module, expression `eval`s, files `load`s, and module `require`s are executed in the order that they are provided on the command line. If any raises an uncaught exception, then the remaining `eval`s, `load`s, and `require`s are skipped. If the first `require` precedes any `eval` or `load` so that the initialization library is skipped,

then the configure-runtime submodule of the required module or the 'configure-runtime property of the required module's library language is used before the module is instantiated; see §18.1.5 "Language Run-Time Configuration".

After running all command-line expressions, files, and modules, Racket or GRacket then starts a read-eval-print loop for interactive evaluation if no command line flags are provided other than configuration options. If any command-line argument is provided that is not a configuration option, then the read-eval-print-loop is not started, unless the -i/--repl flag is provided on the command line to specifically re-enable it. In addition, just before the command line is started, Racket loads the file (find-system-path 'init-file) and GRacket loads the file (find-graphical-system-path 'init-file) is loaded, unless the -q/--no-init-file flag is specified on the command line.

Finally, before Racket or GRacket exits, it calls the procedure that is the current value of executable-yield-handler in the main thread, unless the -V/--no-yield command-line flag is specified. Requiring racket/gui/base sets this parameter call (racket 'yield).

18.1.2 Exit Status

The default exit status for a Racket or GRacket process is non-zero if an error occurs during a command-line eval (via -e, etc.), load (via -f, -r, etc.), or require (via --l, -t, etc.), but only when no read-eval-print loop is started. Otherwise, the default exit status is 0.

In all cases, a call to exit (when the default exit handler is in place) can end the process with a specific status value.

18.1.3 Init Libraries

(require racket/init) package: base

The racket/init library is the default start-up library for Racket. It re-exports the racket, racket/enter and racket/help libraries, and it sets current-print to use pretty-print.

(require racket/gui/init) package: gui-lib

The racket/gui/init library is the default start-up library for GRacket. It re-exports the racket/init and racket/gui/base libraries, and it sets current-load to use text-editor-load-handler.

(require racket/language-info) package: base

The racket/language-info library provides a get-info function that takes any value and returns another function; the returned function takes a key value and a default value, and

it returns `'(#(racket/runtime-config configure #f))` if the key is `'configure-runtime` or the default value otherwise.

See also §17.3.6 "Module-Handling Configuration" in ***The Racket Guide***.

The vector `'#(racket/language-info get-info #f)` is suitable for attaching to a module as its language info to get the same language information as the `racket/base` language.

`(require racket/runtime-config)` package: `base`

The `racket/runtime-config` library provides a `configure` function that takes any value and sets `print-as-expression` to `#t`.

The vector `#(racket/runtime-config configure #f)` is suitable as a member of a list of runtime-configuration specification (as returned by a module's language-information function for the key `'configure-runtime`) to obtain the same runtime configuration as for the `racket/base` language.

18.1.4 Command Line

The Racket and GRacket executables recognize the following command-line flags:

- File and expression options:
  - `-e` ⟨*expr*⟩ or `--eval` ⟨*expr*⟩ : `eval`s ⟨*expr*⟩. The results of the evaluation are printed via `current-print`.
  - `-f` ⟨*file*⟩ or `--load` ⟨*file*⟩ : `load`s ⟨*file*⟩; if ⟨*file*⟩ is `"-"`, then expressions are read and evaluated from standard input.
  - `-t` ⟨*file*⟩ or `--require` ⟨*file*⟩ : `requires` ⟨*file*⟩, and then `requires (submod (file "`⟨*file*⟩`") main)` if available.
  - `-l` ⟨*path*⟩ or `--lib` ⟨*path*⟩ : `requires (lib "`⟨*path*⟩`")`, and then `requires (submod (lib "`⟨*path*⟩`") main)` if available.
  - `-p` ⟨*package*⟩ : `requires (planet "`⟨*package*⟩`")`, and then `requires (submod (planet "`⟨*package*⟩`") main)` if available.
  - `-r` ⟨*file*⟩ or `--script` ⟨*file*⟩ : `load`s ⟨*file*⟩ as a script. This flag is like `-t` ⟨*file*⟩ plus `-N` ⟨*file*⟩ to set the program name and `--` to cause all further command-line elements to be treated as non-flag arguments.

    Despite its name, `--script` is not usually used for Unix scripts. See §21.2 "Scripts" for more information on scripts.
  - `-u` ⟨*file*⟩ or `--require-script` ⟨*file*⟩ : `requires` ⟨*file*⟩ as a script; This flag is like `-t` ⟨*file*⟩ plus `-N` ⟨*file*⟩ to set the program name and `--` to cause all further command-line elements to be treated as non-flag arguments.
  - `-k` ⟨*n*⟩ ⟨*m*⟩ ⟨*p*⟩ : Loads code embedded in the executable from file position ⟨*n*⟩ to ⟨*m*⟩ and from ⟨*m*⟩ to ⟨*p*⟩. (On Mac OS X, ⟨*n*⟩, ⟨*m*⟩, and ⟨*p*⟩ are relative to a `__PLTSCHEME` segment in the executable.) The first range is loaded in every new place, and any modules declared in that range are considered predefined in the sense of `module-predefined?`. This option is normally embedded in a stand-alone binary that also embeds Racket code.

- `-m` or `--main` : Evaluates a call to `main` as bound in the top-level environment. All of the command-line arguments that are not processed as options (i.e., the arguments put into `current-command-line-arguments`) are passed as arguments to `main`. The results of the call are printed via `current-print`.

  The call to `main` is constructed as an expression `(main arg-str ...)` where the lexical context of the expression gives `#%app` and `#%datum` bindings as `#%plain-app` and `#%datum`, but the lexical context of `main` is the top-level environment.

- Interaction options:

  - `-i` or `--repl` : Runs an interactive read-eval-print loop, using either `read-eval-print-loop` (Racket) or `graphical-read-eval-print-loop` (GRacket) after showing `(banner)` and loading `(find-system-path 'init-file)`. In the case of Racket, `(read-eval-print-loop)` is followed by `(newline)`. For GRacket, supply the `-z`/`--text-repl` configuration option to use `read-eval-print-loop` (and `newline`) instead of `graphical-read-eval-print-loop`.
  - `-n` or `--no-lib` : Skips requiring the initialization library (i.e., `racket/init` or `racket/gui/init`, unless it is changed with the `-I` flag) when not otherwise disabled.
  - `-v` or `--version` : Shows `(banner)`.
  - `-K` or `--back` : GRacket, Mac OS X only; leave application in the background.
  - `-V` `--no-yield` : Skips final `executable-yield-handler` action, which normally waits until all frames are closed, etc. in the main eventspace before exiting for programs that use `racket/gui/base`.

- *Configuration options*:

  - `-c` or `--no-compiled` : Disables loading of compiled byte-code `".zo"` files, by initializing `current-compiled-file-paths` to `null`. Use judiciously: this effectively ignores the content of all `"compiled"` subdirectories, so that any used modules are compiled on the fly—even `racket/base` and its dependencies—which leads to prohibitively expensive run times.
  - `-q` or `--no-init-file` : Skips loading `(find-system-path 'init-file)` for `-i`/`--repl`.
  - `-z` or `--text-repl` : GRacket only; changes `-i`/`--repl` to use `textual-read-eval-print-loop` instead of `graphical-read-eval-print-loop`.
  - `-I` ⟨*path*⟩ : Sets `(lib "`⟨*path*⟩`")` as the path to `require` to initialize the namespace, unless namespace initialization is disabled. Using this flag can effectively set the language for the read-eval-print loop and other top-level evaluation.
  - `-X` ⟨*dir*⟩ or `--collects` ⟨*dir*⟩ : Sets ⟨*dir*⟩ as the path to the main collection of libraries by making `(find-system-path 'collects-dir)` produce ⟨*dir*⟩. If ⟨*dir*⟩ is an empty string, then `(find-system-path 'collects-dir)` returns `"."`, but `current-library-collection-paths` is initialized to the empty list, and `use-collection-link-paths` is initialized to `#f`.

- `-S` ⟨*dir*⟩ or `--search` ⟨*dir*⟩ : Adds ⟨*dir*⟩ to the default library collection search path after the main collection directory. If the `-S`/`--dir` flag is supplied multiple times, the search order is as supplied.
- `-R` ⟨*paths*⟩ or `--compiled` ⟨*paths*⟩ : Sets the initial value of the `current-compiled-file-roots` parameter, overriding any `PLTCOMPILEDROOTS` setting. The ⟨*paths*⟩ argument is parsed in the same way as `PLTCOMPILEDROOTS` (see `current-compiled-file-roots`).
- `-G` ⟨*dir*⟩ or `--config` ⟨*dir*⟩ : Sets the directory that is returned by `(find-system-path 'config-dir)`.
- `-A` ⟨*dir*⟩ or `--addon` ⟨*dir*⟩ : Sets the directory that is returned by `(find-system-path 'addon-dir)`.
- `-U` or `--no-user-path` : Omits user-specific paths in the search for collections, C libraries, etc. by initializing the `use-user-specific-search-paths` parameter to `#f`.
- `-N` ⟨*file*⟩ or `--name` ⟨*file*⟩ : sets the name of the executable as reported by `(find-system-path 'run-file)` to ⟨*file*⟩.
- `-J` ⟨*name*⟩ or `--wm-class` ⟨*name*⟩ : GRacket, Unix only; sets the `WM_CLASS` program class to ⟨*name*⟩ (while the `WM_CLASS` program name is derived from the executable name or a `-N`/`--name` argument).
- `-j` or `--no-jit` : Disables the native-code just-in-time compiler by setting the `eval-jit-enabled` parameter to `#f`.
- `-d` or `--no-delay` : Disables on-demand parsing of compiled code and syntax objects by setting the `read-on-demand-source` parameter to `#f`.
- `-b` or `--binary` : Requests binary mode, instead of text mode, for the process's input, out, and error ports. This flag currently has no effect, because binary mode is always used.
- `-W` ⟨*levels*⟩ or `--warn` ⟨*levels*⟩ : Sets the logging level for writing events to the original error port. The possible ⟨*level*⟩ values are the same as for the `PLTSTDERR` environment variable. See §15.5 "Logging" for more information.
- `-L` ⟨*levels*⟩ or `--syslog` ⟨*levels*⟩ : Sets the logging level for writing events to the system log. The possible ⟨*level*⟩ values are the same as for the `PLTSYSLOG` environment variable. See §15.5 "Logging" for more information.

- Meta options:
  - `--` : No argument following this flag is itself used as a flag.
  - `-h` or `--help` : Shows information about the command-line flags and start-up process and exits, ignoring all other flags.

If at least one command-line argument is provided, and if the first one after any configuration option is not a flag, then a `-u`/`--require-script` flag is implicitly added before the first non-flag argument.

If no command-line arguments are supplied other than configuration options, then the `-i`/`--repl` flag is effectively added.

For GRacket on Unix, the follow flags are recognized when they appear at the beginning of the command line, and they count as configuration options (i.e., they do not disable the read-eval-print loop or prevent the insertion of `-u`/`--require-script`):

- `-display` ⟨*display*⟩ : Sets the X11 display to use.
- `-geometry` ⟨*arg*⟩, `-bg` ⟨*arg*⟩, `-background` ⟨*arg*⟩, `-fg` ⟨*arg*⟩, `-foreground` ⟨*arg*⟩, `-fn` ⟨*arg*⟩, `-font` ⟨*arg*⟩, `-iconic`, `-name` ⟨*arg*⟩, `-rv`, `-reverse`, `+rv`, `-selectionTimeout` ⟨*arg*⟩, `-synchronous`, `-title` ⟨*arg*⟩, `-xnllanguage` ⟨*arg*⟩, or `-xrm` ⟨*arg*⟩ : Standard X11 arguments that are mostly ignored but accepted for compatibility with other X11 programs. The `-synchronous` flag behaves in the usual way.
- `-singleInstance` : If an existing GRacket is already running on the same X11 display, if it was started on a machine with the same hostname, and if it was started with the same name as reported by `(find-system-path 'run-file)`—possibly set with the `-N`/`--name` command-line argument—then all non-option command-line arguments are treated as filenames and sent to the existing GRacket instance via the application file handler (see `application-file-handler`).

Similarly, on Mac OS X, a leading switch starting with `-psn_` is treated as a special configuration option. It indicates that Finder started the application, so the current input, output, and error output are redirected to a GUI window.

Multiple single-letter switches (the ones preceded by a single `-`) can be collapsed into a single switch by concatenating the letters, as long as the first switch is not `--`. The arguments for each switch are placed after the collapsed switches (in the order of the switches). For example,

`-ifve` ⟨*file*⟩ ⟨*expr*⟩

and

`-i -f` ⟨*file*⟩ `-v -e` ⟨*expr*⟩

are equivalent. If a collapsed `--` appears before other collapsed switches in the same collapsed set, it is implicitly moved to the end of the collapsed set.

Extra arguments following the last option are available from the `current-command-line-arguments` parameter.

18.1.5 Language Run-Time Configuration

See also §17.3.6 "Module-Handling Configuration" in ***The Racket Guide***.

A module can have a `configure-runtime` submodule that is `dynamic-require`d before the module itself when a module is the main module of a program. Normally, a `configure-runtime` submodule is added to a module by the module's language (i.e., by the `#%module-begin` form among a `module`'s initial bindings).

Alternatively or in addition, an older protocol is in place. When a module is implemented using `#lang`, the language after `#lang` can specify configuration actions to perform when a module using the language is the main module of a program. The language specifies run-time configuration by

- attaching a `'module-language` syntax property to the module as read from its source (see `module` and `module-compiled-language-info`);
- having the function indicated by the `'module-language` syntax property recognize the `'configure-runtime` key, for which it returns a list of vectors; each vector must have the form `(vector mp name val)` where *mp* is a module path, *name* is a symbol, and *val* is an arbitrary value; and
- having each function called as `((dynamic-require mp name) val)` configure the run-time environment, typically by setting parameters such as `current-print`.

A `'configure-runtime` query returns a list of vectors, instead of directly configuring the environment, so that the indicated modules to be bundled with a program when creating a stand-alone executable; see §2 "`raco exe`: Creating Stand-Alone Executables" in *raco: Racket Command-Line Tools*.

For information on defining a new `#lang` language, see `syntax/module-reader`.

18.2 Libraries and Collections

A *library* is `module` declaration for use by multiple programs. Racket further groups libraries into *collections*. Typically, collections are added via *packages* (see *Package Management in Racket*); the package manager works outside of the Racket core, but it configures the core run-time system through collection links files.

Libraries in collections are referenced through `lib` paths (see `require`) or symbolic shorthands. For example, the following module uses the `"getinfo.rkt"` library module from the `"setup"` collection, and the `"cards.rkt"` library module from the `"games"` collection's `"cards"` subcollection:

```
#lang racket
(require (lib "setup/getinfo.rkt")
```

```
          (lib "games/cards/cards.rkt"))
....
```

This example is more compactly and more commonly written using symbolic shorthands:

```
#lang racket
(require setup/getinfo
         games/cards/cards)
....
```

When an identifier *id* is used in a `require` form, it is converted to `(lib` *rel-string*`)` where *rel-string* is the string form of *id*.

A *rel-string* in `(lib` *rel-string*`)` consists of one or more path elements that name collections, and then a final path element that names a library file; the path elements are separated by `/`. If *rel-string* contains no `/`s, then `/main.rkt` is implicitly appended to the path. If *rel-string* contains `/` but does not end with a file suffix, then `.rkt` is implicitly appended to the path.

Libraries also can be distributed via PLaneT packages. Such libraries are referenced through a `planet` module path (see `require`) and are downloaded by Racket on demand, instead of referenced through collections.

The translation of a `planet` or `lib` path to a `module` declaration is determined by the module name resolver, as specified by the `current-module-name-resolver` parameter.

18.2.1 Collection Search Configuration

For the default module name resolver, the search path for collections is determined by the `current-library-collection-links` parameter and the `current-library-collection-paths` parameter:

- The most primitive collection-based modules are located in `"collects"` directory relative to the Racket executable. Libraries for a collection are grouped within a directory whose name matches the collection name. The path to the `"collects"` directory is normally included in `current-library-collection-paths`.
- Collection-based libraries also can be installed other directories, perhaps user-specific, that are structured like the `"collects"` directory. Those additional directories can be included in the `current-library-collection-paths` parameter either dynamically, through command-line arguments to `racket`, or by setting the `PLTCOLLECTS` environment variable; see `find-library-collection-paths`.
- Collection links files provide a mapping from top-level collection names to directories, plus additional `"collects"`-like directories (that have subdirectories with names that

match collection names). Each collection links file to be searched is referenced by the current-library-collection-links parameter; the parameter references the file, and not the file's content, so that changes to the file can be detected and affect later module resolution. See also find-library-collection-links.

- The current-library-collection-links parameter's value can also include hash tables that provide the same content as collection links files: a mapping from collection names in symbol form to a list of paths for the collection, or from #f to a list of "collects"-like paths.
- Finally, the current-library-collection-links parameter's value includes #f to indicate the point in the search process at which the module-name resolver should check current-library-collection-paths relative to the files and hash tables in current-library-collection-links.

To resolve a module reference *rel-string*, the default module name resolver searches collection links in current-library-collection-links from first to last to locate the first directory that contains *rel-string*, splicing a search through in current-library-collection-paths where in current-library-collection-links contains #f. The filesystem tree for each element in the link table and search path is effectively *spliced* together with the filesystem trees of other path elements that correspond to the same collection. Some Racket tools rely on unique resolution of module path names, so an installation and configuration should not allow multiple files to match the same collection and file combination.

The value of the current-library-collection-links parameter is initialized by the racket executable to the result of (find-library-collection-links), and the value of the current-library-collection-paths parameter is initialized to the result of (find-library-collection-paths).

18.2.2 Collection Links

Collection links files are used by collection-file-path, collection-path, and the default module name resolver to locate collections before trying the (current-library-collection-paths) search path. The collection links files to use are determined by the current-library-collection-links parameter, which is initialized to the result of find-library-collection-links.

A collection links file is read with default reader parameter settings to obtain a list. Every element of the list must be a link specification with one of the forms (list *string path*), (list *string path regexp*), (list 'root *path*), (list 'root *path regexp*), (list 'static-root *path*), (list 'static-root *path regexp*). A *string* names a top-level collection, in which case *path* is a path that can be used as the collection's path (directly, as opposed to a subdirectory of *path* named by *string*). A 'root entry, in contrast, acts like an path in (current-library-collection-paths). A

'static-root entry is like a 'root entry, but where the immediate content of the directory is assumed not to change unless the collection links file changes. If *path* is a relative path, it is relative to the directory containing the collection links file. If *regexp* is specified in a link, then the link is used only if (regexp-match? *regexp* (version)) produces a true result.

A single top-level collection can have multiple links in a collection links file, and any number of 'root entries can appear. The corresponding paths are effectively spliced together, since the paths are tried in order to locate a file or sub-collection.

The raco link command-link tool can display, install, and remove links in a collection links file. See §9 "raco link: Library Collection Links" in *raco: Racket Command-Line Tools* for more information.

18.2.3 Collection Paths and Parameters

```
(find-library-collection-paths [pre-extras
                                post-extras]) → (listof path?)
  pre-extras : (listof path-string?) = null
  post-extras : (listof path-string?) = null
```

Produces a list of paths, which is normally used to initialize current-library-collection-paths, as follows:

- The path produced by (build-path (find-system-path 'addon-dir) (get-installation-name) "collects") is the first element of the default collection path list, unless the value of the use-user-specific-search-paths parameter is #f.
- Extra directories provided in *pre-extras* are included next to the default collection path list, converted to complete paths relative to the executable.
- If the directory specified by (find-system-path 'collects-dir) is absolute, or if it is relative (to the executable) and it exists, then it is added to the end of the default collection path list.
- Extra directories provided in *post-extras* are included last in the default collection path list, converted to complete paths relative to the executable.
- If the PLTCOLLECTS environment variable is defined, it is combined with the default list using path-list-string->path-list, as long as the value of use-user-specific-search-paths is true. If it is not defined or if the value use-user-specific-search-paths is #f, the default collection path list (as constructed by the first three bullets above) is used directly.

  Note that on Unix and Mac OS X, paths are separated by :, and on Windows by ;. Also, path-list-string->path-list splices the default paths at an empty

path, for example, with many Unix shells you can set PLTCOLLECTS to ":`pwd`", "`pwd`:", or "`pwd`" to specify search the current directory after, before, or instead of the default paths, respectively.

```
(find-library-collection-links)
 → (listof (or/c #f (and/c path? complete-path?)))
```

Produces a list of paths and #f, which is normally used to initialized current-library-collection-links, as follows:

- The list starts with #f, which causes the default module name resolver, collection-file-path, and collection-path to try paths in current-library-collection-paths before collection links files.
- As long as the values of use-user-specific-search-paths and use-collection-link-paths are true, the second element in the result list is the path of the user–specific collection links file, which is (build-path (find-system-path 'addon-dir) (get-installation-name) "links.rktd").
- As long as the value of use-collection-link-paths is true, the rest of the list contains the result of get-links-search-files. Typically, that function produces a list with a single path, (build-path (find-config-dir) "links.rktd").

```
(collection-file-path file
                     collection ...+
                     [#:check-compiled? check-compiled?])
 → path?
  file : path-string?
  collection : path-string?
  check-compiled? : any/c = (regexp-match? #rx"[.]rkt$" file)
(collection-file-path file
                     collection ...+
                     #:fail fail-proc
                     [#:check-compiled? check-compiled?]) → any
  file : path-string?
  collection : path-string?
  fail-proc : (string? . -> . any)
  check-compiled? : any/c = (regexp-match? #rx"[.]rkt$" file)
```

Returns the path to the file indicated by *file* in the collection specified by the *collection*s, where the second *collection* (if any) names a sub-collection, and so on. The search uses the values of current-library-collection-links and current-library-collection-paths.

If *file* is not found, but *file* ends in ".rkt" and a file with the suffix ".ss" exists, then the directory of the ".ss" file is used. If *file* is not found and the ".rkt"/".ss" conversion does not apply, but a directory corresponding to the *collection*s is found, then a path using the first such directory is returned.

If *check-compiled?* is true, then the search also depends on use-compiled-file-paths and current-compiled-file-roots; if *file* is not found, then a compiled form of *file* with the suffix ".zo" is checked in the same way as the default compiled-load handler. If a compiled file is found, the result from collection-file-path reports the location that *file* itself would occupy (if it existed) for the found compiled file.

Finally, if the collection is not found, and if *fail-proc* is provided, then *fail-proc* is applied to an error message (that does not start "collection-file-path:" or otherwise claim a source), and its result is the result of collection-file-path. If *fail-proc* is not provided and the collection is not found, then the exn:fail:filesystem exception is raised.

Changed in version 6.0.1.12 of package base: Added the *check-compiled?* argument.

```
(collection-path collection ...+) → path?
  collection : path-string?
(collection-path collection
                 ...+
                 #:fail fail-proc) → any
  collection : path-string?
  fail-proc : (string? . -> . any)
```

> **NOTE:** This function is deprecated; use collection-file-path, instead. Collection splicing implies that a given collection can have multiple paths, such as when multiple packages provide modules for a collection.

Like collection-file-path, but without a specified file name, so that a directory indicated by *collection*s is returned.

When multiple directories correspond to the collection, the first one found in the search sequence (see §18.2.1 "Collection Search Configuration") is returned.

```
(current-library-collection-paths)
 → (listof (and/c path? complete-path?))
(current-library-collection-paths paths) → void?
  paths : (listof (and/c path-string? complete-path?))
```

Parameter that determines a list of complete directory paths for finding libraries (as referenced in require, for example) through the default module name resolver and for finding paths through collection-path and collection-file-path. See §18.2.1 "Collection Search Configuration" for more information.

```
(current-library-collection-links)
    (listof (or/c #f
                  (and/c path? complete-path?)
 →                  (hash/c (or/c (and/c symbol? module-path?) #f)
                            (listof (and/c path? complete-path?))))))
(current-library-collection-links paths) → void?
          (listof (or/c #f
                        (and/c path-string? complete-path?)
  paths :               (hash/c (or/c (and/c symbol? module-path?) #f)
                                (listof (and/c path-string? complete-path?))))))
```

Parameter that determines collection links files, additional paths, and the relative search order of `current-library-collection-paths` for finding libraries (as referenced in `require`, for example) through the default module name resolver and for finding paths through `collection-path` and `collection-file-path`. See §18.2.1 “Collection Search Configuration” for more information.

```
(use-user-specific-search-paths) → boolean?
(use-user-specific-search-paths on?) → void?
  on? : any/c
```

Parameter that determines whether user-specific paths, which are in the directory produced by `(find-system-path 'addon-dir)`, are included in search paths for collections and other files. For example, the initial value of `find-library-collection-paths` omits the user-specific collection directory when this parameter’s value is `#f`.

If `-U` or `--no-user-path` argument to `racket`, then `use-user-specific-search-paths` is initialized to `#f`.

```
(use-collection-link-paths) → boolean?
(use-collection-link-paths on?) → void?
  on? : any/c
```

Parameter that determines whether collection links files are included in the result of `find-library-collection-links`.

If this parameter’s value is `#f` on start-up, then collection links files are effectively disabled permanently for the Racket process. In particular, if an empty string is provided as the `-X` or `--collects` argument to `racket`, then not only is `current-library-collection-paths` initialized to the empty list, but `use-collection-link-paths` is initialized to `#f`.

18.3 Interactive Help

`(require racket/help)` package: `base`

The bindings documented in this section are provided by the `racket/help` and `racket/init` libraries, which means that they are available when the Racket executable is started with no command-line arguments. They are not provided by `racket/base` or `racket`.

```
help
(help string ...)
(help id)
(help id #:from module-path)
(help #:search datum ...)
```

For general help, see the main documentation page.

The `help` form searches the documentation and opens a web browser (using the user's selected browser) to display the results.

See `net/sendurl` for information on how the user's browser is launched to display help information.

A simple `help` or `(help)` form opens the main documentation page.

The `(help string ...)` form—using literal strings, as opposed to expressions that produce strings—performs a string-matching search. For example,

```
(help "web browser" "firefox")
```

searches the documentation index for references that include the phrase "web browser" or "firefox."

A `(help id)` form looks for documentation specific to the current binding of `id`. For example,

```
(require net/url)
(help url->string)
```

opens a web browser to show the documentation for `url->string` from the `net/url` library.

For the purposes of `help`, a `for-label` require introduces a binding without actually executing the `net/url` library—for cases when you want to check documentation, but cannot or do not want to run the providing module.

```
(require racket/gui) ; does not work in racket
(require (for-label racket/gui)) ; ok in racket
(help frame%)
```

If `id` has no for-label and normal binding, then `help` lists all libraries that are known to export a binding for `id`.

The (help id #:from module-path) variant is similar to (help id), but using only the exports of module-path. (The module-path module is required for-label in a temporary namespace.)

```
(help frame% #:from racket/gui) ; equivalent to the above
```

The (help #:search datum ...) form is similar to (help string ...), where any non-string form of datum is converted to a string using display. No datum is evaluated as an expression.

For example,

```
(help #:search "web browser" firefox)
```

also searches the documentation index for references that include the phrase "web browser" or "firefox."

18.4 Interactive Module Loading

The racket/rerequire and racket/enter libraries provide support for loading, reloading, and using modules.

18.4.1 Entering Modules

(require racket/enter) package: base

The bindings documented in this section are provided by the racket/enter and racket/init libraries, which means that they are available when the Racket executable is started with no command-line arguments. They are not provided by racket/base or racket.

```
(enter! module-path)
(enter! #f)
(enter! module-path flag ...+)

flag = #:quiet
     | #:verbose-reload
     | #:verbose
     | #:dont-re-require-enter
```

Intended for use in a REPL, such as when racket is started in interactive mode. When a module-path is provided (in the same sense as for require), the corresponding module

is loaded or invoked via dynamic-rerequire, and the current namespace is changed to the body of the module via module->namespace. When #f is provided, then the current namespace is restored to the original one.

Additional *flag*s can customize aspects of enter!:

- The #:verbose, #:verbose-reload, and #:quiet flags correspond to 'all, 'reload, and 'none verbosity for dynamic-rerequire. The default corresponds to #:verbose-reload.
- After switching namespaces to the designated module, enter! automatically requires racket/enter into the namespace, so that enter! can be used to switch namespaces again. In some cases, requiring racket/enter might not be desirable (e.g., in a tool that uses racket/enter); use the #:dont-re-require-enter flag to disable the require.

```
(dynamic-enter!  mod
                [#:verbosity verbosity
                 #:re-require-enter? re-require-enter?]) → void?
  mod : (or/c module-path? #f)
  verbosity : (or/c 'all 'reload 'none) = 'reload
  re-require-enter? : any/c = #t
```

Procedure variant of enter!, where *verbosity* is passed along to dynamic-rerequire and *re-require-enter?* determines whether dynamic-enter! requires racket/enter in a newly entered namespace.

Added in version 6.0.0.1 of package base.

18.4.2 Loading and Reloading Modules

(require racket/rerequire) package: base

The bindings documented in this section are provided by the racket/rerequire library, not racket/base or racket.

```
(dynamic-rerequire  module-path
                    [#:verbosity verbosity]) → void?
  module-path : module-path?
  verbosity : (or/c 'all 'reload 'none) = 'reload
```

Like (dynamic-require *module-path* 0), but with reloading support. The dynamic-rerequire function is intended for use in an interactive environment, especially via enter!.

If invoking *module-path* requires loading any files, then modification dates of the files are recorded. If the file is modified, then a later dynamic-rerequire re-loads the module from source; see also §1.1.10.4 "Module Redeclarations". Similarly if a later dynamic-rerequire transitively requires a modified module, then the required module is re-loaded. Re-loading support works only for modules that are first loaded (either directly or indirectly through transitive requires) via dynamic-rerequire.

When enter! loads or re-loads a module from a file, it can print a message to (current-error-port), depending on *verbosity*: 'verbose prints a message for all loads and re-loads, 'reload prints a message only for re-loaded modules, and 'none disables printouts.

18.5 Debugging

Racket's built-in debugging support is limited to context (i.e., "stack trace") information that is printed with an exception. In some cases, disabling the JIT compiler can affect context information. The errortrace library supports more consistent (independent of the JIT compiler) and precise context information. The racket/trace library provides simple tracing support. Finally, the DrRacket programming environment provides much more debugging support.

18.5.1 Tracing

(require racket/trace) package: base

The bindings documented in this section are provided by the racket/trace library, not racket/base or racket.

The racket/trace library mimics the tracing facility available in Chez Scheme.

(trace *id* ...)

Each *id* must be bound to a procedure in the environment of the trace expression. Each *id* is set!ed to a new procedure that traces procedure calls and returns by printing the arguments and results of the call via current-trace-notify. If multiple values are returned, each value is displayed starting on a separate line.

When traced procedures invoke each other, nested invocations are shown by printing a nesting prefix. If the nesting depth grows to ten and beyond, a number is printed to show the actual nesting depth.

The trace form can be used on an identifier that is already traced. In this case, assuming that the variable's value has not been changed, trace has no effect. If the variable has been changed to a different procedure, then a new trace is installed.

Tracing respects tail calls to preserve loops, but its effect may be visible through continuation marks. When a call to a traced procedure occurs in tail position with respect to a previous traced call, then the tailness of the call is preserved (and the result of the call is not printed for the tail call, because the same result will be printed for an enclosing call). Otherwise, however, the body of a traced procedure is not evaluated in tail position with respect to a call to the procedure.

The result of a `trace` expression is `#<void>`.

Examples:

```
> (define (f x) (if (zero? x) 0 (add1 (f (sub1 x)))))

> (trace f)

> (f 10)
>(f 10)
> (f 9)
> >(f 8)
> > (f 7)
> > >(f 6)
> > > (f 5)
> > > >(f 4)
> > > > (f 3)
> > > > >(f 2)
> > > > > (f 1)
> > > >[10] (f 0)
< < < <[10] 0
< < < < < 1
< < < < <2
< < < < 3
< < < <4
< < < 5
< < <6
< < 7
< <8
< 9
<10
10
```

```
(trace-define id expr)
(trace-define (head args) body ...+)
```

The `trace-define` form is short-hand for first defining a function then tracing it. This form supports all `define` forms.

Examples:

```
> (trace-define (f x) (if (zero? x) 0 (add1 (f (sub1 x)))))

> (f 5)
>(f 5)
> (f 4)
> >(f 3)
> > (f 2)
> > >(f 1)
> > > (f 0)
< < < 0
< < <1
< < 2
< <3
< 4
<5
5
```

Examples:

```
> (trace-define ((+n n) x) (+ n x))

> (map (+n 5) (list 1 3 4))
>(+n 5)
<#<procedure>
'(6 8 9)
```

```
(trace-define-syntax id expr)
(trace-define-syntax (head args) body ...+)
```

The `trace-define-syntax` form is short-hand for first defining a macro then tracing it. This form supports all `define-syntax` forms.

For example:

Examples:

```
> (trace-define-syntax fact
    (syntax-rules ()
      [(_ x) 120]))

> (fact 5)
>(fact #<syntax:11:0 (fact 5)>)
<#<syntax:11:0 120>
120
```

By default, `trace` prints out syntax objects when tracing a macro. This can result in too much output if you do not need to see, e.g., source information. To get more readable output, try this:

Examples:

```
> (require (for-syntax racket/trace))

> (begin-for-syntax
    (current-trace-print-args
      (let ([ctpa (current-trace-print-args)])
        (lambda (s l kw l2 n)
          (ctpa s (map syntax->datum l) kw l2 n))))
    (current-trace-print-results
      (let ([ctpr (current-trace-print-results)])
        (lambda (s l n)
         (ctpr s (map syntax->datum l) n)))))

> (trace-define-syntax fact
  (syntax-rules ()
    [(_ x) #'120]))

> (fact 5)
>(fact '(fact 5))
<'#'120
#<syntax:14:0 120>
```

```
(trace-lambda [#:name id] args expr)
```

The `trace-lambda` form enables tracing an anonymous function. This form will attempt to infer a name using `syntax-local-infer-name`, or a name can be specified using the optional `#:name` argument. A syntax error is raised if a name is not given and a name cannot be inferred.

Example:

```
> ((trace-lambda (x) 120) 5)
>(eval:16:0 5)
<120
120
```

```
(trace-let id ([arg expr] ...+) body ...+)
```

The `trace-let` form enables tracing a named let.

Example:

```
> (trace-let f ([x 5])
    (if (zero? x)
        1
        (* x (f (sub1 x)))))
```

```
>(f 5)
> (f 4)
> >(f 3)
> > (f 2)
> > >(f 1)
> > > (f 0)
< < < 1
< < <1
< < 2
< <6
< 24
<120
120
```

```
(untrace id ...)
```

Undoes the effects of the `trace` form for each *id*, `set!`ing each *id* back to the untraced procedure, but only if the current value of *id* is a traced procedure. If the current value of a *id* is not a procedure installed by `trace`, then the variable is not changed.

The result of an `untrace` expression is `#<void>`.

```
(current-trace-notify) → (string? . -> . any)
(current-trace-notify proc) → void?
  proc : (string? . -> . any)
```

A parameter that determines the way that trace output is displayed. The string given to *proc* is a trace; it does not end with a newline, but it may contain internal newlines. Each call or result is converted into a string using `pretty-print`. The parameter's default value prints the given string followed by a newline to `(current-output-port)`.

```
(trace-call id proc #:<kw> kw-arg ...) → any/c
  id : symbol?
  proc : procedure?
  kw-arg : any/c
```

Calls *proc* with the arguments supplied in `args`, and possibly using keyword arguments. Also prints out the trace information during the call, as described above in the docs for `trace`, using *id* as the name of *proc*.

```
(current-trace-print-args) →  (-> symbol?
                                  list?
                                  (listof keyword?)
                                  list?
                                  number?
                                  void?)
```

```
(current-trace-print-args trace-print-args) → void?
                      (-> symbol?
                          list?
  trace-print-args :      (listof keyword?)
                          list?
                          number?
                          void?)
```

The value of this parameter is invoked to print out the arguments of a traced call. It receives the name of the function, the function's ordinary arguments, its keywords, the values of the keywords, and a number indicating the depth of the call.

```
                                  (-> symbol?
                                      list?
(current-trace-print-results) →       number?
                                      any)
(current-trace-print-results trace-print-results) → void?
                        (-> symbol?
                            list?
  trace-print-results :     number?
                            any)
```

The value of this parameter is invoked to print out the results of a traced call. It receives the name of the function, the function's results, and a number indicating the depth of the call.

```
(current-prefix-in) → string?
(current-prefix-in prefix) → void?
  prefix : string?
```

This string is used by the default value of `current-trace-print-args` indicating that the current line is showing the a call to a traced function.

It defaults to `">"`.

```
(current-prefix-out) → string?
(current-prefix-out prefix) → void?
  prefix : string?
```

This string is used by the default value of `current-trace-print-results` indicating that the current line is showing the result of a traced call.

It defaults to `"<"`.

18.6 Kernel Forms and Functions

`#lang racket/kernel` package: `base`

The racket/kernel library is a cross-phase persistent module that provides a minimal set of syntactic forms and functions.

"Minimal" means that racket/kernel includes only forms that are built into the Racket compiler and only functions that are built into the run-time system. Currently, the set of bindings is not especially small, nor is it particularly well-defined, since the set of built-in functions can change frequently. Use racket/kernel with care, and beware that its use can create compatibility problems.

The racket/kernel module exports all of the bindings in the grammar of fully expanded programs (see §1.2.3.1 "Fully Expanded Programs"), but it provides #%plain-lambda as lambda and λ, #%plain-app as #%app, and #%plain-module-begin as #%module-begin. Aside from #%datum (which expands to quote), racket/kernel provides no other syntactic bindings.

The racket/kernel module also exports many of the function bindings from racket/base, and it exports a few other functions that are not exported by racket/base because racket/base exports improved variants. The exact set of function bindings exported by racket/kernel is unspecified and subject to change across versions.

(require racket/kernel/init) package: base

The racket/kernel/init library re-provides all of racket/kernel. It also provides #%top-interaction, which makes racket/kernel/init useful with the -I command-line flag for racket.

Bibliography

[C99] ISO/IEC, "ISO/IEC 9899:1999 Cor. 3:2007(E))." 2007.

[Culpepper07] Ryan Culpepper, Sam Tobin-Hochstadt, and Matthew Flatt, "Advanced Macrology and the Implementation of Typed Scheme," Workshop on Scheme and Functional Programming, 2007.

[Danvy90] Olivier Danvy and Andre Filinski, "Abstracting Control," LISP and Functional Programming, 1990.

[Felleisen88a] Matthias Felleisen, "The theory and practice of first-class prompts," Principles of Programming Languages, 1988.

[Felleisen88] Matthias Felleisen, Mitch Wand, Dan Friedman, and Bruce Duba, "Abstract Continuations: A Mathematical Semantics for Handling Full Functional Jumps," LISP and Functional Programming, 1988.

[Flatt02] Matthew Flatt, "Composable and Compilable Macros: You Want it When?," International Conference on Functional Programming (ICFP), 2002.

[Friedman95] Daniel P. Friedman, C. T. Haynes, and R. Kent Dybvig, "Exception system proposal," web page, 1995. `http://www.cs.indiana.edu/scheme-repository/doc.proposals.exceptions.html`

[Gasbichler02] Martin Gasbichler and Michael Sperber, "Processes vs. User-Level Threads in Scsh," Workshop on Scheme and Functional Programming, 2002.

[Gunter95] Carl Gunter, Didier Remy, and Jon Rieke, "A Generalization of Exceptions and Control in ML-like Languages," Functional Programming Languages and Computer Architecture, 1995.

[Haynes84] Christopher T. Haynes and Daniel P. Friedman, "Engines Build Process Abstractions," Symposium on LISP and Functional Programming, 1984.

[Hayes97] Barry Hayes, "Ephemerons: a New Finalization Mechanism," Object-Oriented Languages, Programming, Systems, and Applications, 1997.

[Hieb90] Robert Hieb and R. Kent Dybvig, "Continuations and Concurrency," Principles and Practice of Parallel Programming, 1990.

[L'Ecuyer02] Pierre L'Ecuyer, Richard Simard, E. Jack Chen, and W. David Kelton, "An Object-Oriented Random-Number Package With Many Long Streams and Substreams," Operations Research, 50(6), 2002.

[Queinnec91] Queinnec and Serpette, "A Dynamic Extent Control Operator for Partial Continuations," Principles of Programming Languages, 1991.

[Shan04] Ken Shan, "Shift to Control," Workshop on Scheme and Functional Programming, 2004.

[Sperber07] Michael Sperber, R. Kent Dybvig, Matthew Flatt, and Anton van Straaten (editors), "The Revised^6 Report on the Algorithmic Language Scheme." 2007. `http://www.r6rs.org/`

[Sitaram90] Dorai Sitaram and Matthias Felleisen, "Control Delimiters and Their Hierarchies," *Lisp and Symbolic Computation*, 1990.

[Sitaram93] Dorai Sitaram, "Handling Control," Programming Language Design and Implementation, 1993.

[SRFI-42] Sebastian Egner, "SRFI-42: Eager Comprehensions," SRFI, 2003. `http://srfi.schemers.org/srfi-42/`

[Strickland12] T. Stephen Strickland, Sam Tobin-Hochstadt, Matthew Flatt, and Robert Bruce Findler, "Chaperones and Impersonators: Run-time Support for Reasonable Interposition," Object-Oriented Programming, Systems, and Languages (OOPSLA), 2012. http://www.eecs.northwestern.edu/~robby/pubs/papers/oopsla2012-stff.pdf

Index